

Session 7 - Leetcode Problem "Two Sum"

=====

in python we do not have Array..

instead we have a list

[1,2,3,4,5] - Array

[1,'sumit',3.5] - List

from python 3.5 version onwards we can give hints while creating lists..

Python 3 is the language in which we will code

<https://leetcode.com/> is the platform from where we will pick the questions.

you do not need a premium account.

1st test case

nums = [2,7,11,15]

target = 9

output = [0,1]

2nd test case

nums = [3,2,4], target = 6

output = [1,2]

3rd test case

Input: nums = [3,3], target = 6

Output: [0,1]

To practise coding in python 3 we can use google colab

<https://colab.research.google.com>

class Solution:

def twoSum(self, nums: List[int], target: int) -> List[int]:

input

[1,5,8,4,9,3,18,2]

outer loop - 1
inner loop - 5,8,4,9,3,18,2

outer loop - 5
inner loop - 8,4,9,3,18,2

outer loop - 8
inner loop - 4,9,3,18,2

outer loop - 4
inner loop - 9,3,18,2

outer loop - 9
inner loop - 3,18,2

outer loop - 3
inner loop - 18,2

outer loop - 18
inner loop - 2

.
target
27

output
[4,6]

1st approach - brute force..

$O(n^2)$

code :

```
from typing import List
class Solution:
    def twoSum(self, nums: list, target: int) -> list:
        for x in range(0,len(nums)-1):
            for y in range(x+1,len(nums)):
                if (nums[x] + nums[y] == target):
                    return [x,y]
```

```
s1 = Solution()
s1.twoSum([1,5,8,4,9,3,18,2],20)
```

! python --version

Python 3.8.10

2nd Approach -

[1,5,8,4,9,3,18,2]
(1,4)

lets sort the array - $O(n \log n)$

[1,2,3,4,5,8,9,18] - target 14
(4,6)

14 is sum of (5,9)

. .

we would scan each element at the max once to get the answer

if the sum of numbers at 2 pointers is more than the target than we shift the 2nd pointer one step backwards

if the sum of numbers at 2 pointers is less than the target than we shift the 1st pointer one step forward.

$O(n \log n) + O(n) \sim O(n \log n)$

[1,5,8,4,9,3,18,2]

correct answer should be (0,6)

[1,2,3,4,5,8,9,18] - target 19

if I consider just sorted array (0,7) which is not the correct answer.

```
if(new_nums[i] + new_nums[j] == target):  
    return [nums.index(new_nums[i]),nums.index(new_nums[j])]
```

nums [3,3]

target 6

[3 ,0, 5, 2 ,3, 10] target 6

[0,2,3,3,5,10]

1st element is 3

2nd element is 3

code:

```
from typing import List
class Solution:
    def twoSum(self, nums: List[int], target: int) -> List[int]:
        new_nums = nums.copy()
        new_nums.sort()
        i = 0
        j = len(nums) - 1

        while (i < j):
            if(new_nums[i] + new_nums[j] == target):
                if (new_nums[i] == new_nums[j]):
                    return [nums.index(new_nums[i]),nums.index(new_nums[j],nums.index(new_nums[i])+1)]
                else:
                    return [nums.index(new_nums[i]),nums.index(new_nums[j])]
            elif (new_nums[i] + new_nums[j] < target):
                i = i + 1
            else:
                j = j - 1

s1 = Solution()
s1.twoSum([1,5,8,4,9,3,18,2],27)
```

3rd Approach -

dictionary

in java we have a hash map

in python we have dictionary

it stores key and value

```
employee_list = []
```

```
dict = {'sumit' : 9900100901, 'kapil' : 99001000902}
```

```
dict['satish'] = 9090909090
```

```
print(dict)
```

[1,5,8,4,9,3,18,2] 27 is my target element

27-18 = 9

dict =

1->0

5->1
8->2
4->3
9->4
3->5

(4,6)

space complexity = $O(n)$
time complexity = $O(n)$

code:

```
class Solution:
    def twoSum(self, nums: List[int], target: int) -> List[int]:
        dict = {}
        for i in range(0, len(nums)):
            complement = target - nums[i]
            if complement in dict:
                return [i, dict[complement]]
            dict[nums[i]] = i
```

