

[Chap.3-4] Machine-level Representation of Programs

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Contents

- Introduction
- Program encodings
- Data formats
- Intel processors
- Accessing information
- Primitive instructions
- Data movement instructions
- Arithmetic and logic instructions
- **Control instructions**
- Procedures
- ...

Control Instructions



■ C control structures

- Conditionals, loops, switches, ...
 - if, switch, for, while, do, ...

■ Machine-level control structures

- Jump instructions
 - Pass control to some other part of the program, possibly contingent on the result of some test

Control Instructions

■ Condition codes (CC)

- Single-bit registers
 - CF, ZF, SF, OF
- Implicitly set by arithmetic or logical operations
 - CF (carry flag)
 - ✓ Set if carry out from most significant bit
 - ✓ Used to detect unsigned overflow
 - ZF (zero flag)
 - ✓ Set if the most recent operation yielded 0
 - SF (sign flag)
 - ✓ Set if the most recent operation yielded a negative value
 - OF (overflow flag)
 - ✓ Set if the most recent operation caused 2's-complement overflow

Control Instructions

■ Condition codes (CC)

- Most arithmetic/logical instructions affects CCs
- The **leaq** instruction does not alter any CCs
- Some instructions affects the CCs in their own way
 - For logical operations such as **xor**,
 - ✓ The CF and OF are set to 0
 - For shift operations,
 - ✓ The CF is set to the last bit shifted out and OF is set to 0
 - For **inc** and **dec**,
 - ✓ The OF and ZF are set by the normal rule, leaving CF unchanged
 - Etc...

Control Instructions

■ Condition codes (CC)

- Explicit setting by compare (**CMP**) instructions
 - Example) **cmpq b, a**
 - ✓ Computes **(a – b)** without saving the result
 - ✓ CF set if carry out from most significant bit
 - Used for unsigned comparisons
 - ✓ ZF set if **a == b**
 - ✓ SF set if **(a – b) < 0**
 - ✓ OF set if two's complement overflow
 - **(a>0 && b<0 && (a-b)<0) || (a<0 && b>0 && (a-b)>0)**

Instruction		Based on	Description
CMP	s1, s2	s2 – s1	Compare
cmpb			Compare byte
cmpw			Compare word
cmpd			Compare double word
cmpq			Compare quad word

Control Instructions

■ Condition codes (CC)

- Explicit setting by test (**TEST**) instructions
 - Example) **testq b, a**
 - ✓ Sets CCs based on the value of **a** and **b**
 - Useful to have one of the operands as a mask
 - ✓ Computes **a & b** without setting destination
 - ✓ ZF set when **a & b == 0**
 - ✓ SF set when **a & b < 0**
 - ✓ CF and OF are cleared to **0**

Instruction		Based on	Description
TEST	s1, s2	s1 & s2	Test
testb			Test byte
testw			Test word
testl			Test double word
testq			Test quad word

Control Instructions

■ Accessing CCs

▪ **set** instructions

- Set a single byte destination (register or memory) to 0 or 1 depending on some combination of CCs

Instruction		Synonym	Effect	Description
sete	D	setz	$D \leftarrow ZF$	Equal / Zero
setne	D	setnz	$D \leftarrow \sim ZF$	Not Equal / Not Zero
sets	D		$D \leftarrow SF$	Negative
setns	D		$D \leftarrow \sim SF$	Nonnegative
setg	D	setnle	$D \leftarrow \sim(SF \wedge OF) \ \& \ \sim ZF$	Greater (Signed >)
setge	D	setnl	$D \leftarrow \sim(SF \wedge OF)$	Greater or Equal (Signed >=)
setl	D	setnge	$D \leftarrow (SF \wedge OF)$	Less (Signed <)
setle	D	setng	$D \leftarrow (SF \wedge OF) \mid ZF$	Less or Equal (Signed <=)
seta	D	setnbe	$D \leftarrow \sim CF \ \& \ \sim ZF$	Above (Unsigned >)
setae	D	setnb	$D \leftarrow \sim CF$	Above or Equal (Unsigned >=)
setb	D	setnae	$D \leftarrow CF$	Below (Unsigned <)
setbe	D	setna	$D \leftarrow CF \mid ZF$	Below or Equal (Unsigned <=)

