

CSE 477: Introduction to Computer Security

Lecture – 4

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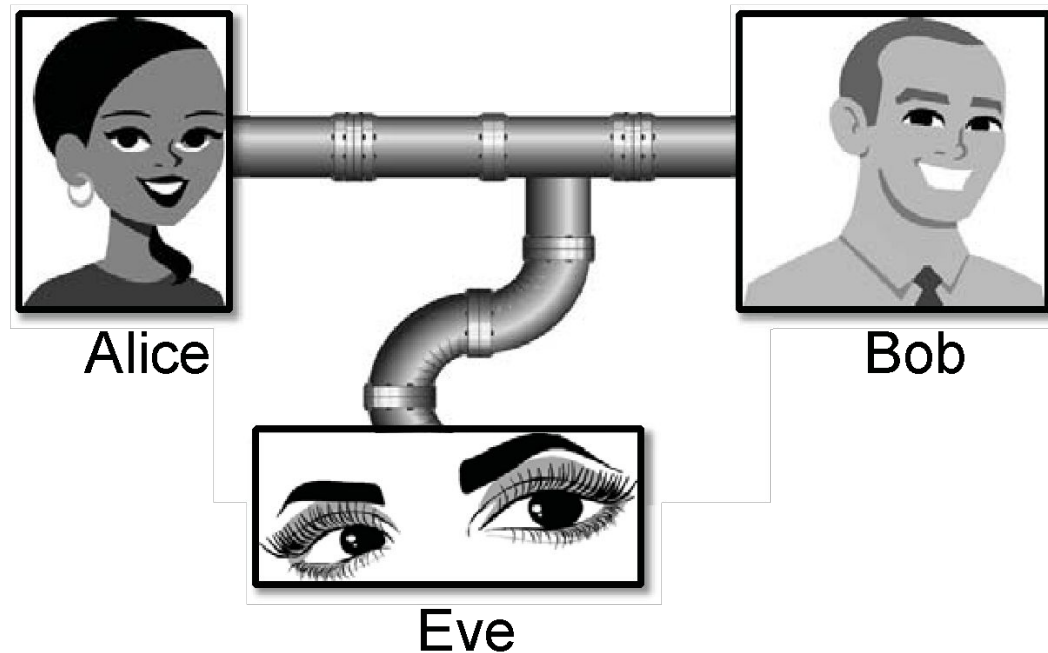
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Outline

- Basic crypto concepts

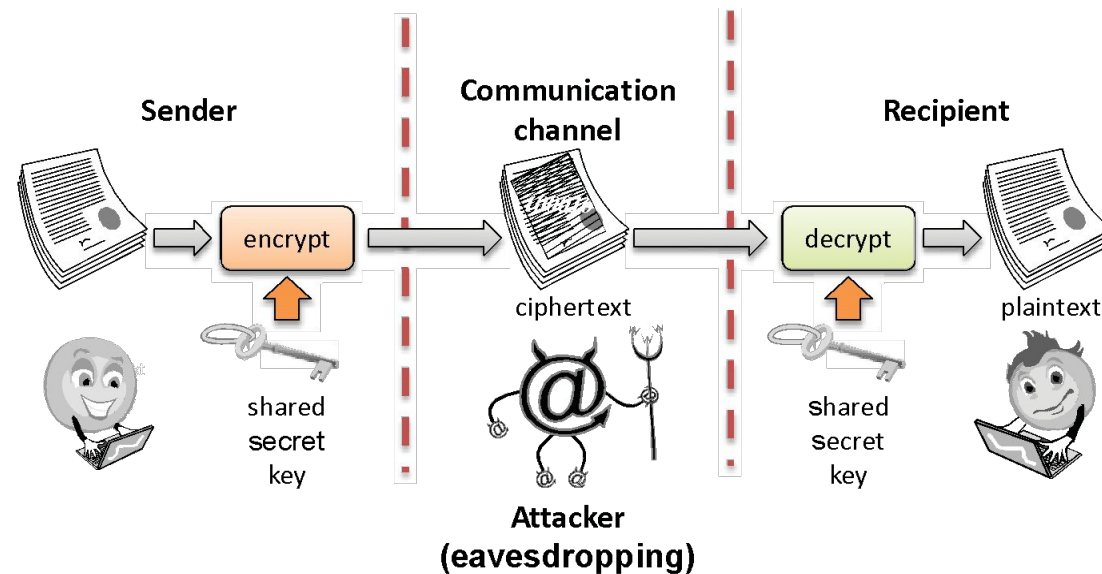
Encryption

- A means to allow two parties, customarily called Alice and Bob, to establish confidential communication over an insecure channel that is subject to eavesdropping



