### **Lecture 13**

- **4.1.4 Simplified Turing Machine Models**
- **4.1.5 Generating Languages Using Turing Machines**



### **Lecture 13**

- **4.1.4 Simplified Turing Machine Models**
- **4.1.5 Generating Languages Using Turing Machines**



# **Simplified TM Models**

- Stack machine
- Counter machines
- TM with a limited number of states and tape symbols
- Universal Turing Machine





Turing machine



### Turing machine



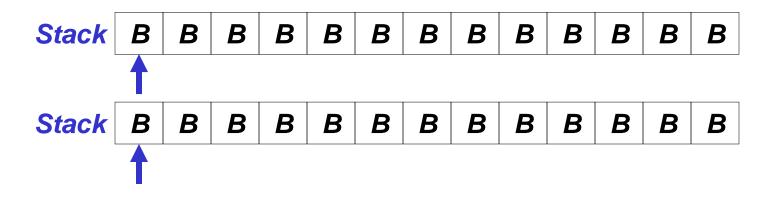
### Turing machine



### Turing machine

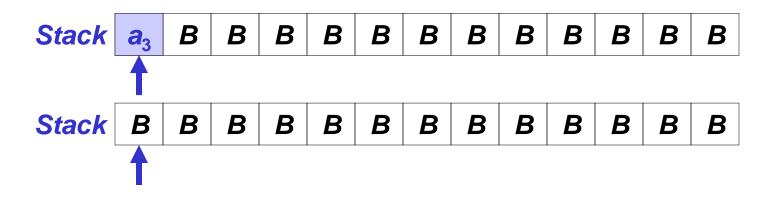


### Turing machine



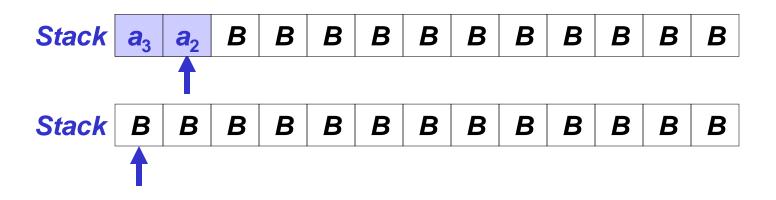


### Turing machine





### Turing machine

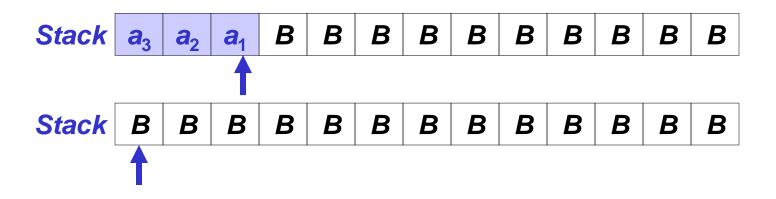




### Turing machine

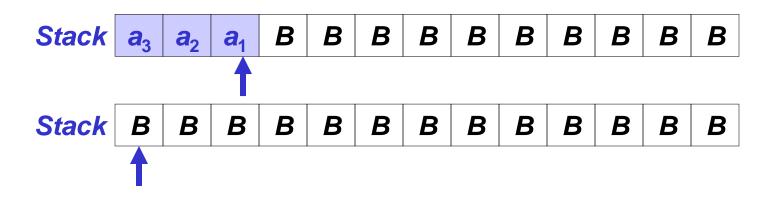


### Turing machine



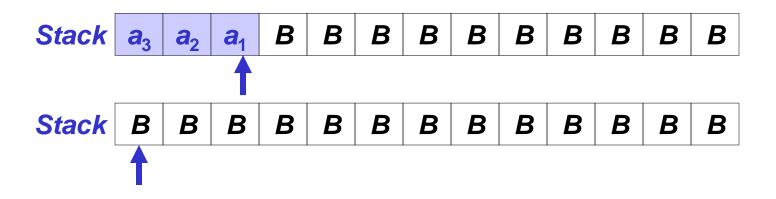


### Turing machine



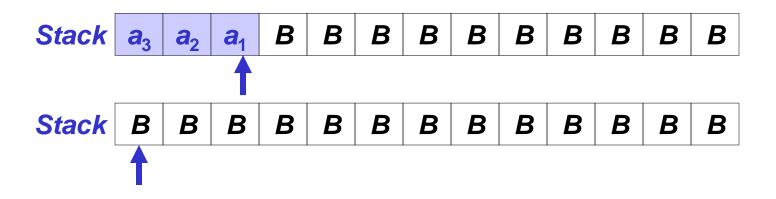


### Turing machine



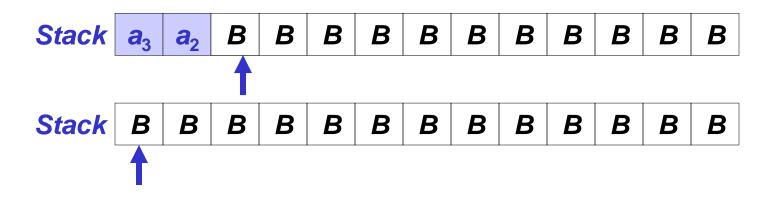


### Turing machine





### Turing machine

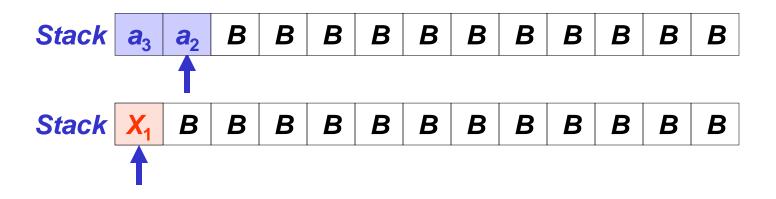




### Turing machine



### Turing machine

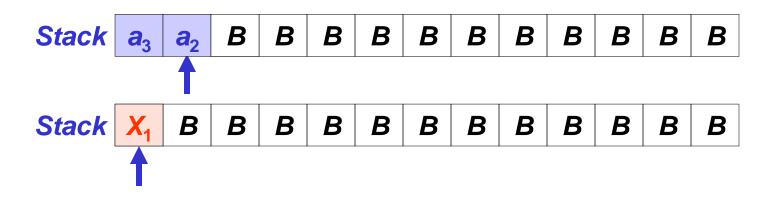




### Turing machine

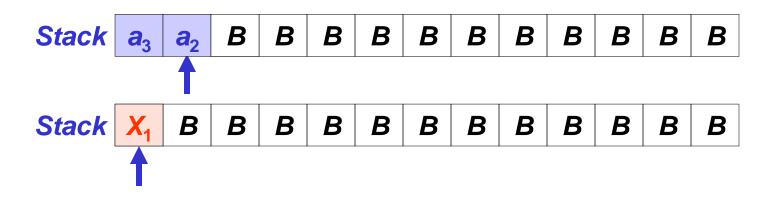


### Turing machine



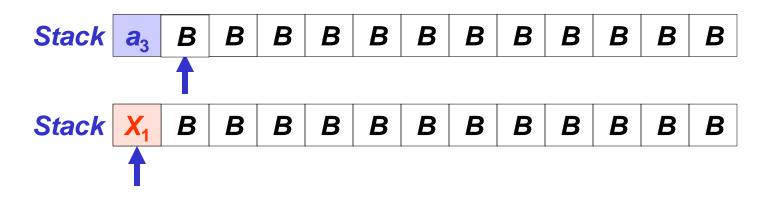


### Turing machine



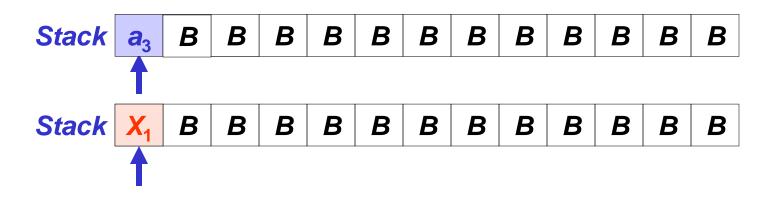


### Turing machine



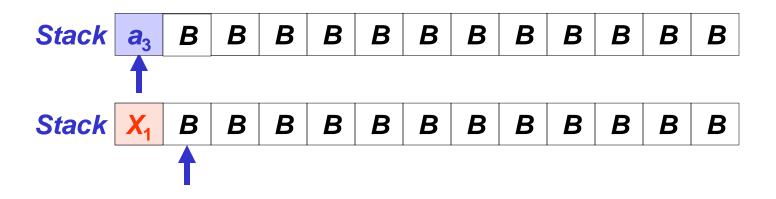


### Turing machine



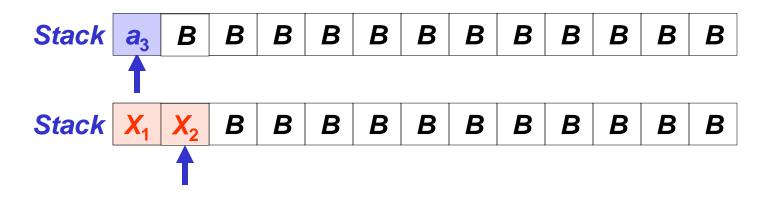


### Turing machine



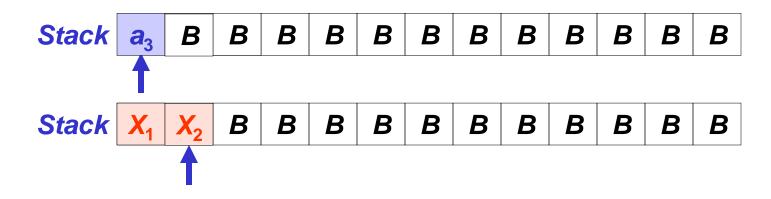


### Turing machine



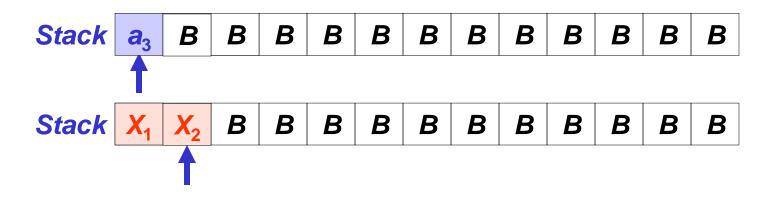


### Turing machine



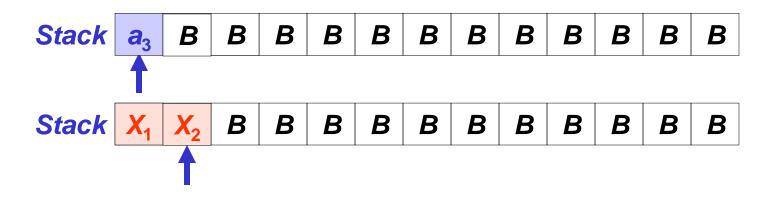


### Turing machine



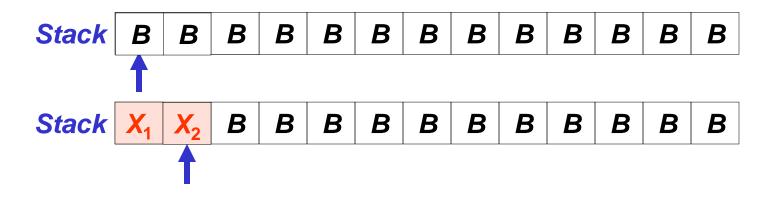


### Turing machine



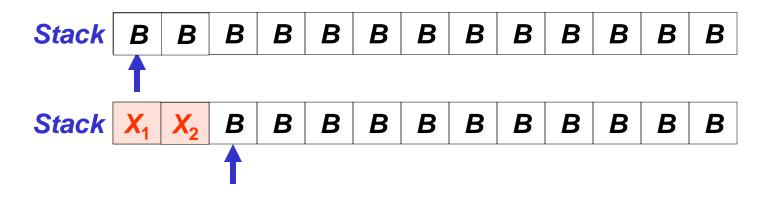


### Turing machine



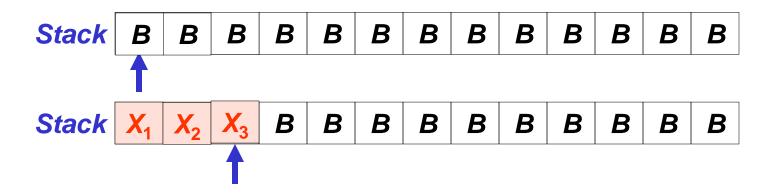


### Turing machine



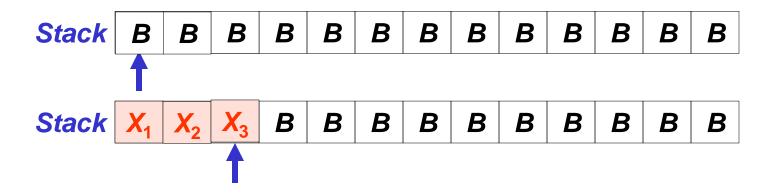


### Turing machine



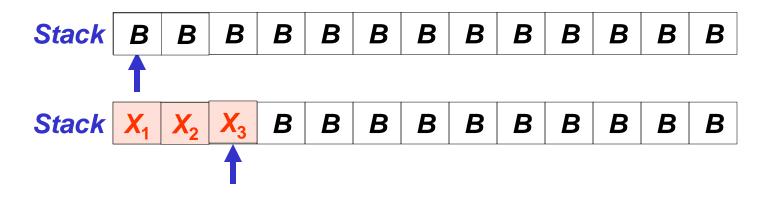


### Turing machine



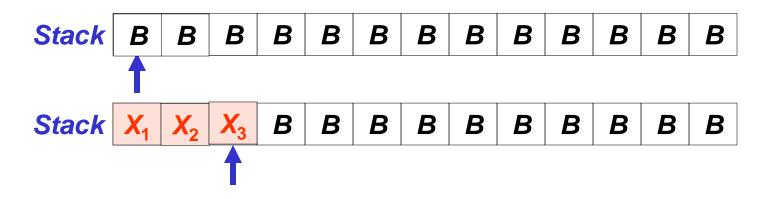


### Turing machine



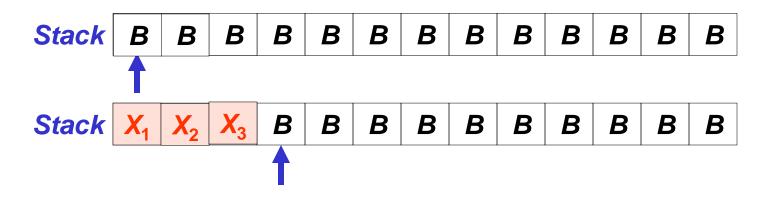


### Turing machine



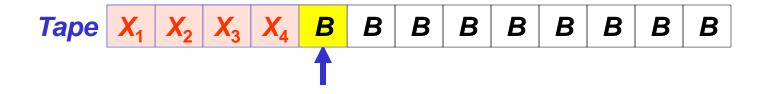


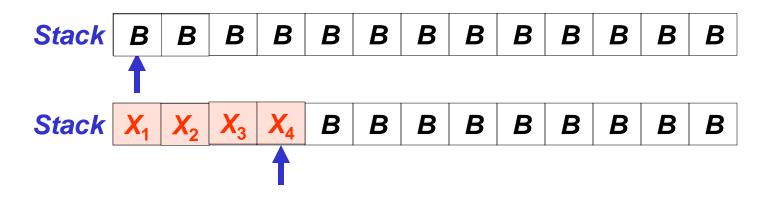
### Turing machine





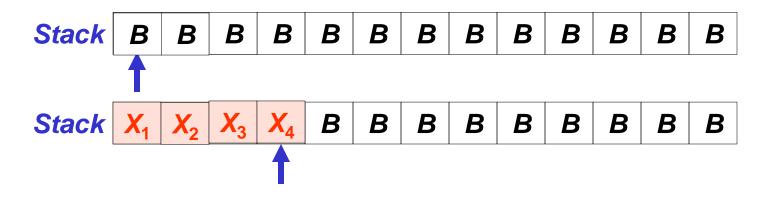
### Turing machine





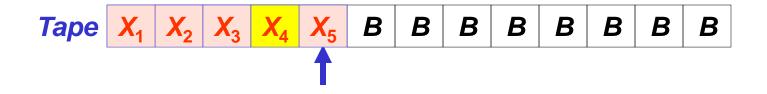


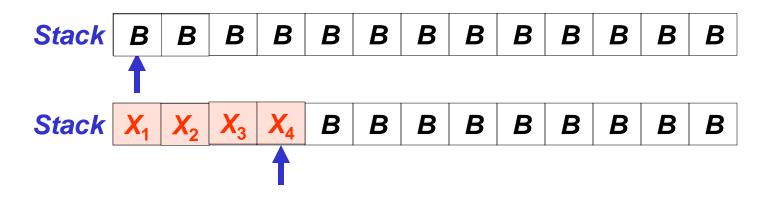
### Turing machine





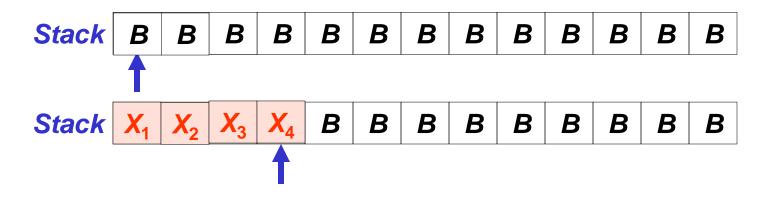
### Turing machine





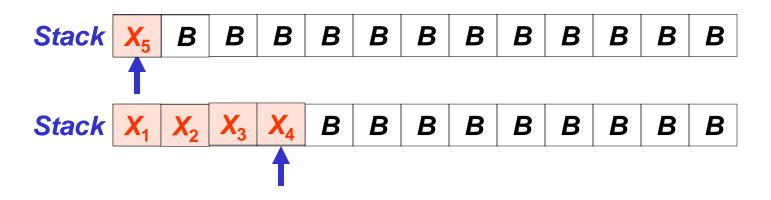


### Turing machine



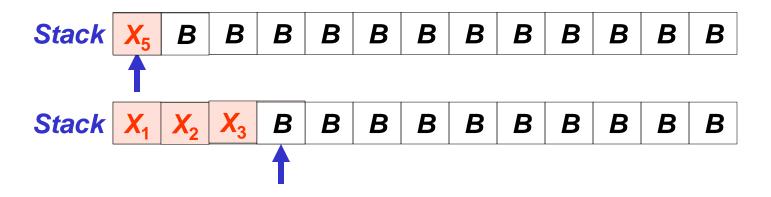


### Turing machine



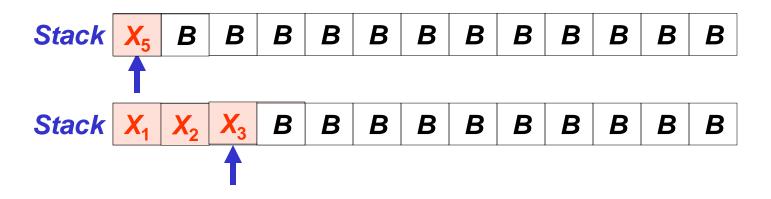


### Turing machine





### Turing machine

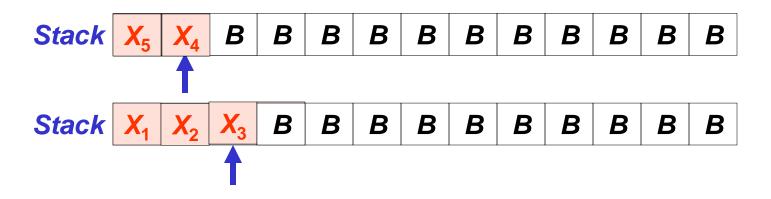




### Turing machine



### Turing machine

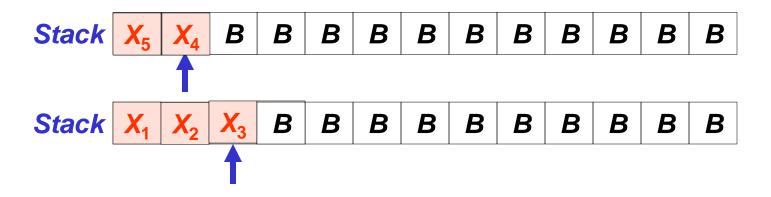




### Turing machine



### Turing machine





### Turing machine



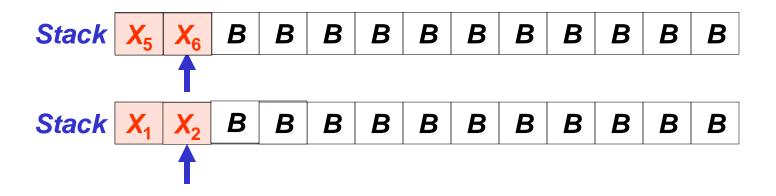
### Turing machine



### Turing machine

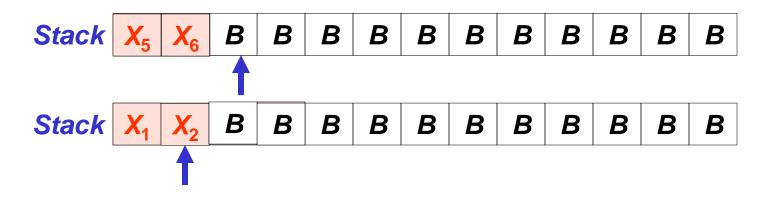