Control Instructions

■ Accessing CCs

▪ **set** instructions

- Have a single-byte register or a single-byte memory location as its destination
- Does not alter remaining 3 bytes
- Typically uses **movzbl** or **movzbq** to finish job

Control Instructions

■ Accessing CCs

- Example) **set** instructions

```
[C code]
int comp(long a, long b){
    return a < b;
}
```

```
[Assembly code]
    int comp(long, long)
    a in %rdi, b in %rsi
comp:
    cmpq %rsi,%rdi      compare a:b
    setl %al            set low-order byte of %rax to 0 or 1
    movzbl %al,%eax     clear rest of %eax (and rest of %rax)
    ret
```

Note inverted ordering!

Control Instructions

■ Jump instructions

■ Unconditional jump

- Direct jump
 - ✓ Jump target encoded in the instruction (use label in assembly)
 - Eg) **jmp .L1**
- Indirect jump
 - ✓ Jump target in register or memory location
 - ✓ Use * followed by an operand specifier
 - Eg) **jmp *%eax**
jmp *(%eax)

```
movq $0,%rax
jmp .L1
movq (%rax),%rdx
.L1:
popq %rdx
```

■ Conditional jump

- Can only be direct

Control Instructions

■ Jump instructions

Instruction	Synonym	Condition	Description
jmp		1	Unconditional
je	jz	ZF	Equal / Zero
jne	jnz	~ZF	Not Equal / Not Zero
js		SF	Negative
jns		~SF	Nonnegative
jg	jnle	~(SF ^ OF) & ~ZF	Greater (Signed >)
jge	jnl	~(SF ^ OF)	Greater or Equal (Signed >=)
jl	jnge	(SF ^ OF)	Less (Signed <)
jle	jng	(SF ^ OF) ZF	Less or Equal (Signed <=)
ja	jnb	~CF & ~ZF	Above (Unsigned >)
jae	jnb	~CF	Above or Equal (Unsigned >=)
jb	jnae	CF	Below (Unsigned <)
jbe	jna	CF ZF	Below or Equal (Unsigned <=)

Control Instructions

■ Jump instructions

- Example) PC-relative addressing → ± 128 byte 내의 범위에서 한정

```
movq %rdi,%rax
jmp .L2
.L3:
sarq %rax
.L2:
testq %rax,%rax
jg .L3
rep; ret
```

Assembly
code

Disassembled
version

```
0: 48 89 f8
3: eb 03
5: 48 d1 f8
8: 48 85 c0
b: 7f f8
d: f3 c3
```

+3

-8

```
mov %rdi,%rax
jmp 8 <loop+0x8>
sar %rax
test %rax,%rax
jg 5 <loop+0x5>
repz retq
```

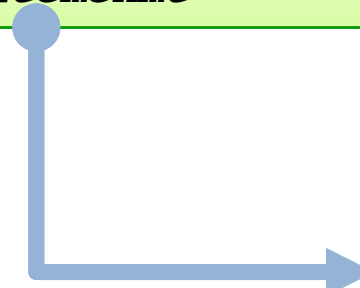
Control Instructions

■ Translating conditional branches

- Uses conditional and unconditional branches
- Typical translations

```
if (test-expr)
    then-statement
else
    else-statement
```

goto version



```
t = test_expr;
if (!t)
    goto false;
then-statement
goto done;
false:
    else-statement
done:
```

Control Instructions

■ Translating conditional branches

■ Example)

[Original C code]

```
long lt_cnt = 0;
long ge_cnt = 0;

long adiff_se(long x, long y)
{
    long result;
    if (x < y) {
        lt_cnt++;
        result = y - x;
    }
    else {
        ge_cnt++;
        result = x - y;
    }
    return result;
}
```

[goto version]

```
long lt_cnt = 0;
long ge_cnt = 0;

long gdiff_se(long x, long y)
{
    long result;
    if (x >= y)
        goto x_ge_y;
    lt_cnt++;
    result = y - x;
    return result;
x_ge_y:
    ge_cnt++;
    result = x - y;
    return result;
}
```

Control Instructions

■ Translating conditional branches

■ Example)