Encryption

- The message M is called the **plaintext**
- Alice will convert plaintext M to an encrypted form using an encryption algorithm E that outputs a **ciphertext C for M**



Encryption

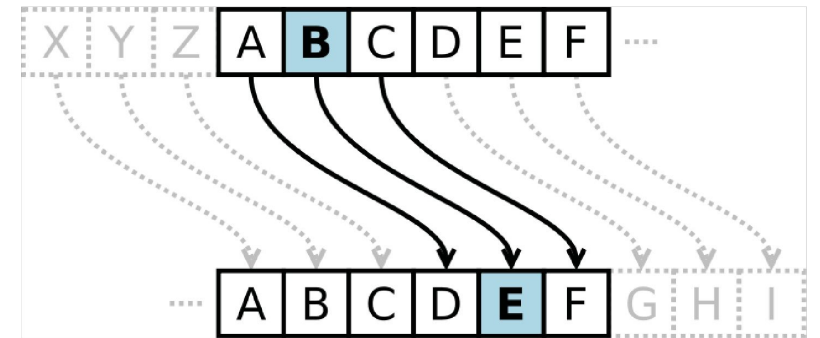
- As equations:
 - $C = E(M)$, $M = D(C)$
- The encryption and decryption algorithms are chosen so that it is infeasible for someone other than Alice and Bob to determine plaintext M from ciphertext C
- Thus, ciphertext C can be transmitted over an insecure channel that can be eavesdropped by an adversary
- The decryption algorithm must use some secret information known to Bob, and possibly also to Alice, but no other party
 - using an auxiliary input a secret number or string called ***decryption key***
 - the decryption algorithm itself can be implemented by standard, publicly available software and only the decryption key needs to remain secret
- Similarly, the encryption algorithm uses as auxiliary input an ***encryption key***, which is associated with the decryption key
- If it is feasible to derive the decryption key from the encryption key, the encryption key should be kept secret as well

Cryptosystem

- A ***cryptosystem*** consists of seven components:
 - The set of possible plaintexts
 - The set of possible ciphertexts
 - The set of encryption keys
 - The set of decryption keys
 - The correspondence between encryption keys and decryption keys
 - The encryption algorithm to use
 - The decryption algorithm to use

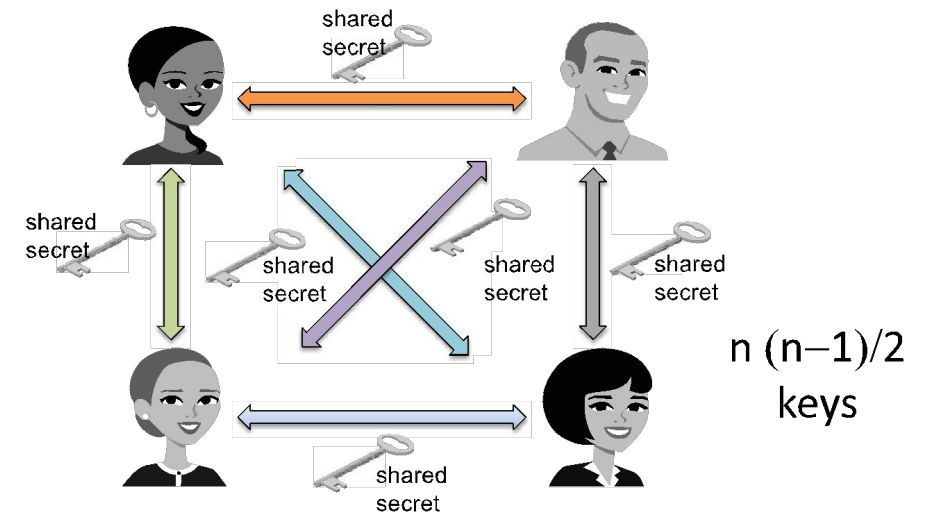
Cryptosystem: Caesar cipher

- Replace each letter with the one “three over” in the alphabet
 - An example of a shift cipher
- Can be denoted using the following formula:
 - $s(c, k)$, here c represents one of 23 letters in Latin Alphabet and k represents the key
 - If $k > 0$, forward shift (encryption) and if $k < 0$, backward shift (decryption)
 - Example: $s(A, 3) = D$, $s(D, -3) = A$
- For Caesar cipher, $\{k=3\}$ is the set of encryption key and $\{k=-3\}$ is the set of decryption key
- Uses modulo operation in a sort of circular shift fashion when shift goes beyond the limit



