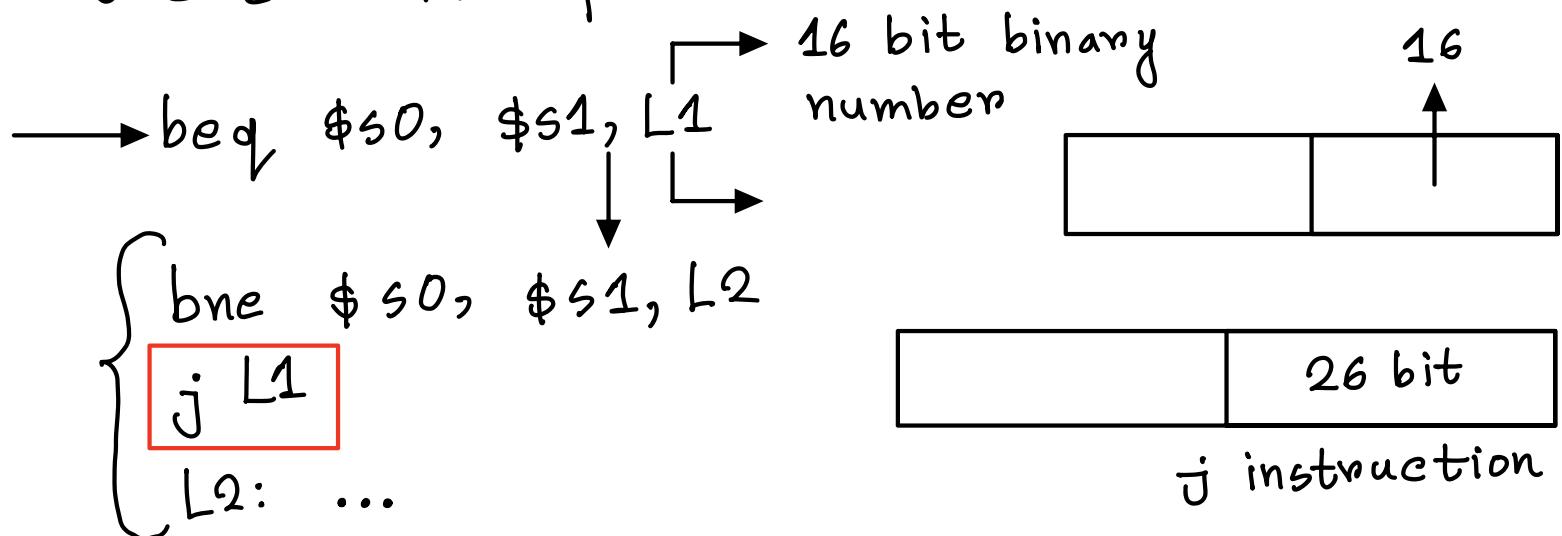


Branching far away (Slide 54)

If branch target is too far to encode with 16-bit offset, assembler rewrites the code. Example:



Slide 58 (important) (must know):

Register Usage

- \$a0 – \$a3: arguments (reg's 4 – 7)
- \$v0, \$v1: result values (reg's 2 and 3)
- \$t0 – \$t9: temporaries
 - Can be overwritten by callee
- \$s0 – \$s7: saved
 - Must be saved/restored by callee
- \$gp: global pointer for static data (reg 28)
- \$sp: stack pointer (reg 29)
- \$fp: frame pointer (reg 30)
- \$ra: return address (reg 31)

Functions:

Main () {

 int x = 0;

 int y = 5;

 int z = addition(x, y);

}

int addition (int a, int b) {

 int c = a + b;

 return c;

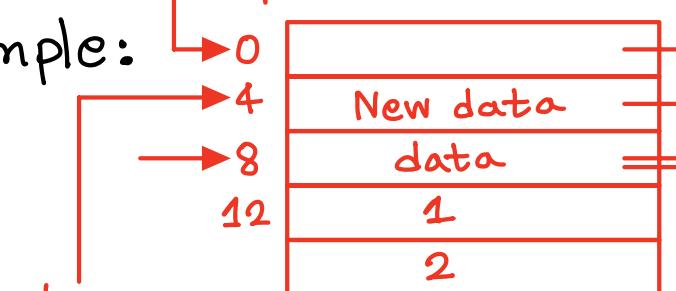
}

→ Caller function

→ Called function

\$SP - 4

\$S1 → Top Data of the Stack



Leaf Procedure Example:

Main () {

 int f = 0;

 f = f + 1; → \$SD

\$SD

 int z = leaf_example(1, 2, 3, 4);

}

 int y = f + z;

 int leaf_example (int g, h, i, j) {

 int f;

 → f = (g + h) - (i + j);

 return f;

Stack → LIFO

\$SP holds
address of
top data.
jal → jump
and link

Arguments

g	→ \$a0
h	→ \$a1
i	→ \$a2
j	→ \$a3

f → \$s0 (hence, need to save \$s0 on stack)

Result in \$v0

z, y in \$s1, \$s2

Answer: Mips Code:

add \$s0, \$zero, \$zero

addi \$s0, \$s0, 1

jal leaf-example

add \$s1, \$v0, \$zero

add \$s2, \$s0, \$s1

leaf-example:

addi \$sp, \$sp, -4

sw \$s0, 0(\$sp) → sw \$s0, 0(\$sp)

add \$t0, \$a0, \$a1 → g+h

add \$t1, \$a2, \$a3 → i+j

sub \$s0, \$t0, \$t1 → (g+h)-(i+j)

add \$v0, \$s0, \$zero

lw \$s0, 0(\$sp)

addi \$sp, \$sp, 4

jr \$ra



$f = \$S0$

$g = \$S2$

addi \$SP, \$SP, -4

SW \$S0, 0(\$SP)

addi \$SP, \$SP, -4

SW \$S2, 0(\$SP)

lw \$S2, 0(\$SP)

addi \$SP, \$SP, 4

lw \$S2, 0(\$SP)

addi \$SP, \$SP, 4

Top এর Data রাখার
জন্য

Slide 68:

