# Introducing a Subroutine Call-Return Mechanism

- How can we have a different delay for each call of the subroutine?
  - We must somehow convey the delay duration to the subroutine from the calling location.
  - There are 2 ways of doing this
    - 1. We simply set the value of the bl register to the desired delay before making the call; the subroutine will use this value as before.
    - 2. We put the value on the stack before making the call and add appropriate code to the subroutine to access it correctly.
  - Method 1 is equivalent to having a global variable
  - Method 2 is in the spirit of passing a parameter to the subroutine.
  - Both methods have advantages and disadvantages depending on your perspective and the abstraction level that you are working on.

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#### Illustrating Method 1

```
;Traffic Lights with a various delay durations using subroutines
 mov al, 84 ; 84 corresponds to Red-Green on the Traffic Lights
 out 01 ; Write to Traffic Lights Port
mov bl, fa ; initialize bl with a value representing the delay
  call 70
  mov al, 48 ; 48 corresponds to Amber-Amber on the Traffic Lights
 out 01 ; Write to Traffic Lights Port mov bl, fe ; initialize bl with a value representing the delay
  call 70
  mov al, 30
               ; 30 corresponds to Green-Red on the Traffic Lights
  mov bl, ef
               ; initialize bl with a value representing the delay
  call 70
 mov al, 48 ; 48 corresponds to Amber-Amber on the Traffic Lights
  out 01
 out 01 ; Write to Traffic Lights Port mov bl, fe ; initialize bl with a value representing the delay
  call 70
jmp start
org 70 ;
loop:
inc bl
 cmp bl, 00 ; check to see if bl has overflowed
jnz loop ; if not continue incrementing and checking
```

## Illustrating Method 2

```
;Traffic Lights with a various delay durations using subroutines
 mov al, 84 ; 84 corresponds to Red-Green on the Traffic Lights
 out 01
              ; Write to Traffic Lights Port
 mov bl, fa ; initialize bl with a value representing the delay
 push bl
               ; place parameter on stack
 call 70
 mov al, 48 ; 48 corresponds to Amber-Amber on the Traffic Lights
 out 01 ; Write to Traffic Lights Port
mov bl, fe ; initialize bl with a value representing the delay
 push bl
              ; place parameter on stack
 call 70
 mov al, 30 ; 30 corresponds to Green-Red on the Traffic Lights
 out 01
 mov bl, ef ; initialize bl with a value representing the delay
 push bl
               ; place parameter on stack
 call 70
 mov al, 48 ; 48 corresponds to Amber-Amber on the Traffic Lights
 out 01
               ; Write to Traffic Lights Port
 mov bl, fe ; initialize bl with a value representing the delay
 push bl
              ; place parameter on stack
 call 70
jmp start
```

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## Illustrating Method 2 - continued

```
org 70 ; Subroutine

pop cl ; this will be the return address.

pop bl ; this is the parameter.

push cl; put return address back on stack

loop:

inc bl

cmp bl, 00 ; check to see if bl has overflowed

jnz loop ; if not continue incrementing and checking

ret

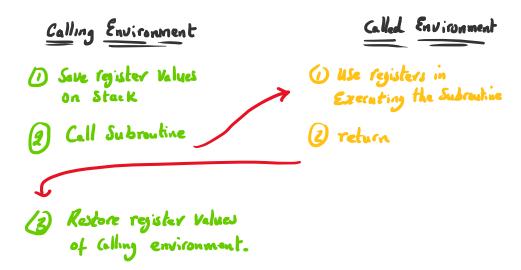
end
```

#### Why Method 2?

- Method 2 looks to be a lot more complicated
  - It is, but it is also more flexible
  - In general, we not only use the stack to pass parameters; we also use it to store the values of all the registers from the calling environment.
  - This, context switch, allows us to reuse all the registers in processing the code of the subroutine without fear of overwriting their values from another context.
  - When the subroutine ends, the context can be switched again – with the stored values of the calling environment registers being restored to the appropriate registers.

