

MIPS ASSEMBLY

LECTURE 08-1

JIM FIX, REED COLLEGE CS2-S20

TODAY'S PLAN

- ▶ ON-LINE COURSE LOGISTICS
- ▶ REVIEW OF MIPS ASSEMBLY, BUT DONE VIA...
- ▶ ... MY SOLUTIONS TO HOMEWORK 07
- ▶ MIPS MEMORY OPERATIONS (WITH EXERCISES)

SOLUTION TO HOMEWORK 07 EXERCISE 1

- ▶ Below is the "kernel" of my MIPS code to multiply using repeated addition:

```
1. multiply:
2.      li      $t0, 0          # product = 0
3. multiply_loop:
4.      beqz    $t2, report     # if y == 0 goto report
5.      addu    $t0, $t0, $t1    # product += x
6.      addiu   $t2, $t2, -1     # y -= 1
7.      b       multiply_loop
8. report:
```

- ▶ It uses registers `t0` to compute the product, repeatedly adding `t1` to it.
- ▶ It starts with a product of 0, performs the addition at **line 5** `t2` times.

C++ SOLUTION TO HOMEWORK 07 EXERCISE 1

► Here is C++ code that mimics that MIPS code. The "kernel" loop is in **green**.

```
1. int main() { int x,y, product;  
2.     std::cin >> x;  
3.     std::cin >> y;  
4.     product = 0;  
5.     while (y != 0) {  
6.         product += x;  
7.         y--;  
8.     }  
9.     std::cout << product;  
10.    return 0;  
11. }
```

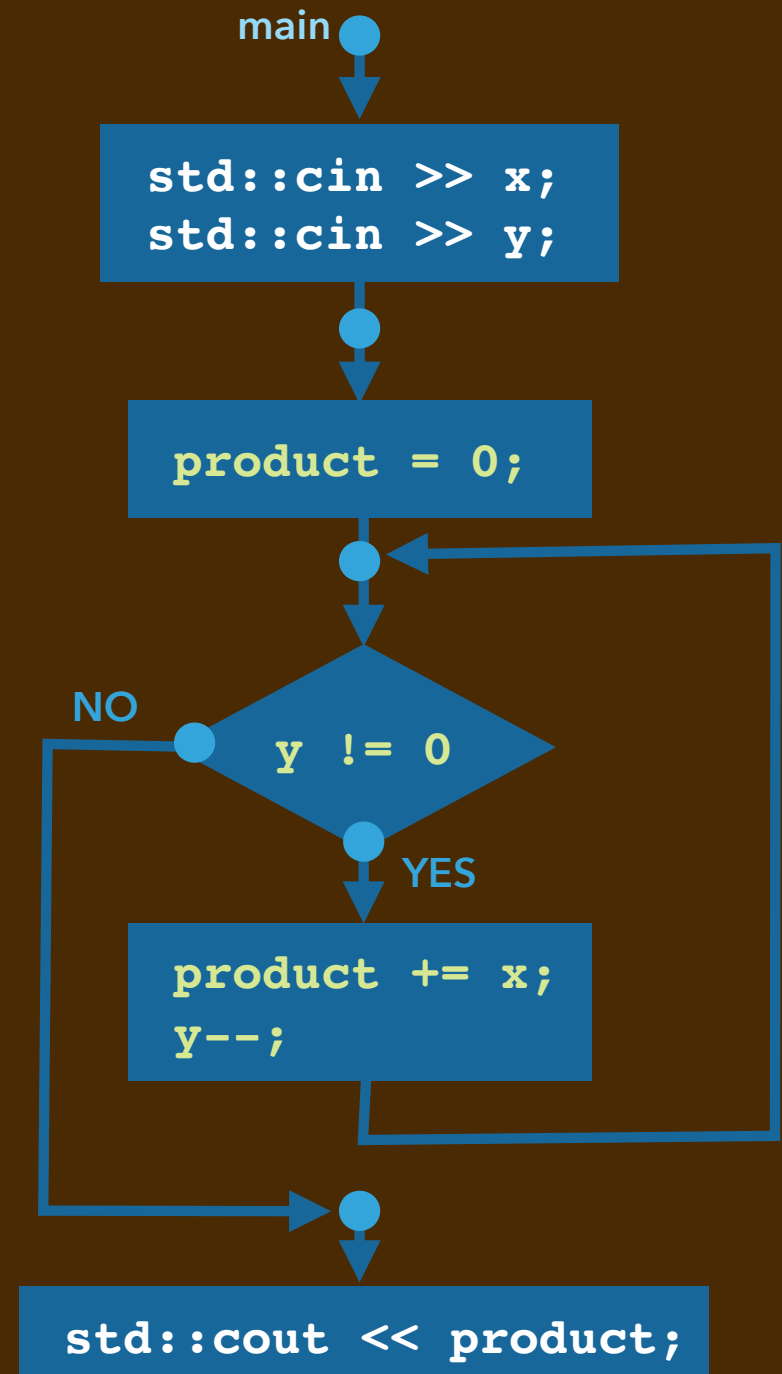
► Let's convert it to MIPS code, maybe how a compiler might...

FLOWCHART OF SOLUTION

```
1. int main() { int x,y, product;  
2.   std::cin >> x;  
3.   std::cin >> y;  
4.   product = 0;  
5.   while (y != 0) {  
6.     product += x;  
7.     y--;  
8.   }  
9.   std::cout << product;  
10.  return 0;  
11. }
```

► Here is the "flow logic" of that code:

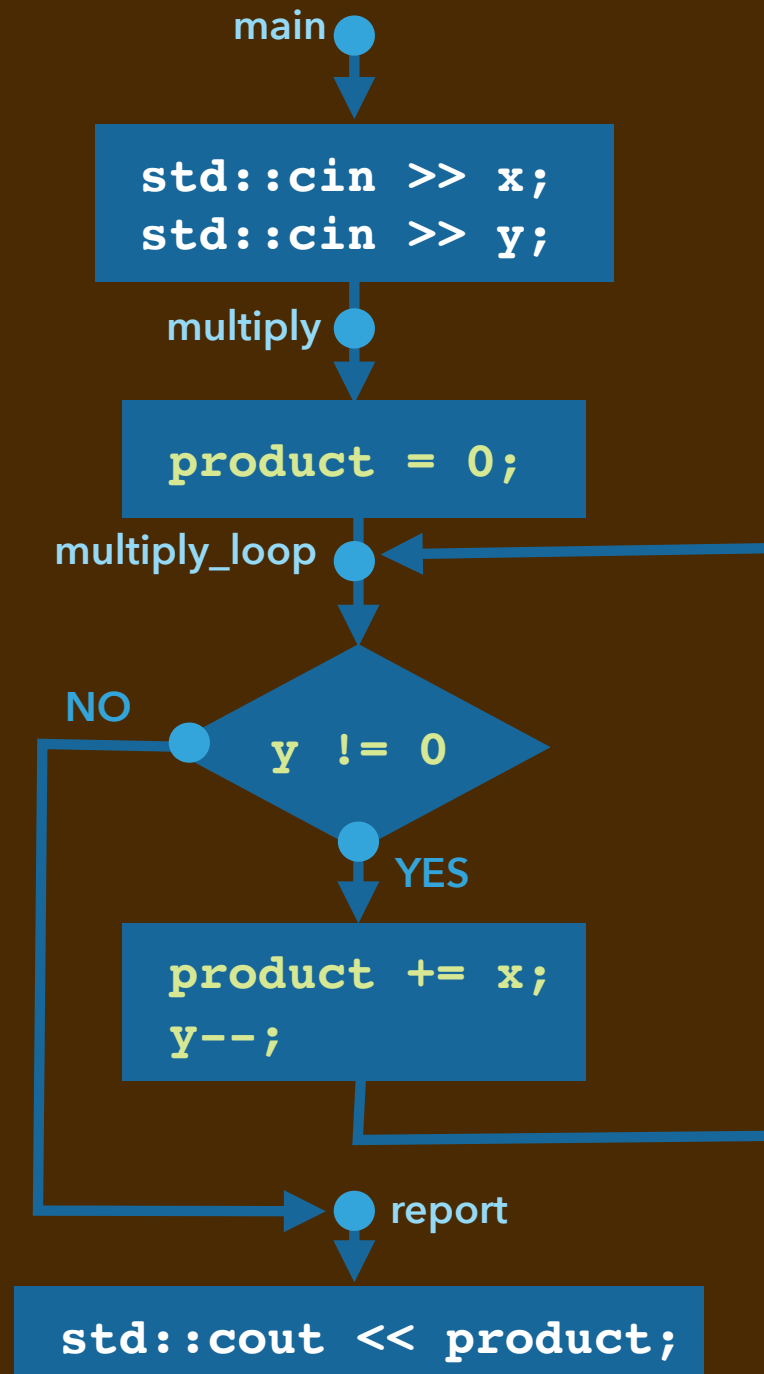
► Let's convert it to MIPS...