Leetcode - Single Number

=====

input
[5,6,7,5,7]

output
6

solution - 1

brute force solution

$O(n^2)$ it involves 2 for loops

time complexity was $O(n^2)$
space complexity was $O(1)$

solution - 2

what is a hashmap

key -> value

[5,6,7,5,7]

dict[5]

key value

5 -> 2

6 -> 1

7 -> 2

in case of python (dictionary)

key -> value

extra space or auxillary space

$O(n)$

time complexity $O(n)$

dict

[5,6,7,5,7]

hash_table[5] = hash_table[5] + 1

5 -> 2

6 -> 1

7 -> 2

class Solution:

def singleNumber(self, nums: List[int]) -> int:

hash_table = defaultdict(int)

for num in nums:

hash_table[num] = hash_table[num] + 1

for num in nums:

if hash_table[num] == 1:

return num

time complexity = $O(n)$

space complexity = $O(n)$

solution 3

=====

[4,1,2,1,2]

xor - when we xor a number even number of times it returns 0

and when we xor a number odd number of time it returns the num itself

result = 0

$result = result \wedge 4 \wedge 1 \wedge 2 \wedge 1 \wedge 2$

class Solution:

def singleNumber(self, nums: List[int]) -> int:

result = 0

for num in nums:

result = result ^ num

return result

time complexity = $O(n)$

space complexity = $O(1)$

Majority Element - Leetcode

=====

Brute force - $O(n^2)$

[2,2,1,1,1,2,2]

[1,1,1,2,2,2,2] $O(n \log n)$

$O(1)$

$O(n \log n) + O(1) \sim O(n \log n)$

$9/2 = 4.5$

$\text{floor}(4.5) = 4$

class Solution:

def majorityElement(self, nums: List[int]) -> int:

nums.sort()

return nums[math.floor(len(nums)/2)]

Time complexity = $O(n \log n)$

space complexity = $O(1)$

Boyer-Moore voting algorithm

Time complexity = $O(n)$

[2,2,1,1,1,2,2]

majority element = 2

count = 1

majority element = 2

count = 2

majority element = 2
count = 1

majority element = 2
count = 0

majority element = 1
count = 1

majority element = 1
count = 0

majority element = 2
count = 1

majority_element = 1
count = 1

[2,2,1,1 | 1,2 | 2]

[3]

majority element = 3
count = 1

[2,2,3,3 | 4,4,2]

majority element = 4
count = 1

[2,2,1,1,1,2,2]

class Solution:

```
def majorityElement(self, nums: List[int]) -> int:
    count = 0
    majority = None
    for num in nums:
        if count == 0:
            majority = num
        if majority != num:
            count = count - 1
        else:
            count = count + 1
    return majority
```

[2,2,1,1,1,2,2]

majority = 2

count = 1

In a single pass we are able to handle this..

O(n)

no extra space required...

Best Time to Buy and Sell Stock

=====

[5,8,9,4,3,10]

min_num = 3

profit = 7

max_profit = 7

[7,1,5,3,6,4]

profit = current_num - min_num 5

max_profit = 5

min_num = 1

class Solution:

def maxProfit(self, prices: List[int]) -> int:

max_profit = 0

min_num = prices[0]

for i in range (1,len(prices)):

if prices[i] < min_num :

min_num = prices[i]

profit = prices[i] - min_num

if profit > max_profit :

max_profit = profit

return max_profit

Best Time to Buy and Sell Stock II

=====

[7,1,5,3,6,4]

5-1 = 4

6-3 = 3

4+3 = 7

[1,2,3,4,5]

total_profit = 0

```

for i in range (1,len(prices)):
    if prices[i] > prices[i-1]:
        total_profit = total_profit + prices[i] - prices[i-1]

```

```

return total_profit

```

buy at 1

sell at 5

5-1 = 4

buy - 1

sell - 2

buy - 2

sell - 3

buy - 3

sell - 4

buy - 4

sell - 5

total_profit = 7

[7,1,5,3,6,4]

.

class Solution:

```

def maxProfit(self, prices: List[int]) -> int:
    total_profit = 0
    for i in range (1,len(prices)):
        if prices[i] > prices[i-1]:
            total_profit = total_profit + prices[i] - prices[i-1]
    return total_profit

```

Remove Duplicates from Sorted Array

=====

input array

[1,1,1,3,3,5,5,5,5,8,8,9]

[1,3,5,8,9,5,5,5,5,8,8,9]

class Solution:

```

def removeDuplicates(self, nums: List[int]) -> int:
    size = len(nums)
    insertIndex = 1
    for i in range(1,size):

```

```

        if nums[i-1] != nums[i] :
            nums[insertIndex] = nums[i]
            insertIndex = insertIndex + 1
    return insertIndex

```

Rotate Array

=====

[1,2,3,4,5,6,7]

k=7

[7,1,2,3,4,5,6]

[6,7,1,2,3,4,5]

