

[Chap.6-1] The Memory Hierarchy

Young Ik Eom (yieom@skku.edu, 031-290-7120)
Distributing Computing Laboratory
Sungkyunkwan University
<http://dclab.skku.ac.kr>



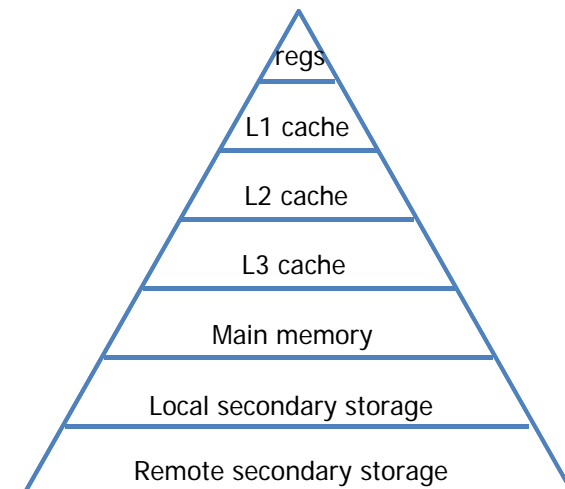
Contents

- Introduction
- Storage technologies
- Locality
- The memory hierarchy
- Cache memories
- Writing cache-friendly code
- Impact of caches on program performance

Introduction

■ Memory system

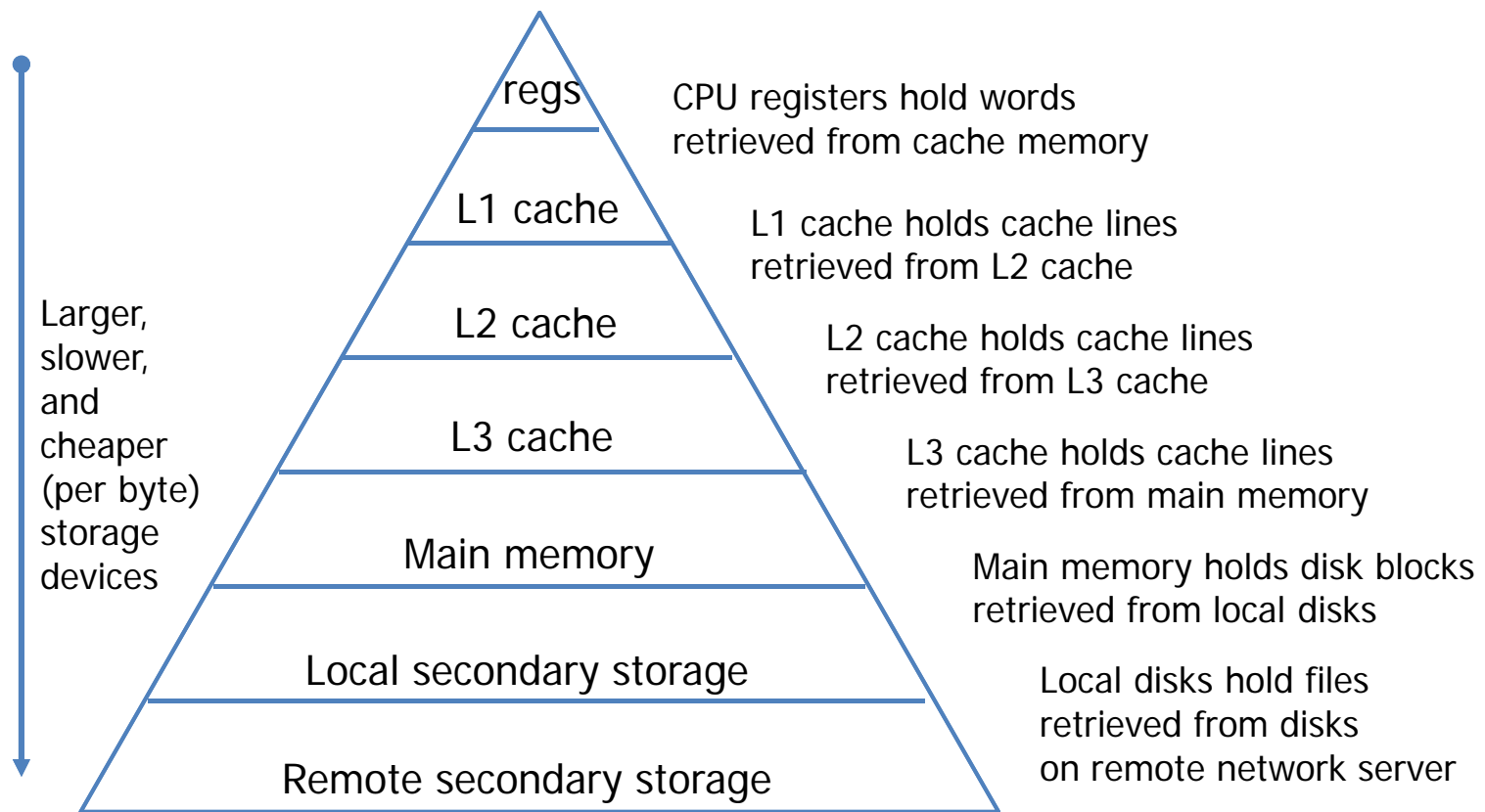
- Hierarchy of storage devices with different capacities, costs, and access times
 - CPU registers (can be accessed in 0 cycles)
 - Cache memories (can be accessed in 4~75 cycles)
 - Main memory (can be accessed in hundreds of cycles)
 - Hard disk (can be accessed in tens of millions of cycles)
 - Etc



Introduction

■ Memory system

■ Hierarchy of storage devices



Introduction

■ Memory system

■ Locality

- Programs with good locality
 - ✓ Tend to access the same set of instructions and data items over and over again
 - ✓ Tend to access more instructions and data items from the upper levels of memory hierarchy than programs with poor locality
 - ✓ Run faster than programs with poor locality

Introduction

■ Memory system

■ Cache memory

- Have the most impact on application program performance
- Focuses on ...
 - ✓ Analyzing programs for locality
 - ✓ Improving locality

Storage Technologies

■ RAM (Random Access Memory)

- SRAM (Static RAM)
 - Faster and more expensive
 - Used for cache memories
- DRAM (Dynamic RAM)
 - Used for main memory and frame buffer of a graphic system
- Both are volatile
 - Lose information if the supply voltage is turned off

Storage Technologies

■ SRAM and DRAM

	Trans. per bit	Access time	Needs refresh?	Sensitive?	Cost	Applications
SRAM	6	1×	No	No	1,000×	Cache M
DRAM	1	10×	Yes	Yes	1×	Main M, Frame buffers

■ SRAM

- Persistent
 - ✓ No refresh is necessary
- Faster than DRAM
- Not sensitive to disturbances such as light and electrical noise
- Lower density
- Consumes more power
- More expensive

Storage Technologies

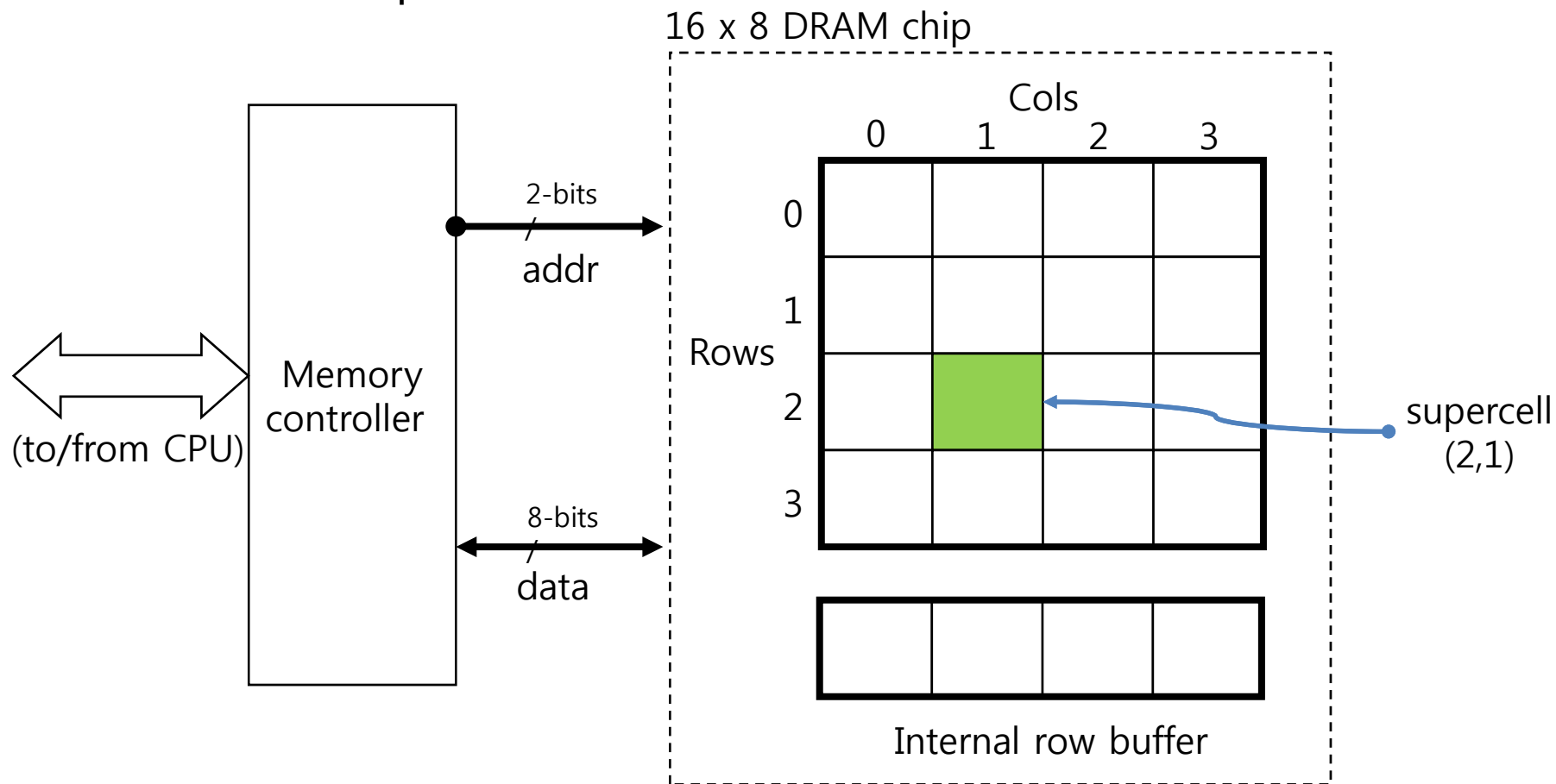
■ Conventional DRAMs

- DRAM chip
 - Partitioned into **d** supercells, each consisting of **w** DRAM cells
 - A **d**×**w** DRAM can store **d**·**w** bits
 - Supercells organized as a rectangular array with **r** rows and **c** columns (**d** = **r**·**c**)
- MC(Memory Controller)
 - Transfers **w** bits at a time to/from DRAM chip
 - For reading supercell (**i**, **j**),
 - ✓ MC sends row address **i** (RAS(Row Access Strobe) request), and then column address **j** (CAS(Column Address Strobe) request)
 - ✓ RAS and CAS shares the same address pins

