10 toolal -4

S. Ree-took Koman Drugo (Sec 1) Poul: no = 08 Sem = 4th

Define defferent Asymptotic motation with

Ans Asymptotic instations are imathematical tools to represent time compliantly of algorithms your asymptotic analysis.

 $8x: 3u^3 + 6u^2 + 6000 = \Theta(u^3)$

Big O Notation: y(u) = O(g(n))if O = y(n) = O(g(n))for all n = n0ex: $3n^3 + 10n + 32 = O(n^3)$

3 Ω Notation : $y(n) = \Omega (g(n))$ if $0 < = C \cdot g(n) < = y(n)$ you all n > = no $ex : n^2 + \log n = \Omega (n^2)$.

(4) 6 mall o Notation : y(n) = o(g(n))why o < = y(n) < o(g(n))for all n > = no $ex: 3n^3 + 1on + 32 = o(n^4)$.

(5) w Notation:
$$y(n) - w(g(n))$$

with, $0 < 0 + g(n) < y(n)$

for all $n > 1/1 + n = w(n)$

(i) what should be time complexity of -

And The sequence is $2, 4, 8, \ldots, 2^{K}$.

till $0^{K} = N$

$$K = \log_2 n$$

$$T(n) = O(\log_2 n)$$

(3) T(n) = \$3T(n-1) of n>0, otherwise 4)

Ans using subotitution.

$$T(n) = 3T(n-1)$$

$$= 3^{2}T(n-2)$$

$$\vdots$$

$$= 3^{n}T(0)$$

(4) T(n) = g et (n-1)-1 if n>0, otherwise 4 9

Ans T(n) = aT(n-1)-1using subotitution,

$$T(n) = 2^{2}(T(n-2)) - 2 - 4$$

=
$$2^{3}(t(n-3)) - 2^{2} - 2 - 2^{\circ}$$

$$= 2^{h} T(0) - 2^{n-1} - 2^{n-2} \dots - 2^{1-2^{n}}$$

$$= 2^{h} - \left(2^{h-1} + 2^{h-2} \dots + 2^{n}\right)$$

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$$= 2^{h} - \left(2^{h} + 2^{h-2} \dots + 2^{h-2}\right)$$

$$= 2^{h} - \left(2^{h} + 2^{h} +$$

int ?, count = 0;

count ++;

for (== 1; ?+ ? <= ~; ?++)

Ans valves to be checked. 4^{2} , 2^{2} , 3^{2} . . . K^{2} K2 = N K = O(Ju). .. time Complexity is O(Tu). (4). Time Complexity of vould junction (unt n) ? unt u, J.K, count = 0; Jos (i= 10/2;iv=10;i++) → O(u) 90x (2=4; 2x=n; 2=2+2) -> (0(Jogh) Jos (K=1; K x=n; K= K+2) (0 (dogen)). Count tt. Outermost doop ours n/2 times. for inner do p with vocable J. 2°, 21, 22. . . 2 K 2K = h Similarity, loop with variable & ours logar times. 0 (n (log2n)2) .. time Complexity is

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(8.) Time Complexity of:
    Janation (unt n) {
            for ( == 1) return;
                      for (2=1 go m) {
                   painty ("");
        Hunction (n-3);
       The recursive for calle itself n/3 times so,
the complexity foor that is a co o (u).
Sno
       Each recurrère vall have a mested loof
       with time complexity o(u2).
       .. itotal time complication is o(13).
 (9.) Time complexity of:
       You'd function (int n) &

for (3=1 to n) &

for (5=1) 5 x= n; 5=5+?)
                      posinty (" * ")
 Ans total number of operations
       = \frac{u^2-1}{u^2-1} hostimus + \frac{y-2}{u^2-1} + \frac{y-2}{3} demis ... 4 time
                 \frac{4+\frac{n}{n-1}}{n-1} \frac{n}{2}+\frac{n}{1} \frac{1}{2} \frac{1}{1}
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=
$$n \left(\frac{1}{1} + \frac{1}{2} + \dots + \frac{1}{n} \right)$$

= $n \log n$
= $o \left(n \log n \right) + dime complexity.$

(6). for the function, nk & com what is the asymptotic rolationship botwoon those functions? Assume that K7/1 & C7/1 are constants. find out value of c & no four which ordered holds. Ans Exponent Junction genous the Jastest hence.

wk 0 0 (0")

let, K= 43 & C= 2.

for h >10, the above relation holds.

Similarly. There can be orginate number of pair a 2 no jour which above relation halds.