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# Ramulator 2.0 Summary

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***Intelligent System  
Laboratory***

# How to measure off-chip power?

## □ Analysis based on the paper below

### BlockHammer: Preventing RowHammer at Low Cost by Blacklisting Rapidly-Accessed DRAM Rows

#### • Section 6. Hardware Complexity Analysis

##### ▪ In case of **on-chip** analysis

#### 6. Hardware Complexity Analysis

We evaluate BlockHammer's (1) chip area, static power, and access energy consumption using CACTI [99] and (2) circuit latency using Synopsys DC [143]. We demonstrate that BlockHammer's physical costs are competitive with state-of-the-art RowHammer mitigation mechanisms.

#### 6.2. Latency Analysis

We implement BlockHammer in Verilog HDL and synthesize our design using Synopsys DC [143] with a 65 nm process technology to evaluate the latency impact on memory accesses. According to our RTL model, which we open source [124], BlockHammer responds to an "Is this ACT RowHammer-safe?" query (1 in Figure 2) in only 0.97 ns. This latency can be hidden because it is one-to-two orders of magnitude smaller than the row access latency (e.g., 45–50 ns) that DRAM standards (e.g., DDRx, LPDDRx, GDDRx) enforce [36, 53, 55].

Mitigation Mechanism	$N_{RH}=32K^*$						$N_{RH}=1K$					
	SRAM KB	CAM KB	Area mm <sup>2</sup>	% CPU	Access Energy (pJ)	Static Power (mW)	SRAM KB	CAM KB	Area mm <sup>2</sup>	% CPU	Access Energy (pJ)	Static Power (mW)
BlockHammer	51.48	1.73	0.14	0.06	20.30	22.27	441.33	55.58	1.57	0.64	99.64	220.99
Dual counting Bloom filters	48.00	-	0.11	0.04	18.11	19.81	384.00	-	0.74	0.30	86.29	158.46
H3 hash functions	-	-	< 0.01	< 0.01	-	-	-	-	< 0.01	< 0.01	-	-
Row activation history buffer	1.73	1.73	0.03	0.01	1.83	2.05	55.58	55.58	0.83	0.34	12.99	62.12
AttackThrottler counters	1.75	-	< 0.01	< 0.01	0.36	0.41	1.75	-	< 0.01	< 0.01	0.36	0.41
PARA [73]	-	-	< 0.01	-	-	-	-	-	< 0.01	-	-	-
ProHIT [137]*	-	0.22	< 0.01	< 0.01	3.67	0.14	×	×	×	×	×	×
MrLoc [161]*	-	0.47	< 0.01	< 0.01	4.44	0.21	×	×	×	×	×	×
CBT [132]	16.00	8.50	0.20	0.08	9.13	35.55	512.00	272.00	3.95	1.60	127.93	535.50
TWICE [84]	23.10	14.02	0.15	0.06	7.99	21.28	738.32	448.27	5.17	2.10	124.79	631.98
Graphene [113]	-	5.22	0.04	0.02	40.67	3.11	-	166.03	1.14	0.46	917.55	93.96

\* ProHIT [137] and MrLoc [161] do not provide a concrete discussion on how to adjust their empirically-determined parameters for different  $N_{RH}$  values. Therefore, we (1) report their values for a fixed design point that each paper provides for  $N_{RH}=2K$  and (2) mark values we cannot estimate using an ×.

Table 4: Per-rank area, access energy, and static power of BlockHammer vs. state-of-the-art RowHammer mitigation mechanisms.

# How to measure off-chip power?

## □ Analysis based on the paper below

### BlockHammer: Preventing RowHammer at Low Cost by Blacklisting Rapidly-Accessed DRAM Rows

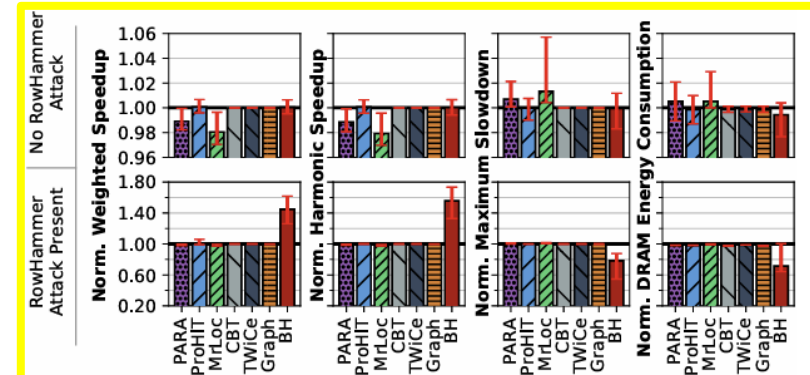
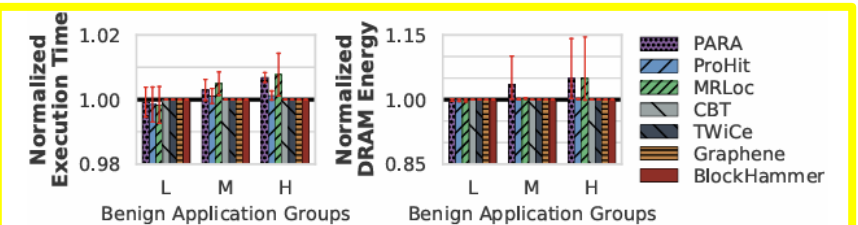
- Section 6. Hardware Complexity Analysis
  - In case of **off-chip** analysis

#### 7. Experimental Methodology

We evaluate BlockHammer's effect on a typical DDR4-based memory subsystem's performance and energy consumption as compared to six prior RowHammer mitigation mechanisms [73, 84, 113, 132, 137, 161]. We use Ramulator [77, 125] for performance evaluation and DRAMPower [18] to estimate DRAM energy consumption. We open-source our infrastructure, which implements both BlockHammer and six state-of-the-art RowHammer mitigation mechanisms [124]. Table 5 shows our system configuration.

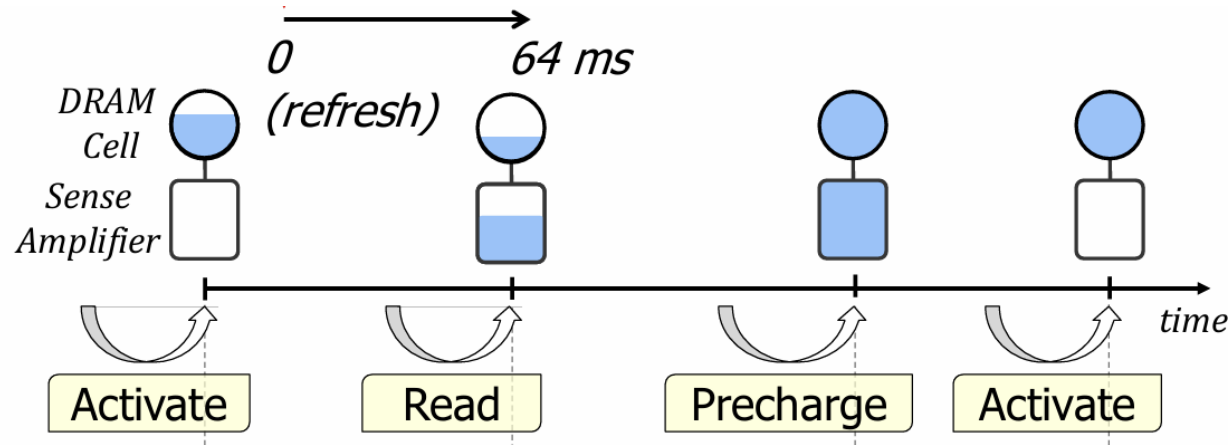
Processor	3.2 GHz, {1,8} core, 4-wide issue, 128-entry instr. window
Last-Level Cache	64-byte cache line, 8-way set-associative, 16 MB
Memory Controller	64-entry each read and write request queues; Scheduling policy: FR-FCFS [122, 164]; Address mapping: MOP [60]
Main Memory	DDR4, 1 channel, 1 rank, 4 bank groups, 4 banks/bank group, 64K rows/bank

Table 5: Simulated system configuration.



# DRAM Operations & States

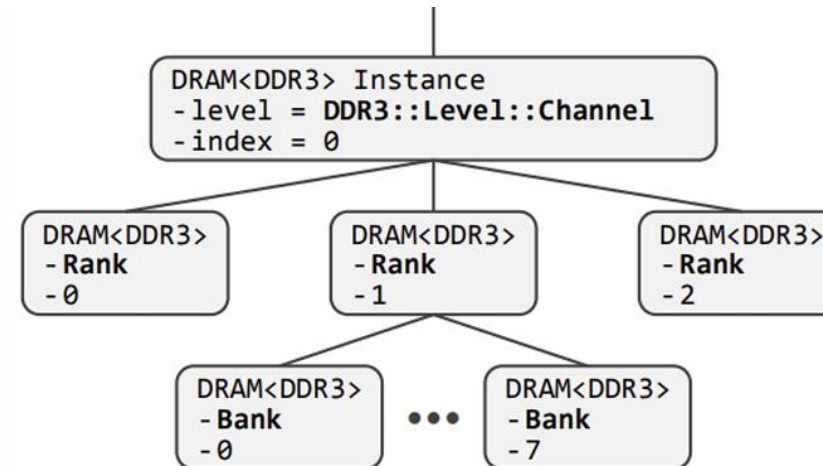
## □ DRAM Operations & States



- Main DRAM states
  - Activate
  - Read/Write
  - Precharge

```
1 // DRAM.h
2 template <typename T>
3 class DRAM {
4     DRAM<T>* parent;
5     vector<DRAM<T>*>
6         children;
7     T::Level level;
8     int index;
9     // more code...
10 };
```

```
1 // DDR3.h/cpp
2 class DDR3 {
3     enum class Level {
4         Channel, Rank, Bank,
5         Row, Column, MAX
6     };
7
8     // more code...
9
10 };
```



# src files $\Leftrightarrow$ DRAM Operation

## □ Simulation Configuration

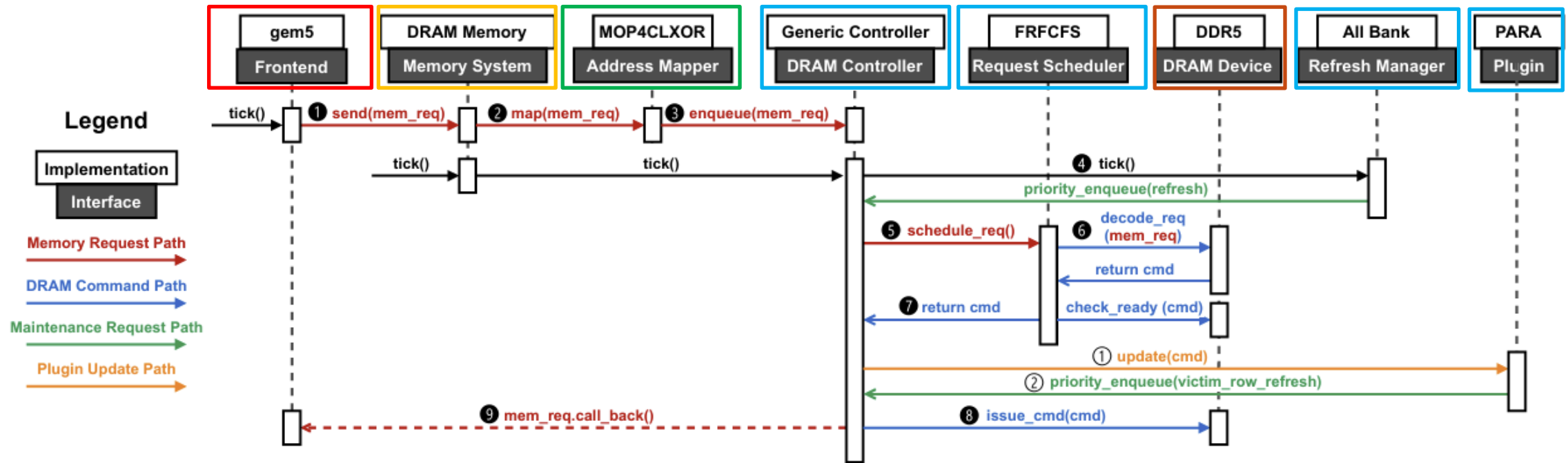
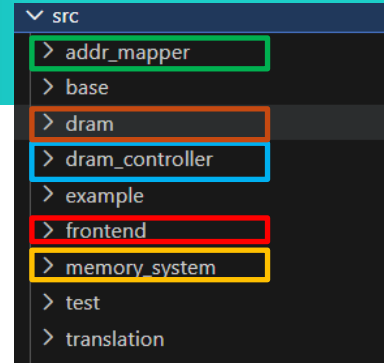


Fig. 1: High-level software architecture of Ramulator 2.0 using an example DDR5 system configuration

# src files <=> DRAM Operation

## □ Simulation Configuration

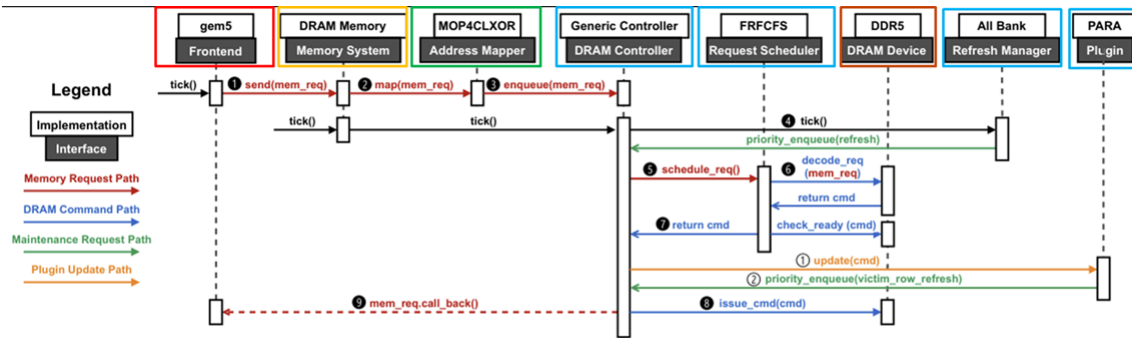
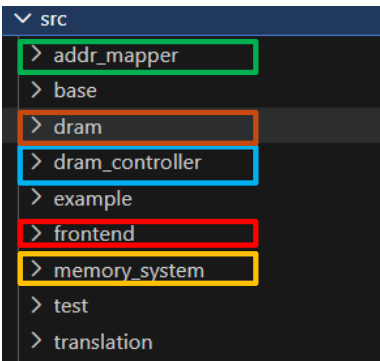


Fig. 1: High-level software architecture of Ramulator 2.0 using an example DDR5 system configuration



1. **Requests are sent:** Front-end(trace file)에서 Memory Request를 보냄
2. **Memory Addresses are Mapped:** Address Mapper가 Request Address를 DRAM 구조에 맞게 변환
3. **Enqueue:** DRAM Ctrlr의 Buffer에 Request를 넣음
4. **DRAM Ctrlr - Ticking Refresh Manager:** Ctrlr가 Refresh Manager를 호출해 high-priority maintenance request(ex. Refresh)을 추가
5. **DRAM Ctrlr - Request Scheduling:** Request Scheduler에게 최적의 Request를 선택하라고 요청
6. **DRAM Device가 Request 확인:** Scheduler가 DRAM Device Model을 참조해 적합한 Command를 Decode
7. **Issue Command:** DRAM Ctrlr가 DRAM Command를 보냄
8. **Updates the behavior and timing information:** DRAM Command Issue시 State & Timing이 Update
9. **Notify the frontend:** Memory Request가 끝나면 callback으로 frontend에 알림

# main function

## main.cpp

```
13 ~ int main(int argc, char* argv[]) {
14     // Parse command line arguments
15     argparse::ArgumentParser program("Ramulator", "2.0");
16     program.add_argument("-c", "--config").metavar("\ndumped YAML configuration")
17     .help("String dump of the yaml configuration.");
18     program.add_argument("-f", "--config_file").metavar("path-to-configuration-file")
19     .help("Path to a YAML configuration file.");
20     program.add_argument("-p", "--param").metavar("KEY=VALUE")
21     .append()
22     .help("Specify parameter to override in the configuration file. Repeat this option to change multiple parameters.");
```

