

[Chap.2-1] Representing and Manipulating Information

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Contents

- Introduction
- Information storage
- Integer representations
- Integer arithmetic
- Floating point
- Summary

Introduction

■ Bit

- Binary digit
- 0/1

■ Grouping bits together

- Can represent the elements of any finite set
- n bits $\rightarrow 2^n$ representations
(different meanings for each data type)
 - Integer
 - ✓ Unsigned, 2's-complement, ...
 - Floating-point number
 - Character
 - Etc
 - By studying the actual number representations, we can understand the ranges of values that can be represented and the properties of the different arithmetic operations

Information Storage

■ Byte → char

- Binary Term
- 8-bits
- Smallest addressable unit of memory★

Information Storage

■ Hexadecimal notation

Hex digit	0	1	2	3	4	5	6	7
Decimal value	0	1	2	3	4	5	6	7
Binary value	0000	0001	0010	0011	0100	0101	0110	0111
Hex digit	8	9	A	B	C	D	E	F
Decimal value	8	9	10	11	12	13	14	15
Binary value	1000	1001	1010	1011	1100	1101	1110	1111

- Single byte: $00_{16} \sim FF_{16}$
- Converting between binary and hexadecimal
 - Straightforward
- Converting between decimal and hexadecimal
 - Requires some computation [Refer to text, Section 2.1.1]

Information Storage

■ Word

- Nominal size of integer and pointer data
 - Virtual address is encoded by a word
 - ✓ Maximum size of the virtual address space is determined by the word size
- Typical word size today
 - 32/64 -bits
 - ✓ 4GB virtual address space

Information Storage

■ Multiple data formats and data sizes

■ Multiple data formats

- Different ways to encode data, different lengths
- Sizes of C numeric data types

C declarations		Bytes	
Signed	Unsigned	32-bit M	64-bit M
[signed] char	unsigned char	1	1
short	unsigned short	2	2
int	unsigned	4	4
long	unsigned long	4	8
int32_t	uint32_t	4	4
int64_t	uint64_t	8	8
char *		4	8
float		4	4
double		8	8

Information Storage

■ Byte ordering

- Whether to choose to store the object in memory ordered from least significant byte to most or from most to least
- Little endian
 - The least significant byte (LSB) comes first
 - Used in most Intel-compatible machines
- Big endian
 - The most significant byte (MSB) comes first
 - Used in most machines from IBM and SUN Microsystems
- Bi-endian
 - Can be configured to operate as either little- or big-endian machines

Information Storage

■ Byte ordering

- Example) Data 0x01234567 at address 0x100

Big endian

	0x100	0x101	0x102	0x103	
...	01	23	45	67	...

Little endian

	0x100	0x101	0x102	0x103	
...	67	45	23	01	...

Information Storage

■ Strings

- Array of characters terminated by a NULL character
- ASCII character code
 - Decimal digits
 - ✓ '0' ~ '9' → 0x30 ~ 0x39
 - Alphabets (lower case)
 - ✓ 'a' ~ 'z' → 0x61 ~ 0x7A
 - NULL character
 - ✓ NULL → 0x00

ASCII Code Table

Dec	Hx	Oct	Char	Dec	Hx	Oct	Html	Chr	Dec	Hx	Oct	Html	Chr	Dec	Hx	Oct	Html	Chr
0	0	000	NUL (null)	32	20	040	 	Space	64	40	100	@	@	96	60	140	`	`
1	1	001	SOH (start of heading)	33	21	041	!	!	65	41	101	A	A	97	61	141	a	a
2	2	002	STX (start of text)	34	22	042	"	"	66	42	102	B	B	98	62	142	b	b
3	3	003	ETX (end of text)	35	23	043	#	#	67	43	103	C	C	99	63	143	c	c
4	4	004	EOT (end of transmission)	36	24	044	$	\$	68	44	104	D	D	100	64	144	d	d
5	5	005	ENQ (enquiry)	37	25	045	%	%	69	45	105	E	E	101	65	145	e	e
6	6	006	ACK (acknowledge)	38	26	046	&	&	70	46	106	F	F	102	66	146	f	f
7	7	007	BEL (bell)	39	27	047	'	'	71	47	107	G	G	103	67	147	g	g
8	8	010	BS (backspace)	40	28	050	((72	48	110	H	H	104	68	150	h	h
9	9	011	TAB (horizontal tab)	41	29	051))	73	49	111	I	I	105	69	151	i	i
10	A	012	LF (NL line feed, new line)	42	2A	052	*	*	74	4A	112	J	J	106	6A	152	j	j
11	B	013	VT (vertical tab)	43	2B	053	+	+	75	4B	113	K	K	107	6B	153	k	k
12	C	014	FF (NP form feed, new page)	44	2C	054	,	,	76	4C	114	L	L	108	6C	154	l	l
13	D	015	CR (carriage return)	45	2D	055	-	-	77	4D	115	M	M	109	6D	155	m	m
14	E	016	SO (shift out)	46	2E	056	.	.	78	4E	116	N	N	110	6E	156	n	n
15	F	017	SI (shift in)	47	2F	057	/	/	79	4F	117	O	O	111	6F	157	o	o
16	10	020	DLE (data link escape)	48	30	060	0	0	80	50	120	P	P	112	70	160	p	p
17	11	021	DC1 (device control 1)	49	31	061	1	1	81	51	121	Q	Q	113	71	161	q	q
18	12	022	DC2 (device control 2)	50	32	062	2	2	82	52	122	R	R	114	72	162	r	r
19	13	023	DC3 (device control 3)	51	33	063	3	3	83	53	123	S	S	115	73	163	s	s
20	14	024	DC4 (device control 4)	52	34	064	4	4	84	54	124	T	T	116	74	164	t	t
21	15	025	NAK (negative acknowledge)	53	35	065	5	5	85	55	125	U	U	117	75	165	u	u
22	16	026	SYN (synchronous idle)	54	36	066	6	6	86	56	126	V	V	118	76	166	v	v
23	17	027	ETB (end of trans. block)	55	37	067	7	7	87	57	127	W	W	119	77	167	w	w
24	18	030	CAN (cancel)	56	38	070	8	8	88	58	130	X	X	120	78	170	x	x
25	19	031	EM (end of medium)	57	39	071	9	9	89	59	131	Y	Y	121	79	171	y	y
26	1A	032	SUB (substitute)	58	3A	072	:	:	90	5A	132	Z	Z	122	7A	172	z	z
27	1B	033	ESC (escape)	59	3B	073	;	;	91	5B	133	[[123	7B	173	{	{
28	1C	034	FS (file separator)	60	3C	074	<	<	92	5C	134	\	\	124	7C	174	|	
29	1D	035	GS (group separator)	61	3D	075	=	=	93	5D	135]]	125	7D	175	}	}
30	1E	036	RS (record separator)	62	3E	076	>	>	94	5E	136	^	^	126	7E	176	~	~
31	1F	037	US (unit separator)	63	3F	077	?	?	95	5F	137	_	_	127	7F	177		DEL

