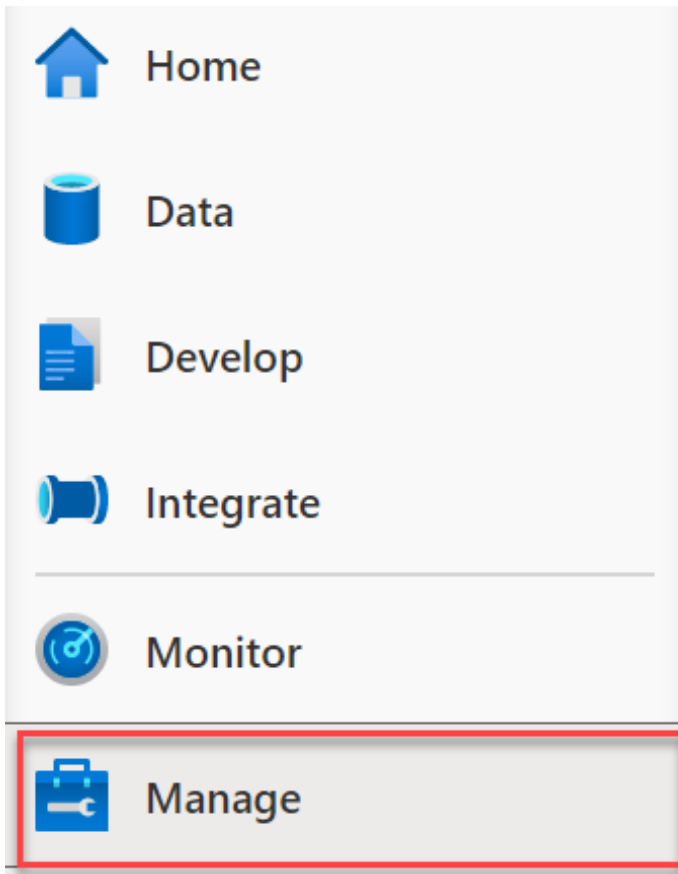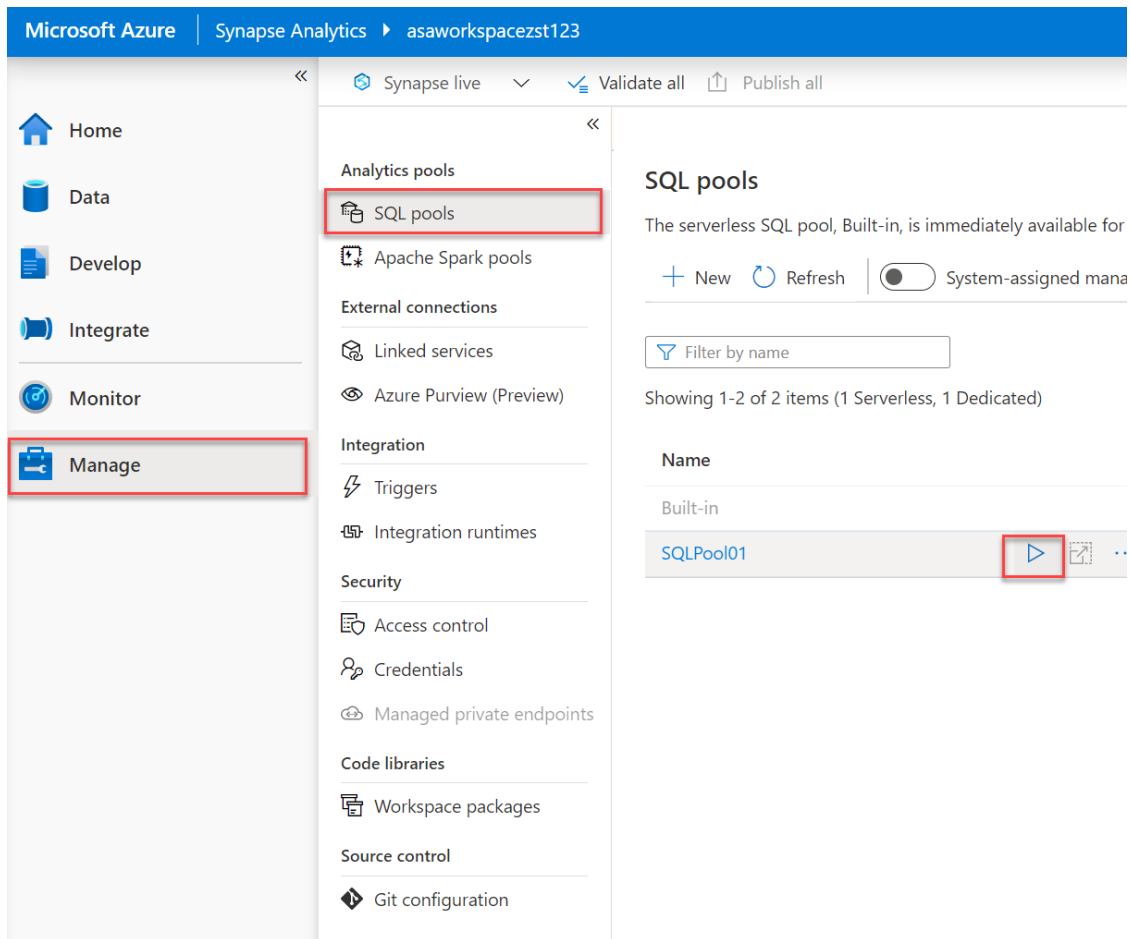# DW Optimization Part 1

Lab pre-requisite

Start the SQL Pool in your lab environment.

1. Open the Synapse Studio workspace and navigate to the **Manage** hub.

The Manage menu item is highlighted.

2. From the center menu, select **SQL pools** from beneath the **Analytics pools** heading. Locate SQLPool01, and select the **Resume** button.

The Manage menu item is selected, with SQL pools selected from the center menu. The resume button is selected next to the SQLPool01 item.

Exercise 1 - Explore query performance and improve table structure

Task 1 - Identify performance issues related to tables

1. In `Synapse Studio`, open a new SQL script and run the following statement (make sure you run queries on **SQLPool01** as opposed to `Built-in`):

```
SELECT
    COUNT_BIG(*)
FROM
    [wwi_perf].[Sale_Heap]
```

The script takes up to 15 seconds to execute and returns a count of ~ 340 million rows in the table.

2. Run the following (more complex) statement:

```
SELECT TOP 1000 * FROM
(
    SELECT
```

```
        S.CustomerId
        ,SUM(S.TotalAmount) as TotalAmount
    FROM
        [wwi_perf].[Sale_Heap] S
    GROUP BY
        S.CustomerId
) T
OPTION (LABEL = 'Lab03: Heap')
```

The script takes up to a 30 seconds to execute and returns the result. There is clearly something wrong with the Sale_Heap table that induces the performance hit.

