Machine-Level Programming III: (Switch Statements) and IA32 Procedures

Lecture 4 - 2015

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Credits to Alexandre David (AAU) and Randy Bryant & Dave O'Hallaron (CMU)

Today

- **Switch statements** *
- **IA 32 Procedures**
  - Stack Structure
  - Calling Conventions
  - Illustrations of Recursion & Pointers
- Arrays
- Structs

```
long switch_eg
   (long x, long y, long z)
{
    long w = 1;
    switch(x) {
    case 1:
        w = y*z;
        break;
    case 2:
        w = y/z;
        /* Fall Through */
    case 3:
        w += z;
        break;
    case 5:
    case 6:
        w -= z;
        break;
    default:
        w = 2;
    }
    return w;
}
```

Switch Statement Example

- **Multiple case labels**
  - Here: 5 & 6
- **Fall through cases**
  - Here: 2
- **Missing cases**
  - Here: 4

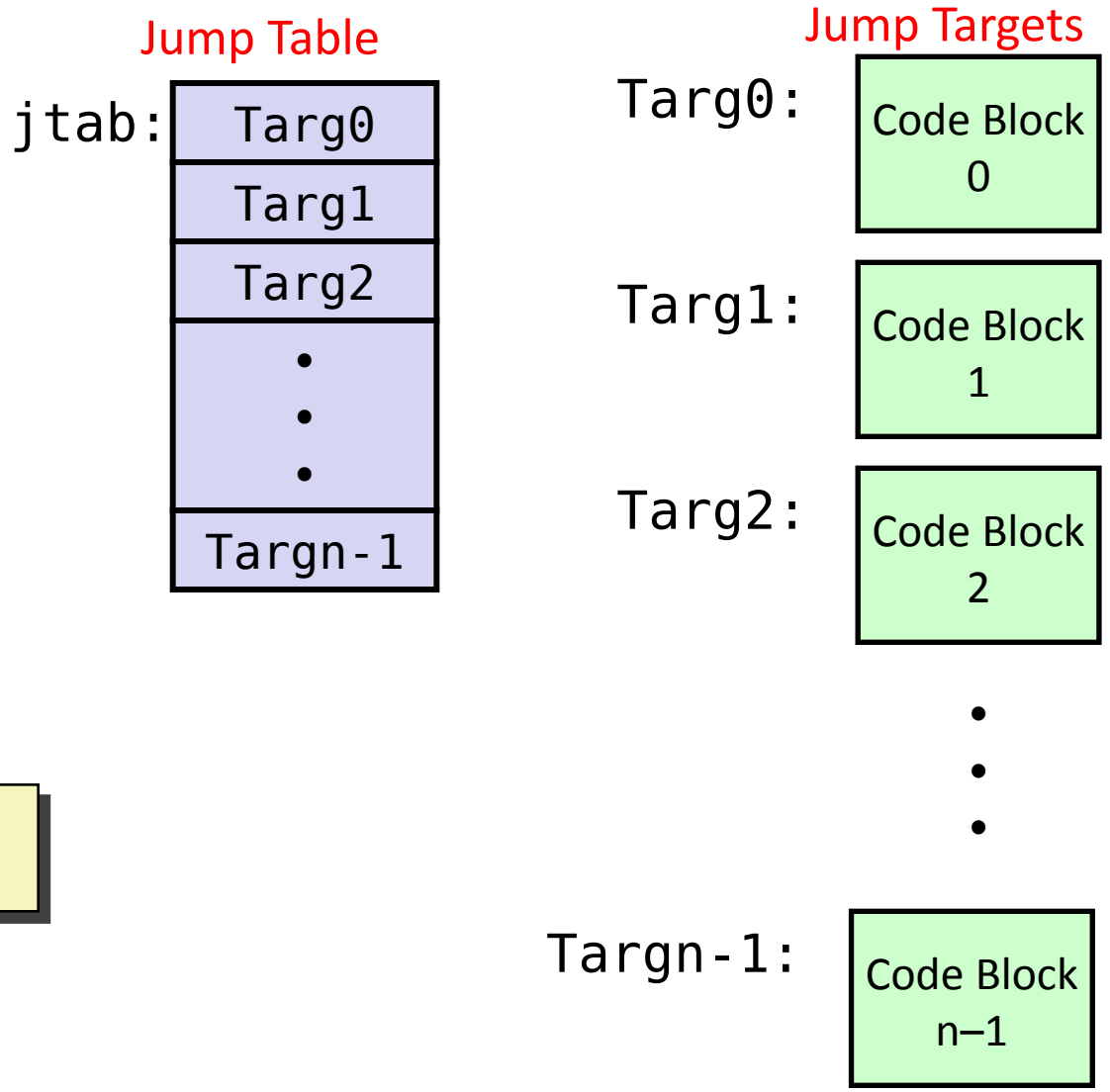
Jump Table Structure

Switch Form

```
switch(x) {
  case val_0:
    Block 0
  case val_1:
    Block 1
    • • •
  case val_n-1:
    Block n–1
}
```

Approximate Translation

```
target = JTab[x];
goto *target;
```



Switch Statement Example (IA32)

```
long switch_eg(long x, long y, long z)
{
    long w = 1;
    switch(x) {
      . . .
    }
    return w;
}
```

What range of values takes default?

Setup:

```
  switch_eg:
    pushl   %ebp            # Setup
    movl    %esp, %ebp      # Setup
    movl    8(%ebp), %eax   # %eax = x
    cmpl    $6, %eax        # Compare x:6
    ja      .L2             # If unsigned > goto default
    jmp     *.L7(,%eax,4)   # Goto *JTab[x]
```

Note that **w** not initialized here

Switch Statement Example (IA32)

```
long switch_eg(long x, long y, long z)
{
    long w = 1;
    switch(x) {
      . . .
    }
    return w;
}
```

Jump table

```
.section    .rodata
  .align 4
.L7:
  .long     .L2 # x = 0
  .long     .L3 # x = 1
  .long     .L4 # x = 2
  .long     .L5 # x = 3
  .long     .L2 # x = 4
  .long     .L6 # x = 5
  .long     .L6 # x = 6
```

Setup:

```
switch_eg:
   pushl   %ebp            # Setup
   movl    %esp, %ebp      # Setup
   movl    8(%ebp), %eax   # eax = x
   cmpl    $6, %eax        # Compare x:6
   ja      .L2             # If unsigned > goto default
   jmp     *.L7(,%eax,4)   # Goto *JTab[x]
```

Indirect jump

Assembly Setup Explanation

- **Table Structure**
  - Each target requires 4 bytes
  - Base address at `.L7`

Jump table

```
.section   .rodata
  .align 4
.L7:
  .long    .L2 # x = 0
  .long    .L3 # x = 1
  .long    .L4 # x = 2
  .long    .L5 # x = 3
  .long    .L2 # x = 4
  .long    .L6 # x = 5
  .long    .L6 # x = 6
```

- **Jumping**
  - Direct: `jmp .L2`
  - Jump target is denoted by label `.L2`
  - Indirect: `jmp *.L7(,%eax,4)`
  - Start of jump table: `.L7`
  - Must scale by factor of 4 (labels have 32-bits = 4 Bytes on IA32)
  - Fetch target from effective Address `.L7 + eax*4`
    - Only for $0 \leq x \leq 6$

IA32 Object Code

Try it yourself!

- **Setup**
  - Label `.L2` becomes address `0x8048422`
  - Label `.L7` becomes address `0x8048660`

Assembly Code

```
switch_eg:
  . . .
   ja      .L2             # If unsigned > goto default
   jmp     *.L7(,%eax,4)   # Goto *JTab[x]
```

Disassembled Object Code

```
08048410 <switch_eg>:
 . . .
 8048419:77 07                    ja     8048422 <switch_eg+0x12>
 804841b:ff 24 85 60 86 04 08     jmp    *0x8048660(,%eax,4)
```

IA32 Object Code (cont.)

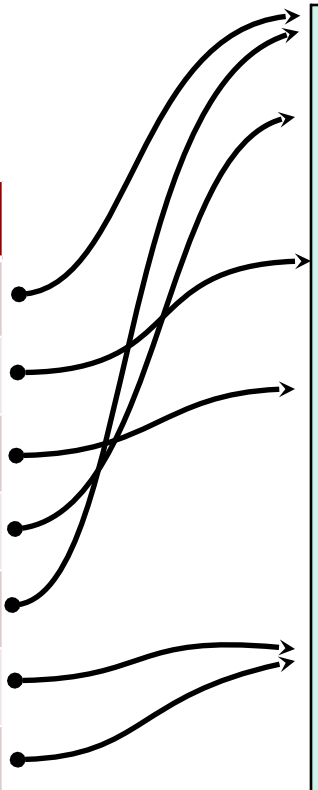
- **Jump Table**
  - Doesn't show up in disassembled code
  - Does show up with gcc -S
  - Can inspect using GDB
  - `gdb switch`
  - `(gdb) x/7xw 0x8048660`
    - Examine 7 hexadecimal format "words" (4-bytes each)
    - Use command "`help x`" to get format documentation

PP3.29

Matching Disassembled Targets

Value
0x8048422
0x8048432
0x804843b
0x8048429
0x8048422
0x804844b
0x804844b

Jump table



```
8048422:    mov     $0x2,%eax
8048427:    jmp     8048453 <switch_eg+0x43>
8048429:    mov     $0x1,%eax
804842e:    xchg    %ax,%ax
8048430:    jmp     8048446 <switch_eg+0x36>
8048432:    mov     0x10(%ebp),%eax
8048435:    imul    0xc(%ebp),%eax
8048439:    jmp     8048453 <switch_eg+0x43>
804843b:    mov     0xc(%ebp),%edx
804843e:    mov     %edx,%eax
8048440:    sar     $0x1f,%edx
8048443:    idivl   0x10(%ebp)
8048446:    add     0x10(%ebp),%eax
8048449:    jmp     8048453 <switch_eg+0x43>
804844b:    mov     $0x1,%eax
8048450:    sub     0x10(%ebp),%eax
8048453:    pop     %ebp
8048454:    ret
```

Summarizing

- **C Control**
  - if-then-else
  - do-while
  - while, for
  - switch
- **Assembler Control**
  - Conditional jump
  - Conditional move
  - Indirect jump
  - Compiler generates code sequence to implement more complex control
- **Standard Techniques**
  - Loops converted to do-while form
  - Large switch statements use jump tables
  - Sparse switch statements may use decision trees

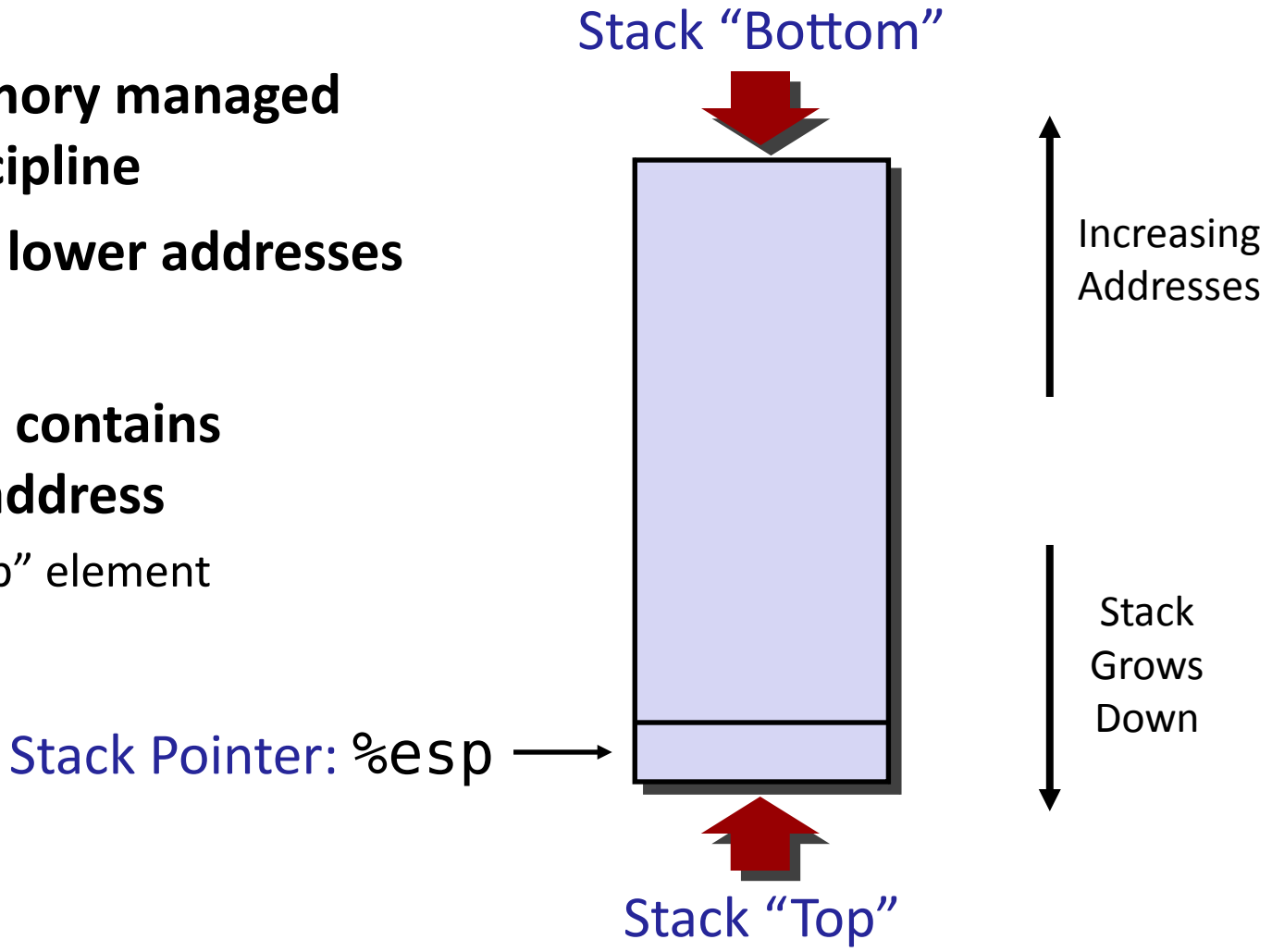
Today

- Switch statements
- **IA 32 Procedures**
  - Stack Structure
  - Calling Conventions
  - Illustrations of Recursion & Pointers
- Arrays
- (Structs)