### Turing machine



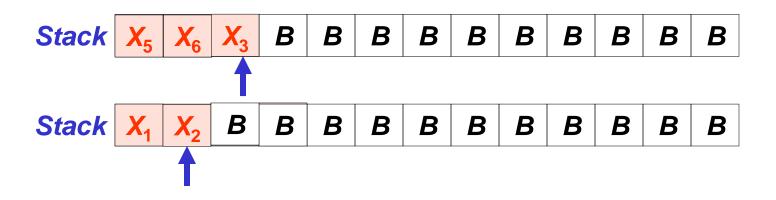


### Turing machine



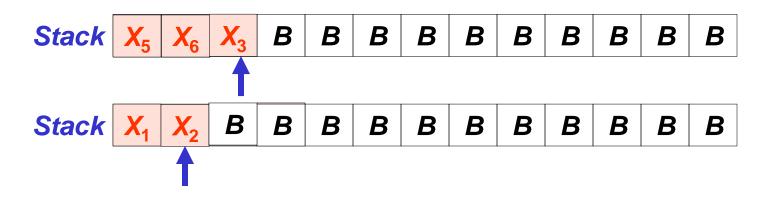


### Turing machine



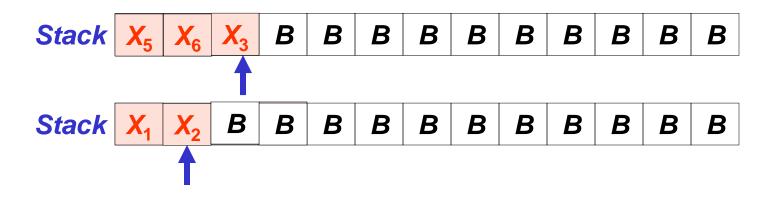


### Turing machine



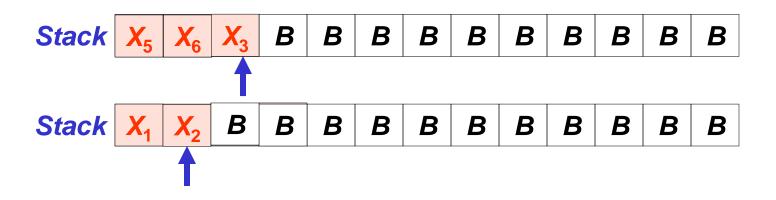


### Turing machine





### Turing machine

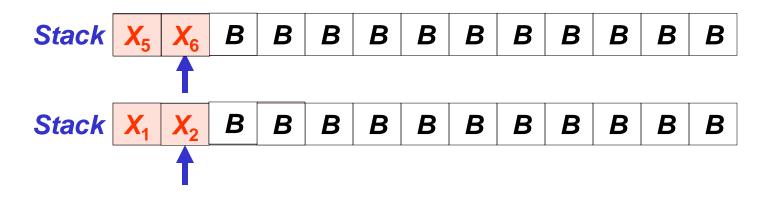




### Turing machine

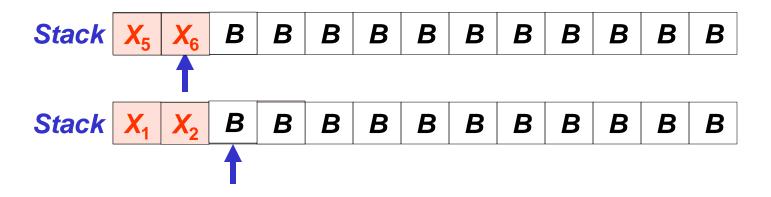


### Turing machine



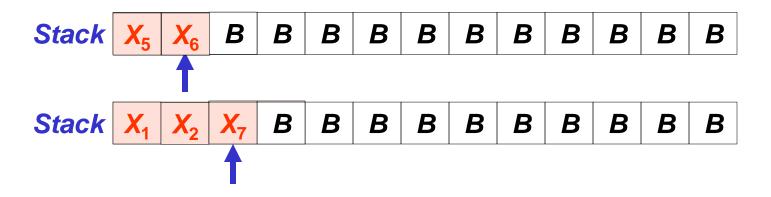


### Turing machine



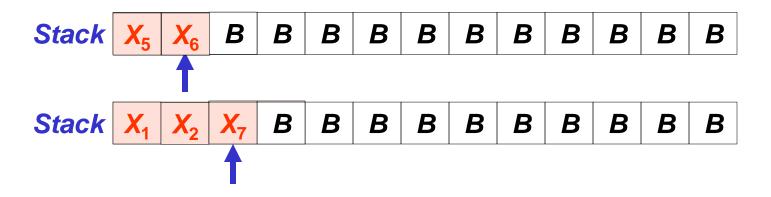


### Turing machine



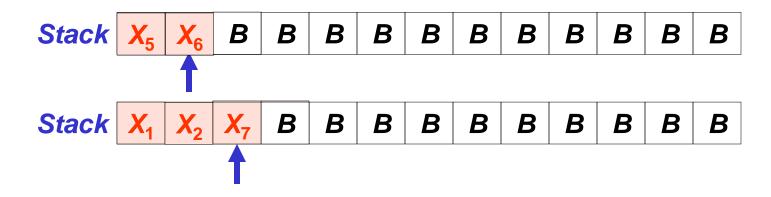


### Turing machine





### Turing machine

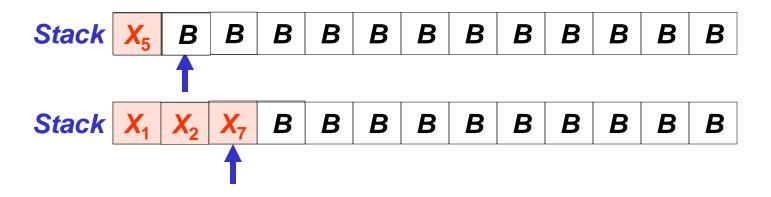




### Turing machine

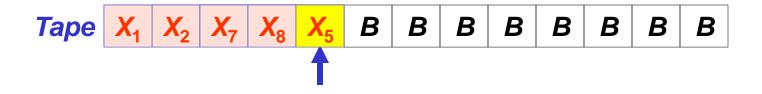


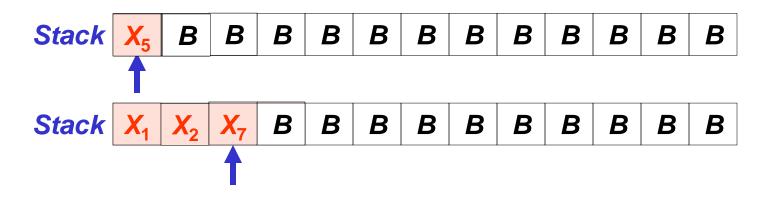
### Turing machine





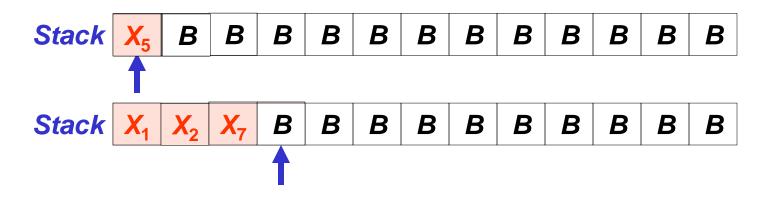
### Turing machine





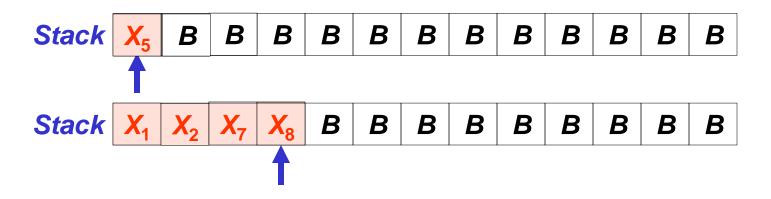


### Turing machine





### Turing machine











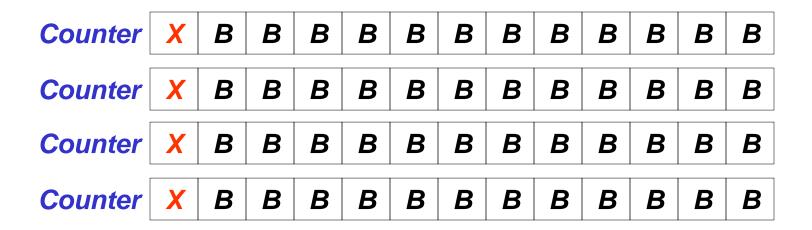
#### Stack machine

#### Counter machine



#### Stack machine

#### Counter machine







Stack machine



Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B



Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B

**Counter Machine** 



### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### **Counter Machine**

Counter X B B B B B B B B B B B B B B B

Counter



### Stack machine

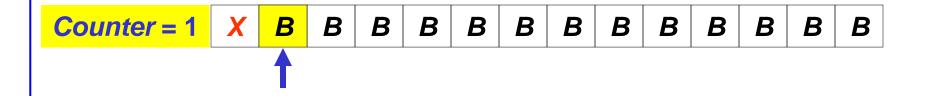
### **Counter Machine**



Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

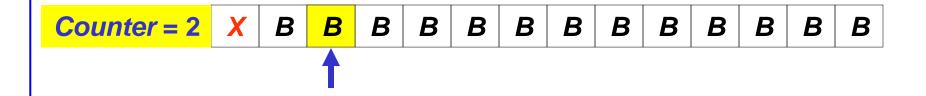
### **Counter Machine**





### Stack machine

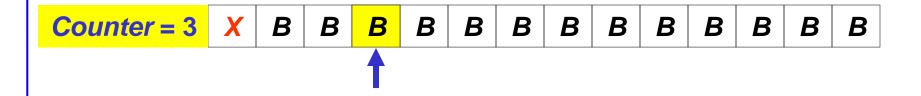
### **Counter Machine**





### Stack machine

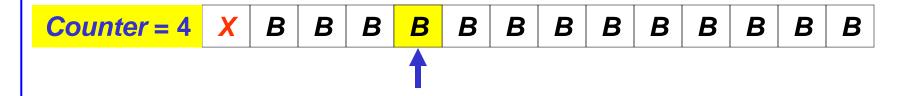
#### **Counter Machine**





### Stack machine

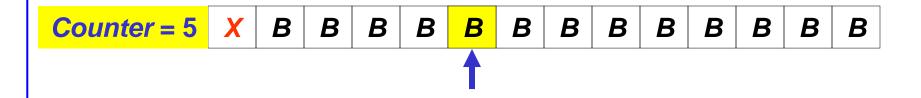
### **Counter Machine**





### Stack machine

### **Counter Machine**



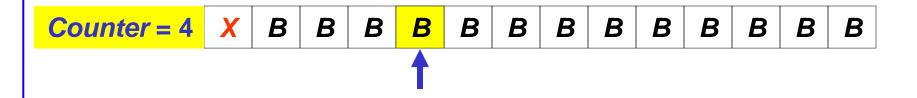
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### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### **Counter Machine**

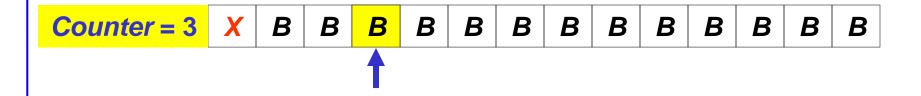




### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### **Counter Machine**





### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Counter

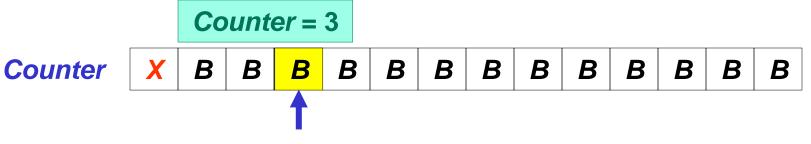
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### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine



[ q, p, Counter to k]



#### Stack machine

#### Counter machine



### Stack machine

### Counter machine



### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Counter = 3



[ q, p, 2]



### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Counter = 3



[ q, p, 1]



### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Counter = 3



[ q, p, 0]



### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

## Counter machine

Counter В B B B B B B B B B B B B B B B

Counter B B B B B B B B B B B B B B B B



### Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B

### Counter machine

**k** distinct stack symbols

Counter B B B B B B B B B B B B B B B B

B B B B B B B B B B B B B B Counter B B



### Stack machine

### Counter machine

k distinct stack symbols

$$Z_0, Z_1, Z_2, ..., Z_{k-1}$$

**Counter** B B B B B B B B B B B B B B B B

B B B B B B B B B B B **Counter** B B B B B



### Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B

### Counter machine

k distinct stack symbols

$$Z_0, Z_1, Z_2, ..., Z_{k-1}$$

symbol  $Z_i$  is coded as integer i,  $0 \le i \le k-1$ 

**Counter** 



Counter





### Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B

### Counter machine

k distinct stack symbols

$$Z_0, Z_1, Z_2, ..., Z_{k-1}$$

symbol  $Z_i$  is coded as integer i,  $0 \le i \le k-1$ 

**Counter** 



Counter





### Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B

### Counter machine

k distinct stack symbols

$$Z_0, Z_1, Z_2, ..., Z_{k-1}$$

symbol  $Z_i$  is coded as integer i,  $0 \le i \le k-1$ 

Counter = 
$$j$$
  $X$   $j = i_m + k i_{m-1} + k^2 i_{m-2} + k^3 i_{m-3} + ... + k^{m-1} i_1$   $B$   $B$ 



### Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B

## Counter machine

Counter В B B B B B B B B B B B B B B B

B Counter B B B B B B B B B B B B B B B



### Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B

### Counter machine

Putting symbol  $Z_r$  on the top of the stack

Counter B B B B B B B B B B B B B B B B

B B B B B B B B B B B B B B **Counter** B B



### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Putting symbol  $Z_r$  on the top of the stack Before: stack symbols' value = j

Counter B B B B B B B B B B B B B B B B

B B B B B B B B B **Counter** B B B B B B B



### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Putting symbol  $Z_r$  on the top of the stack

Before: stack symbols' value = j

After: stack symbols' value = j k + r

Counter B B B B B B B B B B B B B B B B

B B B B B B B B B **Counter** B B B B B B B

### Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B

### Counter machine

Putting symbol  $Z_r$  on the top of the stack

Before: stack symbols' value = j

After: stack symbols' value = j k + r

Counter B B B B B B B B B B B B B B B B

[ q, ..., Counter k, ..., Counter r , ...]

B B Counter B B B B B B B B B B B B B B



### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Putting symbol  $Z_r$  on the top of the stack

Before: stack symbols' value = j

After: stack symbols' value = j k + r

Counter B B B B B B B B B B B B B B B B

[ q, ..., Counter k, ..., Counter r , ...]

B B Counter B B B B B B B B B B B B B B



### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Putting symbol  $Z_r$  on the top of the stack

Before: stack symbols' value = j

After: stack symbols' value = j k + r

Counter X j B B B B

[ q, ..., Counter k, ..., Counter r , ...]

Counter B B B B B B B B B B B B B B B B



### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Putting symbol  $Z_r$  on the top of the stack