Note the OPTION clause used in the statement. This comes in handy when you're looking to identify your query in the sys.dm_pdw_exec_requests DMV.

```
SELECT   *
FROM     sys.dm_pdw_exec_requests
WHERE    [label] = 'Lab03: Heap';
```

3. Check the structure of the Sale_Heap table, by right-clicking on it in the Data hub and selecting New SQL script and then CREATE. Take a look at the script used to create the table:

```
CREATE TABLE [wwi_perf].[Sale_Heap]
(
    [TransactionId] [uniqueidentifier]  NOT NULL,
    [CustomerId] [int]  NOT NULL,
    [ProductId] [smallint]  NOT NULL,
    [Quantity] [smallint]  NOT NULL,
    [Price] [decimal](9,2)  NOT NULL,
    [TotalAmount] [decimal](9,2)  NOT NULL,
    [TransactionDateId] [int]  NOT NULL,
    [ProfitAmount] [decimal](9,2)  NOT NULL,
    [Hour] [tinyint]  NOT NULL,
    [Minute] [tinyint]  NOT NULL,
    [StoreId] [smallint]  NOT NULL
)
WITH
(
    DISTRIBUTION = ROUND_ROBIN,
    HEAP
)
```

You can immediately spot at least two reasons for the performance hit:

- The ROUND_ROBIN distribution
- The HEAP structure of the table

NOTE

In this case, when we are looking for fast query response times, the heap structure is not a good choice as we will see in a moment. Still, there are cases where using a heap table can help performance rather than hurting it. One such example is when we're looking to ingest large amounts of data into the SQL pool.

4. Run the same script as the one you've run at step 2, but this time with the EXPLAIN WITH_RECOMMENDATIONS line before it:

```
EXPLAIN WITH_RECOMMENDATIONS
SELECT TOP 1000 * FROM
(
    SELECT
        S.CustomerId
        ,SUM(S.TotalAmount) as TotalAmount
    FROM
        [wwi_perf].[Sale_Heap] S
    GROUP BY
        S.CustomerId
) T
```

The EXPLAIN WITH_RECOMMENDATIONS clause returns the query plan for an Azure Synapse Analytics SQL statement without running the statement. Use EXPLAIN to preview which operations will require data movement and to view the estimated costs of the query operations. By default, you will get the execution plan in XML format, which you can export to other formats like CSV or JSON. **Do not** select Query Plan from the toolbar as it will try do download the query plan and open it in SQL Server Management Studio.

Your query should return something similar to:

```
<?xml version="1.0" encoding="utf-8"?>
<dsql_query number_nodes="1" number_distributions="60"
number_distributions_per_node="60">
<sql>SELECT TOP 1000 * FROM (SELECTS.CustomerId,SUM(S.TotalAmount) as
TotalAmountFROM[wwi_perf].[Sale_Heap] SGROUP BYS.CustomerId ) T</sql>
<materialized_view_candidates>
    <materialized_view_candidates with_constants="False">CREATE
MATERIALIZED VIEW View1 WITH (DISTRIBUTION = HASH([Expr0])) AS SELECT
[S].[CustomerId] AS [Expr0],SUM([S].[TotalAmount]) AS [Expr1] FROM
[wwi_perf].[Sale_Heap] GROUP BY
[S].[CustomerId]</materialized_view_candidates>
</materialized_view_candidates>
<dsql_operations total_cost="4.6605624" total_number_operations="5">
    <dsql_operation operation_type="RND_ID">
    <identifier>TEMP_ID_5</identifier>
    </dsql_operation>
    <dsql_operation operation_type="ON">
    <location permanent="false" distribution="AllDistributions" />
    <sql_operations>
        <sql_operation type="statement">CREATE TABLE
```

```
[qtabledb].[dbo].[TEMP_ID_5] ([CustomerId] INT NOT NULL, [col]
DECIMAL(38, 2) NOT NULL ) WITH(DISTRIBUTED_MOVE_FILE='');</sql_operation>
    </sql_operations>
    </dsql_operation>
    <dsql_operation operation_type="SHUFFLE_MOVE">
    <operation_cost cost="4.6605624" accumulative_cost="4.6605624"
average_rowsize="13" output_rows="89626.2" GroupNumber="11" />
    <source_statement>SELECT [T1_1].[CustomerId] AS [CustomerId],
[T1_1].[col] AS [col] FROM (SELECT SUM([T2_1].[TotalAmount]) AS [col],
[T2_1].[CustomerId] AS [CustomerId] FROM
[SQLPool01].[wwi_perf].[Sale_Heap] AS T2_1 GROUP BY [T2_1].[CustomerId])
AS T1_1 OPTION (MAXDOP 4, MIN_GRANT_PERCENT = [MIN_GRANT],
DISTRIBUTED_MOVE(N''))</source_statement>
    <destination_table>[TEMP_ID_5]</destination_table>
    <shuffle_columns>CustomerId;</shuffle_columns>
    </dsql_operation>
    <dsql_operation operation_type="RETURN">
    <location distribution="AllDistributions" />
    <select>SELECT [T1_1].[CustomerId] AS [CustomerId], [T1_1].[col] AS
[col] FROM (SELECT TOP (CAST ((1000) AS BIGINT)) SUM([T2_1].[col]) AS
[col], [T2_1].[CustomerId] AS [CustomerId] FROM
[qtabledb].[dbo].[TEMP_ID_5] AS T2_1 GROUP BY [T2_1].[CustomerId]) AS
T1_1 OPTION (MAXDOP 4, MIN_GRANT_PERCENT = [MIN_GRANT])</select>
    </dsql_operation>
    <dsql_operation operation_type="ON">
    <location permanent="false" distribution="AllDistributions" />
    <sql_operations>
        <sql_operation type="statement">DROP TABLE
[qtabledb].[dbo].[TEMP_ID_5]</sql_operation>
    </sql_operations>
    </dsql_operation>
</dsql_operations>
</dsql_query>
```

Notice the details of the internal layout of the MPP system:

```
<dsql_query number_nodes="1" number_distributions="60"
number_distributions_per_node="60">
```

This layout is given by the current Date Warehouse Units (DWU) setting. In the setup used for the example above, we were running at DW500c which means that there is a single physical node to service the 60 distributions, giving a number of, again, 60 distributions per physical node. Depending on your own DWU settings, these numbers will vary.

Here is an example layout we would get if we were running at DW2000c.

```
<dsql_query number_nodes=""4"" number_distributions=""60""
number_distributions_per_node=""15"">
```

In this case, we have four physical nodes to service the 60 distributions, giving a number of 15 distributions per physical node.

The query plan indicates data movement is required. This is indicated by the SHUFFLE_MOVE distributed SQL operation. Data movement is an operation where parts of the distributed tables are moved to different nodes during query execution. This operation is required where the data is not available on the target node, most commonly when the tables do not share the distribution key. The most common data movement operation is shuffle. During shuffle, for each input row, Synapse computes a hash value using the join columns and then sends that row to the node that owns that hash value. Either one or both sides of join can participate in the shuffle. The diagram below displays shuffle to implement join between tables T1 and T2 where neither of the tables is distributed on the join column col2.
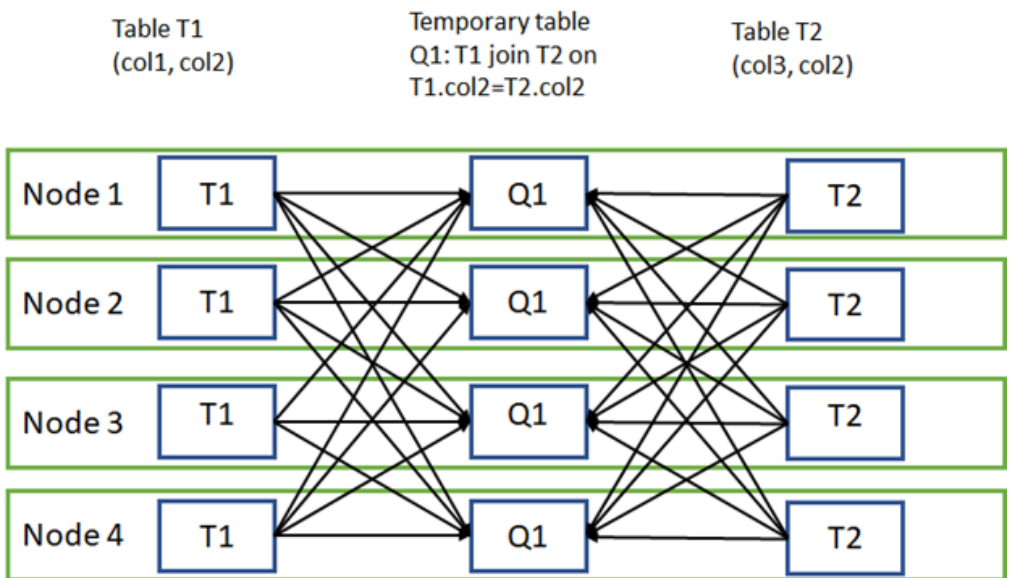
Shuffle move conceptual representation