IA32 Stack

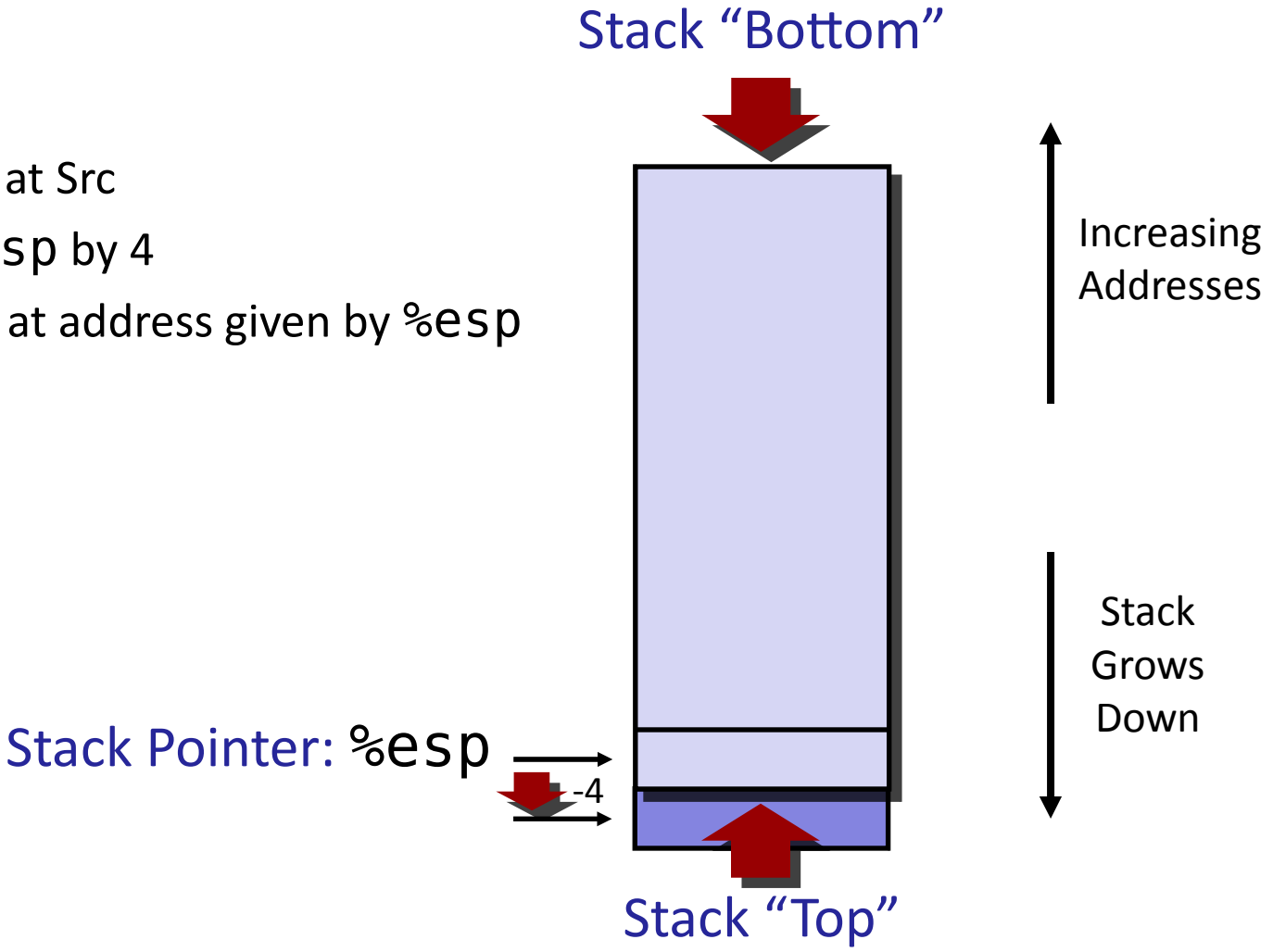
- **Region of memory managed with stack discipline**
- **Grows toward lower addresses**

- **Register `%esp` contains lowest stack address**
  - address of "top" element

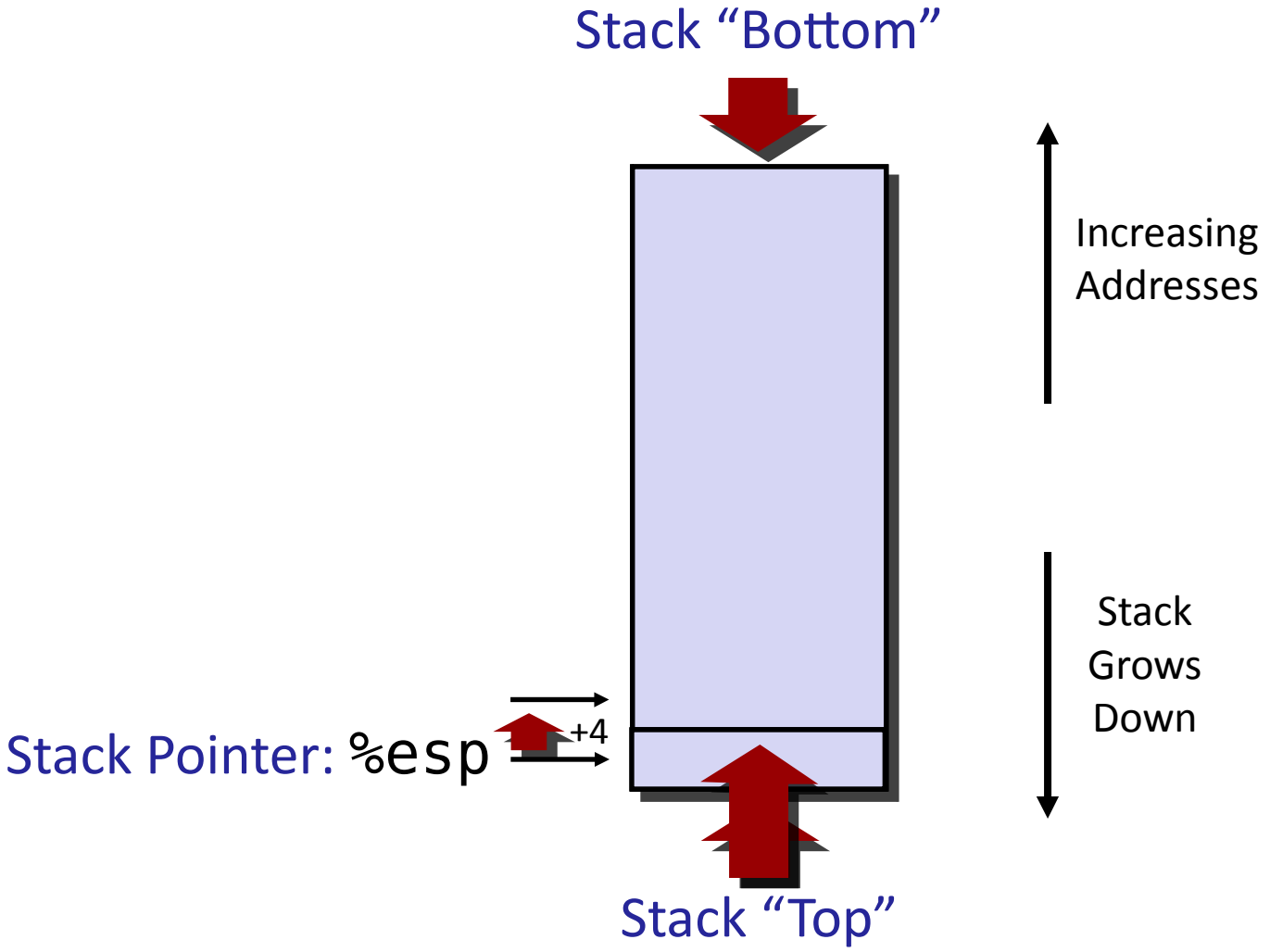


IA32 Stack: Push

- **`pushl Src`**
  - Fetch operand at Src
  - Decrement `%esp` by 4
  - Write operand at address given by `%esp`



IA32 Stack: Pop



PP3.30

Procedure Control Flow

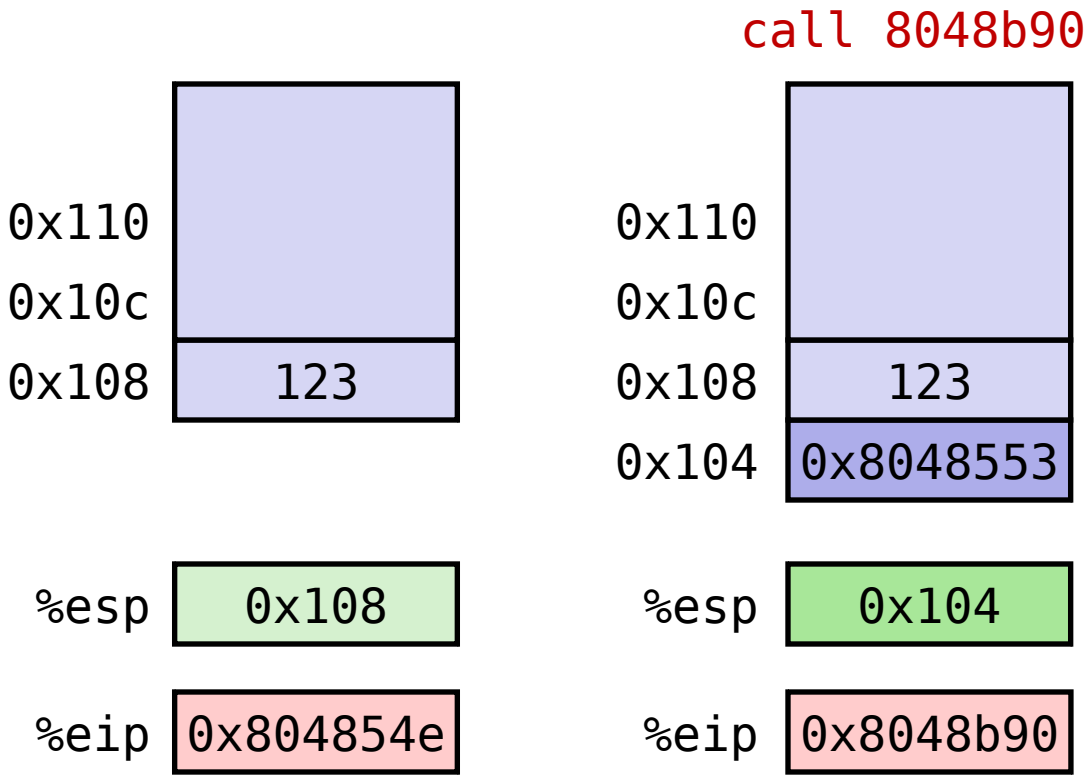
- **Use stack to support procedure call and return**
- **Procedure call: `call label`**
  - Push return address on stack
  - Jump to label
- **Return address:**
  - Address of the next instruction right after call
  - Example from disassembly

```
804854e:   e8 3d 06 00 00    call   8048b90 <main>
8048553:   50                pushl  %eax
```

  - Return address = `0x8048553`
- **Procedure return: `ret`**
  - Pop address from stack
  - Jump to address

Procedure Call Example

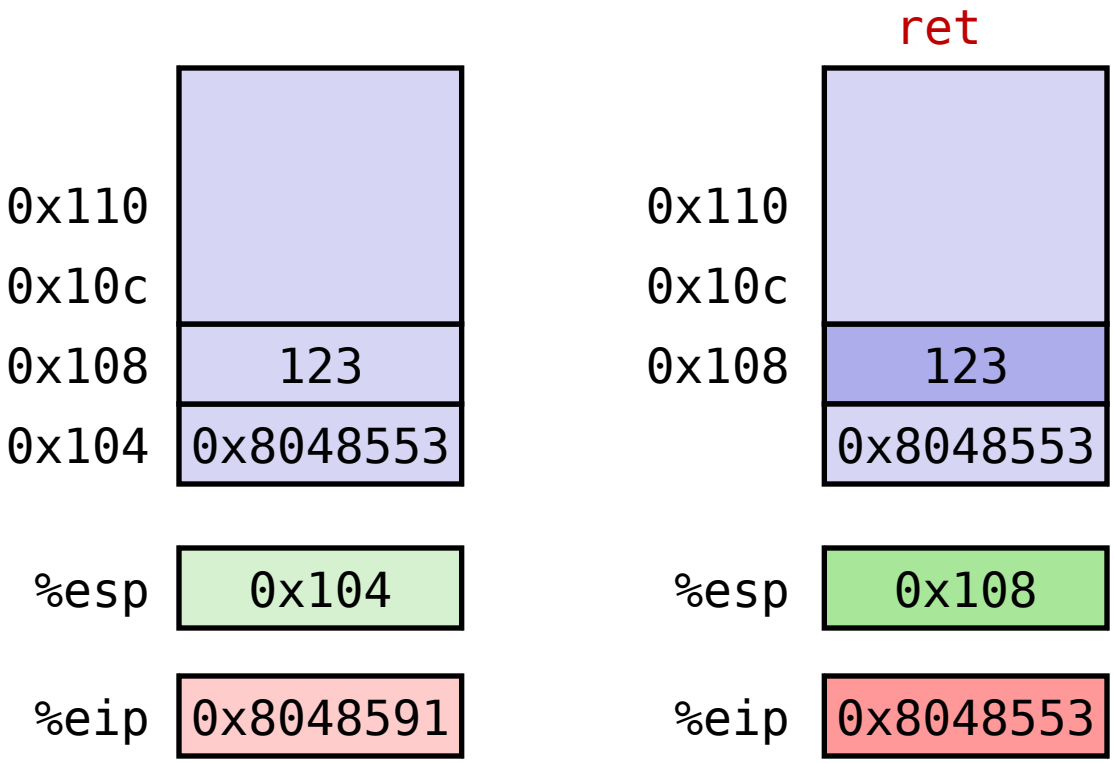
```
804854e:  e8 3d 06 00 00   call   8048b90 <main>
8048553:  50               pushl  %eax
```



%eip: program counter

Procedure Return Example

```
8048591: c3                    ret
```



`%eip:` program counter

Stack-Based Languages

- **Languages that support recursion**
  - e.g., C, Pascal, Java
  - Code must be “Reentrant”
    - Multiple simultaneous instantiations of single procedure
  - Need some place to store state of each instantiation
    - Arguments
    - Local variables
    - Return pointer
- **Stack discipline**
  - State for given procedure needed for limited time
    - From when called to when return
  - Callee returns before caller does
- **Stack allocated in Frames**
  - state for single procedure instantiation

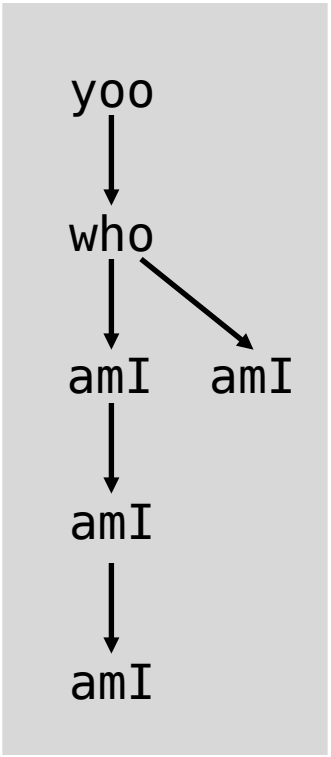
Call Chain Example

```
yoo(…)
{
   •
   •
   who();
   •
   •
}
```

```
who(…)
{
   • • •
   amI();
   • • •
   amI();
   • • •
}
```

```
amI(…)
{
   •
   •
   amI();
   •
   •
}
```