Source: www.LookupTables.com

** EBCDIC 코드 참고

ASCII Code Table



Friendly	Numeric	Description	Friendly	Numeric	Description	Friendly	Numeric	Description	Friendly	Numeric	Description
‘		left single quote	R	R	uppercase letter	·	·	middle dot			
’		right single quote	S	R	uppercase letter	¸	¸	cedilla			
‚		single low-9 quote	T	R	uppercase letter	¹	¹	superscript one			
“		left double quote	U	R	uppercase letter	º	º	masculine			
”		right double quote	V	R	uppercase letter	»	»	right angle			
„		double low-9 quote	W	V	uppercase letter	¼	¼	one-fourth			
†		dagger	X	X	uppercase letter	½	½	one-half			
‡		double dagger	Y	Y	uppercase letter	¾	¾	three-fourths			
‰		per mill sign	Z	Z	uppercase letter	?	¿	inverted			
‹		left angle quote	[[left square bracket	À	À	A, grave			
›		right angle quote	\	\	backslash	&Acute;	Á	A, acute accent			
♠		black spade suit]]	right square bracket	Â	Â	A, circumflex			
♣		black club suit	^	^	caret	Ã	Ã	A, tilde			
♥		black heart suit	_	_	underscore	Ä	Ä	A, umlaut			
♦		black diamond suit	`	`	grave accent	Å	Å	A, ring			
‾		overline	a	a	lowercase letter	Æ	Æ	AE			
←		leftward arrow	b	b	lowercase letter	Ç	Ç	C, cedilla			
↑		upward arrow	c	c	lowercase letter	È	È	E, grave accent			
→		rightward arrow	d	d	lowercase letter	É	É	E, acute accent			
↓		downward arrow	e	e	lowercase letter	Ê	Ê	E, circumflex			
™		trademark sign	f	f	lowercase letter	Ë	Ë	E, umlaut			
			horizontal tab	g	g	lowercase letter	&lgrave;	Ì	l, grave accent			
	