CONVERSION TO MIPS

```
1. main:
2.     std::cin >> x;
3.     std::cin >> y;
4. multiply:
5.     product = 0;
6. multiply_loop:
7.     while (y != 0) {
8.         product += x;
9.         y--;
10.    }
11. report:
12.    std::cout << product;
13. end_main:
14.    return 0;
```

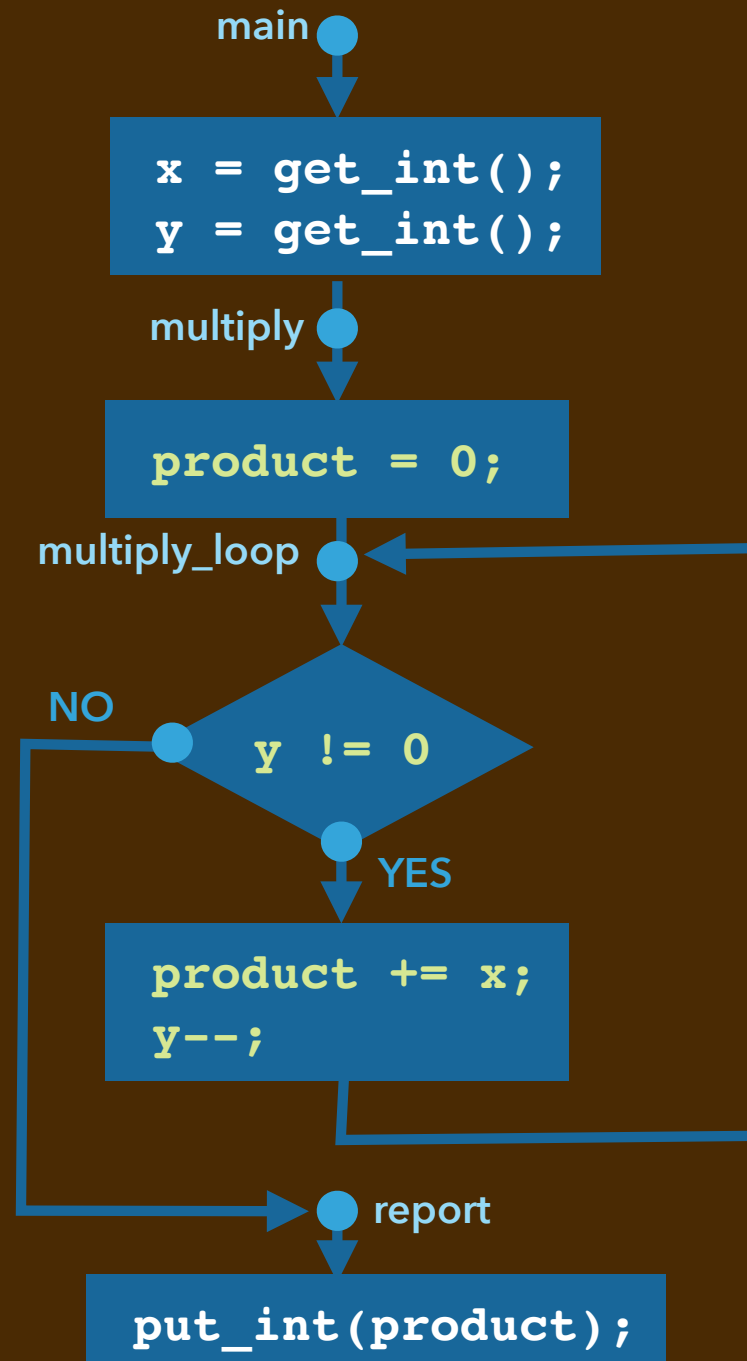
► Let's label the flow's targets.



CONVERSION TO MIPS

```
1. main:
2.     x = get_int();
3.     y = get_int();
4. multiply:
5.     product = 0;
6. multiply_loop:
7.     while (y != 0) {
8.         product += x;
9.         y--;
10.    }
11. report:
12.    put_int(product);
13. end_main:
14.    return 0;
```

► Substitute system calls.



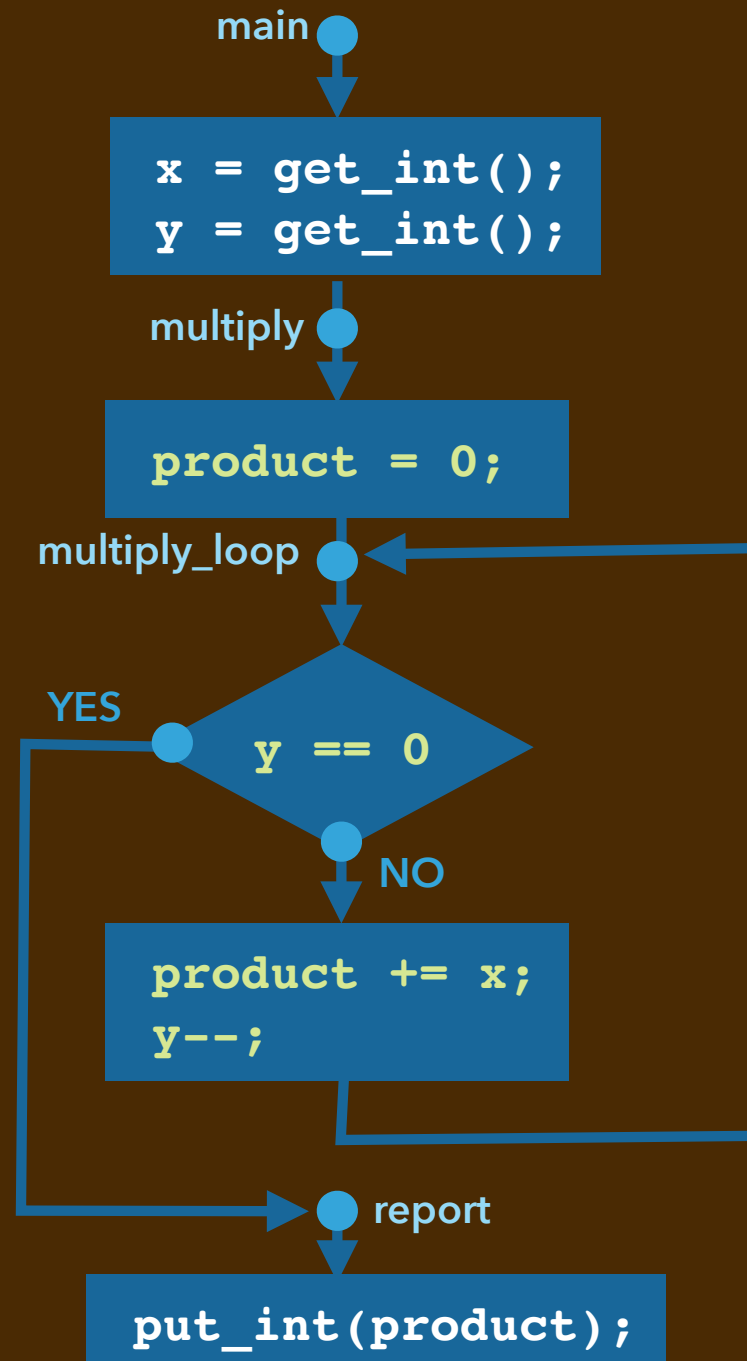
CONVERSION TO MIPS

```

1. main:
2.     x = get_int();
3.     y = get_int();
4. multiply:
5.     product = 0;
6. multiply_loop:
7.     if (y == 0) goto report;
8.     product += x;
9.     y--;
10.    goto multiply_loop;
11. report:
12.    put_int(product);
13. end_main:
14.    return 0;