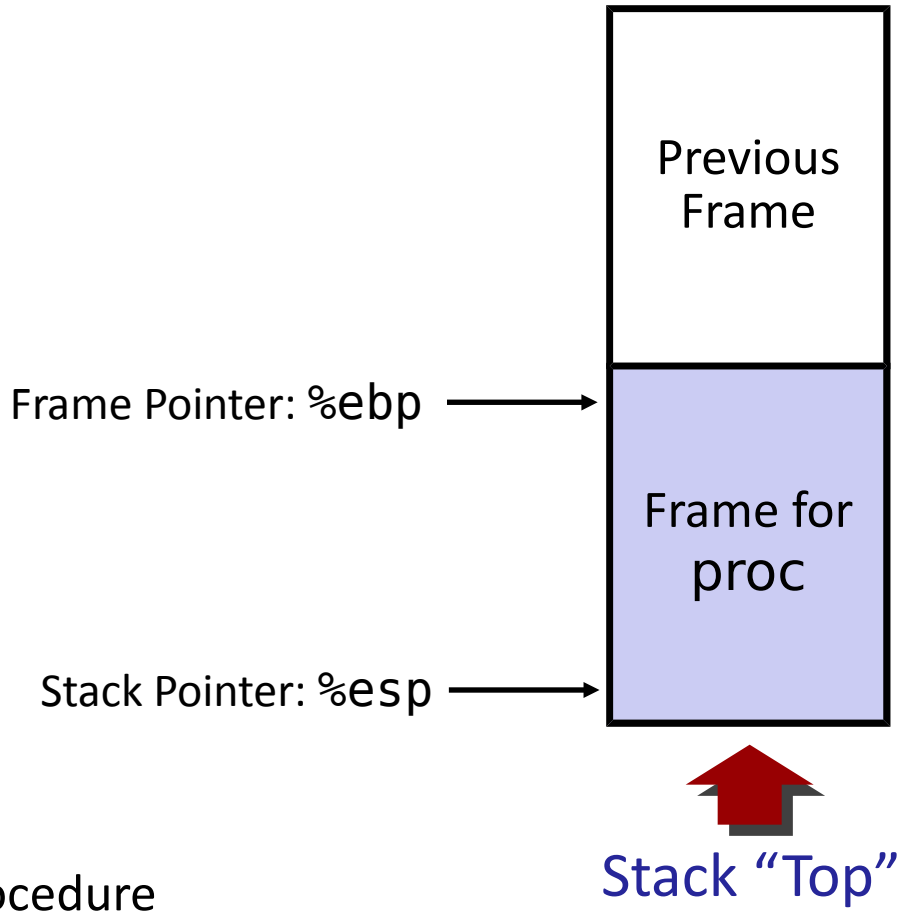
Example
Call Chain



Procedure `amI()` is recursive

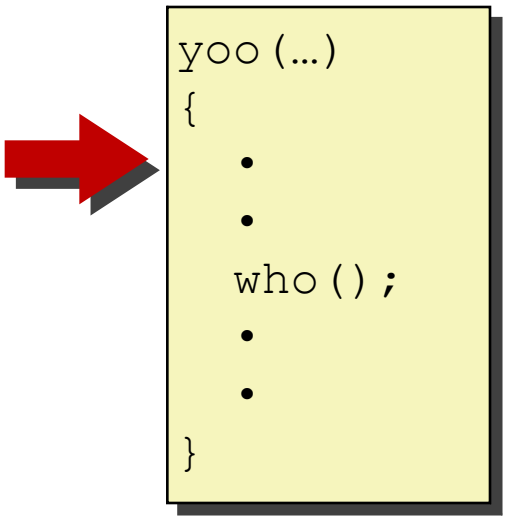
Stack Frames

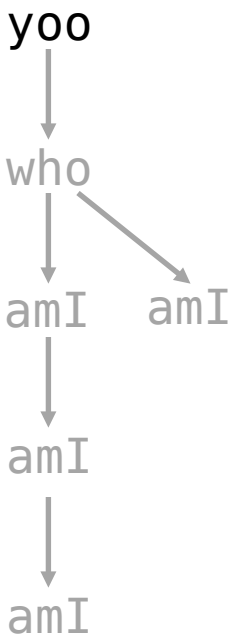
- **Contents**
  - Local variables
  - Return information
  - Temporary space

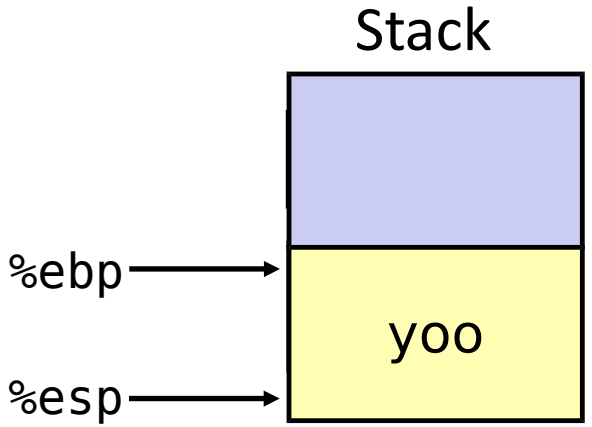


- **Management**
  - Space allocated when enter procedure
    - "Set-up" code
  - Deallocated when return
    - "Finish" code

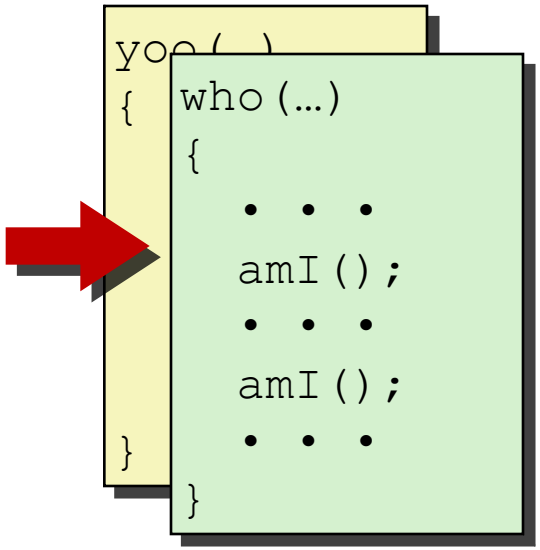
Example

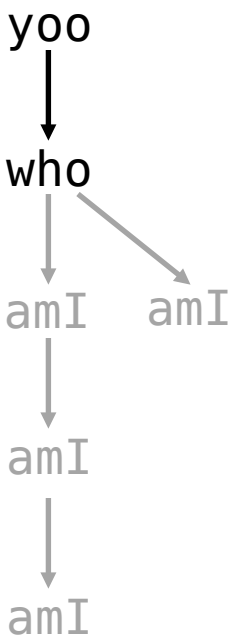


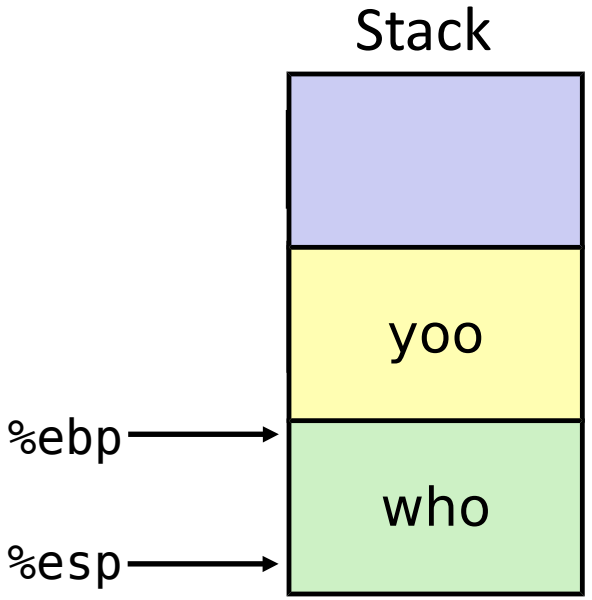




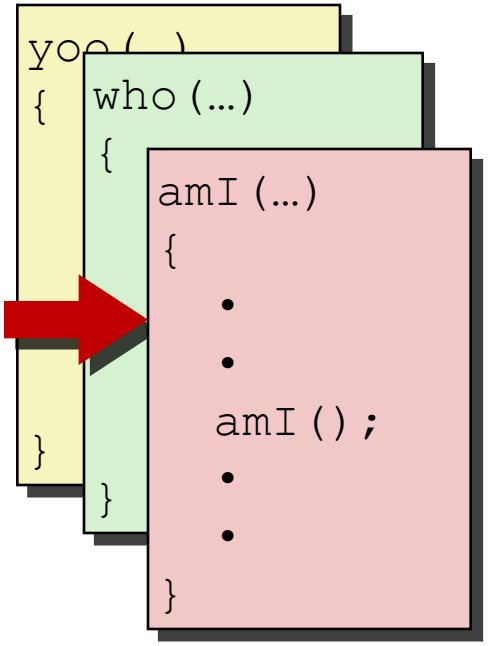
Example

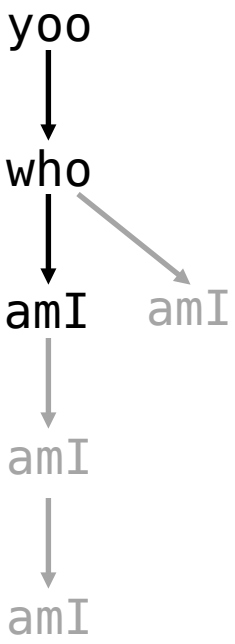


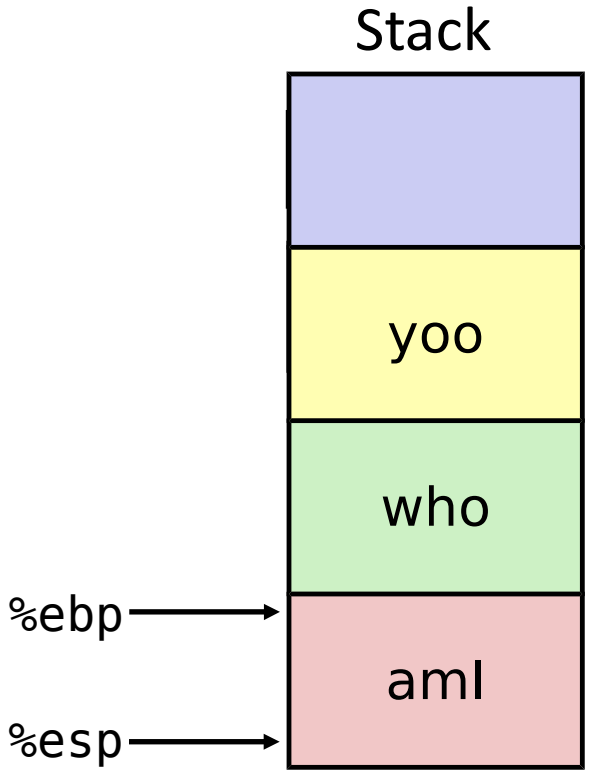




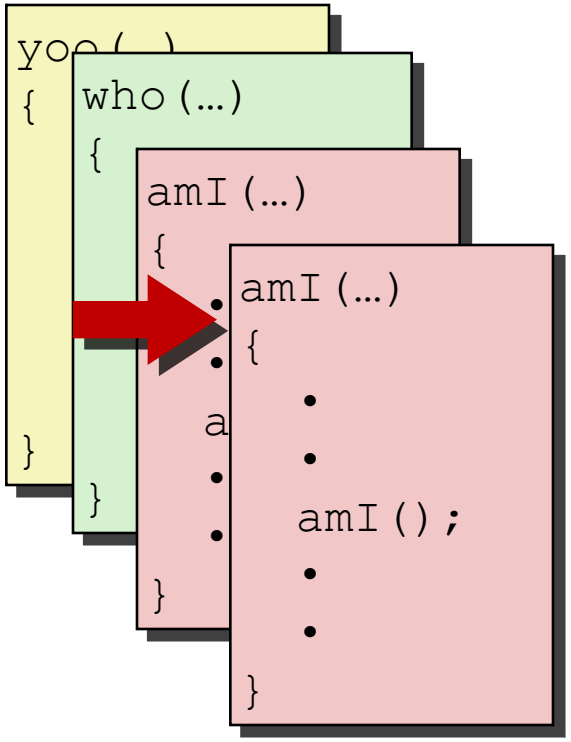
Example


yoo( )
{
who(…)
{
amI(…)
{
•
•
amI();
•
•
}
}
}


yoo
who
amI
amI
amI
amI


Stack
yoo
who
amI
%ebp
%esp

Example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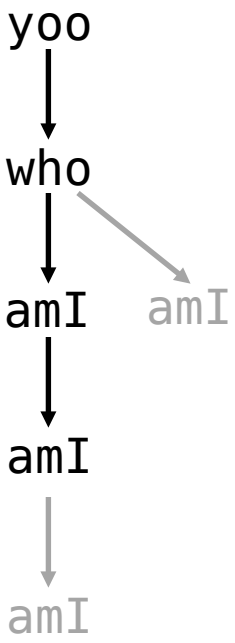
{
who(…)
{
amI(…)
{
amI(…)
{
amI();
}
}
}
}


yoo
who
amI
amI
amI
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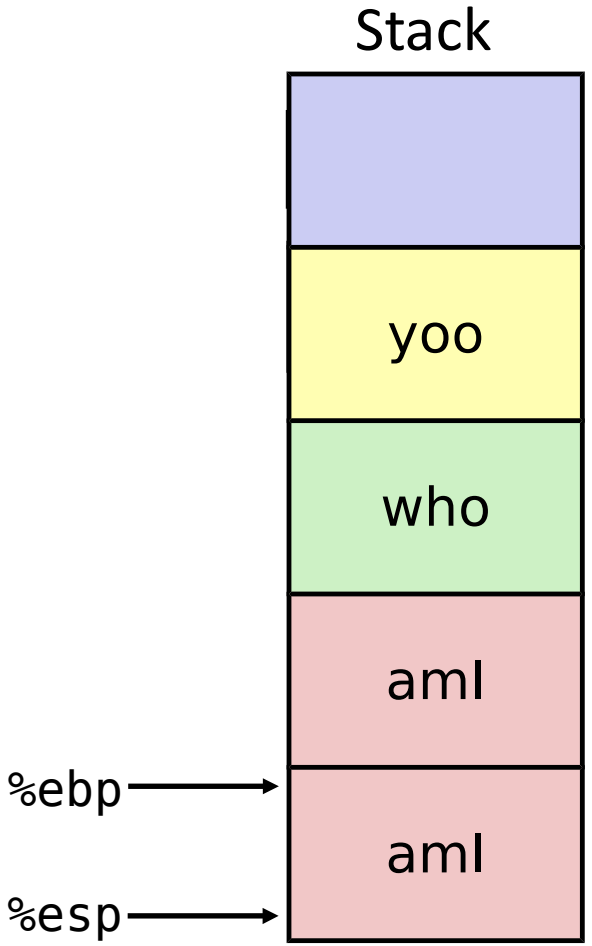
Stack
yoo
who
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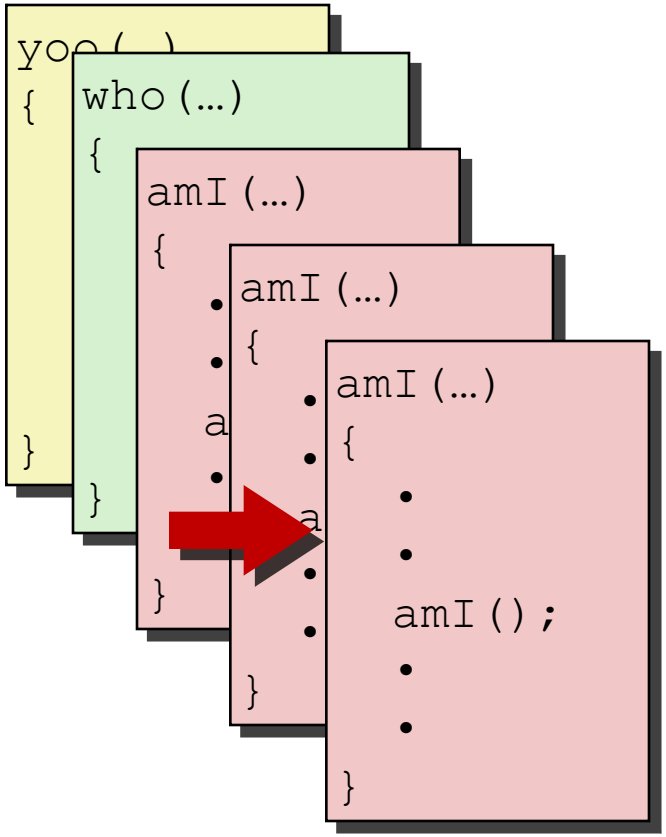
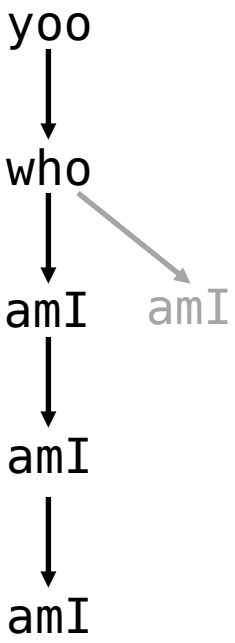
%esp