:

```
88     // Connect the frontend and the memory system together,
89     // this recursively calls the "setup" function in all instantiated components
90     // so that they can get each other's parameters (if needed) after their initialization
91     frontend->connect_memory_system(memory_system);
92     memory_system->connect_frontend(frontend);
93
94     // Get the relative clock ratio between the frontend and memory system
95     int frontend_tick = frontend->get_clock_ratio();
96     int mem_tick = memory_system->get_clock_ratio();
97
98     int tick_mult = frontend_tick * mem_tick;
99
100     for (uint64_t i = 0;; i++) {
101         if (((i % tick_mult) % mem_tick) == 0) {
102             frontend->tick();
103         }
104
105         if (frontend->is_finished()) {
106             break;
107         }
108
109         if ((i % tick_mult) % frontend_tick == 0) {
110             memory_system->tick();
111         }
112     }
113
114     // Finalize the simulation. Recursively print all statistics from all components
115     frontend->finalize();
116     memory_system->finalize();
117
118     return 0;
119 }
```

## main.cpp

### 1. Argument 받는 부분

- Options

1. -c: command line dump
2. -f: YAML document
3. -p: overriding parameters in a YAML document

### 2. Long for loop를 통한 tick() 기반 simul

1. frontend(core)가 발행한 예상 instructions들을 모두 처리시 is\_finished()가 true가 됨



# yaml file

## □ example\_config.yaml

```
1  Frontend:
2    impl: SimpleO3
3    clock_ratio: 8
4    num_expected_insts: 500000
5  traces:
6    - example_inst.trace
7
8  Translation:
9    impl: RandomTranslation
10   max_addr: 2147483648
11
12
13 MemorySystem:
14   impl: GenericDRAM
15   clock_ratio: 3
16
17   DRAM:
18     impl: DDR4
19     org:
20       preset: DDR4_8Gb_x8
21       channel: 1
22     rank: 2
23     timing:
24       preset: DDR4_2400R
25
26   Controller:
27     impl: Generic
28     Scheduler:
29       impl: FRFCFS
30     RefreshManager:
31       impl: AllBank
32     RowPolicy:
33       impl: ClosedRowPolicy
34       cap: 4
35     plugins:
36
37   AddrMapper:
38     impl: RoBaRaCoCh
```

## 1. Frontend Interface(IFrontEnd) 부분

- trace file에서 Instruction 읽고, Memory Request 생성
- impl: SimpleO3  
⇒ Simple Out-of-Order (O3) CPU
- clock\_ratio: 8  
⇒ global CLK 대비 Frontend CLK 속도
- num\_expected\_insts: 500000  
⇒ Simulation이 해당 instruction 수에 도달 시 종료
- traces  
⇒ Instruction trace file(include memory access Inst)
- impl: RandomTranslation  
⇒ Physical Memory ↔ Virtual Memory 변환  
⇒ System의 Page Table 등을 간단히 Modeling
- max\_addr: 2147483648  
⇒ Translation 시 address overflow 방지



# yaml file

## □ example\_config.yaml

```
1  Frontend:
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31      impl: AllBank
32    RowPolicy:
33      impl: ClosedRowPolicy
34      cap: 4
35    plugins:
36
37  AddrMapper:
38    impl: RoBaRaCoCh
```

## 2. MemorySystem Interface 부분

- Frontend의 Request를 받아 DRAM Ctrlr을 통해 처리
- Latency, en/dequeue, Timing Constraints 처리
- impl: GenericDRAM  
⇒ 기본 DRAM 기반 System, Ctrlr와 DRAM을 통합
- clock ratio: 3  
⇒ global CLK 대비 MemorySystem CLK 속도  
⇒ 현재: DRAM이 CPU보다 느린 System (= 3:8)

# yaml file

## example\_config.yaml 2. MemorySystem Interface 부분

```
1 Frontend: 13 MemorySystem:
2   impl: Simple03 14   impl: GenericDRAM
3   clock_ratio: 8 15   clock_ratio: 3
4   num_expected_insts: 500000 16
5   traces: 17 DRAM:
6     - example_inst.trace 18   impl: DDR4
7 8 Translation: 19   org:
9   impl: RandomTranslation 20     preset: DDR4_8Gb_x8
10   max_addr: 2147483648 21     channel: 1
11 22     rank: 2
23     timing:
24     preset: DDR4_2400R
25
26 Controller:
27   impl: Generic
28   Scheduler:
29     impl: FRFCFS
30   RefreshManager:
31     impl: AllBank
32   RowPolicy:
33     impl: ClosedRowPolicy
34     cap: 4
35   plugins:
36
37 AddrMapper:
38   impl: RoBaRaCoCh
```

```
26 inline static const std::map<std::string, std::vector<int>> timing_presets = {
27   // name rate nBL nCL nRCD nRP nRAS nRC nWR nRTP nCWL nCCDS nCCDL nRRDS nRRDL nWTRS nW
28   {"DDR4_1600J", {1600, 4, 10, 10, 10, 28, 38, 12, 6, 9, 4, 5, -1, -1, 2,
29   {"DDR4_1600K", {1600, 4, 11, 11, 11, 28, 39, 12, 6, 9, 4, 5, -1, -1, 2,
30   {"DDR4_1600L", {1600, 4, 12, 12, 12, 28, 40, 12, 6, 9, 4, 5, -1, -1, 2,
31   {"DDR4_1866L", {1866, 4, 12, 12, 12, 32, 44, 14, 7, 10, 4, 5, -1, -1, 3,
32   {"DDR4_1866M", {1866, 4, 13, 13, 13, 32, 45, 14, 7, 10, 4, 5, -1, -1, 3,
33   {"DDR4_1866N", {1866, 4, 14, 14, 14, 32, 46, 14, 7, 10, 4, 5, -1, -1, 3,
34   {"DDR4_2133N", {2133, 4, 14, 14, 14, 36, 50, 16, 8, 11, 4, 6, -1, -1, 3,
35   {"DDR4_2133P", {2133, 4, 15, 15, 15, 36, 51, 16, 8, 11, 4, 6, -1, -1, 3,
36   {"DDR4_2133R", {2133, 4, 16, 16, 16, 36, 52, 16, 8, 11, 4, 6, -1, -1, 3,
37   {"DDR4_2400P", {2400, 4, 15, 15, 15, 39, 54, 18, 9, 12, 4, 6, -1, -1, 3,
38   {"DDR4_2400R", {2400, 4, 16, 16, 16, 39, 55, 18, 9, 12, 4, 6, -1, -1, 3,
39   {"DDR4_2400U", {2400, 4, 17, 17, 17, 39, 56, 18, 9, 12, 4, 6, -1, -1, 3,
40   {"DDR4_2400T", {2400, 4, 18, 18, 18, 39, 57, 18, 9, 12, 4, 6, -1, -1, 3,
```

### • DRAM Section

### • impl: DDR4

⇒ tick() 시 Timing Check / Command Issue 실행.

### • org: DDR4 8Gb x8

⇒ 현재 **DRAM preset** - 8Gb 용량, x8bit data bus

⇒ **Channel/Rank 설정** 시 기본 Preset설정을 Override 함

### • timing: DDR4 2400R

⇒ **Timing preset** - nRCD등의 Timing Constraint 정의

⇒ 이를 이용해 tick() 시 Latency 계산

```
class DDR4 : public IDRAM, public Implementation {
    RAMULATOR_REGISTER_IMPLEMENTATION(IDRAM, DDR4, "DDR4", "DDR4 Device Model")

public:
    inline static const std::map<std::string, Organization> org_presets = {
        // name density DQ Ch Ra Bg Ba Ro Co
        {"DDR4_2Gb_x4", {2<<10, 4, {1, 1, 4, 4, 1<<15, 1<<10}}},
        {"DDR4_2Gb_x8", {2<<10, 8, {1, 1, 4, 4, 1<<14, 1<<10}}},
        {"DDR4_2Gb_x16", {2<<10, 16, {1, 1, 2, 4, 1<<14, 1<<10}}},
        {"DDR4_4Gb_x4", {4<<10, 4, {1, 1, 4, 4, 1<<16, 1<<10}}},
        {"DDR4_4Gb_x8", {4<<10, 8, {1, 1, 4, 4, 1<<15, 1<<10}}},
        {"DDR4_4Gb_x16", {4<<10, 16, {1, 1, 2, 4, 1<<15, 1<<10}}},
        {"DDR4_8Gb_x4", {8<<10, 4, {1, 1, 4, 4, 1<<17, 1<<10}}},
        {"DDR4_8Gb_x8", {8<<10, 8, {1, 1, 4, 4, 1<<16, 1<<10}}},
        {"DDR4_8Gb_x16", {8<<10, 16, {1, 1, 2, 4, 1<<16, 1<<10}}},
        {"DDR4_16Gb_x4", {16<<10, 4, {1, 1, 4, 4, 1<<18, 1<<10}}},
        {"DDR4_16Gb_x8", {16<<10, 8, {1, 1, 4, 4, 1<<17, 1<<10}}},
        {"DDR4_16Gb_x16", {16<<10, 16, {1, 1, 2, 4, 1<<17, 1<<10}}},
    };
};
```

```
class DDR4 : public IDRAM, public Implementation {
    void tick() override {

    void init() override {
        RAMULATOR_DECLARE_SPECS();
        set_organization();
        set_timing_vals();

        set_actions();
        set_preqs();
        set_rowwhits();
        set_rowopens();
        set_powers();

        create_nodes();
    };
};
```

# yaml file

## □ example\_config.yaml

```
1 Frontend:
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31       impl: AllBank
32     RowPolicy:
33       impl: ClosedRowPolicy
34       cap: 4
35     plugins:
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```

## 2. MemorySystem Interface 부분

- **Controller Section**
- impl: Generic
  - ⇒ Generic - 기본 Ctrlr
  - ⇒ Request Queue/Scheduling 등 관리
- Scheduler - impl: FRFCFS
  - ⇒ FRFCFS - First-Ready First-Come-First-Serve
  - ⇒ 준비된 Request를 Queue에서 꺼내 우선 처리
- RefreshManager - impl: AllBank
  - ⇒ AllBank - 모든 Bank simultaneous Refresh
- RowPolicy - impl: ClosedRowPolicy
  - ⇒ ClosedRowPolicy - 사용 후 Row 즉시 닫음(Precharge)
  - ⇒ cap:4 - 열려있는 Row 최대 수 제한
- plugins
  - ⇒ 현재 Ramulator에서는 Row Hammering 완화 기법을 plugin으로 제공해줌

# yaml file

## □ example\_config.yaml

```
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```

## 2. MemorySystem Interface 부분

- AddrMapper Section

- impl: RoBaRaCoCh

⇒ Row-Bank-Rank-Column-Channel Mapping Scheme

⇒ Requested Address 변환(Physical → DRAM Vector)

[Physical → DRAM Vector Example]

⇒ Physical Address: 0x12345678

⇒ DRAM Vector:

[Channel:0, Rank:1, Bank:2, Row:128, Column:512]

## □ example\_inst.trace & trace.cpp & core.cpp

- simpleO3 CPU model기준
  - Ramulator는 “Memory” Simulator
  - Memory 명령어만 취급하기에, 3가지로만 Instruction을 분리한다.

1. Not Memory Operation

2. Load

3. Store

- 따라서, simpleO3 기반 trace파일:
  - 1st column은 Not Memory Operation Cycle 수 (or ticks 수)
  - 2nd column은 load operation address
  - 3rd column은 store operation address

example_inst.trace		
1	3	20734016
2	1	20846400
3	6	20734208
4	1	20846400
5	8	20841280 20841280
6	0	20734144
7	2	20918976 20734016
8	1	20846400
9		

- 각 line은 load(& store)동작을 나타낸다.
- 1<sup>st</sup> line : 3cycle동안 stall → load
- 5<sup>th</sup> line : 8cycle동안 stall → load → store

# trace file

## □ example\_inst.trace & trace.cpp & core.cpp

The screenshot displays three code files: `trace.cpp`, `core.cpp`, and `example_inst.trace`. In `trace.cpp`, the `while` loop processes tokens from the trace file. Tokens are split by spaces, and the first token is converted to an integer representing a bubble count. The second token is converted to a long integer representing a load address. The third token is converted to a long integer representing a store address. These values are then used to push back the trace. In `core.cpp`, the `SimpleO3Core::tick()` function processes the trace. It first issues non-memory instructions, then sends loads to the LLC, and finally sends writebacks to the LLC. The `example_inst.trace` file shows a sequence of instructions with their bubble counts, load addresses, and store addresses. A red box highlights the first instruction (bubble 3, load 20734016, store 20841280). A green box highlights the second instruction (bubble 1, load 20846400, store 20734016). A blue box highlights the third instruction (bubble 6, load 20734208, store 20841280). A red arrow points from the first instruction in the trace file to the `SimpleO3Core::tick()` function in `core.cpp`. A green arrow points from the second instruction in the trace file to the `SimpleO3Core::tick()` function in `core.cpp`. A blue arrow points from the third instruction in the trace file to the `SimpleO3Core::tick()` function in `core.cpp`.

```
src > frontend > impl > processor > simpleO3 > trace.cpp > ...
std::string line;
while (std::getline(trace_file, line)) {
    std::vector<std::string> tokens;
    tokenize(tokens, line, " ");

    int num_tokens = tokens.size();
    if (num_tokens != 2 & num_tokens != 3) {
        throw ConfigurationError("Trace {} format invalid!", file_path_str);
    }

    int bubble_count = std::stoi(tokens[0]);
    Addr_t load_addr = std::stoll(tokens[1]);

    bool has_store = num_tokens == 3 ? true : false;
    if (has_store) {
        Addr_t store_addr = std::stoll(tokens[2]);
        m_trace.push_back({bubble_count, load_addr, store_addr});
    } else {
        m_trace.push_back({bubble_count, load_addr, -1});
    }
}

trace_file.close();
m_trace.length = m_trace.size();

src > frontend > impl > processor > simpleO3 > core.cpp > {} Ramulator > tick()
void SimpleO3Core::tick() {
    m_clk++;

    s_insts_retired += m_window.retire();
    if (!reached_expected_num_insts) {
        if (s_insts_retired >= m_num_expected_insts) {
            reached_expected_num_insts = true;
            s_cycles_recorded = m_clk;
        }
    }

    // First, issue the non-memory instructions.
    int num_inserted_insts = 0;
    while (m_num_bubbles > 0) {
        if (num_inserted_insts == m_window.m_ipc) {
            return;
        }
        if (m_window.is_full()) {
            return;
        }
        m_window.insert(true, -1);
        num_inserted_insts++;
        m_num_bubbles--;
    }

    // Second, try to send the load to the LLC
    if (m_load_addr != -1) {
        if (num_inserted_insts == m_window.m_ipc) {
            return;
        }
        if (m_window.is_full()) {
            return;
        }
    }

    Request load_request(m_load_addr, Request::Type::Read, m_id, m_callback);
    if (!m_translation->translate(load_request)) {
        return;
    }

    if (m_llc->send(load_request)) {
        m_window.insert(false, load_request.addr);
        m_load_addr = -1;
        if (m_writeback_addr != -1) {
            // If there is still writeback, return without getting the next trace line
            // The write back will be issued in the next cycle
            // TODO: Should we allow both load and writeback to issue at the same cycle?
            return;
        }
    } else {
        return;
    }

    // Third, try to send the writeback to the LLC
    if (m_writeback_addr != -1) {
        Request writeback_request(m_writeback_addr, Request::Type::Write, m_id, m_callback);
        if (!m_translation->translate(writeback_request)) {
            return;
        }
        if (m_llc->send(writeback_request)) {
            return;
        }
    }

    auto inst = m_trace.get_next_inst();
    m_num_bubbles = inst.bubble_count;
    m_load_addr = inst.load_addr;
    m_writeback_addr = inst.store_addr;
}
```

example\_inst.trace

Line	Bubble	Load Address	Store Address
1	3	20734016	
2	1	20846400	
3	6	20734208	
4	1	20846400	
5	8	20841280	20841280
6	0	20734144	
7	2	20918976	20734016
8	1	20846400	
9			