Let's dive now into the details provided by the query plan to understand some of the problems our current approach has. The following table contains the description of every operation mentioned in the query plan:

Operation	Operation Type	Description
1	RND_ID	Identifies an object that will be created. In our case, it's the TEMP_ID_76 internal table.
2	ON	Specifies the location (nodes or distributions) where the operation will occur. AllDistributions means here the operation will be performed on each of the 60 distributions of the SQL pool. The operation will be a SQL operation (specified via <sql_operations>) that will create the

		TEMP_ID_76 table.
3	SHUFFLE_MOVE	The list of shuffle columns contains only one column which is CustomerId (specified via <suffle_columns>). The values will be distributed to the hash owning distributions and saved locally in the TEMP_ID_76 tables. The operation will output an estimated number of 41265.25 rows (specified via <operation_cost>). According to the same section, the average resulting row size is 13 bytes.
4	RETURN	Data resulting from the shuffle operation will be collected from all distributions (see <location>) by querying the internal temporary table TEMP_ID_76.
5	ON	The TEMP_ID_76 will be deleted from all distributions.

It becomes clear now what is the root cause of the performance problem: the inter-distribution data movements. This is actually one of the simplest examples given the small size of the data that needs to be shuffled. You can imagine how much worse things become when the shuffled row size becomes larger.

You can learn more about the structure of the query plan generated by the EXPLAIN statement here.

5. Besides the EXPLAIN statement, you can also understand the plan details using the sys.dm_pdw_request_steps DMV.

Query the sys.dm_pdw_exec_requests DMW to find your query id (this is for the query you executed previously at step 2):

```
SELECT
    *
FROM
    sys.dm_pdw_exec_requests
WHERE
    [label] = 'Lab03: Heap'
```

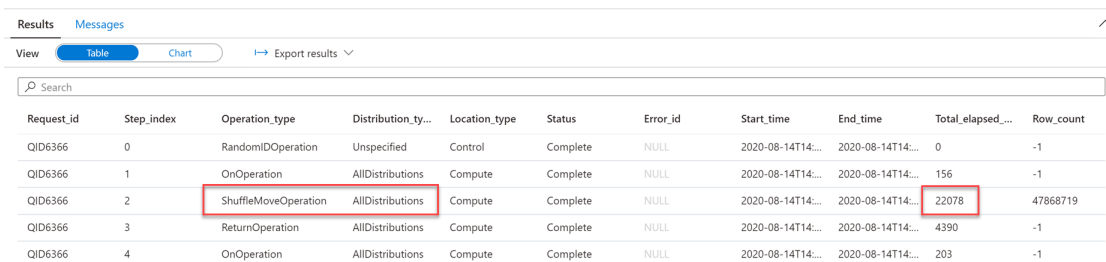The result contains, among other things, the query id (Request_id), the label, and the original SQL statement:

Request_id	Session_id	Status	Submit_time	Start_time	End_compile_ti...	End_time	Total_elapsed_...	Label	Error_id	Database_id	Command	F
QID6366	SID164	Completed	2020-08-14T14:...	2020-08-14T14:...	2020-08-14T14:...	2020-08-14T14:...	28906	Lab03: Heap	NULL	35	SELECT TOP 1000 * FROM (SELECT S.Cus...	s

Retrieving the query id

6. With the query id (`QID6366` in this case) you can now investigate the individual steps of the query:

```sql
SELECT
    *
FROM
    sys.dm_pdw_request_steps
WHERE
    request_id = 'QID6366'
ORDER BY
    step_index
```

The steps (indexed 0 to 4) are matching operations 2 to 6 from the query plan. Again, the culprit stands out: the step with index 2 describes the inter-partition data movement operation. By looking at the `TOTAL_ELAPSED_TIME` column one can clearly tell the largest part of the query time is generated by this step.

Results	Messages										∧
View	Table	Chart	⊢→ Export results ∨								
🔍 Search											
Request_id	Step_index	Operation_type	Distribution_ty...	Location_type	Status	Error_id	Start_time	End_time	Total_elapsed_...	Row_count	
QID6366	0	RandomIDOperation	Unspecified	Control	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	0	-1	
QID6366	1	OnOperation	AllDistributions	Compute	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	156	-1	
QID6366	2	ShuffleMoveOperation	AllDistributions	Compute	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	22078	47868719	
QID6366	3	ReturnOperation	AllDistributions	Compute	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	4390	-1	
QID6366	4	OnOperation	AllDistributions	Compute	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	203	-1	

Query execution steps

7. Get more details on the problematic step using the following SQL statement:

```sql
SELECT
    *
FROM
    sys.dm_pdw_sql_requests
WHERE
    request_id = 'QID6366'
    AND step_index = 1
```

The results of the statement provide details about what happens on each distribution within the SQL pool.

Request_id	Step_index	Pdw_node_id	Distribution_id	Status	Error_id	Start_time	End_time	Total_elapsed_...	Row_count	Spid	Command
QID6366	1	48	1	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	31	-1	700	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_1]
QID6366	1	48	2	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	31	-1	782	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_2]
QID6366	1	48	3	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	15	-1	739	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_3]
QID6366	1	48	4	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	46	-1	707	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_4]
QID6366	1	48	5	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	46	-1	726	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_5]
QID6366	1	48	6	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	62	-1	786	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_6]
QID6366	1	48	7	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	46	-1	728	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_7]
QID6366	1	48	8	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	62	-1	716	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_8]
QID6366	1	48	9	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	31	-1	787	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_9]
QID6366	1	48	10	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	62	-1	735	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_10]
QID6366	1	48	11	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	93	-1	737	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_11]
QID6366	1	48	12	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	46	-1	685	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_12]
QID6366	1	48	13	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	79	-1	798	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_13]
QID6366	1	48	14	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	15	-1	691	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_14]
QID6366	1	48	15	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	62	-1	751	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_15]
QID6366	1	48	16	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	46	-1	698	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_16]
QID6366	1	48	17	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	79	-1	703	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_17]
QID6366	1	48	18	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	79	-1	754	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_18]
QID6366	1	48	19	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	31	-1	749	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_19]
QID6366	1	48	20	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	62	-1	733	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_20]
QID6366	1	48	21	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	79	-1	725	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_21]
QID6366	1	48	22	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	46	-1	756	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_22]
QID6366	1	48	23	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	46	-1	719	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_23]
QID6366	1	48	24	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	46	-1	710	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_24]
QID6366	1	48	25	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	31	-1	784	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_25]
QID6366	1	48	26	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	62	-1	693	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_26]
QID6366	1	48	27	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	93	-1	720	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_27]
QID6366	1	48	28	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	46	-1	687	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_28]
QID6366	1	48	29	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	78	-1	732	CREATE TABLE tempdb.[dbo].[QTable_cf7b39ad821f49fa8f07ed1240eacece_29]
QID6366	1	48	30	Complete	NULL	2020-08-14T14:...	2020-08-14T14:...	31	-1	708	CREATE TABLE ...