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#### Why Method 2? - Context Switching



```
mov al, 30
                                 } use registers in calling Env
mov bl, 31
mov cl, 32
mov dl, 33
push al ; save the context of the calling env | Save State of calling push bl
push bl
push cl
push dl
                                                Call Subsoutine
call 70
pop dl ; restore the context of thecalling env \ restore state of
pop cl
                                                I calling Env
pop bl
pop al
mov [c0], al | Print State | mov [c2], c1 | mov [c3], d1 |
add a1, b1 Reuse Registers to Implement add a1, a1 functionality of Submittine mov [d0], a1
ret
```

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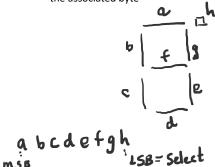
#### **Creating a Translation Table**

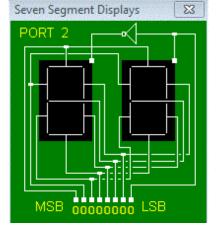
- The ASCII table allows us to translate between symbols and numbers. We can view the ASCII table as a translation table.
- To translate between numbers and the 7segment display (for example) we also need to create an appropriate translation table.

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## Creating a Translation Table

- Consider the 7-segment display:
- We could label all of the LEDs, comprising the 7 segments, in order from the most significant bit (MSB) to the least significant bit (LSB) in the associated byte





Select=1 => R.H.S. Display active Select=0 => 2.H.S. Display active

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## Creating a Translation Table

- To display 0, for example, on a display requires
- a, b, c, d, e, g = 1 and f = 0
- The value of h will determine the display



## Creating a Translation Table

Image	а	b	С	d	е	f	g	h	LHS	RHS
0	1	1	1	1	1	0	1	0 1	FA	FB
1 1	0	0	0	0	1	0	1	0 1	0A	OB
2 💆	1	0	1	1	0	1	1	0 1	В6	В7
3 🗒	1	0	0	1	1	1	1	0 1	9E	9F
4 4	0	1	0	0	1	1	1	0 1	4E	4F
5 =	1	1	0	1	1	1	0	0 1	DC	DD
6	1	1	1	1	1	1	0	0 1	FC	FD
7 ~	1	0	0	0	1	0	1	0 1	8A	8B
8 🗒	1	1	1	1	1	1	1	0 1	FE	FF
9 🖺	1	1	0	0	1	1	1	0 1	CE	CF

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## Creating a Translation Table

- In Samphire, we can use the Assembler
   Directives, DB and ORG to place a specific byte
   into memory starting at a specific memory
   location.
- Thus,

org 50

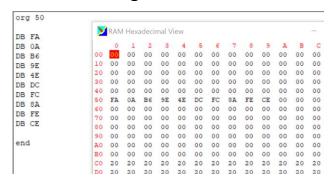
DB FA

Will place the byte FA into memory at address 50.

Each subsequent DB directive will place a byte in subsequent memory addresses.

## Creating a Translation Table

 Thus, we can place the complete translation table for the LHS of the 7-segment display into RAM starting at address 50, as follows:



Note that there is no need to put both tables into memory, since the RHS display values can be calculated by adding 1 to the LHS display values

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#### Seven Segment Display

```
;The following code displays a numberic key from the keyboard on the 7-segment display
loop:
 in 00 ; kbd input put in al
sub al, 30 ; subtract ascii value of 0 to get numberic value of key
 add al, 50 ; index into the translation table by adding the numberic kbd key
               ; value to the base address of the translation table.
 mov al, [al] ; copy the 7-seg code at address given in al into al
              ; write to 7-seg display
 out 02
 jmp loop
org 50
DB FA ; 7-seg code for symbol for 0
DB 0A ; 7-seg code for symbol for 1
DB B6 ; 7-seg code for symbol for 2
DB 9E ; 7-seg code for symbol for 3
DB 4E ; 7-seg code for symbol for 4
DB DC ; 7-seg code for symbol for 5
DB FC ; 7-seg code for symbol for 6
DB 8A ; 7-seg code for symbol for
DB FE; 7-seg code for symbol for 8
DB CE; 7-seg code for symbol for 9
end
```