Example


yoo( )
{
who(…)
{
amI(…)
{
amI(…)
{
amI(…)
{
amI();
}
}
}
}


yoo
who
amI
amI
amI
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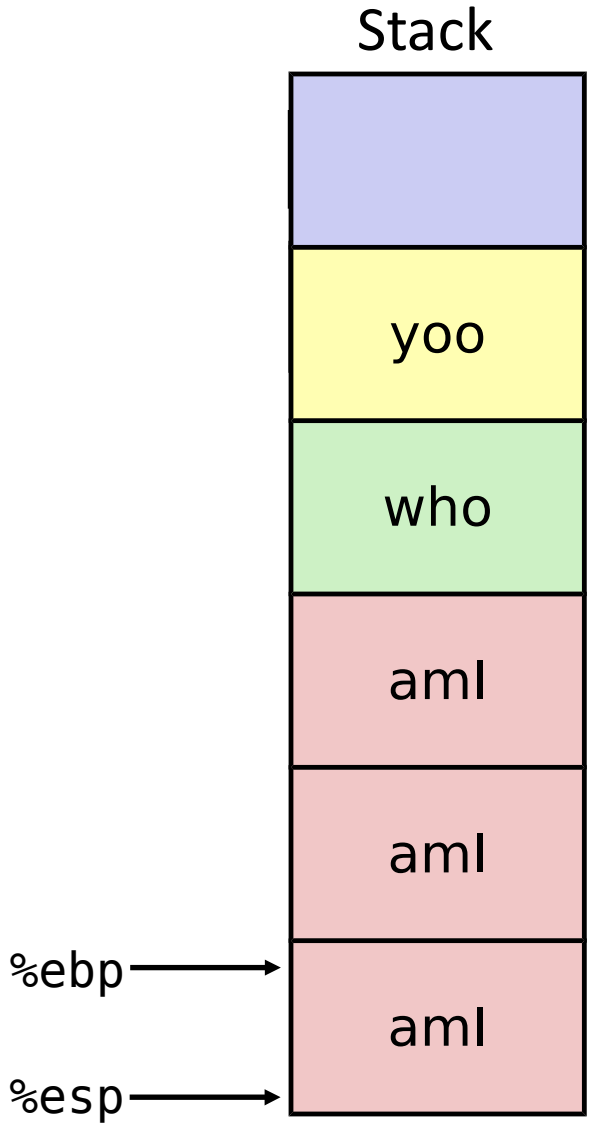
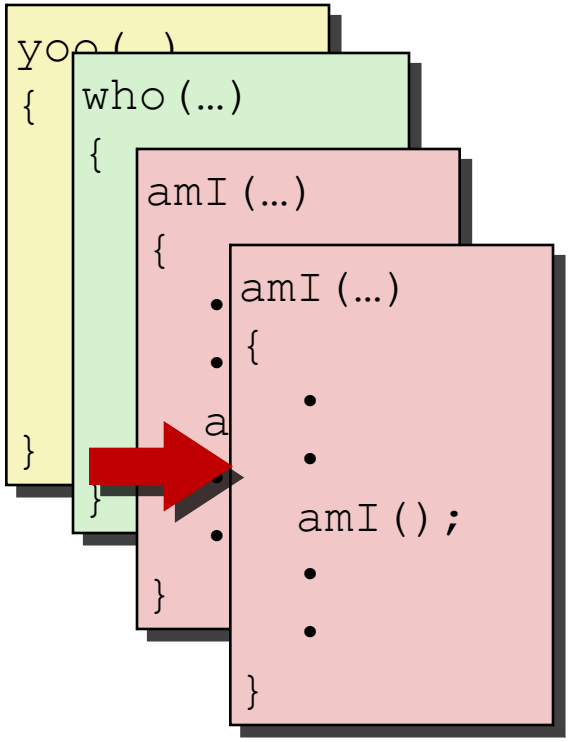

Stack
yoo
who
amI
amI
amI
%ebp
%esp

Example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amI(…)
{
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}
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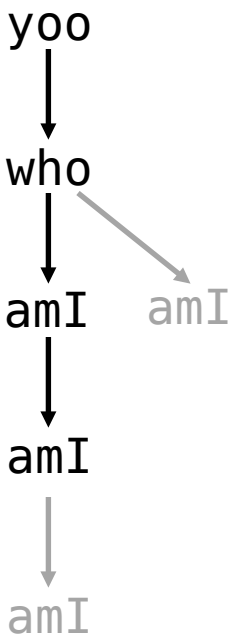
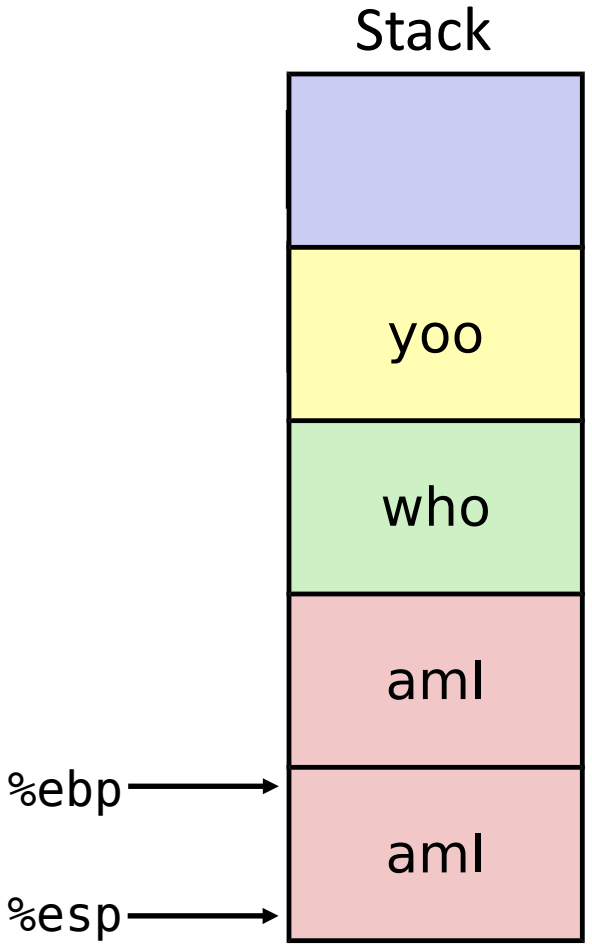
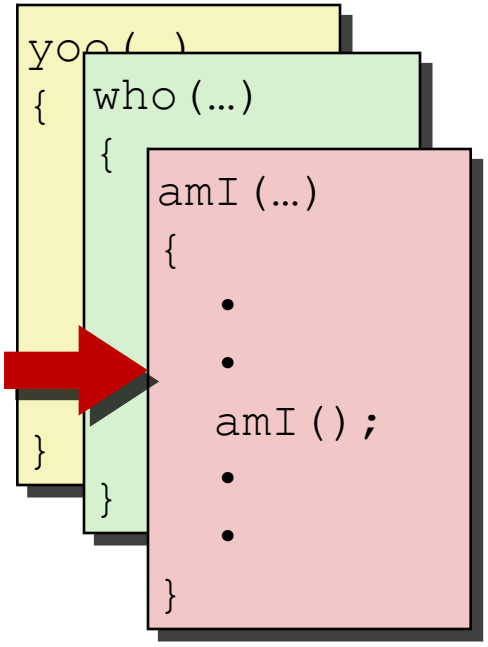

yoo
who
amI
amI
amI
amI


Stack
yoo
who
aml
aml
%ebp
%esp

Example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amI(…)
{
•
•
amI();
•
•
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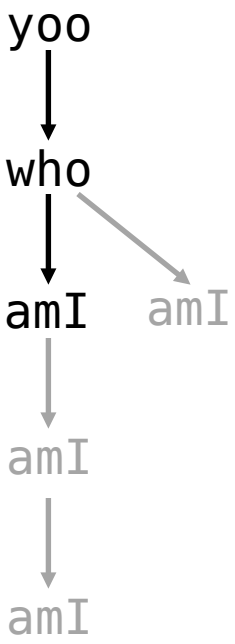
}
}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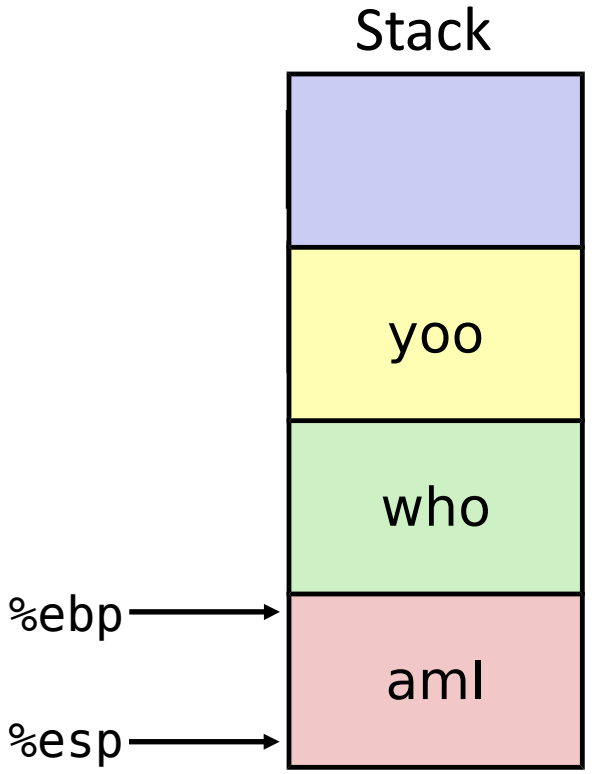
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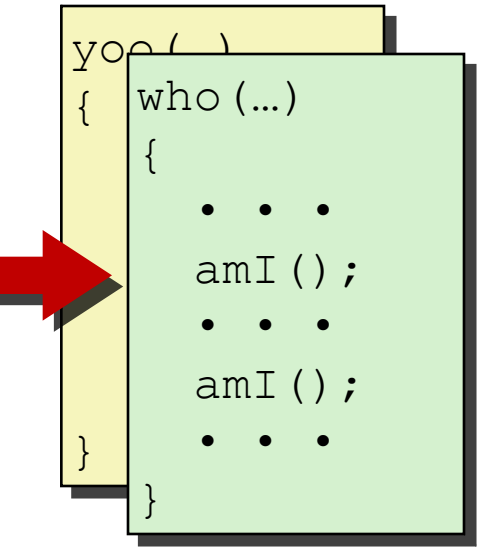
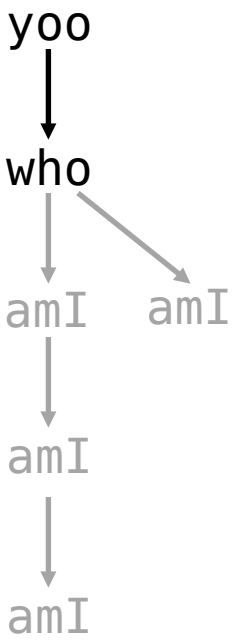
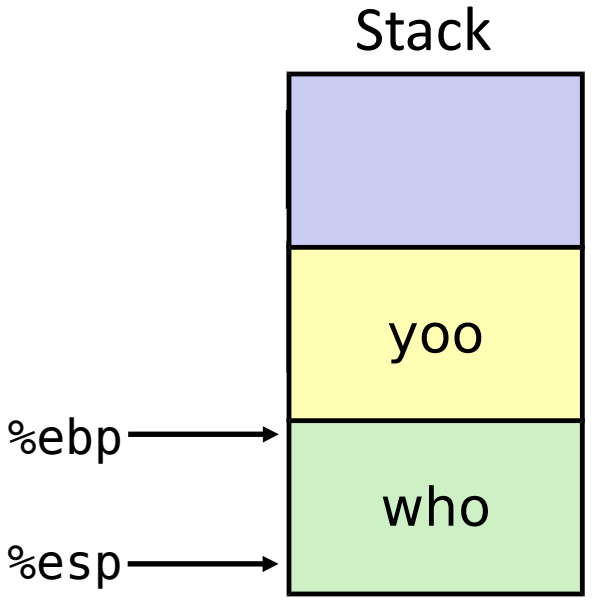
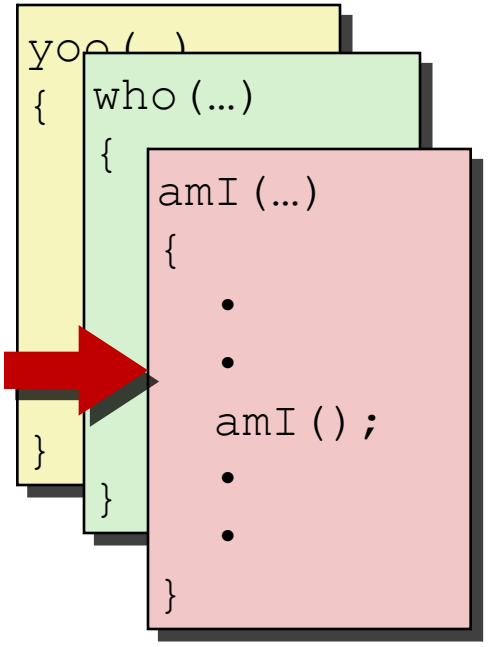
Example