Storage Technologies

■ Conventional DRAMs

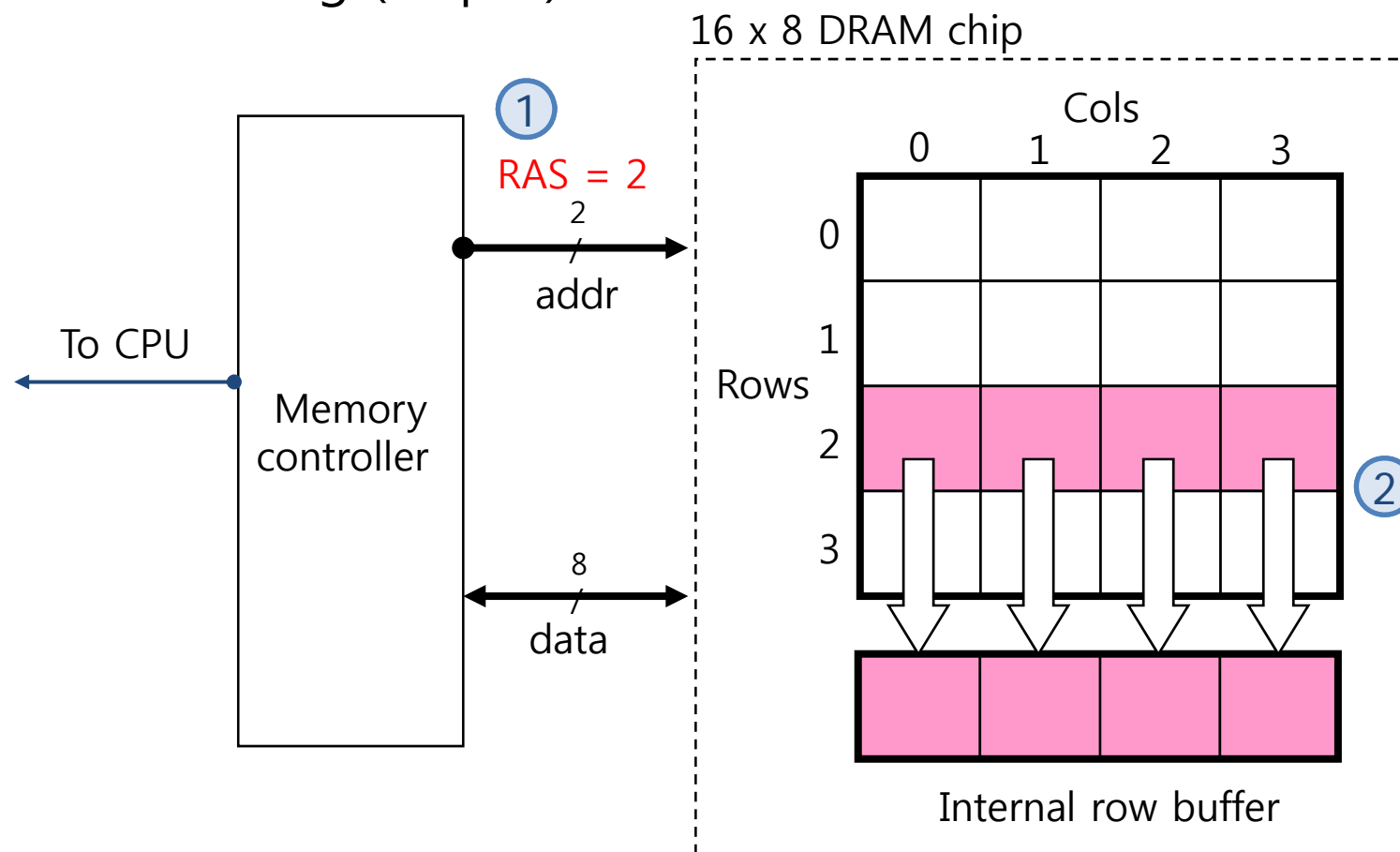
■ DRAM chip



Storage Technologies

■ Conventional DRAMs

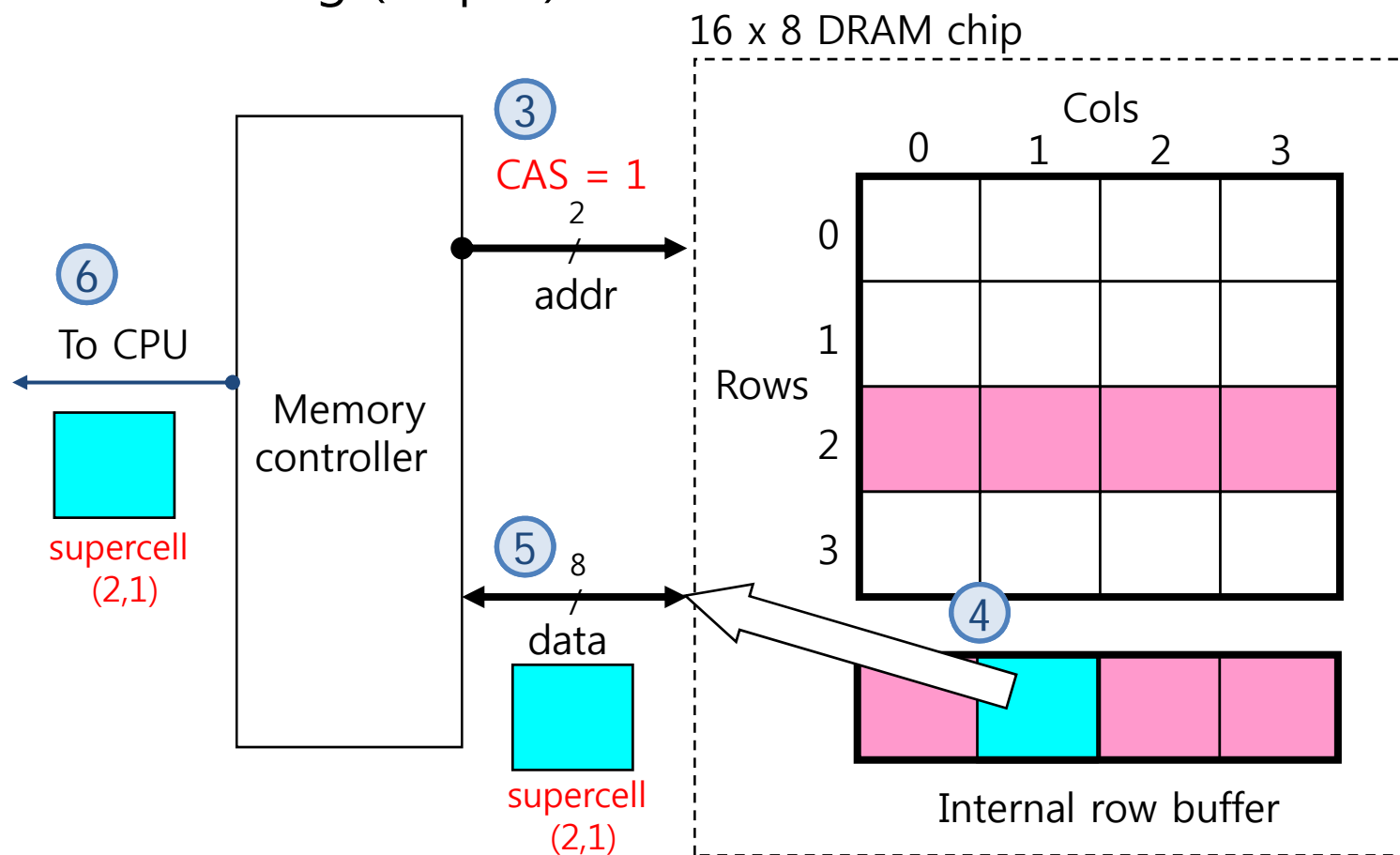
■ Reading (step-1)



Storage Technologies

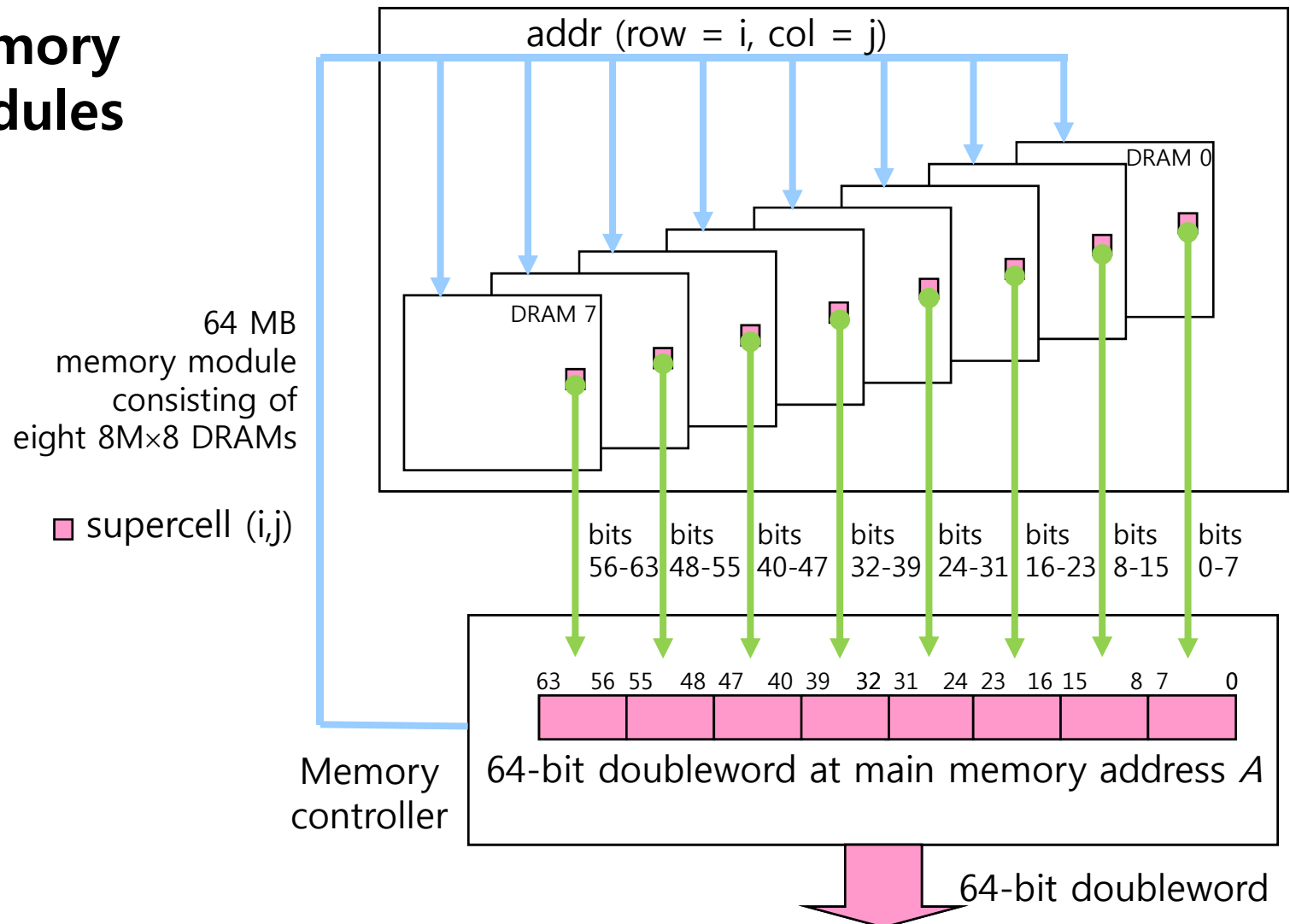
■ Conventional DRAMs

▪ Reading (step-2)



Storage Technologies

■ Memory modules



Storage Technologies

■ Enhanced DRAMs

- FPM DRAM (Fast page mode)
 - Allows reuse of row addresses (quick access)
 - ✓ Eg) RAS → CAS → CAS → CAS
- EDO DRAM (Extended data out)
- SDRAM (Synchronous)
 - Uses a conventional clock signal instead of asynchronous control
- DDR SDRAM (Double data-rate synchronous)
 - Double edge clocking sends two bits per cycle per pin
 - DDR2, DDR3 (used in Intel Core i7)
- VRAM (Video)
 - Used in frame buffers of graphics system

Storage Technologies

■ Nonvolatile memories

- Retain their information even when they are powered off
- ROMs
 - PROM, EPROM, EEPROM
- Flash memory (SSD)
- Hard disk (HDD)
- NVRAM(Non-Volatile RAM)
 - PRAM, STT-MRAM

Storage Technologies

■ Nonvolatile memories: ROMs (Read-Only Memories)

- ROM
 - Programmed during production
- PROM (Programmable ROM)
 - Can be programmed exactly once
- EPROM (Erasable PROM)
 - Can be erased and reprogrammed on the order of 1,000 times
 - Erasing by shining ultraviolet light
 - Programming is done by using a special device
- EEPROM (Electrically Erasable PROM)
 - Can be programmed in-place on printed circuit cards
 - ✓ No separate programming device
 - Can be reprogrammed on the order of 10,000 times

Storage Technologies

■ Nonvolatile memories: Flash memories

- Fast and durable nonvolatile storages for many electronic devices, including smart phones, tablets, laptops, desktops, and digital cameras
- Asymmetric read-write speed
- No in-place writes
- Wears out after about 100,000 erasings

- NAND flash
- NOR flash

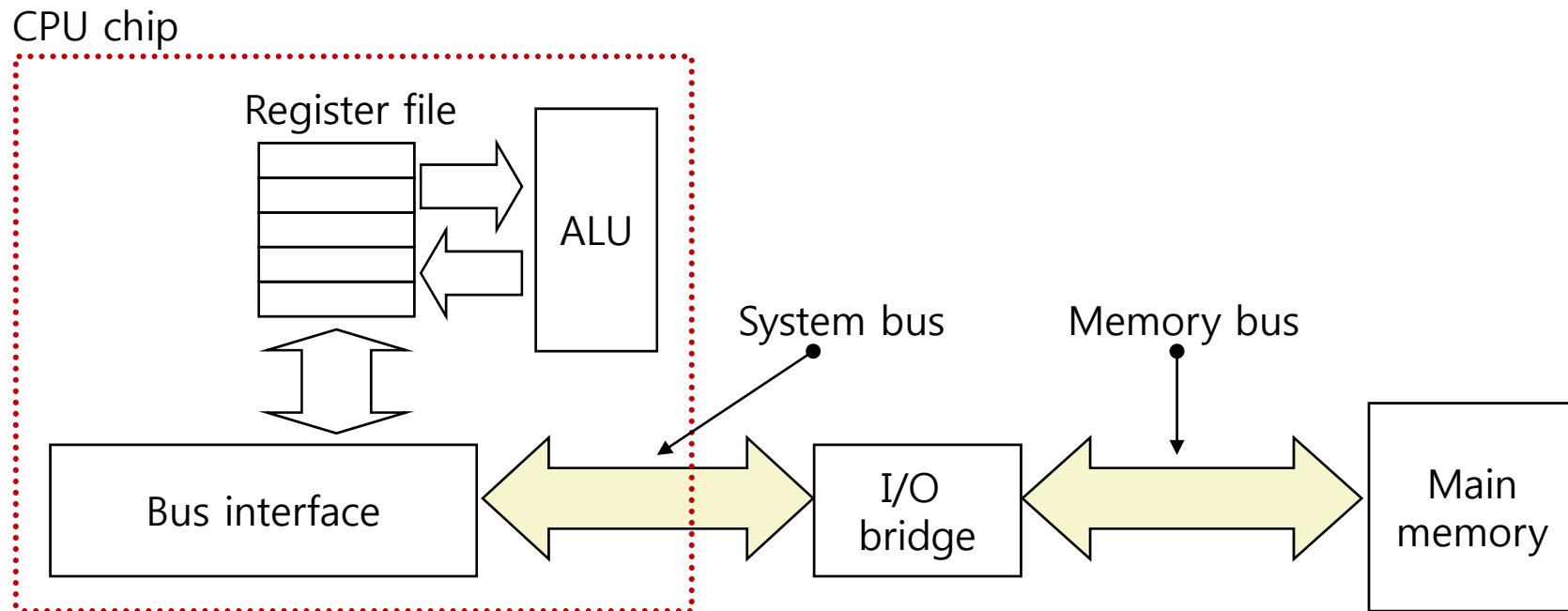
- SSD (Solid State Disk)