- Trace file을 arg로 받아, Frontend (ex: simpleO3.cpp)에서 처리되어 메모리 접근 request를 생성
- simpleO3 CPU model기준: trace의 token[0]: bubble / token[1]: load address / token[2]: store address

# trace file result

## ❑ ./ramulator2 -f ./example\_config.yaml

```
root@947e591ed45c:/workspace# ./ramulator2 -f ./example_config.yaml
Frontend:
  impl: SimpleO3
  memory_access_cycles_recorded_core_0: 61
  cycles_recorded_core_0: 216815
  llc_mshr_unavailable: 0
  llc_read_misses: 37
  llc_read_access: 133336
  llc_write_misses: 8
  llc_write_access: 33334
  llc_eviction: 0
  num_expected_insts: 500000
  Translation:
    impl: RandomTranslation

MemorySystem:
  impl: GenericDRAM
  total_num_other_requests: 0
  total_num_write_requests: 0
  total_num_read_requests: 6
  memory_system_cycles: 81306
  DRAM:
    impl: DDR4
  AddrMapper:
    impl: RoBaRaCoCh

Controller:
  impl: Generic
  id: channel 0
  avg_read_latency_0: 46.5
  read_queue_len_avg_0: 0.00232455181
  write_queue_len_0: 0
  queue_len_0: 245
  num_other_reqs_0: 0
  num_write_reqs_0: 0
  read_latency_0: 279
  priority_queue_len_avg_0: 0.000688756059
  row_hits_0: 2
  priority_queue_len_0: 56
  row_misses_0: 4
  row_conflicts_0: 0
  read_row_misses_0: 4
  queue_len_avg_0: 0.00301330769
  read_row_conflicts_core_0: 0
  read_row_hits_0: 2
  write_queue_len_avg_0: 0
  read_row_conflicts_0: 0
  write_row_misses_0: 0
  write_row_conflicts_0: 0
  read_queue_len_0: 189
  write_row_hits_0: 0
  read_row_hits_core_0: 2
  read_row_misses_core_0: 4
  num_read_reqs_0: 6
  Scheduler:
    impl: FRFCFS
  RefreshManager:
    impl: AllBank

RowPolicy:
  impl: ClosedRowPolicy
  num_close_reqs: 0
```

- 기본 yaml file setting을 이용.  
=> trace file의 Inst 실행결과가 분석됨
- 1. Frontend:  
CPU ↔ Memory<Cache> 접근 분석
- 2. MemorySystem<DRAM> 관점 분석
  - Request 수
  - Cycle 수
  - Latency
  - Queue 길이
  - Row Hit / Miss



# trace file result analysis

## □ ./ramulator2 -f ./example\_config.yaml

- 분석 시 사용한 trace file은 오른쪽의 example\_inst이다.
- **예상 simulation 결과**는 다음과 같아야 한다 → *이를 실제 Results와 비교한다.*
- 8번의 load (address: 0x20846400가 3번 나온다)
- 2번의 write (1<sup>st</sup> Inst의 address, 5<sup>th</sup> Inst의 address에 Write [이미 load된 address 접근 → cache hit])

### ▪ DRAM 관점

- 이를 통해서 **6번의 DRAM Read Access**를 요구하는 것을 예측 가능
  - 나머지 중복 load들은 cache line에 hit되어 DRAM Access X
- Write의 경우 이미 Load된 address에 write하므로 cache이용 : **0번의 DRAM Write Access**

### ▪ SRAM(cache) 관점

- 적은 memory 주소만 반복 접근 하므로 cache eviction X (이전에 load된 addr들은 다 hit 됨)
- **Initial Read Misses: 6번 / Initial Write Misses: 0번**

example_inst.trace			
1	3	20734016	
2	1	20846400	
3	6	20734208	
4	1	20846400	
5	8	20841280	20841280
6	0	20734144	
7	2	20918976	20734016
8	1	20846400	
9			

# trace file result analysis

## ❑ ./ramulator2 -f ./example\_config.yaml

```
root@947e591ed45c:/workspace# ./ramulator2 -f ./example_config.yaml
Frontend:
  impl: SimpleO3
  memory_access_cycles_recorded_core_0: 61
  cycles_recorded_core_0: 216815
  llc_mshr_unavailable: 0
  llc_read_misses: 37
  llc_read_access: 133336
  llc_write_misses: 8
  llc_write_access: 33334
  llc_eviction: 0
  num_expected_insts: 500000
  Translation:
    impl: RandomTranslation

MemorySystem:
  impl: GenericDRAM
  total_num_other_requests: 0
  total_num_write_requests: 0
  total_num_read_requests: 6
  memory_system_cycles: 81306
  DRAM:
    impl: DDR4
  AddrMapper:
    impl: RoBARacoch

Controller:
  impl: Generic
  id: Channel 0
  avg_read_latency_0: 46.5
  read_queue_len_avg_0: 0.00232455181
  write_queue_len_0: 0
  queue_len_0: 245
  num_other_reqs_0: 0
  num_write_reqs_0: 0
  read_latency_0: 279
  priority_queue_len_avg_0: 0.000688756059
  row_hits_0: 2
  priority_queue_len_0: 56
  row_misses_0: 4
  row_conflicts_0: 0
  read_row_misses_0: 4
  queue_len_avg_0: 0.00301330769
  read_row_conflicts_core_0: 0
  read_row_hits_0: 2
  write_queue_len_avg_0: 0
  read_row_conflicts_0: 0
  write_row_misses_0: 0
  write_row_conflicts_0: 0
  read_queue_len_0: 189
  write_row_hits_0: 0
  read_row_hits_core_0: 2
  read_row_misses_core_0: 4
  num_read_reqs_0: 6
  Scheduler:
    impl: FRFCFS
  RefreshManager:
    impl: AllBank

RowPolicy:
  impl: ClosedRowPolicy
  num_close_reqs: 0
```

```
52: const SimpleO3Core::Trace::Inst& SimpleO3Core::Trace::get_next_inst(
53:     const Inst& inst = m_trace[m_curr_trace_idx];
54:     m_curr_trace_idx = (m_curr_trace_idx + 1) % m_trace_length;
55:     return inst;
56: }
```

짧은 trace지만 반복 실행 되기에 많은 access 발생

### 1. Frontend

- Last Level Cache(LLC) 분석
  - Core 0의 Memory(Cache 포함) 접근 Cycles
  - LLC의 Read Request Cache Miss 수
  - LLC의 Read Access 수
  - Miss Status Handling Register(MSHR) unavailable cnt  
→ MSHR가 가득 차서 Request 거부된 횟수
  - LLC의 Write Request Cache Miss 수
  - LLC의 Write Access 수
  - LLC의 Eviction 수  
→ Cache Line 교체된 수 (DRAM Write 유발)

*이전의 예상치와 다르게 LLC의 miss수가 훨씬 많음*

- dependency check, reorder buffer 등 때문일 수 있음
- LLC operation 코드 분석이 필요 → OoO CPU의 특성 때문임
- `bool SimpleO3LLC::send(Request req)`

# trace file result analysis

## ❑ ./ramulator2 -f ./example\_config.yaml

```
root@947e591ed45c:/workspace# ./ramulator2 -f ./example_config.yaml
Frontend:
  impl: SimpleO3
  memory_access_cycles_recorded_core_0: 61
  cycles_recorded_core_0: 216815
  llc_mshr_unavailable: 0
  llc_read_misses: 37
  llc_read_access: 133336
  llc_write_misses: 8
  llc_write_access: 33334
  llc_eviction: 0
  num_expected_insts: 500000
  Translation:
    impl: RandomTranslation

MemorySystem:
  impl: GenericDRAM
  total_num_other_requests: 0
  total_num_write_requests: 0
  total_num_read_requests: 6
  memory_system_cycles: 81306
  DRAM:
    impl: DDR4
    AddrMapper:
      impl: RoBaracoch

Controller:
  impl: Generic
  id: Channel 0
  avg_read_latency_0: 46.5
  read_queue_len_avg_0: 0.00232455181
  write_queue_len_0: 0
  queue_len_0: 245
  num_other_reqs_0: 0
  num_write_reqs_0: 0
  read_latency_0: 279
  priority_queue_len_avg_0: 0.000688756059
  row_hits_0: 2
  priority_queue_len_0: 56
  row_misses_0: 4
  row_conflicts_0: 0
  read_row_misses_0: 4
  queue_len_avg_0: 0.00301330769
  read_row_conflicts_core_0: 0
  read_row_hits_0: 2
  write_queue_len_avg_0: 0
  read_row_conflicts_0: 0
  write_row_misses_0: 0
  write_row_conflicts_0: 0
  read_queue_len_0: 189
  write_row_hits_0: 0
  read_row_hits_core_0: 2
  read_row_misses_core_0: 4
  num_read_reqs_0: 6
  Scheduler:
    impl: FRFCFS
  RefreshManager:
    impl: AllBank

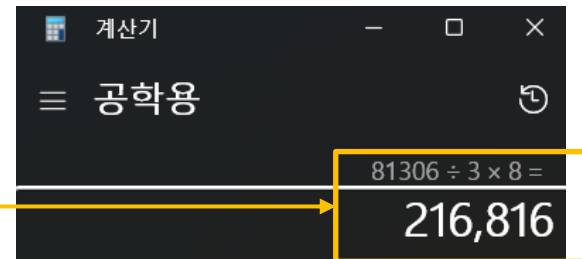
RowPolicy:
  impl: ClosedRowPolicy
  num_close_reqs: 0
```

## 2. MemorySystem Interface

### • DRAM 분석

- Read/Write 외 other requests 수 *이전의 예상치와 일치*
- write requests 수 → 현재 0 (cache eviction 없었기에)
- read requests 수 → 현재 6 (initial cache line load)
- Memory system cycles  
→ Memory CLK과 Proc CLK이 다르기에 (= 3:8)

**DRAM CLK : Processor CLK = 3:8**



```
108 inline static constexpr ImplDef m_requests = {
109     "read", "write", "all-bank-refresh", "open-row", "close-row"
110 };
111
112 inline static const ImplLUT m_request_translations = LUT (
113     m_requests, m_commands, {
114         {"read", "RD"}, {"write", "WR"}, {"all-bank-refresh", "REFab"},
115         {"open-row", "ACT"}, {"close-row", "PRE"}
116     }
117 );
```

Other requests: read/write를 제외한 open-row와 같은 request등을 의미

# trace file result analysis

## ❑ ./ramulator2 -f ./example\_config.yaml

```
root@947e591ed45c:/workspace# ./ramulator2 -f ./example_config.yaml
Frontend:
  impl: SimpleO3
  memory_access_cycles_recorded_core_0: 61
  cycles_recorded_core_0: 216815
  llc_mshr_unavailable: 0
  llc_read_misses: 37
  llc_read_access: 133336
  llc_write_misses: 8
  llc_write_access: 33334
  llc_eviction: 0
  num_expected_insts: 500000
  Translation:
    impl: RandomTranslation

MemorySystem:
  impl: GenericDRAM
  total_num_other_requests: 0
  total_num_write_requests: 0
  total_num_read_requests: 6
  memory_system_cycles: 81306
DRAM:
  impl: DDR4
AddrMapper:
  impl: RoBarAcocH

Controller:
  impl: Generic
  id: Channel 0
  avg_read_latency_0: 46.5
  read_queue_len_avg_0: 0.00232455181
  write_queue_len_0: 0
  queue_len_0: 245
  num_other_reqs_0: 0
  num_write_reqs_0: 0
  read_latency_0: 279
  priority_queue_len_avg_0: 0.000688756059
  row_hits_0: 2
  priority_queue_len_0: 56
  row_misses_0: 4
  row_conflicts_0: 0
  read_row_misses_0: 4
  queue_len_avg_0: 0.00301330769
  read_row_conflicts_core_0: 0
  read_row_hits_0: 2
  write_queue_len_avg_0: 0
  read_row_conflicts_0: 0
  write_row_misses_0: 0
  write_row_conflicts_0: 0
  read_queue_len_0: 189
  write_row_hits_0: 0
  read_row_hits_core_0: 2
  read_row_misses_core_0: 4
  num_read_reqs_0: 6
Scheduler:
  impl: FRFCFS
RefreshManager:
  impl: AllBank

RowPolicy:
  impl: ClosedRowPolicy
  num_close_reqs: 0
```

## 2. MemorySystem Interface

- DRAM Controller 분석
  - abcd

# trace file result

## □ power는 어떻게 측정하느냐면

- yaml file에 아래 drampower\_enable 옵션을 넣음
- dram.h -> DDR4.cpp 의 power함수 enable됨

```
17  ▾ DRAM:
18      impl: DDR4
19  ▾  org:
20      preset: DDR4_8Gb_x8
21      channel: 1
22      rank: 2
23  ▾  timing:
24      preset: DDR4_2400R
25      drampower_enable: true
26      # power_debug: true # option: debug log
27  ▾  voltage:
28      preset: Default # option: voltage preset
29  ▾  current:
30      preset: Default # option: current preset
31
```

```
DRAM:
  impl: DDR4
  active_cycles_rank1: 9361
  pre_background_energy_rank1: 3850.9652475000003
  total_background_energy_rank0: 4424.0171849999997
  idle_cycles_rank1: 68489
  total_cmd_energy: 3203.0342735999993
  total_cmd_energy_rank0: 1601.8769927999997
  total_energy_rank0: 6025.8941777999989
  act_background_energy_rank0: 572.15229750000003
  pre_background_energy_rank0: 3851.8648874999999
  active_cycles_rank0: 9345
  act_background_energy_rank1: 573.13190550000002
  total_energy: 12051.1486116
  idle_cycles_rank0: 68505
  total_energy_rank1: 6025.2544338000007
  total_background_energy_rank1: 4424.0971530000006
  total_background_energy: 8848.1143379999994
  total_cmd_energy_rank1: 1601.1572807999996
AddrMapper:
  impl: RoBaRaCoCh
```

# trace file result

## □ power는 어떻게 측정하느냐면

- power\_debug option enable 시
- 오른쪽과 같은 power log 확인 가능

```
17  ▾  DRAM:
18      impl: DDR4
19  ▾  org:
20      preset: DDR4_8Gb_x8
21      channel: 1
22      rank: 2
23  ▾  timing:
24      preset: DDR4_2400R
25      drampower_enable: true
26      # power_debug: true # option: debug log
27  ▾  voltage:
28      preset: Default # option: voltage preset
29  ▾  current:
30      preset: Default # option: current preset
31
```

```
root@947e591ed45c:/workspace# ./ramulator2 -f ./example_config_power_en.yaml
[Power] Rank0 -----ACT----- @ 19
[Power] Rank0 Rank is idle. idle_cycles: 19 active_start_cycle: 19 @ 19
[Power] Rank0 Bank0 Incrementing ACT counter. @ 19
[Power] Rank1 -----ACT----- @ 20
[Power] Rank1 Rank is idle. idle_cycles: 20 active_start_cycle: 20 @ 20
[Power] Rank1 Bank2 Incrementing ACT counter. @ 20
[Power] Rank0 -----ACT----- @ 23
[Power] Rank0 Bank3 Incrementing ACT counter. @ 23
[Power] Rank1 -----ACT----- @ 24
[Power] Rank1 Bank3 Incrementing ACT counter. @ 24
[Power] Rank0 Bank0 Incrementing RD counter. @ 35
[Power] Rank0 Bank3 Incrementing RD counter. @ 39
[Power] Rank0 Bank0 Incrementing RD counter. @ 43
[Power] Rank1 Bank2 Incrementing RD counter. @ 49
[Power] Rank1 Bank3 Incrementing RD counter. @ 53
[Power] Rank0 Bank0 Incrementing RD counter. @ 59
[Power] Rank0 -----PREA----- @ 9364
[Power] Rank0 Incrementing PRE counter. @ 9364
[Power] Rank0 Rank is not idle. active_cycles: 9345 idle_start_cycle: 9364 @ 9364
[Power] Rank0 -----REFab----- @ 9380
[Power] Rank0 Refresh starts. idle_cycles: 35 @ 9380
[Power] Rank1 -----PREA----- @ 9381
[Power] Rank1 Incrementing PRE counter. @ 9381
[Power] Rank1 Rank is not idle. active_cycles: 9361 idle_start_cycle: 9381 @ 9381
[Power] Rank1 -----REFab----- @ 9397
[Power] Rank1 Refresh starts. idle_cycles: 36 @ 9397
[Power] Rank0 -----REFab_end----- @ 9812
[Power] Rank0 Refresh ends. idle_start_cycle: 9812 @ 9812
[Power] Rank1 -----REFab_end----- @ 9829
[Power] Rank1 Refresh ends. idle_start_cycle: 9829 @ 9829
[Power] Rank0 -----REFab----- @ 18728
[Power] Rank0 Refresh starts. idle_cycles: 8951 @ 18728
[Power] Rank1 -----REFab----- @ 18729
[Power] Rank1 Refresh starts. idle_cycles: 8936 @ 18729
[Power] Rank0 -----REFab_end----- @ 19160
[Power] Rank0 Refresh ends. idle_start_cycle: 19160 @ 19160
[Power] Rank1 -----REFab_end----- @ 19161
[Power] Rank1 Refresh ends. idle_start_cycle: 19161 @ 19161
[Power] Rank0 -----REFab----- @ 28092
```

