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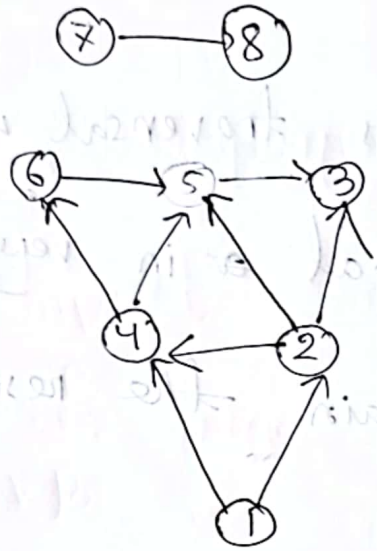
Sec: 16

LAB assignment 5

code explanation

Task-1 (b)



node	indegree
1	0
2	1 + 0
3	2 + 0
4	2 + 1 + 0
5	3 - 2 + 0
6	1 + 0
7	0
8	1 - 1 + 0

3
5
6
4
8
2
7
1

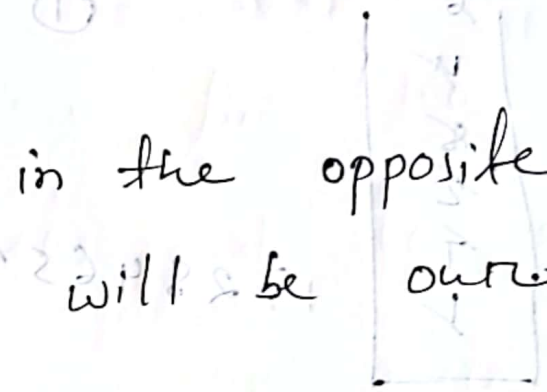
1 7 2 8 4 6 5 3

where first we have count all the indegrees of every node. the if any node's indegree is 0 we add to the queue and visit all his neighbours And updated the indegree count of that visited nodes (-1). the queue is our answer.

1. a)

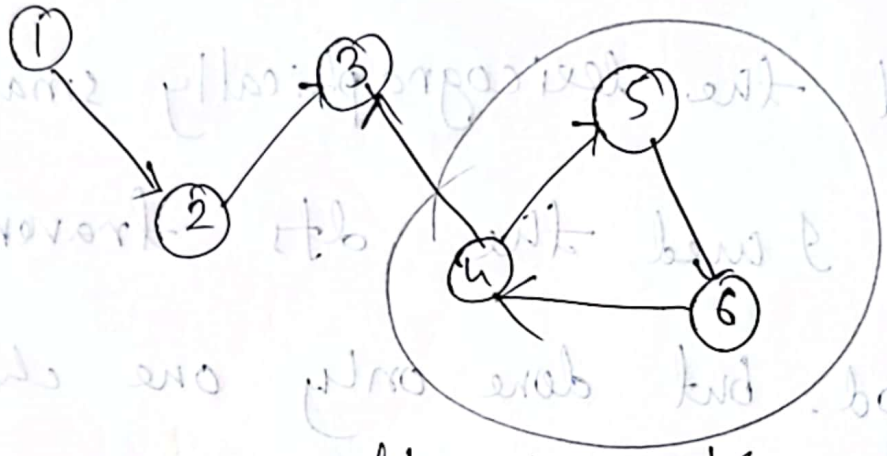
Here ~~we~~ for the dfs traversal we would just do a normal dfs traveral ~~or~~ in very unvisited node and maintain the recursion stack.

The recursion satack in the opposite ordere [: : -1] will be our answer.



- But we will face problems while detecting the "Impossible" cases.

An impossible case :



here we can it's a impossible case.

we can detect impossible cases if any circle is found. So, I have taken a rec-stack ~~LIST~~ ⇒ boolian list to detect any circle. if we find any circle we just return 'Impossible'.

Task 2 |

to get the lexicographically smaller path : I used the dfs traversal method. but done only one change before deque anything for the queue we just made the queue sorted ; so that we can maintain lexicographically smaller path

queue = deque(sorted(queue))

=>

Tasks:

find scc:

fun: dfs:

① it's a function give us the dfs recursion stack

② the transpose-graph give us the same graph just in the opposite direction.

③ then we run dfs-scc in the the transpose-graph in ~~the~~ according to the recursion stack. and we find-scc to detect every cycle.