

Chapter 1 - elementary NT

4 example sheets per course

Chapter 2 - the reals

Chapter 3 - sets and functions

(here  $\mathbb{N}$  is  
positive integers)

Chapter 4 - countability

Chapter 0 - proof

A proof is a logical argument that establishes a conclusion.

- Prove things 1) to be sure they're true  
2) to understand why they're true

### Examples of proofs and non-proofs

1. Claim  $n^3 - n$  is always a multiple of 3 for  $n \in \mathbb{N}$ .

Proof  $n^3 - n = n(n-1)(n+1)$

(product of 3 consecutive numbers)

One of  $n-1, n, n+1$  is a multiple of 3, as they are 3 consecutive integers.

So  $(n-1)n(n+1)$  is a multiple of 3.  $\square$

2. Claim For any positive integer  $n$ , if  $n^2$  is even then  $n$  is even.

(Non) Proof Given  $n \in \mathbb{N}$  that is even, we have  $n = 2k$ ,  $k \in \mathbb{N}$ .

So  $n^2 = (2k)^2 = 4k^2 = 2(2k^2)$  so  $n^2$  is even.  $\square$

This is a wrong proof: wrong way round

"if A then B" is not the same as "if B then A"

3. Claim For any  $n \in \mathbb{N}$ , if  $n^2$  is a multiple of 9 then so is  $n$ .

(Non) Proof For any  $n \in \mathbb{N}$  that's a multiple of 9,  $n = 9k$  for some  $k \in \mathbb{N}$ . So  $n^2 = 81k^2 = 9(9k^2)$

so  $n^2$  is a multiple of 9.  $\square$

But this time the claim isn't even true.

To show "if A then B" is false, it's sufficient to find one counterexample where A is true and B is false.

Back to "if  $n^2$  is even then n is even"

Proof Suppose  $n^2$  is even. Then for some  $k \in \mathbb{N}$ ,  $n^2 = 2k$ .

This statement may be harder to work with e.g.  $n = \sqrt{2}k$  isn't as easy to work with.

So suppose that n is odd. So  $n = 2k+1$  for some  $k \in \mathbb{N}$ .

$$n^2 = (2k+1)^2 = 4k^2 + 4k + 1 = 2(2k^2 + 2k) + 1$$

which is odd. # "contradiction"

Thus n is even. □

Note the claim was true because of properties of odd numbers.

To show "if A then B" we showed that there is no case where A is true and B is false.

Another viewpoint : To show  $A \Rightarrow B$  (A implies B)

This is the same as  $\text{NOT } B \Rightarrow \text{NOT } A$

Claim The solution of the real equation  $x^2 - 5x + 6 = 0$  is  $x = 2$  or  $x = 3$ .

This is really 2 assertions:

- 1)  $x = 2$  and  $x = 3$  are solutions of  $x^2 - 5x + 6 = 0$
- 2) there are no other solutions.

Equivalently

- 1)  $x = 2$  or  $x = 3 \Rightarrow x^2 - 5x + 6 = 0$
- 2)  $x^2 - 5x + 6 = 0 \Rightarrow x = 2$  or  $x = 3$

so claim is  $x = 2$  or  $x = 3 \iff x^2 - 5x + 6 = 0$ .

Proof If  $x = 2$  or  $x = 3$

then  $x - 2 = 0$  or  $x - 3 = 0$

so  $(x-2)(x-3) = 0$  so  $x^2 - 5x + 6 = 0$ .  $\checkmark$

If  $x^2 - 5x + 6 = 0$  then  $(x-2)(x-3) = 0$

so  $x - 2$  is 0 or  $x - 3$  is 0

so  $x = 2$  or  $x = 3$   $\checkmark$

□

Alternative Proof For any real  $x$ ,

$$x^2 - 5x + 6 = 0 \iff (x-2)(x-3) = 0$$

$$\iff x-2 = 0 \text{ or } x-3 = 0$$

$$\iff x = 2 \text{ or } x = 3$$

□

Claim Every positive real number is at least 1.

(Non)-Proof Let  $x$  be the smallest positive real - then we try to show it's 1. Suppose for contradiction that  $x \neq 1$ . Then  $x < 1$  or  $x > 1$ .

If  $x < 1$ :  $0 < x^2 < x$  which is impossible as it contradicts the definition that  $x$  is the smallest positive real. #

If  $x > 1$  then  $\sqrt{x} < x$  contradicting definition. #

So  $x = 1$ .

□

The "proof" assumes that there is a smallest positive real.

This is nonsense, so so are the statements that follow from this assumption.

Moral: Every line in a proof must be justified.

## Chapter 1 : Elementary Number Theory

Intuitively, the natural numbers (written  $\mathbb{N}$ ) consist of

$1, 1+1, 1+1+1, 1+1+1+1, \dots$



How to make precise?

What we assume The natural numbers ( $\mathbb{N}$ ) is a set ~~of~~ containing an element "1" and an operation "+1" satisfying

- 1) For all  $n$ ,  $n+1 \neq 1$
  - 2) If  $m \neq n$  then  $m+1 \neq n+1$
  - 3) For any property  $p(n)$ : if  $p(1)$  is true and for every  $n$ , if  $p(n)$  is true then  $p(n+1)$  is true, then  $p(n)$  is true for all  $n$ .
- } capture "all on list are different"  
"induction axiom"

These 3 are the Peano axioms.

To see why 3) captures our intuitive notion of "that list is everything" we can use induction - take  $p(n) = "n$  is on the list".

Can write 2 for  $1+1$  and have operation "+2" defined by  $n+2 = (n+1) + 1$

In fact we can define " $+k$ " for any  $k \in \mathbb{N}$  by:

$$n + (k+1) = (n+k) + 1 \quad \text{for each } n \in \mathbb{N}$$

By induction this defines  $n+k$  (with operation " $+k$ ")

Take  $p(k) = "+k \text{ is defined}"$ .

$+1$  is defined.

If  $+k$  is defined, then using  $(n+k)+1 = n+(k+1)$  we have  $k+1$  is defined.  $\square$

Similarly we can define multiplication, powers, etc.

Can check the usual rules of arithmetic:

$$1) \forall a, b \quad a+b = b+a \quad (+ \text{ is commutative})$$

$$2) \forall a, b \quad ab = ba \quad (\times \text{ is commutative})$$

$$3) \forall a, b, c \quad a + (b+c) = (a+b) + c \quad (+ \text{ is associative})$$

$$4) \forall a, b, c \quad a(bc) = (ab)c \quad (\times \text{ is associative})$$

$$5) \forall a, b, c \quad a(b+c) = (ab) + (ac) \quad (\times \text{ is distributive over } +)$$

Define  $a < b$  if  $a+c = b$  for some  $c \in \mathbb{N}$  (no 0).

$$6) \forall a, b : \text{ if } a < b \text{ then } a+c < b+c$$

$$7) \forall a, b : a < b \Rightarrow ac < bc$$

$$8) \forall a, b, c : a < b, b < c \Rightarrow a < c$$

$$9) \forall a : \text{ NOT } a < a$$

A more useful form of induction:

Induction says that if  $P(1)$  and  $\forall n, P(n) \Rightarrow P(n+1)$  then  $P(n) \forall n$ .

Strong induction : if  $P(1)$  and  $\forall n$  we have  $P(m) \forall m \leq n \Rightarrow P(n+1)$  then  $P(n) \forall n$ .

To deduce this from ordinary induction, apply ordinary induction to  $Q(n)$  where  $Q(n)$  is " $P(m) \forall m \leq n$ ".

Remarks: 1) Technically you don't need to check  $n=1$  case separately (as implied by the condition if interpreted suitably) — technically none are false below  $n=1$  (but always check anyway)

correct view of induction → 2) Normally to prove  $P(n) \forall n$  we take  $n$  and show  $P(n)$ . Strong induction allows you to assume it's true for smaller cases if it helps (e.g.  $P(m)$  for  $m < n$ ).  
 induction — really works "down", not "up"

2 equivalent forms of strong induction

1. If  $P(n)$  is false for some  $n$  then for some  $n$  we must have  $P(n)$  false but  $P(m)$  true  $\forall m \leq n$ .

"If there is a counterexample, then there is a minimal counterexample"

2. If  $P(n)$  is true for some  $n$  then there is a least  $n$  with  $P(n)$ .  
 "Well-ordering principle"

(there's something else unrelated called the well-ordering theorem)

## The Integers

The integers, written  $\mathbb{Z}$ , consist of all symbols  $n, -n$  ( $n$  a natural number) and 0.

Can define  $+$ ,  $\times$  etc on  $\mathbb{Z}$  from  $\mathbb{N}$ .

We can check all the previous algebraic rules, plus

$$\forall a \quad a+0 = a \quad (\text{identity for } +)$$

$$\forall a \exists b \text{ s.t. } a+b = 0 \quad (\text{inverses for } +)$$

Define  $a < b$  if  $\exists c \in \mathbb{N}$  with  $a+c = b$ .

All previous rules still hold except one change

$$\forall a, b, c: \text{ If } a < b, \underline{c > 0} \text{ then } ac < bc$$

## The Rationals

The rationals, written  $\mathbb{Q}$ , consist of all expressions

$\frac{a}{b}$  where  $a, b \in \mathbb{Z}$  with  $b \neq 0$ , with  $\frac{a}{b}$  regarded as the same as  $\frac{c}{d}$  if  $ad = bc$ .

Define  $\frac{a}{b} + \frac{c}{d} = \frac{ad+bc}{bd}$  (can check: does not matter how  $\frac{a}{b}, \frac{c}{d}$  written to be well defined).

Note cannot define operation on  $\mathbb{Q}$  by

$$\frac{a}{b} \mapsto \frac{a^2}{b^3} \quad \text{as e.g. } \frac{1}{2} \mapsto \frac{1}{8}$$

$$\frac{2}{4} \mapsto \frac{4}{64} = \frac{1}{16} \neq \frac{1}{8}$$

Similarly can check algebraic rules. Also

$$\forall a \in \mathbb{Q}, a \neq 0, \exists b \in \mathbb{Q} \ ab = 1 \quad (\text{inverses for } \cdot)$$

~~Definition of  $\mathbb{Z}$~~   $\mathbb{Z}$  of ~~Integers~~

Define  $\frac{a}{b} < \frac{c}{d}$  (where  $b, d > 0$ ) if  $ad < bc$ .

(can check  $\mathbb{Z}$  rules hold).

Note: Can view  $\mathbb{Z}$  as "living inside"  $\mathbb{Q}$  by identifying  $a \in \mathbb{Z}$  with  $\frac{a}{1}$  in  $\mathbb{Q}$

Structure of  $\mathbb{N}$  under  $+$  is easy: start at 1, keep doing  $+1$ .  
What about under  $\cdot$ ?

↑  
(meaning multiplication,  $\times$ )

Primes  
(not definition) For  $n \in \mathbb{N}$ , the multiples of  $n$  are all integers  $kn$  for some  $k \in \mathbb{Z}$ . e.g.  $2n, 3n, 16n, n, -5n, 0$

If  $m$  is a multiple of  $n$ , say  $n$  divides  $m$  or  
 $n$  is a divisor / factor of  $m$   
or  $n | m$ .

Definition (Prime) A natural number  $n \geq 2$  is prime if its only divisors are 1,  $n$ .

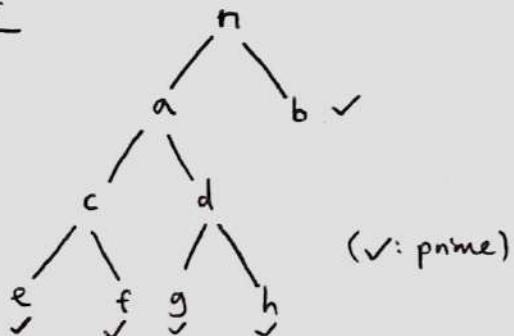
(We don't allow 1 to be prime, otherwise e.g.  $15 = 3 \times 5, 1 \times 3 \times 5, 1^2 \times 3 \times 5, \dots$ )

If  $n \geq 2$  is not prime, then it is composite.

Aim: "Break up" any number into a unique product of primes.  
e.g.  $63 = 3 \times 3 \times 7$

Proposition 1 Every natural number with  $n > 2$  is expressible as a product of primes.

Proof



This will eventually terminate  
This may be hard to write out as a proper proof so use induction.

Induction on  $n$ :  $n = 2 \checkmark$

Given  $n > 2$ :

If  $n$  is prime  $\checkmark$

If not ( $n$  is composite), then  $n = ab$ , for some  $1 < a, b < n$

By induction hypothesis we've assumed that

$$a = p_1 p_2 p_3 \dots p_k, \quad b = q_1 q_2 q_3 \dots q_l \quad \text{for } p_i, q_i \text{ prime}$$

Then  $n = p_1 p_2 p_3 \dots p_k q_1 q_2 q_3 \dots q_l$ , a product of primes.

✓  
□

Remark Can define an empty product (i.e. of no primes) to be 1.

If so, then Proposition 1 could start  $n \geq 1$ .

Theorem There are infinitely many primes.

Proof Suppose there are finitely many primes  $p_1, p_2, \dots, p_k$ .

All  $p_i$  divide  $p_1 p_2 \dots p_k$  so none divide  $p_1 p_2 \dots p_k + 1 = n$ .  
So  $n$  has no prime factor.  $\blacksquare$  □

There is no formula for the  $n^{\text{th}}$  prime.

Want: Prime factorisation of a number is unique (up to reordering)

We need  $p \mid ab \Rightarrow p \mid a$  or  $p \mid b$  for  $p$  prime

### Highest Common Factors

For  $a, b \in \mathbb{N}$ , we have  $c \in \mathbb{N}$  is the HCF of  $a$  and  $b$  if  
 $\underline{c \mid a, c \mid b}$  and if  $\underline{d \mid a, d \mid b}$  then  $d \mid c$ .

(" $c$  is a common factor of  $a$  and  $b$ , and every common factor divides  $c$ ")

e.g. 18 and 12 :      18's factors 1, 2, 3, 6, 9, 18  
                                 12's factors 1, 2, 3, 4, 6, 12

Common factors are 1, 2, 3, 6

So the HCF (if it exists) is the greatest of all common factors  
but if  $a$  and  $b$  had common factors 1, 2, 3, 4, 6 then  $a$  and  $b$  would not have a HCF.

Aim Show that  $\forall a, b$ , HCF  $(a, b)$  exists.

Proposition 3 (Division Algorithm) Let  $n, k \in \mathbb{N}$ . Then we can write  $n = qk + r$  where  $q, r \in \mathbb{Z}$  with  $0 \leq r \leq k-1$ .

Proof Induction on  $n$ :  $n=1 = 0k+1$  ✓

Given  $n > 1$ :

Have  $n-1 = qk+r$  for some  $q, r \in \mathbb{Z}$ ,  $0 \leq r \leq k-1$ .

If  $r < k-1$  then  $n = qk + (r+1)$  ✓

If  $r = k-1$  then  $n = (q+1)k$  ✓ ✓ □

Euclid's Algorithm For finding  $\text{hcf}(a, b)$  (say  $a > b$ )

This proves hcf exists and gives a quick way to find it.

### Algorithm

Write  $a = q_1 b + r_1$  ( $q_i, r_i \in \mathbb{Z}, 0 \leq r_i < b$ )

Write  $b = q_2 r_1 + r_2$

$$r_1 = q_3 r_2 + r_3$$

...

Continue until you get

$$r_{n-1} = q_{n+1} r_n + r_{n+1} \text{ with}$$

$$r_{n+1} = 0$$

Then hcf is  $r_n$ .