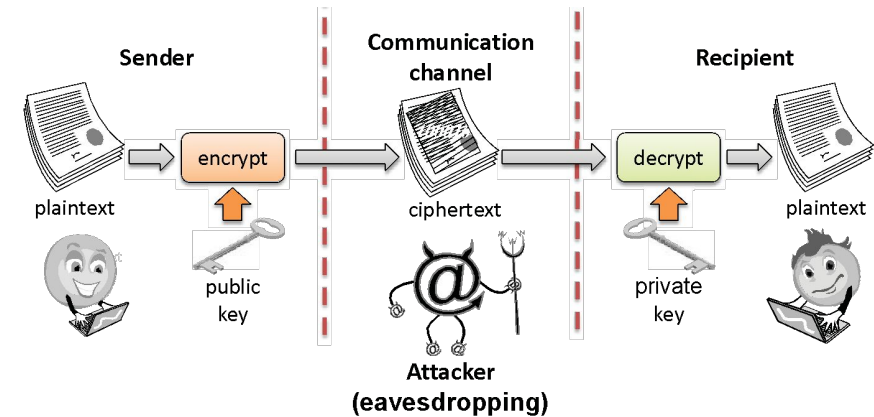
Symmetric key distribution

- Symmetric crypto-systems are quite fast and efficient
- However their main problem is key distribution
- It requires each pair of communicating parties to share a (separate) secret key
- If there are n parties, it means it requires a total of $n(n - 1)/2$ keys



Public key cryptography

- Bob has two keys: a **private key**, S_B , which Bob keeps secret, and a **public key**, P_B , which Bob broadcasts widely
- In order for Alice to send an encrypted message to Bob, she needs only obtain Bob's public key, P_B , and use that to encrypt her message, M , and send the result, $C = E_{P_B}(M)$, to Bob.
- Bob then uses his secret key to decrypt the message as $M = D_{S_B}(C)$
- That is, separate keys are used for encryption and decryption

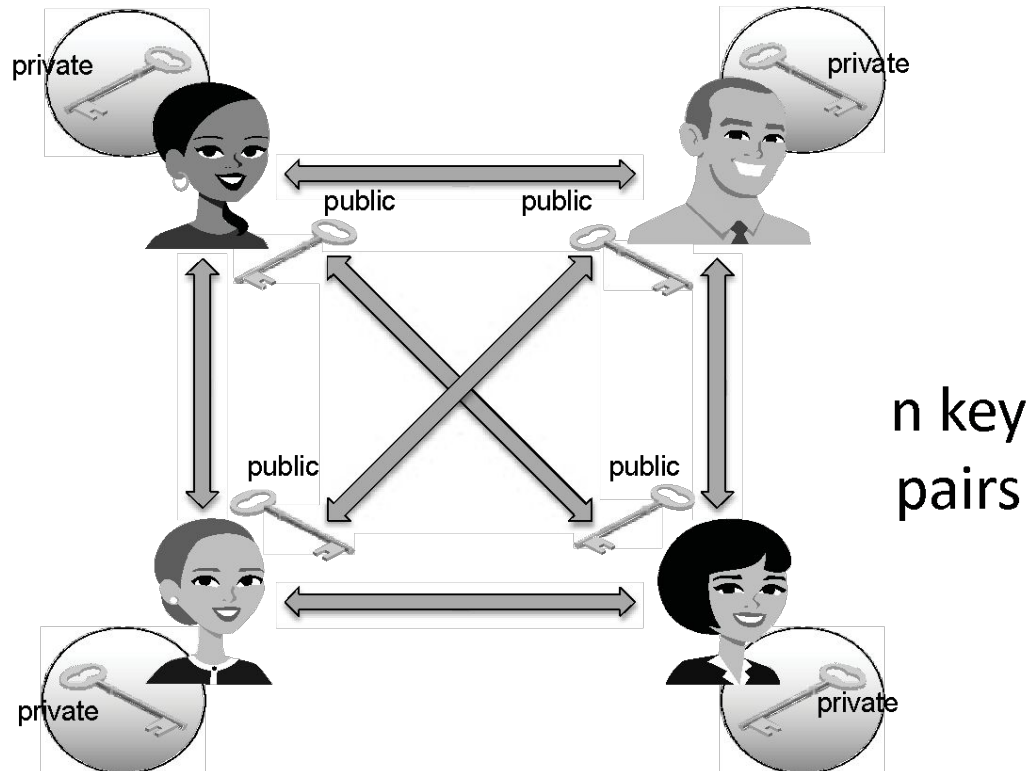


Public key distribution

- How many keys do we need to distribute for n users??