Query execution step details

8. Finally, you can use the following SQL statement to investigate data movement on the distributed databases:

```
SELECT
    *
FROM
    sys.dm_pdw_dms_workers
WHERE
    request_id = 'QID6366'
    AND step_index = 2
ORDER BY
    distribution_id
```

The results of the statement provide details about data being moved at each distribution. The ROWS_PROCESSED column is especially useful here to get an estimate of the magnitude of the data movement happening when the query is executed.

Query execution step data movement

Task 2 - Improve table structure with hash distribution and columnstore index

1. Create an improved version of the table using CTAS (Create Table As Select):

```
CREATE TABLE [wwi_perf].[Sale_Hash]
WITH
(
    DISTRIBUTION = HASH ( [CustomerId] ),
    CLUSTERED COLUMNSTORE INDEX
)
AS
SELECT
    *
FROM
    [wwi_perf].[Sale_Heap]
```

The query will take up to 5 minutes to complete.

NOTE

CTAS is a more customizable version of the SELECT…INTO statement. SELECT…INTO doesn't allow you to change either the distribution method or the index type as part of the operation. You create the new table by using the default distribution type of ROUND_ROBIN, and the default table structure of CLUSTERED COLUMNSTORE INDEX.

With CTAS, on the other hand, you can specify both the distribution of the table data as well as the table structure type.

2. Run the query again to see the performance improvements:

```
SELECT TOP 1000 * FROM
(
    SELECT
        S.CustomerId
        ,SUM(S.TotalAmount) as TotalAmount
    FROM
        [wwi_perf].[Sale_Hash] S
    GROUP BY
        S.CustomerId
) T
```

3. Run the following EXPLAIN statement again to get the query plan (do not select `Query Plan` from the toolbar as it will try do download the query plan and open it in SQL Server Management Studio):

```
EXPLAIN
SELECT TOP 1000 * FROM
(
    SELECT
        S.CustomerId
        ,SUM(S.TotalAmount) as TotalAmount
    FROM
        [wwi_perf].[Sale_Hash] S
    GROUP BY
        S.CustomerId
) T
```

The resulting query plan is clearly much better than the previous one, as there is no more inter-distribution data movement involved.

```
<?xml version="1.0" encoding="utf-8"?>
<dsql_query number_nodes="1" number_distributions="60"
number_distributions_per_node="60">
<sql>SELECT TOP 1000 * FROM (SELECTS.CustomerId,SUM(S.TotalAmount) as
TotalAmountFROM[wwi_perf].[Sale_Hash] SGROUP BYS.CustomerId ) T</sql>
<dsql_operations total_cost="0" total_number_operations="1">
    <dsql_operation operation_type="RETURN">
    <location distribution="AllDistributions" />
    <select>SELECT [T1_1].[CustomerId] AS [CustomerId], [T1_1].[col] AS
[col] FROM (SELECT TOP (CAST ((1000) AS BIGINT))
SUM([T2_1].[TotalAmount]) AS [col], [T2_1].[CustomerId] AS [CustomerId]
FROM [SQLPool01].[wwi_perf].[Sale_Hash] AS T2_1 GROUP BY
[T2_1].[CustomerId]) AS T1_1 OPTION (MAXDOP 4)</select>
    </dsql_operation>
</dsql_operations>
</dsql_query>
```

4. Try running a more complex query and investigate the execution plan and execution steps. Here is an example of a more complex query you can use:

```
SELECT
    AVG(TotalProfit) as AvgMonthlyCustomerProfit
FROM
(
    SELECT
        S.CustomerId
        ,D.Year
        ,D.Month
        ,SUM(S.TotalAmount) as TotalAmount
        ,AVG(S.TotalAmount) as AvgAmount
        ,SUM(S.ProfitAmount) as TotalProfit
        ,AVG(S.ProfitAmount) as AvgProfit
    FROM
        [wwi_perf].[Sale_Partition01] S
        join [wwi].[Date] D on
            D.DateId = S.TransactionDateId
    GROUP BY
        S.CustomerId
        ,D.Year
        ,D.Month
) T
```

Task 3 - Improve further the structure of the table with partitioning

Date columns are usually good candidates for partitioning tables at the distributions level. In the case of your sales data, partitioning based on the TransactionDateId column seems to be a good choice.

Your SQL pool already contains two versions of the Sale table that have been partitioned using TransactionDateId. These tables are [wwi_perf].[Sale_Partition01] and [wwi_perf].[Sale_Partition02]. Below are the CTAS queries that have been used to create these tables.

> **Note**
>
> These queries have already been run on the SQL pool. If you want to test the CTAS queries yourself, make sure you replace the table names with new ones.

```
CREATE TABLE [wwi_perf].[Sale_Partition01]
WITH
(
    DISTRIBUTION = HASH ( [CustomerId] ),
    CLUSTERED COLUMNSTORE INDEX,
    PARTITION
    (
        [TransactionDateId] RANGE RIGHT FOR VALUES (
            20190101, 20190201, 20190301, 20190401, 20190501, 20190601,
20190701, 20190801, 20190901, 20191001, 20191101, 20191201)
    )
)
AS
```

```
SELECT
    *
FROM
    [wwi_perf].[Sale_Heap]
OPTION  (LABEL  = 'CTAS : Sale_Partition01')

CREATE TABLE [wwi_perf].[Sale_Partition02]
WITH
(
    DISTRIBUTION = HASH ( [CustomerId] ),
    CLUSTERED COLUMNSTORE INDEX,
    PARTITION
    (
        [TransactionDateId] RANGE RIGHT FOR VALUES (
            20190101, 20190401, 20190701, 20191001)
    )
)
AS
SELECT *
FROM
    [wwi_perf].[Sale_Heap]
OPTION  (LABEL  = 'CTAS : Sale_Partition02')
```

Notice the two partitioning strategies we've used here. The first partitioning scheme is month-based and the second is quarter-based. You will explore in Lab 04 the subtle differences between these and understand the potential performance implications resulting from these choices.

Exercise 2 - Improve query performance

Task 1 - Improve COUNT performance

1. The following query attempts to find the TOP 100 of customers that have the most sale transactions:

    ```
    SELECT COUNT( DISTINCT CustomerId) from wwi_perf.Sale_Heap
    ```

    Query takes up to 20 seconds to execute. That is expected, since distinct counts are one of the most difficult to optimize types of queries.

2. Run the HyperLogLog approach:

    ```
    SELECT APPROX_COUNT_DISTINCT(CustomerId) from wwi_perf.Sale_Heap
    ```

    Query takes about half the time to execute.

