

Finite Samplings on $[0,1]^n$

Task: Find a set of points $\{x_k\}_{k=1}^P$ in $[0,1]^n$ satisfying

- ① $\{x_k\}_{k=1}^P$ are uniformly distributed.
- ② $\{x_k\}_{k=1}^r$ are uniformly distributed for all $r < P$.
- ③ The process is repeatable.
- ④ The process is independent of computing platform.

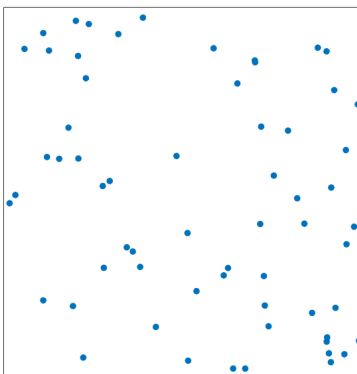
Motivation :

- I have a total budget of P function evaluations and sampling these points will be my entire algorithm.
- my algorithm needs P points to initiate (e.g. Nelder mead) and it would be good to start with a good sampling.
- Every so often, I would like to check P points in my domain to help me escape local minimizer basins.
- I need non-stochastic methods for repeatability and testing purposes.
- If I try for P points, but only get $r < P$ points, I want the sampling to still be uniform (useful).

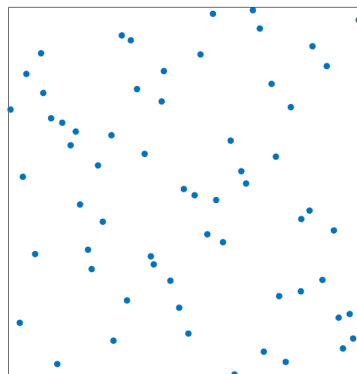
There are many strategies and research is ongoing. We will consider only three options - not all of which satisfy our desired Properties.

(example : $n=2$, $P=24$)

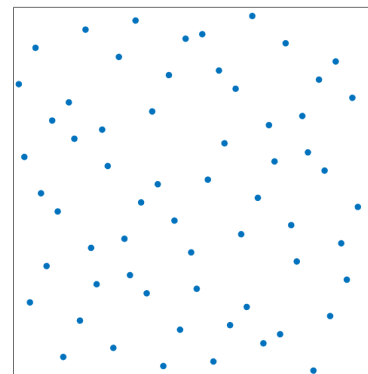
Random Generator



Latin Hypercube



Halton Sequence



Random Generator

Given : n, P, S

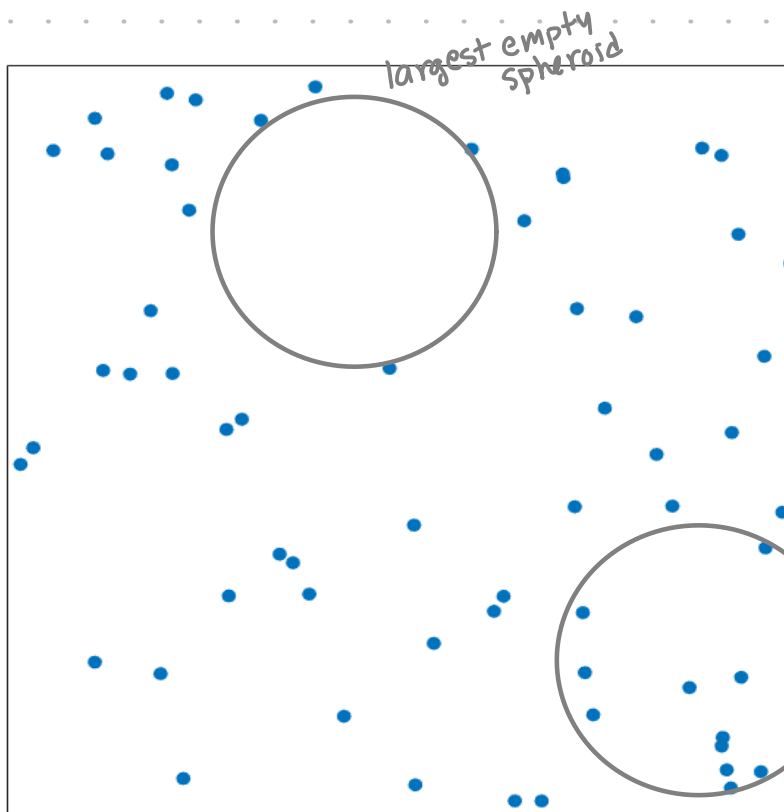
Set random number generator seed to S .

Request $n \times P$ random values uniformly drawn from $(0,1)$.

Arrange the values (by predetermined order) into $n \times P$ array A .

The ordered columns of A represent $\{x_k\}_{k=1}^P \in (0,1)$.

- Repeatable through seed S
- May not be platform independent
- Uniform in large P limit



62 points on $[0,1]^2$

- distributed across domain
- not visually uniform with large empty spaces and point clusters.