Note : terminates because numbers decrease

Terminates  $M \leq b$  steps since  $b > r_1 > r_2 > \dots$

so we must get a "hcf". It must exist.

We then show that this "hcf" satisfies the required properties:

### Example

$$a = 372, b = 162$$

$$372 = 2 \times 162 + 48$$

repeat:

$$162 = 3 \times 48 + 18$$

$$48 = 2 \times 18 + 12$$

$$18 = 1 \times 12 + 6$$

$$12 = 2 \times 6$$

output 6

Theorem 4 The output of the algorithm on  $a, b$  is  $\text{hcf}(a, b)$ .

Proof (i)  $r_n | r_{n-1}$  (as  $r_{n+1} = 0$ ), from the last line

so  $r_n | r_{n-2}$  (by 2nd last line)

so ~~so~~  $r_n | r_i \forall i$  inductively.

So  $r_n | b$  (2nd line)

so  $r_n | a$ . So  $r_n$  is a common factor.

(ii) Given  $d$  with  $d | a, d | b$ :

We have  $d | r_1$  (1st line)

so  $d | r_2$  (2nd line)

and inductively  $d | r_i \forall i$

so  $d | r_n$  as required.  $\square$

e.g. HCF of 87, 52 - use Euclid's algorithm.

$$(1) \quad 87 = 1 \times 52 + 35$$

$$87 = 3 \times 29$$

$$(2) \quad 52 = 1 \times 35 + 17$$

$$52 = 2^2 \times 13$$

$$(3) \quad 35 = 2 \times 17 + 1$$

$$(4) \quad 17 = 17 \times 1$$

$$\text{hcf}(87, 52) = 1$$

We have that  $\text{hcf}(87, 52) = 1$  or write as  $(87, 52) = 1$  (coprime)

Can we write 1 as a multiple of 87 + a multiple of 52?

$$87x + 52y = 1 \quad \text{for some } x, y \in \mathbb{Z}$$

We have  $1 = 1 \times 35 - 2 \times 17$  from (3)

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$$\Rightarrow 1 = 1 \times 35 - 2(52 - 35) \quad \text{from (2)}$$

$$\Rightarrow 1 = -2 \times 52 + 3 \times 35$$

$$\Rightarrow 1 = -2 \times 52 + 3(87 - 52) \quad \text{from (1)}$$

$$\Rightarrow 1 = 3 \times 87 - 5 \times 52$$

so we have a solution  $x = 3, y = -5$ .

Theorem 5  $\forall a, b \in \mathbb{N} \quad \exists x, y \in \mathbb{Z}$  such that  $ax + by = \text{hcf}(a, b)$   
 "can write hcf as a linear combination of a and b"

Proof 1 Run Euclid on  $a$  and  $b$ , say with output  $r_n$ .

Have  $r_n = x r_{n-1} + y r_{n-2}$  (from 2nd last line)  
 $r_{n-2} = q_n r_{n-1} + r_n$  of alg.

But  $r_{n-1}$  is expressible as  $x r_{n-2} + y r_{n-3}$ ,  $x, y \in \mathbb{Z}$   
 (3rd last line)

then substitute for  $r_n$

(note  $x, y$  not the same  
from line to line)

Continue: we obtain that  $\forall i: r_n = x r_i + y r_{i-1}, x, y \in \mathbb{Z}$   
 (inductively)

So  $r_n = ax + by$  for some  $x, y \in \mathbb{Z}$  (from line 1 of Euclid).  $\square$

Note: Euclid is showing  $x, y$  exist and how to get them.

Proof 2 Let  $h$  be the least positive linear combination of  $a$  and  $b$ .

Claim  $h = \text{hcf}(a, b)$

Proof (ii) Given  $d \mid a, d \mid b, d \mid xa + yb \quad \forall x, y \in \mathbb{Z}$   
so  $d \mid h$ .

(i) Suppose  $h \nmid a$ . Then  $a = qh + r, q, r \in \mathbb{Z}$   
and  $0 < r < h$ .

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Then  $r = a - qh = a - q(xa + yb)$ , also a linear combination. But  $r < h$  and  $h$  is the least positive LC.  $\times$   
so  $h \mid a$  and similarly  $h \mid b$ .  $\square$

This doesn't show how to find  $x, y$  though.

Application: Solving integer linear equations

Suppose  $a, b \in \mathbb{N}$ . When can we solve

$$\underline{ax = b}, \quad x \in \mathbb{Z}?$$

If  $x \in \mathbb{Q}$  allowed, then always  
otherwise, it's iff  $a \mid b$ .

Now consider  $ax + by = c \quad x, y \in \mathbb{Z}$

e.g.  $320x + 72y = 33$  not possible: LHS even  
RHS odd

$$87x + 52y = 33 ?$$

Possible:  $87x + 52y = 1$  has a solution so  
just multiply by 33.

Corollary 6 Let  $a, b, c \in \mathbb{N}$ . Then  $ax+by=c$  has an integer solution iff  $\text{hcf}(a, b) | c$ .

Proof Let  $h = \text{hcf}(a, b)$ .

$\Rightarrow$  Have  $ax+by=c$  for  $x, y \in \mathbb{Z}$   
But  $h|a, h|b$  so  $h|ax+by$ .

$\Leftarrow$  Have  $h = ax+by$  for some  $x, y \in \mathbb{Z}$

Multiply up by  $\frac{c}{h}$  (we know  $h|c$  so  $\frac{c}{h} \in \mathbb{Z}$ ) .

$$c = a\left(\frac{xc}{h}\right) + b\left(\frac{yc}{h}\right) \text{ giving a solution. } \square$$

Remark Corollary 6 is sometimes called Bézout's theorem.

Lemma 7 Let  $p$  be prime and  $a, b \in \mathbb{N}$ .

Then  $p|ab \Rightarrow p|a$  or  $p|b$ .

Proof Suppose  $p \nmid a$  then we want to show  $p|b$ .

Then  $\text{hcf}(p, a) = 1$ . ( $p$  prime)

So  $px+ay=1$ ,  $x, y \in \mathbb{Z}$ .

Then  $p|pbx+aby=b$ ,

where  $b$  is a multiple of  $p$  because

$p|pbx$  and  $p|ab$  so  $\cancel{pbx+aby} \Rightarrow p|pbx+aby$   
 $\Rightarrow p|b$ .  $\square$

- Remarks
- 1) Similarly  $p|a_1 a_2 \dots a_n \Rightarrow p|a_i$  for some  $i$ .  
(Inductively)
  - 2) We need  $p$  to be prime.

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### Theorem 8 (Fundamental Theorem of Arithmetic)

Every number  $n \in \mathbb{N}$  is expressible as a product  
( $n > 1$ )  
of primes, uniquely up to reordering.

#### Proof

Existence: Proposition 1 proved this ✓ (lecture 3)

Uniqueness: Induction on  $n$ :  $n=2$  ✓

Given  $n > 2$ : Suppose  $n = p_1 \dots p_k = q_1 \dots q_l$   
where  $p_i, q_i$  are prime. Must show  $k=l$  and  
after reordering,  $p_i = q_i \forall i$ .

We have  $p_1 | n \Rightarrow p_1 | q_1 \dots q_l \Rightarrow p_1 | q_i$  for some  $i$ .

Reorder to assume  $p_1 | q_1$ .

(+) But  $q_1$  is prime, so  $p_1 = q_1$ .

(\*) Then  $\frac{n}{p_1} = p_2 \dots p_k = q_2 \dots q_l$ . Repeat argument.  
(inductively)

Thus  $k=l$  and  $p_2 = q_2, p_3 = q_3, \dots, p_k = q_k$

(induction). □

Remark This argument works: we've checked it's true for  $n=2$ .

Although line (\*) is written below, we can think of it as assuming the uniqueness property holds for all values less than  $n$  so true for  $\frac{n}{p_1}$ .

Then write  $n = p_1 \dots p_k = q_1 \dots q_l$  and have  $p_1 = q_1$ .

More straightforward way: Once we get to line (+) we divide by  $p_1$  to give (\*). Then we can repeat the same argument.  
So it works as a kind of "downwards" inductive proof.

with  
 $\frac{n}{p_1}$  instead  
of  $n$ .

Digression In FTA we took the "things that cannot be broken up" (primes) and broke up every number as a product of these uniquely.

Consider  $\mathbb{Z}[\sqrt{-3}] = a + b\sqrt{-3} = a + b\sqrt{3}i$  for  
Can add / multiply 2 such elements  $a, b \in \mathbb{Z}$ .  
and get something of the same form.

So we can talk about "divides" / "multiple of" in  $\mathbb{Z}[\sqrt{-3}]$

$$\text{But } 4 = 2 \times 2 = (1 + \sqrt{-3})(1 - \sqrt{-3})$$

All these terms  $\xrightarrow{\text{can't be broken up}}$ : "unique factorisation" fails in  $\mathbb{Z}[\sqrt{-3}]$

FLT can be "proved" if unique factorisation is assumed in certain situations.

### Applications of FTA

#### 1) Factors

What are the factors of  $2^3 \times 3^7 \times 11^n$ ?

All of the form  $2^a 3^b 11^c$  ( $0 \leq a \leq 3$ ,  $0 \leq b \leq 7$ ,  $0 \leq c \leq 1$ )  
 $\xleftarrow{\quad}$   $a, b, c \in \mathbb{Z}$

are factors. There are no others - if  $7 \mid n$  then you'd have a factorisation involving 7. The factorisation must be unique so this is impossible.

So the factors of  $n = p_1^{a_1} \cdots p_k^{a_k}$  are precisely all numbers  $p_1^{b_1} \cdots p_k^{b_k}$  for  $0 \leq b_i \leq a_i$ .

#### 2) HCFs

Common factors of  $2^3 \cdot 3^7 \cdot 5 \cdot 11^3$  and  $2^4 \cdot 3^2 \cdot 11 \cdot 13$

All  $2^a 3^b 11^c$ ,  $0 \leq a \leq 3$ ,  $0 \leq b \leq 2$ ,  $0 \leq c \leq 1$   
 are common factors.  $2^3 \cdot 3^2 \cdot 11$  is HCF.

In general, HCF of  $p_1^{a_1} \dots p_k^{a_k}$  and  $p_1^{b_1} \dots p_k^{b_k}$   
 $(a_i, b_i > 0)$  is

$$\frac{p_1^{\min(a_1, b_1)} \dots p_k^{\min(a_k, b_k)}}{}$$

### 3) LCMs

Common multiples of  $2^3 3^7 5 11^3$  and  $2^4 3^2 11 \cdot 13$

are all  $2^a \cdot 3^b \cdot 5^c \cdot 11^d \cdot 13^e$ . (anything)

$$a \geq 4, b \geq 7, c \geq 1, d \geq 3, e \geq 1$$

So  $2^4 \cdot 3^7 \cdot 5^4 \cdot 11^3 \cdot 13$  is the LCM.

In general LCM of  $p_1^{a_1} \dots p_k^{a_k}$  and  $p_1^{b_1} \dots p_k^{b_k}$   
 $(a_i, b_i > 0)$  is

$$\frac{p_1^{\max(a_1, b_1)} \dots p_k^{\max(a_k, b_k)}}{}$$

Consequence: take  $\text{hcf}(x, y) \text{lcm}(x, y) = xy$   
 (since  $\min(a, b) + \max(a, b) = a+b$ ).

### Modular Arithmetic

Let  $n \geq 2$  be a natural number. The integers mod n ( $\mathbb{Z}_n$ ) consist of the integers with two regarded as the same if they differ by a multiple of  $n$ .

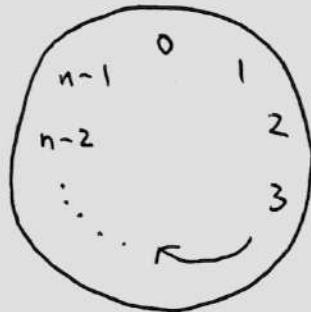
e.g. in  $\mathbb{Z}_7$  2 is the same as 16.

If  $x, y$  are the same in  $\mathbb{Z}_n$ , can write  $\frac{x \equiv y \pmod{n}}{x \equiv y \pmod{n}}$   
 or  $x \equiv y$  in  $\mathbb{Z}_n$

Thus  $x \equiv y \pmod{n} \iff x-y$  is a multiple of  $n$ .  
 $\iff x = y + kn, k \in \mathbb{Z}$ .

Note that no two of  $0, 1, \dots, n-1$  are congruent mod  $n$   
and every  $x$  is congruent to one of them mod  $n$   
(Division algorithm)

Can view  $\mathbb{Z}_n$  as



"correct picture" of  $\mathbb{Z}_n$

Do + and  $\times$  make sense in  $\mathbb{Z}_n$ ?

(Note: even/odd do not make sense in general, e.g.  $2 \equiv 9 \pmod{7}$ )

We want: if  $a \equiv a' \pmod{n}$ ,  $b \equiv b' \pmod{n} \Rightarrow a+b \equiv a'+b' \pmod{n}$   
and  $ab \equiv a'b' \pmod{n}$

$$\text{Have } a' = a+kn \quad (k \in \mathbb{Z})$$

$$b' = b+jn \quad (j \in \mathbb{Z})$$

$$a'+b' = a+b+(k+j)n \quad \text{so } a'+b' \equiv a+b \pmod{n} \quad \checkmark$$

$$a'b' = (a+kn)(b+jn) = ab + n(aj + bk + kjn)$$

$$\Rightarrow ab \equiv ab \pmod{n} \quad \checkmark$$

Laws of arithmetic are inherited from  $\mathbb{Z}$ .

$$\text{e.g. } a+b \equiv b+a \pmod{n}$$

Some things we have already done are expressible in terms of M.A.:

e.g. "p | ab  $\Rightarrow$  p | a or p | b" is  $ab \equiv 0 \pmod{p} \iff a \equiv 0 \pmod{p}$   
or  $b \equiv 0 \pmod{p}$ .

## Inverses

For  $a, b \in \mathbb{Z}_n$ , say  $b$  is an inverse of  $a$  if  $\underline{ab \equiv 1 \pmod{n}}$ .

e.g. in  $\mathbb{Z}_{10}$  inverse of 3 could be 7.

There is no inverse of 4: 1 mod 10 must make it odd.

Note 1. If inverse exists, then it's unique mod  $n$

Suppose in  $\mathbb{Z}_n$ ,  $ab = ac = 1$

$$\text{Then } b(ab) = b(ac) \Rightarrow 1(b) = 1(c).$$

2) If  $a$  inverted in  $\mathbb{Z}_n$ : can write  $a^{-1}$  for inverse.

## Important

1) Can "cancel" an invertible.

$$\begin{aligned} \text{If } a^{-1} \text{ exists then } ab &\equiv ac \pmod{n} \\ (\text{in } \mathbb{Z}_n) & \Rightarrow b \equiv c \pmod{n} \end{aligned}$$

2) In general you cannot cancel.

$$\text{e.g. in } \mathbb{Z}_{10} \quad 4 \times 5 \equiv 2 \times 5 \pmod{n}$$

$$\not\Rightarrow 4 \equiv 2 \pmod{n}$$

$\mathbb{Z}_p$  is very "well behaved" for  $p$  prime.

Proposition 9 Let  $p$  be prime. Then every  $a \not\equiv 0 \pmod{p}$  is invertible mod  $p$ .

[Equivalently in  $\mathbb{Z}_p$ ,  $\Leftrightarrow a \neq 0 \Rightarrow \exists b : ab = 1$ ]

Proof 1 Have  $(a, p) = 1$ .

So  $ax + py = 1$  for some  $x, y \in \mathbb{Z}$ .