Task 2 - Use materialized views

As opposed to a standard view, a materialized view pre-computes, stores, and maintains its data in a Synapse SQL pool just like a table. Here is a basic comparison between standard and materialized views:

Comparison	View	Materialized View
View definition	Stored in Azure data warehouse.	Stored in Azure data warehouse.
View content	Generated each time when the view is used.	Pre-processed and stored in Azure data warehouse during view creation. Updated as data is added to the underlying tables.
Data refresh	Always updated	Always updated
Speed to retrieve view data from complex queries	Slow	Fast
Extra storage	No	Yes
Syntax	CREATE VIEW	CREATE MATERIALIZED VIEW AS SELECT

1. Execute the following query to get an approximation of its execution time:

```sql
SELECT TOP 1000 * FROM
(
    SELECT
        S.CustomerId
        ,D.Year
        ,D.Quarter
        ,SUM(S.TotalAmount) as TotalAmount
    FROM
        [wwi_perf].[Sale_Partition02] S
        join [wwi].[Date] D on
            S.TransactionDateId = D.DateId
    GROUP BY
        S.CustomerId
        ,D.Year
        ,D.Quarter
) T
```

2. Execute this query as well (notice the slight difference):

```sql
SELECT TOP 1000 * FROM
(
    SELECT
        S.CustomerId
        ,D.Year
        ,D.Month
        ,SUM(S.ProfitAmount) as TotalProfit
    FROM
        [wwi_perf].[Sale_Partition02] S
        join [wwi].[Date] D on
            S.TransactionDateId = D.DateId
    GROUP BY
        S.CustomerId
        ,D.Year
```

```
        ,D.Month
    ) T
```

3. Create a materialized view that can support both queries above:

```sql
CREATE MATERIALIZED VIEW
    wwi_perf.mvCustomerSales
WITH
(
    DISTRIBUTION = HASH( CustomerId )
)
AS
SELECT
    S.CustomerId
    ,D.Year
    ,D.Quarter
    ,D.Month
    ,SUM(S.TotalAmount) as TotalAmount
    ,SUM(S.ProfitAmount) as TotalProfit
FROM
    [wwi_perf].[Sale_Partition02] S
    join [wwi].[Date] D on
        S.TransactionDateId = D.DateId
GROUP BY
    S.CustomerId
    ,D.Year
    ,D.Quarter
    ,D.Month
```

4. Run the following query to get an estimated execution plan (do not select Query Plan from the toolbar as it will try do download the query plan and open it in SQL Server Management Studio):

```sql
EXPLAIN
SELECT TOP 1000 * FROM
(
    SELECT
        S.CustomerId
        ,D.Year
        ,D.Quarter
        ,SUM(S.TotalAmount) as TotalAmount
    FROM
        [wwi_perf].[Sale_Partition02] S
        join [wwi].[Date] D on
            S.TransactionDateId = D.DateId
    GROUP BY
        S.CustomerId
        ,D.Year
        ,D.Quarter
) T
```

The resulting execution plan shows how the newly created materialized view is used to optimize the execution. Note the FROM [SQLPool01].[wwi_perf].[mvCustomerSales] in the <dsql_operations> element.

```xml
<?xml version="1.0" encoding="utf-8"?>
<dsql_query number_nodes="1" number_distributions="60"
number_distributions_per_node="60">
<sql>SELECT TOP 1000 * FROM
(SELECTS.CustomerId,D.Year,D.Quarter,SUM(S.TotalAmount) as
TotalAmountFROM[wwi_perf].[Sale_Partition02] Sjoin [wwi].[Date] D
onS.TransactionDateId = D.DateIdGROUP BYS.CustomerId,D.Year,D.Quarter )
T</sql>
<dsql_operations total_cost="0" total_number_operations="1">
    <dsql_operation operation_type="RETURN">
    <location distribution="AllDistributions" />
    <select>SELECT [T1_1].[CustomerId] AS [CustomerId], [T1_1].[Year] AS
[Year], [T1_1].[Quarter] AS [Quarter], [T1_1].[col] AS [col] FROM (SELECT
TOP (CAST ((1000) AS BIGINT)) [T2_1].[CustomerId] AS [CustomerId],
[T2_1].[Year] AS [Year], [T2_1].[Quarter] AS [Quarter], [T2_1].[col1] AS
[col] FROM (SELECT ISNULL([T3_1].[col], CONVERT (BIGINT, 0, 0)) AS [col],
[T3_1].[CustomerId] AS [CustomerId], [T3_1].[Year] AS [Year],
[T3_1].[Quarter] AS [Quarter], [T3_1].[col1] AS [col1] FROM (SELECT
SUM([T4_1].[cb]) AS [col], SUM([T4_1].[TotalAmount]) AS [col1],
[T4_1].[CustomerId] AS [CustomerId], [T4_1].[Year] AS [Year],
[T4_1].[Quarter] AS [Quarter] FROM (SELECT [T5_1].[CustomerId] AS
[CustomerId], [T5_1].[TotalAmount] AS [TotalAmount], [T5_1].[cb] AS [cb],
[T5_1].[Quarter] AS [Quarter], [T5_1].[Year] AS [Year] FROM
[SQLPool01].[wwi_perf].[mvCustomerSales] AS T5_1) AS T4_1 GROUP BY
[T4_1].[CustomerId], [T4_1].[Year], [T4_1].[Quarter]) AS T3_1) AS T2_1
WHERE ([T2_1].[col] != CAST ((0) AS BIGINT))) AS T1_1 OPTION (MAXDOP
4)</select>
    </dsql_operation>
</dsql_operations>
</dsql_query>
```

5. The same materialized view is also used to optimize the second query. Get its execution plan:

```sql
EXPLAIN
SELECT TOP 1000 * FROM
(
    SELECT
        S.CustomerId
        ,D.Year
        ,D.Month
        ,SUM(S.ProfitAmount) as TotalProfit
    FROM
        [wwi_perf].[Sale_Partition02] S
        join [wwi].[Date] D on
            S.TransactionDateId = D.DateId
```

```
    GROUP BY
        S.CustomerId
        ,D.Year
        ,D.Month
) T
```

The resulting execution plan shows the use of the same materialized view to optimize execution:

```xml
<?xml version="1.0" encoding="utf-8"?>
<dsql_query number_nodes="1" number_distributions="60"
number_distributions_per_node="60">
<sql>SELECT TOP 1000 * FROM
(SELECTS.CustomerId,D.Year,D.Month,SUM(S.ProfitAmount) as
TotalProfitFROM[wwi_perf].[Sale_Partition02] Sjoin [wwi].[Date] D
onS.TransactionDateId = D.DateIdGROUP BYS.CustomerId,D.Year,D.Month )
T</sql>
<dsql_operations total_cost="0" total_number_operations="1">
    <dsql_operation operation_type="RETURN">
    <location distribution="AllDistributions" />
    <select>SELECT [T1_1].[CustomerId] AS [CustomerId], [T1_1].[Year] AS
[Year], [T1_1].[Month] AS [Month], [T1_1].[col] AS [col] FROM (SELECT TOP
(CAST ((1000) AS BIGINT)) [T2_1].[CustomerId] AS [CustomerId],
[T2_1].[Year] AS [Year], [T2_1].[Month] AS [Month], [T2_1].[col1] AS
[col] FROM (SELECT ISNULL([T3_1].[col], CONVERT (BIGINT, 0, 0)) AS [col],
[T3_1].[CustomerId] AS [CustomerId], [T3_1].[Year] AS [Year],
[T3_1].[Month] AS [Month], [T3_1].[col1] AS [col1] FROM (SELECT
SUM([T4_1].[cb]) AS [col], SUM([T4_1].[TotalProfit]) AS [col1],
[T4_1].[CustomerId] AS [CustomerId], [T4_1].[Year] AS [Year],
[T4_1].[Month] AS [Month] FROM (SELECT [T5_1].[CustomerId] AS
[CustomerId], [T5_1].[TotalProfit] AS [TotalProfit], [T5_1].[cb] AS [cb],
[T5_1].[Month] AS [Month], [T5_1].[Year] AS [Year] FROM
[SQLPool01].[wwi_perf].[mvCustomerSales] AS T5_1) AS T4_1 GROUP BY
[T4_1].[CustomerId], [T4_1].[Year], [T4_1].[Month]) AS T3_1) AS T2_1
WHERE ([T2_1].[col] != CAST ((0) AS BIGINT))) AS T1_1 OPTION (MAXDOP
4)</select>
    </dsql_operation>
</dsql_operations>
</dsql_query>
```

Note

Even if the two queries have different aggregation levels, the query optimizer is able to infer the use of the materialized view. This happens because the materialized view covers both aggregation levels (`Quarter` and `Month`) as well as both aggregation measures (`TotalAmount` and `ProfitAmount`).