Before: stack symbols' value = j

After: stack symbols' value = j k + r

 Counter
 X
 j-1
 B
 B
 B
 B
 B
 B

[ q, ..., Counter k, ..., Counter r , ...]

Counter B B B B B B B B B B B B B B B B

### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Putting symbol  $Z_r$  on the top of the stack

Before: stack symbols' value = j

After: stack symbols' value = j k + r

 Counter
 X
 j - 1
 B
 B
 B
 B
 B
 B

[ q, ..., Counter k, ..., Counter r , ...]

k B B B Counter B B B B B B B B B B B



### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine

Putting symbol  $Z_r$  on the top of the stack

Before: stack symbols' value = j

After: stack symbols' value = j k + r

 Counter
 X
 j - 2
 B
 B
 B
 B
 B
 B
 B

[ q, ..., Counter k, ..., Counter r , ...]

k B B B Counter B B B B B B B B B B B

## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine

Putting symbol  $Z_r$  on the top of the stack

Before: stack symbols' value = j

After: stack symbols' value = j k + r

 Counter
 X
 j - 2
 B
 B
 B
 B
 B
 B
 B

[ q, ..., Counter k, ..., Counter r , ...]

 Counter
 X
 2k = k + k
 B
 B
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## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine

Putting symbol  $Z_r$  on the top of the stack

Before: stack symbols' value = j

After: stack symbols' value = j k + r

[ q, ..., Counter k, ..., Counter r , ...]

 Counter
 X
 2k = k + k
 B
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## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine

Putting symbol  $Z_r$  on the top of the stack

Before: stack symbols' value =

After: stack symbols' value = j k + r

[ q, ..., Counter k, ..., Counter r , ...]

 Counter
 X
 3k = 2k + k
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## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine

Putting symbol  $Z_r$  on the top of the stack

Before: stack symbols' value = j

After: stack symbols' value = j k + r

Counter B B B B B B B B B B B B B B B B

[ q, ..., Counter k, ..., Counter r , ...]

Counter X j k = (j-1)k + k B B



### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine

Putting symbol  $Z_r$  on the top of the stack

Before: stack symbols' value = j

After: stack symbols' value = j k + r

Counter B B B B B B B B B B B B B B B B

[ q, ..., Counter k, ..., Counter r , ...]

Counter X j k = (j-1)k + k B B



### Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B

### Counter machine

Putting symbol  $Z_r$  on the top of the stack

Before: stack symbols' value = j

After: stack symbols' value = j k + r

Counter B B B B B B B B B B B B B B B B

[ q, ..., Counter k, ..., Counter r , ...]

Counter X jk+r B



## Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B

### Counter machine



[ q, ..., Counter k, ..., Counter r, ...]



## Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B

#### Counter machine

Removing symbol  $Z_r$  from the top of the stack



[ q, ..., Counter k, ..., Counter r, ...]