That is  $ax = 1 - py \Rightarrow ax \equiv 1 \pmod{p}$  as required.  $\square$

Proof 2 In  $\mathbb{Z}_p$  consider all multiples of  $a$

$a \times 0, a \times 1, a \times 2, \dots, a(p-1)$

We must show one is 1.

All the multiples of  $a$  are distinct in  $\mathbb{Z}_p$  as

$ia = ja \Rightarrow (i-j)a = 0 \Rightarrow i-j \equiv 0 \pmod{p}$  or  
 $a \equiv 0 \pmod{p}$  but  $p$  is prime so  $i = j$ .

There are  $p$  distinct multiples mod  $p$ , so one must be  $1 \pmod{p}$ .  $\square$

What about general  $n$ ?

Proposition 9' Let  $n > 2$ . Then  $a$  is invertible mod  $n$  iff  $a$  is coprime to  $n$ :  $(a, n) = 1$ .

Proof  $(a, n) = 1 \Leftrightarrow ax + ny = 1$  for some  $x, y \in \mathbb{Z}$ .

$$\Leftrightarrow ax = 1 - ny$$

$$\Leftrightarrow ax \equiv 1 \pmod{n} \quad \square$$

The Euler  $\phi$ -function is defined for each  $n \in \mathbb{N}$  by

$\phi(n) = \text{number of } x \ (1 \leq x \leq n) \text{ with } (x, n) = 1$ .

(same as number of invertibles mod  $n$ ).

e.g. p prime:  $\phi(p) = p - 1$   
 $\phi(p^2) = p^2 - p$  (-p from  $p, 2p, 3p, \dots, pp$ )

$p, q$  distinct primes:  $\phi(pq) = pq - p - q + 1$

↑                      ↑                      ↙  
 multiples    multiples    from  $pq$   
 of  $p$             of  $q$

How do powers behave in  $\mathbb{Z}_p$ ?

e.g. powers of 2 in  $\mathbb{Z}_7$ :  $2^1 = 2, 2^2 = 4, 2^3 = 1$   
 (so it keeps repeating  $2, 4, 1, 2, 4, 1, \dots$ )  
 Note  $2^6 \equiv 1 \pmod{7}$ .

Powers of 2 in  $\mathbb{Z}_{11}$ :

$$2^1 = 2, 2^2 = 4, 2^3 = 8, 2^4 = 5, 2^5 = 10, 2^6 = 9,  
 2^7 = 7, 2^8 = 3, 2^9 = 6, 2^{10} = 1$$

Then repeats as we've got back to 1.

Theorem 10 (Fermat's Little Theorem) Let  $p$  be prime.

Then in  $\mathbb{Z}_p \pmod{p}$ ,  $\underline{a^{p-1} \equiv 1 \pmod{p}} \quad \forall a \neq 0$ .

Proof In  $\mathbb{Z}_p$  consider  $a \times 1, a \times 2, \dots, a \times (p-1)$

They are distinct in  $\mathbb{Z}_p$  ( $a_i = a_j \Rightarrow i = j$   
 as  $a$  is invertible).

And they are nonzero. ~~as~~  $a_i = 0 \Rightarrow a = 0$  or  $i = 0$  \*

They hence are  $1, 2, \dots, p-1$  in some order.

Multiply them all:  $a^{p-1} (p-1)! = (p-1)! \text{ in } \mathbb{Z}_p$

$(p-1)!$  is invertible as its factors all are, so

$a^{p-1} \equiv 1 \pmod{p}$

□

What about for general  $n$ ?

Theorem 10' (Fermat-Euler theorem)

Let  $n > 2$ . Then in  $\mathbb{Z}_n$ , every invertible  $a$  has  $a^{\phi(n)} = 1$ .

Proof Let the units in  $\mathbb{Z}_n$  be  $x_1, x_2, \dots, x_{\phi(n)}$  -  
↑ word for invertibles

Consider  $ax_1, ax_2, \dots, ax_{\phi(n)}$  : all distinct  
( $ax_i = ax_j \Rightarrow x_i = x_j$  as  $a$  invertible)  
and invertible.

So they are all the invertibles  $x_1, x_2, \dots, x_{\phi(n)}$  in some order.

Multiply them:

$$a^{\phi(n)} x_1 x_2 \dots x_{\phi(n)} = \underbrace{x_1 x_2 \dots x_{\phi(n)}}_{\text{invertible}}$$
$$\Rightarrow \underline{a^{\phi(n)} = 1} \text{ in } \mathbb{Z}_n. \quad \square$$

Question We know  $(p-1)! \not\equiv 0(p)$ . What is it?

$$\text{e.g. } p=5: 4! = 24 \equiv -1(5)$$

$$p=7: 6! = 720 \equiv -1(7)$$

Lemma 11 Let  $p$  be prime. Then in  $\mathbb{Z}_p$ ,  $x^2 = 1$   
 $\Rightarrow x = 1$  or  $x = -1$ .

Proof Work in  $\mathbb{Z}_p$ .  $x^2 = 1 \Rightarrow x^2 - 1 = 0$   
 $\Rightarrow (x-1)(x+1) = 0$

$$(x-1)(x+1) \equiv 0 \pmod{p} \Rightarrow x-1 \equiv 0 \text{ or } x+1 \equiv 0$$
$$\Rightarrow x = \pm 1. \quad \square$$

Remark Turns out that a nonzero polynomial of degree  $k$  in  $\mathbb{Z}_p$  always has at most  $k$  roots in  $\mathbb{Z}_p$ .

Theorem 12 (Wilson's Theorem) Let  $p$  be prime. Then  $(p-1)! \equiv -1 \pmod{p}$ .

Proof May assume  $p > 2$  (as obviously true for  $p=2$ ).

In  $\mathbb{Z}_p$  consider  $1, 2, 3, \dots, p-1$ .

Can pair up each  $a$  with its inverse  $a^{-1}$  (for  $a \neq a^{-1}$ )

But  $a = a^{-1} \Rightarrow a^2 = 1 \Leftrightarrow a = 1 \text{ or } -1$ .

Thus  $1, 2, \dots, p-1$  consists of some pairs  $a, a^{-1}$ , and 1 and  $-1$ .

$$\text{Multiply: } (p-1)! = 1^{p-3/2} \cdot 1 \cdot -1 = -1. \quad \square$$

Is  $-1$  a square in  $\mathbb{Z}_p$ ?

e.g. in  $\mathbb{Z}_5$ :  $x=2$  has  $x^2 = -1$  ✓

$\mathbb{Z}_7$ :  $0^2 = 0, 1^2 = 1, 2^2 = 4, 3^2 = 2$  ✗

(don't need to check 4, 5, 6 as  $-x$  squares to  $x^2$ )

$\mathbb{Z}_{13}$ :  $x=5$ :  $5^2 = -1$

$\mathbb{Z}_{19}$ : ✗ (can check)

Proposition 13 Let  $p$  be an odd prime. Then  $-1$  is a square number mod  $p$  iff  $p \equiv 1 \pmod{4}$ .

Proof First consider  $p = 4k+3$  ( $k \in \mathbb{N}_0$ ). Suppose for contradiction that  $x^2 = -1$  (in  $\mathbb{Z}_p$ )

We have  $x^{4k+2} = 1$  (FLT)

But  $x^{4k+2} = (x^2)^{2k+1} = (-1)^{2k+1} = -1$  ✗

For  $p = 4k+1$ , have  $(4k)! = -1$  (Wilson)

Compute  $(4k)! = 1 \times 2 \times \dots \times 2k (2k+1)(2k+2)\dots(4k)$

with  $(2k)!^2 = 1 \times 2 \times \dots \times 2k \times 1 \times 2 \times \dots \times 2k$

We have  $4k = -1, 4k-1 = -2, \dots, 2k+1 = -2k$

So  $(2k)!^2 = (4k)! (-1)^{2k} = (4k)! = -1$ . ✓

□

## Solving Congruence Equations

1) Solve  $7x \equiv 4 \pmod{30}$

Finding a solution: Have  $(7, 30) = 1$  so can run Euclid on 7 and 30 to find  $13 \times 7 - 3 \times 30 = 1$

so  $13 \times 7 \equiv 1 \pmod{30}$

so  $52 \times 7 \equiv 4 \pmod{30}$

So  $x = 52$  is a solution.

Other solutions:  $x' \equiv 52 \pmod{30}$  also works

No more solutions: If  $x'$  is a solution, want to show  $x' \equiv 52 \pmod{30}$ .

Have  $7x \equiv 4 \pmod{30}$ ,  $7x' \equiv 4 \pmod{30}$

so  $7x \equiv 7x' \pmod{30} \Rightarrow x \equiv x' \pmod{30}$

as 7 is invertible. □

Shorter method:  $7x \equiv 4 \pmod{30}$

$$\Leftrightarrow 13 \times 7x \equiv 13 \times 4 \pmod{30} \quad (13 \text{ invertible})$$

$\uparrow$   
inverse

$$\Leftrightarrow x \equiv 52 \pmod{30}$$

This is faster if you spot the inverse.

2) Solve  $10x \equiv 12 \pmod{34}$ .

$$10x \equiv 12 \pmod{34} \Leftrightarrow 10x = 12 + 34y, y \in \mathbb{Z}$$

$$\Leftrightarrow \cancel{5x} \equiv 5x = 6 + 17y$$

$$\Leftrightarrow 5x \equiv 6 \pmod{17}$$

and then solve as before.

A simultaneous congruence?

Do we expect a solution to

$$x \equiv 6 \pmod{17}$$
$$x \equiv 2 \pmod{19}$$

We guess yes, as 17 and 19 are coprime so "mod 17" and "mod 19" should be independent.

How about:

$$x \equiv 6 \pmod{34} \rightarrow x \text{ even}$$
$$x \equiv 11 \pmod{36} \rightarrow x \text{ odd}$$

34 and 36 aren't coprime.

Theorem 14 (Chinese Remainder Theorem) Let  $u, v$  be coprime.

Then for any  $a, b$ , there is an  $x$  with  $x \equiv a(u)$ ,  $x \equiv b(v)$ .  
Moreover the solution is unique mod  $uv$ .

Proof Existence: Have  $su + tv = 1$  for some  $s, t \in \mathbb{Z}$   
(as  $u, v$  coprime).

$$\begin{aligned} \text{Now, } su &\equiv 0(u), \quad su \equiv 1(v) \\ tv &\equiv 1(u), \quad tv \equiv 0(v) \end{aligned}$$

Hence  $x = a(tv) + b(su)$  has  $x \equiv a(u)$ ,  $b(v)$  ✓

Uniqueness: Certainly any  $x' \equiv x(uv)$  is also a solution.

Now suppose  $x' \equiv a(u)$ ,  $x' \equiv b(v)$ .

$$\text{so } x' \equiv x(u), \quad x' \equiv x(v)$$

$$\Rightarrow u|x' - x, \quad v|x' - x$$

but  $u, v$  are coprime, so  $\underline{uv|x' - x}$

$$\text{So } x' \equiv x(uv). \quad \checkmark$$

□

Remark Similarly, if  $u_1, u_2, \dots, u_k$  are pairwise coprime  
then  $\forall a_1, \dots, a_k \exists x$  such that  $x \equiv a_1(u_1), \dots$   
 $x \equiv a_i(u_i), \dots, x \equiv a_k(u_k)$  by induction.

An application of Fermat-Euler (RSA coding)

Normally to send a coded message



Pick 2 large distinct primes  $p, q$ . Let  $n = pq$ .

Fix a "coding exponent"  $e$ . To encode a message  $x \in \mathbb{Z}_n$ ,  
raise to power  $e$  in  $\mathbb{Z}_n$ .  $x \rightarrow x^e \pmod{n}$

How to decode? Seek  $d$  such that  $x^{ed} = x$ .

We know  $x^{\phi(n)} = 1$  in  $\mathbb{Z}_n$

$$\text{so } x^{k\phi(n)} = 1 \quad \forall k \in \mathbb{Z}$$

$$x^{k\phi(n)+1} = x$$

We want  $d$  with  $ed$  of form  $k\phi(n) + 1$  (some  $k \in \mathbb{Z}$ )  
i.e.  $ed \equiv 1 \pmod{\phi(n)}$

Easy to solve by running Euclid on  $e$  and  $\phi(n)$  assuming  
 $\text{hcf}(e, \phi(n)) = 1$ .

To encode:  $x \rightarrow x^e$  in  $\mathbb{Z}_n$ , so need to know  $n, e$

To decode:  $y \rightarrow y^d$  in  $\mathbb{Z}_n$ , so need  $n, d$  i.e. need  $n, e, \phi(n)$

$$\phi(n) = n-p-q+1$$

So need to know how to factorise  $n$  (hard)

This is the end of Chapter 1 (otherwise) (elementary NT)

## Chapter 2: The Reals

### The Need for the Reals

Have  $\mathbb{N}$  contained in  $\mathbb{Z}$  contained in  $\mathbb{Q}$ : why not stop there?

Proposition 1 There is no rational  $x$  with  $x^2 = 2$ .

(may assume  $x > 0$ , as  $(-x)^2 = x^2$ )

Proof 1 Suppose  $x^2 = 2$  for  $x = \frac{a}{b}$ ,  $a, b \in \mathbb{N}$

$$\text{So } \frac{a^2}{b^2} = 2 \Rightarrow a^2 = 2b^2$$

but exponent of 2 in prime factorisation  $a^2$  is even  
and in  $b^2$  it's odd, contradicting unique factorisation.  $\square$

Note Same proof shows that if  $\exists x \in \mathbb{Q}$  with  $x^2 = n$  ( $n \in \mathbb{N}$ )  
then  $n$  is a square number. (Each exponent in  
prime factorisation of  $n$  must be even)

Proof 2 Suppose  $x^2 = 2$ ,  $x = \frac{a}{b}$  where  $a, b \in \mathbb{N}$

So for any  $c, d \in \mathbb{Z}$ :  $cx+d$  is of the form  $\frac{e}{b}$  for  
some  $e \in \mathbb{Z}$ .

And so  $cx+d > 0 \Rightarrow cx+d > \frac{1}{b}$

But  $0 < x-1 < 1$  (as  $1 < x < 2$ )

↑ from  $x^2 = 2$

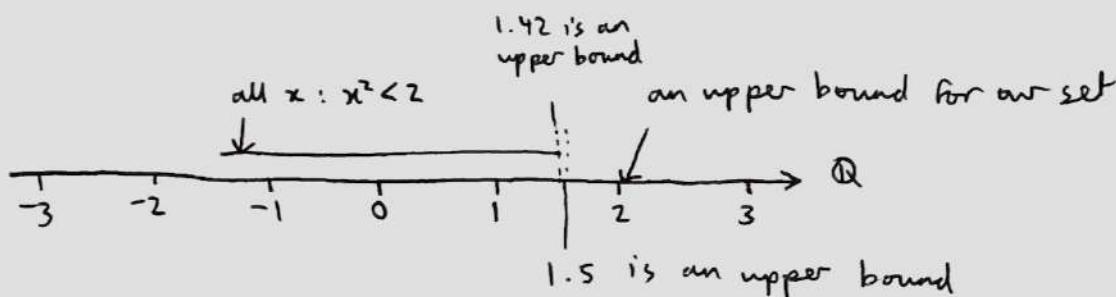
so  $0 < (x-1)^n < \frac{1}{b}$  for sufficiently large  $n$

This is a contradiction because  $(x-1)^n$  is of the

form  $cx+d$  where  $c, d \in \mathbb{Z}$  (using  $x^2 = 2$ ).  $\# \quad \square$

So " $\mathbb{Q}$  has a gap"

How to express this mentioning only  $\mathbb{Q}$ ?