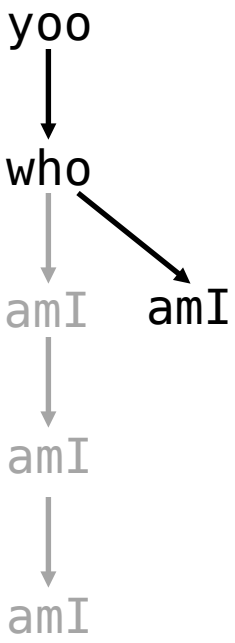


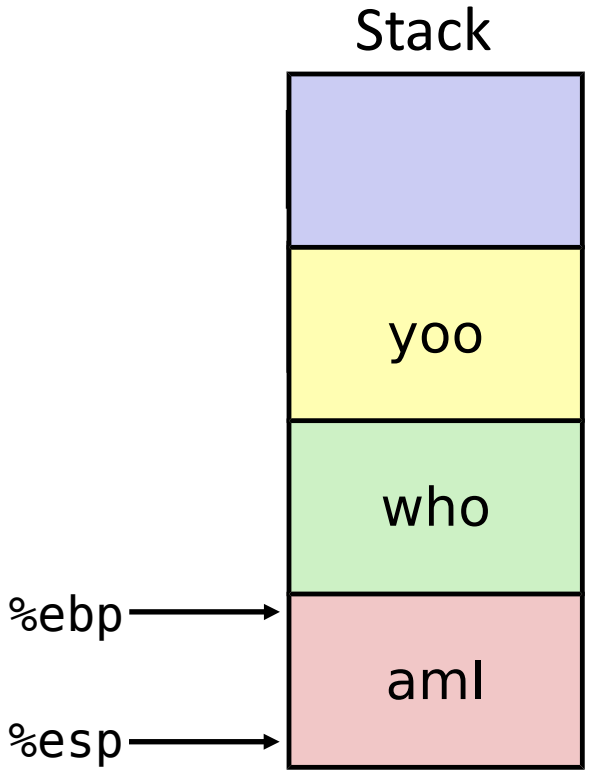




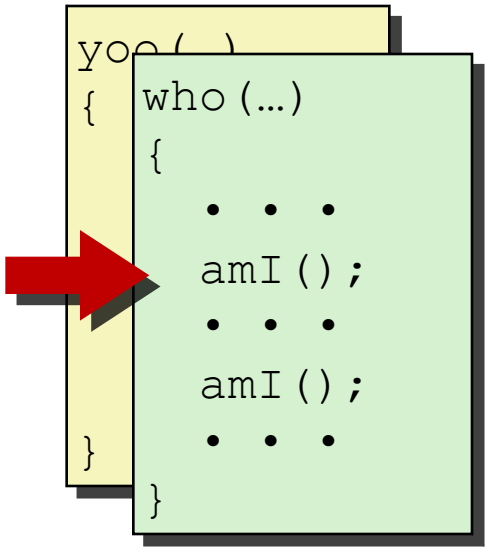
Example

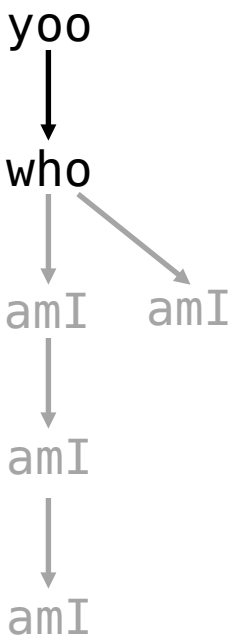

yoo( )
{
who(…)
{
amI(…)
{
•
•
amI();
•
•
}
}
}

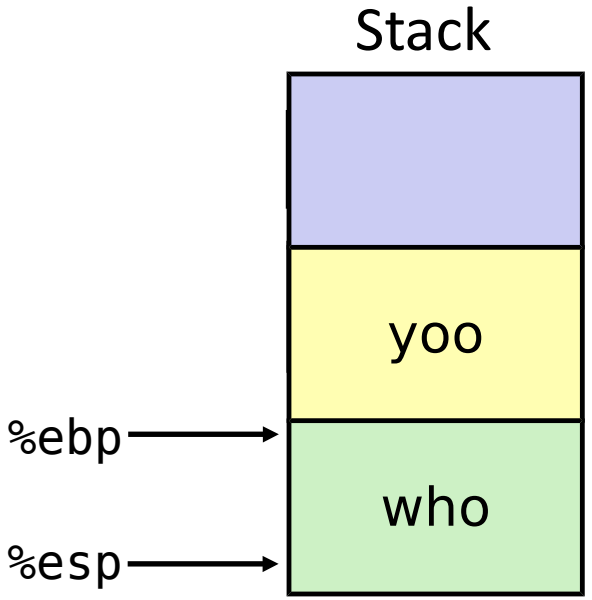

yoo
who
amI
amI
amI
amI


Stack
yoo
who
amI
%ebp
%esp

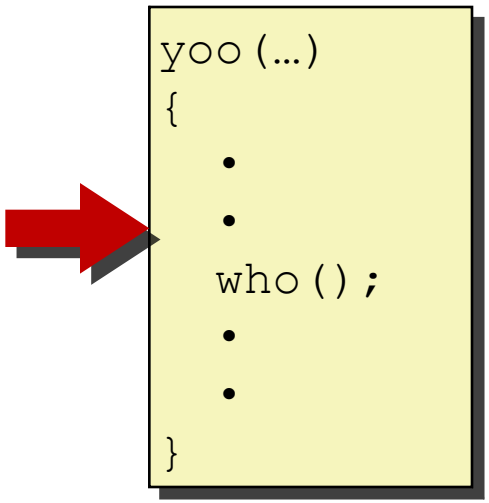
Example

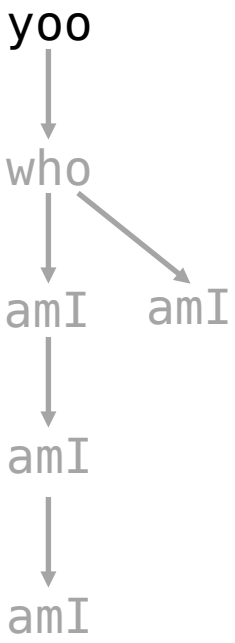


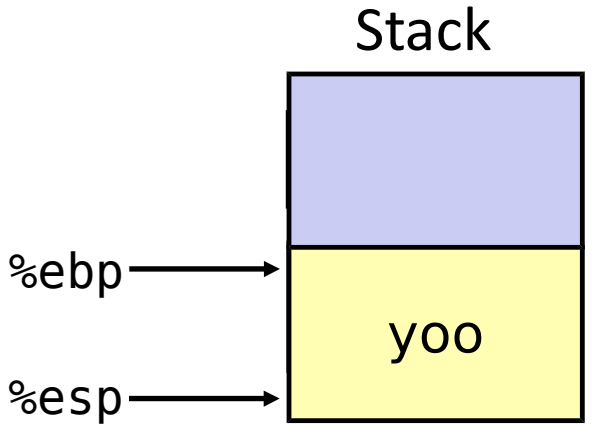




Example



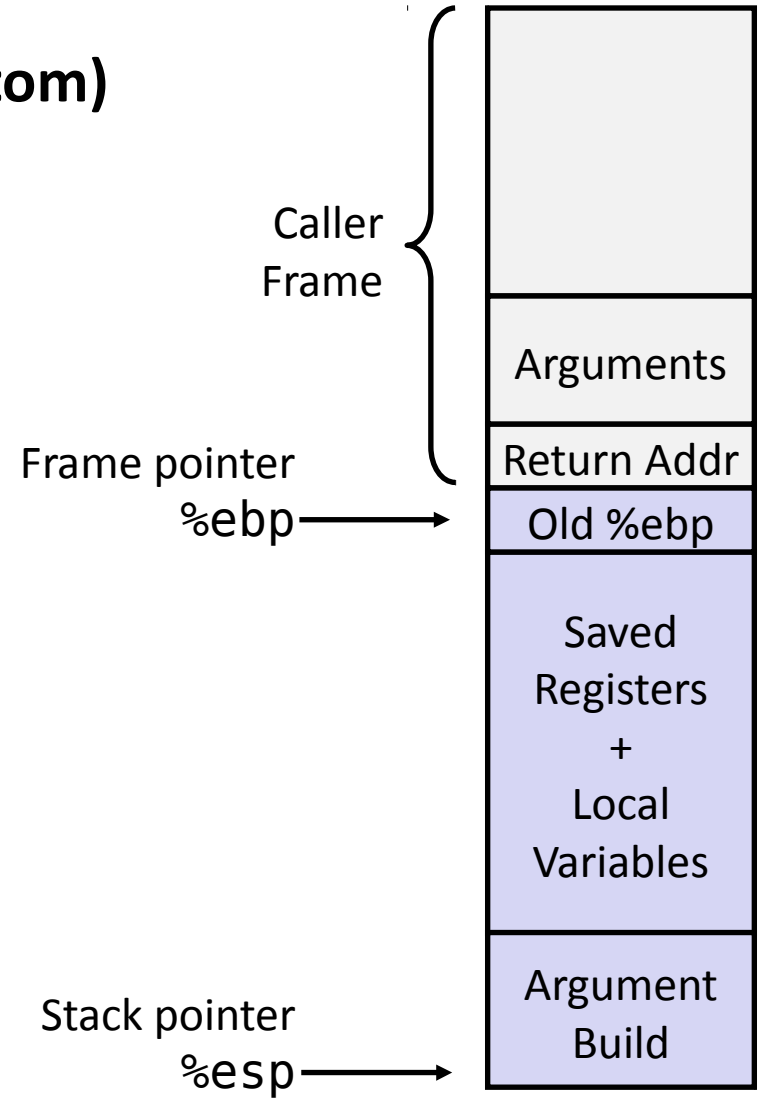




IA32/Linux Stack Frame

- **Current Stack Frame ("Top" to Bottom)**
  - "Argument build:"
    Parameters for function about to call
  - Local variables
    If can't keep in registers
  - Saved register context
  - Old frame pointer

- **Caller Stack Frame**
  - Return address
    - Pushed by `call` instruction
  - Arguments for this call



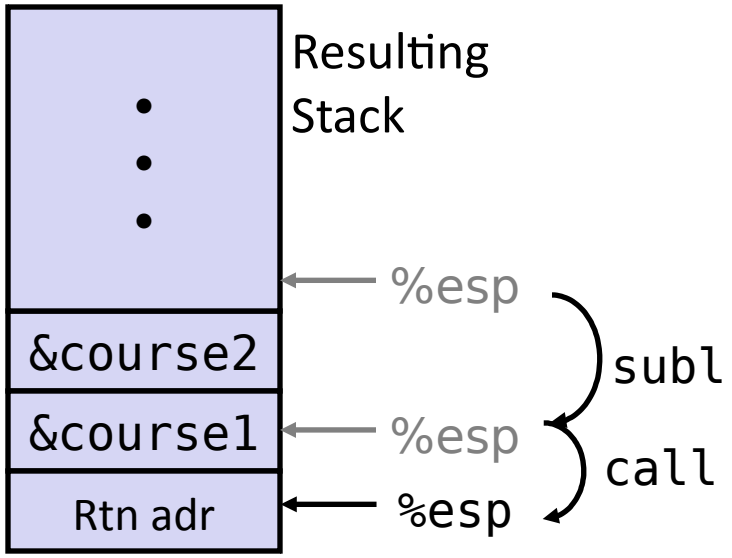
Revisiting swap

```
int course1 = 15213;
int course2 = 18243;

void call_swap() {
  swap(&course1, &course2);
}
```

Calling swap from call_swap

```
call_swap:
    • • •
    subl    $8, %esp
    movl    course2, 4(%esp)
    movl    course1, (%esp)
    call    swap
    • • •
```

```
void swap(int *xp, int *yp)
{
  int t0 = *xp;
  int t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```



Revisiting swap

```
void swap(int *xp, int *yp)
{
  int t0 = *xp;
  int t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```

```
swap:
    pushl %ebp
    movl  %esp, %ebp
    pushl %ebx

    movl  8(%ebp), %edx
    movl  12(%ebp), %ecx
    movl  (%edx), %ebx
    movl  (%ecx), %eax
    movl  %eax, (%edx)
    movl  %ebx, (%ecx)

    popl  %ebx
    popl  %ebp
    ret
```

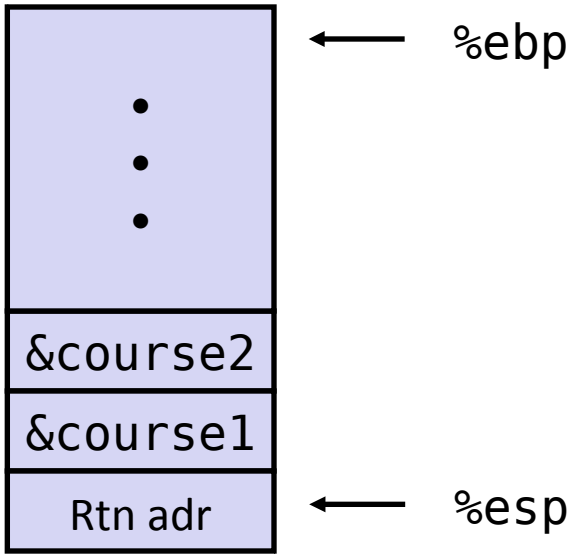
Set Up

Body

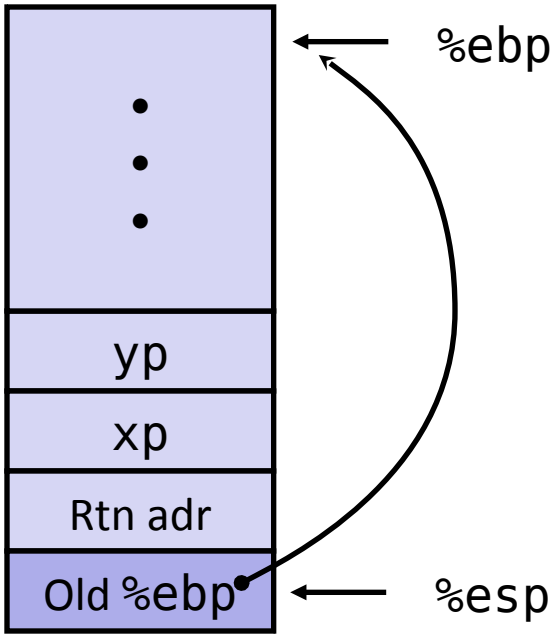
Finish

swap Setup #1

Entering Stack



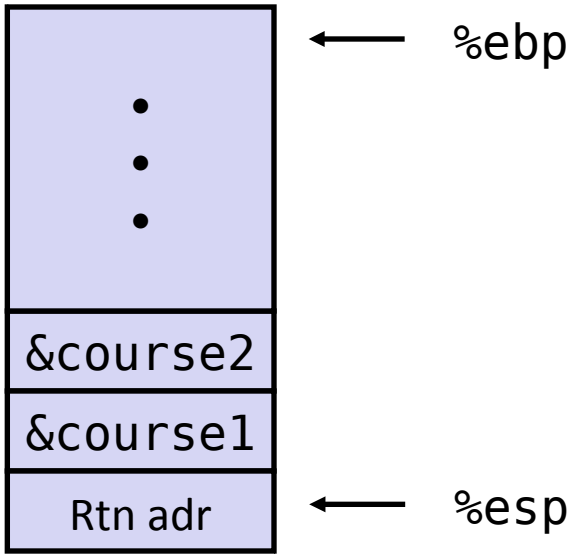
Resulting Stack



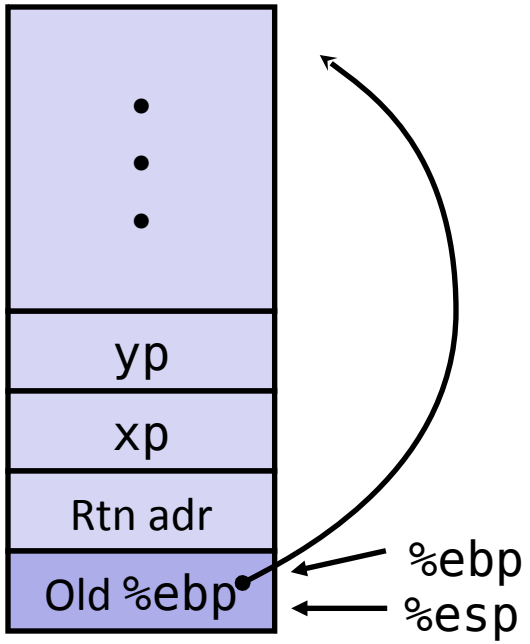

```
swap:
    pushl %ebp
    movl %esp,%ebp
    pushl %ebx
```

