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**Question 0.0.1.** *Kyle: What is the intuition behind the discretization of the box size and why is equivalent to the product definition ( $|B| = \prod_i |I_i|$ ).*

**Answer.** *It will only hold in the case of a box. Recalling the definition of a box  $B = I_1 \times \dots \times I_n$ , if  $I_i = [a_i, b_i]$  we have*

$$|B| = \prod_i |I_i| = \prod_i (b_i - a_i).$$

*Focusing on the single dimensional case (which generalizes via cross products), take  $N \in \mathbb{N}_0$  and note that  $k/N \in \mathbb{Z}/N$  is in  $I$  iff  $a \leq \frac{k}{N} \leq b$ , which is the same as  $\lceil aN \rceil \leq k \leq \lfloor bN \rfloor$ . By simple counting it follows that*

$$\left| I \cap \frac{\mathbb{Z}}{N} \right| = \lfloor bN \rfloor - \lceil aN \rceil + 1.$$

*This can be bounded to*

$$(bN - 1) - (aN + 1) + 1 \leq \lfloor bN \rfloor - \lceil aN \rceil + 1 \leq bN - (aN - 1) + 1$$

$$N(b - a) - 1 \leq \lfloor bN \rfloor - \lceil aN \rceil + 1 \leq N(b - a) + 2$$

*Therefore*

$$(b - a) - \frac{1}{N} \leq \frac{1}{N} \cdot \left| I \cap \frac{\mathbb{Z}}{N} \right| \leq (b - a) + \frac{2}{N}$$

*which in the limit gives the desired result  $|I| = b - a = \lim_{n \rightarrow \infty} \frac{1}{N} \left| I \cap \frac{\mathbb{Z}}{N} \right|$ .*

**Remark.** Charlie: What about for an elementary set  $E = \bigsqcup_i B_i$  defining its “size” as a supremum over all disjoint decompositions? For example:

$$|E| = \sup \left\{ \sum_i^N |B_i| : B_i \cap B_j = \emptyset, B_i \in E \right\}.$$

This could possibly be extended to open sets as well? Open boxes with rational corners would provide a nice basis for  $\mathbb{R}^d$ .

**Remark.** *In addition to the previous proof, we found an alternative example via an explicit construction of a sequence of simple functions  $f_n$  that converge to  $f$ . From this point forward, we will assume that every binary expansion will be*

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non-terminating as to ensure uniqueness of the expansions. Note this can be done since one can substitute the last 1 in a terminating sequence with  $0\bar{1}$ . Suppose that  $x \in [0, 1]$  has a binary expansion

$$x = 0.b_1 b_2 b_3 \dots$$

Clearly the sequence of partial sums

$$f_n(x) = \sum_{j=1}^n 2b_j 3^{-j}$$

converges pointwise to  $f$ . Thus if we can show they are simple functions, we obtain  $f$  is measurable and the previous answer's conclusion follows.

Consider the case when  $b_1$  is 1. Then possible values  $x$  can take on are those satisfying

$$\frac{1}{2} = 0.0111 \dots < x \leq 0.111 \dots = 1.$$

Thus in the place of  $b_1$  in  $f_n$ , we can substitute  $\mathbb{1}_{(\frac{1}{2}, 1]}(x)$ . When  $b_2$  is 1, then  $b_1$  can be either 0 or 1, meaning the possible values for  $x$  must satisfy either of

$$\begin{aligned} \frac{1}{4} = 0.00111 \dots &< x \leq 0.0111 \dots = \frac{1}{2} \\ \frac{3}{4} = 0.10111 \dots &< x \leq 0.1111 \dots = 1 \end{aligned}$$

Since these are disjoint intervals, we can again do a substitution of  $b_2$  in  $f_n$  for  $\mathbb{1}_{(\frac{1}{4}, \frac{1}{2}]}(x) + \mathbb{1}_{(\frac{3}{4}, 1]}(x)$ . In general, the digit  $b_i$  can be written as a sum of disjoint indicator functions. More precisely, it is the sum over all indicator functions corresponding to all the possible bit flippings of  $0.b_1 \dots b_{i-1}$ , as those bits dictate where each interval starts, giving

$$b_i = \sum_{k=1}^{2^{i-1}} \mathbb{1}_{(\frac{2k-1}{2^i}, \frac{2k}{2^i}]}(x).$$

Thus we can write in total

$$f_n(x) = \sum_{j=1}^n \frac{2}{3^j} \cdot \left( \sum_{k=1}^{2^{j-1}} \mathbb{1}_{(\frac{2k-1}{2^j}, \frac{2k}{2^j}]}(x) \right).$$

In this form it is then clear that each  $f_n$  is a simple function, which was to be shown.