In  $\mathbb{Q}$  there is no least upper bound (working in  $\mathbb{Q}$ : " $\Sigma$ " does not exist)

This is how we say inside  $\mathbb{Q}$  that " $\mathbb{Q}$  has a gap"

What we assume about the reals:

The reals are a set written  $\mathbb{R}$  with elements 0 and 1 ( $0 \neq 1$ ) equipped with operations + and  $\times$  and an ordering " $<$ " such that

1. + is commutative and associative (identity 0), every  $x$  has inverse
2.  $\times$  is commutative and associative (identity 1) and every nonzero  $x$  has an inverse
3.  $\times$  is distributive over + :  $a(b+c) = (ab) + (ac)$
4.  $\forall a, b$ , exactly one of  $a < b$ ,  $a = b$ ,  $a > b$  holds  
and  $a < b$ ,  $b < c \Rightarrow a < c$
5.  $\forall a, b, c$ :  $a < b \Rightarrow a+c < b+c$   
 $c > 0$ :  $a < b \Rightarrow ac < bc$

6.

For any set  $S$  of reals that is non-empty and bounded above,  $S$  has a least upper bound. (Least upper bound axiom)

$S$  is bounded above if  $\exists x \in \mathbb{R}$  with  $x \geq y \quad \forall y \in S$ .

Such an  $x$  is an upper bound for  $S$ .

Say  $x$  is the least upper bound of  $S$  if  $x$  is an upper bound for  $S$  and every other upper bound of  $S$  is at least  $x$ .

Remarks

1. From rules 1-5, can check e.g.  $0 < 1$   
If not then  $1 < 0 \Rightarrow 0 < -1$   
so  $0 < 1$  (multiplying by "positive"  $-1$ ) \*
2. May view  $\mathbb{Q}$  as contained in  $\mathbb{R}$  by identifying  $\frac{a}{b} \in \mathbb{Q}$  with  $ab^{-1} \in \mathbb{R}$
3. Least upper bound axiom (6) is false in  $\mathbb{Q}$ .
4. Why "non-empty and bounded above" in (6)?  
  - If  $S$  is not bounded above then it has no upper bound so no least one
  - If  $S$  is empty then every  $x \in \mathbb{R}$  is an upper bound so no least one.
5. Can construct  $\mathbb{R}$  from  $\mathbb{Q}$  and check axioms.

Examples of sets and least upper bounds

1.  $S = \{x \in \mathbb{R} : 0 \leq x \leq 1\}$   $[0, 1]$

L.U.B of  $S$   
is  $\sup(S)$



$x \leq 2 \quad \forall x \in S$        $x \neq \frac{3}{4} \quad \forall x \in S$       but " $x \leq \frac{3}{4}$ " is not an upper bound.

Least upper bound for  $S$  is 1.

1 is an upper bound:  $x \leq 1 \quad \forall x \in S$

Every upper bound  $y$  must have  $y \geq 1$  as  $1 \in S$ .

2.  $S = \{x \in \mathbb{R} : 0 < x < 1\}$   
 $(0, 1)$



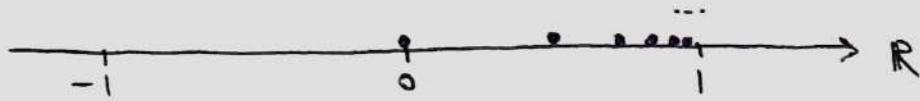
$\sup S = 1$  because 1 is an upper bound ( $x \leq 1 \quad \forall x \in S$ ).

~~No upper bound  $c < 1$  as  $1 + \frac{c}{2} \in S$~~

No upper bound  $c < 1$ : ( $c \geq \frac{1}{2}$  as  $\frac{1}{2} \in S$ )

If  $c \leq 1$  then  $0 < c < 1$  so  $\frac{1+c}{2} \in S$ ,  $\frac{1+c}{2} > c$ . \*

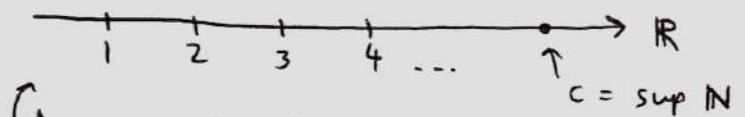
$$3. S = \left\{ 1 - \frac{1}{n} : n \in \mathbb{N} \right\}$$



1 is an upper bound. Is there an upper bound less than 1?

Proposition 2 (Axiom of Archimedes)  $\mathbb{N}$  is not bounded above in  $\mathbb{R}$ .

Proof



Assume not. Then  $\exists c \in \mathbb{R} : c = \sup \mathbb{N}$ .

So  $c-1$  is not an upper bound for  $n$  ( $c$  is L.U.B.).

~~So~~  $\exists n \in \mathbb{N} : n > c-1$ .

But then  $n+1 \in \mathbb{N}, n+1 > c$  ~~\*~~

□

Corollary 3 For each  $t > 0$ ,  $\exists n \in \mathbb{N}$  with  $\frac{1}{n} < t$ .

Proof Have some  $n \in \mathbb{N}$  with  $n > \frac{1}{t}$  (by prop. 2).

So  $\frac{1}{n} < t$ . □

Note that these results show that  $\mathbb{R}$  doesn't contain "infinitely big" or "infinitesimally small" elements.

Back to example 3: do have  $\sup S = 1$  because suppose  $c < 1$  is an upper bound. Then  $1 - \frac{1}{n} < c \quad \forall n \in \mathbb{N} \Rightarrow 1 - c < \frac{1}{n} \quad \forall n$  but  $1 - c > 0$  whilst  $\frac{1}{n}$  can be made arbitrarily small (see Corollary 3) so impossible. □

Warning: If  $S$  has a greatest element (like  $[0, 1]$ ) then that greatest element is  $\sup S$ . So  $\sup S \in S$ .

But if  $S$  has no greatest element e.g.  $(0, 1)$  then  $\sup S \notin S$ .

Note If  $S$  is a nonempty set of reals that is bounded below ( $\exists x$  s.t.  $x \leq y \forall y \in S$ )

Then the set  $-S = \{-y : y \in S\}$  is non-empty and bounded above, so has a least upper bound  $c$ .

So  $-c$  is the greatest lower bound of  $S$  - the infimum of  $S$  or  $\inf S$ .

In particular, if  $S$  is non-empty and bounded above and below, then it has a sup. and inf.

Theorem 4  $\exists x \in \mathbb{R}$  with  $x^2 = 2$ .

Proof let  $S = \{x \in \mathbb{R} : x^2 < 2\}$

Have  $S$  nonempty (e.g.  $1 \in S$ ) and bounded above (e.g. 2 is an upper bound)

So  $S$  has a sup,  $c$  such that  $1 \leq c \leq 2$ .

Claim  $c^2 = 2$

Proof of claim Suppose not. If  $c^2 < 2$

For  $0 < t < 1$ , have  $(c+t)^2 = c^2 + 2ct + t^2 \leq c^2 + 5t$   
as  $2ct \leq 2c$ ,  $t^2 < t$

and  $c^2 + 5t < 2$  for  $t < \frac{2-c^2}{5}$

so contradicts  $c$  being an upper bound for  $S$ .

If  $c^2 > 2$ : have  $0 < t < 1$ ,  $(c-t)^2 = c^2 - 2ct + t^2$   
 $\geq c^2 - 4t > 2$  for  $t$  small ( $t < \frac{c^2-2}{4}$ )  
contradicting  $c$  being <sup>least</sup> upper bound.  
 $c-t < c$  also an upper bound  $\times$

$$\text{So } c^2 = 2.$$

Remark same proof shows that  $\sqrt[n]{x}$  exists  $\forall n \in \mathbb{N}, x \in \mathbb{R}, x > 0$ .

A real that is not rational is called irrational

e.g.  $\sqrt{2}, \sqrt{3}, \sqrt{6}, \sqrt{15}$ .

Claim  $2 + 3\sqrt{5}$  is irrational

Proof If  $2 + 3\sqrt{5} = \frac{a}{b}$  for some  $a, b \in \mathbb{Z}$

$$\text{then } \sqrt{5} = \frac{a-2b}{3b} \in \mathbb{Q}. \quad \times$$

Also "the rationals are dense": if  $a, b \in \mathbb{R}, a < b$   
then  $\exists c \in \mathbb{Q}: a < c < b$ .

May assume  $a, b \geq 0$

Choose  $n \in \mathbb{N}$  with  $\frac{1}{n} < b-a$

Among  $\frac{0}{n}, \frac{1}{n}, \frac{2}{n}, \dots$  there is a final one that is  $\leq a$ , say  
 $\frac{q}{n}$  (else  $a$  would be an upper bound for  $\{\frac{0}{n}, \frac{1}{n}, \frac{2}{n}, \dots\}$   
contradicting axiom of Archimedes)

$$\text{So } a < \frac{q+1}{n} < b \quad \checkmark$$

Also "the irrationals are dense":  $\forall a, b \in \mathbb{R}, a < b,$

$\exists$  irrational  $c$  with  $a \leq c < b$ .

Indeed  $\exists$  rational  $c$  with  $a\sqrt{2} \leq c < b\sqrt{2}$ , so

$$a < \frac{c}{\sqrt{2}} < b. \quad \checkmark$$

What should " $1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots = 2$ " mean?

How about " $0.33333\dots = \frac{1}{3}$ "?

Presumably, that  $1, 1 + \frac{1}{2}, 1 + \frac{1}{2} + \frac{1}{4}, \dots$  should "tend to" 2

$0.3, 0.33, 0.333, \dots$  should "tend to"  $\frac{1}{3}$

This isn't the same as " $x_n$  are getting closer and closer to  $c$ "

It's not " $\forall \varepsilon > 0 \ \exists n \text{ with } c - \varepsilon < x_n < c + \varepsilon$ "

We want the sequence to get and stay within  $\varepsilon$  of  $c$ .

We say that  $x_1, x_2, x_3, \dots$  tends to  $c$  if

$\forall \varepsilon > 0 \ \exists N \text{ s.t. } \forall n > N \text{ have } c - \varepsilon < x_n < c + \varepsilon.$   
 $(|x_n - c| < \varepsilon)$

In  $\mathbb{R}$   $|a|$  of  $a \in \mathbb{R}$  is  $a$  if  $a \geq 0$   
 $-a$  if  $a < 0$

If  $x_n$  tends to  $c$ , can write  $x_n \rightarrow c$  (as  $n \rightarrow \infty$ )

Examples

1.  $\frac{1}{2}, \frac{1}{2} + \frac{1}{4}, \frac{1}{2} + \frac{1}{4} + \frac{1}{8}, \dots$

This is  $x_1, x_2, x_3, \dots$  where  $x_n = 1 - \frac{1}{2^n}$  (inductively)

Claim  $x_n \rightarrow 1$

Proof Given  $\varepsilon > 0$

Choose  $N \in \mathbb{N}$  with  $N > \frac{1}{\varepsilon}$  (or  $\frac{1}{N} < \varepsilon$ )

Then  $\forall n > N : |x_n - 1| = \frac{1}{2^n} < \frac{1}{n} \leq \frac{1}{N} < \varepsilon$

✓ □

2. The constant sequence  $c, c, c, c, \dots$  ( $x_n = c \quad \forall n$ )

Claim  $x_n \rightarrow c$

Proof Given  $\varepsilon > 0$ :

Have  $|x_n - c| < \varepsilon \quad \forall n$

□

3.  $x_n = (-1)^n : -1, 1, -1, 1, -1, 1, \dots$

Claim There is no  $c \in \mathbb{R}$  with  $x_n \rightarrow c$

Proof Suppose  $x_n \rightarrow c$ . Choose  $\varepsilon = 1$ :

So  $\exists N \in \mathbb{N}$  s.t.  $\forall n > N$  have  $|x_n - c| < 1$ .

$|1 - c|$  and  $|(-1) - c| < 1$

So by A inequality  $|1 - (-1)| < 2$  which is a contradiction.

□

4.  $x_n = \begin{cases} \frac{1}{n} & \text{if } n \text{ odd} \\ 0 & \text{if } n \text{ even} \end{cases}$

$$1, 0, \frac{1}{3}, 0, \frac{1}{5}, 0, \frac{1}{7}, \dots$$

PTO

Claim  $x_n \rightarrow 0$

Proof Given  $\varepsilon > 0$ :

Choose  $N \in \mathbb{N}$  with  $\frac{1}{N} < \varepsilon$

Then  $\forall n \geq N: x_n = \frac{1}{n}$  or 0, so  $|x_n - 0| \leq \frac{1}{n} \leq \frac{1}{N} < \varepsilon$

□

Notes - If  $x_n \rightarrow c$  for some  $c$ , then  $x_1, x_2, x_3, \dots$  is convergent.  
or  $(x_n)$  or  $(x_n)_{n=1}^{\infty}$

Otherwise, it is divergent.

- Same idea as example 3: shows limits are unique.

Suppose  $x_n \rightarrow c, x_n \rightarrow d$ .

Suppose  $c \neq d$  and choose  $\varepsilon = \frac{1}{2}|c-d|$

Then  $\exists N \in \mathbb{N}$  with  $|x_n - c| < \varepsilon \quad \forall n \geq N$

and  $\exists M \in \mathbb{N}$  with  $|x_n - d| < \varepsilon \quad \forall n \geq M$

Now for any  $n \geq \max(M, N)$ , have

$|x_n - c|, |x_n - d| < \varepsilon$  so  $|c - d| < 2\varepsilon$  \* as  
 $|c - d| = 2\varepsilon$ . □

- Given a series  $\sum_{n=1}^{\infty} x_n$ ,  $\sum_{n=1}^k x_n$  is the " $k^{th}$  partial sum"

Limits do behave as we expect.

e.g. If  $x_n \leq d \quad \forall n$  and  $x_n \rightarrow c$ , then  $c \leq d$ .

Suppose  $c > d$ . Choose  $\varepsilon = |c-d|$ , then

$\exists N \in \mathbb{N}$  s.t.  $\forall n \geq N$  have  $|x_n - c| < \varepsilon$

but  $|x_n - c| < \varepsilon \Rightarrow x_n > d$  \*

□

Warning If  $x_n < d$  and  $x_n \rightarrow c$ , need not have  $c < d$ :

e.g.  $\frac{1}{2}, \frac{1}{2} + \frac{1}{4}, \frac{1}{2} + \frac{1}{4} + \frac{1}{8}, \dots \rightarrow 1$

and all  $x_n < 1$  so here " $c = d$ ".

How about  $x_n + y_n$ ?

Proposition 5 If  $x_n \rightarrow c$  and  $y_n \rightarrow d$ , then  $x_n + y_n \rightarrow c + d$ .

Ideas: "late  $x_n$  are close to  $c$ , late  $y_n$  are close to  $d$ "

Proof Given  $\varepsilon > 0$ :

- Have  $x_n \rightarrow c$  so  $\exists N \in \mathbb{N}$  s.t.  $|x_n - c| < \frac{\varepsilon}{2} \forall n > N$ .

Also  $y_n \rightarrow d$  so  $\exists M \in \mathbb{N}$  s.t.  $|y_n - d| < \frac{\varepsilon}{2} \forall n > M$ .

∴ for all  $n > \max(M, N)$  have

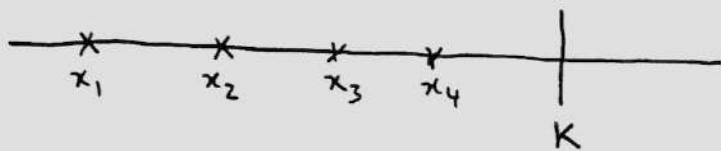
$$|(x_n + y_n) - (c + d)| \leq |x_n - c| + |y_n - d| < \frac{\varepsilon}{2} + \frac{\varepsilon}{2} = \varepsilon.$$

$\Delta$  ineq. □

Remark If we'd used " $|x_n - c| < \varepsilon$ " instead of " $|x_n - c| < \frac{\varepsilon}{2}$ ", then at the end would have got  $|(x_n + y_n) - (c + d)| < 2\varepsilon$  instead which is still fine, as  $\varepsilon$  could be replaced with  $\frac{\varepsilon}{2}$ .