Storage Technologies

■ Bus structure

■ Bus

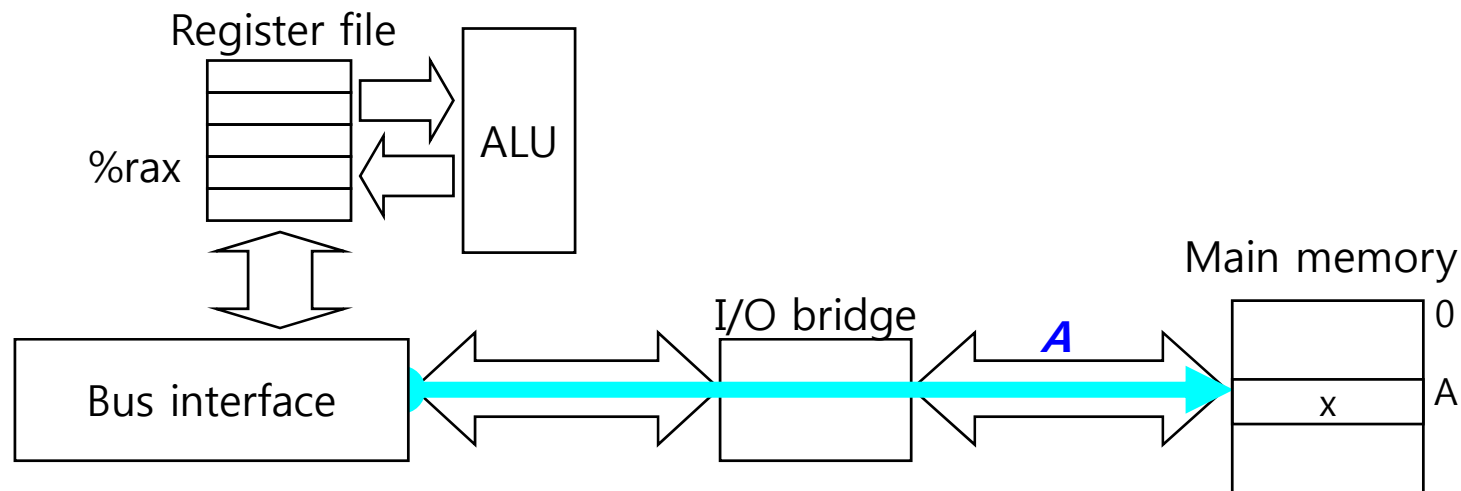
- Collection of parallel wires that carry address, data, and control signals
- Typically shared by multiple devices



Storage Technologies

■ Bus structure

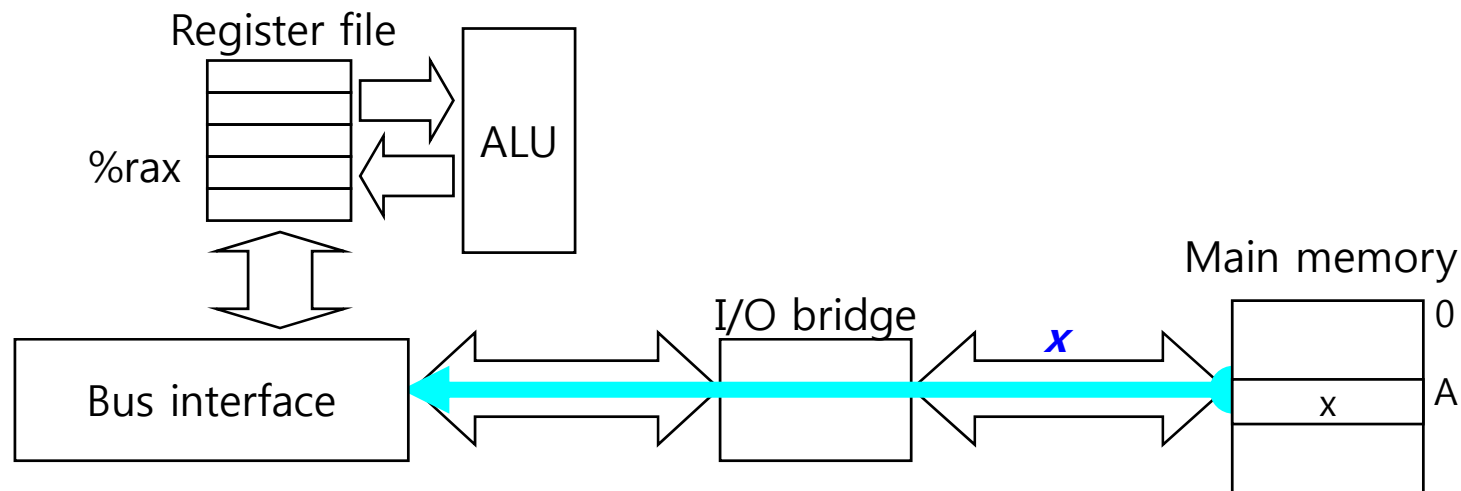
- Read transaction (**movq A,%rax**)
 - (Step-1) CPU places the address **A** on the system bus



Storage Technologies

■ Bus structure

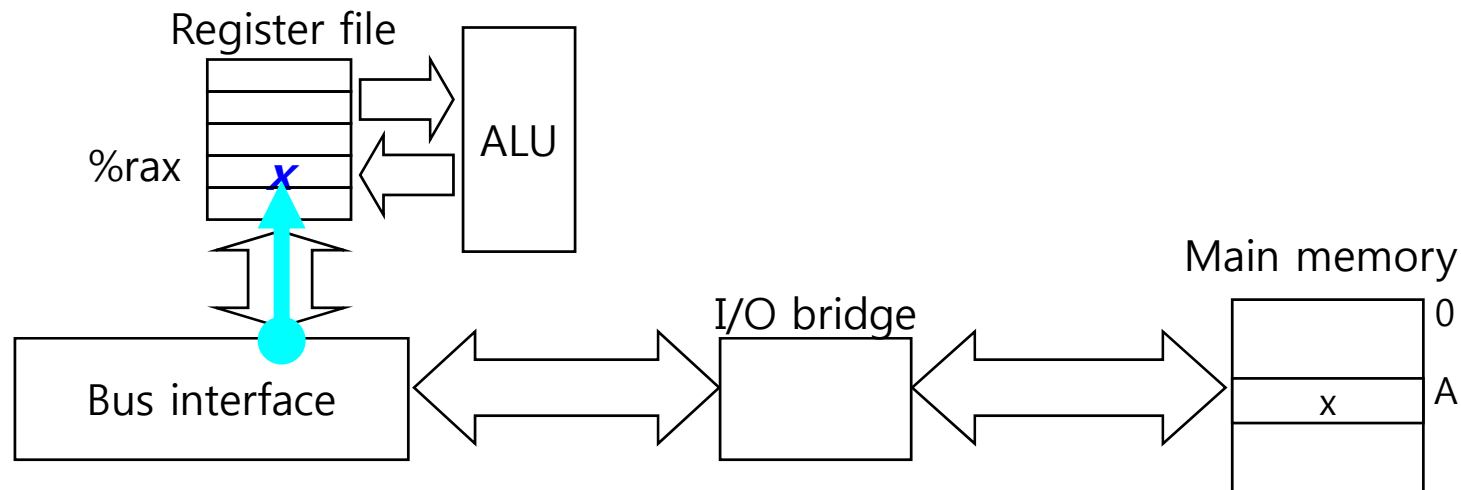
- Read transaction (**movq A,%rax**)
 - (Step-2) Main memory reads **A** from the memory bus, retrieves word **x**, and places it on the bus



Storage Technologies

■ Bus structure

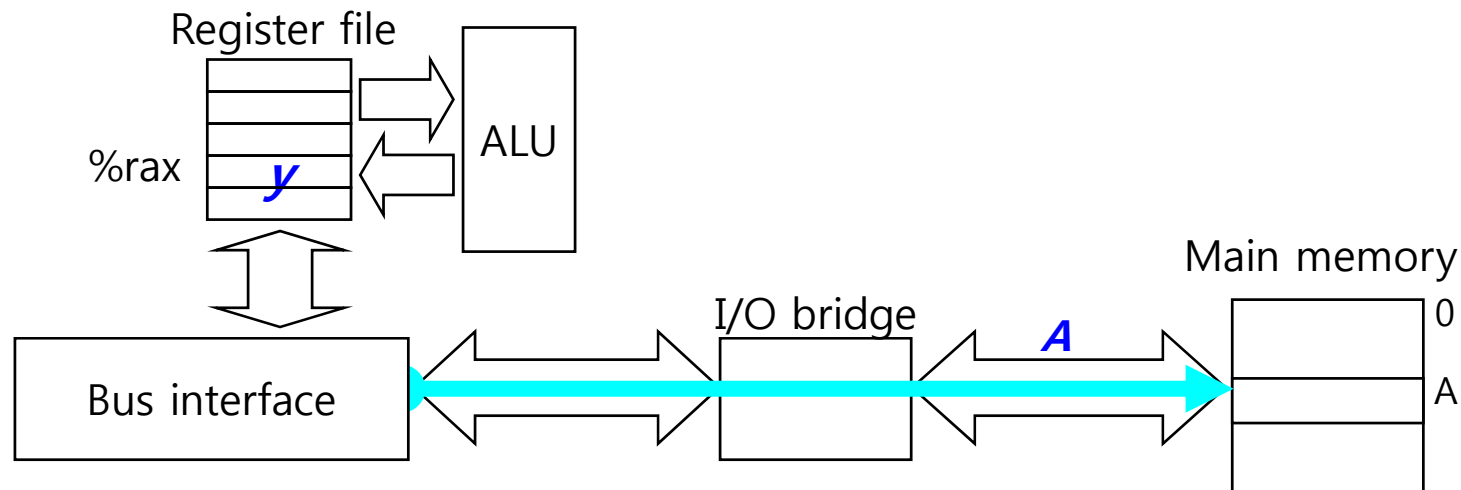
- Read transaction (**movq A,%rax**)
 - (Step-3) CPU reads word **x** from the bus and copies it into register **%rax**



Storage Technologies

■ Bus structure

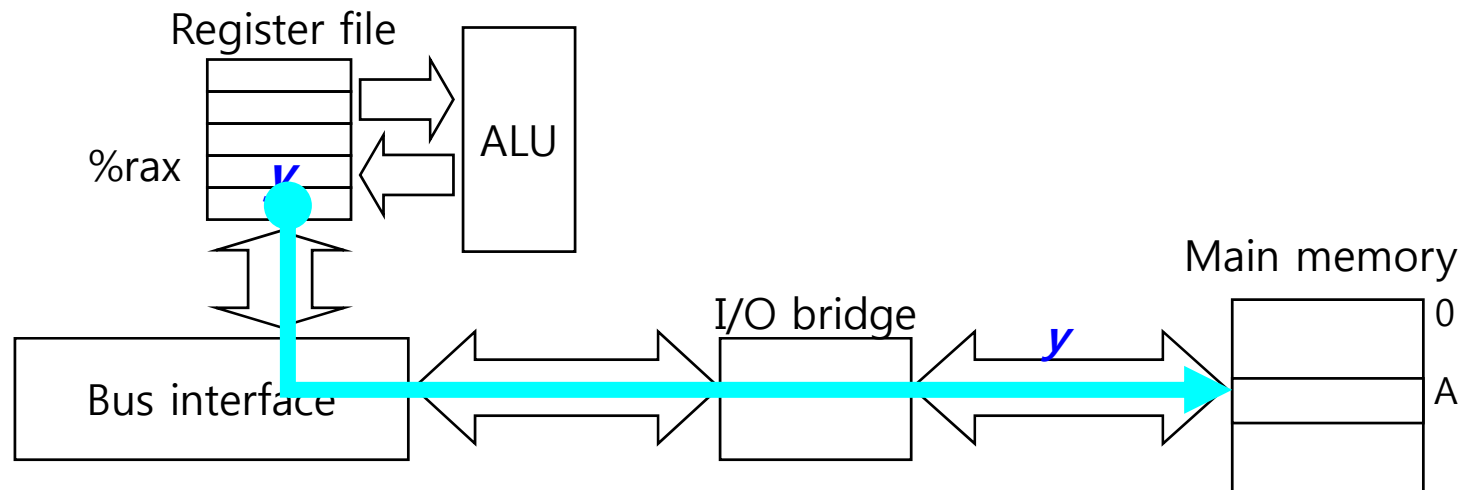
- Write transaction (**movq %rax,A**)
 - (Step-1) CPU places address **A** on bus; Main memory reads it and waits for the corresponding data word to arrive