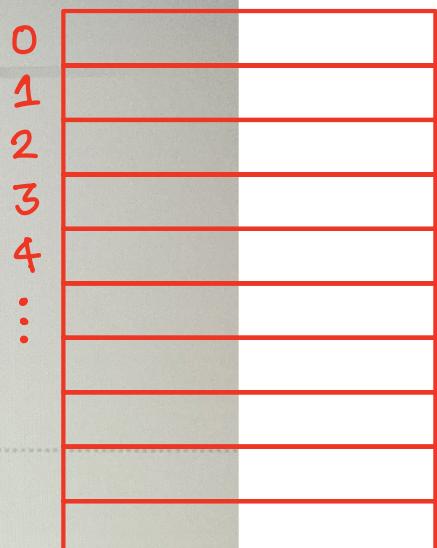
Slides → 68, 69, 70

Character Data

- Byte-encoded character sets
 - ASCII: 128 characters
 - 95 graphic, 33 control
 - Latin-1: 256 characters
 - ASCII, +96 more graphic characters

Unicode: 32-bit character set

- Used in Java, C++ wide characters, ...
- Most of the world's alphabets, plus symbols
- UTF-8, UTF-16: variable-length encodings



Byte/Halfword Operations

- Could use bitwise operations
- MIPS byte/halfword load/store
 - String processing is a common case

lb rt, offset(rs) lh rt, offset(rs)

- Sign extend to 32 bits in rt

lbu rt, offset(rs) lhu rt, offset(rs)

- Zero extend to 32 bits in rt

sb rt, offset(rs) sh rt, offset(rs)

- Store just rightmost byte/halfword

lb rt, offset(rs) → Fetching 8 bit data

000... 0000...	8 bit
----------------	-------

lh rt, offset(rs)

ASCII characters
are always unsigned.

String Copy Example

- C code (naïve):

- Null-terminated string

```
void strcpy (char x[], char y[])
```

```
{ int i;
```

```
    i = 0;
```

```
    while ((x[i] = y[i]) != '\0')
```

```
        i += 1;
```

- Addresses of x, y in \$a0, \$a1

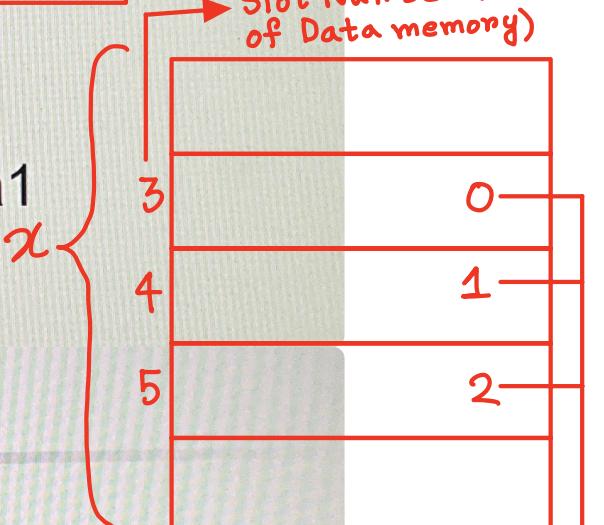
- i in \$s0

String Copy Example

- MIPS code:

strcpy:	
addi	\$sp, \$sp, -4 # adjust stack for 1 item
sw	\$s0, 0(\$sp) # save \$s0
add	\$s0, \$zero, \$zero # i = 0
L1: add	\$t1, \$s0, \$a1 # addr of y[i] in \$t1
lbu	\$t2, 0(\$t1) # \$t2 = y[i]
add	\$t3, \$s0, \$a0 # addr of x[i] in \$t3
sb	\$t2, 0(\$t3) # x[i] = y[i]
beq	\$t2, \$zero, L2 # exit loop if y[i] == 0
addi	\$s0, \$s0, 1 # i = i + 1
j	L1 # next iteration of loop
L2: lw	\$s0, 0(\$sp) # restore saved \$s0
addi	\$sp, \$sp, 4 # pop 1 item from stack
jr	\$ra # and return

Slot Number (Address of Data memory)



Indexing
of *x* []
array.

sw \$s0, 0(\$sp)

while ((x[i] = y[i]) != '\0')

L1: add \$t1, \$s0, \$a1

A[B[i]] = x

lbu \$t2, 0(\$t1)

> A[B[i]] = x

Base address of A and B are in \$s1 and \$s2, x and i are in \$s3 and \$s4

Ans. sll \$t0,\$s4,2

add \$t1, \$s2,\$t0

lw \$t2,0(\$t1)

sll \$t0,\$t2,2

add \$t1, \$s1,\$t2

sw \$s3,0(\$t1)

I

A[B[i]] = x

B[i] = 5

A[5] = x

32 bit architecture = $32/8 = 4$ slots

64 bit architecture = $64/8 = 8$ slots

256 bit architecture = $256/8 = 32$ slots

A[6] in 64 bit architecture

lw \$t0, 48(\$s1)