	line feed	h	h	lowercase letter	í	Í	i, acute accent			
	 	space	i	i	lowercase letter	î	Î	i, circumflex			
"	!	exclamation mark	j	j	lowercase letter	ï	Ï	i, umlaut			
	"	double quotation	k	k	lowercase letter	Ð	Ð	Eth, Icelandic			
	#	number sign	l	l	lowercase letter	Ñ	Ñ	N, tilde			
	$	dollar sign	m	m	lowercase letter	Ò	Ò	O, grave			
&	%	percent sign	n	n	lowercase letter	Ó	Ó	O, acute accent			
	&	ampersand	o	o	lowercase letter	Ô	Ô	O, circumflex			
	'	apostrophe	p	p	lowercase letter	Õ	Õ	O, tilde			
	(left parenthesis	q	q	lowercase letter	Ö	Ö	O, umlaut			
)	right parenthesis	r	r	lowercase letter	×	×	multiplication			
	*	asterisk	s	s	lowercase letter	Ø	Ø	O, slash			
	+	plus sign	t	t	lowercase letter	Ù	Ù	U, grave accent			
	,	comma	u	u	lowercase letter	Ú	Ú	U, acute accent			
	-	hyphen	v	v	lowercase letter	Û	Û	U, circumflex			
	.	period	w	w	lowercase letter	Ü	Ü	U, umlaut			
&frac;	/	slash	x	x	lowercase letter	Ý	Ý	Y, acute accent			
	0	0	y	y	lowercase letter	Þ	Þ	THORN			
	1	1	z	z	lowercase letter	ß	ß	sharps			
	2	2	{	{	left curly brace	à	à	a, grave accent			
	3	3	|		vertical bar	á	á	a, acute accent			
	4	4	}	}	right curly brace	â	â	a, circumflex			
	5	5	~	~	tilde	ã	ã	a, tilde			
	6	6		-	en dash	ä	ä	a, umlaut			
	7	7	€	--	em dash	å	å	a, ring			
	8	8			nonbreaking space	æ	æ	ae			
	9	9	‚	!	inverted exclamation	ç	ç	c, cedilla			
	:	colon	ƒ	¢	cent sign	è	è	e, grave accent			
<	;	semicolon	„	£	pound sterling	é	é	e, acute accent			
	<	less-than sign	…	¢	general currency sign	ê	ê	e, circumflex			
	=	equals sign	†	¥	yen sign	ë	ë	e, umlaut			
>	>	greater-than sign	‡		broken vertical bar	ì	ì	i, grave accent			
	?	question mark	ˆ	§	section sign	í	í	i, acute accent			
	@	@ sign	‰	§	section sign	î	î	i, circumflex			
	A	A uppercase letter	Š	§	section sign	ï	ï	i, umlaut			
	B	B uppercase letter	‹	§	section sign	&ieth;	ð	eth, Icelandic			
	C	C uppercase letter	Œ	§	section sign	ñ	ñ	n, tilde			
	D	D uppercase letter		§	section sign	ò	ò	o, grave accent			
	E	E uppercase letter	Ž	§	section sign	ó	ó	o, acute accent			
	F	F uppercase letter		§	section sign	ô	ô	o, circumflex			
	G	G uppercase letter		§	section sign	õ	õ	o, tilde			
	H	H uppercase letter	‘	§	section sign	ö	ö	o, umlaut			
	I	I uppercase letter	’	§	section sign	÷	÷	division sign			
	J	J uppercase letter	“	§	section sign	ø	ø	o, slash			
	K	K uppercase letter	”	§	section sign	ù	ù	u, grave accent			
	L	L uppercase letter	•	§	section sign	ú	ú	u, acute accent			
	M	M uppercase letter	–	§	section sign	û	û	u, circumflex			
	N	N uppercase letter	—	§	section sign	ü	ü	u, umlaut			
	O	O uppercase letter	˜	§	section sign	ý	ý	y, acute accent			
	P	P uppercase letter	™	§	section sign	þ	þ	thorn, Icelandic			
	Q	Q uppercase letter	š	§	section sign	ÿ	ÿ	y, umlaut			

Source: www.LookupTables.com

Information Storage

■ Code

- Different machine types use different and incompatible instructions and encodings
- Even identical processors running different OSs have differences in their coding conventions and hence are not binary compatible
- Example) C code and its corresponding machine codes

```
int sum(int x, int y) {  
    return x + y;  
}
```

Linux 32: 55 89 e5 8b 45 0c 03 45 08 c9 c3

Windows: 55 89 e5 8b 45 0c 03 45 08 5d c3

Sun: 81 c3 e0 08 90 02 00 09

Linux 64: 55 48 89 e5 89 7d fc 89 75 f8 03 45 fc c9 c3

- [Refer to text, Chap. 3]

Information Storage

■ Boolean algebra

- Defined over 2-element set $\{0, 1\}$
- Basic operations
 - \sim (NOT)
 - $\&$ (AND)
 - $|$ (OR)
 - \wedge (XOR)

Information Storage

■ Bit-level operations in C

- \sim , $\&$, $|$, \wedge
- Example)

C expression	Binary expression	Binary result	Hexadecimal result
$\sim 0x41$	$\sim [0100\ 0001]$	$[1011\ 1110]$	0xBE
$\sim 0x00$	$\sim [0000\ 0000]$	$[1111\ 1111]$	0xFF
$0x69 \& 0x55$	$[0110\ 1001] \& [0101\ 0101]$	$[0100\ 0001]$	0x41
$0x69 0x55$	$[0110\ 1001] [0101\ 0101]$	$[0111\ 1101]$	0x7D

- Can be used for masking
 - Eg) $x \& 0xFF$ (low-order byte of a word x)
when $x = 0x1234ABCD$, $x \& 0xFF == 0x000000CD$

Information Storage

■ Logical operations in C

- `!`, `&&`, `||`
- Treat any non-zero argument as representing TRUE and argument 0 as representing FALSE
- Return 1 or 0 (TRUE for 1, FALSE for 0)
- Example)

Expression	Result
<code>!0x41</code>	<code>0</code>
<code>!0x00</code>	<code>1</code>
<code>!!0x41</code>	<code>1</code>
<code>0x69 && 0x55</code>	<code>1</code>
<code>0x69 0x55</code>	<code>1</code>

* Early Evaluation

- Logical operators do not evaluate their second argument if the result of the expression can be determined by evaluating the first argument
Eg) `a && 5/a`, `p && *p++`

Information Storage

■ Shift operations in C

- `<<, >>`
- `x << k` (left shift)
 - `x` is shifted `k` bits to the left
 - $[x_{n-1}, x_{n-2}, \dots, x_0] \rightarrow [x_{n-k-1}, x_{n-k-2}, \dots, x_0, 0, \dots, 0]$
- `x >> k` (logical right shift)
 - $[x_{n-1}, x_{n-2}, \dots, x_0] \rightarrow [0, \dots, 0, x_{n-1}, x_{n-2}, \dots, x_k]$
- `x >> k` (arithmetic right shift, default for signed data in typical C environments)
 - $[x_{n-1}, x_{n-2}, \dots, x_0] \rightarrow [x_{n-1}, \dots, x_{n-1}, x_{n-1}, x_{n-2}, \dots, x_k]$
- Example)