Storage Technologies

■ Bus structure

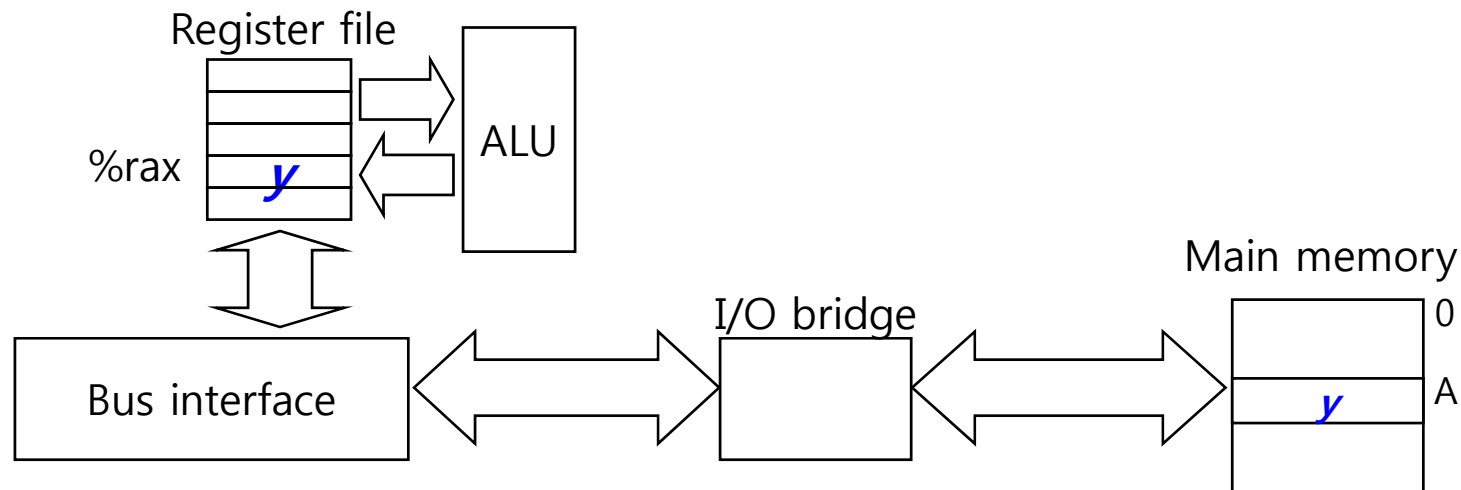
- Write transaction (**movq %rax,A**)
 - (Step-2) CPU places data word **y** on the bus



Storage Technologies

■ Bus structure

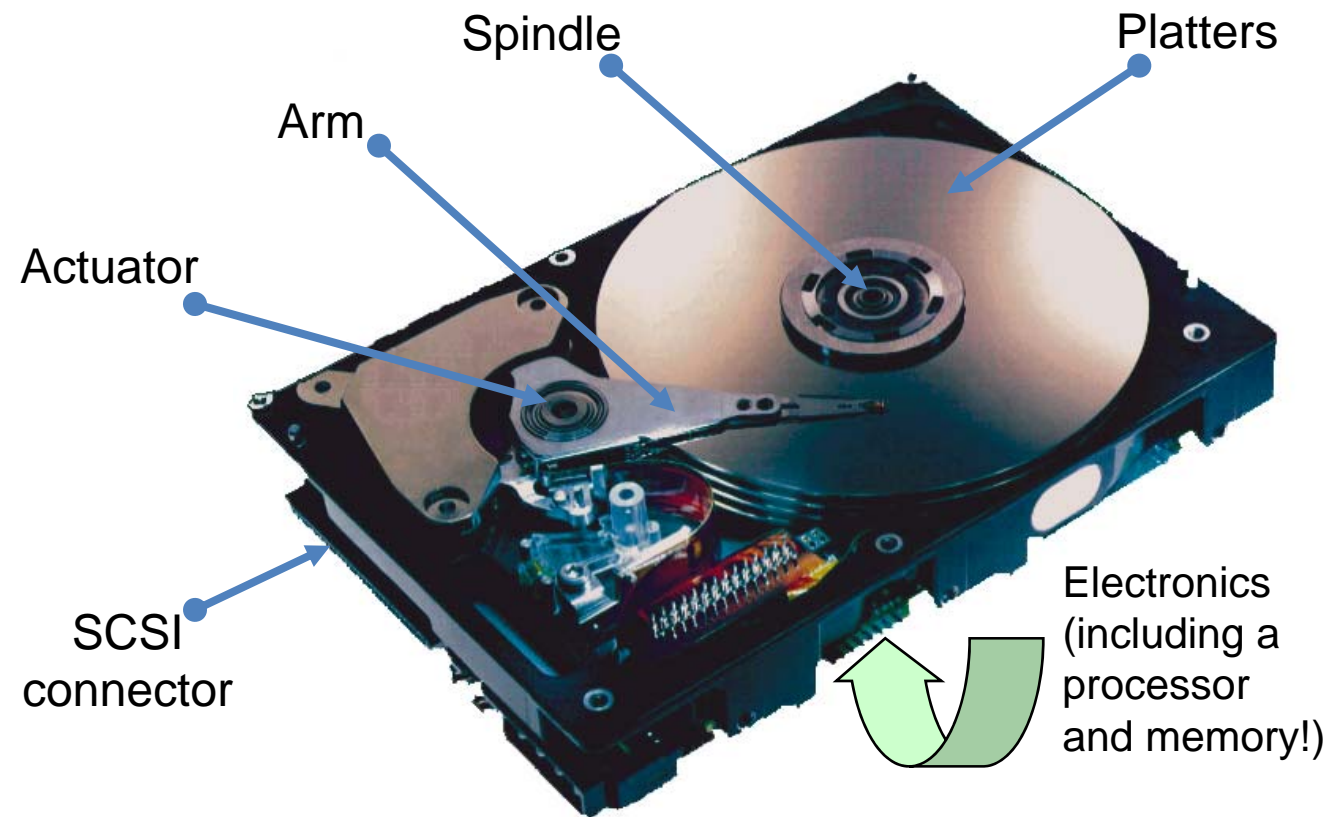
- Write transaction (**movq %rax,A**)
 - (Step-3) Main memory reads data word **y** from the bus and stores it at address A



Storage Technologies

■ HDDs (Hard Disk Drives)

▪ Inside the HDDs

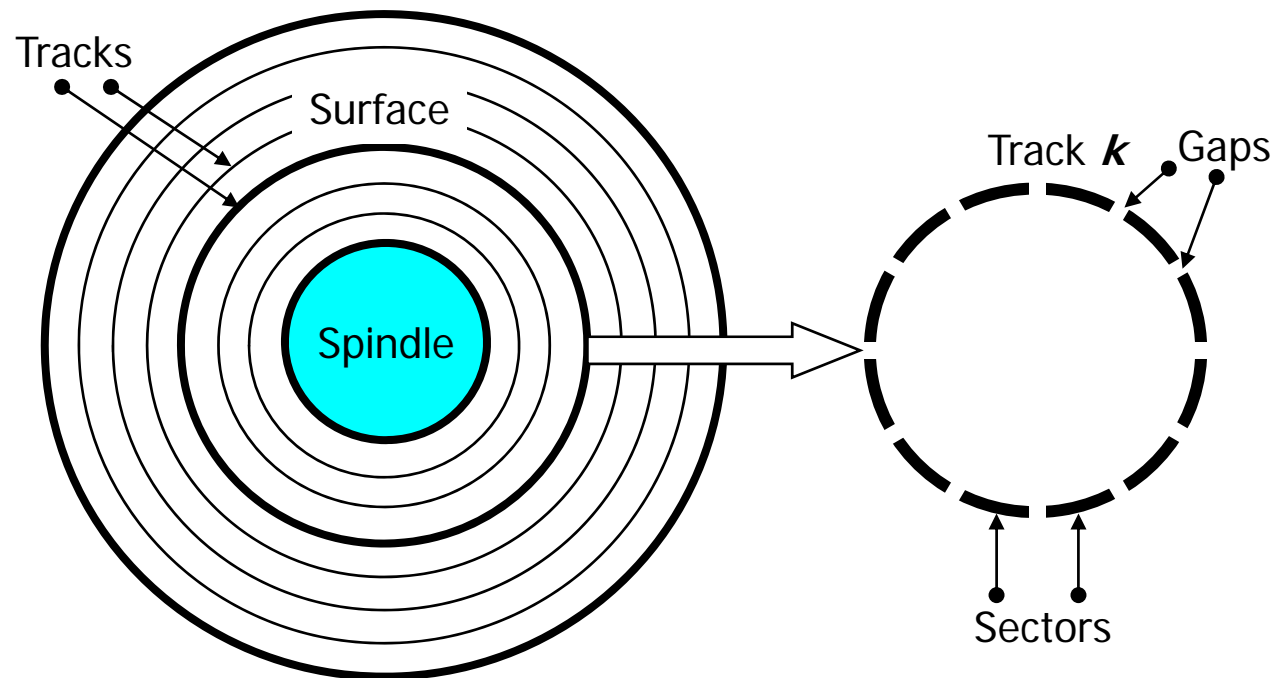


[Image courtesy of Seagate Technology]

Storage Technologies

■ HDDs

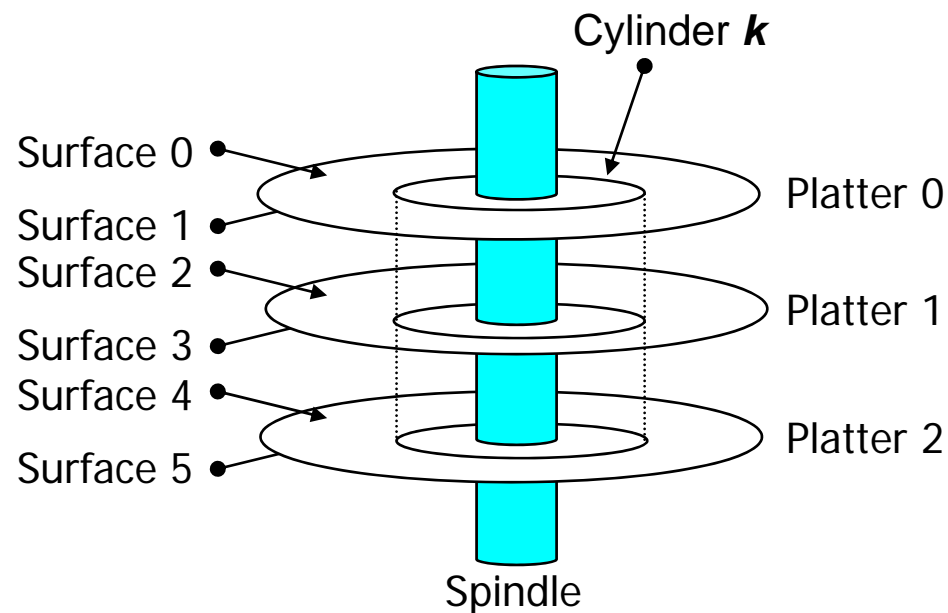
- Disks consist of **platters**, each with two **surfaces**
- Each surface consists of concentric rings called **tracks**
- Each track consists of **sectors** separated by **gaps**



Storage Technologies

■ HDDs

- Aligned tracks form a **cylinder**



Storage Technologies

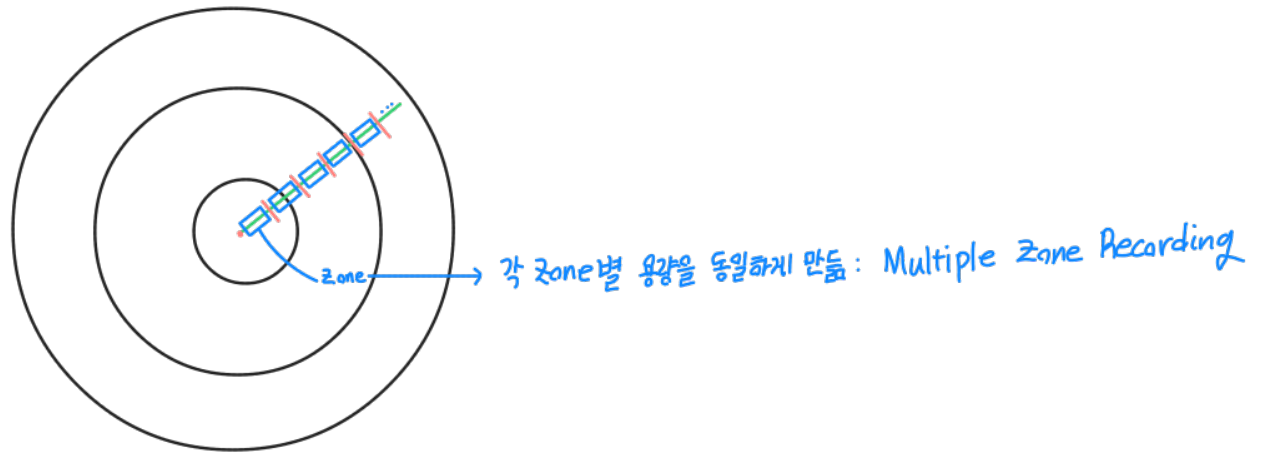
■ Disk capacity

- Maximum number of bits that can be stored
 - Vendors express capacity in units of gigabytes (GB), where $1\text{GB} \cong 10^9$ Bytes
 - Capacity factors
 - ✓ Recording density (bits/in): number of bits that can be squeezed into a 1 inch segment of a track
 - ✓ Track density (tracks/in): number of tracks that can be squeezed into a 1 inch radial segment
 - ✓ Areal density (bits/in²): product of recording and track density