B B B B B B B Counter B B B B B B B B B



## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Removing symbol  $Z_r$  from the top of the stack Before: stack symbols' value =



[ q, ..., Counter k, ..., Counter r, ...]



## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine

Removing symbol  $Z_r$  from the top of the stack

Before: stack symbols' value = j

After: stack symbols' value =  $\lfloor j/k \rfloor$ 

Counter B B B B B B B B B B B B B B B B

[ q, ..., Counter k, ..., Counter r, ...]



## Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B

### Counter machine

Removing symbol  $Z_r$  from the top of the stack

Before: stack symbols' value = j

After: stack symbols' value =  $\lfloor j/k \rfloor$ 

Counter B B B B B B B B B B B B B B B B

[ q, ..., Counter k, ..., Counter r, ...]



## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Removing symbol  $Z_r$  from the top of the stack

Before: stack symbols' value =

After: stack symbols' value =  $\lfloor j/k \rfloor$ 

Counter X j B B B B

[ q, ..., Counter k, ..., Counter r, ...]



## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Removing symbol  $Z_r$  from the top of the stack

Before: stack symbols' value =

After: stack symbols' value =  $\lfloor j/k \rfloor$ 

Counter X j-k B B B B B

[ q, ..., Counter k, ..., Counter r, ...]



## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine

Removing symbol  $Z_r$  from the top of the stack

Before: stack symbols' value =

After: stack symbols' value =  $\lfloor j/k \rfloor$ 

Counter X j-k B B B B B

[ q, ..., Counter k, ..., Counter r, ...]



## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Removing symbol  $Z_r$  from the top of the stack

Before: stack symbols' value =

After: stack symbols' value =  $\lfloor j/k \rfloor$ 

 Counter
 X
 j - 2 k
 B
 B
 B
 B
 B
 B
 B
 B

[ q, ..., Counter k, ..., Counter r, ...]

## Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B

### Counter machine

Removing symbol  $Z_r$  from the top of the stack

Before: stack symbols' value =

After: stack symbols' value =  $\lfloor j/k \rfloor$ 

 Counter
 X
 j - 2 k
 B
 B
 B
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 B

[ q, ..., Counter k, ..., Counter r, ...]



## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine

Removing symbol  $Z_r$  from the top of the stack

Before: stack symbols' value =

After: stack symbols' value =  $\lfloor j/k \rfloor$ 

[ q, ..., Counter k, ..., Counter r, ...]

## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine

Removing symbol  $Z_r$  from the top of the stack

Before: stack symbols' value = j

After: stack symbols' value =  $\lfloor j/k \rfloor$ 

[ q, ..., Counter k, ..., Counter r, ...]

 Counter
 X
 3
 B
 B
 B
 B
 B
 B
 B
 B
 B
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## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Removing symbol  $Z_r$  from the top of the stack

Before: stack symbols' value =

After: stack symbols' value =  $\lfloor j/k \rfloor$ 

Counter B B B B B B B B B B B B B B B B

[ q, ..., Counter k, ..., Counter r, ...]



## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine



[ q, ..., Counter k, ..., Counter r, ..., Mod k counter ]



## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack



[ q, ..., Counter k, ..., Counter r, ..., Mod k counter ]



## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$

Counter B B B B B B B B B B B B B B B B

[ q, ..., Counter k, ..., Counter r, ..., Mod k counter ]



## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 



[ q, ..., Counter k, ..., Counter r, ..., Mod k counter ]



### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 



[ q, ..., Counter k, ..., Counter r, . Mod k counter r ]



## Stack machine

### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 



[ q, ..., Counter k, ..., Counter r, . 0 :r]



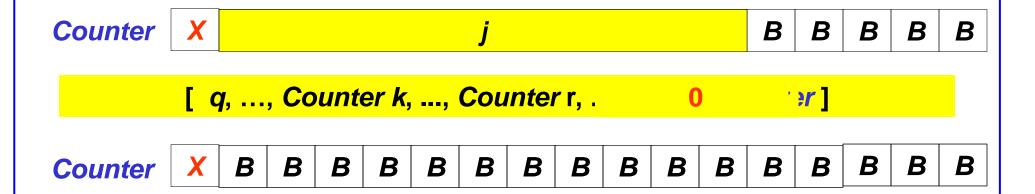
### Stack machine

Stack Z<sub>i1</sub> Z<sub>i2</sub> Z<sub>i3</sub> Z<sub>i4</sub> -- Z<sub>im</sub> B B B B B B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 





### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 



[ q, ..., Counter k, ..., Counter r, . 0 :r]



### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 



[ q, ..., Counter k, ..., Counter r, . 1



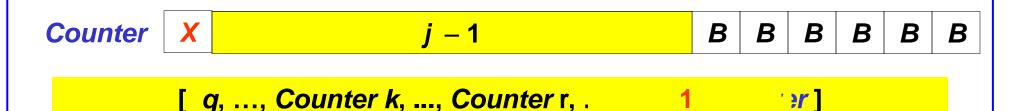
## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 



 Counter
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### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 



[ q, ..., Counter k, ..., Counter r, . 1

 Counter
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## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 



[ q, ..., Counter k, ..., Counter r, . 2

 Counter
 X
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## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 



[ q, ..., Counter k, ..., Counter r, . 2

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## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
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[ q, ..., Counter k, ..., Counter r, . 2

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### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine

Determining symbol at the top of the stack

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[ q, ..., Counter k, ..., Counter r, . 3

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### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine

Determining symbol at the top of the stack

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## Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

### Counter machine

Determining symbol at the top of the stack

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[ 
$$q, ..., Counter k, ..., Counter r, .  $k-3$$$



#### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 



[ 
$$q, ..., Counter k, ..., Counter r, .  $k-2$$$



#### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 



[ q, ..., Counter k, ..., Counter r, . <math>k-1



#### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 



[ q, ..., Counter k, ..., Counter r, . 0



#### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 





#### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 





#### Stack machine

Stack | Z<sub>i1</sub> | Z<sub>i2</sub> | Z<sub>i3</sub> | Z<sub>i4</sub> | -- | Z<sub>im</sub> | B | B | B | B | B | B

#### Counter machine

Determining symbol at the top of the stack

$$j = i_m + ki_{m-1} + k^2i_{m-2} + k^3i_{m-3} + \dots + k^{m-1}i_1 = i_m + k(i_{m-1} + ki_{m-2} + \dots + k^{m-2}i_1)$$
  
value of the symbol at the top of the stack=  $j \mod k$ 



[ q, ..., Counter k, ..., Counter r, . 2

Counter X j







Counter	X	В	В	В	В	В	В	В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В



### Machine with 4 counters

Counter	X	В	В	В	В	В	В	В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В



#### Machine with 4 counters

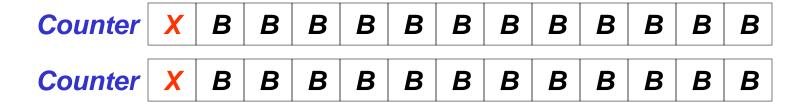
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В





#### Machine with 4 counters

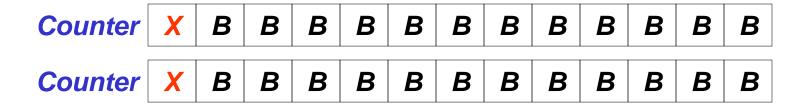
Counter	X			i		В	В	В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В





#### Machine with 4 counters

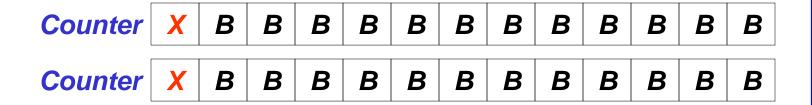
Counter	X			i		В	В	В	В	В	В	В	В
Counter	X			J	i			В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В





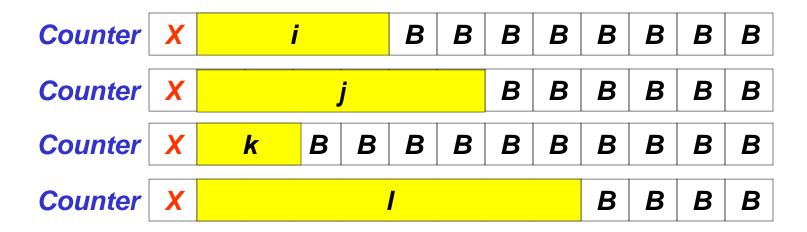
#### Machine with 4 counters

Counter	X			i		В	В	В	В	В	В	В	В
Counter	X			J	i			В	В	В	В	В	В
Counter	X		k	В	В	В	В	В	В	В	В	В	В
Counter	X	В	В	В	В	В	В	В	В	В	В	В	В





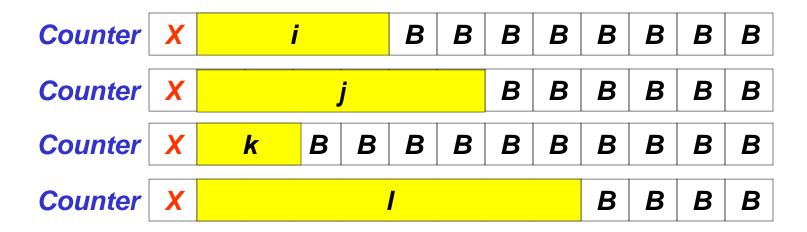
#### Machine with 4 counters

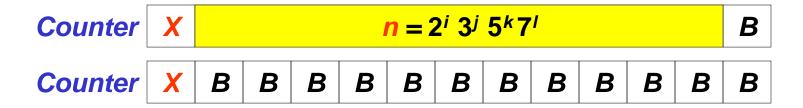






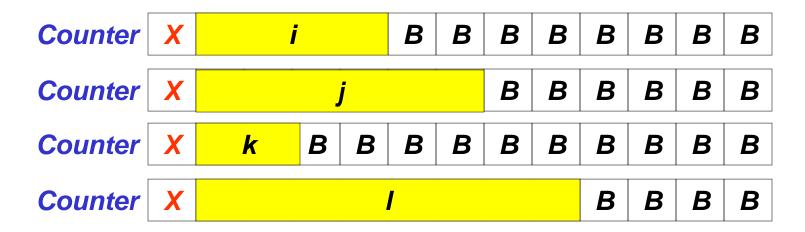
#### Machine with 4 counters







#### Machine with 4 counters



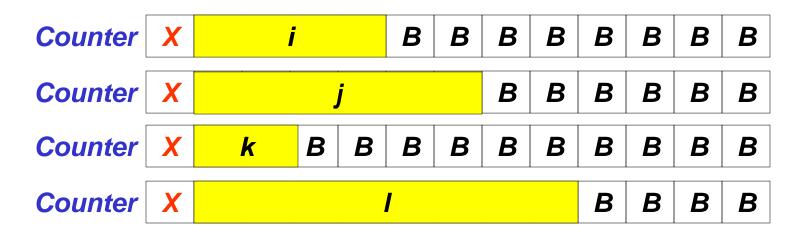
Counter X 
$$n = 2^i \ 3^j \ 5^k 7^l$$
 B