Public key distribution

- Only one key is needed for each recipient!

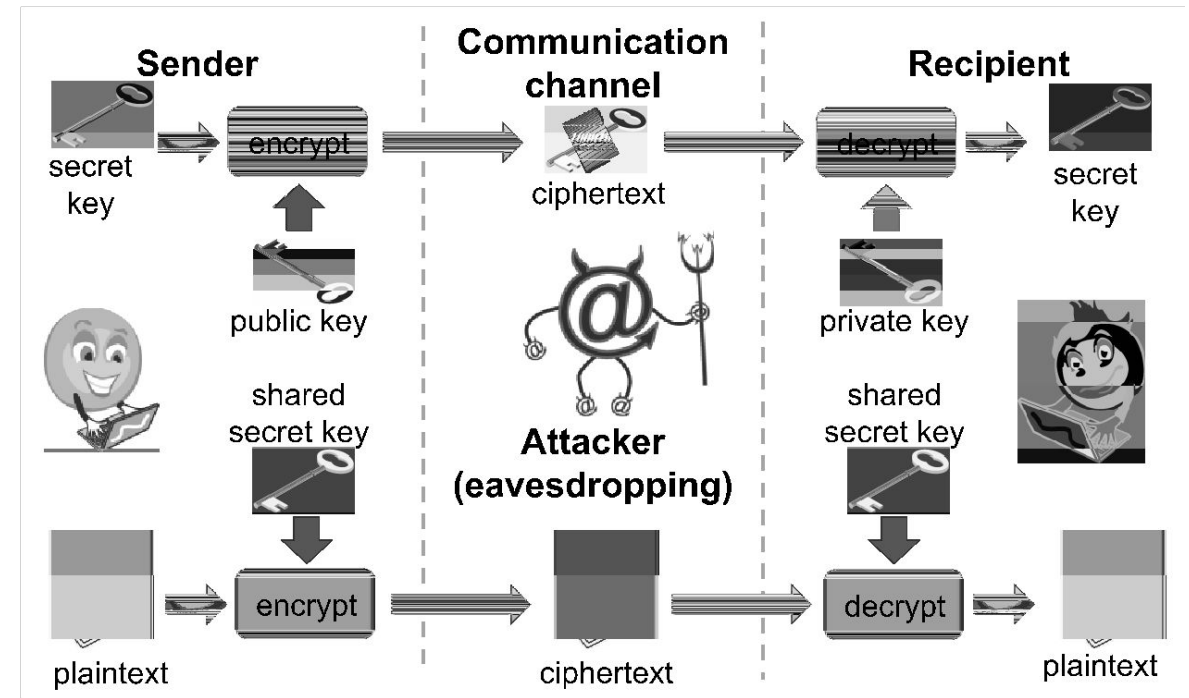


Combining symmetric and public key systems

- The main disadvantage of public-key cryptosystems is that they are much slower than existing symmetric encryption schemes
 - Hence, public-key cryptography is unsuitable for interactive sessions that use a lot of back-and-forth communication
- Also, public-key cryptosystems have larger key lengths than that for symmetric cryptosystems
 - For example, RSA is commonly used with 2,048-bit keys while AES is typically used with 256-bit keys

Combining symmetric and public key systems

- To work around these disadvantages,
 - public-key cryptosystems are often used in practice just to allow Alice and Bob to exchange a shared secret key
- Once exchanged, the shared secret key then can be used for communication encrypted with a symmetric encryption scheme

