EVALUATION SEMANTICS OF COERCION

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An idea for a *new* evaluation dynamics of coercion inspired by the denotational semantics of cubical type theory. The main idea is that coercions should *not* be evaluated in an ordinary environment, but instead an environment where each cell is a line.

$$\llbracket \Gamma \vdash A \rrbracket : \llbracket \Gamma \rrbracket \to \mathsf{Type}$$

$$\llbracket \Gamma \vdash a : A \rrbracket : \prod_{\gamma : \llbracket \Gamma \rrbracket} \llbracket \Gamma \vdash A \rrbracket_{\gamma}$$

$$\mathsf{coe} \llbracket \Gamma \vdash A \rrbracket : \prod_{\gamma : \llbracket \Gamma \rrbracket} \llbracket \Gamma \vdash A \rrbracket_{\gamma}$$

$$\mathsf{coe} \llbracket \Gamma \vdash A \rrbracket : \prod_{\gamma : \llbracket \Gamma \rrbracket} \prod_{r,s:\mathbb{I}} \prod_{a : \llbracket \Gamma \vdash A \rrbracket_{\gamma(r)}} \llbracket \Gamma \vdash A \rrbracket_{\gamma(s)}$$

$$\llbracket \Gamma \vdash \mathsf{coe}_F^{r \to s} M : F(r) \rrbracket \gamma = \mathsf{coe} \llbracket \Gamma . \mathbb{I} \vdash F(i) \rrbracket_{\langle \mathsf{refl}(\gamma), \mathsf{id}_{1} \rangle}^{\lVert \Gamma \vdash r : \mathbb{I} \rrbracket_{\gamma} \to \llbracket \Gamma \vdash s : \mathbb{I} \rrbracket_{\gamma}} \llbracket \Gamma \vdash M : F(r) \rrbracket_{\gamma}$$

$$\mathsf{coe} \llbracket \Gamma \vdash \Pi(A, B) \rrbracket_{\gamma}^{r \to s} f = \lambda a. \mathsf{coe} \llbracket \Gamma . A \vdash B \rrbracket_{\langle \gamma, \lambda i. \mathsf{coe} \llbracket \Gamma \vdash A \rrbracket_{\gamma}^{s \to i} a \rangle} \mathsf{coe} \llbracket \Gamma \vdash A \rrbracket_{\gamma}^{s \to r} a$$

$$\mathsf{coe} \llbracket \Gamma \vdash \Sigma(A, B) \rrbracket_{\gamma}^{r \to s} p = (\mathsf{coe} \llbracket \Gamma \vdash A \rrbracket_{\gamma}^{r \to s} \pi_{1}(p), \mathsf{coe} \llbracket \Gamma . A \vdash B \rrbracket_{\langle \gamma, \lambda i. \mathsf{coe} \llbracket \Gamma \vdash A \rrbracket_{\gamma}^{r \to i} \pi_{1}(p) \rangle}^{r \to s} \pi_{2}(p)$$

$$\mathsf{coe} \llbracket \Gamma, i = 0 \vdash A \rrbracket_{r}^{r \to s} \rangle_{(\mathcal{V})}$$

$$\mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{V}(i)} = \mathsf{if} \forall i : \mathbb{I}, \llbracket i \rrbracket_{\mathcal{$$

$$\mathbf{coe} \llbracket \Gamma \vdash \mathsf{V}_{i}(A,B,E) \rrbracket_{\gamma}^{r \to s}(v) = \begin{cases} \mathbf{coe} \llbracket \Gamma, i = 0 \vdash A \rrbracket_{\langle \gamma, \star \rangle}^{r \to s}(v) & \text{if } \forall j : \mathbb{I}, \llbracket i \rrbracket_{\gamma(j)} = 0 \\ \mathbf{coe} \llbracket \Gamma \vdash B \rrbracket_{\gamma}^{r \to s}(v) & \text{if } \forall j : \mathbb{I}, \llbracket i \rrbracket_{\gamma(j)} = 1 \\ ? & \text{if } \forall j : \mathbb{I}, \llbracket i \rrbracket_{\gamma(j)} = j \end{cases} \\ \mathbf{Vin}_{\llbracket i \rrbracket_{\gamma(s)}}(\mathbf{coe} \llbracket \Gamma.i = 0 \vdash A \rrbracket_{\langle \gamma, \star \rangle}^{r \to s} v, \mathbf{com} \llbracket \Gamma \vdash B \rrbracket_{\gamma}^{r \to s} [\partial(\llbracket i \rrbracket_{\gamma(s)})](\ldots)) & \text{otherwise} \end{cases}$$

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