```

- Replace "structured" conditional statements with GOTOs.



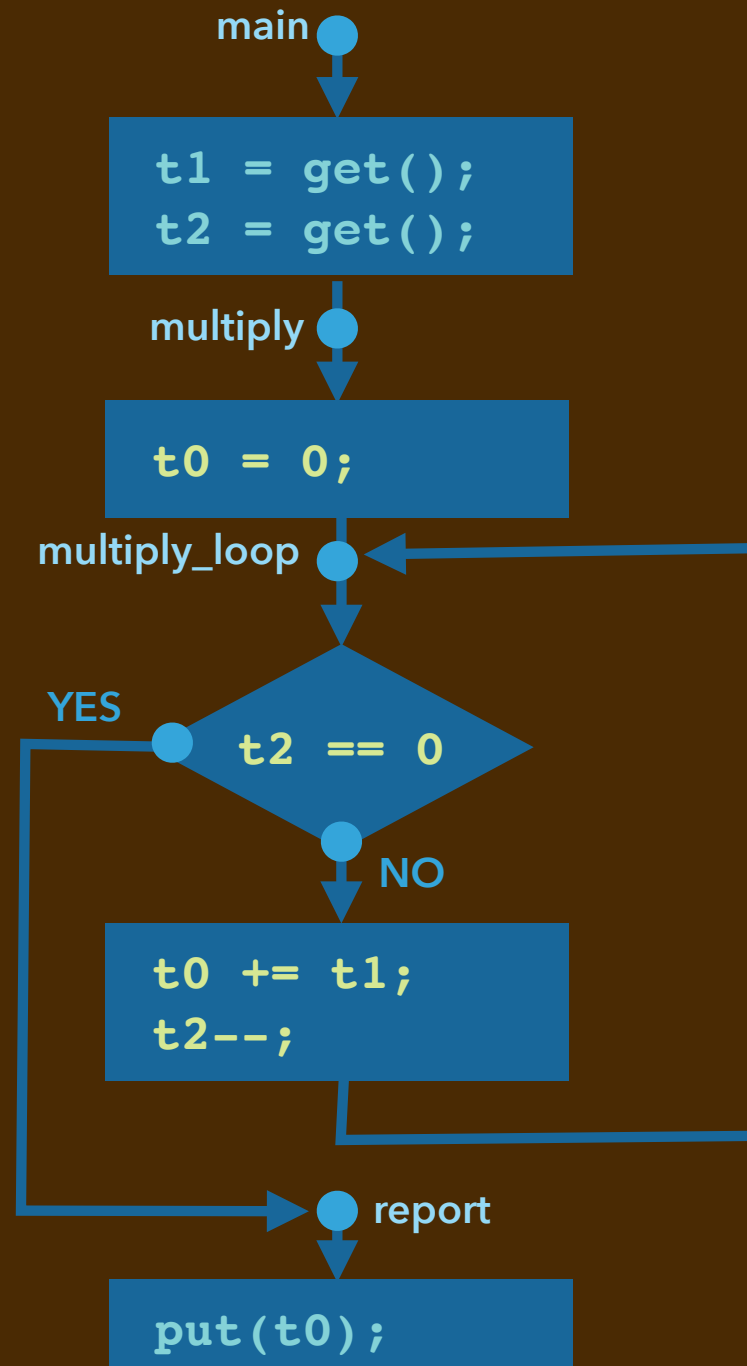
CONVERSION TO MIPS

```

1. main:
2.     t1 = get_int();
3.     t2 = get_int();
4. multiply:
5.     t0 = 0;
6. multiply_loop:
7.     if (t2 == 0) goto report;
8.     t0 += x;
9.     t1--;
10.    goto multiply_loop;
11. report:
12.    put_int(t0);
13. end_main:
14.    return 0;

```

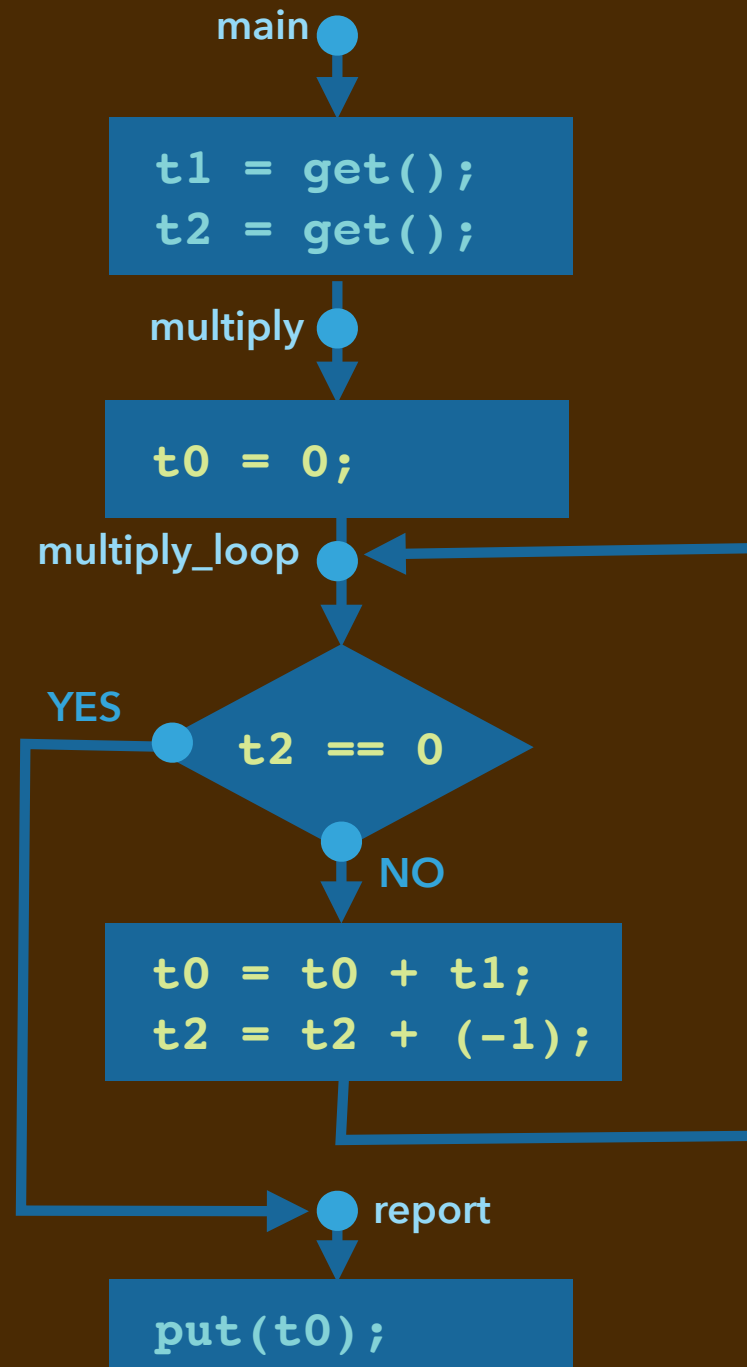
► Choose registers.