# yaml/trace file for **direct Read/Write Request**

## ❑ example\_config\_dram.yaml & example\_dram.trace & readwrite\_trace.cpp

```
! example_config_dram.yaml
1 Frontend:
2   impl: ReadWriteTrace
3   path: example_dram.trace # trace 파일 경로 (상대/절대)
4   clock_ratio: 1
5
6   Translation:
7     impl: RandomTranslation
8     max_addr: 2147483648
9
10
11 MemorySystem:
12   impl: GenericDRAM
13   clock_ratio: 1
14
15   DRAM:
16     impl: DDR4
17     org:
18       preset: DDR4_8Gb_x8
```

```
≡ example_dram.trace
1 R 0,0,0,0,0
2 W 0,0,0,0,128
3 R 0,0,0,1,0
4 W 0,0,0,1,256
5 R 0,0,1,0,0
6 W 0,0,1,0,512
7 R 0,0,0,1024,0
8 W 0,0,0,1024,0
9
```

```
39 void tick() override {
40   const Trace& t = m_trace[m_curr_trace_idx];
41   m_memory_system->send({t.addr_vec, t.is_write ? Request::Type::Write : Request::Type::Read});
42   // m_curr_trace_idx = (m_curr_trace_idx + 1) % m_trace_length;
43   m_curr_trace_idx++; // revised by MinsungKim yonsei ISL lab - 2026.01.09.
44 }
```

- Trace file을 arg로 받아, Frontend (ex: readwrite\_trace.cpp)에서 처리되어 메모리 접근 request를 생성
- **Read-Write Trace:** This frontend expects a trace file containing read (R) or store (W) requests with a *DRAM address vector* (i.e., which channel, rank, bank, row, column). Each line in the trace has the following format: **R/W** **<comma-separated addr vector>**:



# yaml/trace file for direct Read/Write Request

## □ Simulation Result

```
Frontend:
  impl: ReadWriteTrace

MemorySystem:
  impl: GenericDRAM
  total_num_other_requests: 0
  total_num_write_requests: 4
  total_num_read_requests: 4
  memory_system_cycles: 7
  DRAM:
    impl: DDR4
  AddrMapper:
    impl: RoBaRaCoCh

Controller:
  impl: Generic
  id: Channel 0
  avg_read_latency_0: 4.5
  read_queue_len_avg_0: 0.571428597
  write_queue_len_0: 12
  queue_len_0: 16
  num_other_reqs_0: 0
  num_write_reqs_0: 4
  read_latency_0: 18
  priority_queue_len_avg_0: 0
  row_hits_0: 0
  priority_queue_len_0: 0
  row_misses_0: 1
  row_conflicts_0: 0
  read_row_misses_0: 1
  queue_len_avg_0: 2.28571439
  read_row_conflicts_core_0: 0
  read_row_hits_0: 0
  write_queue_len_avg_0: 1.71428573
  read_row_conflicts_0: 0
  write_row_misses_0: 0
  write_row_conflicts_0: 0
  read_queue_len_0: 4
  write_row_hits_0: 0
  read_row_hits_core_0: 0
  read_row_misses_core_0: 0
  num_read_reqs_0: 4
  Scheduler:
    impl: FRFCFS
  RefreshManager:
    impl: AllBank

RowPolicy:
  impl: ClosedRowPolicy
  num_close_reqs: 0
```

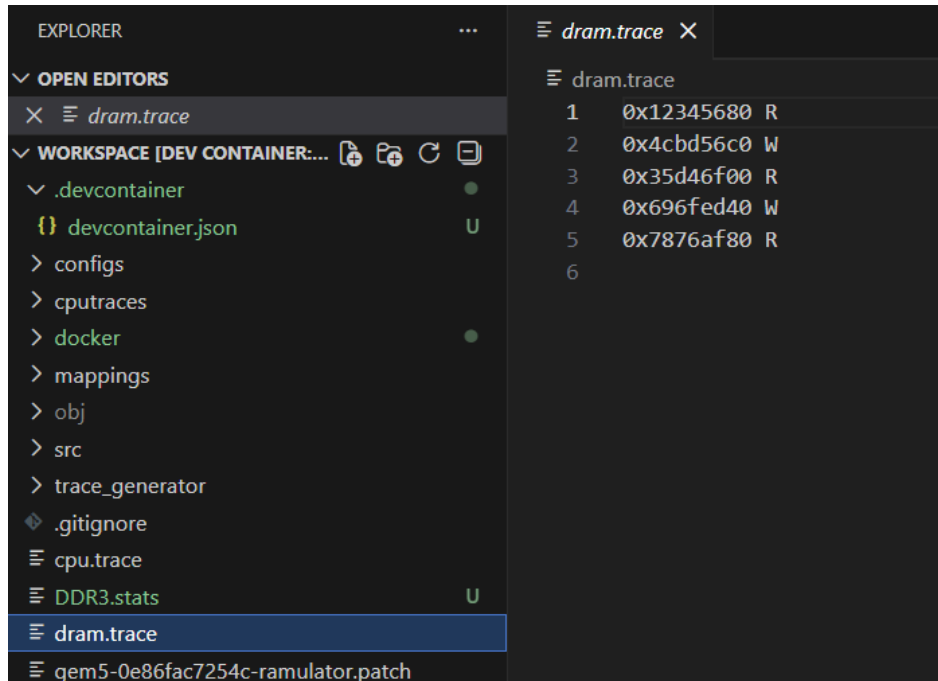
결과가 나오긴하는데 말이 안되는 수치를 가짐

Currently, Ramulator 2 implements the following three frontends:

- **Memory Trace Frontend:** This frontend reads an input trace file that contains *main memory requests* of an application. It *sequentially* sends these requests to the memory system. **This mode does not model any system in sufficient detail to perform timing simulations.** Because of that, the Memory Trace Frontend is **better suited for testing/debugging the functionality** of newly added features. In the main tasks of this assignment, we will *not* use this frontend. Still, feel free to use this mode to test and debug your newly implemented features. Currently, Ramulator 2 implements the following two memory trace frontends:
- **Ramulator2.0은 Read/Write Request의 functionality만 검증**
- **Cycle Accurate하지 않음**

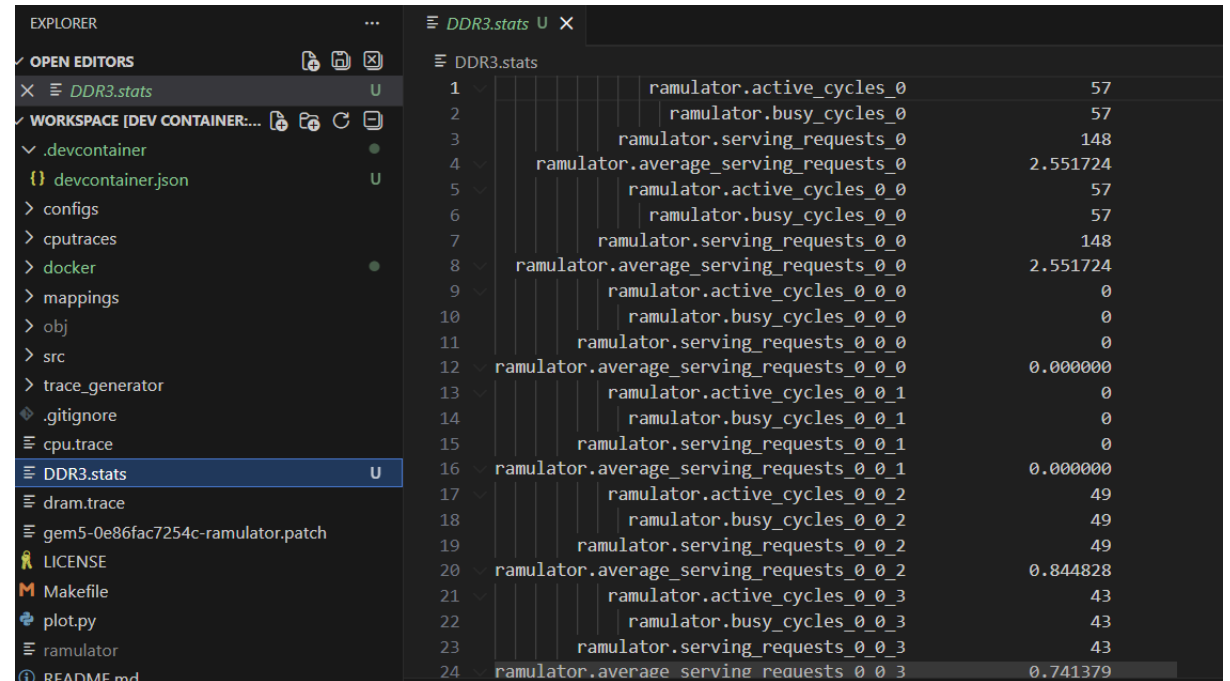
# yaml/trace file for **direct Read/Write Request**

## □ Simulation Result of **Ramulator1.0**



The screenshot shows the VS Code interface. The Explorer on the left lists the project files, including `devcontainer.json`, `configs`, `cpustraces`, `docker`, `mappings`, `obj`, `src`, `trace_generator`, `.gitignore`, `cpu.trace`, `DDR3.stats`, `dram.trace`, and `gem5-0e86fac7254c-ramulator.patch`. The `dram.trace` file is selected and open in the Editor. The content of `dram.trace` is as follows:

```
1 0x12345680 R
2 0x4cbd56c0 W
3 0x35d46f00 R
4 0x696fed40 W
5 0x7876af80 R
6
```



The screenshot shows the VS Code interface. The Explorer on the left lists the project files, including `devcontainer.json`, `configs`, `cpustraces`, `docker`, `mappings`, `obj`, `src`, `trace_generator`, `.gitignore`, `cpu.trace`, `DDR3.stats`, `dram.trace`, `gem5-0e86fac7254c-ramulator.patch`, `LICENSE`, `Makefile`, `plot.py`, `ramulator`, and `README.md`. The `DDR3.stats` file is selected and open in the Editor. The content of `DDR3.stats` is as follows:

```
1 ramulator.active_cycles_0 57
2 ramulator.busy_cycles_0 57
3 ramulator.serving_requests_0 148
4 ramulator.average_serving_requests_0 2.551724
5 ramulator.active_cycles_0_0 57
6 ramulator.busy_cycles_0_0 57
7 ramulator.serving_requests_0_0 148
8 ramulator.average_serving_requests_0_0 2.551724
9 ramulator.active_cycles_0_0_0 0
10 ramulator.busy_cycles_0_0_0 0
11 ramulator.serving_requests_0_0_0 0
12 ramulator.average_serving_requests_0_0_0 0.000000
13 ramulator.active_cycles_0_0_1 0
14 ramulator.busy_cycles_0_0_1 0
15 ramulator.serving_requests_0_0_1 0
16 ramulator.average_serving_requests_0_0_1 0.000000
17 ramulator.active_cycles_0_0_2 49
18 ramulator.busy_cycles_0_0_2 49
19 ramulator.serving_requests_0_0_2 49
20 ramulator.average_serving_requests_0_0_2 0.844828
21 ramulator.active_cycles_0_0_3 43
22 ramulator.busy_cycles_0_0_3 43
23 ramulator.serving_requests_0_0_3 43
24 ramulator.average_serving_requests_0_0_3 0.741379
```

Ramulator1.0은 납득할 수준의 결과가 나옴

Ramulator2.0의 전체 시스템을 끝까지 리뷰 후 Ramulator1.0의 DRAM subsystem만 simulation한 결과를 바탕으로 다시 분석 예정

## □ simpleO3 Core Operations

- class SimpleO3 – init()

```
28 void init() override {
29     m_clock_ratio = param<uint>("clock_ratio").required();
30
31     // Core params
32     std::vector<std::string> trace_list = param<std::vector<std::string>>("traces").desc("A list of traces.").required();
33     m_num_cores = trace_list.size();
34
35     int ipc = param<int>("ipc").desc("IPC of the SimpleO3 core.").default_val(4);
36     int depth = param<int>("inst_window_depth").desc("Instruction window size of the SimpleO3 core.").default_val(128);
37
38     // LLC params
39     int llc_latency = param<int>("llc_latency").desc("Aggregated latency of the LLC.").default_val(47);
40     int llc_linesize_bytes = param<int>("llc_linesize").desc("LLC cache line size in bytes.").default_val(64);
41     int llc_associativity = param<int>("llc_associativity").desc("LLC set associativity.").default_val(8);
42     int llc_capacity_per_core = parse_capacity_str(param<std::string>("llc_capacity_per_core").desc("LLC capacity per core.").default_val("2MB"));
43     int llc_num_mshr_per_core = param<int>("llc_num_mshr_per_core").desc("Number of LLC MSHR entries per core.").default_val(16);
44
45     // Simulation parameters
46     m_num_expected_insts = param<int>("num_expected_insts").desc("Number of instructions that the frontend should execute.").required();
47
48     // Create address translation module
49     m_translation = create_child_ifc<ITranslation>();
50
51     // Create the LLC
52     m_llc = new SimpleO3LLC(llc_latency, llc_capacity_per_core * m_num_cores, llc_linesize_bytes, llc_associativity, llc_num_mshr_per_core * m_num_cores);
53     // m_llc->deserialize(serialization_filename);
54     // m_llc->serialize(serialization_filename);
55
56     // Create the cores
57     for (int id = 0; id < m_num_cores; id++) {
58         SimpleO3Core* core = new SimpleO3Core(id, ipc, depth, m_num_expected_insts, trace_list[id], m_translation, m_llc);
59         core->m_callback = [this](Request& req){return this->receive(req);};
60         m_cores.push_back(core);
61     }
62
63     m_logger = Logging::create_logger("SimpleO3");
64
65     // Register the stats
66     register_stat(m_num_expected_insts).name("num_expected_insts");
67     register_stat(m_llc->s_llc_eviction).name("llc_eviction");
68     register_stat(m_llc->s_llc_read_access).name("llc_read_access");
69     register_stat(m_llc->s_llc_write_access).name("llc_write_access");
70     register_stat(m_llc->s_llc_read_misses).name("llc_read_misses");
71     register_stat(m_llc->s_llc_write_misses).name("llc_write_misses");
72     register_stat(m_llc->s_llc_mshr_unavailable).name("llc_mshr_unavailable");
73
74     for (int core_id = 0; core_id < m_cores.size(); core_id++) {
75         // register_stat(m_cores[core_id]->s_insts_retired).name("cycles_retired_core_{}", core_id);
76         register_stat(m_cores[core_id]->s_cycles_recorded).name("cycles_recorded_core_{}", core_id);
77         register_stat(m_cores[core_id]->s_mem_access_cycles).name("memory_access_cycles_recorded_core_{}", core_id);
78     }
79 }
80
```