[Assembly code]

```
long adiff_se(long x, long y)
```

```
x in %rdi, y in %rsi
```

adiff_se:

```
    cmpq %rsi,%rdi
```

Compare x:y

```
    jge .L2
```

If >= goto x_ge_y

```
    addq $1,lt_cnt(%rip)
```

lt_cnt++

```
    movq %rsi,%rax
```

```
    subq %rdi,%rax
```

result = y - x

```
    ret
```

Return

.L2:

x_ge_y:

```
    addq $1,ge_cnt(%rip)
```

ge_cnt++

```
    movq %rdi,%rax
```

```
    subq %rsi,%rax
```

result = x - y

```
    ret
```

Return

Control Instructions

■ Conditional move instructions

- Conditional transfer of data
 - Either copy a value to a register or do nothing, depending on the values of the CCs
- Introduced from Pentium Pro (1995)
- **gcc** generates code using conditional moves, if possible (64-bit versions of Linux and Windows)

Control Instructions

■ Conditional move instructions

- Two operands (16/32/64-bits long)
 - First operand: source register or memory location
 - Second operand: destination register

Instruction	Synonym	Move condition	Description
<code>cmovz</code> <i>S, R</i>	<code>cmovz</code>	ZF	Equal / zero
<code>cmovne</code> <i>S, R</i>	<code>cmovnz</code>	$\neg ZF$	Not equal / not zero
<code>cmovs</code> <i>S, R</i>		SF	Negative
<code>cmovns</code> <i>S, R</i>		$\neg SF$	Nonnegative
<code>cmovg</code> <i>S, R</i>	<code>cmovnl</code>	$\neg(SF \wedge OF) \wedge \neg ZF$	Greater (signed >)
<code>cmovge</code> <i>S, R</i>	<code>cmovnl</code>	$\neg(SF \wedge OF)$	Greater or equal (signed >=)
<code>cmovl</code> <i>S, R</i>	<code>cmovnge</code>	$SF \wedge OF$	Less (signed <)
<code>cmovle</code> <i>S, R</i>	<code>cmovng</code>	$(SF \wedge OF) \vee ZF$	Less or equal (signed <=)
<code>cmova</code> <i>S, R</i>	<code>cmovnbe</code>	$\neg CF \wedge \neg ZF$	Above (unsigned >)
<code>cmovae</code> <i>S, R</i>	<code>cmovnb</code>	$\neg CF$	Above or equal (Unsigned >=)
<code>cmovb</code> <i>S, R</i>	<code>cmovnae</code>	CF	Below (unsigned <)
<code>cmovbe</code> <i>S, R</i>	<code>cmovna</code>	$CF \vee ZF$	below or equal (unsigned <=)

Control Instructions

■ Conditional move instructions

▪ Typical translations

```
v = test_expr ? then_expr : else_expr
```

```
if (!test_expr)  
    goto false;  
    v = then_expr;  
    goto done;  
false:  
    v = else_expr;  
done:
```

Conditional assignment version

```
v = then_expr;  
ve = else_expr;  
t = test_expr;  
if (!t) v = ve;
```

Control Instructions

■ Conditional move instructions

■ Example)

[Original C code]

```
long adiff(long x, long y)
{
    long result;
    if (x < y)
        result = y - x;
    else
        result = x - y;
    return result;
}
```

[conditional assignment version]

```
long cdiff(long x, long y)
{
    long rval = y - x;
    long eval = x - y;
    long ntest = x >= y;
    if (ntest) rval = eval;
    return rval;
}
```

Control Instructions

■ Conditional move instructions

■ Example)

[Assembly code]

```
long adiff(long x, long y)
x in %rdi, y in %rsi
```

adiff:

```
movq %rsi,%rax
```

```
subq %rdi,%rax
```

```
movq %rdi,%rdx
```

```
subq %rsi,%rdx
```

```
cmpq %rsi,%rdi
```

```
cmovge %rdx,%rax
```

```
ret
```

```
rval = y - x
```

```
eval = x - y
```

```
Compare x : y
```

```
if >=, rval = eval
```

```
Return rval
```

Control Instructions

■ Conditional move instructions

- Not all conditional expressions can be compiled using conditional moves
 - The abstract code evaluates both *then-expr* and *else-expr* regardless of the test outcome
 - If one of those two expressions could possibly generate an error condition or a side effect,
 - ✓ This could lead to invalid behavior