swap Setup #2

Entering Stack



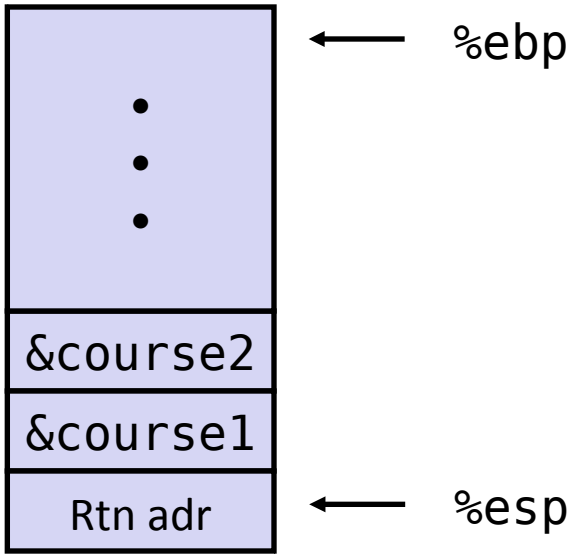
Resulting Stack



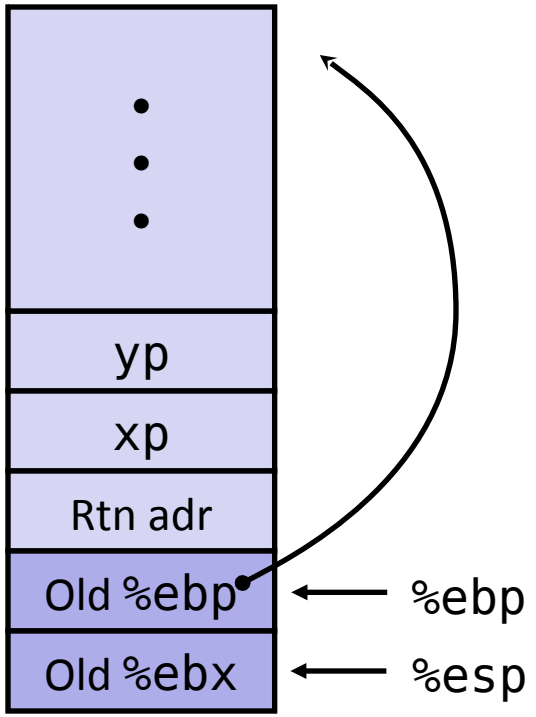

```
swap:
    pushl %ebp
    movl %esp,%ebp
    pushl %ebx
```

swap Setup #3

Entering Stack

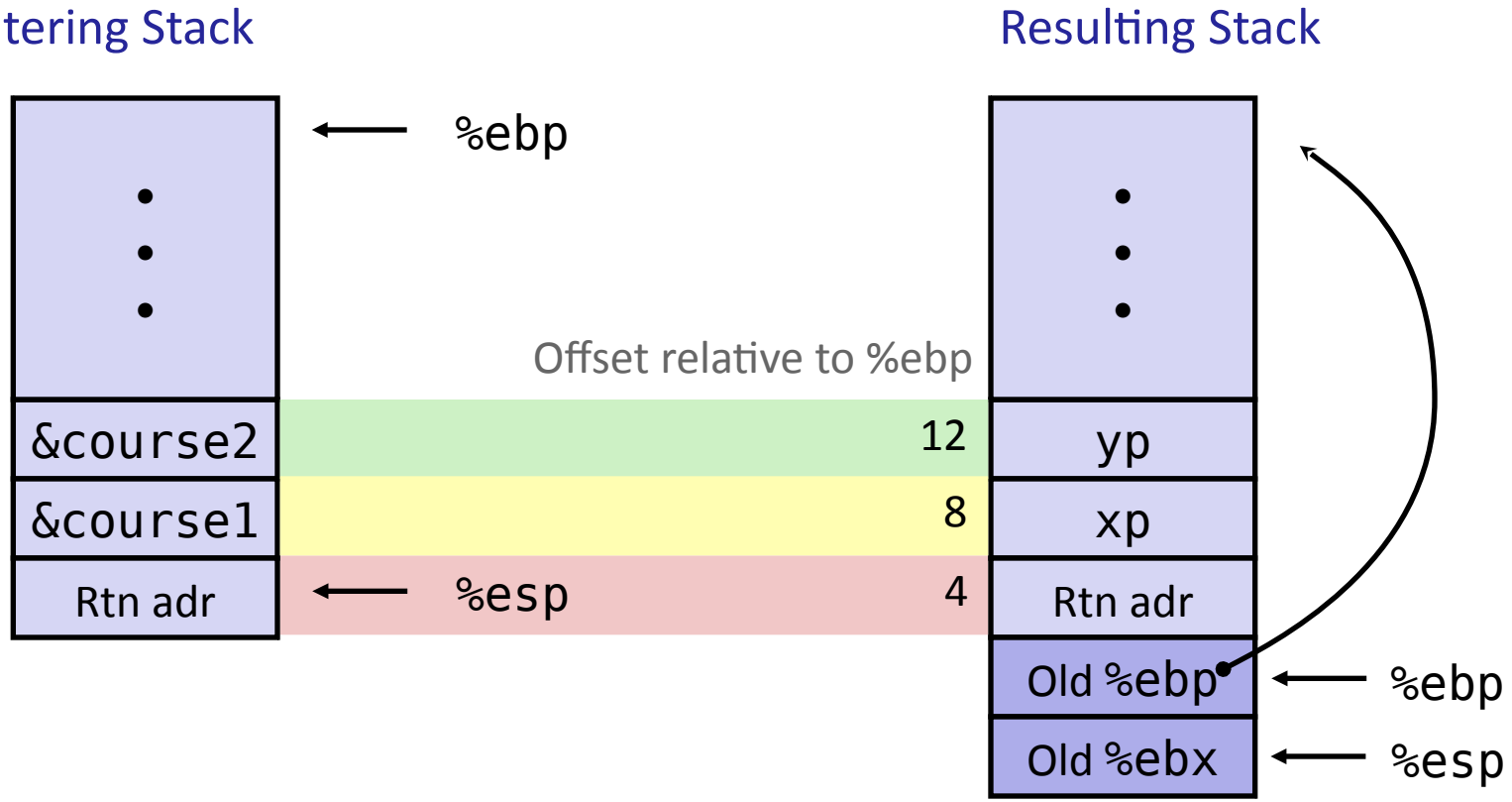


Resulting Stack




```
swap:
    pushl %ebp
    movl %esp,%ebp
    pushl %ebx
```

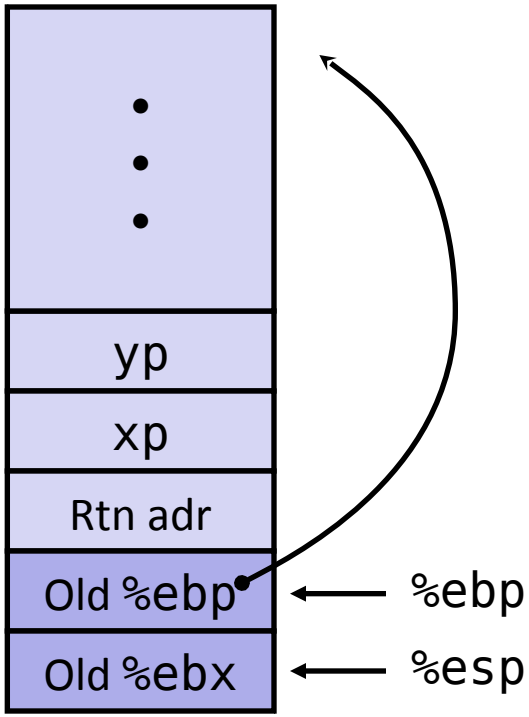
swap Body




```
movl 8(%ebp),%edx    # get xp
movl 12(%ebp),%ecx   # get yp
. . .
```

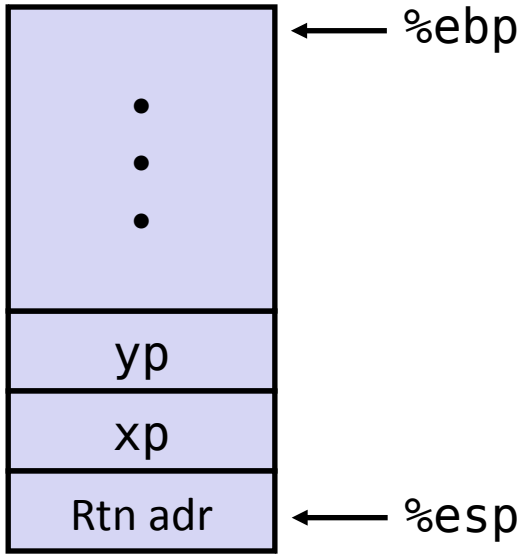
swap Finish

Stack Before Finish




```
popl    %ebx
popl    %ebp
```

Resulting Stack



- Observation
  - Saved and restored register %ebx
  - Not so for %eax, %ecx, %edx

Disassembled swap

```
08048384 <swap>:
 8048384:  55                    push   %ebp
 8048385:  89 e5                 mov    %esp,%ebp
 8048387:  53                    push   %ebx
 8048388:  8b 55 08              mov    0x8(%ebp),%edx
 804838b:  8b 4d 0c              mov    0xc(%ebp),%ecx
 804838e:  8b 1a                 mov    (%edx),%ebx
 8048390:  8b 01                 mov    (%ecx),%eax
 8048392:  89 02                 mov    %eax,(%edx)
 8048394:  89 19                 mov    %ebx,(%ecx)
 8048396:  5b                    pop    %ebx
 8048397:  5d                    pop    %ebp
 8048398:  c3                    ret
```

Calling Code

```
 80483b4:  movl   $0x8049658,0x4(%esp) # Copy &course2
 80483bc:  movl   $0x8049654,(%esp)    # Copy &course1
 80483c3:  call   8048384 <swap>       # Call swap
 80483c8:  leave                       # Prepare to return
 80483c9:  ret                         # Return
```

Pointer Code

Generating Pointer

```
/* Compute x + 3 */
int add3(int x) {
  int localx = x;
  incrk(&localx, 3);
  return localx;
}
```

Referencing Pointer

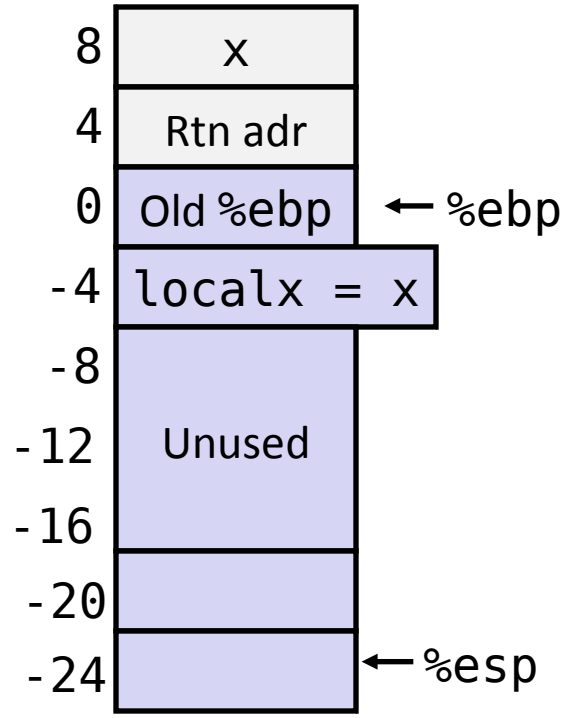
```
/* Increment value by k */
void incrk(int *ip, int k) {
  *ip += k;
}
```

- **add3 creates pointer and passes it to `incrk`**

Creating and Initializing Local Variable

```
int add3(int x) {
  int localx = x;
  incrk(&localx, 3);
  return localx;
}
```

- **Variable localx must be stored on stack**
  - Because: Need to create pointer to it
- **Compute pointer as -4(%ebp)**



First part of add3

```
add3:
  pushl %ebp
  movl  %esp, %ebp
  subl  $24, %esp       # Alloc. 24 bytes
  movl  8(%ebp), %eax
  movl  %eax, -4(%ebp)  # Set localx to x
```

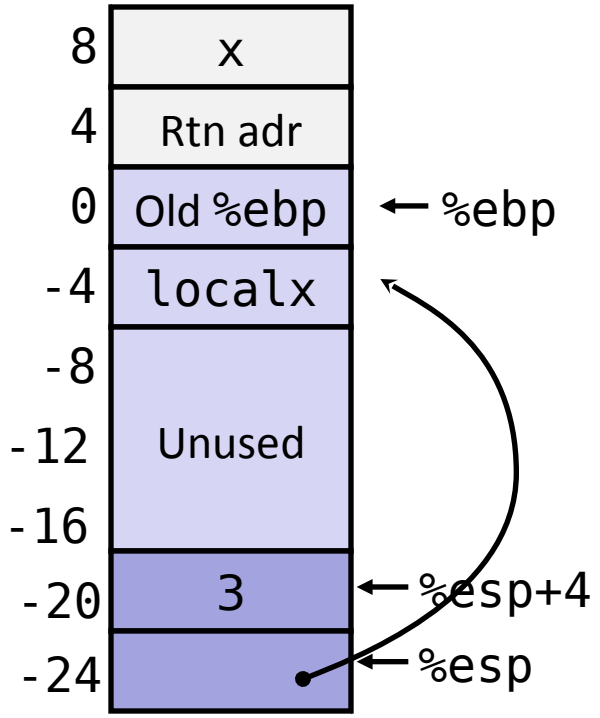
Creating Pointer as Argument

```
int add3(int x) {
  int localx = x;
  incrk(&localx, 3);
  return localx;
}
```

- Use leal instruction to compute address of localx

Middle part of add3

```
  movl $3, 4(%esp)    # 2nd arg = 3
  leal -4(%ebp), %eax# &localx
  movl %eax, (%esp)   # 1st arg = &localx
  call incrk
```



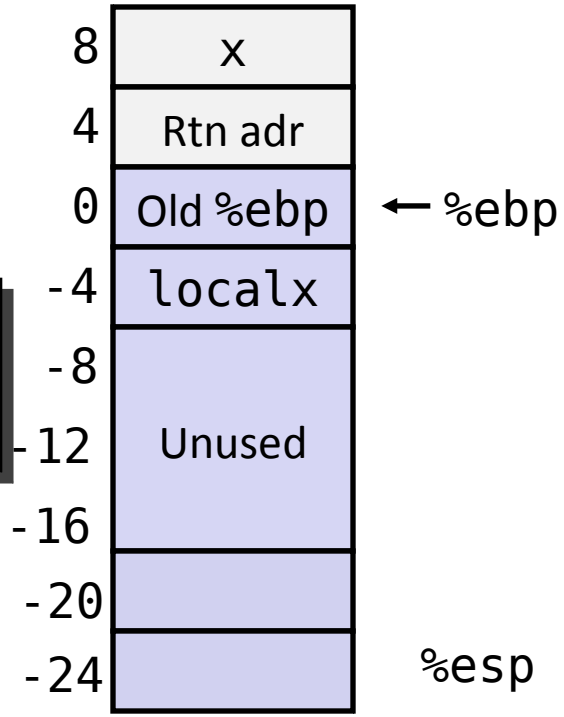
Retrieving local variable

```
int add3(int x) {
  int localx = x;
  incrk(&localx, 3);
  return localx;
}
```

- Retrieve localx from stack as return value

Final part of add3

```
movl -4(%ebp), %eax # Return val= localx
leave
ret
```



Today

- Switch statements *
- IA 32 Procedures
  - Stack Structure
  - Calling Conventions
  - Illustrations of Recursion & Pointers
- Arrays
- (Structs)