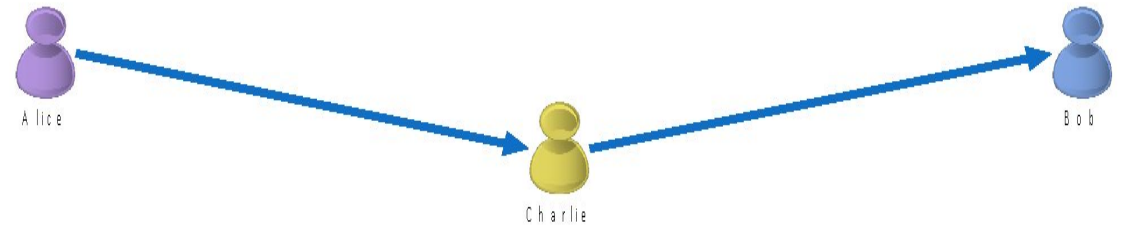


Digital signature

- Public key crypto systems allow the reversing the order in which encryption and decryption is carried out
 - $E_{P_B}(D_{S_B}(M)) = M$
- This concept is leveraged to create digital signatures
- To sign a message, M , Alice just decrypts it with her private key, S_A , creating $C = D_{S_A}(M)$
- Anyone can encrypt this message using Alice's public key, as $M' = E_{P_A}(C)$ and then compare if $M = M'$
 - If they match, the signature is valid
- Indeed, no one but Alice, who has private key S_A , could have produced such an object C , so that $M' = E_{P_A}(C)$
- Only disadvantage, the signature is as long as the message!
 - Compare it with real life signatures!

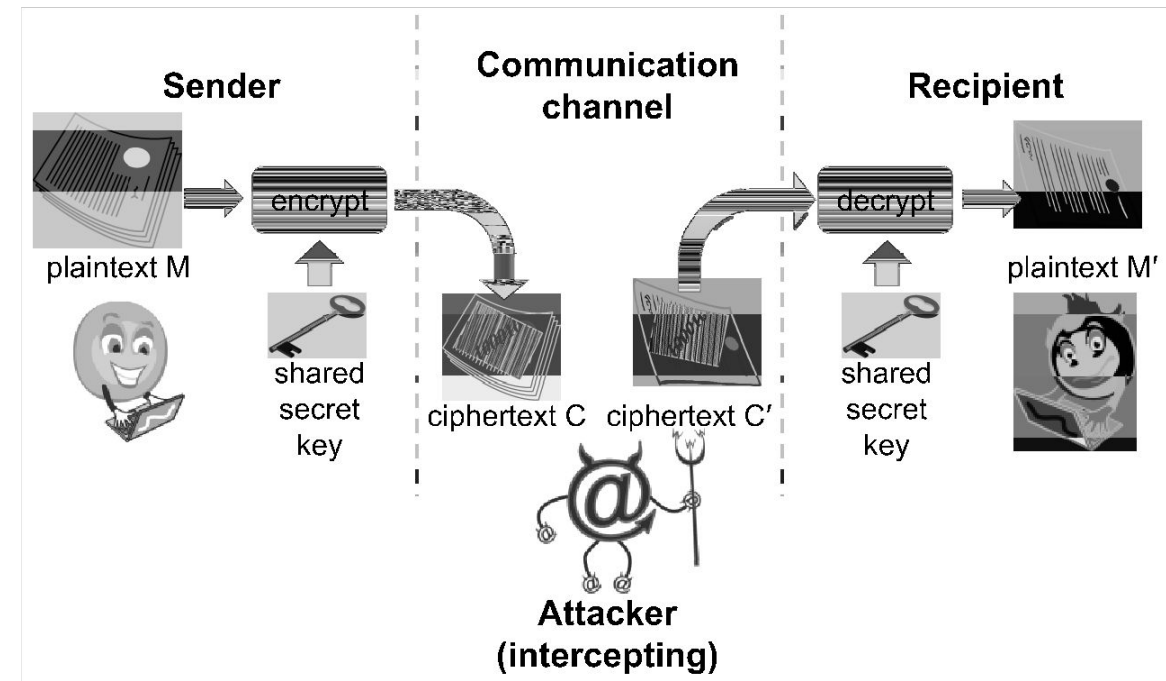
Man In The Middle (MITM) Attack

- Charlie is in the middle between Alice and Bob
- Charlie can:
 - View traffic
 - Change traffic
 - Add traffic
 - Delete traffic
- Charlie could be:
 - Internet service provider
 - Virtual Private Network (VPN) provider
 - WIFI provider such as a coffee shop
 - An attacker re-routing your connection
 - An incompetent admin (it happens)



Simple attacks on crypto systems: MITM

- The straightforward use of a cryptosystem presented, which consists of simply transmitting the ciphertext, assures confidentiality
- However, it does not guarantee the authenticity and integrity of the message if the adversary can intercept and modify the ciphertext
- How the recipient can be assured that the message s(he) receives is the intended one?