CONVERSION TO MIPS

```
1. main:
2.     t1 = get_int();
3.     t2 = get_int();
4. multiply:
5.     t0 = 0;
6. multiply_loop:
7.     if (t2 == 0) goto report;
8.     t0 = t0 + t1;
9.     t1 = t1 + (-1);
10.    goto multiply_loop;
11. report:
12.    put_int(t0);
13. end_main:
14.    return 0;
```

- Convert assignments to machine-like statements.



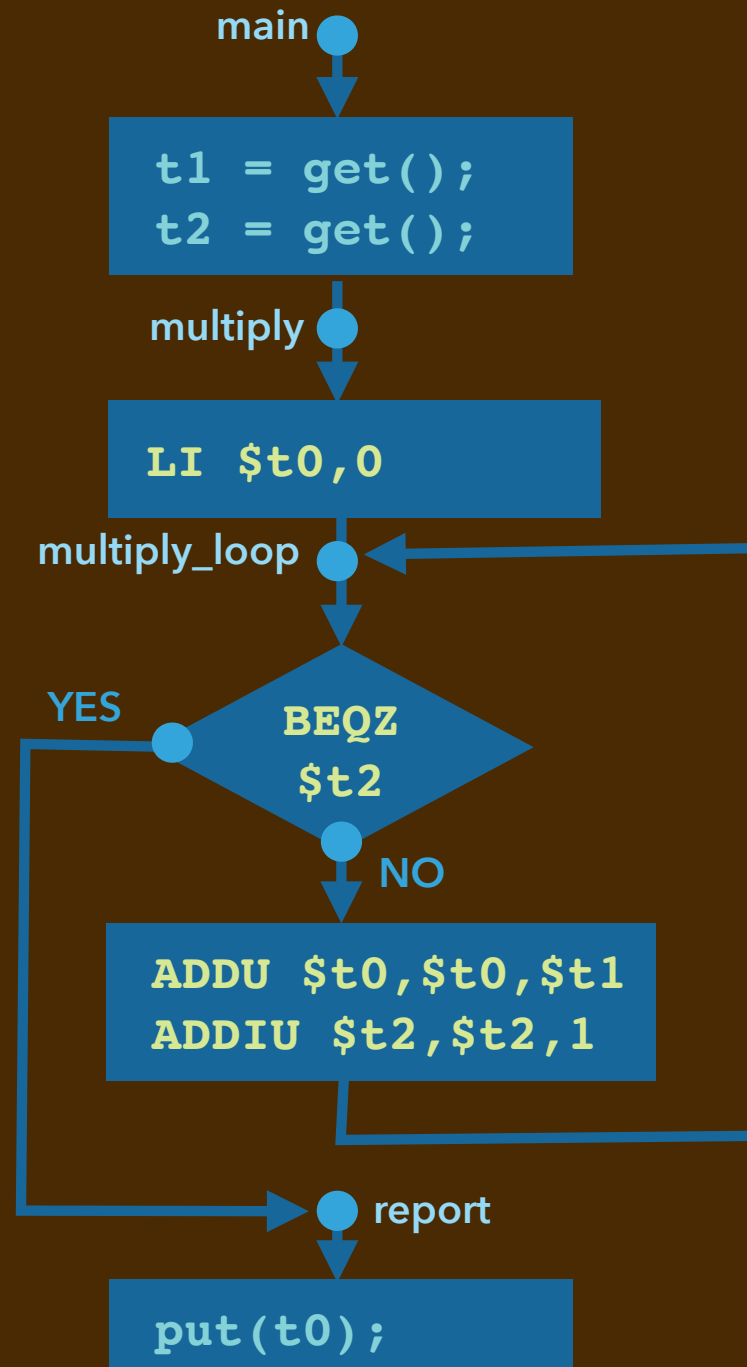
CONVERSION TO MIPS

```

1. main:
2.     t1 = get_int();
3.     t2 = get_int();
4. multiply:
5.     LI $t0,0
6. multiply_loop:
7.     BEQZ $t2,report
8.     ADDU $t0,$t0,$t1
9.     ADDIU $t2,$t2,-1
10.    B multiply_loop
11. report:
12.    put_int(t0);
13. end_main:
14.    return 0;

```

- Change C statements to MIPS instructions.



LET'S REVIEW SOME OF THESE INSTRUCTIONS...

```
1. multiply:
2.     LI $t0,0
3. multiply_loop:
4.     BEQZ $t2,report
5.     ADDU $t0,$t0,$t1
6.     ADDIU $t2,$t2,1
7.     B multiply_loop
8. report:
```

LET'S REVIEW SOME OF THESE INSTRUCTIONS...

```
1. multiply:
2.    LI $t0,0
3. multiply_loop:
4.    BEQZ $t2,report
5.    ADDU $t0,$t0,$t1
6.    ADDIU $t2,$t2,-1
7.    B multiply_loop
8. report:
```

LOAD IMMEDIATE

► Format:

LI *destination register, value*

► Meaning: load a register with a specific value.