A sequence  $x_1, x_2, \dots$  is increasing if  $x_{n+1} \geq x_n \forall n$ .

Theorem 6 If  $x_1, x_2, \dots$  is increasing and bounded above, then it converges.



Remark If we lived in  $\mathbb{Q}$  then this would be false:

e.g.  $1, 1.4, 1.41, 1.414, 1.4142 \rightarrow \sqrt{2}$   
but  $\sqrt{2} \notin \mathbb{Q}$ .

Proof Let  $c = \sup(x_1, x_2, \dots)$

Claim  $x_n \rightarrow c$

Proof Given  $\varepsilon > 0$ :

$\exists n$  s.t.  $x_n > c - \varepsilon$  (else  $c - \varepsilon$  is an upper bound  $< c$ )

so  $\forall n > N$ :

$$c - \varepsilon < x_N \leq x_n < c$$

$$\text{so } |x_N - c| < \varepsilon$$

□

So "a bounded monotone sequence is convergent"

↑  
increasing or decreasing

i) Proposition 7 i)  $\sum_{n=1}^{\infty} \frac{1}{n}$  diverges

ii)  $\sum_{n=1}^{\infty} \frac{1}{n^2}$  converges

Note: There is no closed form for  $1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n}$   
or  $1 + \frac{1}{4} + \frac{1}{9} + \dots + \frac{1}{n^2}$

Idea  $1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \frac{1}{5} + \frac{1}{6} + \frac{1}{7} + \frac{1}{8} + \frac{1}{9} + \dots$

$$\cancel{1} + \cancel{\frac{1}{2}} + \cancel{\frac{1}{4}} + \cancel{\frac{1}{8}} + \cancel{\frac{1}{16}} + \cancel{\frac{1}{32}}$$

$$\geq 1 + \underbrace{\frac{1}{2}}_{\frac{1}{2}} + \underbrace{\frac{1}{4} + \frac{1}{4}}_{\frac{1}{2}} + \underbrace{\frac{1}{8} + \frac{1}{8} + \frac{1}{8} + \frac{1}{8}}_{\frac{1}{2}} + \frac{1}{16} + \dots$$

Proof (i) Have  $\frac{1}{3} + \frac{1}{4} > \frac{1}{2}$

$$\frac{1}{5} + \frac{1}{6} + \frac{1}{7} + \frac{1}{8} > \frac{1}{2}$$

In general  $\frac{1}{2^n+1} + \frac{1}{2^{n+2}} + \dots + \frac{1}{2^{n+1}} > \frac{2^n}{2^{n+1}} = \frac{1}{2}$

Hence partial sums of  $\sum_{n=1}^{\infty} \frac{1}{n}$  are unbounded.

So it is not convergent.

(ii) Idea  $1 + \frac{1}{2^2} + \frac{1}{3^2} + \frac{1}{4^2} + \frac{1}{5^2} + \frac{1}{6^2} + \frac{1}{7^2} + \frac{1}{8^2} + \dots$

$$\leq 1 + \underbrace{\frac{1}{2^2} + \frac{1}{2^2}}_{\frac{2}{2^2}} + \underbrace{\frac{1}{4^2} + \frac{1}{4^2} + \frac{1}{4^2} + \frac{1}{4^2}}_{\frac{4}{4^2}} + \underbrace{\frac{1}{8^2} + \dots}_{\frac{8}{8^2}}$$

$$= 1 + \frac{1}{2} + \frac{1}{4} + \dots$$

We have  $\frac{1}{2^2} + \frac{1}{2^2} \leq \frac{2}{2^2} = \frac{1}{2} \quad \frac{1}{4^2} + \frac{1}{5^2} + \frac{1}{6^2} + \frac{1}{7^2} \leq \frac{4}{4^2} = \frac{1}{4}$

In general  $\frac{1}{(2^n)^2} + \frac{1}{(2^n+1)^2} + \dots + \frac{1}{(2^{n+1}-1)^2} \leq \frac{2^n}{(2^n)^2} = \frac{1}{2^n}$

So partial sums are bounded. (by  $1 + \frac{1}{2} + \frac{1}{4} + \dots = 2$ )  $\forall n$ .

□

### Remarks

1)  $1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \dots$  is called the harmonic series

2)  $\sum_{n=1}^{\infty} \frac{1}{n^2} = \frac{\pi^2}{6}$  - proved in Part II Linear Analysis

3) Decimal Expansions - what should " $0.a_1 a_2 a_3 \dots$ " mean?  
It should be the limit (if it exists) of

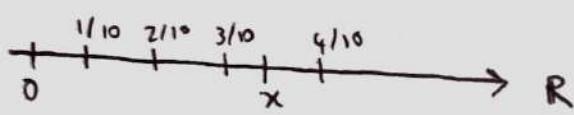
$$0.a_1, \cancel{a_2} 0.a_1 a_2, 0.a_1 a_2 a_3, \dots$$

So ~~too~~ define  $0.a_1 a_2 a_3 \dots = \sum_{n=1}^{\infty} \frac{a_n}{10^n}$

(converges as all terms  $> 0$  and partial sums bounded  
e.g. by 1)

Conversely given  $x \in \mathbb{R}$ ,  $0 < x < 1$ , want to write

$$x = 0.a_1 a_2 a_3 \dots \text{ for some } a_1, a_2, \dots \in \{0, 1, \dots, 9\}$$



Choose the greatest  $a_1$  in  $\{0, 1, \dots, 9\}$  with  $\frac{a_1}{10} \leq x$

$$\text{so } 0 \leq x - \frac{a_1}{10} < \frac{1}{10}.$$

Now greatest  $a_2 \in \{0, 1, \dots, 9\}$  s.t.  $\frac{a_1}{10} + \frac{a_2}{100} \leq x$

$$\text{so } 0 \leq x - \frac{a_1}{10} - \frac{a_2}{100} < \frac{1}{100}.$$

Inductively: we obtain  $a_1, a_2, a_3, \dots \in \{0, 1, \dots, 9\}$

such that  $0 \leq x - \sum_{n=1}^k \frac{a_n}{10^n} < \frac{1}{10^k} \forall k$

Thus  $\sum_{n=1}^{\infty} \frac{a_n}{10^n} = x$ , i.e.  $0.a_1 a_2 a_3 \dots = x$ .

Remarks

1. A decimal expansion  $0.a_1 a_2 a_3 \dots$  is recurrent if  $a_{n+k} = a_n \quad \forall n > N$ , some  $N$  and  $k$ .

e.g.  $0.3178426426426\dots$

Can check  $x = 0.a_1 a_2 \dots$  is recurrent iff  $x$  is rational.

2. Decimal expansions need not be unique:

$$0.3700000\dots = 0.3699999\dots$$

3. That's the only way to have non-unique expansions

If  $0.a_1 a_2 \dots$ ,  $0.b_1 b_2 \dots$ , are the ~~two~~ decimal expansions (different) of the same number then  $\exists N \in \mathbb{N}$  where

$$a_n = b_n \quad \forall n < N$$

$$a_N = b_N - 1, \quad a_n = 9, b_n = 0 \quad \forall n > N$$

or vice versa.

The number e

$$\text{Define } e = 1 + \frac{1}{1!} + \frac{1}{2!} + \frac{1}{3!} + \dots$$

$\uparrow$        $\uparrow$        $\uparrow$        $\uparrow$   
 1      1       $\geq \frac{1}{2}$        $\geq \frac{1}{4}$

Partial sums are bounded by  $1 + 1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots = 3$ ,  
 and all terms are  $> 0$ . ( $\frac{1}{n!} < \frac{1}{2^{n-1}} \quad \forall n \geq 2$ , inductively)

$$\text{If we write } 0! = 1 \text{ then } e = \sum_{n=0}^{\infty} \frac{1}{n!}$$

Say  $x \in \mathbb{R}$  is algebraic if it is a root of a non-zero polynomial with integer coefficients.

i.e.  $a_d x^d + a_{d-1} x^{d-1} + \dots + a_1 x + a_0 = 0$   
 where  $a_i \in \mathbb{Z}$ , for some  $d \in \mathbb{Z}$  ( $a_d \neq 0$ ).

Every rational is algebraic :  $\frac{p}{q}$  is a root of  $qx - p = 0$

$\sqrt{2}$  is algebraic :  $x^2 - 2 = 0$

$\sqrt{2} + 1$  is algebraic :  $((\sqrt{2} + 1) - 1)^2 = 2$  so  $(x-1)^2 - 2 = 0$ .

Are all reals algebraic?

Proposition 8  $e$  is irrational.

Proof Suppose rational:  $e = \frac{p}{q}$ ,  $p, q \in \mathbb{Z}$ ,  $q > 1$   
 $\Rightarrow q! e = (\text{some integer})$

$$\text{so } \sum_{n=0}^{\infty} \frac{q!}{n!} \in \mathbb{Z}$$

$$\text{So } q! + \frac{q!}{1!} + \frac{q!}{2!} + \dots + \frac{q!}{q!} + \frac{q}{(q+1)!} + \frac{q}{(q+2)!} + \dots$$

an integer

$$\text{Also } \frac{q!}{(q+1)!} = \frac{1}{q+1}, \quad \frac{q!}{(q+2)!} = \frac{1}{(q+1)(q+2)} \leq \frac{1}{(q+1)^2}$$

$$\text{Generally } \frac{q!}{(q+n)!} \leq \frac{1}{(q+1)^n}$$

$$\text{so } \sum_{n=q+1}^{\infty} \frac{q^n}{n!} \leq \frac{1}{q+1} + \frac{1}{(q+1)^2} + \frac{1}{(q+2)^2} + \dots = \frac{1}{q} < 1$$

as  $q \geq 2$ .

So this is not an integer, contradicting the assumption

that  $\sum_{n=0}^{\infty} \frac{q^n}{n!}$  is an integer. \*

1

A real that is not algebraic is called transcendental.

Theorem 9 The number  $c = \sum_{n=1}^{\infty} \frac{1}{10^n}$  is transcendental.

To prove, we need 2 facts about polynomials.

Fact 1 For any polynomial  $P$ ,  $\exists$  a constant  $k$  such that

$$|P(x) - P(y)| \leq k|x-y| \quad \forall 0 \leq x, y \leq 1$$

$$\underline{\text{proof}} \quad P(x) = a_d x^d + \dots + a_0$$

$$P(x) - P(y) = a_d (x^d - y^d) + a_{d-1} (x^{d-1} - y^{d-1}) + \dots + a_1 (x - y)$$

$$= (x-y) \left[ a_d (x^{d-1} + x^{d-2}y + \dots + y^{d-1}) + \dots + a_1 \right]$$

$$\text{So } |P(x) - P(y)| \leq |x-y| \left[ (|a_d| + |a_{d-1}| + \dots + |a_1|) d \right] \text{ as } 0 \leq x, y \leq 1.$$

Fact 2 A nonzero polynomial of degree  $d$  has  $\leq d$  roots.

Proof Given a polynomial  $p$  of degree  $d$ :

If  $p$  has no roots it's certainly true

If  $p$  has a root, have  $a$  a root:

$$p(x) = (x-a) q(x) \quad \text{for some } q \text{ of degree } d-1.$$

So each root of  $p$  is either  $a$  or a root of  $q$   
but  $q$  has at most  $d-1$  roots inductively.  $\square$

Proof of Theorem Write  $c_n = \sum_{k=0}^n \frac{1}{10^k}$  so  $c_n \rightarrow c$ .

Suppose for contradiction that polynomial  $P$  has  $c$  as a root.

Then  $\exists k$  with  $|P(c_n)| \leq k|x-y| \quad \forall 0 \leq x, y \leq 1$

Say  $P$  has degree  $d$ :  $P(x) = a_d x^d + \dots + a_0 \quad (a_i \in \mathbb{Z}, a_d \neq 0)$

$$\text{Now } |c - c_n| = \sum_{k=n+1}^{\infty} \frac{1}{10^k} \leq \frac{2}{10^{(n+1)!}}$$

$$\text{Also } c_n = \frac{a}{10^n!} \quad \text{for some } a \in \mathbb{Z} \quad \text{so } P(c_n) = \frac{b}{10^{dn!}} \quad (b \in \mathbb{Z})$$

(since  $P\left(\frac{s}{t}\right) = \frac{q}{t^d}$  for some  $q \in \mathbb{Z}$  whenever  $s, t \in \mathbb{Z}$ )

But for  $n$  large enough,  $c_n$  is not a root of  $p$  as finitely many roots. So  $|P(c_n)| > \frac{1}{10^{dn!}}$  i.e.  $|P(c_n) - P(c)| > \frac{1}{10^{dn!}}$

Thus  $\frac{1}{10^{dn!}} \leq k \frac{2}{10^{(n+1)!}}$ , a contradiction for  $n$  sufficiently large.  $\square$

(I may write this up better)

Written up in LaTeX - see Overleaf doc

Remarks

1. Same proof shows that any real  $x$  such that  $\forall n \exists \text{ rational } \frac{p}{q} \text{ with } 0 < |x - \frac{p}{q}| < \frac{1}{q^n}$  is transcendental.

" $x$  has very good rational approximations"

2. Such  $x$  are called Liouville numbers so could view Thm 9 as "every Liouville number is transcendental"
3. Theorem 9 does not show  $e$  is transcendental (but it is).
4. Another proof of existence of transcendentals is in Ch 4.

### The Complex Numbers

Some polynomials have no real roots e.g.  $x^2 + 1 = 0$

We'll try to "force" an  $x$  with  $x^2 = -1$

The complex numbers  $\mathbb{C}$  consist of  $\mathbb{R}^2$  with operations

+ and  $\times$ , defined by  $(a, b) + (c, d) = (a+c, b+d)$

and  $(a, b) \cdot (c, d) = (ac - bd, ad + bc)$

Can view  $\mathbb{R}$  as contained in  $\mathbb{C}$  by identifying  $a \in \mathbb{R}$  with  $(a, 0)$

Let  $i = (0, 1)$  Then  $i^2 = (0, 1) \cdot (0, 1) = (-1, 0) = -1$

Note that every  $z \in \mathbb{C}$  is of the form  $a+bi$ ,  $a, b \in \mathbb{R}$

Remarks 1.  $\mathbb{C}$  obeys all usual algebraic rules (note  $\mathbb{C}$  is a field)

2. Every nonzero polynomial (even with complex coefficients) has a root in  $\mathbb{C}$ .

This is the Fundamental Theorem of Algebra

### Chapter 3 - Sets and Functions

Definition A set is a collection of mathematical objects

e.g.  $\mathbb{R}$ ,  $\mathbb{N}$ ,  $\{1, 5, 9\}$ ,  $[0, 1]$

Two sets with the same members are the same.

If  $\forall x, x \in A \Leftrightarrow x \in B$  then  $A = B$ .

(Note  $\{3, 4, 4, 6\} = \{3, 4, 6\}$ )

Subsets Given set  $A$  and property  $P(x)$ , can form  $\{x \in A : P(x)\}$ , the subset of  $A$  with property  $P$ .

$B$  is a subset of  $A$  if  $\forall x: x \in B \Rightarrow x \in A$ .

Written  $B \subset A$  or  $B \subseteq A$ .