1.1. Core Parameters를 yaml file에서 가져옴 → clock\_ratio / traces (ipc, depth 하드 코딩)

1.2. LLC Parameters (latency, cache line size, 몇 assosiative cahce인지, MSHR 수, cache size) 하드 코딩

1.3. Address Translation / LLC / Cores(trace file의 수만큼 생성 됨) Instance 생성

1.4. LLC의 동작 통계 변수들은 registration class에 등록

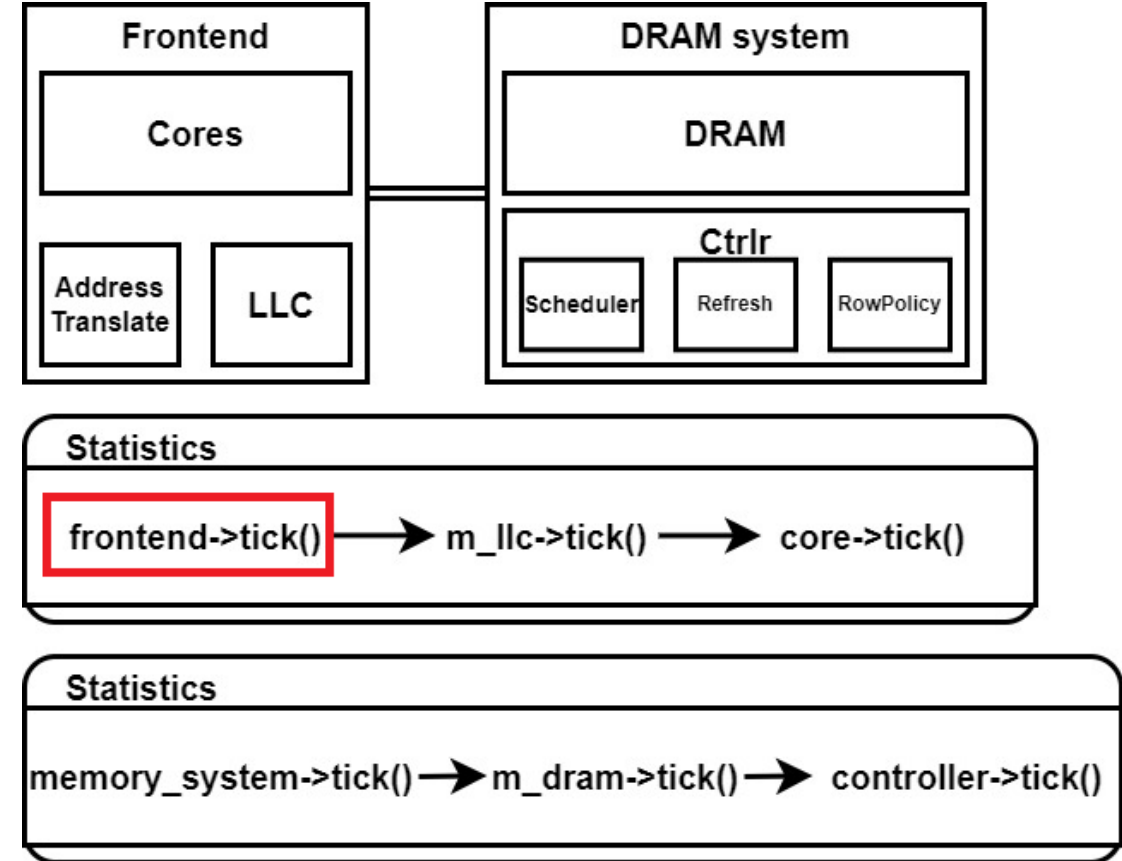
# src/frontend folder

## □ simpleO3 Core Operations

- class SimpleO3 – tick() → **frontend->tick()**

```
81 void tick() override {  
82     m_clk++;  
83  
84     if(m_clk % 10000000 == 0) {  
85         m_logger->info("Processor Heartbeat {} cycles.", m_clk);  
86     }  
87  
88     m_llc->tick();  
89     for (auto core : m_cores) {  
90         core->tick();  
91     }  
92 }
```

- m\_clk (global clk) 증가
- 아래의 순서로 frontend 동작
- Last Level Cache tick() → Core tick()

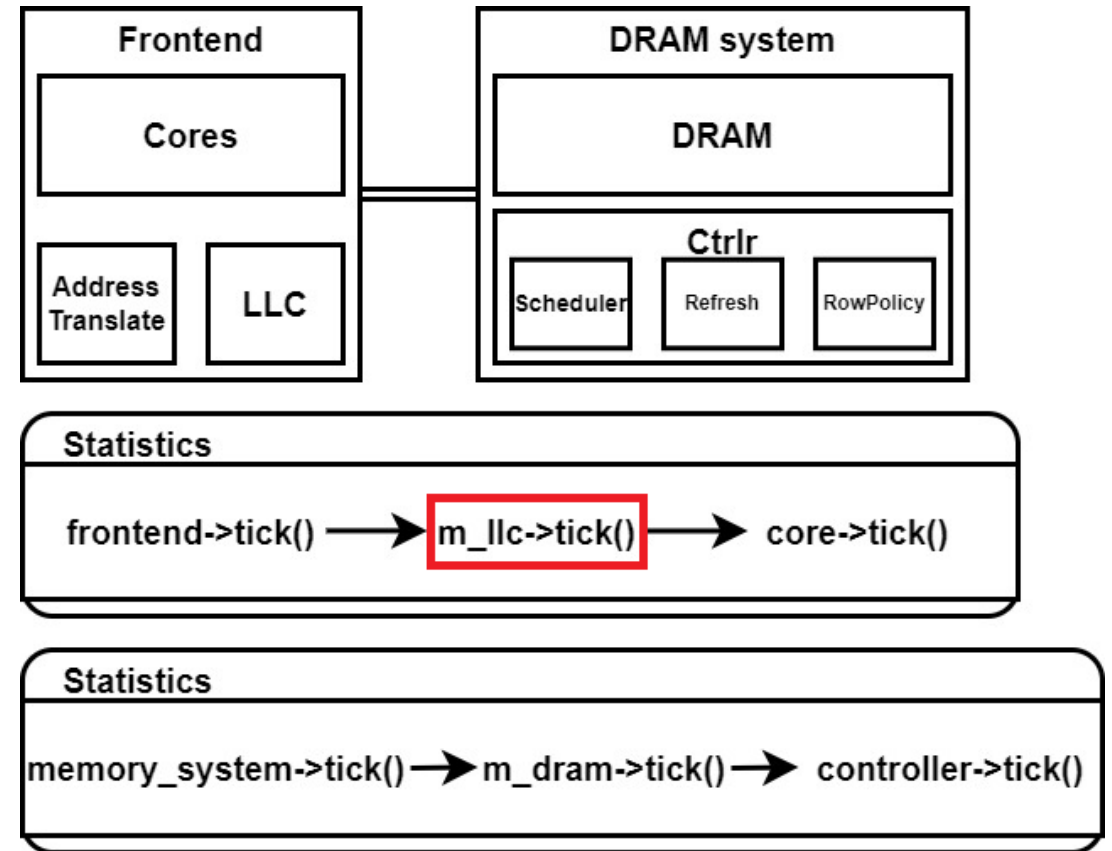


## □ simpleO3 Core Operations

- class SimpleO3 – tick() → **m\_llc->tick()**

```
20 void SimpleO3LLC::tick() {
21     m_clk++;
22
23     // Send miss requests to the memory system when LLC latency is met
24     // TODO: Optimization by assuming in-order issue?
25     auto it = m_miss_list.begin();
26     while (it != m_miss_list.end()) {
27         if (m_clk >= it->first) {
28             if (!m_memory_system->send(it->second)) {
29                 it++;
30             }
31             else {
32                 it = m_miss_list.erase(it);
33             }
34         } else {
35             it++;
36         }
37     }
38
39     // call hit request callback when LLC latency is met
40     it = m_hit_list.begin();
41     while (it != m_hit_list.end()) {
42         if (m_clk >= it->first) {
43             std::vector<Request> _req_v{it->second};
44             m_receive_requests[it->second.addr] = _req_v;
45
46             it->second.callback(it->second);
47             it = m_hit_list.erase(it);
48         }
49         else {
50             it++;
51         }
52     }
53 }
```

```
207 void SimpleO3LLC::evict_line(CacheSet_t& set, CacheSet_t::iterator victim_it) {
208     DEBUG_LOG(OSIMPLEO3LLC, m_logger, "Evicting {}.", victim_it->addr);
209     s_llc_eviction++;
210
211     // Generate writeback request if victim line is dirty
212     if (victim_it->dirty) {
213         Request writeback_req(victim_it->addr, Request::Type::Write);
214         m_miss_list.push_back(std::make_pair(m_clk + m_latency, writeback_req));
215
216         DEBUG_LOG(OSIMPLEO3LLC, m_logger, "Writeback Request will be issued at Clk={}.", m_clk + m_latency);
217     }
218     set.erase(victim_it);
219 }
220 }
```



- Miss List에서 지연 끝난 Missed Request을 MemorySystem으로 전달
- LLC latency(m\_latency) Simulation & Send missed req to DRAM
- 현재: LLC latency modeling 값 → m\_latency = 47

## □ simpleO3 Core Operations

- class SimpleO3 – tick() → **m\_llc->tick()**

```
20 void SimpleO3LLC::tick() {
21     m_clk++;
22
23     // Send miss requests to the memory system when LLC latency is met
24     // TODO: Optimization by assuming in-order issue?
25     auto it = m_miss_list.begin();
26     while (it != m_miss_list.end()) {
27         if (m_clk >= it->first) {
28             if (!m_memory_system->send(it->second)) {
29                 it++;
30             }
31             else {
32                 it = m_miss_list.erase(it);
33             }
34         } else {
35             it++;
36         }
37     }
38
39     // call hit request callback when LLC latency is met
40     it = m_hit_list.begin();
41     while (it != m_hit_list.end()) {
42         if (m_clk >= it->first) {
43             std::vector<Request> req_v(it->second);
44             m_receive_requests[it->second.addr] = req_v;
45             it->second.callback(it->second);
46             it = m_hit_list.erase(it);
47         }
48         else {
49             it++;
50         }
51     }
52 }
53 }
```

```
64 if (auto line_it = check_set_hit(set, req.addr); line_it != set.end()) {
65     // Hit in the set
66     DEBUG_LOG(DSIMPLEO3LLC, m_logger,
67         "[clk={}] Request Source: {}, Type: {}, Addr: {}, Index: {}, Tag: {}. Hit, will finish at clk={}",
68         m_clk, req.source_id, req.type_id, req.addr, get_index(req.addr), get_tag(req.addr), m_clk, m_clk + m_latency);
69
70     // Update the LRU status
71     set.push_back({req.addr, get_tag(req.addr), line_it->dirty || (req.type_id == Request::Type::Write), true});
72     set.erase(line_it);
73
74     // Add to the hit list to callback when finished
75     m_hit_list.push_back(std::make_pair(m_clk + m_latency, req));
76     return true;
77 } else {
78 }
```

```
56 // Create the cores
57 for (int id = 0; id < m_num_cores; id++) {
58     SimpleO3Core* core = new SimpleO3Core(id, ipc, depth, m_num_expected_insts, trace_list);
59     core->m_callback = [this](Request& req){return this->receive(req);};
60     m_cores.push_back(core);
61 }
```

```
void receive(Request& req) {
    m_llc->receive(req);

    // TODO: LLC latency for the core to receive the request?
    for (auto r : m_llc->m_receive_requests[req.addr]) {
        r.arrive = req.arrive;
        r.depart = req.depart;
        m_cores[r.source_id]->receive(r);
    }
    m_llc->m_receive_requests[req.addr].clear();
};
```

- Hit List에서 자연 끝난 Request들을 m\_receive\_requests배열에 저장
  - 여러 Core가 같은 주소를 가지는 Requests를 기다리고 있을 때, 완료되면 모든 Core에게 통지할 수 있도록 Mapping 하는 것
    - ex. m\_receive\_requests[0x1000] = [Request(addr=0x1000, source\_id=0)]
- callback이 호출됨 (이 callback은 SimpleO3::receive를 가리킴)

## □ simpleO3 Core Operations

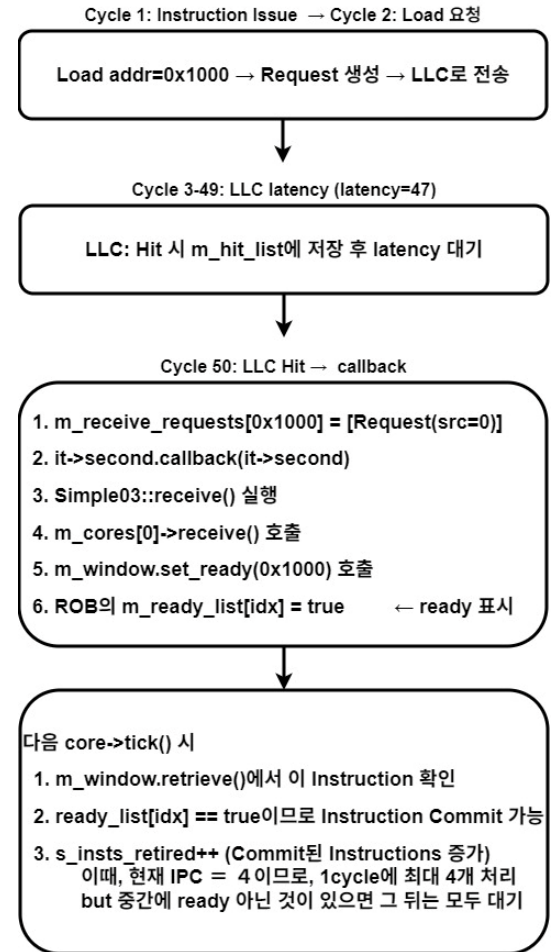
- class SimpleO3 – tick() → m\_llc->tick() → **callback**

```
56 // Create the cores
57 for (int id = 0; id < m_num_cores; id++) {
58     SimpleO3Core* core = new SimpleO3Core(id, ipc, depth, m_num_expected_insts, trace_list);
59     core->m_callback = [this](Request& req){return this->receive(req);};
60     m_cores.push_back(core);
61 }
```

```
void receive(Request& req) {
    m_llc->receive(req); // cache 내 MSLR 처리
    // TODO: LLC latency for the core to receive the request?
    for (auto r : m_llc->m_receive_requests[req.addr]) {
        r.arrive = req.arrive;
        r.depart = req.depart;
        m_cores[r.source_id]->receive(r);
    }
    m_llc->m_receive_requests[req.addr].clear();
}
```

```
184 void SimpleO3Core::receive(Request& req) {
185     m_window.set_ready(req.addr);
186
187     if (req.arrive != -1 && req.depart > m_last_mem_cycle) {
188         if (!reached_expected_num_insts) {
189             s_mem_access_cycles += (req.depart - std::max(m_last_mem_cycle, req.arrive));
190             m_last_mem_cycle = req.depart;
191         }
192     }
193 }
```

```
90 void SimpleO3Core::InstWindow::set_ready(Addr_t addr) {
91     if (m_load == 0) return;
92
93     int index = m_tail_idx;
94     for (int i = 0; i < m_load; i++) {
95         if (m_addr_list[index] == addr) {
96             m_ready_list[index] = true;
97         }
98         index++;
99         if (index == m_depth) {
100             index = 0;
101         }
102     }
103 }
```



- Out of order Processor는 Instruction의 Dependency를 고려하여 Instruction Re-Ordering 진행
- ROB에 쌓인 이미 완료된 Instruction은 이제 sequential하게 기존 Instruction 순서에 맞게 commit됨
- 항상 Processor는 User입장에서 Sequential하게 보여야 함 - The MicroArchitecture of SuperScalar Processors(1995)