Basic Data Types

Integral

- Stored & operated on in general (integer) registers
- Signed vs. unsigned depends on instructions used

Intel	ASM	Bytes	C
byte	`b`	1	`[unsigned] char`
word	`w`	2	`[unsigned] short`
double word	`l`	4	`[unsigned] int`
quad word	`q`	8	`[unsigned] long int` (x86-64)

Floating Point

- Stored & operated on in floating point registers

Intel	ASM	Bytes	C
Single	`s`	4	`float`
Double	`l`	8	`double`
Extended	`t`	10/12/16	`long double`

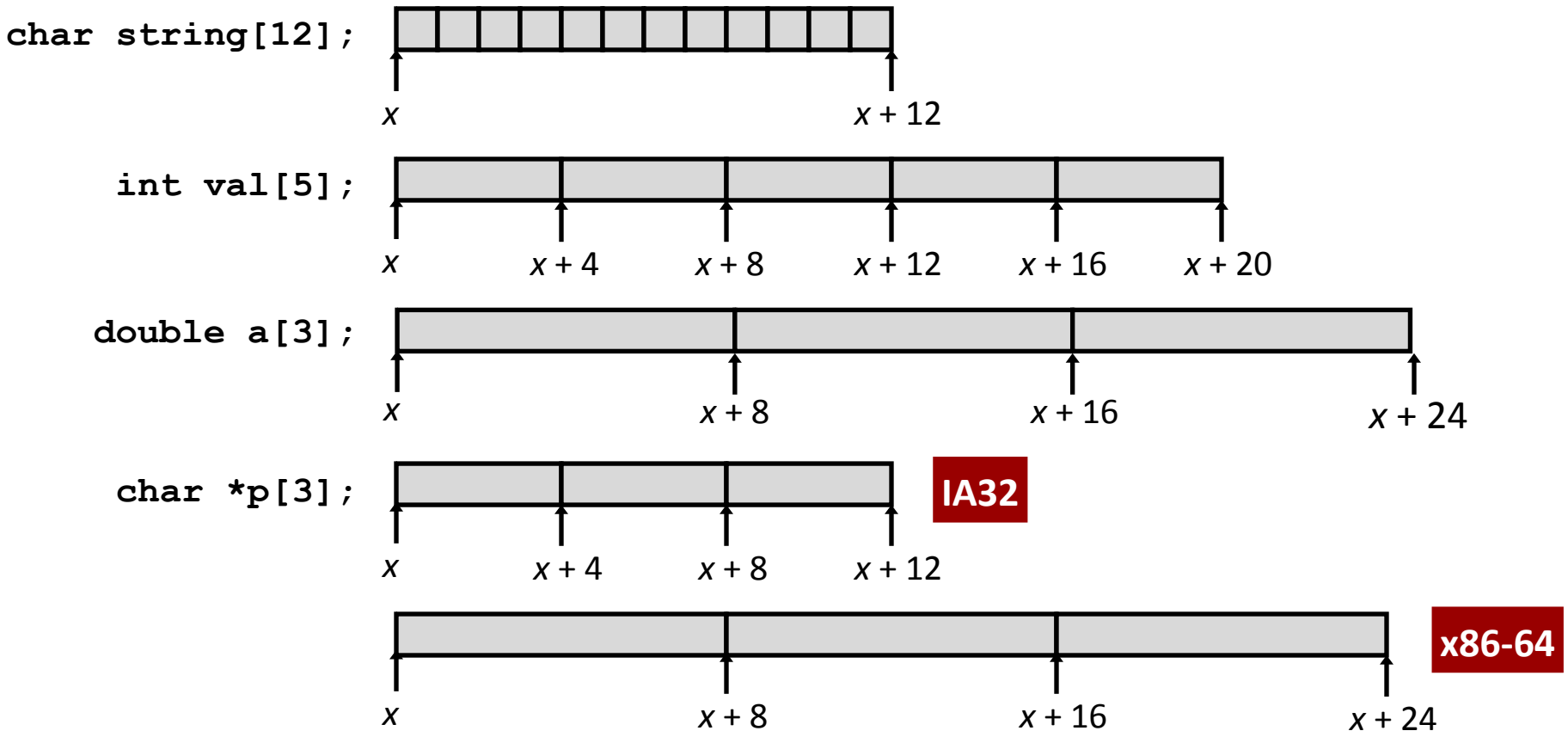
Array Allocation

PP3.35

- **Basic Principle**

  T **`A[`***L***`];`**

  - Array of data type *T* and length *L*
  - Contiguously allocated region of *L* * **`sizeof`**`(`*T*`)` bytes

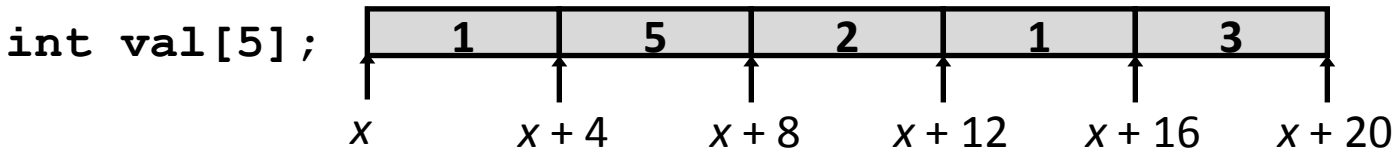


Array Access

- **Basic Principle**

  T **`A[`***L***`];`**

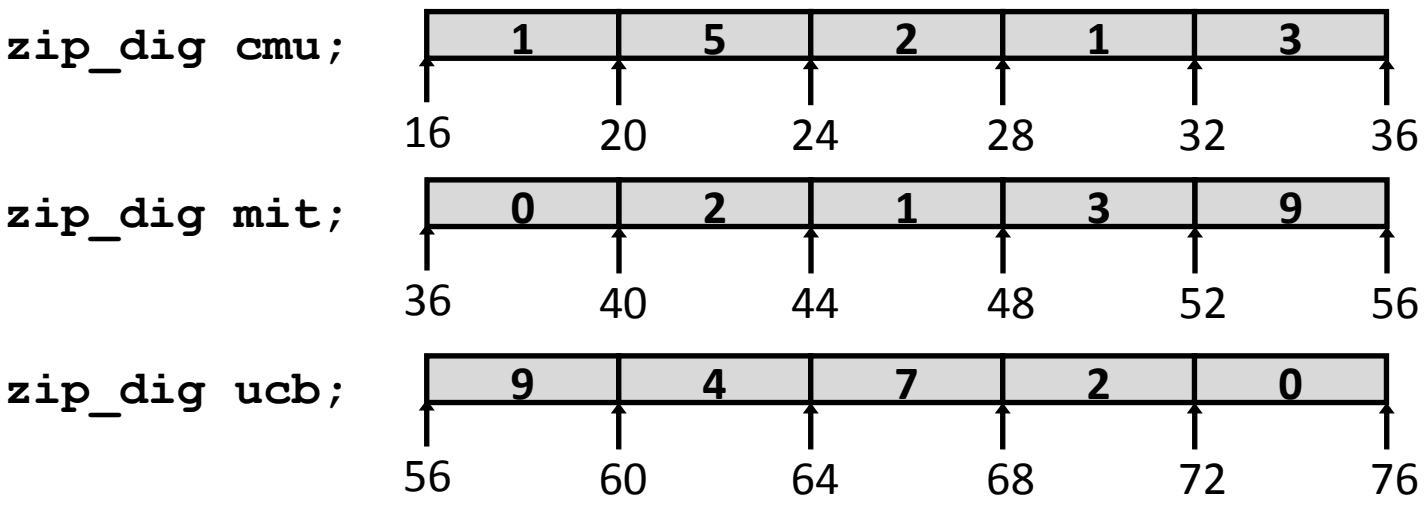
  - Array of data type *T* and length *L*
  - Identifier **`A`** can be used as a pointer to array element 0: Type *T**



Reference	Type	Value
`val[4]`	`int`	3
`val`	`int *`	*x*
`val+1`	`int *`	*x* + 4
`&val[2]`	`int *`	*x* + 8
`val[5]`	`int`	??
`*(val+1)`	`int`	5
`val + i`	`int *`	*x* + 4 *i*

Array Example

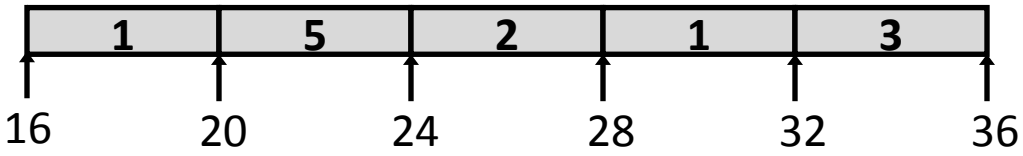
```
#define ZLEN 5
typedef int zip_dig[ZLEN];

zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```



- **Declaration "`zip_dig cmu`" equivalent to "`int cmu[5]`"**
- **Example arrays were allocated in successive 20 byte blocks**
  - Not guaranteed to happen in general

Array Accessing Example

```
zip_dig cmu;
```




```
int get_digit
  (zip_dig z, int dig)
{
  return z[dig];
}
```

IA32

```
  # %edx = z
  # %eax = dig
   movl (%edx,%eax,4),%eax  # z[dig]
```

- Register `%edx` contains starting address of array
- Register `%eax` contains array index
- Desired digit at `4*%eax + %edx`
- Use memory reference `(%edx,%eax,4)`

Array Loop Example (IA32)

```
void zincr(zip_dig z) {
  int i;
  for (i = 0; i < ZLEN; i++)
    z[i]++;
}
```

```
  # edx = z
   movl  $0, %eax            #   %eax = i
.L4:                         # loop:
   addl  $1, (%edx,%eax,4)   #   z[i]++
   addl  $1, %eax            #   i++
   cmpl  $5, %eax            #   i:5
   jne   .L4                 #   if !=, goto loop
```

Pointer Loop Example (IA32)

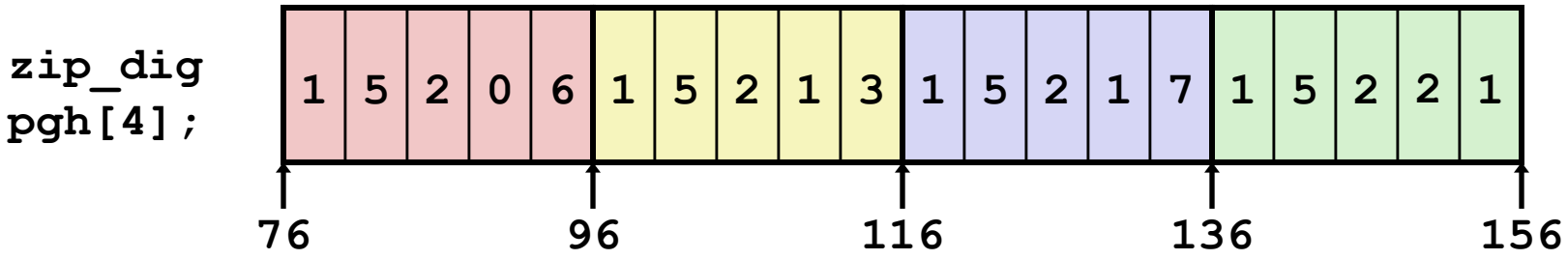
```
void zincr_p(zip_dig z) {
  int *zend = z+ZLEN;
  do {
    (*z)++;
    z++;
  } while (z != zend);
}
```

```
void zincr_v(zip_dig z) {
  void *vz = z;
  int i = 0;
  do {
    (*((int *) (vz+i)))++;
    i += ISIZE;
  } while (i != ISIZE*ZLEN);
}
```

```
  # edx = z = vz
  movl  $0, %eax            #   i = 0
.L8:                        # loop:
  addl  $1, (%edx,%eax)     #   Increment vz+i
  addl  $4, %eax            #   i +=  4
  cmpl  $20, %eax           #   Compare i:20
  jne   .L8                 #   if !=, goto loop
```

Nested Array Example

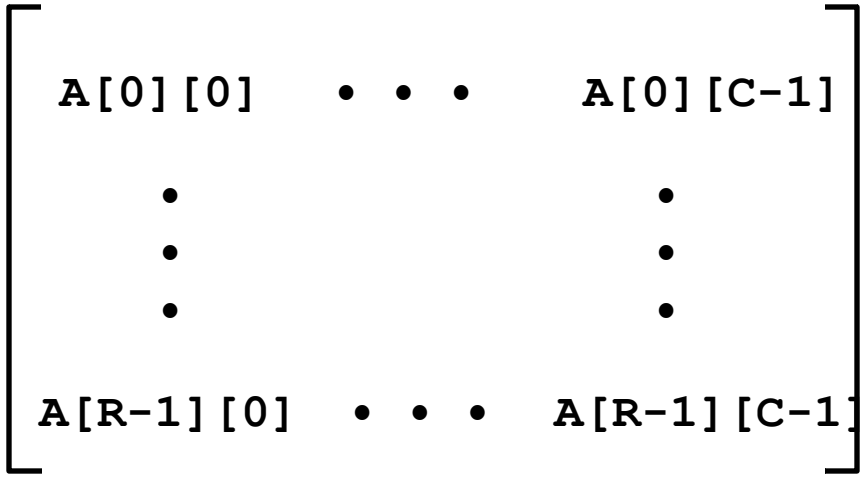
```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
  {{1, 5, 2, 0, 6},
   {1, 5, 2, 1, 3 },
   {1, 5, 2, 1, 7 },
   {1, 5, 2, 2, 1 }};
```



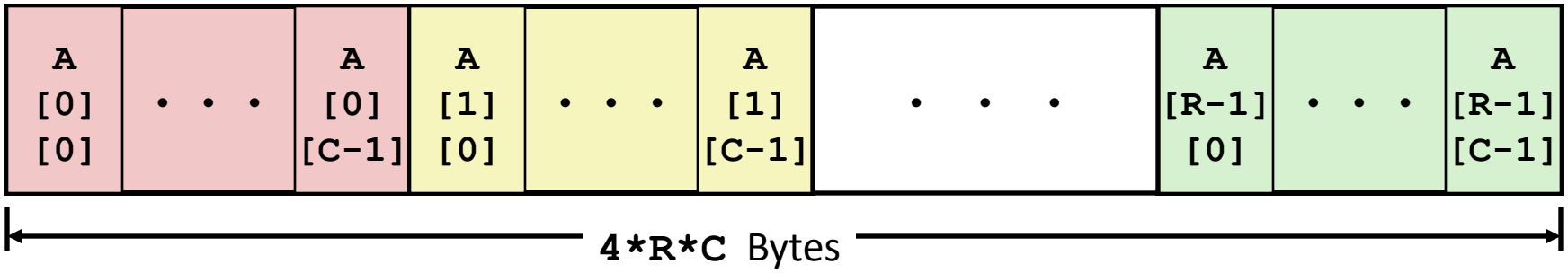
- **"`zip_dig pgh[4]`" equivalent to "`int pgh[4][5]`"**
  - Variable **`pgh`**: array of 4 elements, allocated contiguously
  - Each element is an array of 5 **`int`**'s, allocated contiguously
- **"Row-Major" ordering of all elements guaranteed**

Multidimensional (Nested) Arrays

- **Declaration**
  - *T* **A**[*R*][*C*];
  - 2D array of data type *T*
  - *R* rows, *C* columns
  - Type *T* element requires *K* bytes
- **Array Size**
  - *R* * *C* * *K* bytes
- **Arrangement**
  - Row-Major Ordering



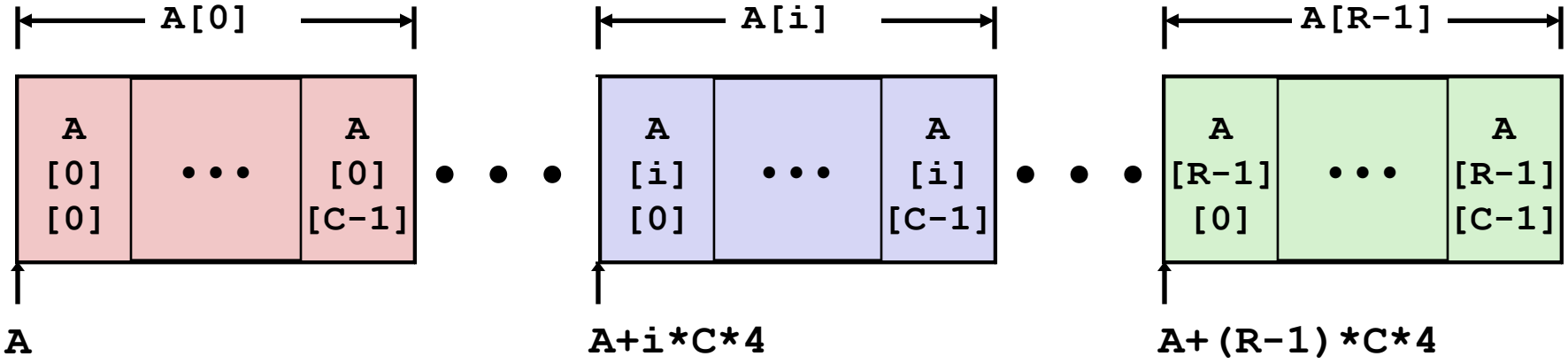
`int A[R][C];`



Nested Array Row Access

- **Row Vectors**
  - **`A[i]`** is array of *C* elements
  - Each element of type *T* requires *K* bytes
  - Starting address **`A`** **`+`** *i* * (*C* * *K*)