Operation	Values	
Argument <code>x</code>	[01100011]	[10010101]
<code>x << 4</code>	[00110000]	[01010000]
<code>x >> 4</code> (logical)	[00000110]	[00001001]
<code>x >> 4</code> (arithmetic)	[00000110]	[11111001]

Integer Representations

- Encoding unsigned integers
- Encoding signed integers
 - Signed-magnitude encoding
 - 1's-complement encoding
 - 2's-complement encoding

Integer Representations

■ Encoding unsigned integers

$$X = [x_{w-1}, x_{w-2}, \dots, x_0] \quad x = 0000\ 0111\ 1101\ 0011_2$$

$$B2U(X) = \sum_{i=0}^{w-1} x_i \cdot 2^i$$

$$\begin{aligned} B2U(X) &= 2^{10} + 2^9 + 2^8 + 2^7 \\ &\quad + 2^6 + 2^4 + 2^1 + 2^0 \\ &= 1024 + 512 + 256 + 128 \\ &\quad + 64 + 16 + 2 + 1 \\ &= 2003 \end{aligned}$$

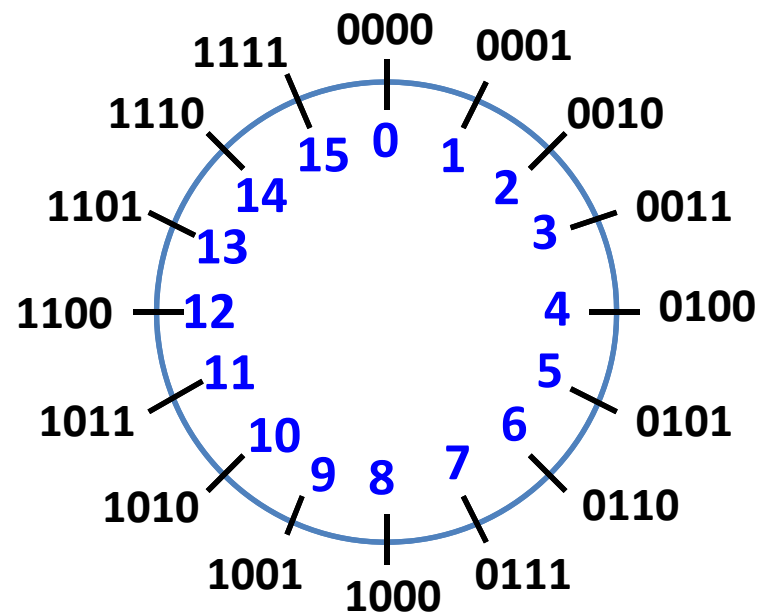
- What is the range for unsigned values with w bits?
 - $0 \sim 2^w - 1$

