

OS Zusammenfassung

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Abstract

This lecture notes are for me to learn LaTeX and review the operating system lecture

1 Introduction

1.1 Overview

The Operating System

- **Provides abstraction layer**
Manages and hides hardware details
Low-level interfaces to access hardware
Multiplexes hardware to multiple programs
makes hardware use efficient for applications
- **Provides protection** from
users/processes using up all resources
processes writing into other processes memory
- **is a resource manager**
Manages and multiplexes hardware resources
Decides between conflicting requests for resource use
Strives for efficient and fair resource use
- **is a control program**
Controls execution of programs
Prevents errors and improper use of the computer

There are no universally accepted definitions

1.2 Hardware

CPU (Central Processing Unit)

- Fetches instructions from memory and executes them
- Internal registers store data and metadata during execution
- **User Mode (x86: "Ring 3" or CPL3)**
Only non-privileged instructions, no hardware management in this mode for protection
- **Kernel Mode (x86: "Ring 0" or CPL0)**
All instruction allowed, including privileged instructions

RAM (Random Access Memory) keeps currently executing instructions and data

Caching

- Ram delivers instruction/data slower than the CPU can execute
- Memory references typically follow principle of locality
- **Caching** helps mitigating this **Memory Wall**
Information in use are copied from slower to faster storage. When needed, check whether it is in faster storage before going down in the Memory Hierarchy, then copy it to cache to be used from there

Access times

- | | |
|------------------------------|---|
| • CPU registers | 1 CPU cycle |
| • L1 cache per core | 4 CPU cycles |
| • L2 cache per pair of cores | 12 CPU cycles |
| • L3 cache | 28 CPU cycles (25 GiB/s) ‘ |
| • DDR3-Ram | 28 CPU cycle for LLC + 50ns (12 GiB/s) |

CPU Cache Organization

- Caches divided up into cache lines (often 64 bytes each)
- Separation of data and instructions in faster caches
- **Cache hit**: Data already in cache
- **Cache miss**: Data has to be fetched from lower level first
- Types of Cache misses
 - **Compulsory Miss**: first reference miss, data has never been accessed
 - **Capacity Miss**: cache not large enough for Working Set of process
 - **Conflict Miss**: cache still has space, but collision due to placement strategy

Device Control

- Device controller accepts command from the OS via device driver
- Control by writing into device register and read status by reading it
- Data transfer by writing/reading device memory
- Port-mapped I/O (PMIO) special CPU instructions to access port-mapped registers and memory
- Memory-mapped I/O (MMIO) same address space for Ram and device memory

Devices can signal the CPU through interrupts

1.3 OS Invocation

Operating System Kernel does **not** always run in the background

Tree occasions invoke the Kernel and switch to kernel-mode

- | | |
|----------------|--|
| • System calls | User-mode process requires higher privileges |
| • Interrupts | CPU-external device sends a signal |
| • Exceptions | CPU signals an unexpected condition |

System Calls

The main Idea behind System calls is the necessity to protect processes from one another. So processes are running in User-Mode. The OS provides services, which the applications can invoke in System Calls/syscalls, in order to get the action performed by the OS, on behalf of application

Syscall interface between applications and OS provides a limited number of well-defined entry points to the kernel

Application Program Interfaces (API) brings another level of abstraction between applications and Programmers (API invokes Syscalls invokes Kernel-Mode operations)

One single entry point to the kernel for all System calls, the **trap**. Trap switches CPU to kernel mode and enters the kernel in the same, predefined way for every syscall. The system call dispatcher in the kernel acts as a multiplexer for all syscalls.

syscalls identified by a number, passed as parameter, **system call table** maps **system call number** to kernel funktion, dispatcher decides where to jump based on the number and table.

Programs have the System call number compiled in! Never reuse old numbers in future versions of kernel

Intrrupts

Interrupts are used by devices to signal predefined conditions to the OS, they are managed by the Programmable Interrupt Controller. When masked the interrupts are only delivered, when unmasked.

interrupt vector: table pinned in memory containing the addresses of all service routines **interrupt**

service routine: takes the control in order to handle a specific interrupt. Saves the state of the interrupted process

- Instruction pointer
- Stack pointer
- Status word

Exceptions

- Generated by the CPU itself, if an unusual condition makes it impossible to continue processing
- CPU interrupts program and relegate control to the kernel
- Kernel determines the reason for exceptions
- If kernel can resolve the problem, it does so and continue the **faulting instruction**
- Otherwise process get killed

Interrupts can happen in **any** context, Exceptions always occur **synchronous to** and **in the context** of a process.

2 OS Concepts

Early on, programs were load directly into **physical memory**. If the program was too large, the programmer had to manually partition his program into **overlays**. OS could swap between disk and memory.

Problems: Buggy programs trash other programs, malicious jobs can read other program's operations, Jobs can take all memory for themselves,...

address spaces: every job has his own address space, so they can't reach other jobs addresses. Jobs only use virtual addresses.

MMU (memory management unit): translates virtual address to physical address

- allows kernel-only virtual addresses
- can enforce read-only virtual addresses
- can enforce execute disable

Not all addresses need to be mapped at all times. If a virtual address is not mapped, the MMU throws a **page fault** exception. Handled by loading the faulting address and then continuing the program.

over-commitment: more memory than physically available. Page faults also issued by MMU on illegal memory access.

A **process** is a program in execution, associated with a process control block (PCB) and with a virtual **address space (AS)**. AS is the only memory a program can name and starts at 0 for every program. AS are layed out in sections. Memory acces between those sections is illegal and causes a page fault, called **sementation fault**. Segmentation faults results in the process getting killed by the OS.

A section has the following layout:

- Stack: Function history and local variables
- Data: Constatnts, static variables, global variables, strings
- Text: Program code

Threads represents execution states of a program

- Instruction pointer (IP) register stores currently executed instruction
- Stack pointer (SP) register stores the address of the top of the stack
- Program status word (PSW) contains flags about executeion history
- ...

Two things to consider when designing an OS:

Mechanism: Implementation of what is done

Policy: The rules which decide when what is done and how much

Operating System need to handle multiple processes and threads in order to provide multi-tasking. The **scheduler** decides which job to run next, while the **dispatcher** performs the task-switching. Schedulers provide fairness while trying to reach goals after setted priorities.

Persistent Data is for users is stored in flies and directories. A file is associated with file name and offset with bytes. Directories associate directory names with eigher directory names or file names.

The **file system** is an ordered collection of blocks, what can be operated on by programmers, with operations like open, read, seek, ...

Processes communicate directly through a special **named pipe** file

Directories form a **directory tree/file hierarchy**. The **root directory** is the topmost directory of a directory tree. Files can be accessed by their **path name**.

OS abstract the view of information storage to file systems. Drivers hide specific hardware device. OS increases the performance of I/O devices by

- **Buffering**: Store data temporarily while it is being transferred
- **Caching**: Store parts of data in faster storage for performance
- **Spooling**: Overlap of output of one job with input of other jobs

3 Processes

3.1 Process Abstraction

Multiprogramming is the art of switching quickly between processes
Process abstraction

Glossary

physical address address of real memory. 4

process control block (PCB) Informations about allocated resources of a process. 4

virtual address address in process' address space. 4