Storage Technologies

■ Disk capacity

- Multiple zone recording
 - Modern disks partition tracks into disjoint subsets called **recording zones**
 - ✓ Each track in a zone has the same number of sectors, determined by the circumference of innermost track in the zone
 - ✓ Each zone has a different number of sectors per track



Storage Technologies

■ Disk capacity

$$\text{Capacity} = (\# \text{ bytes/sector}) \times (\text{avg. } \# \text{ sectors/track}) \times (\# \text{ tracks/surface}) \times (\# \text{ surfaces/platter}) \times (\# \text{ platters/disk})$$

Example)

- 512 bytes/sector
- 300 sectors/track (on average)
- 20,000 tracks/surface
- 2 surfaces/platter
- 5 platters/disk

$$\text{Capacity} = 512 \times 300 \times 20,000 \times 2 \times 5 = 30,720,000,000 = 30.72 \text{ GB}$$

Storage Technologies

■ Disk operation

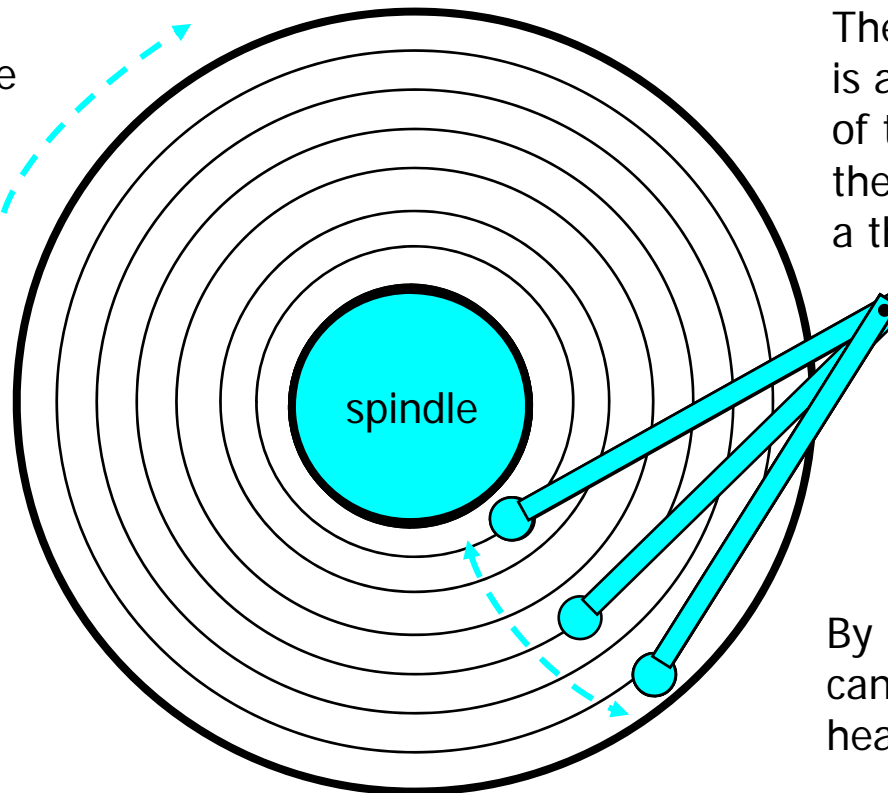
▪ Average time to access some target sector

- $T_{\text{access}} = T_{\text{avg-seek}} + T_{\text{avg-rotation}} + T_{\text{avg-transfer}}$
- Seek time ($T_{\text{avg-seek}}$)
 - ✓ Time to position heads over cylinder containing target sector
 - ✓ Typical $T_{\text{avg-seek}}$ is 3~9ms (Max 20ms)
- Rotational latency ($T_{\text{avg-rotation}}$)
 - ✓ Time waiting for first bit of target sector to pass under head
 - ✓ $T_{\text{avg-rotation}} = 1/2 \times 1/\text{RPMs} \times 60\text{sec}/1\text{min}$
 - ✓ Typical $T_{\text{avg-rotation}} = 1\sim 4\text{ms}$
- Transfer time ($T_{\text{avg-transfer}}$)
 - ✓ Time to read the bits in the target sector
 - ✓ $T_{\text{avg-transfer}} = 1/\text{RPM} \times 1/(\text{avg. \# sectors/track}) \times 60\text{secs}/1\text{min}$

Storage Technologies

■ Disk operation

The disk surface spins at a fixed rotational rate

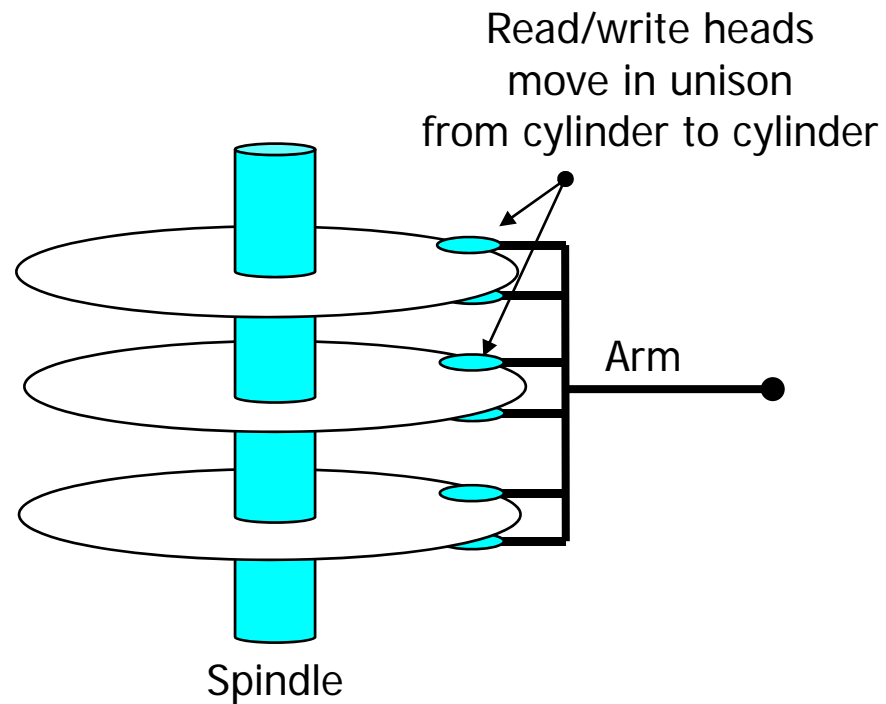


The read/write head is attached to the end of the arm and flies over the disk surface on a thin cushion of air