► NOTES:

- destination can be any MIPS register
- value must be a 32-bit constant, including negative values using 2's complement encoding
- "immediate" because bits of value are in the instruction's code

LET'S REVIEW SOME OF THESE INSTRUCTIONS...

```
1. multiply:
2.     LI $t0,0
3. multiply_loop:
4.     BEQZ $t2,report
5.     ADDU $t0,$t0,$t1
6.     ADDIU $t2,$t2,-1
7.     B multiply_loop
8. report:
```

ADD (UNTRAPPED)

► Format:

ADDU *destination, source1, source2*

► Meaning: add the contents of two registers, store sum the sum in a register

► NOTES:

- destination can be any MIPS register
- sources can be any registers
- "untrapped" means errors (e.g. overflow) can be ignored by the machine
- there's SUBU and bitwise AND, OR, XOR also

LET'S REVIEW SOME OF THESE INSTRUCTIONS...

```
1. multiply:
2.     LI $t0,0
3. multiply_loop:
4.     BEQZ $t2,report
5.     ADDU $t0,$t0,$t1
6.     ADDIU $t2,$t2,-1
7.     B multiply_loop
8. report:
```

ADD IMMEDIATE (UNTRAPPED)

► Format:

ADDIU *destination, source, value*

► Meaning: compute the some of a register's contents to a value, store the sum in a register

► NOTES:

- source and destination can be any MIPS registers
- value is a 16-bit constant, including negative values. Assumes 2's complement encoding.
- (there's no SUBIU instruction)

LET'S REVIEW SOME OF THESE INSTRUCTIONS...

```
1. multiply:
2.     LI $t0,0
3. multiply_loop:
4.     BEQZ $t2,report
5.     ADDU $t0,$t0,$t1
6.     ADDIU $t2,$t2,-1
7.     B multiply_loop
8. report:
```

BRANCH IF EQUAL TO ZERO


► Format:

BEQZ *register, target-label*

► Meaning: go to the labeled instruction if a register's value is zero. Continue below if not.

LET'S REVIEW SOME OF THESE INSTRUCTIONS...

```
1. multiply:
2.    LI $t0,0
3. multiply_loop:
4.    BEQZ $t2,report
5.    ADDU $t0,$t0,$t1
6.    ADDIU $t2,$t2,-1
7.    B multiply_loop
8. report:
```



BRANCH IF EQUAL TO ZERO


► Format:

BEQZ *register, target-label*

► Meaning: go to the labeled instruction **if a register's value is zero**. Continue below if not.

LET'S REVIEW SOME OF THESE INSTRUCTIONS...

```
1. multiply:
2.    LI $t0,0
3. multiply_loop:
4.    BEQZ $t2,report
5.    ADDU $t0,$t0,$t1
6.    ADDIU $t2,$t2,-1
7.    B multiply_loop
8. report:
```



BRANCH IF EQUAL TO ZERO

► Format:

BEQZ *register, target-label*

► Meaning: go to the labeled instruction **if** a register's value is zero. Continue below **if not**.

LET'S REVIEW SOME OF THESE INSTRUCTIONS...

```
1. multiply:
2.    LI $t0,0
3. multiply_loop:
4.    BEQZ $t2,report
5.    ADDU $t0,$t0,$t1
6.    ADDIU $t2,$t2,-1
7.    B multiply_loop
8. report:
```

```
product = 0;
while (y != 0) {
    product += x;
    y--;
}
```

BRANCH IF EQUAL TO ZERO

► Format:

BEQZ *register, target-label*

► Meaning: go to the labeled instruction if a register's value is zero. Continue below if not.

► NOTE:

- if "guarding" a while loop, the **condition** specifies when the loop should *stop*, the opposite of the condition for continuing

LET'S REVIEW SOME OF THESE INSTRUCTIONS...

```
1. multiply:
2.     LI $t0,0
3. multiply_loop:
4.     BEQZ $t2,report
5.     ADDU $t0,$t0,$t1
6.     ADDIU $t2,$t2,-1
7.     B multiply_loop
8. report:
```

BRANCH IF EQUAL TO ZERO

► Format:

BEQZ *register, target-label*

► Meaning: go to the labeled instruction if a register's value is zero. Continue below if not.

► NOTES:

- branch target can be above or below
- there is BEQZ, BNEZ, BLTZ, BGTZ, BLEZ, BGEZ
- these are =, ≠, <, >, ≤, ≥ tests with 0

LET'S REVIEW SOME OF THESE INSTRUCTIONS...

```
1. multiply:
2.    LI $t0,0
3. multiply_loop:
4.    BEQZ $t2,report
5.    ADDU $t0,$t0,$t1
6.    ADDIU $t2,$t2,-1
7.    B multiply_loop
8. report:
```

BRANCH (UNCONDITIONALLY)

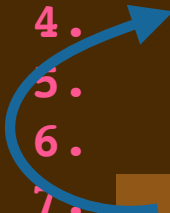
► Format:

B *target-label*

► Meaning: go to the labeled instruction

LET'S REVIEW SOME OF THESE INSTRUCTIONS...

```
1. multiply:
2.    LI $t0,0
3. multiply_loop:
4.    BEQZ $t2,report
5.    ADDU $t0,$t0,$t1
6.    ADDIU $t2,$t2,-1
7.    B multiply_loop
8. report:
```



BRANCH (UNCONDITIONALLY)

► Format:

B *target-label*

► Meaning: **go to** the labeled instruction

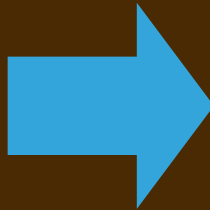
WHAT ABOUT INPUT AND OUTPUT?

- ▶ There are several different *system calls* you can make that SPIM understands.
 - get an integer input from the console
 - print an integer output to the console
 - print a null-terminated character sequence out to the console
- Each has a number that identifies it.
- You tell the system which should run by setting register v0 to that number.

SYSTEM CALLS

- ▶ Getting an integer input is *system call #5*.

```
t1 = get_int();
```



```
LI $v0,5  
syscall  
MOVE $t1,$v0
```

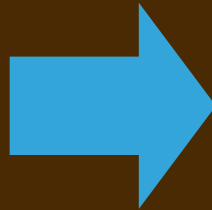
- ▶ Meaning of the MIPS code:

- We load register **v0** with the system call number.
- We make the system call.
 - A special sequence of instructions gets executed by the system to read input from the console. The integer read gets placed in register **v0**.
- We copy (i.e. "move") its value to register **t1**, our storage for the first multiplicand.

SYSTEM CALLS

- ▶ Outputting an integer is *system call #1*.

```
put_int(t0);
```



```
LI $v0,1  
MOVE $a0,$t0  
syscall
```

- ▶ Meaning of the MIPS code:

- We load register **v0** with the system call number.
- We load register **a0** with the "argument", the value we want the system to output.
 - This has us copy (again, "move") the register for the product into **a0**.
- We make the system call.
 - A sequence of instructions gets executed to output the value in **a0** to the console.

THE (POORLY NAMED) MOVE INSTRUCTION

MOVE

► Format:

MOVE *destination-register , source-register*

► Meaning:

copy the contents of the source register into the destination register

► NOTE: the value in the source register *doesn't change*

- For example the MIPS code

MOVE \$t0, \$t1

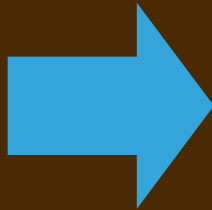
would be like the C++ code

product = x;

RETURNING FROM MAIN

- ▶ Our (main) program must return an integer value to the system.

```
return 0;
```



```
LI $v0,0  
JR $ra
```

- ▶ We'll see (on Friday) how to write functions and procedures and use them.
- ▶ These two instructions hint at how that works. In short:
 - Functions, by convention, return a value by setting register **v0** to that value.
 - A register named **ra** stores information about who called **main**.
 - This is called the **return address** of the "caller" code.
 - The **JR** instruction "jumps back" to that instruction.
- ▶ *For now, just make sure your main code always ends with these two lines.*

THE FULL SOLUTION TO HOMEWORK 07 EXERCISE 1

```
1.  main:
2.      li      $v0,5
3.      syscall
4.      move    $t1,$v0
5.      li      $v0,5
6.      syscall
7.      move    $t2,$v0
8.  multiply:
9.      li      $t0,0
10. multiply_loop:
11.     beqz     $t2,report
12.     addu     $t0,$t0,$t1
13.     addiu    $t2,$t2,-1
14.     b        multiply_loop
15. report:
16.     li      $v0, 1
17.     move     $a0, $t0
18.     syscall
19.     li      $v0, 0
20.     jr      $ra
```

