

pthread_mutex_lock();	pthread_mutex_trylock();	pthread_mutex_unlock();
if (mutex == 1)	0 →	if (mutex == 0)
mutex--;	EBUSY (16)	release a blocked thread;
else	EINVAL (22)	else
block the thread;		mutex++;

CPU
 ↙
 the producer-consumer problem

Semaphore $e = n, f = 0$; # of buffers in buffer
 binary_semaphore b ; ← initialized to 1 item type in_item, out_item,
 int in = 0, out = 0; ← position of buffers for p & c

producer:

```

while(1) {
    produce an in_item;
    P(e);
    bP(b);
    buffer[in] = in_item;
    in = (in + 1) % n;
    bV(b);
    V(f);
}
  
```

consumer:

```

while(1) {
    P(f);
    bP(b);
    out_item = buffer[out];
  
```

```

out = (out+1)%n;
bV(b);
V(e);
}

```

reader's - writer's problem

Solution Favors reader:

```

binary_semaphore mutex;
binary_semaphore rw_mutex;
int reader_count;
reader:

```

```

while(1){
    bp(mutex);
    reader_count++;
    if (reader_count==1)
        bp(rw_mutex);

```

```

    bV(mutex);
    .....

```

```

/* reading occurs here */
.....

```

```

    bp(mutex);
    reader_count--;
    if (reader_count==0)
        bV(rw_mutex);

```

```

    bV(mutex);

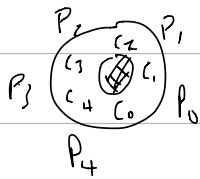
```

writer:

```
while(1){  
    bP(vw_mutex);  
    .....  
    /* writing occurs here */  
    .....  
    bV(vw_mutex);  
}
```

$R1 \rightarrow W1 \rightarrow R2 \rightarrow K3$

Dining philosopher problem



Semaphore chopsticks $[S] = \{1, 1, 1, 1\}$;

```
while(1){
```

.....

feeling hungry;

$p(chopstick[i]);$

$p(chopstick[(i+1)\%5]);$

.....

/* eating here */

.....

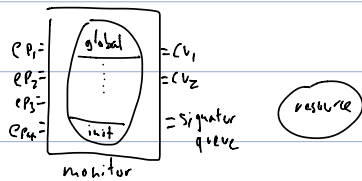
$V(chopstick[i]);$

$V(chopstick[(i+1)\%5]);$

/* thinking here */

}

monitor (ADT) monitor classes in C++ (NO SYNCHRO)



bounded_buffer problem:

monitor Bounded_Buffer {

char buffer[n];

int nextin = 0, nextout = 0, full_out = 0;

condition notempty, notFull;

deposit(char c) {

}

remove(char c) {

}

}