Control Instructions

■ Conditional move instructions

- Not all conditional expressions can be compiled using conditional moves

[C code]

```
long cread(long *xp){  
    return (xp ? *xp : 0);  
}
```

[Assembly code]

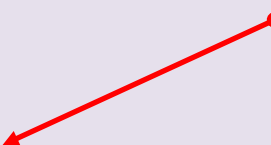
```
long cread(long *xp)  
xp in %rdi
```

cread:

```
    movl $0,%edx  
    movq (%rdi),%rax  
    testq %rdi,%rdi  
    cmovle %rdx,%rax  
    ret
```

```
ve = 0  
v = *xp  
Test x  
If x==0, v = ve  
Return v
```

Null pointer
dereferencing
when xp == 0



Control Instructions

■ Conditional move instructions

- The code using conditional data transfers can outperform the code based on conditional control transfers
 - We should understand modern processor architectures
 - ✓ Piplining
 - ✓ Branch prediction
 - ✓ Out-of-order execution
 - ✓ Etc...
- But, using conditional moves does not always improve the code efficiency

Control Instructions

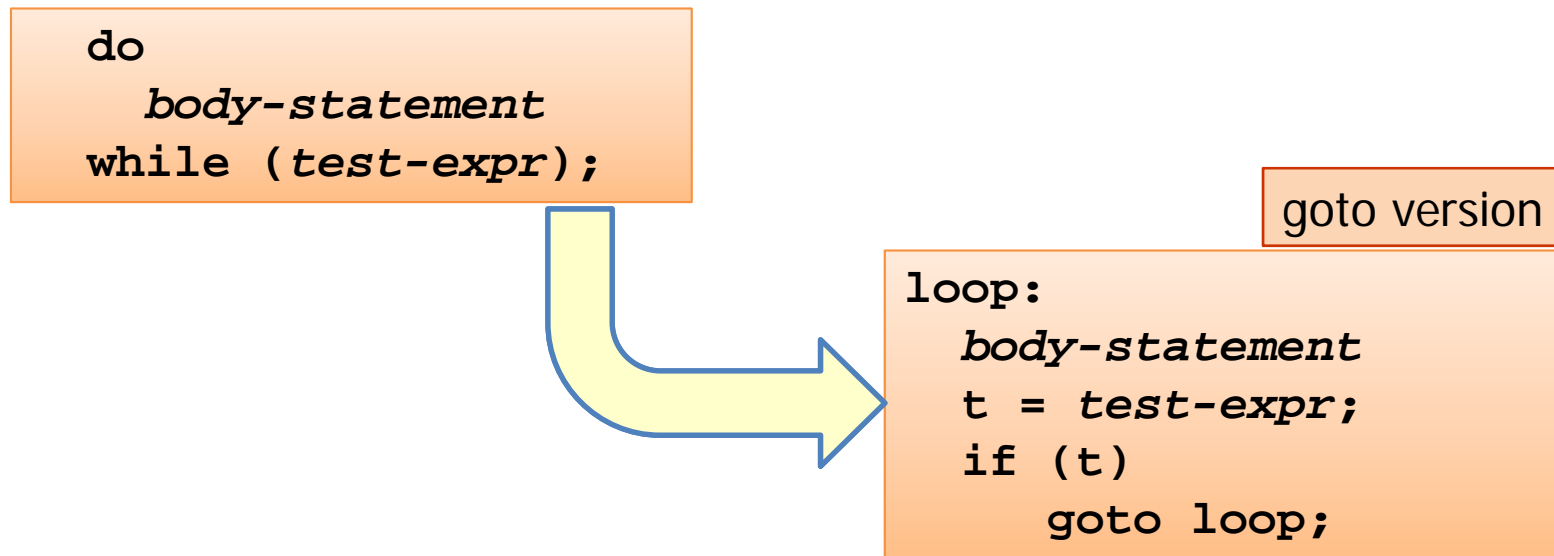
■ Types of loops

- do-while
 - while
 - for
-
- **gcc** and other compilers generate loop code based on the two basic loop patterns

Control Instructions

■ Translating do-while loops

- Uses combinations of conditional tests and jumps
- Used as a base form in generating loop code (by most compilers)