Integer Representations

■ Encoding unsigned integers

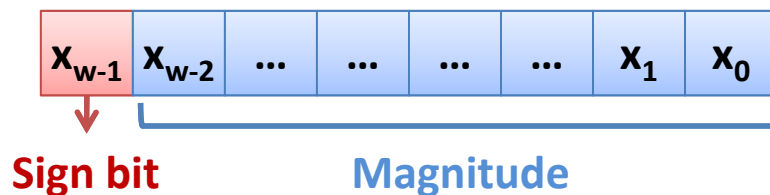
$$X = [x_{w-1}, x_{w-2}, \dots, x_0]$$

$$B2U(X) = \sum_{i=0}^{w-1} x_i \cdot 2^i$$

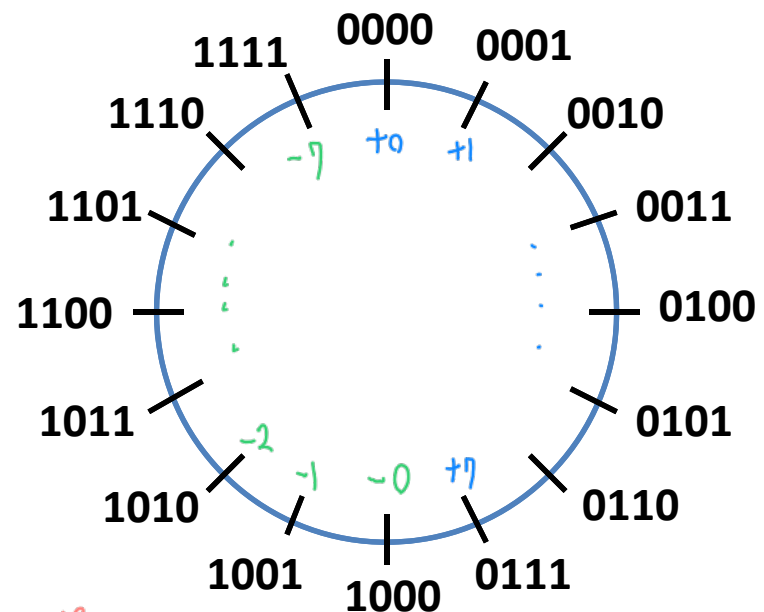


Integer Representations

■ Encoding signed integers: sign-magnitude



$$B2S(X) = (-1)^{x_{w-1}} \cdot \left(\sum_{i=0}^{w-2} x_i \cdot 2^i \right)$$



- Two zeros
 - [000...00], [100...00]
- Used for floating-point numbers

+ 사칙연산에 어려움 \Rightarrow 정수 표현에 잘 이용X

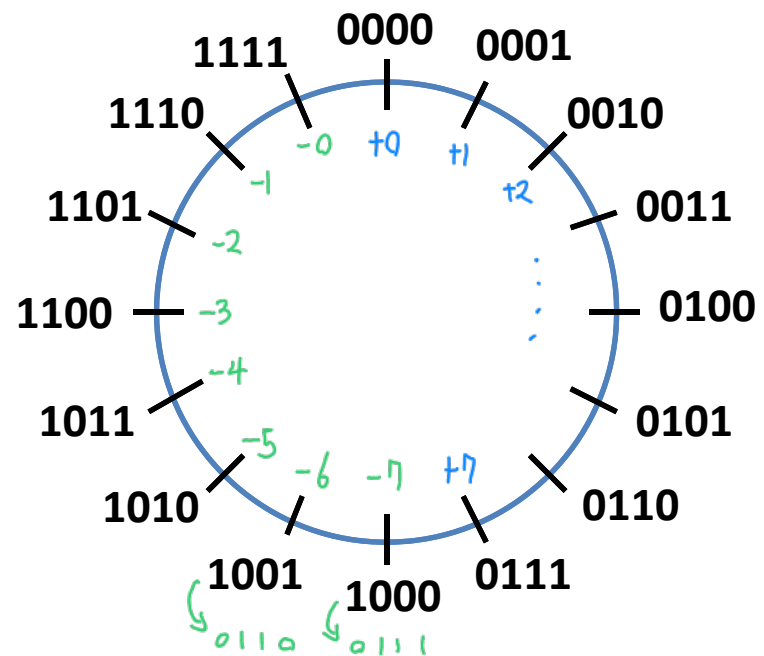
Integer Representations

■ Encoding signed integers: 1's-complement



$$B2O(X) = -x_{w-1}(2^{w-1} - 1) + \left(\sum_{i=0}^{w-2} x_i \cdot 2^i \right)$$

- Easy to find $-n$
- Two zeros
 - $[000...00]$, $[111...11]$
- No longer used



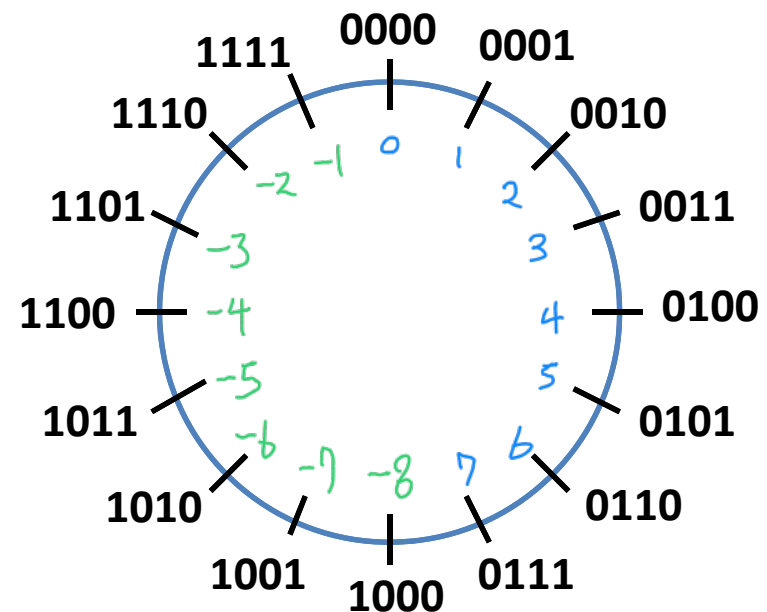
Integer Representations

■ Encoding signed integers: 2's-complement



$$B2T(B) = -x_{w-1} \cdot 2^{w-1} + \left(\sum_{i=0}^{w-2} x_i \cdot 2^i \right)$$

- Unique zero
- Easy for hardware
 - leading 0 ≥ 0 , leading 1 < 0
- Used by almost all modern machines



Integer Representations

■ Notes) 2's-complement representation

- Following holds for 2's-complement

$$\sim x + 1 == -x$$

▪ Observation

- $\sim x + x == 1111...11_2 == -1$
- $\sim x + x + (-x + 1) == -1 + (-x + 1)$
- $\sim x + 1 == -x$

Integer Representations

■ Value ranges of w-bit integers

■ Unsigned values

- UMin = 0 [000...00]
- UMax = $2^w - 1$ [111...11]

■ 2's-complement values

- TMin = -2^{w-1} [100...00]
- TMax = $2^{w-1} - 1$ [011...11]

(Values for w = 16)

	Decimal	Hex	Binary
UMax	65535	FF FF	11111111 11111111
TMax	32767	7F FF	01111111 11111111
TMin	-32768	80 00	10000000 00000000
-1	-1	FF FF	11111111 11111111
0	0	00 00	00000000 00000000