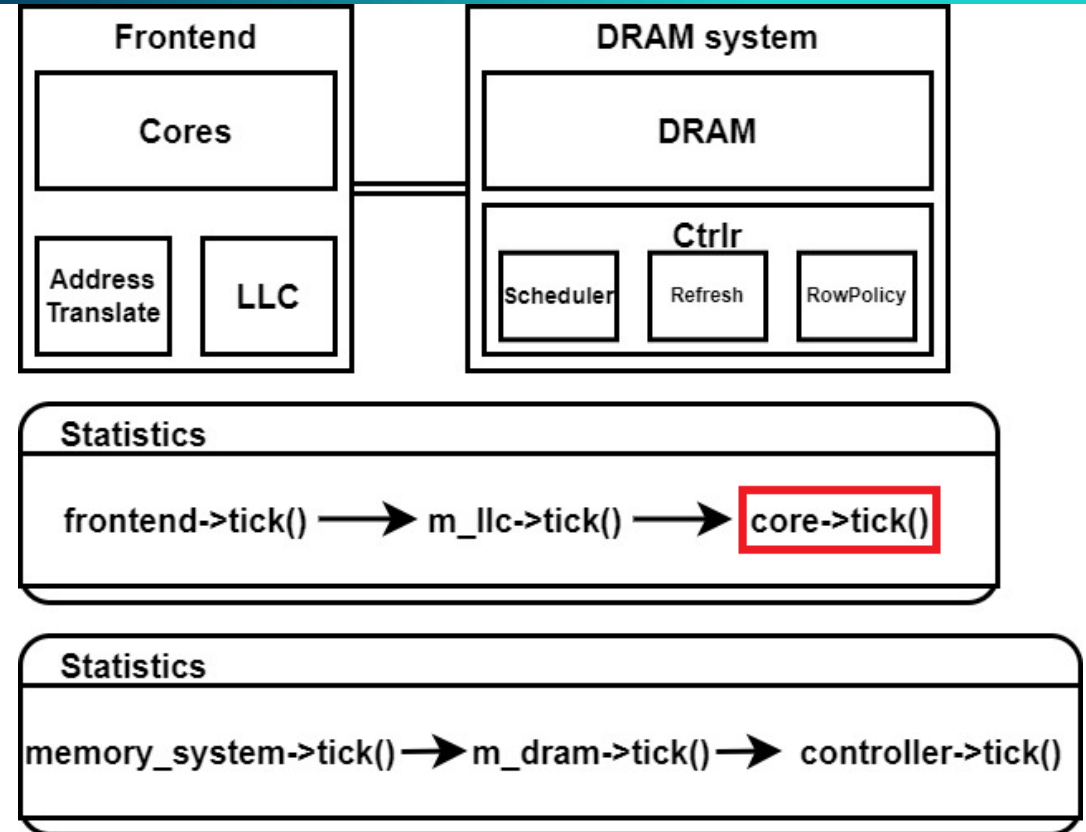


## □ simpleO3 Core Operations

- class SimpleO3 – tick() → **core->tick()**

```
148 void SimpleO3Core::tick() {
149     m_clk++;
150
151     // 목표 명령어 개수 도달 시 시간 기록
152     s_insts_retired += m_window.retire();
153     if (!reached_expected_num_insts) {
154         if (s_insts_retired >= m_num_expected_insts) {
155             reached_expected_num_insts = true;
156             s_cycles_recorded = m_clk;
157         }
158     }
159 }
```

```
96 int SimpleO3Core::InstWindow::retire() {
97     if (m_load == 0) return 0; // ROB 비어있으면 0 반환
98
99     int num_retired = 0;
100     while (m_load > 0 && num_retired < m_ipc) { // IPC 제한 내에서
101         if (!m_ready_list.at(m_tail_idx)) // tail이 준비 안 됐으면
102             break; // 즉시 중단 (In-Order Retirement!)
103
104         m_tail_idx = (m_tail_idx + 1) % m_depth; // tail 이동
105         m_load--; // ROB 사용량 감소
106         num_retired++; // commit한 명령어 개수 증가
107     }
108     return num_retired; // 이 cycle에 commit된 명령어 수 반환
109 }
```



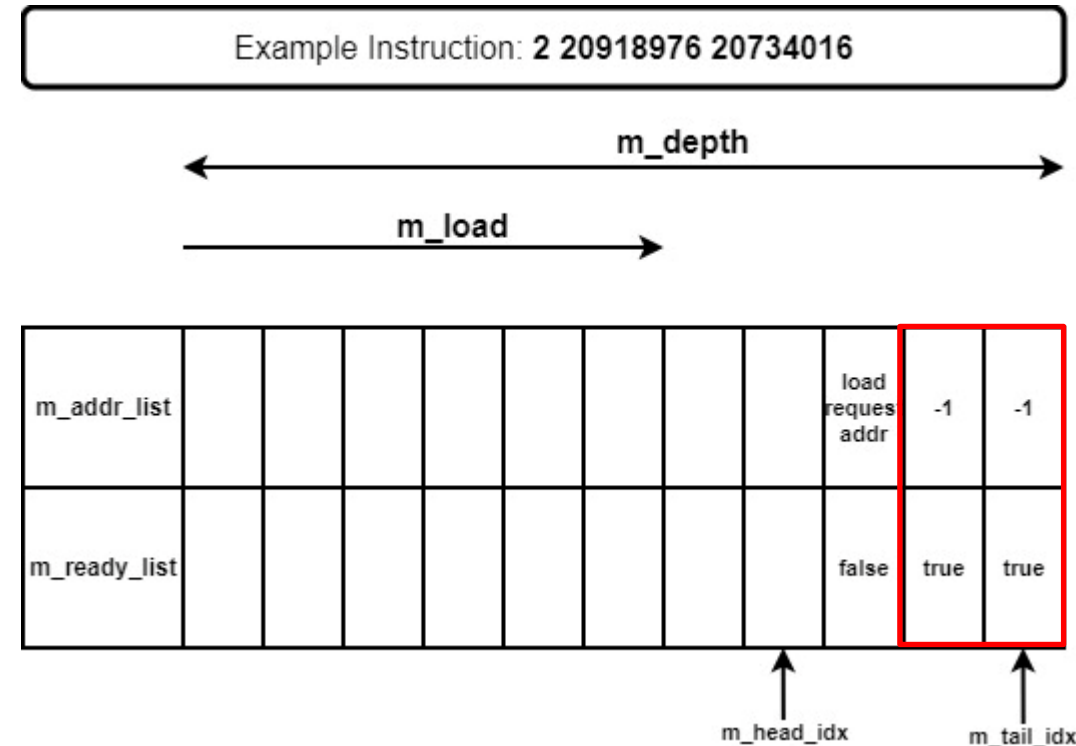
- ROB에서 명령어가 실행될 준비가 되면 ready flag를 **true**로 바꿈
- 이 경우 instruction이 실행되기 위한 data가 존재한다는 뜻
- 따라서, **ready = true**이면 해당 Inst를 commit 시킴 → Inst 처리됨

## □ simpleO3 Core Operations

- class SimpleO3 – tick() → **core->tick()**

```
160 // First, issue the non-memory instructions
161 int num_inserted_insts = 0;
162 while (m_num_bubbles > 0) {
163     if (num_inserted_insts == m_window.m_ipc) { // IPC 제한 도달
164         return;
165     }
166     if (m_window.is_full()) { // ROB 가득 참
167         return;
168     };
169     m_window.insert(true, -1); // ready=true (메모리 operation 아니기에)
170     num_inserted_insts++;
171     m_num_bubbles--;
172 }
```

```
80 // ROB에 새로운 명령어를 삽입
81 void SimpleO3Core::InstWindow::insert(bool ready, Addr_t addr) {
82     m_ready_list.at(m_head_idx) = ready;
83     m_addr_list.at(m_head_idx) = addr;
84
85     m_head_idx = (m_head_idx + 1) % m_depth;
86     m_load++;
87 }
```



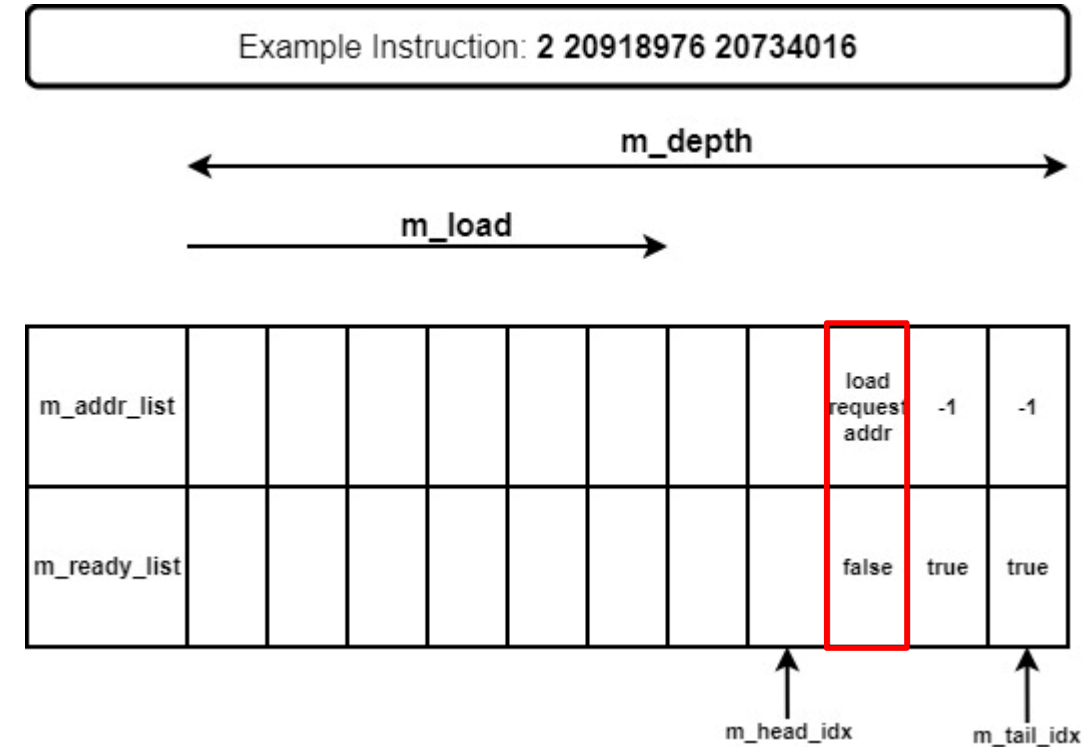
- ROB에 Non-Memory Operation을 최대 IPC이하의 개수만큼 채워 넣음
- 위 Example Inst. (= 2 20918976 20734016)의 경우 Non-Memory Op가 2개이기에 위 그림의 **빨간 박스 부분**이 완성됨

## □ simpleO3 Core Operations

- class SimpleO3 – tick() → **core->tick()**

```
174 // Second, try to send the load to the LLC
175 if (m_load_addr != -1) {
176     if (num_inserted_insts == m_window.m_ipc) { // IPC 제한
177         return;
178     }
179     if (m_window.is_full()) { // ROB 가득 참
180         return;
181     };
182
183     // Load 요청 생성 (가상 주소)
184     Request load_request(m_load_addr, Request::Type::Read, m_id, m_callback);
185
186     // 주소 변환 (가상 → 물리)
187     if (!m_translation->translate(load_request)) {
188         return; // 변환 실패
189     };
190
191     // LLC로 요청 발송
192     if (m_llc->send(load_request)) {
193         m_window.insert(false, load_request.addr); // ready=false (메모리 대기)
194         m_load_addr = -1; // 처리 완료
195         if (m_writeback_addr != -1) { // Store 있으면 다음 사이클에 발행
196             // If there is still writeback, return without getting the next trace line
197             // The write back will be issued in the next cycle
198             // TODO: Should we allow both load and writeback to issue at the same cycle?
199             return;
200         }
201     } else {
202         return;
203     }
204 }
```

ROB = InstWindow



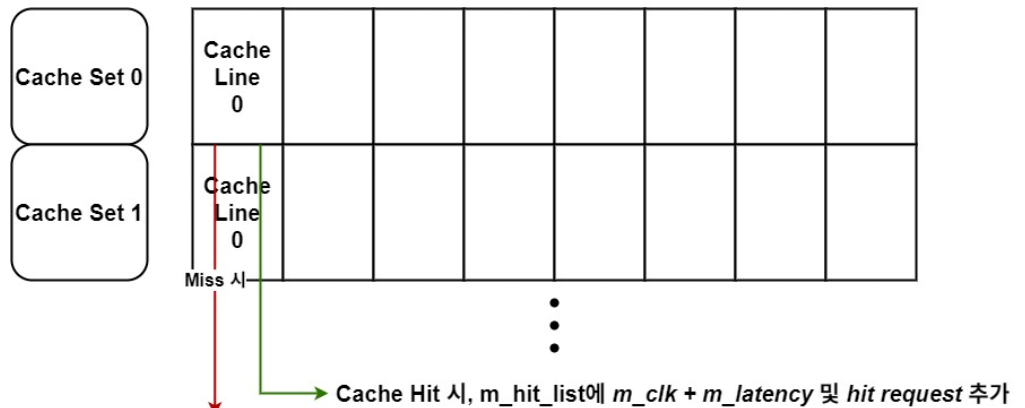
- Load Request가 존재하면, Load 요청의 주소를 translate
- 그 후, LLC로 Request를 보냄

## □ simpleO3 Core Operations

- class SimpleO3 – tick() → core->tick() → **m\_llc->send(load\_request)**

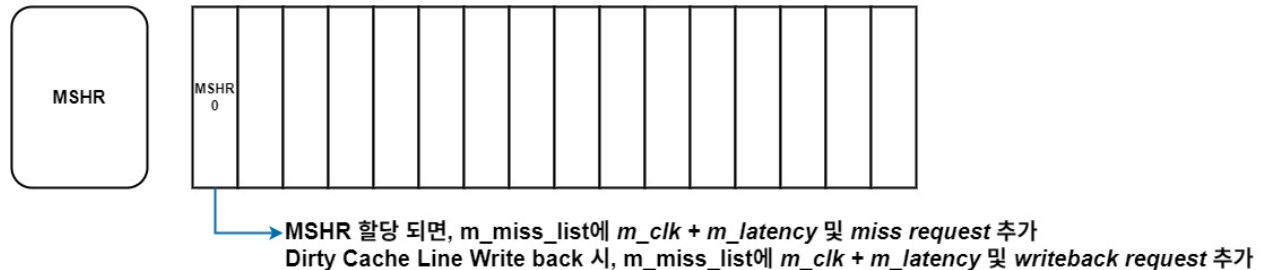
```
174 // Second, try to send the load to the LLC
175 if (m_load_addr != -1) {
176     if (num_inserted_insts == m_window.m_ipc) { // IPC 제한
177         return;
178     }
179     if (m_window.is_full()) { // ROB 가득 참
180         return;
181     };
182
183     // Load 요청 생성 (가상 주소)
184     Request load_request(m_load_addr, Request::Type::Read, m_id, m_callback);
185
186     // 주소 변환 (가상 → 물리)
187     if (!m_translation->translate(load_request)) {
188         return; // 변환 실패
189     };
190
191     // LLC로 요청 발송
192     if (m_llc->send(load_request)) {
193         m_window.insert(false, load_request.addr); // ready=false (메모리 대기)
194         m_load_addr = -1; // 처리 완료
195         if (m_writeback_addr != -1) { // Store 있으면 다음 사이클에 발행
196             // If there is still writeback, return without getting the next trace line
197             // The write back will be issued in the next cycle
198             // TODO: Should we allow both load and writeback to issue at the same cycle?
199             return;
200         }
201     } else {
202         return;
203     }
204 }
```

check\_set\_hit(set, req.addr): Tag 일치 && line\_it->ready=true이면 Hit



check\_mshr\_hit(req.addr)

- MSHR Hit이면 → 기존 DRAM 요청에 병합 (MSHR Hit는 Cache Line단위로 Check)
- m\_receive\_requests[mshr\_it->first].push\_back(req);
- MSHR Miss이면 → 새 DRAM 요청 발송
- MSHR entry가 부족 or 해당 Set의 Cache line이 Unavailable일 때 stall
- 새로운 Cache Line 가져올 시, 기존이 dirty이면 cache line write back 발생