6. Check the materialized view overhead:

```
DBCC PDW_SHOWMATERIALIZEDVIEWOVERHEAD ( 'wwi_perf.mvCustomerSales' )
```

The results show that BASE_VIEW_ROWS are equal to TOTAL_ROWS (and hence OVERHEAD_RATIO is 1). The materialized view is perfectly aligned with the base view. This situation is expected to change once the underlying data starts to change.

7. Update the original data the materialized view was built on:

```
UPDATE
    [wwi_perf].[Sale_Partition02]
SET
    TotalAmount = TotalAmount * 1.01
    ,ProfitAmount = ProfitAmount * 1.01
WHERE
    CustomerId BETWEEN 100 and 200
```

8. Check the materialized view overhead again:

```
DBCC PDW_SHOWMATERIALIZEDVIEWOVERHEAD ( 'wwi_perf.mvCustomerSales' )
```

OBJECT_ID	BASE_VIEW_ROWS	TOTAL_ROWS	OVERHEAD_RATIO
544720993	118109924	118121863	1.00010108380054499

Materialized view overhead after update

There is now a delta stored by the materialized view which results in TOTAL_ROWS being greater than BASE_VIEW_ROWS and OVERHEAD_RATIO being greater than 1.

9. Rebuild the materialized view and check that the overhead ration went back to 1:

```
ALTER MATERIALIZED VIEW [wwi_perf].[mvCustomerSales] REBUILD
```

```
DBCC PDW_SHOWMATERIALIZEDVIEWOVERHEAD ( 'wwi_perf.mvCustomerSales' )
```

OBJECT_ID	BASE_VIEW_ROWS	TOTAL_ROWS	OVERHEAD_RATIO
544720993	118109924	118109924	1.00000000000000000

Materialized view overhead after rebuild

Task 3 - Use result set caching

1. Check if result set caching is on in the current SQL pool:

```
SELECT
    name
    ,is_result_set_caching_on
```

```
FROM
    sys.databases
```


Check result set caching settings at the database level

If `False` is returned for your SQL pool, run the following query to activate it (you need to run it on the `master` database and replace `<sql_pool>` with the name of your SQL pool):

```
ALTER DATABASE [<sql_pool>]
SET RESULT_SET_CACHING ON
```

Important

The operations to create result set cache and retrieve data from the cache happen on the control node of a Synapse SQL pool instance. When result set caching is turned ON, running queries that return large result set (for example, >1GB) can cause high throttling on the control node and slow down the overall query response on the instance. Those queries are commonly used during data exploration or ETL operations. To avoid stressing the control node and cause performance issue, users should turn OFF result set caching on the database before running those types of queries.

2. After activating result set caching, run a query and immediately check if it hit the cache (change the database back to your SQL Pool):

```
SELECT
    D.Year
    ,D.Quarter
    ,D.Month
    ,SUM(S.TotalAmount) as TotalAmount
    ,SUM(S.ProfitAmount) as TotalProfit
FROM
```

```
    [wwi_perf].[Sale_Partition02] S
    join [wwi].[Date] D on
        S.TransactionDateId = D.DateId
GROUP BY
    D.Year
    ,D.Quarter
    ,D.Month
OPTION (LABEL = 'Lab03: Result set caching')

SELECT
    result_cache_hit
FROM
    sys.dm_pdw_exec_requests
WHERE
    request_id =
    (
        SELECT TOP 1
            request_id
        FROM
            sys.dm_pdw_exec_requests
        WHERE
            [label] = 'Lab03: Result set caching'
        ORDER BY
            start_time desc
    )
```

As expected, the result is `False`. Still, you can identify that, while running the query, Synapse has also cached the result set. Run the following query to get the execution steps:

```
SELECT
    step_index
    ,operation_type
    ,location_type
    ,status
    ,total_elapsed_time
    ,command
FROM
    sys.dm_pdw_request_steps
WHERE
    request_id =
    (
        SELECT TOP 1
            request_id
        FROM
            sys.dm_pdw_exec_requests
        WHERE
            [label] = 'Lab03: Result set caching'
        ORDER BY
```

```
                    start_time desc
    )
```

The execution plan reveals the building of the result set cache:

STEP_INDEX	OPERATION_TYPE	LOCATION_TYPE	STATUS	TOTAL_ELAPSED_TIME	COMMAND
0	RandomIDOperation	Control	Complete	0	TEMP_ID_41
1	OnOperation	Compute	Complete	140	CREATE TABLE [qtabledb].[dbo].[TEMP_ID_41] ([Month] TINYINT NOT NULL, [Quarter] TINYINT NOT NULL, [Year] SMALLINT NOT NULL, [col] DECIMAL(38, 2), [col1] D
2	ShuffleMoveOperation	Compute	Complete	625	SELECT [T1_1].[Month] AS [Month], [T1_1].[Quarter] AS [Quarter], [T1_1].[Year] AS [Year], [T1_1].[col1] AS [col], [T1_1].[col2] AS [col1], [T1_1].[col] AS [col2] FROM (SEL
3	PartitionMoveOperation	Compute	Complete	156	SELECT [T1_1].[Year] AS [Year], [T1_1].[Quarter] AS [Quarter], [T1_1].[Month] AS [Month], [T1_1].[col] AS [col], [T1_1].[col1] AS [col1] FROM (SELECT [T2_1].[Year] AS [Ye
4	ReturnOperation	Control	Complete	156	select * from [DWResultCacheDb].dbo.[iq_[00D973C]-5510-440D-A345-5939E7175E05]] OPTION(MAXDOP 1)
5	OnOperation	Compute	Complete	93	DROP TABLE [qtabledb].[dbo].[TEMP_ID_41]

The building of the result set cache

3. You can control at the user session level the use of the result set cache. The following query shows how to deactivate and activate the result cache:

```sql
SET RESULT_SET_CACHING OFF

SELECT
    D.Year
    ,D.Quarter
    ,D.Month
    ,SUM(S.TotalAmount) as TotalAmount
    ,SUM(S.ProfitAmount) as TotalProfit
FROM
    [wwi_perf].[Sale_Partition02] S
    join [wwi].[Date] D on
        S.TransactionDateId = D.DateId
GROUP BY
    D.Year
    ,D.Quarter
    ,D.Month
OPTION (LABEL = 'Lab03: Result set caching off')

SET RESULT_SET_CACHING ON

SELECT
    D.Year
    ,D.Quarter
    ,D.Month
    ,SUM(S.TotalAmount) as TotalAmount
    ,SUM(S.ProfitAmount) as TotalProfit
FROM
    [wwi_perf].[Sale_Partition02] S
    join [wwi].[Date] D on
        S.TransactionDateId = D.DateId
GROUP BY
    D.Year
    ,D.Quarter
    ,D.Month
```

```
OPTION (LABEL = 'Lab03: Result set caching on')

SELECT TOP 2
    request_id
    ,[label]
    ,result_cache_hit
FROM
    sys.dm_pdw_exec_requests
WHERE
    [label] in ('Lab03: Result set caching off', 'Lab03: Result set
caching on')
ORDER BY
    start_time desc
```

The result of SET RESULT_SET_CACHING OFF is visible in the cache hit test results (The
result_cache_hit column returns 1 for cache hit, 0 for cache miss, and *negative values*
for reasons why result set caching was not used.):

Result cache on and off

4. At any moment, you can check the space used by the results cache:

```
DBCC SHOWRESULTCACHESPACEUSED
```


Check the size of the result set cache

5. Clear the result set cache using:

```
DBCC DROPRESULTSETCACHE
```

6. Finally, disable result set caching on the database using the following query (you need to run it on the `master` database and replace ` with the name of your SQL pool): (It can take up to a minute to disable Result Set Caching.)

```
ALTER DATABASE [<sql_pool>]
SET RESULT_SET_CACHING OFF
```

Important

Make sure you disable result set caching on the SQL pool. Failing to do so will have a negative impact on the remainder of this lab, as it will skew execution times and defeat the purpose of several upcoming exercises.