Same size spheroid
but containing
12 points.

Latin Hypercube Sampling

Given : n, p, s

Set random number generator seed to s .

Create n random permutations of $[1, 2, 3, \dots, p]$.

Arrange these row vectors, by predetermined order, into $n \times p$ array B .

Create (by previous algorithm) $n \times p$ array R .

Compute $A = \frac{1}{p}(B - R)$.

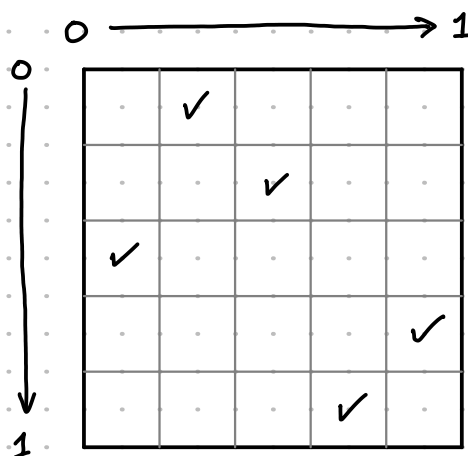
The ordered columns of A represent $\{x_k\}_{k=1}^p \in (0, 1)$.

This algorithm partitions $[0, 1]^n$ into a $p \times p \times p \times \dots \times p$ hyper Rubik's cube and guarantees that each column contains one point, each row contains one point, each hyperrow contains one point. Then in each small subcube, a random point is chosen.

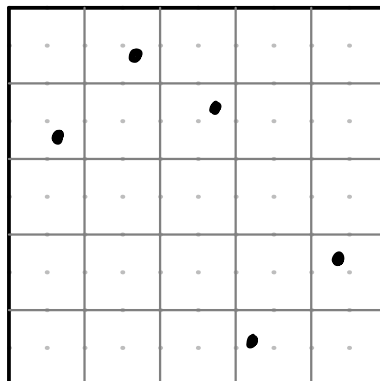
Example : ($n=2, p=5$) Five points on unit square

$$B = \begin{bmatrix} 2 & 3 & 5 & 1 & 4 \\ 3 & 1 & 4 & 2 & 5 \end{bmatrix}$$

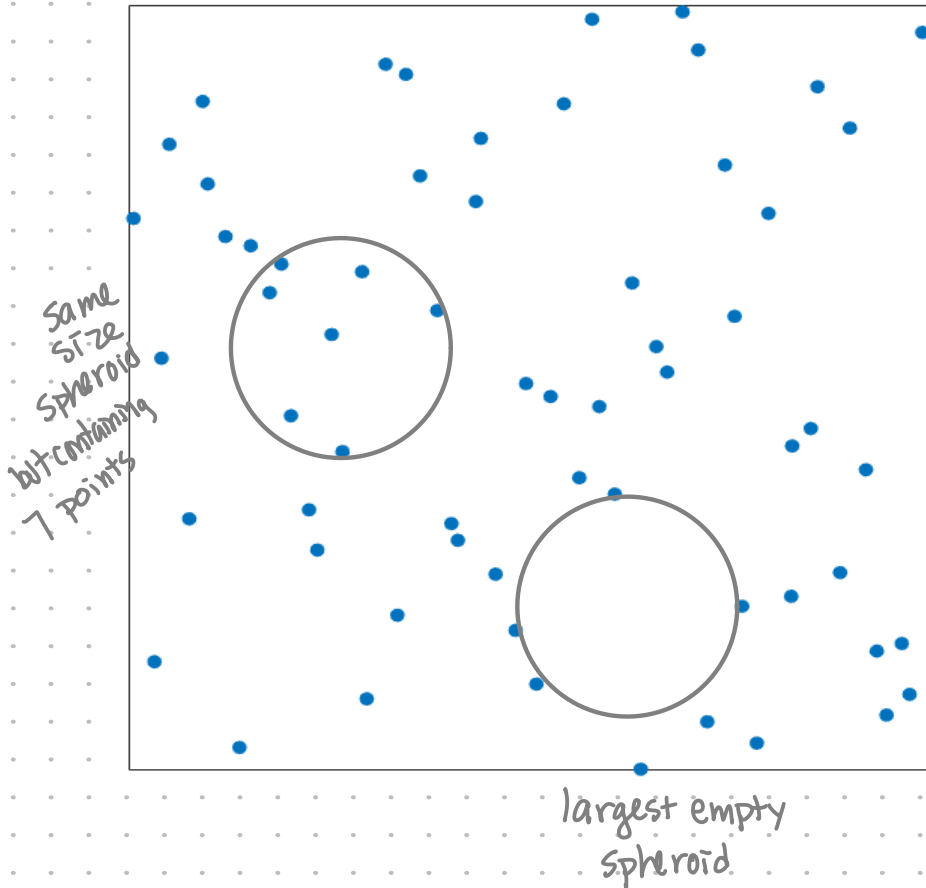
represents :



The columns of $A = \frac{B-R}{p}$
are then the points :
(for some random values R)



- repeatable through seeds
- may not be platform independent
- better uniformity enforced by B
- Samples individual coordinates uniformly.
- any coordinate hyperplane projection exhibits the same uniformity.