Counter X B B B B B B B B B B B B

1) 
$$(i, j, k, l) = \pm 1$$



#### Machine with 4 counters



#### Machine with 2 counters

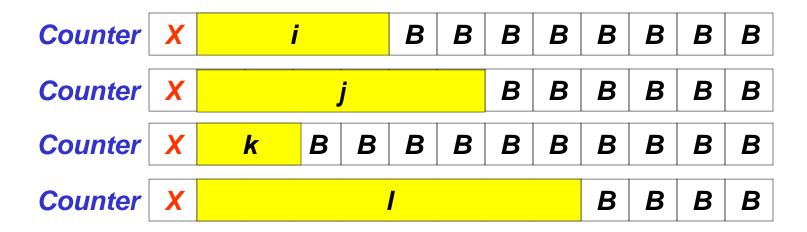
CounterX
$$n = 2^i \ 3^j \ 5^k 7^l$$
BCounterXBBBBBBBBBB

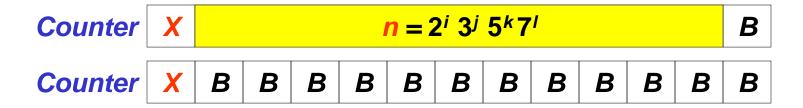
1) 
$$(i, j, k, l) = \pm 1$$

- multiply or divide value *n* by an appropriate value 2, 3, 5 or 7



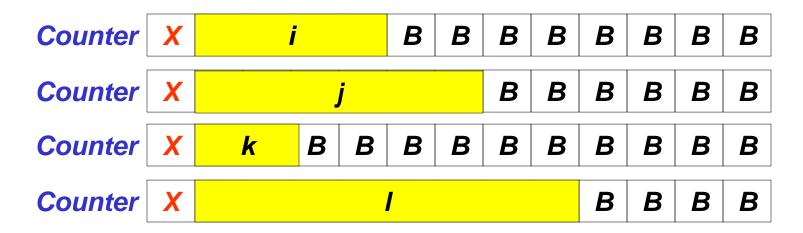
#### Machine with 4 counters







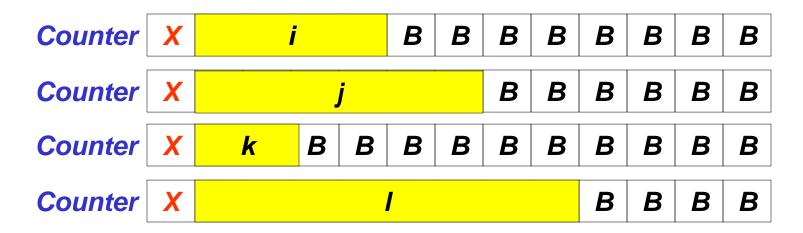
#### Machine with 4 counters



2) 
$$(i, j, k, l) == 0$$
?



#### Machine with 4 counters

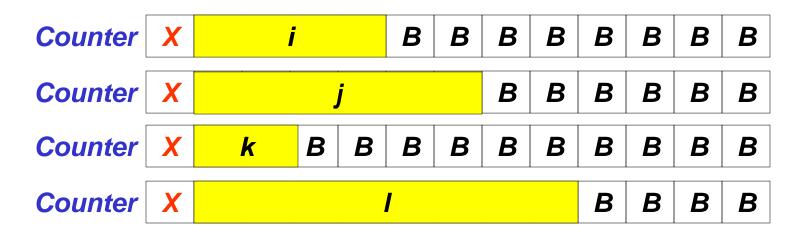


CounterX
$$n = 2^i \ 3^j \ 5^k 7^l$$
BCounterXBBBBBBBBBB

2) 
$$(i, j, k, l) == 0$$
?  
- divide  $n$  by 2, 3, 5 or 7



### Machine with 4 counters



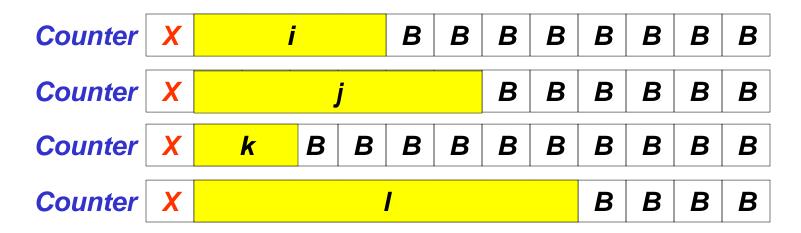
CounterX
$$n = 2^i \ 3^j \ 5^k 7^l$$
BCounterXBBBBBBBBBB

2) 
$$(i, j, k, l) == 0$$
?

- divide *n* by 2, 3, 5 or 7
- (division remainder == 0) ?



### Machine with 4 counters



2) 
$$(i, j, k, l) == 0$$
?

- divide *n* by 2, 3, 5 or 7
- (division remainder == 0) ?

$$(i, j, k, l) \neq 0$$





Acceptance of any language



- Acceptance of any language
- Unlimited number of tape symbols



- Acceptance of any language
- Unlimited number of tape symbols
  - Single tape



- Acceptance of any language
- Unlimited number of tape symbols
  - Single tape
  - Three states



- Acceptance of any language
- Unlimited number of tape symbols
  - Single tape
  - Three states
    - —Two non accepting states



- Acceptance of any language
- Unlimited number of tape symbols
  - Single tape
  - Three states
    - —Two non accepting states
    - —One accepting state





Unlimited number of states



- Unlimited number of states
  - tape symbols {0, 1, *B*}



- Unlimited number of states
  - tape symbols {0, 1, B}
  - 2<sup>k-1</sup>+1 < cardinal number of tape symbol set < 2<sup>k</sup>



## Unlimited number of states

- tape symbols {0, 1, B}
- $2^{k-1}+1 < cardinal\ number\ of\ tape\ symbol\ set < 2^k$
- tape symbols are coded by binary strings of length k



## Unlimited number of states

- tape symbols {0, 1, B}
- 2<sup>k-1</sup>+1 < cardinal number of tape symbol set < 2<sup>k</sup>
- tape symbols are coded by binary strings of length k
- k tape cells contain k-bit binary code



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- $2^{k-1}+1 < cardinal\ number\ of\ tape\ symbol\ set < 2^k$
- tape symbols are coded by binary strings of length k
- k tape cells contain k-bit binary code
- one state component state TM M<sub>1</sub>



- tape symbols {0, 1, *B*}
- $2^{k-1}+1 < cardinal\ number\ of\ tape\ symbol\ set < 2^k$
- tape symbols are coded by binary strings of length k
- k tape cells contain k-bit binary code
- one state component state TM M<sub>1</sub>
- another state component sets head at the beginning of the binary code



- tape symbols {0, 1, B}
- 2<sup>k-1</sup>+1 < cardinal number of tape symbol set < 2<sup>k</sup>
- tape symbols are coded by binary strings of length k
- k tape cells contain k-bit binary code
- one state component state TM M<sub>1</sub>
- another state component sets head at the beginning of the binary code
- empty cell symbol B



- tape symbols {0, 1, *B*}
- 2<sup>k-1</sup>+1 < cardinal number of tape symbol set < 2<sup>k</sup>
- tape symbols are coded by binary strings of length k
- k tape cells contain k-bit binary code
- one state component state TM M<sub>1</sub>
- another state component sets head at the beginning of the binary code
- empty cell symbol B
  - —TM  $M_2$  changes k empty cell symbols B with binary code for an empty cell



- tape symbols {0, 1, *B*}
- 2<sup>k-1</sup>+1 < cardinal number of tape symbol set < 2<sup>k</sup>
- tape symbols are coded by binary strings of length k
- k tape cells contain k-bit binary code
- one state component state TM M<sub>1</sub>
- another state component sets head at the beginning of the binary code
- empty cell symbol B
  - —TM  $M_2$  changes k empty cell symbols B with binary code for an empty cell
  - —TM  $M_1$  simulation continues





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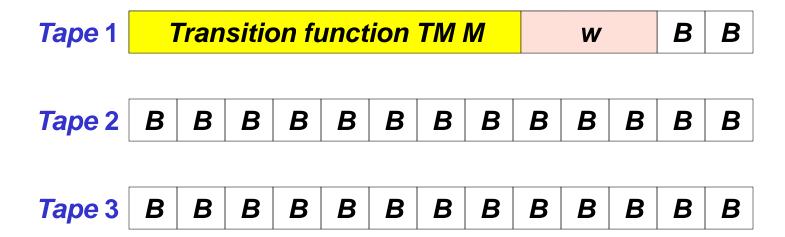


