

## LAB 3: Principles of Reliable Data Transfer

### Instructions:

- This is a **team** assignment that can be done in teams of at-most 2 persons. Sharing of code or reports across teams or copying code from the Internet or other sources will be considered plagiarism and will incur zero marks for the entire LAB component in the course and reporting of the cases to SSAC.
- This is a **take-home** assignment. The final report and code for all the questions have to be submitted on Google classrooms.
- **Format for final submission:**
  1. Only one person in the team is required to submit the solution. The other team member may submit a comment stating the name of their team mate who has submitted on their behalf. The report must contain the names of all members.
  2. You need to upload into the Google form:
    - A typeset report in pdf format. The file should be named **<name>\_Lab3\_report.pdf**, where name is of the person submitting the solutions as it appears on google classroom. The quality of the report carries significant weightage. Your report should be complete, correct, clear and concise.
    - A Python file named **<name>\_Protocol\_rdt22.py** containing your implementation of the rdt2.2 protocol as asked in question 2. The code should contain comments where appropriate to explain your logic.
    - A Python file named **<name>\_Protocol\_rdt3.py** containing your implementation of the rdt3.0 protocol as asked in question 3.
    - Do **not** include any other files in your submission (such as the template files provided by the instructor). Your implementation of the protocols must adhere to the interface provided in the template. It should be possible for anyone to run the simulation with the existing template, by only adding your protocol files to the existing template folder and changing the name in the “import <protocol-file>” line in the testbench.

1. Go through the Python code provided in the Template, and try to understand the behavior of each block. The file `Channel.py` implements a model for an unreliable channel over which packets can be corrupted or lost. This model has the following parameters:

**Pc:** The probability of a packet being corrupted

**Pl:** The probability of a packet being lost

**Delay:** The time it takes for a packet to travel over the channel and reach the receiver.

The file `Protocol_rdt1.py` implements the trivial protocol rdt1.0 which works only if the channel is assumed to be ideal. Run the simulation first with  $P_c=0$ , and then  $P_c=0.5$ . Check that the protocol fails in the second case, and list the failure symptoms.

[Note: The channel parameters such as  $P_c$  and  $P_l$  can be set in the file `Testbench.py`]

2. The file `Protocol_rdt2.py` implements the simple ACK/NAK based protocol rdt2.0 that can work when data packets can get corrupted. Check that this protocol indeed works by setting  $P_c>0$  for the data-channel.

[Note: The protocol to be used can be specified in the file `Testbench.py` by modifying the line `from Protocol_rdt<version> import *`]

3. For the testbench using `Protocol_rdt2.py`, modify the code such that:
  - a) The sending application generates a fixed total number of messages (say 1000), with a fixed time interval between each message (say 3 units of time)
  - b) As soon as the protocol at the sending-side (`rdt_Sender`) receives positive acknowledgements for all of the 1000 messages, the simulation ends, and a quantity “ $T_{avg}$ ” is printed as output, where  $T_{avg}$  is the time between sending a packet and receiving a positive acknowledgement for it, averaged across all packets. This is, in essence the Average Round-Trip Time (RTT avg).

(Note that you need to run each simulation as long as necessary for all 1000 messages to be acknowledged.)

You would expect that  $T_{avg}$  should increase with  $P_c$ , but in what manner? (linearly? exponentially? geometrically?). Obtain a plot of  $T_{avg}$  versus  $P_c$  for ( $0 \leq P_c \leq 0.9$ ) and explain what trend you observe and why.

4. Protocol\_rdt2.py will not work if ACK/NAK can also get corrupted. Check this by setting  $P_c > 0$  for the ack-channel and state the symptoms you observe.
5. Develop an alternating-bit protocol as described in K&R (rdt\_2.2) which can work even when both the data and ack packets can get corrupted. Test that your protocol works by simulating for a large amount of time or a large number of packets sent, with  $P_c > 0$  for both the data and ack channels. You need to submit only the protocol as a file named `Protocol_rdt22.py`.
6. The protocol rdt2.2 will not work if packets can be lost.
  - a) Check this by setting  $P_l > 0$  in the testbench with your implementation of Protocol\_rdt22. What failure symptoms do you observe?
  - b) Implement an alternating-bit protocol with Timeouts (rdt3.0) that can work when data or ack packets can be corrupted or lost. Set the timeout value to  $3 * \text{Delay}$ . Test that your protocol indeed works by simulating for a large amount of time or a large number of packets sent, with  $P_c > 0$  and  $P_l > 0$  for both the data and ack channels. You need to submit the protocol as a file named "Protocol\_rdt3.py".

[**Hint:** To implement timeouts, you may need to model a Timer in SimPy. A "skeleton" for implementing such a timer is provided here in Appendix A.]

- c) You would expect that  $T_{avg}$  (defined the same way as in question 1) would increase with  $P_l$ . For your implementation of the rdt3.0 protocol, with channel  $\text{Delay}=2$ ,  $P_c=0.2$ ,  $\text{timeout}=3 * \text{Delay}$  and total packets generated =1000, plot  $T_{avg}$  versus  $P_l$ .
- d) **[BONUS QUESTION]** For the scenario where packet loss is possible (assume  $P_c=0$ ), derive an analytical expression for how  $T_{avg}$  should vary with  $P_l$ . Check if the trend observed from simulations matches this.
- e) Protocol rdt3.0 will not work if packets can be re-ordered over the channel. This can be easily modeled by setting the channel delay for each packet to be a randomly chosen number instead of having the same value for all packets. Thus, packet-2 sent after packet-1, might experience a lower delay and arrive ahead of a packet-1 at the receiver. Implement this feature in the model, and check if the rdt3.0 protocol indeed fails. In your report, show the code snippet that models this packet reordering.

## APPENDIX A: SimPy Template for implementing Timeouts

# Here is a "skeleton" code for modeling timers and timeout events for the rdt\_Sender

```
class rdt_Sender(object):

    def __init__(self, env):
        .....

        # additional timer-related variables
        self.timeout_value=10
        self.timer_is_running=False
        self.timer=None

    # This function models a Timer's behavior.
    def timer_behavior(self):
        try:
            # Start
            self.timer_is_running=True
            yield self.env.timeout(self.timeout_value)
            # Stop
            self.timer_is_running=False
            # take some actions
            self.timeout_action()
        except simpy.Interrupt:
            # upon interrupt, stop the timer
            self.timer_is_running=False

    # This function can be called to start the timer
    def start_timer(self):
        assert(self.timer_is_running==False)
        self.timer=self.env.process(self.timer_behavior())

    # This function can be called to stop the timer
    def stop_timer(self):
        assert(self.timer_is_running==True)
        self.timer.interrupt()

    def timeout_action(self):
        # add here the actions to be performed
        # upon a timeout
        ...

    def rdt_send(self, msg):
        # whatever actions should go here
        ....

    def rdt_rcv(self, packet):
        # whatever actions should go here
        ....
```