Control Instructions

■ Translating do-while loops

- Example) Factorial ($n!$ for $n > 0$)

C code

```
long fact_do (long n)
{
    long result = 1;
    do {
        result *= n;
        n = n-1;
    } while (n > 1);
    return result;
}
```

goto version

```
long fact_do_goto (long n)
{
    long result = 1;
loop:
    result *= n;
    n = n-1;
    if (n > 1)
        goto loop;
    return result;
}
```

Control Instructions

■ Translating do-while loops

■ Example)

Goto version

```
long fact_do_goto (long n)
{
    long result = 1;
loop:
    result *= n;
    n = n-1;
    if (n > 1)
        goto loop;
    return result;
}
```

[Registers]

%rdi	n
%rax	result

Assembly code

long fact_do(long n)	
n in %rdi	
fact_do:	
movl \$1,%eax	Set result = 1
.L2:	loop:
imulq %rdi,%rax	Compute result *= n
subq \$1,%rdi	Decrement n
cmpq \$1,%rdi	Compare n:1
jg .L2	If >, goto loop
rep; ret	Return

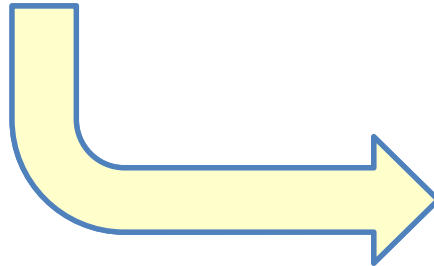
Control Instructions

■ Translating while loops

- **Jump-to-middle** method (`gcc -Og`)

C code

```
while (test-expr)  
    body-statement
```



goto version

```
goto test  
loop:  
    body-statement  
test:  
    if (test-expr)  
        goto loop;
```

Control Instructions

■ Translating while loops

- Example) Factorial ($n!$ for $n \geq 0$)

C code

```
long fact_while (long n)
{
    long result = 1;
    while (n > 1) {
        result *= n;
        n = n-1;
    }
    return result;
}
```

goto version

```
long fact_while_jm (long n)
{
    long result = 1;
    goto test;
loop:
    result *= n;
    n = n-1;
test:
    if (n > 1)
        goto loop;
    return result;
}
```

Control Instructions

■ Translating while loops

- Example) Factorial ($n!$ for $n \geq 0$)

Goto version

```
long fact_while_jm (long n)
{
    long result = 1;
    goto test;
loop:
    result *= n;
    n = n-1;
test:
    if (n > 1)
        goto loop;
    return result;
}
```

[Registers]

%rdi	n
%rax	result

Assembly code

```
long fact_while(long n)
n in %rdi
fact_while:
    movl $1,%eax        Set result = 1
    jmp .L5             Goto test
.L6:                   loop:
    imulq %rdi,%rax      Compute result *= n
    subq $1,%rdi         Decrement n
.L5:                   test:
    cmpq $1,%rdi         Compare n:1
    jg .L6              If >, goto loop
    rep; ret            Return
```

Control Instructions

■ Translating while loops

■ **Guarded-do** method (**gcc -O1**)

C code

```
while (test-expr)
    body-statement
```

→ 차이는 피아프랑스라고 jump to middle 방식보다
while 조건에 있어 실행시간이 더 적은 방식

do-while version

```
if (!test-expr)
    goto done;
do
    body-statement
while (test-expr);
done:
```

goto version

```
if (!test-expr)
    goto done;
loop:
    body-statement
    if (test-expr)
        goto loop;
done:
```


Control Instructions

■ Translating while loops

- Example) Factorial ($n!$ for $n \geq 0$)

C code

```
long fact_while (long n)
{
    long result = 1;
    while (n > 1) {
        result *= n;
        n = n-1;
    }
    return result;
}
```

goto version

```
long fact_while_gd (long n)
{
    long result = 1;
    if (n <= 1)
        goto done;
loop:
    result *= n;
    n = n-1;
    if (n != 1)
        goto loop;
done:
    return result;
}
```

Control Instructions

■ Translating while loops

- Example) Factorial ($n!$ for $n \geq 0$)