Note

The maximum size of result set cache is 1 TB per database. The cached results are automatically invalidated when the underlying query data change.

The cache eviction is managed by SQL Analytics automatically following this schedule:
   – Every 48 hours if the result set hasn't been used or has been invalidated.
   – When the result set cache approaches the maximum size.

Users can manually empty the entire result set cache by using one of these options:
   – Turn OFF the result set cache feature for the database
   – Run DBCC DROPRESULTSETCACHE while connected to the database

Pausing a database won't empty cached result set.

Task 4 - Create and update statistics

The more the SQL pool resource knows about your data, the faster it can execute queries. After loading data into SQL pool, collecting statistics on your data is one of the most important things you can do for query optimization.

The SQL pool query optimizer is a cost-based optimizer. It compares the cost of various query plans, and then chooses the plan with the lowest cost. In most cases, it chooses the plan that will execute the fastest.

For example, if the optimizer estimates that the date your query is filtering on will return one row it will choose one plan. If it estimates that the selected date will return 1 million rows, it will return a different plan.

1. Check if statistics are set to be automatically created in the database:

```
SELECT name, is_auto_create_stats_on
FROM sys.databases
```

2. See statistics that have been automatically created (change the database back to your SQL Pool):

```
SELECT
    *
FROM
```

```
    sys.dm_pdw_exec_requests
WHERE
    Command like 'CREATE STATISTICS%'
```

Notice the special name pattern used for automatically created statistics:

View automatically created statistics

3. Check if there are any statistics created for `CustomerId` from the `wwi_perf.Sale_Has` table:

```
DBCC SHOW_STATISTICS ('wwi_perf.Sale_Hash', CustomerId) WITH HISTOGRAM
```

You should get an error stating that statistics for `CustomerId` does not exist.

4. Create statistics for `CustomerId`:

```
CREATE STATISTICS Sale_Hash_CustomerId ON wwi_perf.Sale_Hash (CustomerId)
```

Display the newly created statistics:

```
DBCC SHOW_STATISTICS([wwi_perf.Sale_Hash], 'Sale_Hash_CustomerId')
```

In the results pane, switch to `Chart` display and set the category column and the legend columns as presented below:
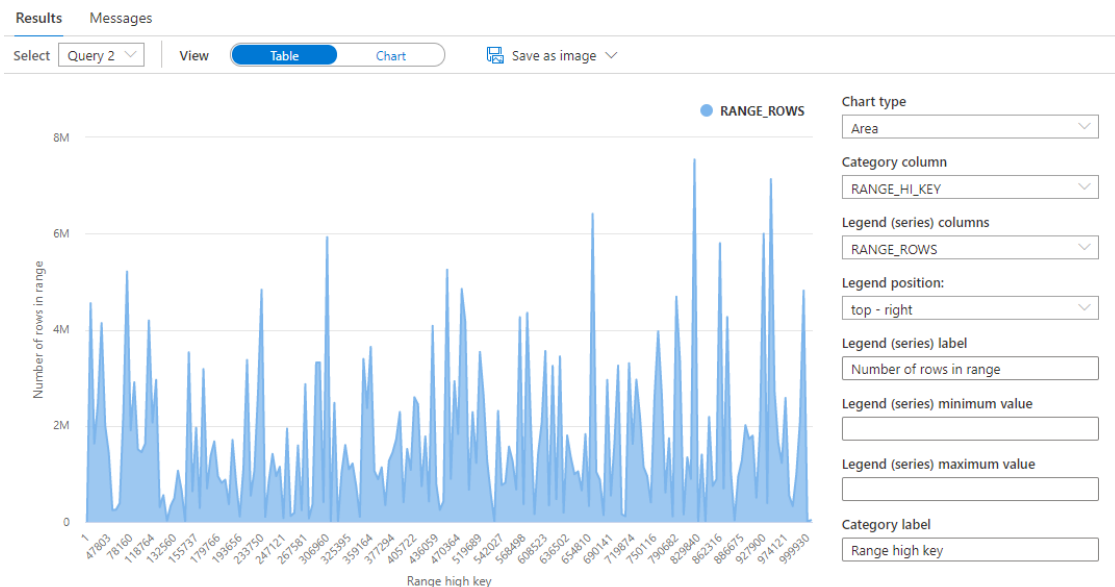
Statistics created for CustomerId

You now have a visual on the statistics created for the `CustomerId` column.

Important

The more SQL pool knows about your data, the faster it can execute queries against it. After loading data into SQL pool, collecting statistics on your data is one of the most important things you can do to optimize your queries.

The SQL pool query optimizer is a cost-based optimizer. It compares the cost of various query plans, and then chooses the plan with the lowest cost. In most cases, it chooses the plan that will execute the fastest.

For example, if the optimizer estimates that the date your query is filtering on will return one row it will choose one plan. If it estimates that the selected date will return 1 million rows, it will return a different plan.

Task 5 - Create and update indexes

Clustered Columnstore Index vs. Heap vs. Clustered and Nonclustered

Clustered indexes may outperform clustered columnstore indexes when a single row needs to be quickly retrieved. For queries where a single or very few row lookup is required to perform with extreme speed, consider a cluster index or nonclustered secondary index. The disadvantage to using a clustered index is that only queries that benefit are the ones that use a highly selective filter on the clustered index column. To improve filter on other columns a nonclustered index can be added to other columns. However, each index which is added to a table adds both space and processing time to loads.

1. Retrieve information about a single customer from the table with CCI:

```
SELECT
    *
FROM
    [wwi_perf].[Sale_Hash]
WHERE
    CustomerId = 500000
```

Take a note of the execution time.

2. Retrieve information about a single customer from the table with a clustered index:

```
SELECT
    *
FROM
    [wwi_perf].[Sale_Index]
WHERE
    CustomerId = 500000
```

The execution time is similar to the one for the query above. Clustered columnstore indexes have no significant advantage over clustered indexes in the specific scenario of highly selective queries.

3. Retrieve information about multiple customers from the table with CCI:

```
SELECT
    *
FROM
    [wwi_perf].[Sale_Hash]
WHERE
    CustomerId between 400000 and 400100
```

and then retrieve the same information from the table with a clustered index:

```
SELECT
    *
FROM
    [wwi_perf].[Sale_Index]
WHERE
    CustomerId between 400000 and 400100
```

Run both queries several times to get a stable execution time. Under normal conditions, you should see that even with a relatively small number of customers, the CCI table starts yielding better results than the clustered index table.

4. Now add an extra condition on the query, one that refers to the StoreId column:

```
SELECT
    *
FROM
    [wwi_perf].[Sale_Index]
WHERE
    CustomerId between 400000 and 400100
    and StoreId between 2000 and 4000
```

Take a note of the execution time.