By moving radially, the arm can position the read/write head over any track

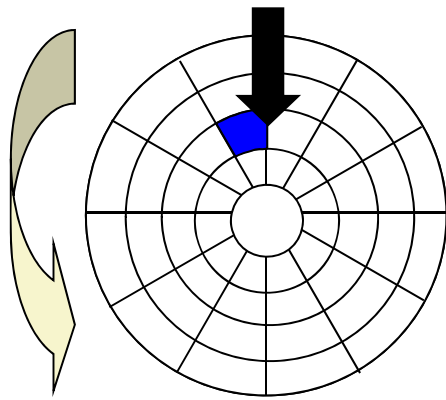
Storage Technologies

■ Disk operation



Storage Technologies

■ Disk operation: A scenario

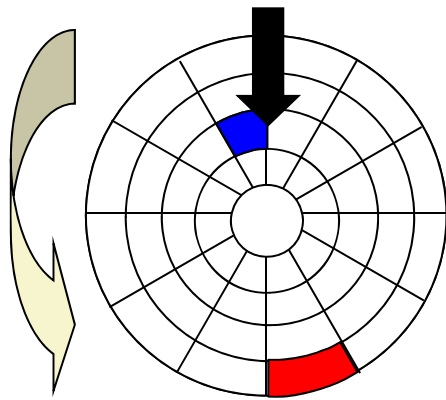


After BLUE read

After reading blue sector

Storage Technologies

■ Disk operation: A scenario

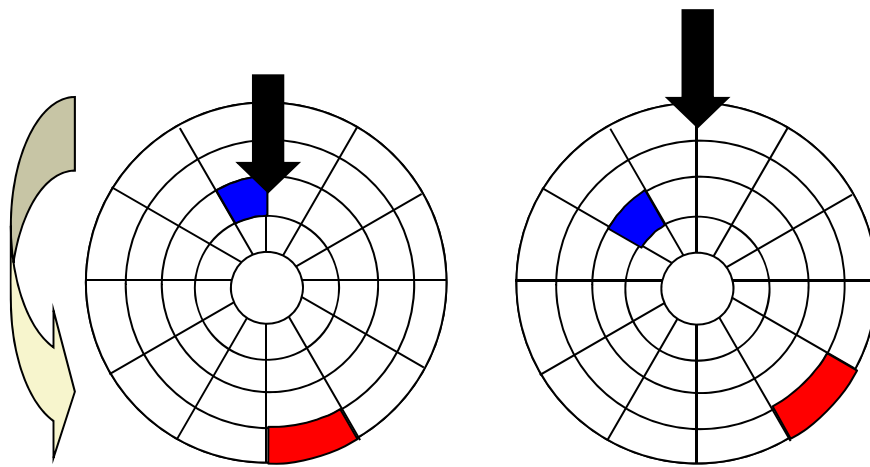


After BLUE read

Red request scheduled next

Storage Technologies

■ Disk operation: A scenario



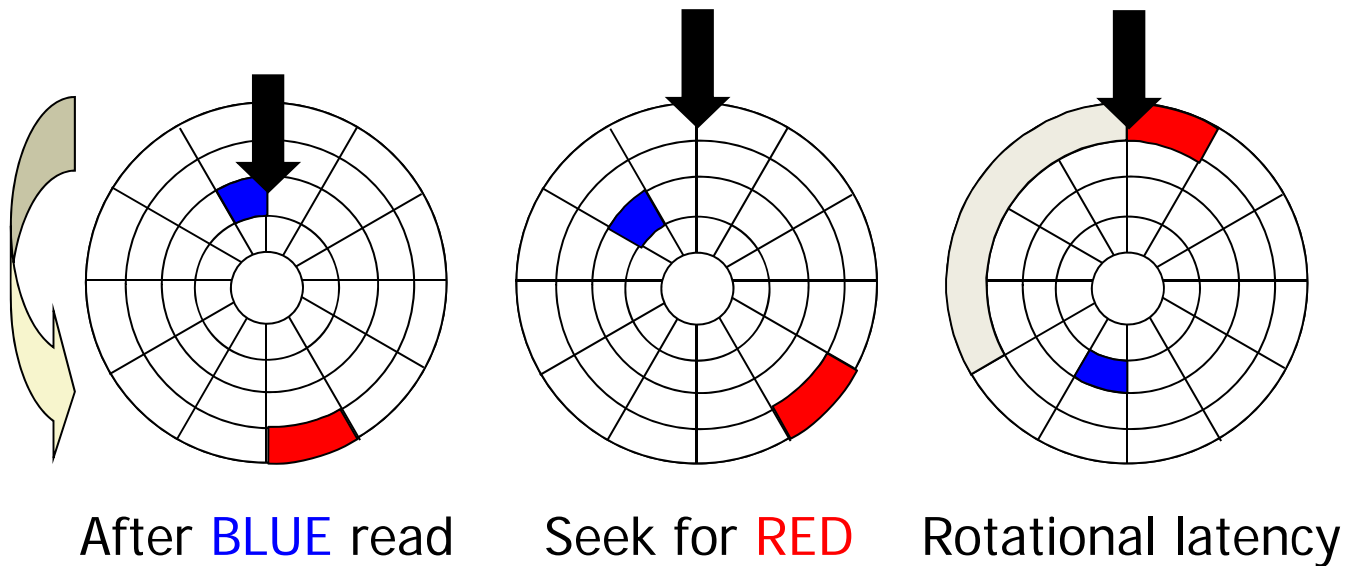
After BLUE read

Seek for RED

Seek to red's track

Storage Technologies

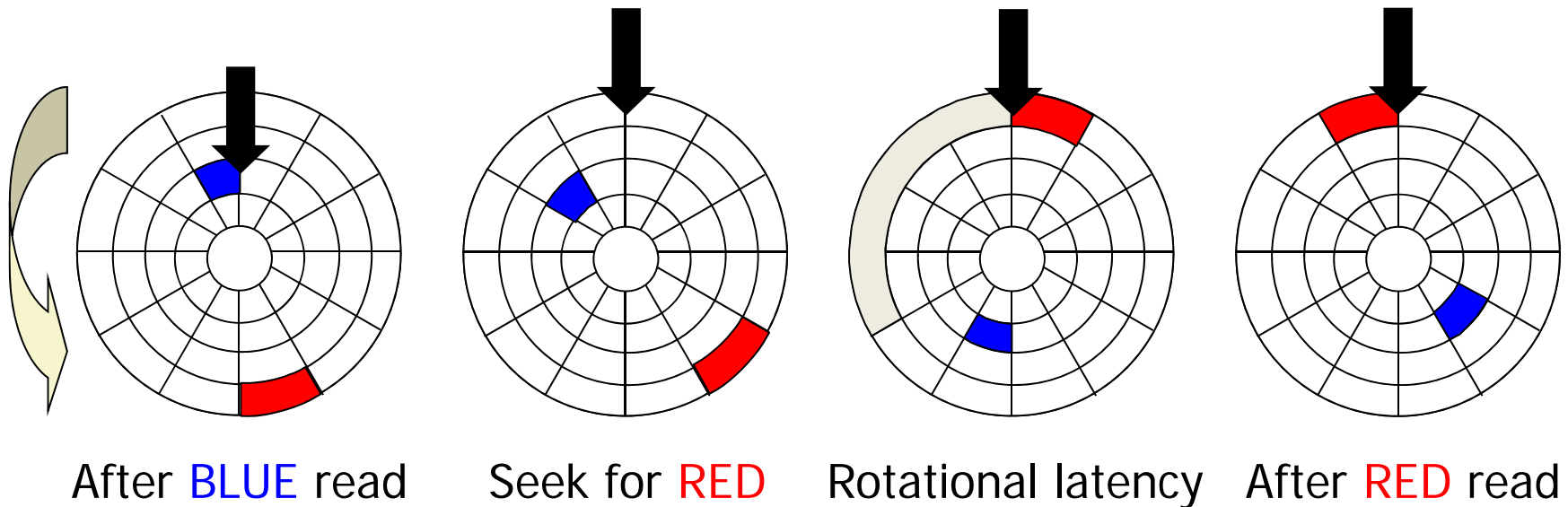
■ Disk operation: A scenario



Wait for red sector to rotate around

Storage Technologies

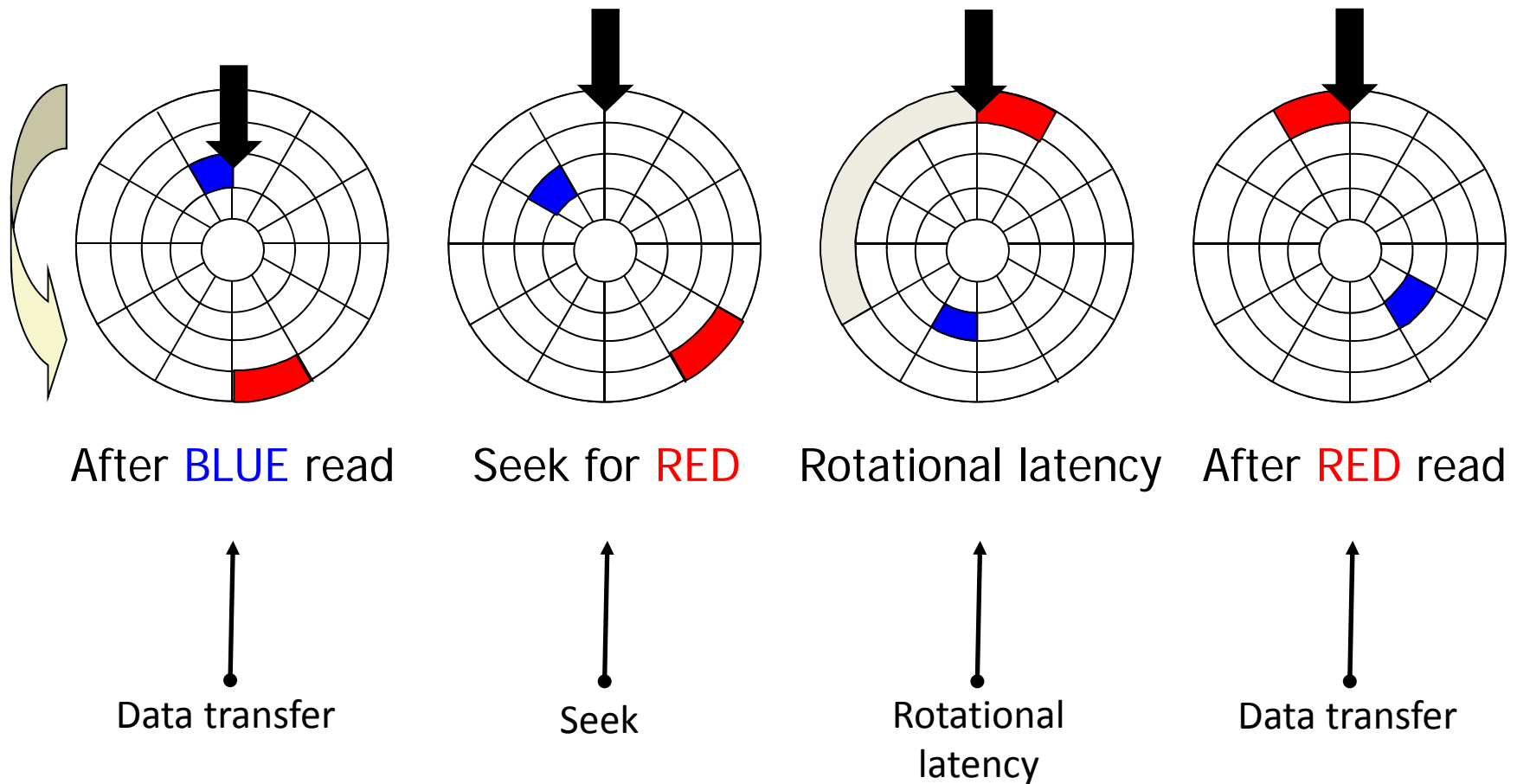
■ Disk operation: A scenario



Complete read of red

Storage Technologies

■ Disk operation: A scenario



Storage Technologies

■ Disk operation

■ Example)

- Given

- ✓ Average seek time = 9ms
- ✓ Rotational rate = 7,200 RPM
- ✓ Avg. # sectors/track = 400

- Derived

- ✓ $T_{\text{avg-rotation}} = 1/2 \times (60\text{secs}/7200\text{RPM}) \times 1000\text{ms/sec} = 4.17\text{ms}$
- ✓ $T_{\text{avg-transfer}} = 1/(400\text{sectors/track}) \times 60/7200\text{RPM} \times 1000\text{ms/sec}$
 $= 0.02 \text{ ms}$
- ✓ $T_{\text{access}} = 9\text{ms} + 4.17\text{ms} + 0.02\text{ms} = 13.19\text{ms}$

Storage Technologies



■ Disk operation

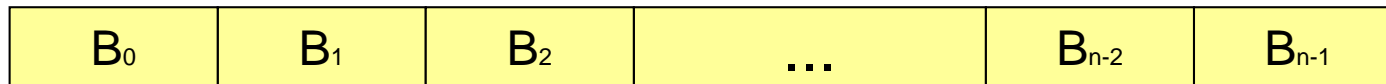
■ Notes)

- Access time dominated by seek time and rotational latency
- SRAM access time is about 4ns/doubleword (256ns for 512B),
DRAM about 60ns (about 4,000ns for 512B)
 - ✓ Disk, which has roughly 10ms access time,
is about 40,000 times slower than SRAM, and
2,500 times slower than DRAM

Storage Technologies

■ Logical disk blocks

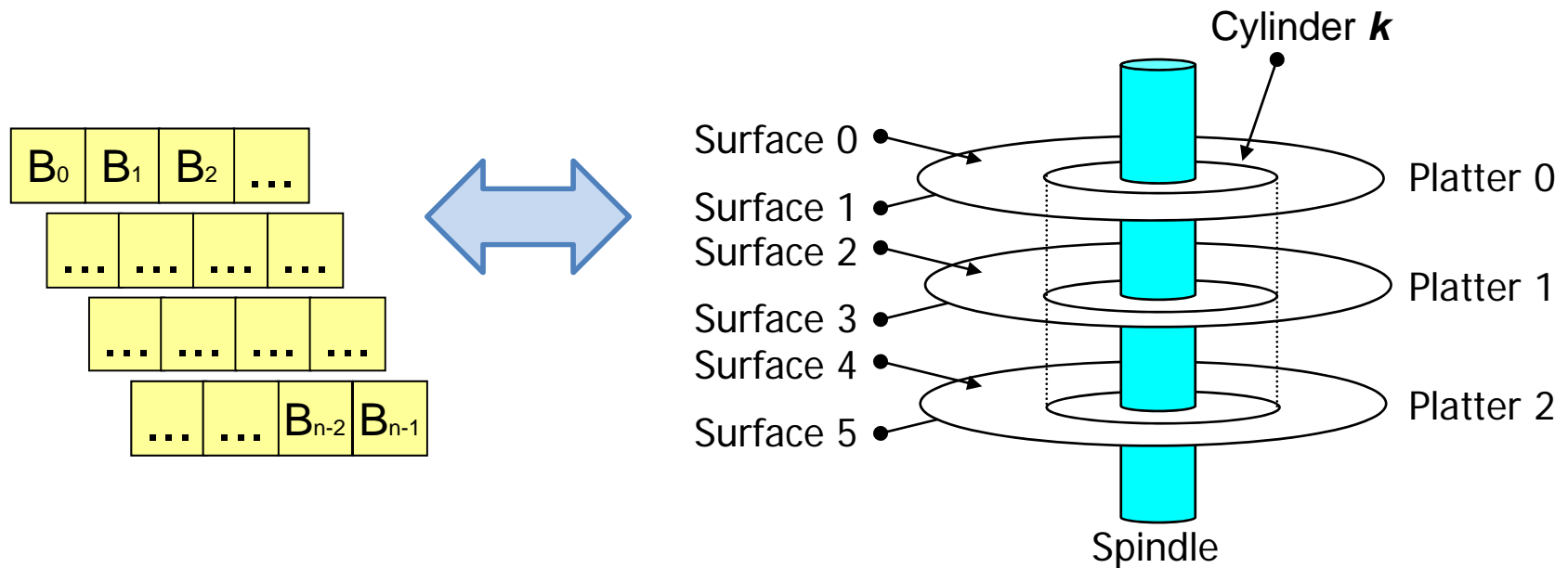
- Modern disks present a simpler abstract view of the complex sector geometry
 - The set of available sectors is modeled as a sequence of **b**-sized logical blocks (0, 1, 2, ...)



Storage Technologies

■ Logical disk blocks

- Mapping between logical blocks and physical sectors
 - Maintained by hardware/firmware device called **disk controller**
 - Converts requests for logical blocks into (track#, surface#, sector#) triples



Storage Technologies

■ Disk formatting

- Filling in the gaps between sectors with information that identifies the sector
- Identifying cylinders with surface defects and taking them out of action (*bad sector $\frac{3}{2}$ out of action*)
- Setting aside a set of cylinders in each zone as spares
 - Called into action when one or more cylinders in the zone goes to bad
- Accounts for the difference in formatted capacity and maximum capacity

Storage Technologies



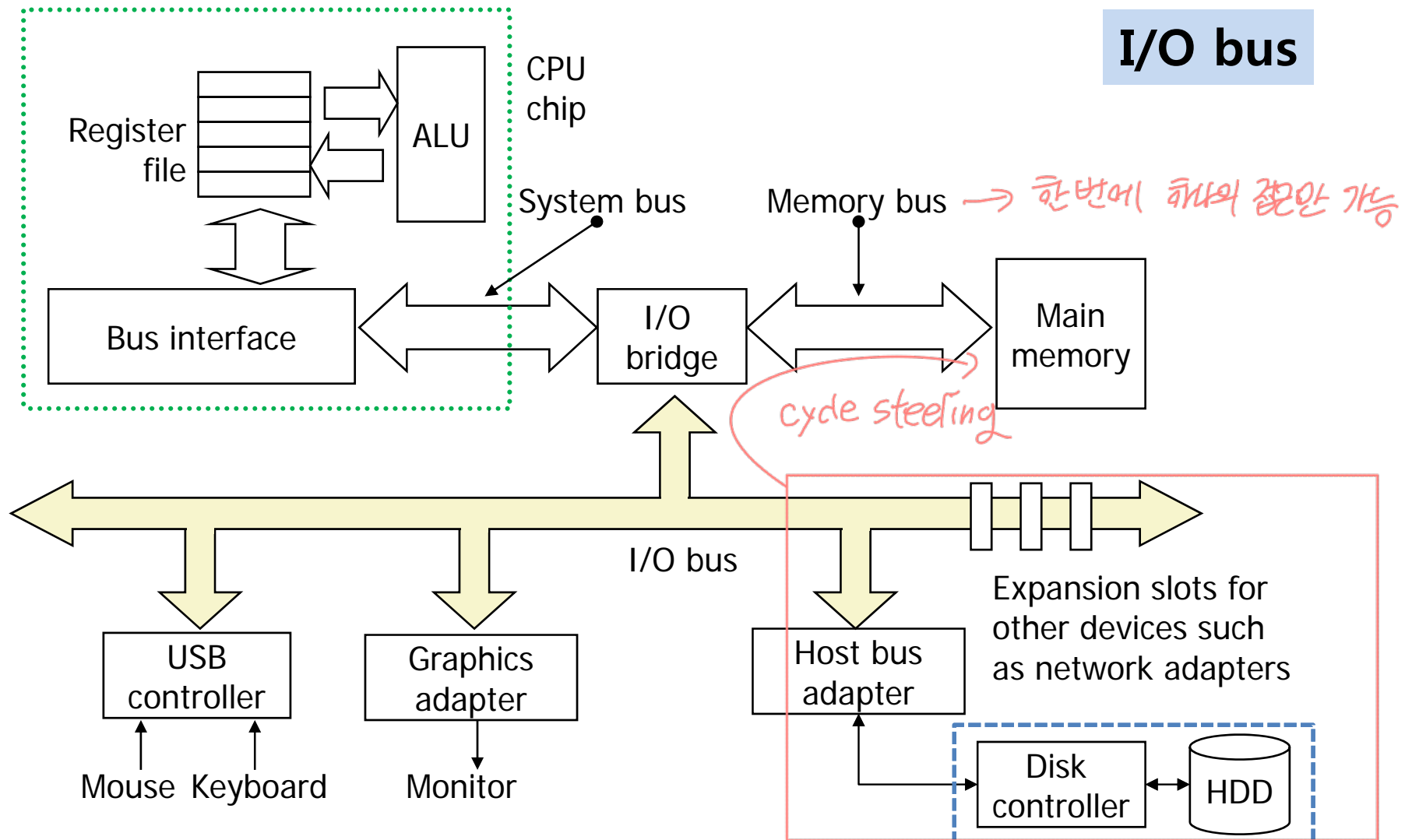
■ I/O bus

- Connects I/O devices to the CPU and memory
- Designed to be independent of the underlying CPU
- Can accommodate a wide variety of 3rd-party I/O devices
- Eg) Intel's PCI (Peripheral Component Interconnect) bus