 Tape 1
 Transition function TM M
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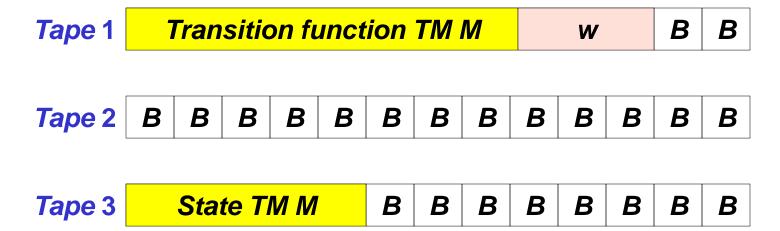
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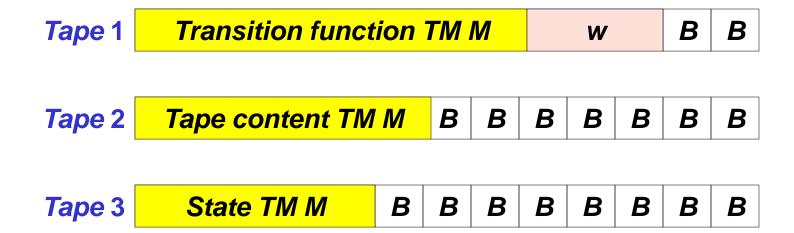




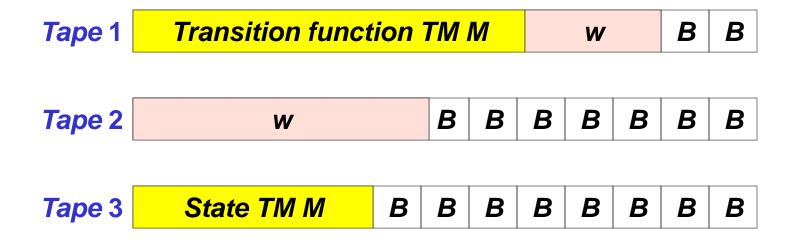




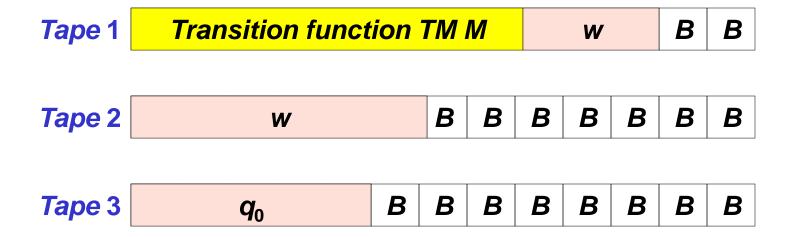






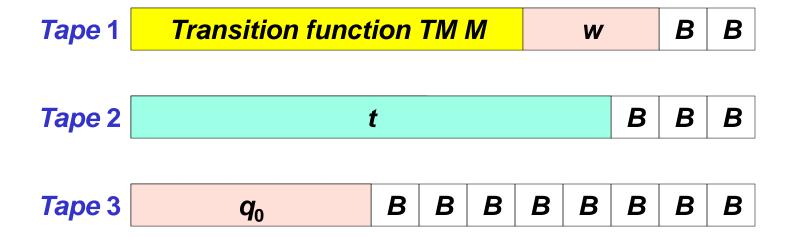






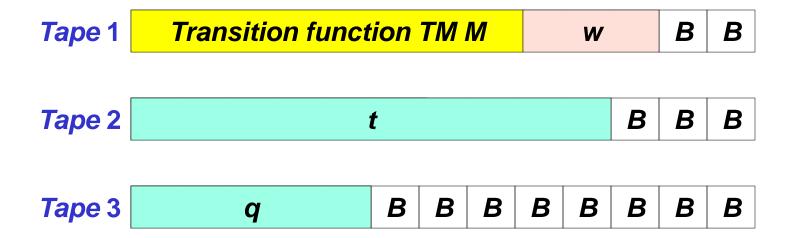


# Universal Turing Machine $M_u$

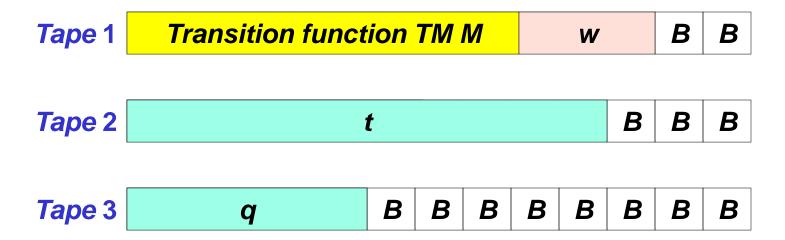




# Universal Turing Machine $M_u$

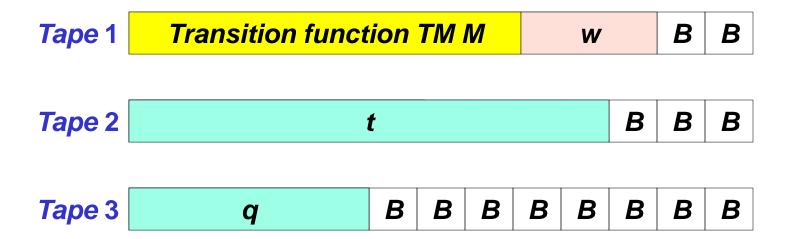






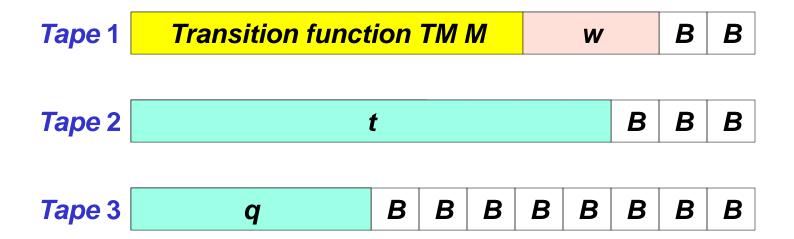
• It is possible to construct universal TM  $M_u$  using:





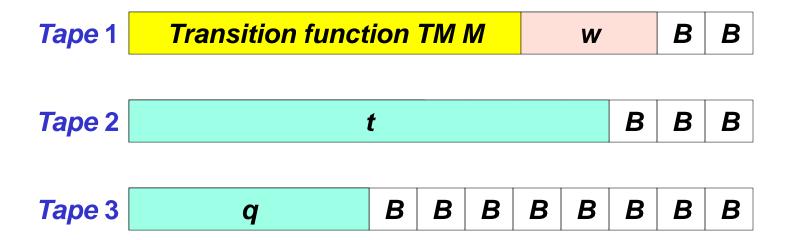
- It is possible to construct universal TM  $M_u$  using:
  - Single tape





- It is possible to construct universal TM  $M_u$  using:
  - Single tape
  - Five states





- It is possible to construct universal TM  $M_u$  using:
  - Single tape
  - Five states
  - Five tape symbols



#### **Lecture 13**

- **4.1.4 Simplified Turing Machine Models**
- **4.1.5 Generating Languages Using Turing Machines**















 Output tape
 # w<sub>1</sub> #
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Output tape #W W<sub>1</sub> # B B B B B B B B B B B B B















Output tape # w<sub>1</sub> # w<sub>2</sub> # w<sub>3</sub> w # B B B B B B B B

























Output tape # w<sub>1</sub> # w<sub>2</sub> # w<sub>3</sub> # w<sub>4</sub> # w<sub>5</sub> # w<sub>6</sub> # w<sub>7</sub> #



Output tape # w<sub>1</sub> # w<sub>2</sub> # w<sub>3</sub> # w<sub>4</sub> # w<sub>5</sub> # w<sub>6</sub> # w<sub>7</sub> w #





















Output tape B B B B B B B B B B B B B B B B B















































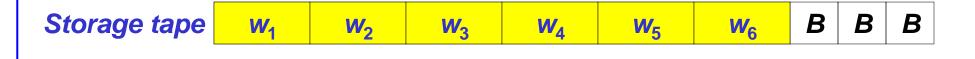


Output tape # W<sub>2</sub> # W<sub>3</sub> # B B B B B B B B



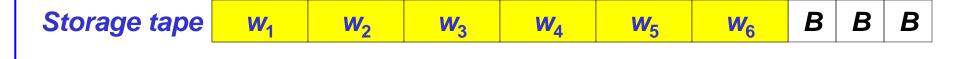






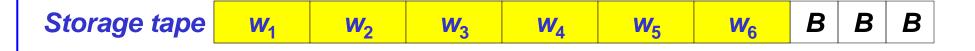
















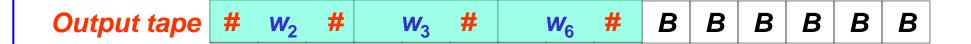




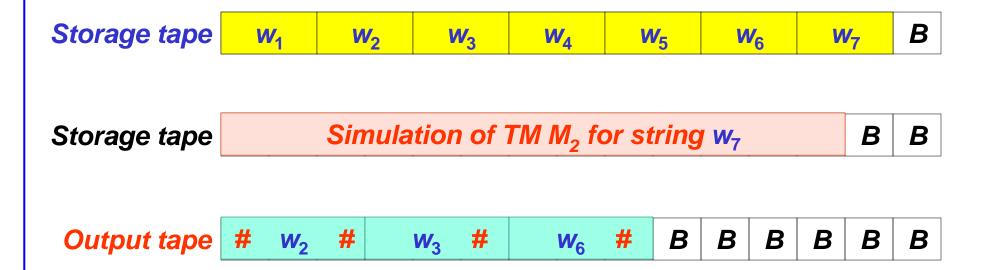






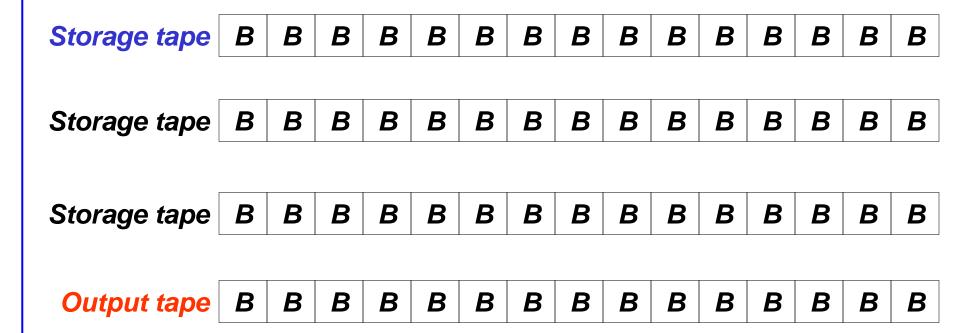






# Possible only for recursive languages







Storage tape	В	В	В	В	В	В	В	В	В	В	В	В	В	В	В
Storage tape	В	В	В	В	В	В	В	В	В	В	В	В	В	В	В
Storage tape	В	В	В	В	В	В	В	В	В	В	В	В	В	В	В
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Output tano	R	R	R	R	R	R	R	R	R	R	R	R	R	R	R