We have  $A = B \iff A \subseteq B, B \subseteq A$

### Unions and Intersections

Union :  $A \cup B = \{x : x \in A \text{ or } x \in B\}$

Intersection :  $A \cap B = \{x : x \in A \text{ and } x \in B\}$

A and B are disjoint if  $A \cap B = \emptyset \leftarrow$  empty set

Intersection is a special case of subset selection

$A \cap B = \{x \in A : x \in B\}$

Difference :  $A \setminus B = \{x \in A : x \notin B\}$

$\cup$  is distributive over  $\cap$  :  $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$

and  $\cap$  over  $\cup$  :  $A \cap (B \cup C) = (A \cap B) \cup (A \cap C) \quad (*)$

Prove (\*) :

LHS  $\subset$  RHS : Given  $x \in A \cap (B \cup C)$

$\Rightarrow x \in A$  and either  $x \in B$  or  $x \in C$

If  $x \in B$  have  $x \in A, x \in B$

so  $x \in A \cap B \Rightarrow x \in \text{RHS}$

Same for  $x \in C$  case.

RHS  $\subset$  LHS : Given  $x \in (A \cap B) \cup (A \cap C)$

WLOG if  $x \in (A \cap B)$  then  $x \in A$  and  $x \in B \cup C$   
so  $x \in \text{LHS}$ . □

Can have bigger unions e.g. infinite union  $\bigcup_{n=1}^{\infty} A_n = \bigcup_{n \in \mathbb{N}} A_n$  ← index set  
Not defined as a limit

Ordered Pairs

For any  $a, b$  can form ordered pair  $(a, b)$

$$(a, b) = (c, d) \iff a = c \text{ and } b = d$$

For sets  $A$  and  $B$ , can form their product  $A \times B$ , the set of all ordered pairs  $\{(a, b) : a \in A, b \in B\}$ .

e.g. can view  $\mathbb{R}^2 = \mathbb{R} \times \mathbb{R}$  as a plane.

[Note] If we wished, could define  $(a, b) = \{\{a\}, \{a, b\}\}$  and check "key point" above.

Power Sets

For any set  $X$ , can form the power set  $\mathcal{P}(X)$  consisting of all the subsets of  $X$ .

$$\mathcal{P}(X) = \{Y : Y \subseteq X\}$$

e.g. if  $X = \{1, 2\}$  then  $\mathcal{P}(X) = \{\emptyset, \{1\}, \{2\}, \{1, 2\}\}$

Note / Warning For a set  $A$ , can form  $\{x \in A : p(x)\}$  but cannot form  $\{x : p(x)\}$

If we could, then consider  $\{x : x \notin x\} = X$

Does  $X \in X$ ? If yes: Then  $X \notin X$  by defn of  $X$

If no:  $X \notin X$  so  $X \in X$  by defn of  $X$   $\ast\ast$

(Russell's paradox)

Similarly, there is no "universal set"  $Y$  with  $x \in Y \ \forall x$ . Otherwise could form  $X$  above by subset selection.

To guarantee a given set exists, must somehow obtain it from known sets.

## Finite Sets

Write  $\mathbb{N}_0 = \{0, 1, 2, \dots\}$

For  $n \in \mathbb{N}_0$ , say set  $A$  has size  $n$  if we can write  
 $A = \{a_1, a_2, \dots, a_n\}$  with  $a_i$  distinct.

Say  $A$  is finite if  $\exists n \in \mathbb{N}_0$  where  $A$  has size  $n$   
(otherwise infinite)

Note  $A$  cannot have size  $m$  and size  $n$  for  $n \neq m$ .

Suppose  $A$  has size  $m$  and size  $n$ ,  $n, m > 0$ .

Then, removing an element, we get a set of size  
 $n-1$  and  $m-1$  - done by induction (on max element).

Proposition 1 A set of size  $n$  has exactly  $2^n$  subsets.

Proof 1 May assume set is  $\{1, 2, \dots, n\}$  wlog.

To specify a subset  $S$  we must say if  $1 \in S$  or  $1 \notin S$ ,  
then if  $2 \in S$  or  $2 \notin S$ , and so on.

So no. of choices is  $\underbrace{2 \times 2 \times \dots \times 2}_{n \text{ times}} = 2^n$ . □

Proof 2 Induction on  $n$ :  $n=0 \checkmark$

Take  $n > 0$ .

1	2	...	$n-1$	$n$
---	---	-----	-------	-----

For  $T \subseteq \{1, 2, \dots, n-1\}$ , how many  $S \subseteq \{1, \dots, n\}$   
have  $S \cap \{1, \dots, n-1\} = T$ ?

Exactly 2:  $T$  and  $T \cup \{n\}$ .

Hence # subsets of  $\{1, 2, \dots, n-1\} \times 2$  is

# subsets of  $\{1, \dots, n\}$  so done by induction. □

If  $A$  has size  $n$  then  $|A| = n$

## Binomial Coefficients

For  $n \in \mathbb{N}_0$ ,  $0 \leq k \leq n$ , write  $\binom{n}{k}$  for number of subsets of an  $n$ -set that are of size  $k$ .  
( $n$ -set: size  $n$ )

$$\binom{n}{k} = \left| \{S \subseteq \{1, 2, \dots, n\} : |S| = k\} \right|$$

Note that  $\binom{n}{0} = 1$ ,  $\binom{n}{n} = 1$ ,  $\binom{n}{1} = n$  ( $n > 0$ )

Note also that  $\binom{n}{0} + \binom{n}{1} + \dots + \binom{n}{n} = 2^n$  because we are counting all subsets.

Also 1.  $\binom{n}{k} = \binom{n}{n-k}$   $n \in \mathbb{N}_0$ ,  $0 \leq k \leq n$

2.  $\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}$

#  $k$ -subsets of  $\{1, \dots, n\}$  without  $n$  is  $\binom{n-1}{k}$

#  $k$ -subsets of  $\{1, \dots, n\}$  is  $\binom{n-1}{k-1}$  as we

choose  $n$  and then pick remaining  $k-1$  from  $n-1$ .

so  $\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}$ .

## Pascal's Triangle

		1	1			
	1	2	1			
	1	3	3	1		
	1	4	6	4	1	
1	5	10	10	5	1	
1	6	15	20	15	6	1

$$\underline{\text{Proposition 2}} \quad \binom{n}{k} = \frac{n(n-1)\dots(n-k+1)}{k!}$$

Proof # of ways to name an element is  $n$ , then  $n-1, \dots$

so # of ways to specify a  $k$ -set is

$$n(n-1)(n-2)\dots(n-k+1)$$

# of times a given  $k$ -set is named:  $k!$

Hence #  $k$ -sets is  $\frac{n(n-1)\dots(n-k+1)}{k!}$  □

Can also write  $\binom{n}{k} = \frac{n!}{k!(n-k)!}$

Notice  $\binom{n}{3} \approx \frac{n^3}{6}$  for large  $n$ .

An application of binomial coefficients:

$$\underline{\text{Theorem 3 (Binomial theorem)}} \quad (a+b)^n = \binom{n}{0} a^n + \binom{n}{1} a^{n-1} b + \dots + \binom{n}{n} b^n$$

$\forall a, b \in \mathbb{R}, n \in \mathbb{N}$

Proof When we expand  $(a+b)^n = (a+b)(a+b)\dots(a+b)$   
we get terms of the form  $a^k b^{n-k}$  ( $0 \leq k \leq n$ )

# terms  $a^k b^{n-k}$  is  $\binom{n}{k}$  so have

$$(a+b)^n = \sum_{k=0}^n \binom{n}{k} a^k b^{n-k}$$

$$= \sum_{k=0}^n \binom{n}{n-k} a^k b^{n-k}$$

□

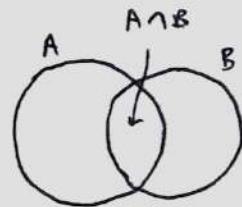
e.g.  $(1+x)^n = 1 + nx + \frac{n(n-1)}{2} x^2 + \binom{n}{3} x^3 + \dots + nx^{n-1} + x^n$

so for  $x$  small,  $(1+x)^n \approx 1 + nx$

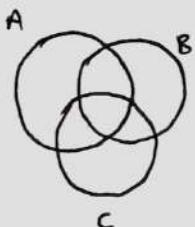
What can we say about the relationships between sizes of unions and intersections of finite sets?

e.g.  $|A \cup B| = |A| + |B| - |A \cap B|$

to make sure  $A \cap B$  not counted twice



Now  $|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |B \cap C| - |C \cap A| + |A \cap B \cap C|$



Notation

$$S_1 \cap S_2 = S_{\{1,2\}}$$

#### Theorem 4

(Inclusion-Exclusion Theorem) Let  $S_1, \dots, S_n$  be finite sets. Then  $|S_1 \cup \dots \cup S_n|$  is

$$\sum_{|A|=1} |S_A| - \sum_{|A|=2} |S_A| + \sum_{|A|=3} |S_A| - \dots + (-1)^{n+1} \sum_{|A|=n} |S_A|$$

with  $S_A$  meaning  $\bigcap_{i \in A} S_i$  and  $\sum_{|A|=R}$  as a sum over all  $A \subseteq \{1, \dots, n\}$  of size R

Proof Let  $x \in \text{LHS}$ : Say  $x$  belongs to  $S_i$  for k of the  $S_i$ . We want  $x$  counted exactly once on RHS.

# A,  $|A|=1$  with  $x \in S_A = k$

# A,  $|A|=2$  with  $x \in S_A = \binom{k}{2}$  ← (must choose 2 of the i that it belongs to)

In general # A,  $|A|=r$  with  $x \in S_A$  is  $\binom{k}{r}$ .

# times x counted on RHS is

$$k - \binom{k}{2} + \binom{k}{3} - \dots + (-1)^{k+1} \binom{k}{k} = 1 - (1 + (-1))^k \quad (\text{by BT})$$

= 1.

□

## Functions

For sets  $A$  and  $B$ , a function from  $A$  to  $B$  is a rule that assigns each  $x \in A$  a unique point  $f(x) \in B$ .

More precisely, a function from  $A$  to  $B$  is a set  $f \subseteq A \times B$  with  $\forall x \in A$  a unique  $y \in B$  with  $(x, y) \in f$ .

If  $(x, y) \in f$  then have  $f(x) = y$ .

Examples

- $f(x) = x^2$  from  $\mathbb{R}$  to  $\mathbb{R}$ : can say "f:  $\mathbb{R} \rightarrow \mathbb{R}$  given by  $f(x) = x^2$ " or " $f: \mathbb{R} \rightarrow \mathbb{R}$ "  
 $x \mapsto x^2$

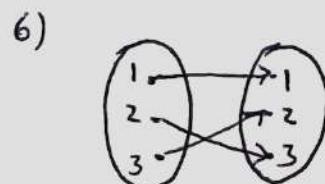
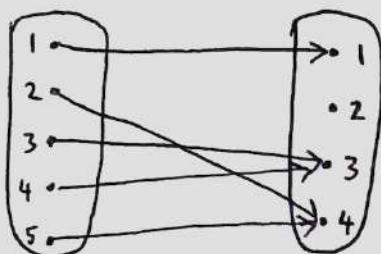
- Not a function from  $\mathbb{R} \rightarrow \mathbb{R}$  with  $f(x) = \frac{1}{x}$ .  
not defined for  $x = 0$

- Not a function  $\mathbb{R} \rightarrow \mathbb{R}$ :  $f(x) = \pm \sqrt{|x|}$  as  $f(2)$  can be  $\sqrt{2}$  and  $-\sqrt{2}$ : not unique

- $f(x) = \begin{cases} 1 & \text{if } x \text{ rational} \\ 0 & \text{if not} \end{cases}$  f:  $\mathbb{R} \rightarrow \mathbb{R}$   
(This is a valid function)

Examples of functions - cont.

$$5) \quad A = \{1, 2, 3, 4, 5\} \quad B = \{1, 2, 3, 4\}$$



Say  $f: A \rightarrow B$  is injective if  $\forall a, a' \in A$ , then  
 $a \neq a' \Rightarrow f(a) \neq f(a')$

or equivalently  $f(a) = f(a') \Rightarrow a = a'$ .

"different points stay different"

Say  $f: A \rightarrow B$  is surjective if  $\forall b \in B, \exists a \in A$  with  
 $f(a) = b$ . "everything in B is hit"

Say  $f: A \rightarrow B$  is bijection if it's injective and surjective.  
 "everything hits, exactly once" or "f pairs up A and B"

For  $f: A \rightarrow B$ : A is the domain, B is the range

The image of f is  $\{f(a) : a \in A\} = \{b \in B : f(a) = b, a \in A\}$

Image is "everything that is hit"

e.g. For  $f: \mathbb{R} \rightarrow \mathbb{R}, x \mapsto x^2$ , image of f is  $\{y \in \mathbb{R} : y \geq 0\}$

↳ Must specify domain and range.

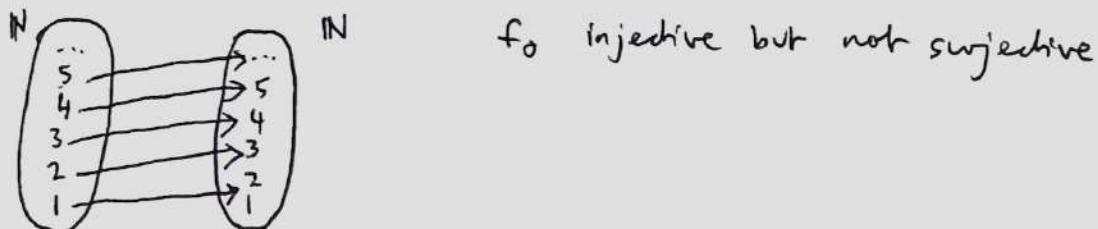
For A, B finite:

1. No surjection  $A \rightarrow B$  if  $|B| > |A|$
2. No injection  $A \rightarrow B$  if  $|A| > |B|$
3. For  $f: A \rightarrow A$ , f injective  $\Leftrightarrow$  f surjective

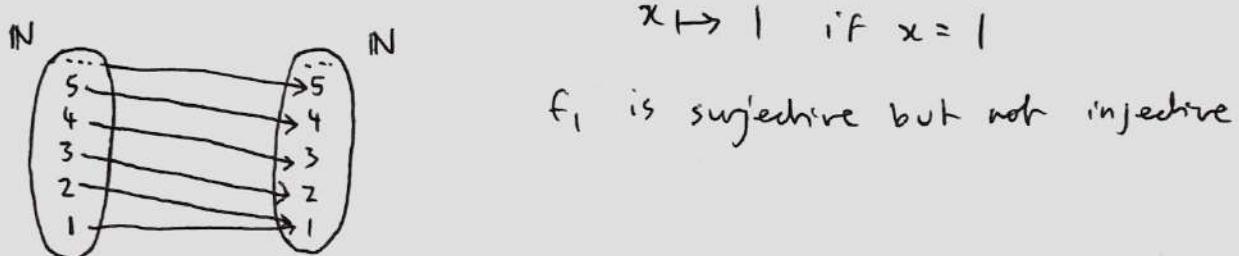
4. No bijection  $A \rightarrow B$  with  $B \subseteq A$ ,  $B \neq A$  (proper subset)

But for infinite sets:

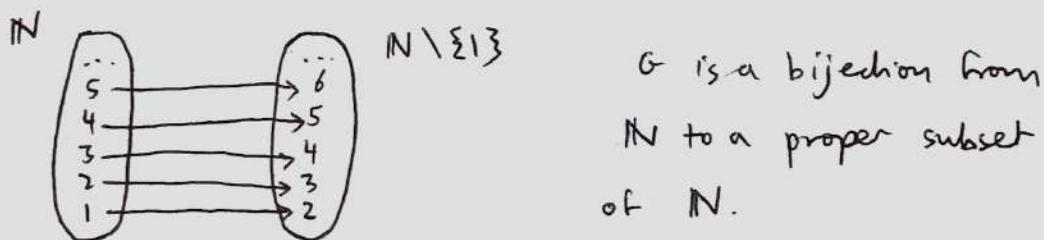
1. Define  $f_0: \mathbb{N} \rightarrow \mathbb{N}$  by  $x \mapsto x+1$



2. Define  $f_1: \mathbb{N} \rightarrow \mathbb{N}$  by  $x \mapsto x-1$  if  $x \neq 1$



3. Define  $g: \mathbb{N} \rightarrow \mathbb{N} \setminus \{1\}$ ,  $x \mapsto x+1$



More examples

1.  $I_x: X \rightarrow X$  with  $x \mapsto x$ , the identity function

2. For any set  $X$  with  $A \subseteq X$ , have indicator/characteristic function  $\chi_A: X \rightarrow \{0, 1\}$

with  $x \mapsto \begin{cases} 1 & \text{if } x \in A \\ 0 & \text{if } x \notin A \end{cases}$

3. A sequence  $x_1, x_2, \dots$  of reals is a function  $\mathbb{N} \rightarrow \mathbb{R}$ :  
 $n \mapsto x_n$

4. Operation  $+$  on  $\mathbb{N}$  is function  $\mathbb{N}^2 \rightarrow \mathbb{N}$  with  $(m, n) \mapsto m+n$

5. A set  $X$  has size  $n$  iff  $\exists$  a bijection  
 $\{1, 2, \dots, n\} \rightarrow X$  with  $i \mapsto a_i$  ( $X = \{a_1, a_2, \dots, a_n\}$ )

### Composition of Functions

Given  $f: A \rightarrow B$  and  $g: B \rightarrow C$ , the composition

$g \circ f : A \rightarrow C$  is defined by  $g \circ f(a) = g(f(a))$ ,  $a \in A$

e.g.  $f: \mathbb{R} \rightarrow \mathbb{R}$ ,  $g: \mathbb{R} \rightarrow \mathbb{R}$   
 $x \mapsto 2x$                      $x \mapsto x+1$

In general

$\circ$  is not  
commutative.

$$f \circ g(x) = 2(x+1)$$

$$g \circ f(x) = 2x+1$$

Function composition is associative.