■ Host bus adapter

- Connects one or more disks to the I/O bus
- Host bus interface
(Communication protocol for the connection)
 - SATA (supports only one drive)
 - SCSI (supports multiple drives, faster and more expensive)

Storage Technologies

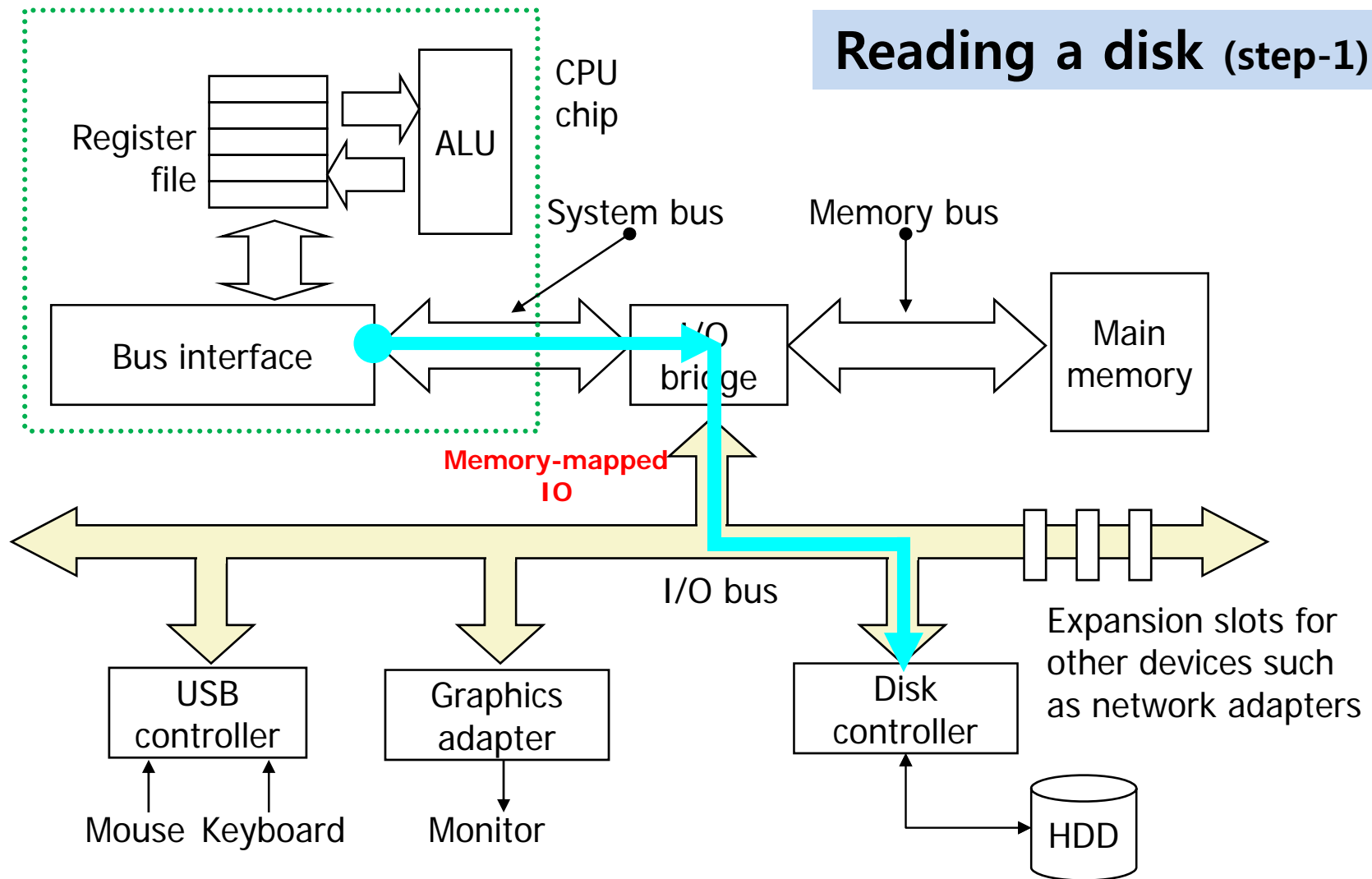


Storage Technologies

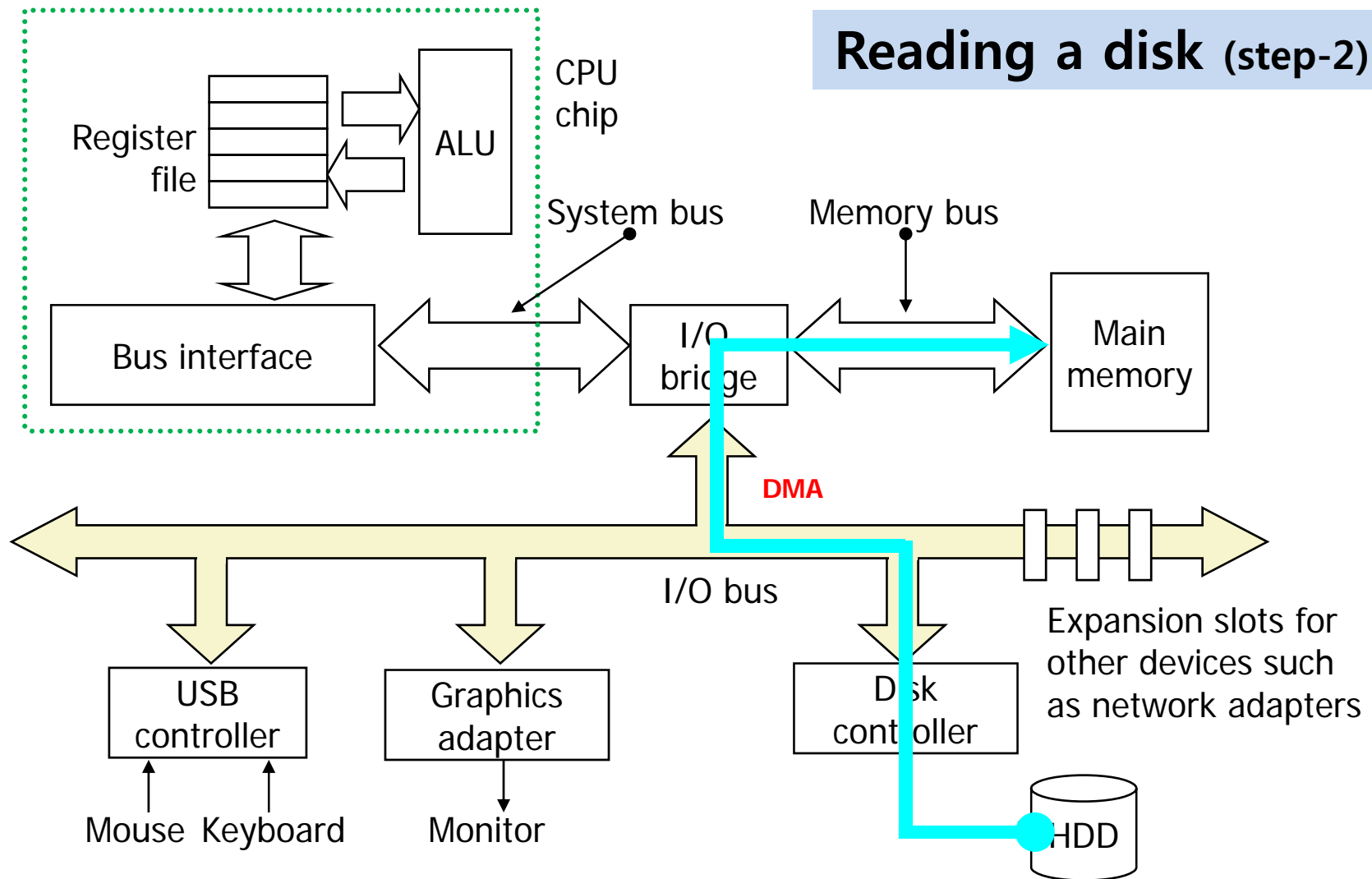
■ I/O bus

- Reading a disk sector
 - [Step-1] CPU initiates a disk read by writing a command, logical block number, and destination memory address to a **port** (address) associated with disk controller (**Memory-mapped I/O**)
 - [Step-2] Disk controller reads the sector and performs a **DMA (Direct Memory Access)** transfer into main memory
 - [Step-3] When the DMA transfer completes, the disk controller notifies the CPU with an **interrupt** (i.e., asserts a special “interrupt” pin on the CPU)

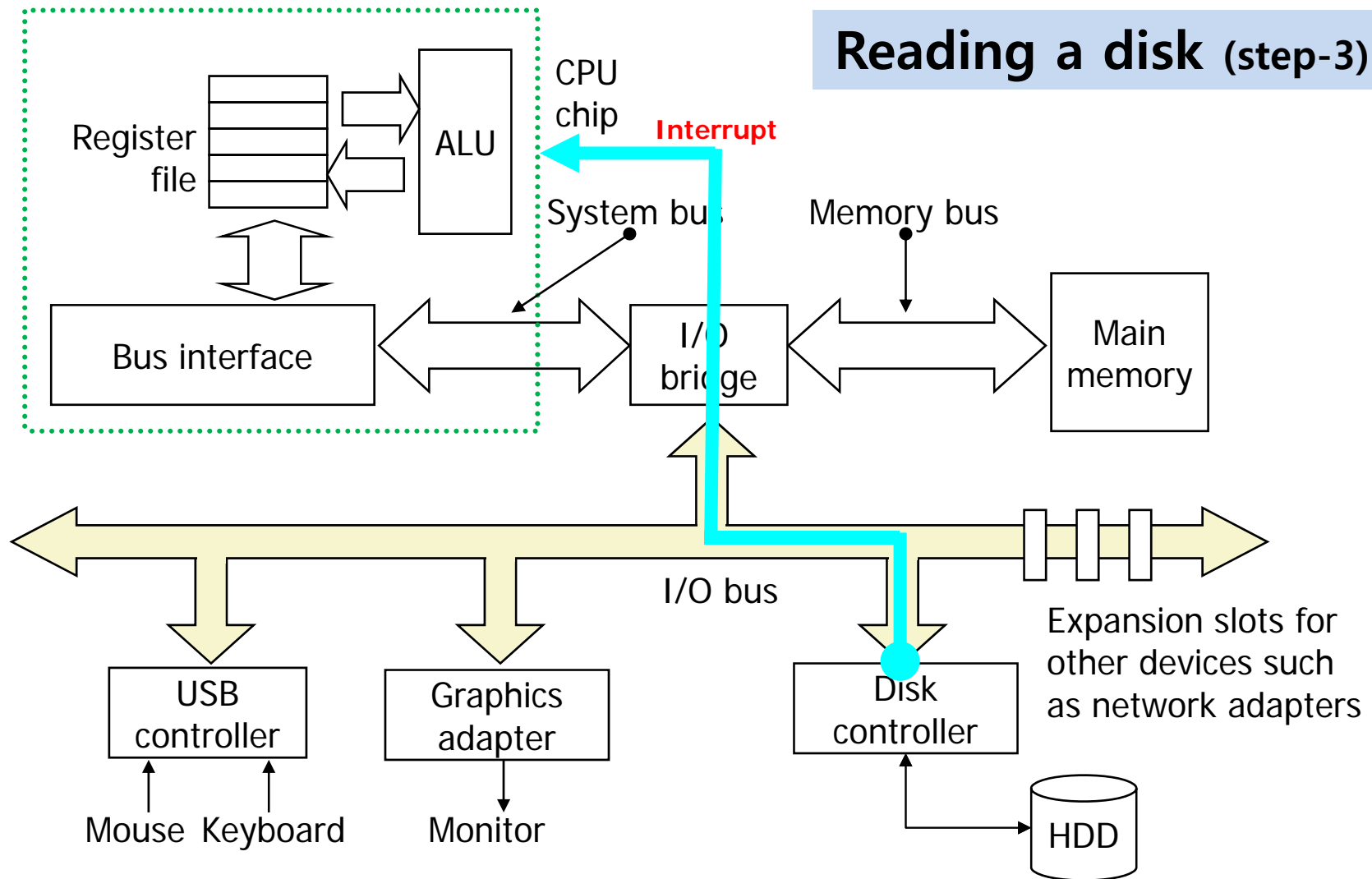
Storage Technologies



Storage Technologies



Storage Technologies



Storage Technologies

■ Commercial disks

■ Seagate Cheetah 15K.4

Geometry attribute	Value
Platters	4
Surfaces (read/write heads)	8
Surface diameter	3.5 in.
Sector size	512 bytes
<u>Zones</u>	<u>15</u>
Cylinders	50,864
Recording density (max)	628,000 bits/in.
Track density	85,000 tracks/in.
Areal density (max)	53.4 Gbits/sq. in.
Formatted capacity	146.8 GB