Goto version

```
long fact_while_gd (long n)
{
    long result = 1;
    if (n <= 1)
        goto done;
loop:
    result *= n;
    n = n-1;
    if (n != 1)
        goto loop;
done:
    return result;
}
```

Assembly code

```
long fact_while(long n)
n in %rdi
fact_while:
    movl $1,%eax        Set result = 1
    cmpq $1,%rdi        Compare n:1
    jle .L7             If <=, goto done
.L6:                   loop:
    imulq %rdi,%rax     Compute result *= n
    subq $1,%rdi        Decrement n
    cmpq $1,%rdi        Compare n:1
    jne .L6             If !=, goto loop
    rep; ret            Return
.L7:                   done:
    ret                 Return
```

Control Instructions

■ Translating for loops

- **Jump-to-middle** method

for version

```
for (init; test; update )  
    body-statement
```

while version

```
init;  
while (test) {  
    body-statement  
    update;  
}
```

goto version

```
init;  
goto test1;  
loop:  
    body-statement  
    update;  
test1:  
    if (test)  
        goto loop;
```

Control Instructions

■ Translating for loops

■ Guarded-do method

for version

```
for (init; test; update )  
    body-statement
```

while version

```
init;  
while (test) {  
    body-statement  
    update;  
}
```

do-while version

```
init;  
if (!test)  
    goto done;  
do {  
    body-statement  
    update;  
} while (test)  
done:
```

goto version

```
init;  
if (!test)  
    goto done;  
loop:  
    body-statement  
    update;  
    if (test)  
        goto loop;  
done:
```

Control Instructions

■ Translating for loops

- Example) Factorial ($n!$ for $n \geq 0$)

C code

```
long fact_for(long n)
{
    long i;
    long result = 1;
    for (i=2; i<=n; i++)
        result *= i;
    return result;
}
```

while version

```
long fact_for_while (long n)
{
    long result = 1;
    long i = 2;
    while (i <= n) {
        result *= i;
        i++;
    }
    return result;
}
```

Control Instructions

■ Translating for loops

- Example) Factorial ($n!$ for $n \geq 0$)

while version

```
long fact_for_while (long n)
{
    long result = 1;
    long i = 2;
    while (i <= n) {
        result *= i;
        i++;
    }
    return result;
}
```

goto version (jm)

```
long fact_for_jm (long n)
{
    long result = 1;
    long i = 2;
    goto test;
loop:
    result *= i;
    i++;
test:
    if (i <= n)
        goto loop;
    return result;
}
```

Control Instructions

■ Translating for loops

- Example) Factorial ($n!$ for $n \geq 0$)

Goto version

```
long fact_for_jm (long n)
{
    long result = 1;
    long i = 2;
    goto test;
loop:
    result *= i;
    i++;
test:
    if (i <= n)
        goto loop;
    return result;
}
```

Assembly code

```
long fact_for(long n)
n in %rdi

fact_for:
    movl $1,%eax        Set result = 1
    movl $2,%edx        Set i = 2
    jmp .L8             Goto test
.L9:                    loop:
    imulq %rdx,%rax      Compute result *= i
    addq $1,%rdx         Increment i
.L8:                    test:
    cmpq %rdi,%rdx       Compare i:n
    jle .L9              If <=, goto loop
    rep; ret             Return
```

Control Instructions

■ Translating switches

- Multi-way branching capability
- Efficient implementation is possible using a data structure called a **jump table**
 - **Jump table**
 - ✓ An array where entry i is the address of a code segment implementing the action the program should take when the switch index equals i
 - More efficient than a long sequence **if-else** in that the time taken to perform the switch is independent of the number of switch cases

Control Instructions

■ Translating switches

- **gcc** selects the method of translating a switch statement based on the number of cases and the sparsity of the case values

Control Instructions

■ Translating switches

- Jump table structure

Switch Form

```
switch (op) {  
  case val_0:  
    Block 0  
  case val_1:  
    Block 1  
    . . .  
  case val_n-1:  
    Block n-1  
}
```

Jump Table

jtab

Targ_0
Targ_1
Targ_2
•
•
•
Targ_n-1

Jump Targets

Targ_0:

Code block
0

Targ_1:

Code block
1

•
•
•

Targ_n-1:

Code block
n-1

Approx. Translation

```
target = jtab[op];  
goto *target;
```

Control Instructions

■ Translating switches: Example)

```
void switch_eg (long x, long n, long *d)
{
    long v = x;
    switch(n) {
        case 100:
            v *= 13;
            break;
        case 102:
            v += 10;
            /* Fall through */
        case 103:
            v += 11;
            break;
        case 104:
        case 106:
            v *= v;
            break;
        default:
            v = 0;
    }
    *d = v;
}
```

```
void switch_eg_impl(long x, long n, long *d)
{
    /* Table of code pointers */
    static void *jt[7] = {
        &l_A, &l_def, &l_B, &l_C,
        &l_D, &l_def, &l_D};
    unsigned long index = n - 100;
    long v;

    if (index > 6)
        goto l_def;
    goto *jt[index]; /* Multiway branch */

l_A: /* case 100 */
    v = x * 13;
    goto done;
l_B: /*case 102 */
    v = x + 10;
    /* Fall through */
l_C: /* case 103 */
    v = x + 11;
    goto done;
l_D: /* case 104, 106 */
    v = x * x;
    goto done;
l_def:
    v = 0;
done:
    *d = v;
}
```

Control Instructions

■ Translating switches: Example)

Assembly code

```
void switch_eg_impl(long x, long n, long *d)
{
    /* Table of code pointers */
    static void *jt[7] = {
        &l_A, &l_def, &l_B, &l_C,
        &l_D, &l_def, &l_D};
    unsigned long index = n - 100;
    long v;

    if (index > 6)
        goto l_def;
    goto *jt[index]; /* Multiway branch */

l_A:    /* case 100 */
    v = x * 13;
    goto done;
l_B:    /*case 102 */
    v = x + 10;
    /* Fall through */
l_C:    /* case 103 */
    v = x + 11;
    goto done;
l_D:    /* case 104, 106 */
    v = x * x;
    goto done;
l_def:
    v = 0;
done:
    *d = v;
}
```

```
void switch_eg(long x, long n, long *d)
x in %rdi, n in %rsi, d in %rdx
```

switch_eg:

subq \$100,%rsi

Compute index = n-100

cmpq \$6,%rsi

Compare index:6

ja .L8
jmp *.L4(,%rsi,8)

If >, goto loc_def
Goto *jt[index]

.L3:

loc_A:

leaq (%rdi,%rdi,2),%rax

3*x

leaq (%rdi,%rax,4),%rdi

v = 13*x

jmp .L2

Goto done

.L5:

loc_B:

addq \$10,%rdi

v = x + 10

.L6:

loc_C:

addq \$11,%rdi

v = x + 11

jmp .L2

Goto done

.L7:

loc_D:

imulq %rdi,%rdi

v = x * x

jmp .L2

Goto done

.L8:

loc_def:

movl \$0,%edi

v = 0

.L2:

done:

movq %rdi,(%rdx)

*dest = v

ret

Return

Control Instructions

■ Translating switches: Example)

■ Jumping

- `jmp *.L4(,%rsi,8)`
 - ✓ Start of jump table denoted by label `.L4`
 - ✓ Register `%rsi` holds `index`
 - Must scale by factor of 8 to get offset into the table
 - ✓ Fetch target from effective address `.L4 + index * 8`

Control Instructions

■ Translating switches: Example) jump table

```
.section .rodata
    .align 8                Align address to multiple of 8
.L4:
    .quad .L3               Case 100: l_A
    .quad .L8               Case 101: l_def
    .quad .L5               Case 102: l_B
    .quad .L6               Case 103: l_C
    .quad .L7               Case 104: l_D
    .quad .L8               Case 105: l_def
    .quad .L7               Case 106: l_D
```

Control Instructions

■ Translating switches

■ Note) **sparse switch**

- Not practical to use jump table
 - ✓ Would require 1000 entries
- Translation into if-else would have maximum 9 tests

```
int div(int x) {  
    switch (x) {  
        case 0: return 0;  
        case 111: return 1;  
        case 232: return 2;  
        case 356: return 3;  
        case 389: return 4;  
        case 560: return 5;  
        case 682: return 6;  
        case 877: return 7;  
        case 899: return 8;  
        case 900: return 9;  
        default: return -1;  
    }  
}
```

7338 X → index mapping이 불가

Summary