Man-In-The-Middle

Simple attacks on crypto systems: MITM

- Bob (M, S) : $(C, S) \rightarrow$ Eve $(C', S') \rightarrow$ Alice (C', S') : (M', S') ??
- Note that M' will be different from the original message M
 - When Alice verifies the digital signature S' , she obtains message M' by encrypting S'
 - Thus, Alice is led to believe that Bob has signed M' instead of M
- Note that in the above attacks the adversary can arbitrarily alter the transmitted ciphertext or signature
- However, the adversary cannot choose, or even figure out, what would be the resulting plaintext since she does not have the ability to decrypt
- Thus, the above attacks are effective only if any arbitrary sequence of bits is a possible message
- This scenario occurs, for example, when a randomly generated symmetric key is transmitted encrypted with a public-key cryptosystem

Simple attacks on crypto systems: brute-force

- How about brute forcing a crypto system?
 - Trying different probable keys over a cipher text to decrypt it to a meaningful text
- Mind you, encryption and decryption functions are open, just the key is secret!
- If the plaintext is an arbitrary binary string, this attack cannot succeed, as there is no way for the attacker to distinguish a valid message
- However, if the plaintext is known to be text in a natural language, then the adversary hopes that only a small subset of the decryption results (ideally just a single plaintext) will be a meaningful text for the language

Brute force attack

- Some knowledge about the possible message being sent will then help the attacker pinpoint the correct plaintext
- Key should be a sufficiently long random value to make exhaustive search attacks unfeasible
- Problem is it is usually easy to recognize that a message is a valid plaintext
 - For example, given a certain ciphertext, if an attacker could decrypt it with a given key and get message NGGNPXNGQNJABAVEIVARORNPU, which she can immediately dismiss
 - But if she gets message ATTACKATDAWNONIRVINEBEACH, then she can be confident she has found the decryption key
- This ability is related to the ***unicity distance*** for a cryptosystem

Brute force attack

- English text is typically represented with 8-bit ASCII encoding
- A message with t characters has n bits, with $n = 8t$
- The total number of possible n -bit (or t -byte) arrays is $2^{8t} = 2^n$
- It is estimated that each character of English text carries about 1.25 bits of information, known as the entropy of English
- The number of t -byte arrays that correspond to English text:
 - $2^{1.25t} = 2^{1.25n/8} \approx 2^{0.16n}$

Brute force attack

- In general, for some constant $0 < \alpha < 1$, there are $2^{\alpha n}$ valid text messages among the 2^n possible plaintexts
 - as not all combinations are valid messages
- The probability that a randomly selected plaintext corresponds to meaningful text is represented with: $\frac{2^{\alpha n}}{2^n} = \frac{1}{2^{(1-\alpha)n}}$
- The fraction of valid messages tends rapidly to zero as n grows

Brute force attack

- Let k be the length (number of bits) of the decryption key
- For a given ciphertext, there are 2^k possible plaintexts, each corresponding to a key
- From the previous discussion, each such plaintext is a valid text message with probability $\frac{1}{2^{(1-\alpha)n}}$
- Hence, the expected number of plaintexts corresponding to valid text messages is $\frac{2^k}{2^{(1-\alpha)n}}$
- As the key length k is fixed, the above number tends rapidly to zero as the ciphertext length n grows

Brute force attack

- We expect that there is a unique valid plaintext for the given ciphertext when $n = \frac{k}{1-\alpha}$
- The above threshold value for n is called the **unicity distance** for the given language and key length
- The **unicity distance** for a cryptosystem is the minimum number of characters of ciphertext that are needed so that there is a single intelligible plaintext associated with it
- Consider a cipher text: FJKFPO that is encrypted with a substitution cipher (discussed later) and brute force might result in the following: ***that is of your season on your that we the top the two oxford that in that of***
 - Which is the correct one? It is impossible to know!

Brute force attack

- For the English language and the 256-bit AES cryptosystem, the unicity distance is about 304 bits or 38 ASCII-encoded characters
 - This is only half a line of text
- We can conclude that brute-force decryption is likely to succeed for messages in natural language that are not too short
- Namely, when a key yields a plaintext that is a meaningful text, the attacker has probably recovered the original message

**OKAY CLASS, DOES ANYONE
HAVE ANY QUESTIONS?**