Integer Representations

■ Value ranges for different word sizes

	w = 8	w = 16	w = 32	w = 64
Umax	255	65,535	4,294,967,295	18,446,744,073,709,551,615
Tmax	127	32,767	2,147,483,647	9,223,372,036,854,775,807
Tmin	-128	-32,768	-2,147,483,648	-9,223,372,036,854,775,808

■ Observations

- $|TMin| = TMax + 1$
(Asymmetric range)
- $UMax = 2 * TMax + 1$

■ In C programming

- `#include <limits.h>`
- `INT_MIN`, `INT_MAX`,
`LONG_MIN`, `LONG_MAX`,
`UINT_MAX`, ...
- Platform-specific values

Integer Representations

■ Value ranges of w-bit integers

- Unsigned integer
 - $0 \sim 2^w - 1$
- Signed-magnitude
 - $-(2^{w-1} - 1) \sim (2^{w-1} - 1)$
- 1's-complement
 - $-(2^{w-1} - 1) \sim (2^{w-1} - 1)$
- 2's-complement
 - $-2^{w-1} \sim (2^{w-1} - 1)$

Integer Representations

■ Type conversion (signed ↔ unsigned)

- Guess the output

```
short    int    v = -12345;  
unsigned short uv = (unsigned short) v;  
printf("v = %d, uv = %u\n", v, uv);
```

⇒ $v = -12345$, $uv = 53191$

```
unsigned u = 4294967295u;    /* UMax_32 */  
int      tu = (int) u;  
printf("u = %u, tu = %d\n", u, tu);
```

⇒ $u = 4294967295$, $tu = -1$

Integer Representations

■ Type conversion (signed ↔ unsigned)

■ Principle

- The effect of casting is to keep the bit values identical but change how these bits are interpreted
 - ✓ Underlying bit representation stays the same

→ 기존 표현을 동일하게 가져가지만
그 해석이 달라진다.

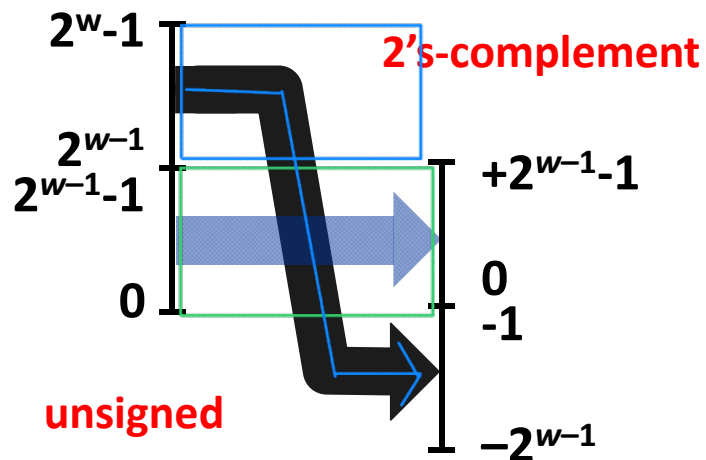
Integer Representations

■ Type conversion (unsigned → signed)

- The same bit pattern is interpreted as a signed number
- Mathematical analyses

$$U2T_w(x) = \begin{cases} x, & x < 2^{w-1} \\ x - 2^w, & x \geq 2^{w-1} \end{cases}$$

unsigned short `x = 2003;`
 short `ix = (short) x;`
 unsigned short `y = 0xbabe;`
 short `iy = (short) y;`



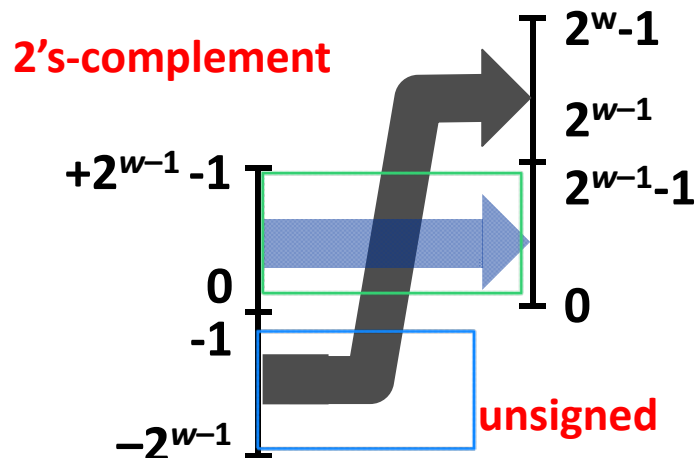
	Decimal	Hex	Binary
x	2003	07 D3	00000111 11010011
ix	2003	07 D3	00000111 11010011
y	47806	BA BE	10111010 10111110
iy	-17730	BA BE	10111010 10111110

Integer Representations

■ Type conversion (signed → unsigned)

- Ordering inversion (negative → big positive)
- Mathematical analyses

$$T2U_w(x) = \begin{cases} x + 2^w, & x < 0 \\ x, & x \geq 0 \end{cases}$$



```
short      ix = 2003;
Unsigned short x = (unsigned short) ix;
short      iy = -2003;
unsigned short y = (unsigned short) iy;
```

	Decimal	Hex	Binary
i x	2003	07 D3	00000111 11010011
x ix	2003	07 D3	00000111 11010011
i y	-2003	F8 2D	11111000 00101101
y iy	63533	F8 2D	11111000 00101101