5. Create a non-clustered index on the StoreId column:

```
CREATE INDEX Store_Index on wwi_perf.Sale_Index (StoreId)
```

The creation of the index should complete in a few minutes. Once the index is created, run the previous query again. Notice the improvement in execution time resulting from the newly created non-clustered index.

Note

Creating a non-clustered index on the wwi_perf.Sale_Index is based on the already existing clustered index. As a bonus exercise, try to create the same type of index on the wwi_perf.Sale_Hash table. Can you explain the difference in index creation time?

Task 6 - Ordered Clustered Columnstore Indexes

By default, for each table created without an index option, an internal component (index builder) creates a non-ordered clustered columnstore index (CCI) on it. Data in each column is compressed into a separate CCI rowgroup segment. There's metadata on each segment's value range, so segments that are outside the bounds of the query predicate aren't read from disk during query execution. CCI offers the highest level of data compression and reduces the size of segments to read so queries can run faster. However, because the index builder doesn't sort data before compressing them into segments, segments with overlapping value ranges could occur, causing queries to read more segments from disk and take longer to finish.

When creating an ordered CCI, the Synapse SQL engine sorts the existing data in memory by the order key(s) before the index builder compresses them into index segments. With sorted data, segment overlapping is reduced allowing queries to have a more efficient segment elimination and thus faster performance because the number of segments to read from disk is smaller. If all data can be sorted in memory at once, then segment overlapping can be avoided. Due to large tables in data warehouses, this scenario doesn't happen often.

Queries with the following patterns typically run faster with ordered CCI:

- The queries have equality, inequality, or range predicates
- The predicate columns and the ordered CCI columns are the same.
- The predicate columns are used in the same order as the column ordinal of ordered CCI columns.
1. Run the following query to show the segment overlaps for the `Sale_Hash` table:

```
select
    OBJ.name as table_name
    ,COL.name as column_name
    ,NT.distribution_id
    ,NP.partition_id
    ,NP.rows as partition_rows
    ,NP.data_compression_desc
    ,NCSS.segment_id
    ,NCSS.version
    ,NCSS.min_data_id
    ,NCSS.max_data_id
    ,NCSS.row_count
from
    sys.objects OBJ
    JOIN sys.columns as COL ON
        OBJ.object_id = COL.object_id
    JOIN sys.pdw_table_mappings TM ON
        OBJ.object_id = TM.object_id
    JOIN sys.pdw_nodes_tables as NT on
        TM.physical_name = NT.name
    JOIN sys.pdw_nodes_partitions NP on
        NT.object_id = NP.object_id
```

```sql
        and NT.pdw_node_id = NP.pdw_node_id
        and substring(TM.physical_name, 40, 10) = NP.distribution_id
    JOIN sys.pdw_nodes_column_store_segments NCSS on
        NP.partition_id = NCSS.partition_id
        and NP.distribution_id = NCSS.distribution_id
        and COL.column_id = NCSS.column_id
where
    OBJ.name = 'Sale_Hash'
    and COL.name = 'CustomerId'
    and TM.physical_name  not like '%HdTable%'
order by
    NT.distribution_id
```

Here is a short description of the tables involved in the query:

Table Name	Description
sys.objects	All objects in the database. Filtered to match only the `Sale_Hash` table.
sys.columns	All columns in the database. Filtered to match only the `CustomerId` column of the `Sale_Hash` table.
sys.pdw_table_mappings	Maps each table to local tables on physical nodes and distributions.
sys.pdw_nodes_tables	Contains information on each local table in each distribution.
sys.pdw_nodes_partitions	Contains information on each local partition of each local table in each distribution.
sys.pdw_nodes_column_store_segments	Contains information on each CCI segment for each partition and distribution column of each local table in each distribution. Filtered to match only the `CustomerId` column of the `Sale_Hash` table.

With this information on hand, take a look at the result:

CCI segment structure on each distribution

Browse through the result set and notice the significant overlap between segments. There is literally overlap in customer ids between every single pair of segments (CustomerId values in the data range from 1 to 1,000,000). The segment structure of this CCI is clearly inefficient and will result in a lot of unnecessary reads from storage.

2. Run the following query to show the segment overlaps for the Sale_Hash_Ordered table:

```
select
    OBJ.name as table_name
    ,COL.name as column_name
    ,NT.distribution_id
    ,NP.partition_id
    ,NP.rows as partition_rows
    ,NP.data_compression_desc
    ,NCSS.segment_id
    ,NCSS.version
    ,NCSS.min_data_id
    ,NCSS.max_data_id
    ,NCSS.row_count
from
    sys.objects OBJ
    JOIN sys.columns as COL ON
        OBJ.object_id = COL.object_id
    JOIN sys.pdw_table_mappings TM ON
        OBJ.object_id = TM.object_id
    JOIN sys.pdw_nodes_tables as NT on
        TM.physical_name = NT.name
    JOIN sys.pdw_nodes_partitions NP on
        NT.object_id = NP.object_id
        and NT.pdw_node_id = NP.pdw_node_id
        and substring(TM.physical_name, 40, 10) = NP.distribution_id
    JOIN sys.pdw_nodes_column_store_segments NCSS on
        NP.partition_id = NCSS.partition_id
        and NP.distribution_id = NCSS.distribution_id
        and COL.column_id = NCSS.column_id
where
    OBJ.name = 'Sale_Hash_Ordered'
```

```
    and COL.name = 'CustomerId'
    and TM.physical_name  not like '%HdTable%'
order by
    NT.distribution_id
```

The CTAS used to create the `wwi_perf.Sale_Hash_Ordered` table was the following:

```
CREATE TABLE [wwi_perf].[Sale_Hash_Ordered]
WITH
(
    DISTRIBUTION = HASH ( [CustomerId] ),
    CLUSTERED COLUMNSTORE INDEX ORDER( [CustomerId] )
)
AS
SELECT
    *
FROM
    [wwi_perf].[Sale_Heap]
OPTION  (LABEL  = 'CTAS : Sale_Hash', MAXDOP 1)
```

Notice the creation of the ordered CCI with MAXDOP = 1. Each thread used for ordered CCI creation works on a subset of data and sorts it locally. There's no global sorting across data sorted by different threads. Using parallel threads can reduce the time to create an ordered CCI but will generate more overlapping segments than using a single thread. Currently, the MAXDOP option is only supported in creating an ordered CCI table using CREATE TABLE AS SELECT command. Creating an ordered CCI via CREATE INDEX or CREATE TABLE commands does not support the MAXDOP option.

The results show significantly less overlap between segments:

TABLE_NAME	COLUMN_NAME	DISTRIBUTION_ID	PARTITION_ID	PARTITION_ROWS	DATA_COMPRESSION_DESC	SEGMENT_ID	VERSION	MIN_DATA_ID	MAX_DATA_ID	ROW_COUNT
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	0	1	10	347705	1048576
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	1	1	347705	692049	1048576
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	2	1	10	999996	1048576
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	3	1	20939	377130	1048576
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	4	1	377130	723058	1048576
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	5	1	10	999996	1048576
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	6	1	63020	410614	1048576
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	7	1	410614	756876	1048576
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	8	1	10	999996	1048576
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	9	1	94801	439760	1048576
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	10	1	439760	785796	1048576
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	11	1	10	999996	1048576
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	12	1	127555	474518	1048576
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	13	1	474518	817481	1048576
Sale_Hash_Ordered	CustomerId	1	72057594261405696	48215156	COLUMNSTORE	14	1	10	999996	1048576

CCI segment structure on each distribution with ordered CCI

Note

You will learn more about the internal organization of the clustered columnstore indexes in the following lab.