THE HOMEWORK 07 EXERCISE 2 SOLUTION LOOKS SIMILAR

```
1.  main:
2.      li      $v0, 5
3.      syscall
4.      move    $t1, $v0
5.      li      $v0, 5
6.      syscall
7.      move    $t2, $v0
8.  divide:
9.      li      $t0, 0
10. divide_loop:
11.     blt      $t1, $t2, report
12.     addiu    $t0, $t0, 1
13.     subu     $t1, $t1, $t2
14.     b        divide_loop
15. report:
...    # now $t0 contains the quotient, $t1 the remainder. We output them with two system calls.
22.     li      $v0, 0
23.     jr      $ra
```

C++ SOLUTION TO HOMEWORK 07 EXERCISE 2

► Here is the C++ for it. Again, the kernel of the code is **green**.

```
1.  int main() { int number, divisor, quotient;
2.      std::cin >> number;
3.      std::cin >> divisor;
4.      quotient = 0;
5.      while (number >= divisor) {
6.          quotient++;
7.          number -= divisor;
8.      }
9.      std::cout << quotient << std::endl;
10.     std::cout << number << std::endl;
11.     return 0;
12. }
```

► It repeatedly subtracts the divisor until only a remainder is left.

► It counts the number of subtractions made. That count is the quotient.

SOME MORE INSTRUCTIONS WORTH HIGHLIGHTING

```

1.  main:
2.      li      $v0, 5
3.      syscall
4.      move    $t1, $v0
5.      li      $v0, 5
6.      syscall
7.      move    $t2, $v0
8.  divide:
9.      li      $t0, 0
10. divide_loop:
11.     blt     $t1, $t2, report
12.     addiu   $t0, $t0, 1
13.     subu    $t1, $t1, $t2
14.     b       divide_loop
15. report:
...   # output $t0 and $t1
22.    li      $v0, 0
23.    jr      $ra

```

BRANCH IF LESS THAN

► Format:

BLT *source1, source2, target-label*

► Meaning: go to the labeled instruction if the first source register is less than the second. Continue below if not.

► NOTES:

- there is BLT, BGT, BLE, BGE, BEQ, BNE
- these are $<$, $>$, \leq , \geq , $=$, \neq

SOME MORE INSTRUCTIONS WORTH HIGHLIGHTING

```
1.  main:
2.      li      $v0, 5
3.      syscall
4.      move    $t1, $v0
5.      li      $v0, 5
6.      syscall
7.      move    $t2, $v0
8.  divide:
9.      li      $t0, 0
10. divide_loop:
11.     blt     $t1, $t2, report
12.     addiu   $t0, $t0, 1
13.     subu    $t1, $t1, $t2
14.     b       divide_loop
15. report:
...    # output $t0 and $t1
22.     li     $v0, 0
23.     jr      $ra
```

SUBTRACT (UNTRAPPED)

► Format:

SUBU *destination, source1, source2*

► Meaning: subtract source2 from source1, store the difference in a register

► NOTES:

- destination is a register
- sources are registers
- "untrapped" means errors (e.g. borrow) can be ignored by the machine
- does not modify the source registers

THERE'S MORE THAN JUST CODE THERE

```
1.      # quotient.asm
2.      #
3.      # This code asks for two integers and then
4.      # outputs the quotient and remainder due to
5.      # their division.
6.
7.      .data
8.  prompt_num:  .ascii "Enter an integer: "
9.  prompt_div:  .ascii "Enter an integer divisor: "
10. fdbk1:       .ascii "Dividing "
11. fdbk2:       .ascii " by "
12. fdbk3:       .ascii " yields a quotient of "
13. fdbk4:       .ascii " with a remainder of "
14. fdbk5:       .ascii ".\n"
15.
16.      .globl main
17.      .text
18. main:
19.      # Key:
20.      #      $t0 - the number to be divided
21.      #      $t1 - the divisor
22.      #      $t2 - the remainder
23.      #      $t3 - the quotient, a count
24.      li      $v0, 5
25.      ...
```

SEGMENTS IN A PROGRAM MEMORY IMAGE

- ▶ In MIPS assembly we specify the contents of a program's **text** and **data segments**:
 - **.text** contains the bytes of the executable code of the program
 - **.data** contains additional sequences of bytes, and for various types of values
 - ✦ SPIM loads all the bytes of a program's image, both the text and the data.
- ▶ For example, "output100.asm" holds two strings and an integer in its data:

```
        .data
report:  .asciiz "The value held in memory is "
dot_eoln: .asciiz ".\n"
value:   .word 100
        .globl main
        .text

main:
    la    $a0, report
    ...
```

SEGMENTS IN A PROGRAM MEMORY IMAGE

► Here is the full "output100.asm" program:

```
1.      .data
2.  report:  .asciiz "The value held in memory is "
3.  dot_eoln: .asciiz ".\n"
4.  value:   .word 100
5.          .globl main
6.          .text
7.  main:
8.      la    $a0,report
9.      li    $v0,4
10.     syscall
11.     la    $t0,value
12.     lw    $a0,($t0)
13.     li    $v0,1
14.     syscall
15.     la    $a0,eoln
16.     li    $v0,4
17.     syscall
18.     li    $v0,0
19.     jr    $ra
```

LOADING A VALUE STORED AT AN ADDRESS

► Here is the full "output100.asm" program:

```
1.      .data
2.  report:  .asciiz "The value held in memory is "
3.  dot_eoln: .asciiz ".\n"
4.  value:   .word 100
5.          .globl main
6.          .text
7.  main:
8.      la    $a0,report
9.      li    $v0,4
10.     syscall
11.     la    $t0,value
12.     lw    $a0,($t0)
13.     li    $v0,1
14.     syscall
15.     la    $a0,eoln
16.     li    $v0,4
17.     syscall
18.     li    $v0,0
19.     jr    $ra
```

LOAD ADDRESS INTO A REGISTER

► Format:

LA *destination*, *program-label*

► Meaning: loads a register with the address of a labelled item in the code

► NOTE: For a string (.asciiz), this is the address of the first letter of its sequence.

LOADING A VALUE STORED AT AN ADDRESS

► Here is the full "output100.asm" program:

```
1.      .data
2.  report:  .asciiz "The value held in memory is "
3.  dot_eoln: .asciiz ".\n"
4.  value:   .word 100
5.          .globl main
6.          .text
7.  main:
8.      la    $a0,report
9.      li    $v0,4
10.     syscall
11.     la    $t0,value
12.     lw    $a0,($t0)
13.     li    $v0,1
14.     syscall
15.     la    $a0,eoln
16.     li    $v0,4
17.     syscall
18.     li    $v0,0
19.     jr    $ra
```

LOAD ADDRESS INTO A REGISTER

► Format:

LA *destination*, *program-label*

► Meaning: loads a register with the address of a labelled item in the code

► NOTE: For an integer (.word), this is the address of the first of its four bytes.