Given  $f: A \rightarrow B$ ,  $g: B \rightarrow C$ ,  $h: C \rightarrow D$

$$h \circ (g \circ f) = (h \circ g) \circ f$$

Indeed for any  $x \in A$ :

$$(h \circ (g \circ f))(x) = h(g \circ f(x)) = h(g(f(x)))$$

$$((h \circ g) \circ f)(x) = (h \circ g)(f(x)) = h(g(f(x)))$$

The two compositions are equal  $\forall x$  so are the same function.

□

Say  $f: A \rightarrow B$  is invertible if  $\exists g: B \rightarrow A$  with  $gf = 1_A$  and  $fg = 1_B$

e.g.  $f: \mathbb{R} \rightarrow \mathbb{R}$ ,  $g: \mathbb{R} \rightarrow \mathbb{R}$

$$x \mapsto 2x+1 \quad x \mapsto \frac{x-1}{2}$$

$$(g \circ f)(x) = g(2x+1) = x \text{ so } g \circ f \text{ is } 1_{\mathbb{R}}$$

$$(f \circ g)(x) = f\left(\frac{x-1}{2}\right) = x \text{ so } f \circ g \text{ is } 1_{\mathbb{R}}$$

So  $f$  is invertible:  $g = f^{-1}$

In fact  $f: A \rightarrow B$  is invertible iff it's a bijection.

$\Rightarrow$  If  $g$  is inverse to  $f$ :  
 surjective:  $\forall b \in B, b = f(g(b))$   
 injective:  $\forall a, a' \in A : f(a) = f(a')$   
 $\Rightarrow g(f(a)) = g(f(a')) \Rightarrow a = a'$  ✓

$\Leftarrow$  Let  $g(b) =$  the unique  $a \in A$  with  $f(a) = b$ , for each  $b \in B$ . ✓

### Equivalence Relations

A relation on a set  $X$  is a subset  $R$  of  $X \times X$ .

Usually write  $a R b$  for  $(a, b) \in R$ .

- e.g. 1. on  $\mathbb{N}$ ,  $a R b$  if  $a \equiv b \pmod{5}$   
 2. on  $\mathbb{N}$ ,  $a R b$  if  $a | b$   
 3. on  $\mathbb{N}$ ,  $a R b$  if  $a \neq b$   
 4. on  $\mathbb{N}$ ,  $a R b$  if  $a = b \pm 1$   
 5. on  $\mathbb{N}$ ,  $a R b$  if either  $a, b \leq 6$  or  $a, b > 6$

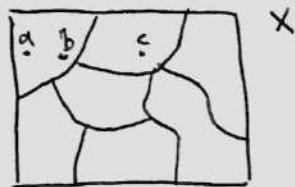
Some properties a relation might have:

1.  $R$  reflexive if  $\forall x \in X : x R x$ .
2.  $R$  symmetric if  $\forall x, y \in X : x R y \Rightarrow y R x$
3.  $R$  transitive if  $\forall x, y, z \in X : x R y, y R z \Rightarrow x R z$

Exercise: check examples - see which have what properties

An equivalence relation is reflexive, symmetric and transitive.

Example



Let  $X$  be a set, let  $\{C_i : i \in I\}$  be a partition of  $X$ :  
each  $C_i$  nonempty, disjoint,  $\bigcup_{i \in I} C_i = X$

Then have  $a R b$  if  $\exists i$  s.t  $a, b \in C_i$  is an equivalence relation.

All equivalence relations are of this form.

For an ER on  $X$ ,  $x \in X$ , the equivalence class of  $x$  is  
 $[x] = \{y \in X : y R x\}$ .

In Example 1 ( $xRy$  if  $x \equiv y \pmod{5}$ )

$\equiv 0(5)$
$\equiv 1(5)$
$\equiv 2(5)$
$\equiv 3(5)$
$\equiv 4(5)$

$\mathbb{N}$

Proposition 5 let  $R$  be an equivalence relation on a set  $X$ . Then the equivalence classes partition  $X$ .

Proof Equivalence classes are all nonempty:  $xRx$ .

$$\bigcup_{x \in X} [x] = X \quad (x \in [x] \quad \forall x \in X)$$

↓  
lowercase  $x$

So need to show ECs are disjoint (or are the same).

Given  $x, y$  with  $[x] \cap [y] \neq \emptyset$ , need  $[x] = [y]$ .

Have  $z \in [x] \cap [y]$ , some  $z$ .

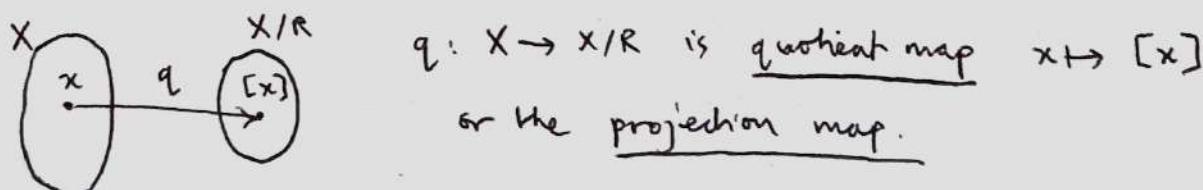
So  $zRx, zRy \Rightarrow xRy$ .  $\forall t, tRx \Rightarrow tRy$  (transitivity)  
 so  $[x] = [y]$ .  $tRy \Rightarrow tRx$   $\square$

Is there an equivalence relation on  $\mathbb{N}$  with 3 ECs: 2 infinite, 1 finite?

Should think of this in terms of partitions then it becomes easy.

Given an ER on set  $X$ , the quotient of  $X$  by  $R$  is

$X/R = \{[x] : x \in X\}$  "the set of countries"



Another example: on  $\mathbb{Z} \times \mathbb{N}$ , define  $(a, b) R (c, d)$  if  $ad = bc$ .

e.g.  $[(1, 2)] = \{(1, 2), (2, 4), (3, 6), \dots\}$

can think of as a copy of  $\mathbb{Q}$  ~~don't do this~~  $\mathbb{Z} \times \mathbb{N} / R$

## Chapter 4 - Countability

Looking at "sizes" of infinite sets e.g.  $\mathbb{N}$  "looks smaller" than  $\mathbb{Z}$

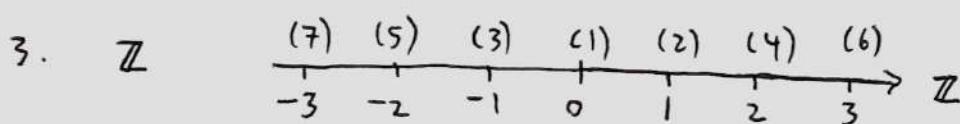
Say a set  $X$  is countable if  $X$  is finite or bijects with  $\mathbb{N}$ .

i.e. "you can list the elements" as  $a_1, a_2, a_3, \dots$  (may terminate).

### Examples

1. Any finite set

2.  $\mathbb{N}$



Can list as  $0, 1, -1, 2, -2, 3, -3, \dots, n, -n, \dots$   
so is countable

Write  $a_n = \frac{n}{2}$  if  $n$  even,  $-\frac{(n-1)}{2}$  if  $n$  odd

So in a sense " $\mathbb{Z}$  is the same size as  $\mathbb{N}$ "

Are all sets countable?

Proposition 2 A set  $X$  is countable iff  $\exists$  an injection  $f: X \rightarrow \mathbb{N}$ .

Proof If  $X$  finite then  $X$  injects into  $\mathbb{N}$

If  $X$  bijects then it injects into  $\mathbb{N}$ .  $\Rightarrow \checkmark$

$\Leftarrow X$  finite  $\Rightarrow X$  countable

Suppose  $X$  is infinite.

Have  $X$  bijects with its image  $f(X)$  under  $f$   
so enough to show that  $f(X)$  is countable.

But in  $\mathbb{N}$   $\exists$  a smallest element, so can write down  $f(X)$  in increasing order so is countable.  $\Leftarrow \checkmark \quad \square$

Warning In  $\mathbb{R}$  let  $X = \left\{ \frac{1}{2}, \frac{2}{3}, \frac{3}{4}, \dots \right\} \cup \{1\}$ ,  $X$  is countable.

But if we counted by "least element" would write down  $\frac{1}{2}, \frac{2}{3}, \frac{3}{4}, \frac{4}{5}, \dots$  and never reach 1.

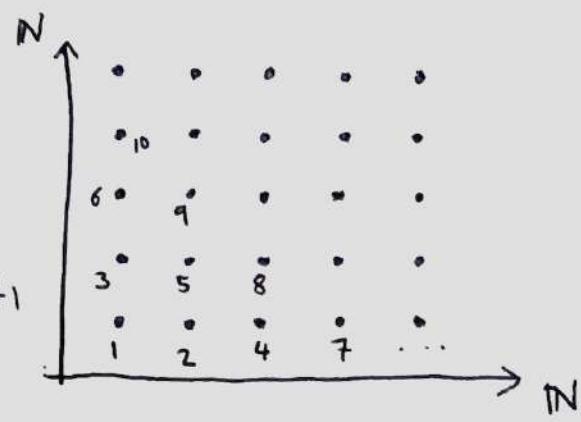
So have to take care e.g. when proving Prop 2.

Theorem 2  $\mathbb{N} \times \mathbb{N}$  is countable.

Proof 1 Define  $a_1 = (1, 1)$

and  $a_n$  inductively by writing

$$a_{n+1} = (p, q) : a_n = \begin{cases} (p-1, q+1) & p \neq 1 \\ (1, p+q) & \text{if } p = 1 \end{cases}$$



This does hit every point  $(x, y) \in \mathbb{N} \times \mathbb{N}$  (e.g. induction on  $x+y$ ) so have listed  $\mathbb{N} \times \mathbb{N}$ .  $\square$

Proof 2 Define  $f: \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$   $(x, y) \mapsto 2^x 3^y$  is injective.  $\square$

Theorem 2' Let  $A_1, A_2, A_3, \dots$  be countable sets.

Then  $A_1 \cup A_2 \cup A_3 \cup \dots$  is countable.

Proof For each  $i$  have  $A_i$  countable, so can list  $A_i$  as  $a_{i1}, a_{i2}, a_{i3}, \dots$  (may terminate)

Define  $f: \bigcup_{n \in \mathbb{N}} A_n \rightarrow \mathbb{N}$

$x \mapsto 2^x 3^y$  where  $x = a_{ij}$ , least such  $i$  (as  $x$  could be in more than one set  $A_i$ ).

This is injective.  $\square$

### Examples

1.  $\mathbb{Q}$  is countable :  $\mathbb{Q} = \mathbb{Z} \cup \frac{1}{2}\mathbb{Z} \cup \frac{1}{3}\mathbb{Z} \cup \dots$

2) Set A (all algebraic numbers) is countable.

Enough to show that the set of all integer polynomials is countable as then A is a countable union of finite sets.

Thus enough to show that for each d, the set of all integer polynomials of degree d is countable.

But this set injects into  $\mathbb{Z}^{d+1}$  ( $a_d x^d + \dots + a_1 x + a_0 \mapsto (a_d, \dots, a_1, a_0)$ ) and  $\mathbb{Z}^n$  is countable  $\forall n$

( $\mathbb{Z} \times \mathbb{Z}$  countable so  $\mathbb{Z} \times \mathbb{Z} \times \dots \times \mathbb{Z}$  countable by induction).

□

A set is ~~or~~ uncountable if it's not countable.

Theorem 3  $\mathbb{R}$  is uncountable.

Proof We'll show  $(0, 1)$  is uncountable.

Given  $r_1, r_2, \dots$  in  $(0, 1)$  and  $s \in (0, 1)$  not on list.

For each  $r_n$ , have decimal expansion  $r_n = 0.r_{n1}r_{n2}r_{n3}\dots$

$$r_1 = 0.r_{11}r_{12}r_{13}\dots$$

$$r_2 = 0.r_{21}r_{22}r_{23}\dots$$

$$r_3 = 0.r_{31}r_{32}r_{33}\dots$$

Define  $s = 0.s_1s_2s_3\dots$  by e.g.

$$s_1 = \begin{cases} 5 & \text{if } r_{11} \neq 5 \\ 6 & \text{if } r_{11} = 5 \end{cases} \quad \text{ensures } s \neq r_1$$

$$s_2 = \begin{cases} 5 & \text{if } r_{22} \neq 5 \\ 6 & \text{if } r_{22} = 5 \end{cases} \quad \text{etc}, \quad s_n = \begin{cases} 5 & \text{if } r_{nn} \neq 5 \\ 6 & \text{if } r_{nn} = 5 \end{cases}$$

Then S can't be on the list.  $s \neq r_n$ : differs in decimal digit n.

□

Remarks

1. This is called a diagonal argument  
(Cantor's diagonal argument)
2.  $\mathbb{R}$  is uncountable and algebraic numbers are countable, so there exists a transcendental number (i.e.  $\mathbb{A} \neq \mathbb{R}$ ).

Indeed, "most" numbers are transcendental  
(i.e.  $\mathbb{R} \setminus \mathbb{A}$  is uncountable).

because if  $\mathbb{R} \setminus \mathbb{A}$  countable, then  $\mathbb{R} = \mathbb{A} \cup (\mathbb{R} \setminus \mathbb{A})$   
would be countable. □

Another uncountable set

Theorem 4 The power set of the naturals,  $\mathbb{P}(\mathbb{N})$  is uncountable.