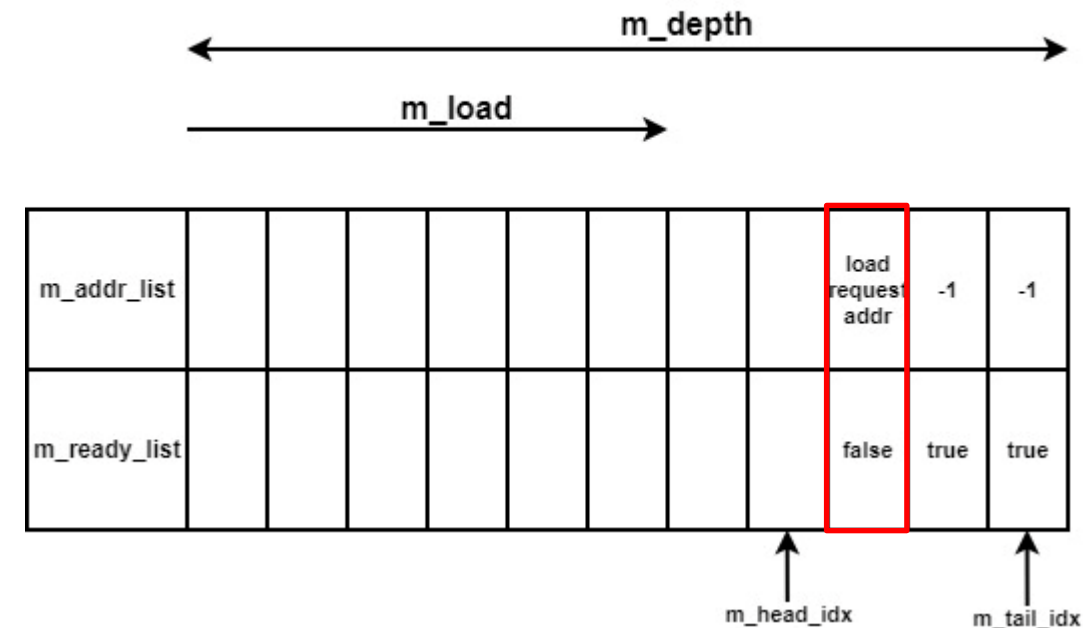
## □ simpleO3 Core Operations

- class SimpleO3 – tick() → **core->tick()**

```
174 // Second, try to send the load to the LLC
175 if (m_load_addr != -1) {
176     if (num_inserted_insts == m_window.m_ipc) { // IPC 제한
177         return;
178     }
179     if (m_window.is_full()) { // ROB 가득 참
180         return;
181     };
182
183     // Load 요청 생성 (가상 주소)
184     Request load_request(m_load_addr, Request::Type::Read, m_id, m_callback);
185
186     // 주소 변환 (가상 → 물리)
187     if (!m_translation->translate(load_request)) {
188         return; // 변환 실패
189     };
190
191     // LLC로 요청 발송
192     if (m_llc->send(load_request)) {
193         m_window.insert(false, load_request.addr); // ready=false (메모리 대기)
194         m_load_addr = -1; // 처리 완료
195         if (m_writeback_addr != -1) { // Store 있으면 다음 사이클에 발행
196             // If there is still writeback, return without getting the next trace line
197             // The write back will be issued in the next cycle
198             // TODO: Should we allow both load and writeback to issue at the same cycle?
199             return;
200         }
201     } else {
202         return;
203     }
204 }
205
206 // ROB에 새로운 명령어를 삽입
207 void SimpleO3Core::InstWindow::insert(bool ready, Addr_t addr) {
208     m_ready_list.at(m_head_idx) = ready;
209     m_addr_list.at(m_head_idx) = addr;
210
211     m_head_idx = (m_head_idx + 1) % m_depth;
212     m_load++;
213 }
```

ROB = InstWindow

Example Instruction: 2 20918976 20734016



- LLC에 Request를 send하였으므로, 해당 Load Request를 Instruction Window에 넣음
- Write Request가 존재 시, 이는 다음 tick()=cycle에 처리
- Write Request가 존재 시, return됨

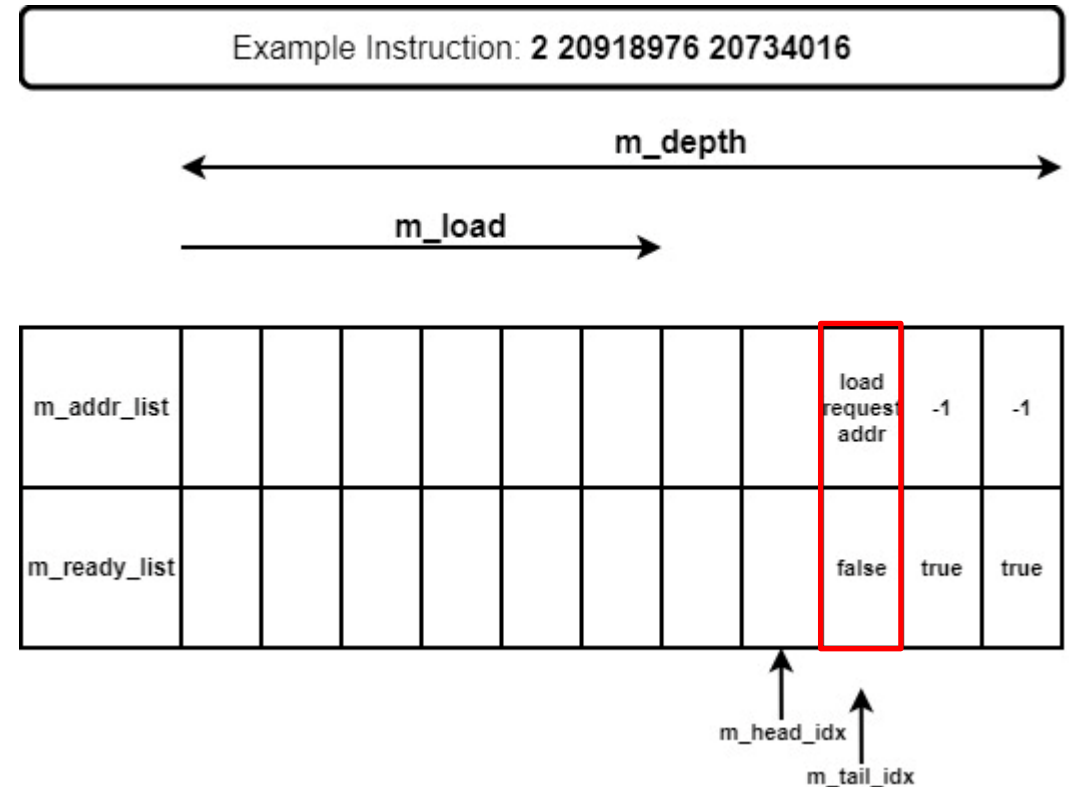
## □ simpleO3 Core Operations

- class SimpleO3 – tick() → **core->tick()**

```
206 // Third, try to send the writeback to the LLC
207 // Store(writeback) 요청 발행
208 if (m_writeback_addr != -1) {
209     Request writeback_request(m_writeback_addr, Request::Type::Write, m_id, m_callback);
210     if (!m_translation->translate(writeback_request)) {
211         return;
212     };
213     if (!m_llc->send(writeback_request)) {
214         return;
215     }
216 }
217
218 // 다음 명령어 미리 로드
219 auto inst = m_trace.get_next_inst();
220 m_num_bubbles = inst.bubble_count;
221 m_load_addr = inst.load_addr;
222 m_writeback_addr = inst.store_addr;
223 }
```

- Write Request가 존재 시, return되었기에, 아래 다음 Instruction load되지 않고, 진행되고 있던 Inst 처리
- core->tick()에서 1. *issue the non-memory instructions* / 2. *try to send the load to the LLC* 건너뛰고 Write 처리
- Write Request도 Read와 마찬가지로, LLC의 send()함수를 실행시킨다.
- 그 후, 다음 Instruction의 정보를 core.h에 update한다.
- 또한, Write Request는 CPU가 issue후 신경 안 써도 되니, ROB에 저장하지 않는다.

ROB = InstWindow



# src/memory\_system folder

## □ memory\_system

- class SimpleO3 – tick() → m\_llc->tick() → **m\_memory\_system – send()**

```
37 void SimpleO3LLC::tick() {
38     m_clk++;
39
40     // Send miss requests to the memory system when LLC latency is met
41     // TODO: Optimization by assuming in-order issue?
42     auto it = m_miss_list.begin();
43     while (it != m_miss_list.end()) {
44         if (m_clk >= it->first) { // 지연 완료되었는가?
45             if (!m_memory_system->send(it->second)) { // 발송 실패 → 다음 사이클 재시도
46                 it++;
47             }
48             else {
49                 it = m_miss_list.erase(it); // 발송 성공 → 제거
50             }
51         } else {
52             it++; // 아직 시간 미도달
53         }
54     }
55 }
```

```
50 bool send(Request req) override {
51     // separete physical address to dram vector, [channel, rank, bank, ...]
52     m_addr_mapper->apply(req);
53     int channel_id = req.addr_vec[0]; // 채널 ID 추출
54     bool is_success = m_controllers[channel_id]->send(req); // 해당 채널 컨트롤러로 요청 전송
55
56     // update statisticss
57     if (is_success) {
58         switch (req.type_id) {
59             case Request::Type::Read: {
60                 s_num_read_requests++;
61                 break;
62             }
63             case Request::Type::Write: {
64                 s_num_write_requests++;
65                 break;
66             }
67             default: {
68                 s_num_other_requests++;
69                 break;
70             }
71         }
72     }
73
74     return is_success;
75 }
```

- Frontend에서 LLC miss가 발생 시, 해당 request들을 m\_miss\_list에 queuing하고,
- cache latency가 끝나면, m\_miss\_list의 request들을 memory\_system으로 send
- memory\_system은 physical address → dram vector 후 알맞은 channel의 ctrlr으로 send



# src/memory\_system folder

## memory\_system

- m\_memory\_system – send() → m\_controllers[channel\_id] – send()

```
50 bool send(Request req) override {
51     // separete physical address to dram vector, [channel, rank, bank, ...]
52     m_addr_mapper->apply(req);
53     int channel_id = req.addr.vec[0]; // 채널 ID 추출
54     bool is_success = m_controllers[channel_id]->send(req); // 해당 채널 컨트롤러로 요청 전송
55
56     // update statisticss
57     if (is_success) {
58         switch (req.type_id) {
59             case Request::Type::Read: {
60                 s_num_read_requests++;
61                 break;
62             }
63             case Request::Type::Write: {
64                 s_num_write_requests++;
65                 break;
66             }
67             default: {
68                 s_num_other_requests++;
69                 break;
70             }
71         }
72     }
73     return is_success;
74 }
75
```

- 각 request의 final command를 지정
- Requests의 통계 update

```
115 bool send(Request& req) override {
116     // Request type별 최종 커맨드 결정
117     req.final_command = m_dram->m_request_translations(req.type_id);
118
119     switch (req.type_id) {
120         case Request::Type::Read: {
121             s_num_read_reqs++;
122             break;
123         }
124         case Request::Type::Write: {
125             s_num_write_reqs++;
126             break;
127         }
128         default: {
129             s_num_other_reqs++;
130             break;
131         }
132     }
133
134     inline static const ImplLUT m_request_translations = LUT (
135         m_requests, m_commands, {
136             {"read", "RD"}, {"write", "WR"}, {"all-bank-refresh", "REFab"},
137             {"open-row", "ACT"}, {"close-row", "PRE"}
138         }
139     );
140 }
```

```
147 // Else, enqueue them to corresponding buffer based on request type id
148 bool is_success = false;
149 req.arrive = m_clk; // 도착 시간 기록
150 if (req.type_id == Request::Type::Read) {
151     is_success = m_read_buffer.enqueue(req);
152 } else if (req.type_id == Request::Type::Write) {
153     is_success = m_write_buffer.enqueue(req);
154 } else {
155     throw std::runtime_error("Invalid request type!");
156 }
157 if (!is_success) {
158     // We could not enqueue the request
159     req.arrive = -1;
160     return false;
161 }
162
163 return true;
164 }
```

# src/memory\_system folder

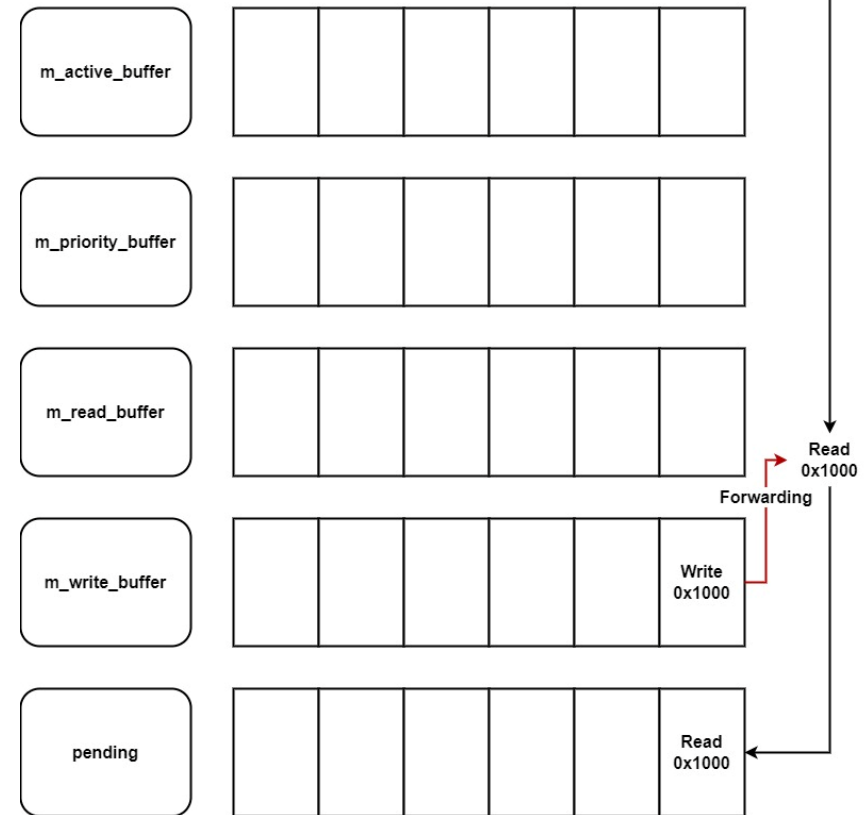
## memory\_system

- m\_memory\_system – send() → m\_controllers[channel\_id] – send()

```
134 // Forward existing write requests to incoming read requests
135 if (req.type_id == Request::Type::Read) {
136     auto compare_addr = [req](const Request& wreq) {
137         return wreq.addr == req.addr;
138     };
139     if (std::find_if(m_write_buffer.begin(), m_write_buffer.end(), compare_addr) != m_write_buffer.end()) {
140         // The request will depart at the next cycle
141         req.depart = m_clk + 1;
142         pending.push_back(req); // DRAM 접근 없이 반환 -> 다음 사이클에 데이터 전달
143         return true;
144     }
145 }
```

- Read Request가 들어왔을 시,
- 기존 write buffer에 Read Request와 같은 주소를 가지는 Req 존재 시,
- Data Forwarding 진행
- Read요청이 원하는 Data를 forwarding을 통해 얻었음.
- 이 Read Request를 Pending Queue에 넣음

```
generic_dram_controller.cpp -> send(Request& req)
clk | 동작
100 | Write(0x1000, X)          → m_write_buffer 삽입
101 | Write 스케줄 대기...      → write mode가 아니라고 가정
102 | Write 스케줄 대기...      → write mode가 아니라고 가정
103 | Read(0x1000)              → m_write_buffer에서 같은 주소 발견
    | → depart = 104 설정
    | → pending queue 삽입
104 | serve_completed_reads() → callback 호출
    | → 올바른 data 반환
    | → latency: 1 cycle
```



# src/memory\_system folder

## memory\_system

- m\_memory\_system – send() → m\_controllers[channel\_id] – send()

```
147 // Else, enqueue them to corresponding buffer based on request type id
148 bool is_success = false;
149 req.arrive = m_clk; // 도착 시간 기록
150 if (req.type_id == Request::Type::Read) {
151     is_success = m_read_buffer.enqueue(req);
152 } else if (req.type_id == Request::Type::Write) {
153     is_success = m_write_buffer.enqueue(req);
154 } else {
155     throw std::runtime_error("Invalid request type!");
156 }
157 if (!is_success) {
158     // We could not enqueue the request
159     req.arrive = -1;
160     return false;
161 }
```

- Data Forwarding 진행할 수 있는 조건 X 시,
- Request Type에 맞춰서 해당하는 buffer에 넣음

generic\_dram\_controller.cpp -> send(Request& req)