LOADING A VALUE STORED AT AN ADDRESS

► Here is the full "output100.asm" program:

```
1.      .data
2.  report:  .asciiz "The value held in memory is "
3.  dot_eoln: .asciiz ".\n"
4.  value:   .word 100
5.
6.      .globl main
7.      .text
8.  main:
9.      la    $a0, report
10.     li    $v0, 4
11.     syscall
12.     la    $t0, value
13.     lw    $a0, ($t0)
14.     li    $v0, 1
15.     syscall
16.     la    $a0, dot_eoln
17.     li    $v0, 4
18.     syscall
19.     li    $v0, 0
20.     jr    $ra
```

LOAD A VALUE STORED AT AN ADDRESS

► Format:

LW *destination*, (*source*)

► Meaning: loads a register with the (4-byte word) value stored at an address

► NOTE: *source* is a register that holds the address., i.e. a pointer to the loaded data.

LOADING A VALUE FROM MEMORY

LOAD A (FOUR BYTE) VALUE FROM AN ADDRESS IN MEMORY

LW *destination*, (*source*)

►NOTES:

- This is sometimes called a "memory-to-register" transfer or "fetch."
- *source* is a register holding the address of the data we're fetching.
- We can think of the source register as a pointer. So **LW** is like the C code:

```
int* source_pointer;  
...  
int destination_value = *source_pointer;
```

REGISTERS VERSUS MEMORY

- ▶ A MIPS processor has only 32 registers for storing calculated values.
 - It can also access a large memory using addresses.
- ▶ MIPS is a *"load/store" computer architecture*.
 - It can only perform calculations on data stored in its registers
 - To modify a value stored in memory it must:
 1. *load* a value from memory by its address into a register
 2. modify that value, computing a new value, within registers
 3. *store* that changed value into memory at that same address

EXAMPLE: INCREMENTING A VALUE IN MEMORY

► Here is MIPS code that adds 10 to a location in memory:

```
1.  LA      $a0,value
2.  LW      $t0,($a0)
3.  ADDIU   $t0,$t0,10
4.  SW      $t0,($a0)
```

STORING A VALUE TO MEMORY

STORE A (FOUR BYTE) VALUE TO AN ADDRESS IN MEMORY

SW *source*, (*destination*)

►NOTES:

- This is sometimes called a "register-to-memory" transfer.
- *destination* is a register holding the address where we're storing the data in the *source* register

WORDS VERSUS BYTES

- ▶ Recall that the fundamental quantity manipulated by a MIPS32 instruction set is 32-bits long.
 - Registers hold 32 bits, i.e 4 bytes, of data.
 - Integers are 32 bits, i.e. 4 bytes, long.
 - Addresses are 32 bits, i.e. 4 bytes, long.
- ▶ So in MIPS32, a **word** is 32 bits, i.e. 4 bytes.
 - The **LW** instruction reads 4 bytes of data from memory.
 - The **SW** instruction write 4 bytes of data out to memory.
- ▶ There are also the **LB** and **SB** instructions for accessing a *single byte* in memory.
 - NOTE: strings are sequences of characters, each character is a byte of data.

LOADING A BYTE AND STORING A BYTE

LOAD A BYTE'S VALUE FROM AN ADDRESS IN MEMORY

LB *destination*, (*source*)

►NOTES:

- This is sometimes called a "memory-to-register" transfer or a "fetch."
- *source* is a register holding the address where we're fetching the data in *destination*

STORE A BYTE'S VALUE TO AN ADDRESS IN MEMORY

SB *source*, (*destination*)

►NOTES:

- This is sometimes called a "register-to-memory" transfer.
- *destination* is a register holding the address where we're storing the data of *source*

PRINTING STRINGS AND CHANGING STRINGS ("HELLOBYE.ASM")

```
1.  .data
2.  hello_ptr: .asciiz "hello\n"
3.  .globl main
4.  .text
5.  main:
6.  print_hello:
7.      li $v0, 4
8.      la $a0, hello_ptr
9.      syscall
10.     j     change_string
11. print_bye:
12.     li $v0, 4
13.     la $a0, hello_ptr
14.     syscall
15. end_main:
16.     li $v0, 0
17.     jr $ra

18. change_string:
19.     la $t0, hello_ptr
20.     li $t1, 'b'
21.     sb $t1, ($t0)
22.     addiu $t0, $t0, 1
23.     li $t1, 'y'
24.     sb $t1, ($t0)
25.     addiu $t0, $t0, 1
26.     li $t1, 'e'
27.     sb $t1, ($t0)
28.     addiu $t0, $t0, 1
29.     li $t1, '\n'
30.     sb $t1, ($t0)
31.     addiu $t0, $t0, 1
32.     li $t1, 0
33.     sb $t1, ($t0)
34.     j print_bye
```

PRINTING STRINGS AND CHANGING STRINGS ("HELLOBYE.ASM")

```
1.  .data
2.  hello_ptr: .asciiz "hello\n"
3.  .globl main
4.  .text
5.  main:
6.  print_hello:
7.    li $v0, 4
8.    la $a0, hello_ptr
9.    syscall
10.   j    change_string
11. print_bye:
12.   li $v0, 4
13.   la $a0, hello_ptr
14.   syscall
15. end_main:
16.   li $v0, 0
17.   jr $ra
18. change_string:
19.   la $t0, hello_ptr
20.   li $t1, 'b'
21.   sb $t1, ($t0)
22.   addiu $t0, $t0, 1
23.   li $t1, 'y'
24.   sb $t1, ($t0)
25.   addiu $t0, $t0, 1
26.   li $t1, 'e'
27.   sb $t1, ($t0)
28.   addiu $t0, $t0, 1
29.   li $t1, '\n'
30.   sb $t1, ($t0)
31.   addiu $t0, $t0, 1
32.   li $t1, 0
33.   sb $t1, ($t0)
34.   j print_bye
```

► The code jumps around to fit neatly on this slide...


PRINTING STRINGS AND CHANGING STRINGS ("HELLOBYE.ASM")

```
1.  .data
2.  hello_ptr: .asciiz "hello\n"
3.  .globl main
4.  .text
5.  main:
6.  print_hello:
7.      li $v0, 4
8.      la $a0, hello_ptr
9.      syscall
10. j    change_string
11. print_bye:
12.     li $v0, 4
13.     la $a0, hello_ptr
14.     syscall
15. end_main:
16.     li $v0, 0
17.     jr $ra
18. change_string:
19.     la $t0, hello_ptr
20.     li $t1, 'b'
21.     sb $t1, ($t0)
22.     addiu $t0, $t0, 1
23.     li $t1, 'y'
24.     sb $t1, ($t0)
25.     addiu $t0, $t0, 1
26.     li $t1, 'e'
27.     sb $t1, ($t0)
28.     addiu $t0, $t0, 1
29.     li $t1, '\n'
30.     sb $t1, ($t0)
31.     addiu $t0, $t0, 1
32.     li $t1, 0
33.     sb $t1, ($t0)
34.     j print_bye
```

► It first prints the string "hello\n" in lines 6-9.