62 points on $[0,1]^2$

- distributed across domain
- better visual uniformity with more diffuse clusters
- smaller largest empty spheroid

Haltom Sequence

```
q=primes(10*n);  
q=q(1:n); ← first n prime numbers  
x=zeros(n,p); ← A matrix initialization  
for k=1:n  
    b=q(k);  
    idx=1:p;  
    f=ones(1,p);  
    r=zeros(1,p);  
    while any(idx>0)  
        f=f/b;  
        r=r+f.*mod(idx,b);  
        idx=floor(idx/b);  
    end  
    x(k,:)=r; ← A matrix of columns  $x_k$   
end
```

A Haltom Sequence is defined on $(0,1)$ with integer seed $q \geq 2$.

The Haltom vector sequence shown here is composed of n Haltom sequences each with a distinct prime number seed.

Haltom Sequence for seed $q=2$:

$$\frac{1}{2}, \frac{1}{4}, \frac{3}{4}, \frac{1}{8}, \frac{5}{8}, \frac{3}{8}, \frac{7}{8}, \frac{1}{16}, \frac{9}{16}, \dots$$

Haltom Sequence for seed $q=3$:

$$\frac{1}{3}, \frac{2}{3}, \frac{1}{9}, \frac{4}{9}, \frac{7}{9}, \frac{2}{9}, \frac{5}{9}, \frac{8}{9}, \frac{1}{27}, \dots$$

⇒ Haltom Vector Sequence for $n=2$ and seeds 2,3:

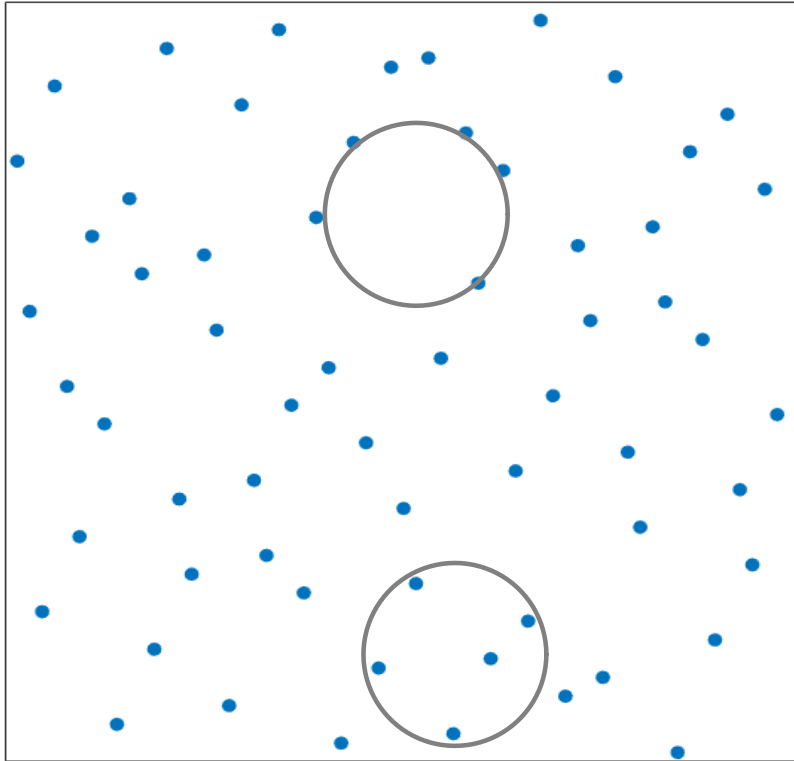
$$\begin{pmatrix} 1/2 \\ 1/3 \end{pmatrix}, \begin{pmatrix} 1/4 \\ 2/3 \end{pmatrix}, \begin{pmatrix} 3/4 \\ 1/9 \end{pmatrix}, \begin{pmatrix} 1/8 \\ 4/9 \end{pmatrix}, \begin{pmatrix} 5/8 \\ 7/9 \end{pmatrix}, \dots$$

Interesting method for finding a sequence with seed q (example: $q=2$)

index k	k base q	reflect bits	decimal
1	1	.1	$1/2$
2	10	.01	$1/4$
3	11	.11	$3/4$
4	100	.001	$1/8$
5	101	.010	$5/8$
6	110	.011	$3/8$
7	111	.111	$7/8$

it's the
Haltom
sequence!

- no random number generation
- platform independent
- maintains good uniformity at all r .



62 points on $[0, 1]^2$

- distributed across domain
- excellent visual uniformity
- exhibits correlation structure if $q_1 \approx q_2 \gg 1$.