Performance attribute	Value
Rotational rate	15,000 RPM
Avg. rotational latency	2 ms
Avg. seek time	4 ms
Sustained transfer rate	58–96 MB/s

Storage Technologies

■ Tool DIXtrac

- Developed at CMU
- Automatically discovers a wealth of low-level information about the geometry and performance of SCSI disks
- Example)
DIXtrac on Seagate disk

Zone number	Sectors per track	Cylinders per zone	Logical blocks per zone
(outer) 0	864	3201	22,076,928
1	844	3200	21,559,136
2	816	3400	22,149,504
3	806	3100	19,943,664
4	795	3100	19,671,480
5	768	3400	20,852,736
6	768	3450	21,159,936
7	725	3650	21,135,200
8	704	3700	20,804,608
9	672	3700	19,858,944
10	640	3700	18,913,280
11	603	3700	17,819,856
12	576	3707	17,054,208
13	528	3060	12,900,096
(inner) 14	—	—	—

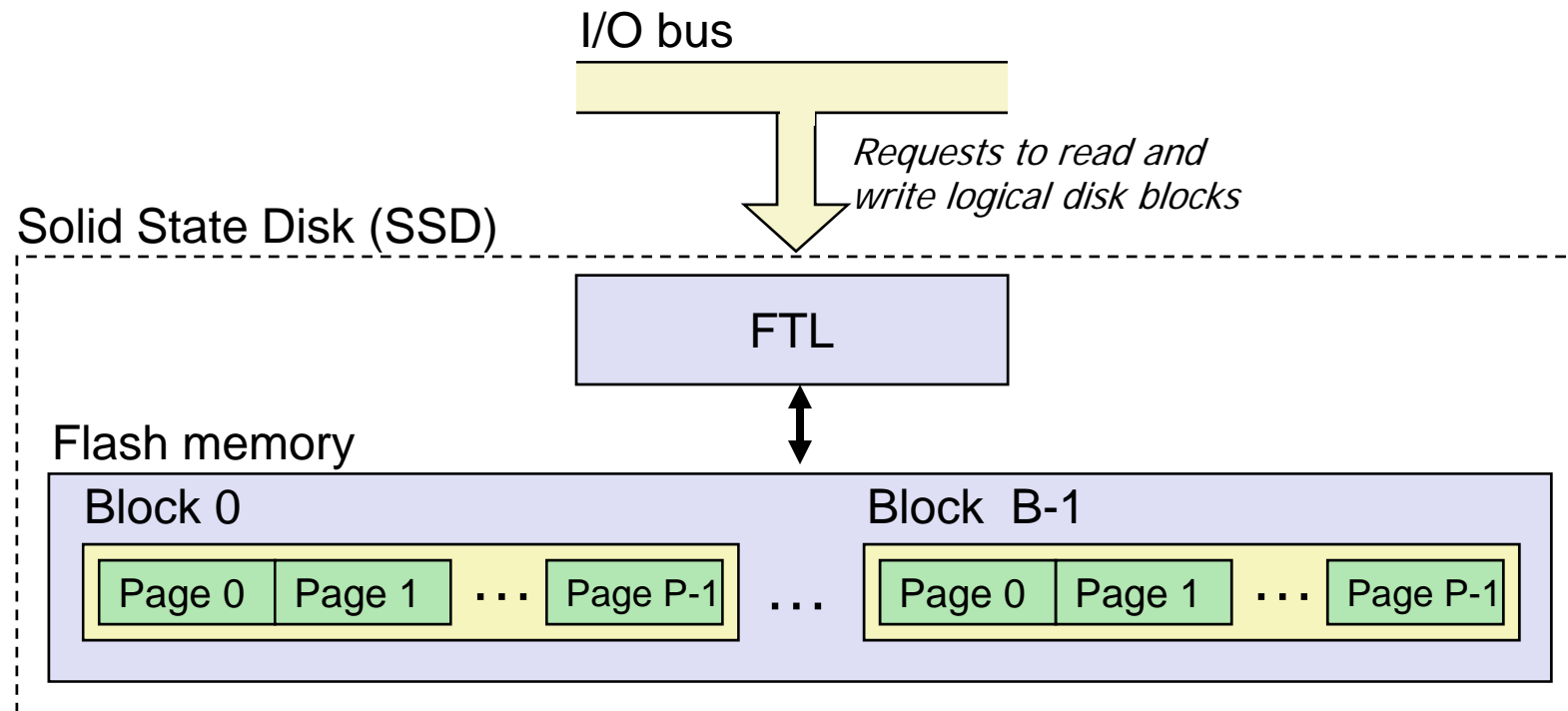
Storage Technologies

■ SSDs (Solid State Disks)

- Based on flash memory
 - Consists of one or more flash memory chips and FTL
 - FTL(Flash Translation Layer) plays the role of disk controller
 - ✓ Translating requests for logical blocks into accesses of the underlying physical device
- Plugs into the standard disk slot on the I/O bus (USB or SATA)
- Behaves like any other disk

Storage Technologies

■ SSDs



Storage Technologies

■ Characteristics of SSDs

- Page size
 - 512B ~ 4KB
- Block size
 - 32 ~ 128 pages (16KB ~ 512KB)
- Data is read and written in units of pages
 - A page can be written only after its block has been erased
- A block wears out after 100,000 repeated writes

Storage Technologies

■ Characteristics of SSDs

- Performance characteristics (Intel SSD 730)

Sequential read tput	550 MB/s	Sequential write tput	470 MB/s
Random read tput (IOPS)	89,000	Random write tput (IOPS)	74,000
Random read tput (MB/s)	365 MB/s	Random write tput (MB/s)	303 MB/s
Avg seq'l read access time	50 μ s	Avg seq'l write access time	60 μ s

- Why are random writes so slow?
 - Erasing a block is slow (around 1ms)
 - Writing a page triggers a copy of all useful pages in the block
 - ✓ Find a new block and erase it, if necessary
 - ✓ Write the page into the new block
 - ✓ Copy other pages from old block to the new block

Storage Technologies

■ Characteristics of SSDs

- Advantages
 - No moving parts → faster, less power, more rugged
- Disadvantages
 - Have the potential to wear out
 - ✓ Mitigated by “wear leveling logic” in FTL
 - About 30× more expensive per byte than in HDD
- Applications
 - Smart phones, laptops, digital cameras
 - Desktops and servers

Storage Technologies

■ Storage trends

SRAM

Metric	1985	1990	1995	2000	2005	2010	2015	2015:1985
\$/MB	2,900	320	256	100	75	60	25	116 ↓
access (ns)	150	35	15	3	2	1.5	1.3	115 ↓

DRAM

Metric	1985	1990	1995	2000	2005	2010	2015	2015:1985
\$/MB	880	100	30	1	0.1	0.06	0.02	44,000 ↓
access (ns)	200	100	70	60	50	40	20	10 ↓
typical size (MB)	0.256	4	16	64	2,000	8,000	16,000	62,500 ↑

HDD

Metric	1985	1990	1995	2000	2005	2010	2015	2015:1985
\$/GB	100,000	8,000	300	10	5	0.3	0.03	3,333,333 ↓
access (ms)	75	28	10	8	5	3	3	25 ↓
typical size (GB)	0.01	0.16	1	20	160	1,500	3,000	300,000 ↑

Storage Technologies

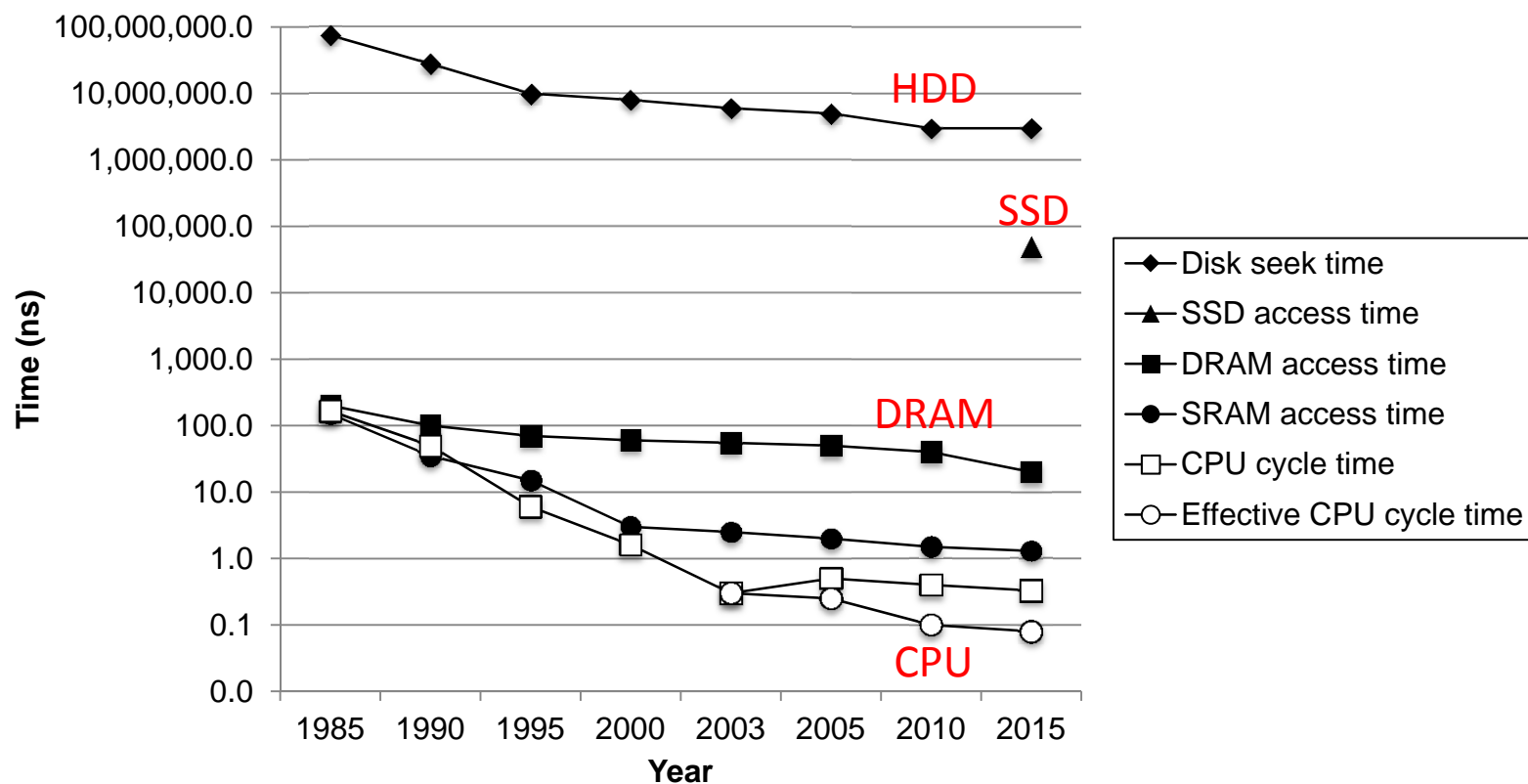
■ CPU clock rates

Inflection point in computer history
when designers hit the “Power Wall”

	1985	1990	1995	2003	2005	2010	2015	2015:1985
CPU	80286	80386	Pentium	P-4	Core 2	Core i7(n)	Core i7(h)	---
Clock rate (MHz)	6	20	150	3,300	2,000	2,500	3,000	500 ↑
Cycle time (ns)	166	50	6	0.3	0.5	0.4	0.33	500 ↓
Cores	1	1	1	1	2	4	4	4 ↑
Effective cycle time (ns)	166	50	6	0.3	0.25	0.1	0.08	2,075 ↓

Storage Technologies

■ The CPU-memory gap



Storage Technologies

■ Notes)

- The gap between DRAM - disk performance and CPU performance is widening
- Locality to the rescue!
 - The key to bridging this CPU-memory gap is a fundamental property of computer programs known as **locality**

Summary