PRINTING STRINGS AND CHANGING STRINGS ("HELLOBYE.ASM")

```
1.  .data
2.  hello_ptr: .asciiz "hello\n"
3.  .globl main
4.  .text
5.  main:
6.  print_hello:
7.      li $v0, 4
8.      la $a0, hello_ptr
9.      syscall
10. j    change_string
11. print_bye:
12.     li $v0, 4
13.     la $a0, hello_ptr
14.     syscall
15. end_main:
16.     li $v0, 0
17.     jr $ra
18. change_string:
19.     la $t0, hello_ptr
20.     li $t1, 'b'
21.     sb $t1, ($t0)
22.     addiu $t0, $t0, 1
23.     li $t1, 'y'
24.     sb $t1, ($t0)
25.     addiu $t0, $t0, 1
26.     li $t1, 'e'
27.     sb $t1, ($t0)
28.     addiu $t0, $t0, 1
29.     li $t1, '\n'
30.     sb $t1, ($t0)
31.     addiu $t0, $t0, 1
32.     li $t1, 0
33.     sb $t1, ($t0)
34.     j print_bye
```

A blue arrow originates from the instruction 'j change_string' on line 10 and points to the label 'change_string:' on line 18, illustrating a jump instruction.

► It then jumps to line 18.

PRINTING STRINGS AND CHANGING STRINGS ("HELLOBYE.ASM")


```
1.  .data
2.  hello_ptr: .asciiz "hello\n"
3.  .globl main
4.  .text
5.  main:
6.  print_hello:
7.    li $v0, 4
8.    la $a0, hello_ptr
9.    syscall
10.   j    change_string
11. print_bye:
12.   li $v0, 4
13.   la $a0, hello_ptr
14.   syscall
15. end_main:
16.   li $v0, 0
17.   jr $ra
```

```
18. change_string:
19.   la $t0, hello_ptr
20.   li $t1, 'b'
21.   sb $t1, ($t0)
22.   addiu $t0, $t0, 1
23.   li $t1, 'y'
24.   sb $t1, ($t0)
25.   addiu $t0, $t0, 1
26.   li $t1, 'e'
27.   sb $t1, ($t0)
28.   addiu $t0, $t0, 1
29.   li $t1, '\n'
30.   sb $t1, ($t0)
31.   addiu $t0, $t0, 1
32.   li $t1, 0
33.   sb $t1, ($t0)
34.   j print_bye
```

► In Lines 19-30, it overwrites the bytes of the string with "bye\n"

PRINTING STRINGS AND CHANGING STRINGS ("HELLOBYE.ASM")

```
1.  .data
2.  hello_ptr: .asciiz "hello\n"
3.  .globl main
4.  .text
5.  main:
6.  print_hello:
7.      li $v0, 4
8.      la $a0, hello_ptr
9.      syscall
10.     j  change_string
11. print_bye:
12.     li $v0, 4
13.     la $a0, hello_ptr
14.     syscall
15. end_main:
16.     li $v0, 0
17.     jr $ra
18. change_string:
19.     la $t0, hello_ptr
20.     li $t1, 'b'
21.     sb $t1, ($t0)
22.     addiu $t0, $t0, 1
23.     li $t1, 'y'
24.     sb $t1, ($t0)
25.     addiu $t0, $t0, 1
26.     li $t1, 'e'
27.     sb $t1, ($t0)
28.     addiu $t0, $t0, 1
29.     li $t1, '\n'
30.     sb $t1, ($t0)
31.     addiu $t0, $t0, 1
32.     li $t1, 0
33.     sb $t1, ($t0)
34.     j print_bye
```



► It then jumps to line 11 to continue the work of **main**.

PRINTING STRINGS AND CHANGING STRINGS ("HELLOBYE.ASM")

```
1.  .data
2.  hello_ptr: .asciiz "hello\n"
3.  .globl main
4.  .text
5.  main:
6.  print_hello:
7.      li $v0, 4
8.      la $a0, hello_ptr
9.      syscall
10.     j     change_string
11. print_bye:
12.     li $v0, 4
13.     la $a0, hello_ptr
14.     syscall
15. end_main:
16.     li $v0, 0
17.     jr $ra
18. change_string:
19.     la $t0, hello_ptr
20.     li $t1, 'b'
21.     sb $t1, ($t0)
22.     addiu $t0, $t0, 1
23.     li $t1, 'y'
24.     sb $t1, ($t0)
25.     addiu $t0, $t0, 1
26.     li $t1, 'e'
27.     sb $t1, ($t0)
28.     addiu $t0, $t0, 1
29.     li $t1, '\n'
30.     sb $t1, ($t0)
31.     addiu $t0, $t0, 1
32.     li $t1, 0
33.     sb $t1, ($t0)
34.     j print_bye
```

► It prints the string at that same address using lines 11-14.

PRINTING STRINGS AND CHANGING STRINGS ("HELLOBYE.ASM")

```
1.  .data
2.  hello_ptr: .asciiz "hello\n"
3.  .globl main
4.  .text
5.  main:
6.  print_hello:
7.      li $v0, 4
8.      la $a0, hello_ptr
9.      syscall
10.     j     change_string
11. print_bye:
12.     li $v0, 4
13.     la $a0, hello_ptr
14.     syscall
15. end_main:
16.     li $v0, 0
17.     jr $ra

18. change_string:
19.     la $t0, hello_ptr
20.     li $t1, 'b'
21.     sb $t1, ($t0)
22.     addiu $t0, $t0, 1
23.     li $t1, 'y'
24.     sb $t1, ($t0)
25.     addiu $t0, $t0, 1
26.     li $t1, 'e'
27.     sb $t1, ($t0)
28.     addiu $t0, $t0, 1
29.     li $t1, '\n'
30.     sb $t1, ($t0)
31.     addiu $t0, $t0, 1
32.     li $t1, 0
33.     sb $t1, ($t0)
34.     j print_bye
```

► But the contents of the string have changed because of 5 **SB** instructions.

PRINTING STRINGS AND CHANGING STRINGS ("HELLOBYE.ASM")

```
1.  .data
2.  hello_ptr: .asciiz "hello\n"
3.  .globl main
4.  .text
5.  main:
6.  print_hello:
7.      li $v0, 4
8.      la $a0, hello_ptr
9.      syscall
10.     j     change_string
11. print_bye:
12.     li $v0, 4
13.     la $a0, hello_ptr
14.     syscall
15. end_main:
16.     li $v0, 0
17.     jr $ra
18. change_string:
19.     la $t0, hello_ptr
20.     li $t1, 'b'
21.     sb $t1, ($t0)
22.     addiu $t0, $t0, 1
23.     li $t1, 'y'
24.     sb $t1, ($t0)
25.     addiu $t0, $t0, 1
26.     li $t1, 'e'
27.     sb $t1, ($t0)
28.     addiu $t0, $t0, 1
29.     li $t1, '\n'
30.     sb $t1, ($t0)
31.     addiu $t0, $t0, 1
32.     li $t1, 0 -
33.     sb $t1, ($t0)
34.     j print_bye
```

► Notice we set the 5th character to 0. This *null terminates* that string.

SUMMARY: LOADING FROM AND STORING TO MEMORY

LOAD A WORD FROM MEMORY

LW *destination* , (*source*)

STORE A WORD TO MEMORY

SW *source* , (*destination*)

LOAD A BYTE FROM MEMORY

LB *destination* , (*source*)

STORE A BYTE TO MEMORY

SB *source* , (*destination*)

(OPTIONAL) EXERCISES FOR WEDNESDAY

- ▶ Change "hello.asm" so that it instead outputs this to the console

```
hello
ello
llo
lo
o
```

- Each line starts from a different place in the character string.

- ▶ Change "hello.asm" so that it instead outputs this to the console

```
hello
elloh
llohe
lohel
ohell
```

- Each line is a result of rotating the string's contents by one character.

- ▶ BONUS: have each instead operate on a string read as input from the console.

- See "string.asm" to see how to input a character string into memory.