Clk | 동작

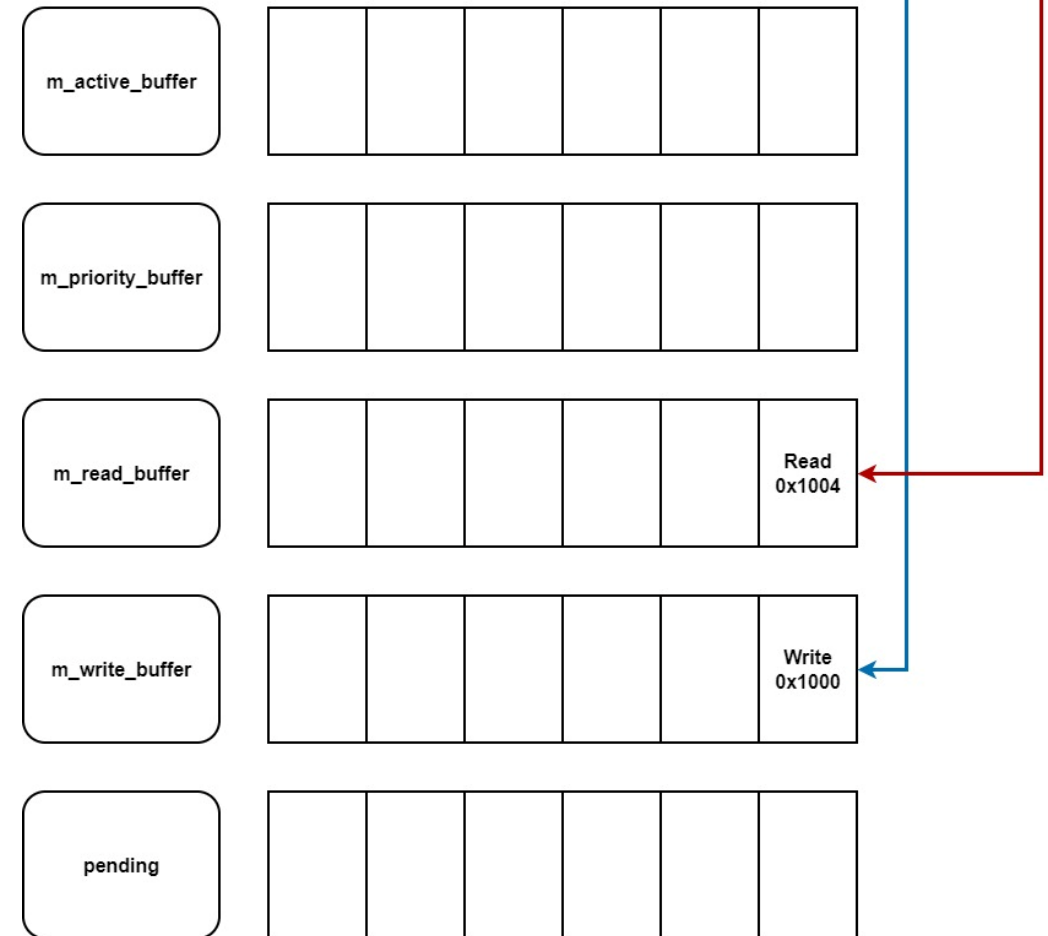
100 | Write(0x1000, X)

101 | Write 스케줄 대기...

102 | Write 스케줄 대기...

103 | Read(0x1004, Y)

- m\_write\_buffer 삽입
- write mode가 아니라고 가정
- write mode가 아니라고 가정
- m\_write\_buffer에서 같은 주소 발견 X

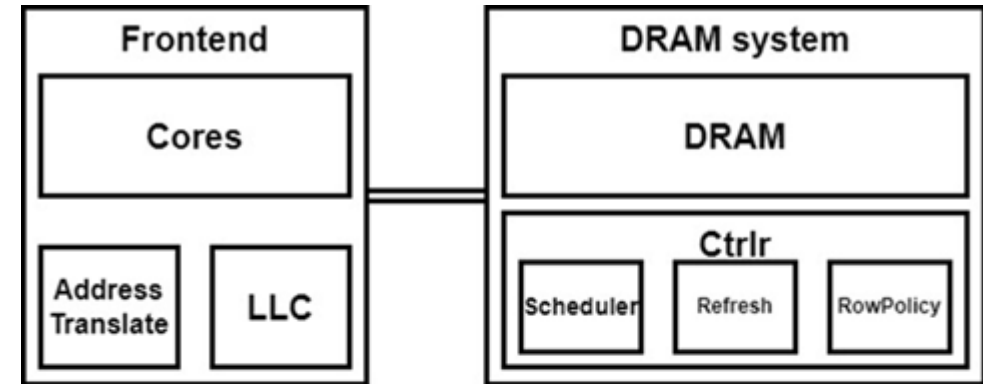


# src/memory\_system folder

## □ memory\_system

```
93 inline static const ImplLUT m_command_meta = LUT<DRAMCommandMeta> (  
94   m_commands, {  
95     // open?   close?   access?   refresh?  
96     {"ACT",    {true,   false,  false,   false}},  
97     {"PRE",    {false,  true,   false,   false}},  
98     {"PREA",   {false,  true,   false,   false}},  
99     {"RD",     {false,  false,  true,    false}},  
100    {"WR",     {false,  false,  true,    false}},  
101    {"RDA",    {false,  true,   true,    false}},  
102    {"WRA",    {false,  true,   true,    false}},  
103    {"REFab",   {false,  false,  false,   true }},  
104    {"REFab_end", {false, true,   false,  false}},  
105  }  
106 );
```

- abcd



### Statistics

frontend->tick() → m\_llc->tick() → core->tick()

### Statistics

memory\_system->tick() → m\_dram->tick() → controller->tick()

# src/addr\_mapper folder

□ **addr\_mapper**

- **abcd**

# src/addr\_mapper folder

□ **addr\_mapper**

- **abcd**

# src/translation folder

## □ translation

- abcd



# src/translation folder

## □ translation

- abcd

# src/translation folder

## □ translation

- abcd

# src/dram folder

□ dram

- abcd

# src/dram folder

□ dram

- abcd

# src/dram folder

□ dram

- abcd

## □ DRAM Controller Operations

- class GenericDRAMController – init()

(src\dram\_controller\impl\generic\_dram\_controller.cpp)

```
54 void init() override {
55     m_wr_low_watermark = param<float>("wr_low_watermark").desc("Threshold for switching back to read mode.").default_val(0.2f);
56     m_wr_high_watermark = param<float>("wr_high_watermark").desc("Threshold for switching to write mode.").default_val(0.8f);
57
58     m_scheduler = create_child_ifce<IScheduler>();
59     m_refresh = create_child_ifce<IRefreshManager>();
60     m_rowpolicy = create_child_ifce<IRowPolicy>();
61
62     if (m_config["plugins"]) {
63         YAML::Node plugin_configs = m_config["plugins"];
64         for (YAML::iterator it = plugin_configs.begin(); it != plugin_configs.end(); ++it) {
65             m_plugins.push_back(create_child_ifce<IControllerPlugin>(*it));
66         }
67     }
68 };
```

1.1. YAML file을 통해 write buffer priority threshold를 초기화. →

→ Scheduler가 사용할 buffer를 결정

1.2. Scheduler, Refresh Manager, Row Policy를 생성.

1.3. Plugin을 load하여 Controller를 확장

```
// 2.2.1 If no request to be scheduled in the priority buffer, check the read and write buffers.
if (!request_found) {
    // Query the write policy to decide which buffer to serve
    set_write_mode();
    auto& buffer = m_is_write_mode ? m_write_buffer : m_read_buffer;
    if (req_it = m_scheduler->get_best_request(buffer); req_it != buffer.end()) {
        request_found = m_dram->check_ready(req_it->command, req_it->addr_vec);
        req_buffer = &buffer;
    }
}
```

## □ DRAM Controller Operations

- class GenericDRAMController – setup()

```
70 void setup(IFrontEnd* frontend, IMemorySystem* memory_system) override {
71     m_dram = memory_system->get_ifce<IDRAM>();
72     m_bank_addr_idx = m_dram->m_levels("bank");
73     m_priority_buffer.max_size = 512*3 + 32;
74
75     m_num_cores = frontend->get_num_cores();
76
77     s_read_row_hits_per_core.resize(m_num_cores, 0);
78     s_read_row_misses_per_core.resize(m_num_cores, 0);
79     s_read_row_conflicts_per_core.resize(m_num_cores, 0);
80
81     register_stat(s_row_hits).name("row_hits_{}", m_channel_id);
82     register_stat(s_row_misses).name("row_misses_{}", m_channel_id);
83     register_stat(s_row_conflicts).name("row_conflicts_{}", m_channel_id);
84     register_stat(s_read_row_hits).name("read_row_hits_{}", m_channel_id);
85     register_stat(s_read_row_misses).name("read_row_misses_{}", m_channel_id);
86     register_stat(s_read_row_conflicts).name("read_row_conflicts_{}", m_channel_id);
87     register_stat(s_write_row_hits).name("write_row_hits_{}", m_channel_id);
88     register_stat(s_write_row_misses).name("write_row_misses_{}", m_channel_id);
89     register_stat(s_write_row_conflicts).name("write_row_conflicts_{}", m_channel_id);
90
91     for (size_t core_id = 0; core_id < m_num_cores; core_id++) {
92         register_stat(s_read_row_hits_per_core[core_id]).name("read_row_hits_core_{}", core_id);
93         register_stat(s_read_row_misses_per_core[core_id]).name("read_row_misses_core_{}", core_id);
94         register_stat(s_read_row_conflicts_per_core[core_id]).name("read_row_conflicts_core_{}", core_id);
95     }
96
97     register_stat(s_num_read_reqs).name("num_read_reqs_{}", m_channel_id);
98     register_stat(s_num_write_reqs).name("num_write_reqs_{}", m_channel_id);
99     register_stat(s_num_other_reqs).name("num_other_reqs_{}", m_channel_id);
100     register_stat(s_queue_len).name("queue_len_{}", m_channel_id);
101     register_stat(s_read_queue_len).name("read_queue_len_{}", m_channel_id);
102     register_stat(s_write_queue_len).name("write_queue_len_{}", m_channel_id);
103     register_stat(s_priority_queue_len).name("priority_queue_len_{}", m_channel_id);
104     register_stat(s_queue_len_avg).name("queue_len_avg_{}", m_channel_id);
105     register_stat(s_read_queue_len_avg).name("read_queue_len_avg_{}", m_channel_id);
106     register_stat(s_write_queue_len_avg).name("write_queue_len_avg_{}", m_channel_id);
107     register_stat(s_priority_queue_len_avg).name("priority_queue_len_avg_{}", m_channel_id);
108
109     register_stat(s_read_latency).name("read_latency_{}", m_channel_id);
110     register_stat(s_avg_read_latency).name("avg_read_latency_{}", m_channel_id);
111 }
```

2.1. DRAM Interface 연결 - (ex. dram.h → DDR4)

2.2. Bank level의 index를 조회

- DRAM Address Vector → ex. addr\_vec[0] = channel, [1] = rank, [2] = bankgroup, [3] = bank, [4] = row, [5] = column
- 이 경우, Bank level index = 3
- 아래는 Row Hit 판별을 위해 Bank level index가 어떻게 사용되는지를 나타냄
- m\_bank\_addr\_idx = 3 → bank = addr\_vec[3] = 2 → Row Hit Check: open\_rows[bank=2]의 row=1024 check → hit?

2.3. priority buffer Size를 하드 코딩



## □ DRAM Controller Operations

- class GenericDRAMController – setup()

```
70 void setup(IFrontEnd* frontend, IMemorySystem* memory_system) override {
71     m_dram = memory_system->get_ifce<IDRAM>();
72     m_bank_addr_idx = m_dram->m_levels("bank");
73     m_priority_buffer.max_size = 512*3 + 32;
74
75     m_num_cores = frontend->get_num_cores();
76
77     s_read_row_hits_per_core.resize(m_num_cores, 0);
78     s_read_row_misses_per_core.resize(m_num_cores, 0);
79     s_read_row_conflicts_per_core.resize(m_num_cores, 0);
80
81     register_stat(s_row_hits).name("row_hits_{}", m_channel_id);
82     register_stat(s_row_misses).name("row_misses_{}", m_channel_id);
83     register_stat(s_row_conflicts).name("row_conflicts_{}", m_channel_id);
84     register_stat(s_read_row_hits).name("read_row_hits_{}", m_channel_id);
85     register_stat(s_read_row_misses).name("read_row_misses_{}", m_channel_id);
86     register_stat(s_read_row_conflicts).name("read_row_conflicts_{}", m_channel_id);
87     register_stat(s_write_row_hits).name("write_row_hits_{}", m_channel_id);
88     register_stat(s_write_row_misses).name("write_row_misses_{}", m_channel_id);
89     register_stat(s_write_row_conflicts).name("write_row_conflicts_{}", m_channel_id);
90
91     for (size_t core_id = 0; core_id < m_num_cores; core_id++) {
92         register_stat(s_read_row_hits_per_core[core_id]).name("read_row_hits_core_{}", core_id);
93         register_stat(s_read_row_misses_per_core[core_id]).name("read_row_misses_core_{}", core_id);
94         register_stat(s_read_row_conflicts_per_core[core_id]).name("read_row_conflicts_core_{}", core_id);
95     }
96
97     register_stat(s_num_read_reqs).name("num_read_reqs_{}", m_channel_id);
98     register_stat(s_num_write_reqs).name("num_write_reqs_{}", m_channel_id);
99     register_stat(s_num_other_reqs).name("num_other_reqs_{}", m_channel_id);
100     register_stat(s_queue_len).name("queue_len_{}", m_channel_id);
101     register_stat(s_read_queue_len).name("read_queue_len_{}", m_channel_id);
102     register_stat(s_write_queue_len).name("write_queue_len_{}", m_channel_id);
103     register_stat(s_priority_queue_len).name("priority_queue_len_{}", m_channel_id);
104     register_stat(s_queue_len_avg).name("queue_len_avg_{}", m_channel_id);
105     register_stat(s_read_queue_len_avg).name("read_queue_len_avg_{}", m_channel_id);
106     register_stat(s_write_queue_len_avg).name("write_queue_len_avg_{}", m_channel_id);
107     register_stat(s_priority_queue_len_avg).name("priority_queue_len_avg_{}", m_channel_id);
108
109     register_stat(s_read_latency).name("read_latency_{}", m_channel_id);
110     register_stat(s_avg_read_latency).name("avg_read_latency_{}", m_channel_id);
111 }
```

2.4. Frontend에서 core 수를 조회하여 m\_num\_cores에 저장

2.5. Read Row Hit Per Core, ... , Read Row Conflict Per Core 등의 배열 초기화 → multi core 지원을 위해 초기화

→ 위 배열들은 통계를 수집하고 출력 용도로만 사용됨

2.6. 다양한 통계 variable를 register하여 DRAM controller의 동작을 분석 → 통계를 수집하고 출력

# src/dram\_controller folder

## □ DRAM Controller Operations

- class GenericDRAMController – send()

```
113 bool send(Request& req) override {
114     req.final_command = m_dram->m_request_translations(req.type_id);
115
116     switch (req.type_id) {
117     case Request::Type::Read: {
118         s_num_read_reqs++;
119         break;
120     }
121     case Request::Type::Write: {
122         s_num_write_reqs++;
123         break;
124     }
125     default: {
126         s_num_other_reqs++;
127         break;
128     }
129 }
130
131 // Forward existing write requests to incoming read requests
132 if (req.type_id == Request::Type::Read) {
133     auto compare_addr = [req](const Request& wreq) {
134         return wreq.addr == req.addr;
135     };
136     if (std::find_if(m_write_buffer.begin(), m_write_buffer.end(), compare_addr) != m_write_buffer.end()) {
137         // The request will depart at the next cycle
138         req.depart = m_clk + 1;
139         pending.push_back(req);
140         return true;
141     }
142 }
143
144 // Else, enqueue them to corresponding buffer based on request type id
145 bool is_success = false;
146 req.arrive = m_clk;
147 if (req.type_id == Request::Type::Read) {
148     is_success = m_read_buffer.enqueue(req);
149 } else if (req.type_id == Request::Type::Write) {
150     is_success = m_write_buffer.enqueue(req);
151 } else {
152     throw std::runtime_error("Invalid request type!");
153 }
154 if (!is_success) {
155     // We could not enqueue the request
156     req.arrive = -1;
157     return false;
158 }
159
160 return true;
161 };
```

# src/dram\_controller folder

□ **dram\_controller**

- **abcd**

# src/base folder

□ **base**

- **abcd**