Storage tape	В	В	В	В	В	В	В	В	В	В	В	В	В	В	В
Storage tape	В	В	В	В	В	В	В	В	В	В	В	В	В	В	В
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Storage tape		( <i>i</i> ,	<i>j</i> )		В	В	В	В	В	В	В	В	В	В	В
'								1		1					
Output tone	_	<b>D</b>		_	D	D		<b>D</b>	<b>D</b>	D	_	_		_	



Storage tape		V	v <sub>i</sub>		В	В	В	В	В	В	В	В	В	В	В
Storage tape	В	В	В	В	В	В	В	В	В	В	В	В	В	В	В
Storage tape	( i, j )				В	В	В	В	В	В	В	В	В	В	В
Output tape	В	В	В	В	В	В	В	В	В	В	В	В	В	В	В



Storage tape		W <sub>i</sub>			В	В	В	В	В	В	В	В	В	В	В
·															
Storage tape	,	Sim	ulati	ion	of T	M M	o fo	r str	ring	w <sub>i</sub> ii	ı j s	teps		В	В
Storage tape	( i, j )			В	B B B B B B B							В	В	В	
'								I							
Output tape	В	В	В	В	В	В	В	В	В	В	В	В	В	В	В



## Recursive enumerable languages

Storage tape B B B B В B B B B B  $W_i$ B Storage tape Simulation of TM  $M_2$  for string  $\mathbf{w}_i$  in  $\mathbf{j}$  steps B Storage tape (i+1, j-1)B B B B B B B B B B B B B Output tape B B B B B B B B B B B B B



## Recursive enumerable languages

Storage tape Simulation of TM M<sub>2</sub> for string w<sub>i</sub> in j steps B

Storage tape (i+1, j-1) B B B B B B B B B B B

B B B Output tape B B B B B B B B B B B B



## Recursive enumerable languages

Storage tape B B B B В B B B B B  $W_{i+1}$ B Storage tape Simulation of TM  $M_2$  for string  $w_{i+1}$  in j-1 steps B

Storage tape (i+1, j-1) B B B B B B B B B B B

B B B B B B B Output tape B B B B B B B B



## Recursive enumerable languages

Storage tape B B B B В B B B B B  $W_{i+1}$ B Storage tape Simulation of TM  $M_2$  for string  $w_{i+1}$  in j-1 steps B Storage tape (i+2, j-2)B B B B B B B B B B B B B B B B Output tape B B B B B B B B B B



## Recursive enumerable languages

Storage tape B B B B В B B B B B  $W_{i+2}$ Simulation of TM  $M_2$  for string  $w_{i+1}$  in j-1 steps B Storage tape B Storage tape (i+2, j-2)B B B B B B B B B B B





### Recursive enumerable languages

Storage tape W<sub>i+2</sub> B B B B B B B B B B

Storage tape | Simulation of TM M<sub>2</sub> for string w<sub>i+2</sub> in j-2 steps | B | B

Storage tape ( i+2, j-2 ) B B B B B B B B B B B

B B B B B B Output tape B B B B B B B B B



## Recursive enumerable languages

Storage tape Simulation of TM  $M_2$  for string  $w_{i+2}$  in j-2 steps B B

 Output tape
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 W<sub>i+2</sub>
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## Recursive enumerable languages

Storage tape Simulation of TM  $M_2$  for string  $w_{i+2}$  in j-2 steps B

 Output tape
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**(1,1)** 



## Recursive enumerable languages

Storage tape W<sub>1</sub> B B B B B B B B B B B

Storage tape Simulation of TM  $M_2$  for string  $w_{i+2}$  in j-2 steps B

 Output tape
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 w<sub>i+2</sub>
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## Recursive enumerable languages

Storage tape W<sub>1</sub> B B B B B B B B B B B

Storage tape Simulation of TM M<sub>2</sub> for string w<sub>1</sub> in 1 step B B

Storage tape (1,1),1+1=2 B B B B B B B B B B B

**(1,1)** 



## Recursive enumerable languages

Storage tape W<sub>1</sub> B B B B B B B B B B B

Storage tape Simulation of TM M<sub>2</sub> for string w<sub>1</sub> in 1 step B B

Storage tape (1, 2), 1+2=3 B B B B B B B B B B B

 Output tape
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 w<sub>i+2</sub>
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(1,1), (1,2)



## Recursive enumerable languages

(1,1), (1,2)



## Recursive enumerable languages

Storage tape W<sub>1</sub> B B B B B B B B B B

Storage tape Simulation of TM M<sub>2</sub> for string w<sub>1</sub> in 2 steps B B

Storage tape (1, 2), 1+2=3 B B B B B B B B B B B

 Output tape
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 w<sub>i+2</sub>
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(1,1), (1,2)



## Recursive enumerable languages

(1,1), (1,2), (2,1)



## Recursive enumerable languages

(1,1), (1,2), (2,1)



## Recursive enumerable languages

Storage tape B B B B B B B B B В  $W_2$ Storage tape B B Simulation of TM M<sub>2</sub> for string w<sub>2</sub> in 1 step Storage tape (2, 1), 2+1=3 B B B B B B B B B B B Output tape # W<sub>i+2</sub> # B B B B B B B B B B B B

(1,1), (1,2), (2,1)



## Recursive enumerable languages

B

B

B

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B

(1,1), (1,2), (2,1), (1,3)

Output tape # W<sub>i+2</sub> #



## Recursive enumerable languages



## Recursive enumerable languages

Storage tape B B B B B B B B B В  $W_1$ Storage tape B B Simulation of TM  $M_2$  for string  $\mathbf{w}_1$  in 3 steps Storage tape (1, 3), 1+3=4 B B B B B B B B B B B

Output tape # W<sub>i+2</sub> # B B B B B B B B B B B B B B B B

(1,1), (1,2), (2,1), (1,3)



### Recursive enumerable languages



## Recursive enumerable languages



## Recursive enumerable languages



## Recursive enumerable languages



## Recursive enumerable languages



## Recursive enumerable languages



## Recursive enumerable languages



## Recursive enumerable languages





Binary alphabet: 
$$\Sigma = \{0, 1\}$$



Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,



Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,

100, 101, 110, 111, 0000, 0001, ...



Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,

100, 101, 110, 111, 0000, 0001, ...



B B Storage tape B B B B B B B B B B B B B



Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,

100, 101, 110, 111, 0000, 0001, ...



B B Storage tape B B B B B B B B B B B B B



Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,

100, 101, 110, 111, 0000, 0001, ...

Storage tape Simulation of TM  $M_2$  for string  $\varepsilon$  B



Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,

100, 101, 110, 111, 0000, 0001, ...

Storage tape Simulation of TM  $M_2$  for string  $\varepsilon$  B



Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,

100, 101, 110, 111, 0000, 0001, ...

Storage tape Simulation of TM M<sub>2</sub> for string 0 B B



Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,





Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,





Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,





Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,





Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,





Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,





Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,







Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,





Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,





Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,







Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,







Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,







Binary alphabet:  $\Sigma = \{0, 1\}$ 

Canonical ordering:  $\varepsilon$ , 0, 1, 00, 01, 10, 11, 000, 001, 010, 011,































































YES – the string is accepted





Input tape B B B B B B B B B B B B B B B

Storage tape B B B B B B B B B B B B B B B

Output tape B B B B B B B B B B B B B B B



Storage tape B B B B B B B B B B B B B B B

Output tape B B B B B B B B B B B B B B B



Storage tape B B B B B B B B B B B B B B B

 Output tape
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Storage tape B B B B B B B B B B B B B B B

Output tape #W<sub>i</sub> W<sub>1</sub> # B B B B B B B B B B B B B



Storage tape B B B B B B B B B B B B B B B

Output tape # w<sub>1</sub> # w<sub>2</sub> # B B B B B B B B B B B



Storage tape B B B B B B B B B B B B B B B

Output tape # W<sub>1</sub> # W<sub>2</sub>W<sub>i</sub> # B B B B B B B B B B B



Storage tape B B B B B B B B B B B B B B B

Output tape # w<sub>1</sub> # w<sub>2</sub> # w<sub>3</sub> # B B B B B B B B



Storage tape B B B B B B B B B B B B B B B

Output tape # W<sub>1</sub> # W<sub>2</sub> # W<sub>3</sub>W<sub>i</sub> # B B B B B B B B



Storage tape B B B B B B B B B B B B B B B

Output tape # w<sub>1</sub> # w<sub>2</sub> # w<sub>3</sub> # B B w<sub>j</sub> # B B B



Storage tape B B B B B B B B B B B B B B B

Output tape # w<sub>1</sub> # w<sub>2</sub> # w<sub>3</sub> # B B w<sub>j</sub>w<sub>i</sub> # B B B





```
B
Storage tape
                B
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                       B
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                                            B
                                                B
                                                   B
                                                       B
                                                          B
                                                              B
             B
```

In canonical ordering j is before i





Storage tape B B B B B B B B B B B B B B B

In canonical ordering j is before i

NO – the string is not accepted





Storage tape B B B B B B B B B B B B B B B

In canonical ordering j is before i

NO – the string is not accepted

If the language is finite and if a string is processed that follows all the strings in  $L(M_2)$  then TM  $M_2$  never stops