```
int A[R][C];
```



Nested Array Row Access Code

```
int *get_pgh_zip(int index)
{
  return pgh[index];
}
```

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
  {{1, 5, 2, 0, 6},
   {1, 5, 2, 1, 3 },
   {1, 5, 2, 1, 7 },
   {1, 5, 2, 2, 1 }};
```

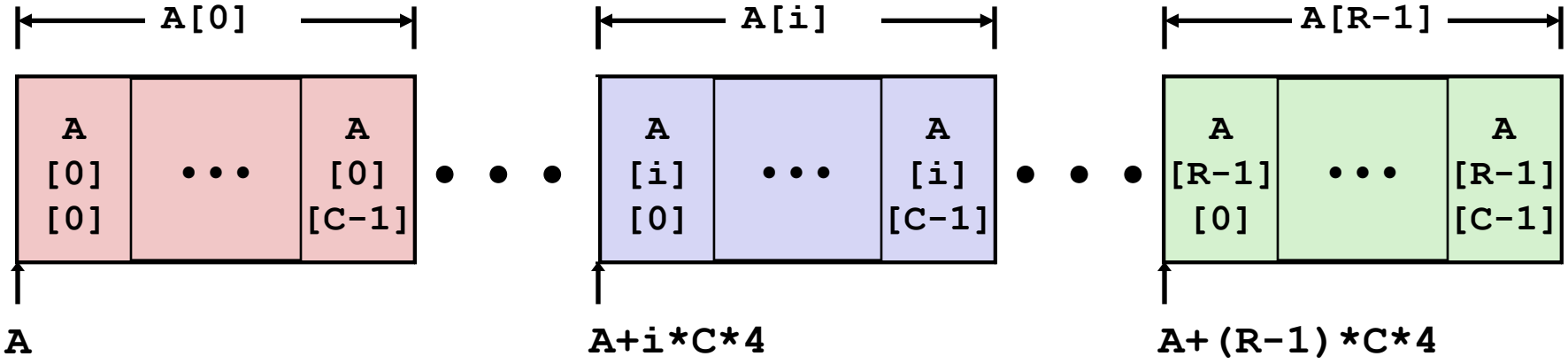
```
  # %eax = index
   leal (%eax,%eax,4),%eax # 5 * index
   leal pgh(,%eax,4),%eax  # pgh + (20 * index)
```

- **Row Vector**
  - **`pgh[index]`** is array of 5 **`int`**'s
  - Starting address **`pgh+20*index`**
- **IA32 Code**
  - Computes and returns address
  - Compute as **`pgh + 4*(index+4*index)`**

Nested Array Row Access

- **Row Vectors**
  - **`A[i]`** is array of *C* elements
  - Each element of type *T* requires *K* bytes
  - Starting address **`A`** **`+`** *i* * (*C* * *K*)

```
int A[R][C];
```



Nested Array Element Access Code

PP3.37
PP3.38

```
int get_pgh_digit
   (int index, int dig)
{
  return pgh[index][dig];
}
```

```
movl   8(%ebp), %eax          # index
leal   (%eax,%eax,4), %eax    # 5*index
addl   12(%ebp), %eax         # 5*index+dig
movl   pgh(,%eax,4), %eax     # offset 4*(5*index+dig)
```

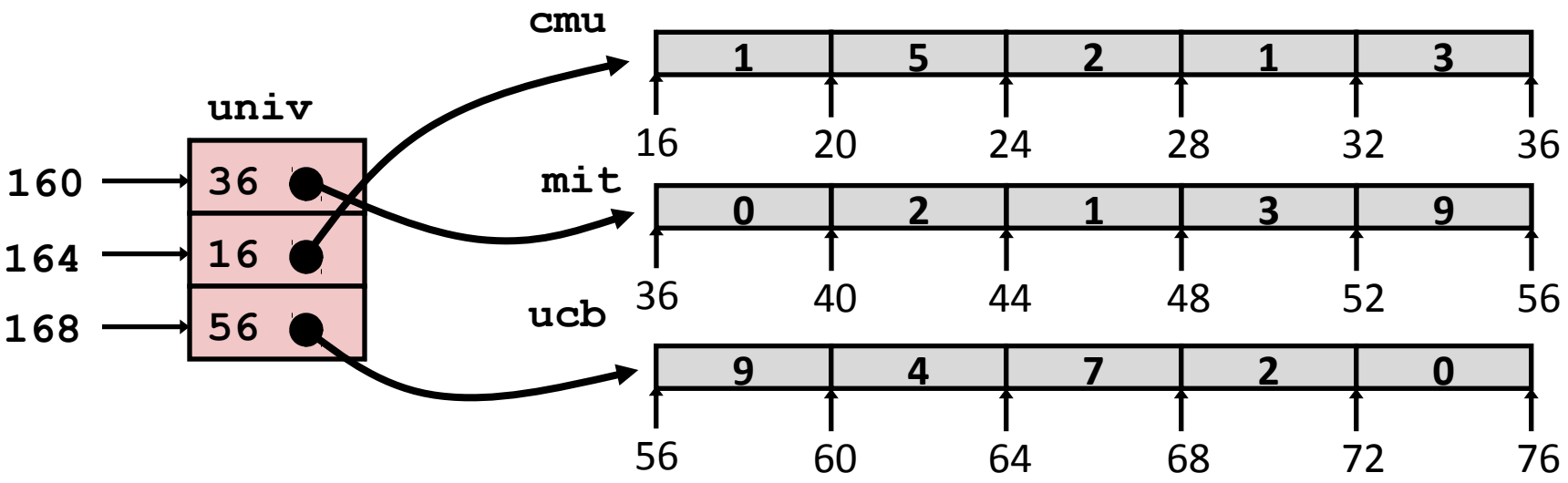
- **Array Elements**
  - `pgh[index][dig]` is `int`
  - Address: `pgh + 20*index + 4*dig`
    - `= pgh + 4*(5*index + dig)`
- **IA32 Code**
  - Computes address `pgh + 4*((index+4*index)+dig)`

Multi-Level Array Example

```
zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

```
#define UCOUNT 3
int *univ[UCOUNT] = {mit, cmu, ucb};
```

- Variable `univ` denotes array of 3 elements
- Each element is a pointer
  - 4 bytes
- Each pointer points to array of `int`'s



Element Access in Multi-Level Array

```
int get_univ_digit
   (int index, int dig)
{
  return univ[index][dig];
}
```

```
   movl   8(%ebp), %eax          # index
   movl   univ(,%eax,4), %edx    # p = univ[index]
   movl   12(%ebp), %eax         # dig
   movl   (%edx,%eax,4), %eax    # p[dig]
```

- **Computation (IA32)**
  - Element access **`Mem[Mem[univ+4*index]+4*dig]`**
  - Must do two memory reads
    - First get pointer to row array
    - Then access element within array

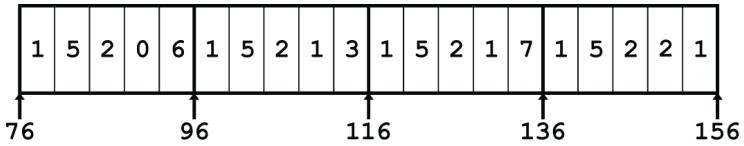
Array Element Accesses - Comparison

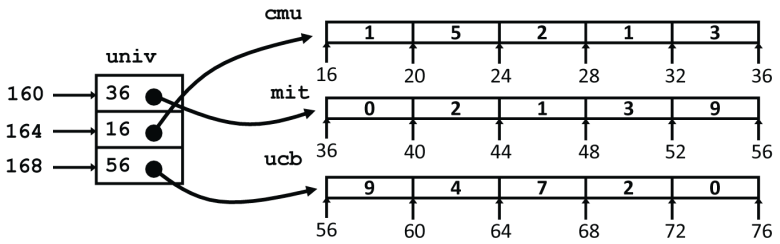
Nested array

```
int get_pgh_digit
   (int index, int dig)
{
  return pgh[index][dig];
}
```

Multi-level array

```
int get_univ_digit
   (int index, int dig)
{
  return univ[index][dig];
}
```





Accesses looks similar in C, but addresses very different:

`Mem[pgh+20*index+4*dig]`

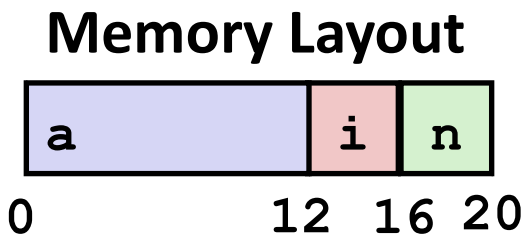
`Mem[Mem[univ+4*index]+4*dig]`

Today

- Switch statements *
- IA 32 Procedures
  - Stack Structure
  - Calling Conventions
  - Illustrations of Recursion & Pointers
- Arrays
- (Structs)

Structure Allocation

```
struct rec {
  int a[3];
  int i;
  struct rec *n;
};
```

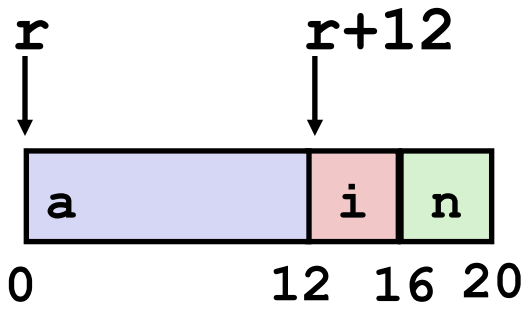


- **Concept**
  - Contiguously-allocated region of memory
  - Refer to members within structure by names
  - Members may be of different types

PP3.39

Structure Access

```
struct rec {
  int a[3];
  int i;
  struct rec *n;
};
```



- **Accessing Structure Member**
  - Pointer indicates first byte of structure
  - Access elements with offsets

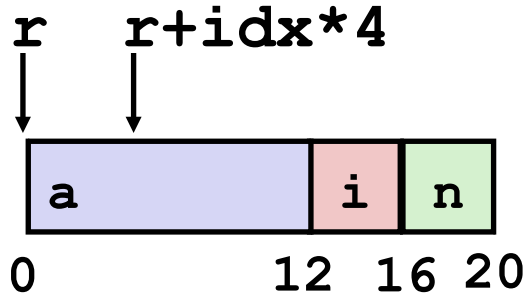
```
void
set_i(struct rec *r,
      int val)
{
  r->i = val;
}
```

IA32 Assembly

```
# %edx = val
# %eax = r
movl %edx, 12(%eax)  # Mem[r+12] = val
```

Generating Pointer to Structure Member

```
struct rec {
  int a[3];
  int i;
  struct rec *n;
};
```



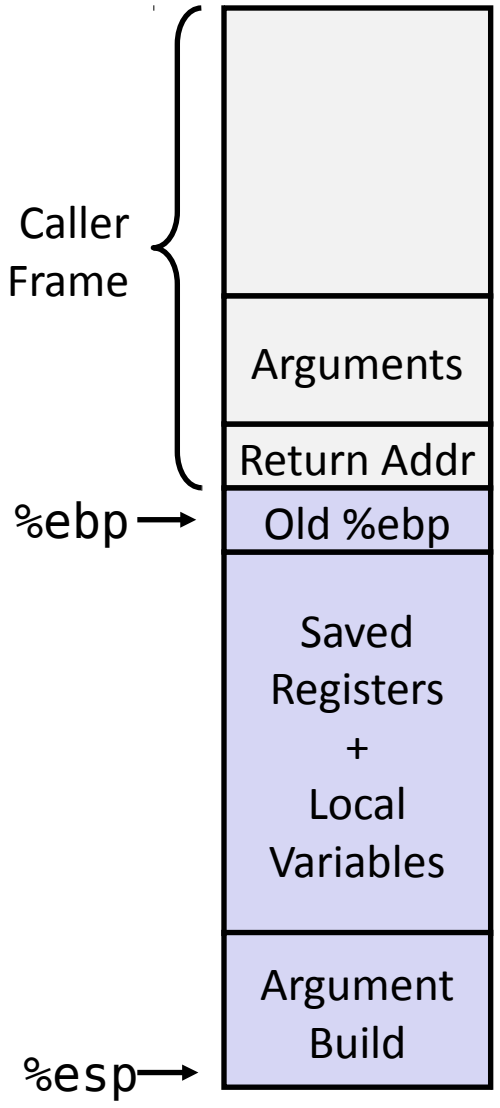
- **Generating Pointer to Array Element**
  - Offset of each structure member determined at compile time
  - Arguments
    - Mem[%ebp+8]: **r**
    - Mem[%ebp+12]: **idx**

```
int *get_ap
 (struct rec *r, int idx)
{
  return &r->a[idx];
}
```

```
movl   12(%ebp), %eax   # Get idx
sall   $2, %eax         # idx*4
addl   8(%ebp), %eax    # r+idx*4
```

IA 32 Procedure Summary

- **Important Points**
  - Stack is the right data structure for procedure call / return
    - If P calls Q, then Q returns before P
- **Recursion (& mutual recursion) handled by normal calling conventions**
  - Can safely store values in local stack frame and in callee-saved registers
  - Put function arguments at top of stack
  - Result return in `%eax`
- **`Pointers are addresses of values`**
  - `On stack or global`



Array, Structs Summary

- **Arrays**
  - One-dimensional
  - Multi-dimensional (nested)
  - Multi-level
  - You can count on your compiler!
- **Structures**
  - Allocation
  - Access