[5,6,7,1,2,3,4]

size = 7

rotate it by k = 9

k%size

9%7 = 2

class Solution:

```

    def reverse(self, nums:list, start:int, end:int) -> None:
        while start < end:
            nums[start] , nums[end] = nums[end], nums[start]
            start = start + 1
            end = end - 1

```

```

    def rotate(self, nums: List[int], k: int) -> None:
        size = len(nums)
        k = k%size
        self.reverse(nums,0,size-1)
        self.reverse(nums,0,k-1)
        self.reverse(nums,k,size-1)

```

Maximum Subarray

=====

[-2,1,-3,4,-1,2,1,-5,4]

.

maxsum = 6

sum = 5

2 important things whenever $\text{sum} < 0$
you make $\text{sum} = 0$ and move ahead

if $\text{sum} > \text{maxsum}$
 $\text{maxsum} = \text{sum}$

[5,4,-1,7,8]

$\text{maxsum} = 23$
 $\text{sum} = 23$

[-2,1,-3,4,-1,2,1,-5,4]

$\text{maxsum} = 6$
 $\text{sum} = 5$

class Solution:

```
def maxSubArray(self, nums: List[int]) -> int:
    sum = 0
    maxsum = nums[0]
    for num in nums:
        if sum < 0:
            sum = 0
        sum = sum + num
        if sum > maxsum:
            maxsum = sum
    return maxsum
```

time complexity = $O(n)$

space complexity = $O(1)$

Zero's and one's

=====

given an array of 0's and 1's

input

[0,1,0,1,0,0,1,1,0,1,0]

[0,0,0,0,0,0,1,1,1,1,1]

$O(n)$

in the first pass we have to take the count of 0's and 1's

and in the second pass we replace the elements...

$O(n)$

[0,0,0,0,0,1,1,1,1,1]

. .

whenever first pointer get a 0 we move it right

whenever second pointer gets a 1 we move it left

if first pointer is pointing on 1 and second pointer is pointing on 0 then we swap the elements

```
class Solution:
    def zerosAndOnes(self,nums: list) -> None :
        size = len(nums)
        first = 0
        second = size - 1
        while first < second:
            if nums[first] == 0:
                first = first + 1
            if nums[second] == 1:
                second = second - 1
            if nums[first] == 1 and nums[second] == 0 :
                nums[first] , nums[second] = nums[second] , nums[first]
        print(nums)
```

```
s = Solution()
s.zerosAndOnes([0,1,0,1,0,0,1,1,1,0,1,0,1,1,1])
```

Leaders in the array
=====

maxelement = 17

input {16, 17, 4, 3, 5, 2}

output = {2,5,17}

output { 17, 5, 2 }

brute force $O(n^2)$

$O(n)$

{16, 17, 4, 3, 5, 2}

size = len(nums) = 6

ans = [nums[size-1]]

ans = [2]

max = 2

```
start index = size-2 = 4
end index = -1
```

```
for i in range (5,-1,-1)
print(i)
```

```
class Solution:
def leaders(self,nums: list) -> list :
    size = len(nums)
    ans = [nums[size-1]]
    max = nums[size-1]
    for i in range(size-2,-1,-1):
        if nums[i] > max:
            max = nums[i]
            ans.append(nums[i])
    return ans
```

```
s = Solution()
ans = s.leaders([16,17,4,3,5,2])
print(ans)
```

Max Consecutive Ones

=====

```
input [1,1,0,1,1,1]
```

```
max_count = 3
count = 3
```

time complexity will be $O(n)$

```
class Solution:
def findMaxConsecutiveOnes(self, nums: List[int]) -> int:
    max_count = 0
    count = 0
    for num in nums:
        if num == 1:
            count = count + 1
        else:
            count = 0
        if count > max_count:
            max_count = count
    return max_count
```

Merge Sorted Array

=====

```
[1,2,3,0,0,0] - n p1
```

.
[2,5,6] - m p2
.

[1,2,2,3,5,6]

time complexity $O(n+m)$

space complexity $O(n+m)$

[1,2,3,2,5,6] - n+m

now you sort this...

$O((n+m)\log(n+m))$

[1,2,3,3,5,6] - n

.
[2,5,6] - m
.

class Solution:

def merge(self, nums1: List[int], m: int, nums2: List[int], n: int) -> None:

p1 = m-1

p2 = n-1

for x in range(n+m-1,-1,-1):

if p2 < 0 :

break

if nums1[p1] > nums2[p2] and p1 >= 0:

nums1[x] = nums1[p1]

p1 = p1 - 1

else:

nums1[x] = nums2[p2]

p2 = p2 - 1