Integer Representations

■ Type conversion (signed ↔ unsigned) in C

- Same as $U2T_w/T2U_w$
- C constants
 - By default, most numbers are considered to be **signed**
 - Use "U" or "u" as a suffix for unsigned constants
 - ✓ 0U, 12345U, 0x1A2Bu
- When conversion occurs?
 - Explicit casting →
 - Implicit casting →
 - ✓ In assignments and procedure calls
 - Printing numeric values with printf()

```
int      tx, ty;
unsigned ux, uy;
tx = (int) ux;
uy = (unsigned) ty;
```

```
int      f(unsigned);
tx = ux;
f(ty);
```


Integer Representations

■ Type conversion (signed ↔ unsigned) in C

■ C rules

- When an operation is performed where one operand is signed and the other is unsigned
 - ✓ C implicitly casts the signed argument to unsigned and performs the operations assuming the numbers are nonnegative
 - Little difference for standard arithmetic operations
 - But non-intuitive results for relational operators (<, >, ==, <=, >=, !=)

Integer Representations

- **Type conversion (signed ↔ unsigned) in C**
 - Example) subtleties in computation with relational operators

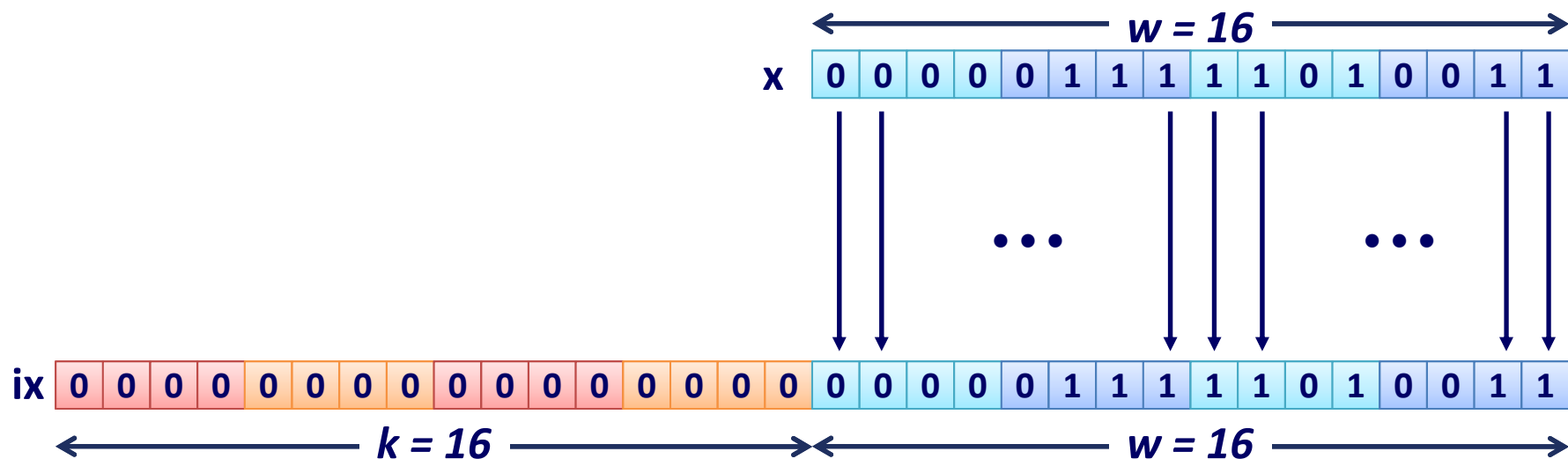
Expression	Type	Evaluation
<code>0 == 0U</code>	unsigned	1
<code>-1 < 0</code>	signed	1
<code>UINT_MAX ← -1 < 0U</code>	unsigned	0 (False)
<code>-1 > -2</code>	signed	1
<code>(unsigned) -1 > -2</code>	unsigned	1
<code>2147483647 > -2147483647-1</code>	signed	1
<code>2147483647U > -2147483647-1</code>	unsigned	0
<code>2147483647 > (int) 2147483648U</code>	signed	1

Integer Representations



- **Expanding the bit representation**
 - Zero extension
 - Sign extension

- ```
unsigned short x = 2003;
unsigned ix = (unsigned) x;
```