Proof Suppose  $\mathbb{P}(\mathbb{N})$  is listed as  $S_1, S_2, S_3, \dots$ . Let  $S = \{n \in \mathbb{N} : n \notin S_n\}$

To ensure  $S \neq S_1$ : take  $1 \in S$  if  $1 \notin S_1$ ,  
 $1 \notin S$  if  $1 \in S_1$ .

To ensure  $S \neq S_2$ : take  $2 \in S$  if  $2 \notin S_2$ ,  $2 \notin S$  if  
 $2 \in S_2$  ...

(can continue  $\forall S_n : n \in \mathbb{N}$ : construct  $S$  not equal to any of the  $S_n$ ).  $\square$

### Remarks

1. This is also a diagonal argument.
2. Alternatively: just inject  $(0, 1)$  into  $\mathbb{P}(\mathbb{N})$ :  
 Given  $x \in (0, 1)$ , write  $x$  in binary as  
 $0.x_1x_2x_3\dots$  (not ending with all 1s) and put  
 $f(x) = \{n : x_n = 1\}.$

In fact our proof of Theorem 4 shows:

Theorem 5 For any set  $X$ , there is no bijection from the set  $X$  to the power set of  $X$ . (In fact: no surjection)

Proof Given any function  $f: X \rightarrow \mathbb{P}(X)$  we show  $f$  is not surjective.

Let  $S = \{x \in X : x \notin f(x)\}$

Then  $S$  does not belong to the image of  $f$  since  $\forall x$  have  $S \neq f(x)$  as they differ at element  $x$ .

1. Similar to Russell's paradox
2. Gives another proof that there is no universal set: if it is, then the power set is contained in it so there'd be a surjection  $\times$

Example on countability:  $\overbrace{(-)(-)(-)(-)}^{\text{reverget to this}} \rightarrow \mathbb{R}$

let  $A_i : i \in I$  be a family of pairwise disjoint open intervals.

Warnings: can't say "next interval"  $\overbrace{(-)(-)(-)(-)(-)(-)\dots}^{\text{as dense in } \mathbb{R}} \rightarrow \mathbb{R}$

Answer Yes it is countable.

Proof 1 Each  $A_i$  contains a rational and the rationals are countable, so the family is countable.  $\checkmark$   $\square$

Proof 2  $\{i \in I : A_i \text{ has length } > 1\}$  countable - injects into  $\mathbb{Z}$   
 $\{i \in I : A_i \text{ has length } > \frac{1}{2}\}$  again countable - injects into  $\mathbb{Z}$

$\forall n : \{i \in I : A_i \text{ has length } > \frac{1}{n}\}$  again countable

Here we've written them all as a countable union of countable sets, so  $\{A_i\}$  is countable.  $\square$

To show a set  $X$  is uncountable we can do one of

1. Run diagonal argument
2. Inject uncountable set into  $X$

To show  $X$  is countable we can

1. List it (usually fiddly)
2. Inject it into  $\mathbb{N}$
3. Use "countable union of countable sets is countable" usually best
4. If in or near  $\mathbb{R}$ , maybe consider  $\mathbb{Q}$ .

Intuitively, think of " $A$  bijects with  $B$ " as saying that  $A$  and  $B$  "have the same size".

" $A$  injects into  $B$ " as saying that " $A$  is at most as large as  $B$ "  
 " $A$  surjects to  $B$ " as saying that " $A$  is at least as large as  $B$ "  
 ( $B$  nonempty)

For these to make sense, we'd want that (for  $A, B \neq \emptyset$ ):

There exists an injection  $f: A \rightarrow B$  iff  $\exists$  a surjection  $g: B \rightarrow A$   
 $\Rightarrow$  Fix  $a_0 \in A$ . Define  $g: B \rightarrow A$  by

$$b \mapsto \begin{cases} \text{the unique } a \in A \text{ with } f(a) = b & \text{if it exists} \\ a_0 & \text{if not} \end{cases}$$

Then  $g$  is surjective.

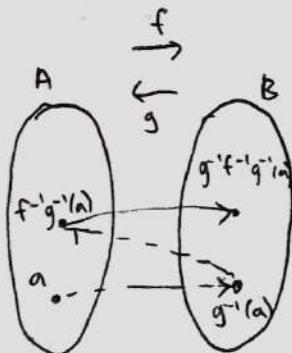
$\Leftarrow$  For each  $a \in A$ , have some  $a' \in B$  with  $g(a') = a$ .  
 (As  $g$  is surjective)

Let  $f(a) = a'$  for each  $a \in A$ . Then  $f$  is injective.

□

We also need "if  $A$  is at most as large as  $B$ " and  
 if "B is at most as large as A" then  $A$  and  $B$  "have the same size".

Theorem 6 (Schröder - Bernstein Theorem) If  $f: A \rightarrow B$  and  
 $g: B \rightarrow A$  are injections, then  $\exists$  a bijection  $h: A \rightarrow B$ .



Proof For  $a \in A$ , write  $g^{-1}(a)$  for the  $b \in B$  (if it exists) such that  $g(b) = a$ .

Similarly for  $f^{-1}(b)$  where  $b \in B$ .

The ancestor sequence of  $a \in A$  is  
 $g'(a), f^{-1}(g'(a)), f g^{-1}(f^{-1}(g'(a)))$  etc.  
 (May terminate)

Similarly for  $b \in B$

Let  $A_0$  be  $\{a \in A : \text{ancestor sequence stops at even time}\}$

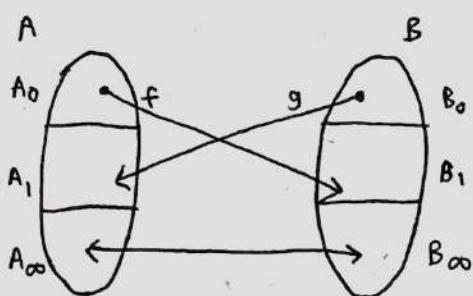
"stops in A"  
↑ includes  $t=0$

Let  $A_1$  be  $\{a \in A : \text{ancestor sequence stops at odd time}\}$

↓ "stops in B"

Let  $A_\infty$  be  $\{a \in A : \text{ancestor sequence does not stop}\}$

Similarly for  $B_0, B_1, B_\infty$ .



Then  $f$  bijects  $A_0$  with  $B_1$ ,

If you start in  $A_0$ , then applying  $f$  puts you in  $B_1$  (even  $\rightarrow$  odd)

So every  $b \in B$  is  $f(a)$  for some  $a \in A_0$ .

Similarly,  $g$  bijects  $B_0$  with  $A_1$ .

And  $f$  or  $g$  bijects  $A_\infty$  with  $B_\infty$ . So the function

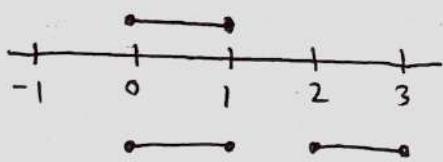
$h: A \rightarrow B$  with

$$a \mapsto \begin{cases} f(a) & \text{if } a \in A_0 \\ g^{-1}(a) & \text{if } a \in A_1 \\ f(a) & \text{if } a \in A_\infty \end{cases}$$

is a bijection.

□

Example Do  $[0,1]$  and  $[0,1] \cup [2,3]$  biject?



Have injection

$$f: [0,1] \rightarrow [0,1] \cup [2,3]$$

$$\text{e.g. } f(x) = x$$

$$\text{and } g: [0,1] \cup [2,3] \rightarrow [0,1]: \text{ e.g. } g(x) = \frac{x}{3}$$

an injection

so by Schröder-Bernstein they biject.

□

Would also be nice to have that for any sets  $A, B$ : either  $A$  injects into  $B$  or  $B$  injects into  $A$ .

This is true, but harder - see Part II "Set Theory and Logic"

Have  $\mathbb{N}, \mathbb{P}(\mathbb{N}), \mathbb{P}\mathbb{P}(\mathbb{N}), \dots$

Does every  $X$  inject into one of these sets?

No: e.g.  $X = \mathbb{N} \cup \mathbb{P}(\mathbb{N}) \cup \mathbb{P}\mathbb{P}(\mathbb{N}) \cup \dots$  is a countable union

Then  $X' = X \cup \mathbb{P}(X) \cup \mathbb{P}\mathbb{P}(X) \cup \dots$

Then  $X'' = X' \cup \mathbb{P}(X') \cup \mathbb{P}\mathbb{P}(X') \cup \dots$

and so on

Then  $Y = X \cup X' \cup X'' \cup X''' \cup \dots$

"and can keep going".

### What happens next

Factorisation  $\rightarrow$  IB Groups, Rings and Modules

Fermat,  $-1$  square mod  $p$  etc.  $\rightarrow$  II Number Theory  
(see quadratic reciprocity)

Analysis  $\rightarrow$  IA Analysis 1

Countability etc  $\rightarrow$  II Set Theory and Logic



### Generators

Working in  $\mathbb{Z}_p^*$  (nonzero elements of  $\mathbb{Z}_p$ ) under multiplication.

The order of  $x \in \mathbb{Z}_p^*$  is the least  $n \in \mathbb{N}$  with  $x^n = 1$ .  
 (Exists as  $x^{p-1} = 1$ ).

e.g. in  $\mathbb{Z}_7^*$ :  $2^1 = 2, 2^2 = 4, 2^3 = 1$   $2$  has  
order  $3$

$$3^1 = 3, 3^2 = 2, 3^3 = 6, 3^4 = 4, 3^5 = 5, \\ 3^6 = 1 \text{ so } 3 \text{ has order } 6.$$

Say  $x$  is a generator or primitive root mod  $p$  if it has order  $p-1$ .

Aim  $\exists$  generator  $g$  with  $\mathbb{Z}_p^* = \{g, g^2, g^3, \dots, g^{p-1} = 1\}$   
 i.e.  $\mathbb{Z}_p^*$  is cyclic (wrt  $x$ ).

If  $x$  has order  $n$ , then  $x^{2n} = 1, x^{3n} = 1, \dots$

Proposition 1 If  $x$  has order  $n$  and  $x^d = 1$ , then  $n$  divides  $d$ .

Proof Suppose not: write  $d = qn + r$  where  $q > 0, 0 < r < n$   
 Then  $x^{qn+r} = 1, x^n = 1 \Rightarrow x^r = 1 \quad \text{**}$  □

Proposition 2 Let  $x$  have order  $a$  and  $y$  have order  $b$ . If  $a$  and  $b$  are coprime, then  $xy$  has order  $ab$ .

Proof Have  $(xy)^{ab} = x^{ab}y^{ab} = 1$ .

Given  $(xy)^d = 1$ , want  $ab | d$ :

Have  $x^d y^d = 1$  (Raise to power  $b$  to kill  $y$ )

So  $x^{db} = 1 \Rightarrow a | db$  but  $a, b$  coprime so  $a | d$ .

Similarly,  $y^{da} = 1 \Rightarrow b | da \Rightarrow b | d$ .

$a, b$  coprime so  $\underline{ab | d}$  as required.

So  $xy$  has order  $ab$ . □

What if  $a, b$  are not coprime?

Proposition 3 Let  $x$  have order  $a$  and  $y$  have order  $b$ .  
Then  $\exists$  an element of order  $\text{lcm}(a, b)$ .

Example first: Suppose  $\text{ord}(x) = 8$ ,  $\text{ord}(y) = 10$ .

Then  $y^2$  has order 5, so  $xy^2$  has order 40  
(prop. 2).

Proof Write  $a = p_1^{a_1} \dots p_k^{a_k} q_1^{c_1} \dots q_j^{c_j}$   
 $b = p_1^{b_1} \dots p_k^{b_k} q_1^{d_1} \dots q_j^{d_j}$

where  $p_1 \dots p_k, q_1 \dots q_j$  distinct primes,  $a_i > b_i > 0 \ \forall i$ ,  
 $d_i > c_i > 0 \ \forall i$ .

Then  $x^{q_1^{c_1} \dots q_j^{c_j}}$  has order  $p_1^{a_1} \dots p_k^{a_k}$   
 $y^{p_1^{b_1} \dots p_k^{b_k}}$  has order  $q_1^{d_1} \dots q_j^{d_j}$

so their product has order  $p_1^{a_1} \dots p_k^{a_k} q_1^{d_1} \dots q_j^{d_j}$   
 $= \underline{\text{lcm}(a, b)}$ . □

Corollary 4 Let greatest order of all elements of  $\mathbb{Z}_p^*$  be  $d$ .  
Then  $\forall x \in \mathbb{Z}_p^*$ ,  $\text{ord}(x)$  divides  $d$ .

Proof If not, then applying Prop. 3 would give an element of  $\text{lcm}(\text{ord}(x), d) > d$ . ✖ □

Remark The above all holds in any finite abelian group.

Proposition 5 In  $\mathbb{Z}_p$  ( $p$  prime), a polynomial of degree  $k$  has at most  $k$  roots.  
coeff. of  $x^k$  is nonzero in  $\mathbb{Z}_p$ .

Proof If polynomial  $f$  has no roots ✓

If  $\exists$  a root  $a \in \mathbb{Z}_p$  then write  $f(x) = (x-a)g(x)$   
where  $g$  has degree  $\leq k-1$ . Now  $g$  has  $\leq k-1$  roots  
(induction) and roots of  $(x-a)g(x)$  are  $a$  and roots of  $g$   
( $p$  prime). □

Theorem 6  $\mathbb{Z}_p^*$  ( $p$  prime) has a generator.

Proof Let  $d$  be the greatest order of all elements of  $\mathbb{Z}_p^*$ :  
want to show that this is  $p-1$ .

$$x^d = 1 \quad \forall x \in \mathbb{Z}_p^* \text{ (corollary 4)}$$

Thus the polynomial  $x^d - 1$  has  $p-1$  roots, so  $d \geq p-1$ .  $\square$

Some things that are easy once we know  $\exists$  a generator  $g$ :

- 1) Fermat's little Theorem: for any  $x \in \mathbb{Z}_p^*$ , have  $x = g^n$  for some  $n$  (def. of generator) so  $x^{p-1} = (g^n)^{p-1} = (g^{p-1})^n = 1^n = 1$ .
- 2) From now on,  $p > 2$ .  
The squares in  $\mathbb{Z}_p^*$  are  $g^2, g^4, g^6, g^8, \dots, g^{\frac{p-1}{2}} = 1$   
 $\downarrow$   
 $\begin{cases} p \text{ odd} \\ p-1 \text{ even} \end{cases}$   
Hence square  $\times$  square = square  
non-square  $\times$  square = non-square  
non-square  $\times$  non-square = square by parity
- 3)  $-1$  a square: have  $g^{\frac{p-1}{2}} = -1$  as  $g^{p-1} = 1$ .  
Thus  $-1$  is a square iff  $\frac{p-1}{2}$  is even.
- 4)  $1 \times 2 \times 3 \times \dots \times (p-1) = g^1 g^2 g^3 \dots g^{\frac{p-1}{2}} = g^{\frac{p(p-1)}{2}}$   
 $= (g^{\frac{p-1}{2}})^p = (-1)^p$  (Wilson's theorem)

- 5) Cubes: cubes are  $g^3, g^6, g^9, \dots$

If  $p = 3k+1$  then cubes are  $g^3, g^6, g^9, \dots, g^{p-1}$

so exactly  $\frac{1}{3}$  of all elements of  $\mathbb{Z}_p^*$  are cubes

If  $p = 3k+2$ : every element is of the form  $g^{3n}$  so is a cube.  
(cyclic group, order  $3k+1$ )

Remarks proved  $\exists$  generator, not how to find it. (Not even understood when 2 is a generator); not known if 2 is a generator mod  $p$  for infinitely many  $p$  (Artin's conjecture)