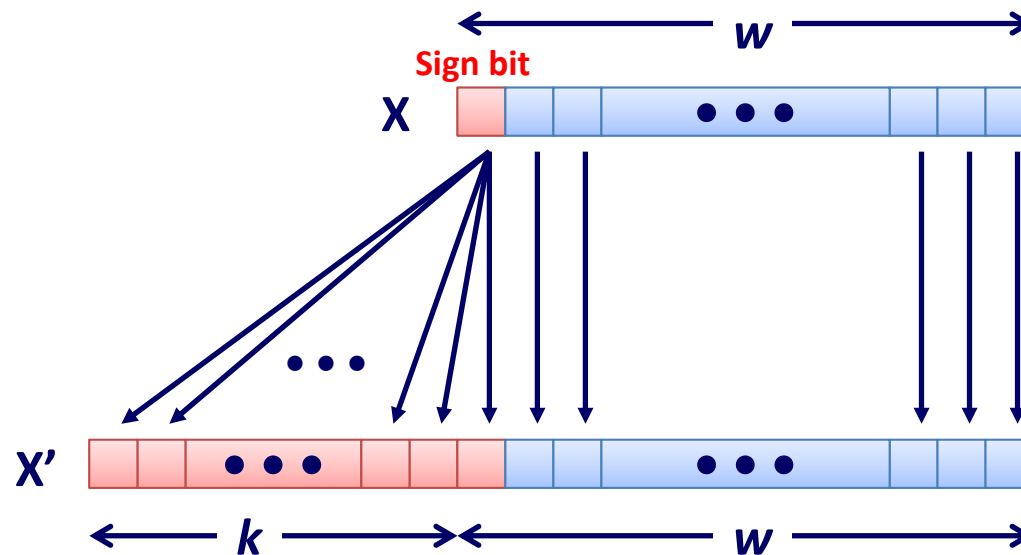


# Integer Representations

## ■ Expansion for signed: $w$ bits $\rightarrow w+k$ bits

### ▪ Sign extension

- Add copies of most significant bit (sign bit)
- Given  $w$ -bit signed integer  $x$ ,  
convert it to  $(w+k)$ -bit integer with the same value



# Integer Representations

## ■ Expansion for signed: Example

- Converting from smaller to larger integer type
- C automatically performs sign extension

| w | Decimal | Binary    | Description                            |
|---|---------|-----------|----------------------------------------|
| 4 | 5       | 0101      | $4 + 1 = 5$                            |
| 5 | 5       | 0 0101    | $4 + 1 = 5$                            |
| 8 | 5       | 0000 0101 | $4 + 1 = 5$                            |
| 4 | -5      | 1011      | $-8 + 2 + 1 = -5$                      |
| 5 | -5      | 1 1011    | $-16 + 8 + 2 + 1 = -5$                 |
| 8 | -5      | 1111 1011 | $-128 + 64 + 32 + 16 + 8 + 2 + 1 = -5$ |

# Integer Representations

## ■ Expansion for signed: Example

- Converting from smaller to larger integer type
- C automatically performs sign extension

```
short int x = 2003;
int ix = (int) x;
short int y = -2003;
int iy = (int) y;
```

|    | Decimal | Hex         | Binary                              |
|----|---------|-------------|-------------------------------------|
| x  | 2003    | 07 D3       | 00000111 11010011                   |
| ix | 2003    | 00 00 07 D3 | 00000000 00000000 00000111 11010011 |
| y  | -2003   | F8 2D       | 11111000 00101101                   |
| iy | -2003   | FF FF F8 2D | 11111111 11111111 11111000 00101101 |

# Integer Representations

## ■ Truncating for unsigned & signed: $w$ bits $\rightarrow k$ bits

- Just drop the high order ( $w-k$ )-bits
- Equivalent to computing  $x \bmod 2^k$  for unsigned and  $\text{U2T}_k(x \bmod 2^k)$  for signed

```
unsigned int x = 0xcafebabe;
unsigned short sx = (unsigned short) x;
int y = 0x2003beef;
short sy = (short) y;
```

|    | Decimal    | Hex         | Binary                              |
|----|------------|-------------|-------------------------------------|
| x  | 3405691582 | CA FE BA BE | 11001010 11111110 10111010 10111110 |
| sx | 47806      | BA BE       | 10111010 10111110                   |
| y  | 537116399  | 20 03 BE EF | 00100000 00000011 10111110 11101111 |
| sy | -16657     | BE EF       | 10111110 11101111                   |



# Integer Representations

## ■ Example-1) Type conversion in C

```
int main ()
{
 unsigned i;
 for (i = 10; i > 0; i--)
 printf("%u\n", i);
}
```

10  
9  
8  
7  
6  
5  
4  
3  
2  
1

```
int main ()
{
 unsigned i;
 for (i = 10; i >= 0; i--)
 printf("%u\n", i);
}
```

*i=0 일때  
마지막으로 i--를 수행하면  
i = UINT\_MAX*

10  
9  
8  
7  
6  
5  
4  
3  
2  
1  
0

# Integer Representations

## ■ Example-2) Type conversion in C

```
int sum_array(int a[], unsigned len)
{
 int i;
 int result = 0;

 for (i = 0; i <= len - 1; i++)
 result += a[i];

 return result;
}
```

when  $len == 0$   
문제 존재

OU

$\Rightarrow OU + (-1) = U\_INT\_MAX$

$-1 \rightarrow \dots \rightarrow U\_INT\_MAX$

Computes the sum of *len* elements of array *a*

# Integer Representations

## ■ Example-3) Type conversion in C

```
int strlonger(char *s, char *t)
{
 return strlen(s) - strlen(t) > 0;
}
```

0U-1U

Compares the length of two strings

```
int strlonger(char *s, char *t)
{
 return strlen(s) > strlen(t);
}
```

Compares the length of two strings

# Integer Representations

## ■ Example-4) Type conversion in C

```
#include <stdio.h>
int main ()
{
 unsigned char c;

 while ((c = getchar()) != EOF)
 putchar(c);
}
```

Copy standard input (or a file) onto standard output (or another file)

# Integer Representations

## ■ Advices on signed vs. unsigned

- There are many tricky situations when you use unsigned integers (hard to debug)
- Do not use unsigned just because numbers are nonnegative
- Use unsigned only when you need collections of bits with no numeric interpretation ("flags")
- Few languages other than C support unsigned integers

# Summary

