

1.1 a) A, B - finite sets

$|A \times B|$ - number of elements in cart. product

$$P(A \times B) = 2^{|A \times B|}$$

b)

$$M = \{1, 2, \dots, 8\}$$

is-divisor-of relation:

$(1, 2), (1, 3), (1, 4), (1, 5), (1, 6), (1, 7), (1, 8)$

$(2, 2), (2, 4), (2, 6), (2, 8)$

$(3, 3), (3, 6)$

$(4, 4), (4, 8)$

$(5, 5), (6, 6), (7, 7), (8, 8)$

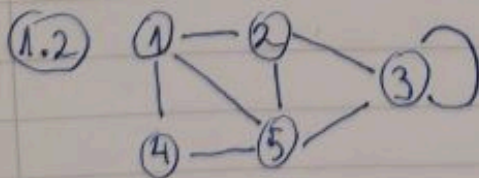
c) can these relationships be modeled as function/relation

- child-mother relationship (function)

is-uncle-of relationship (relation)

relationship between students - enrolment num. (function)

- || - || - courses (relation)



a) $G = (V, E)$

$$V = \{1, 2, 3, 4, 5\}$$

$$E = \{(1, 2), (1, 4), (1, 5), (2, 5), (2, 3), (3, 3), (3, 5), (4, 5)\}$$

// $(a, b) = (b, a)$ - note because it's undirected graph

b)

	1	2	3	4	5
1	0	1	0	1	1
2	1	0	1	0	1
3	0	1	1	0	1
4	1	0	0	0	1
5	1	1	1	1	0

c)

1 $[2, 4, 5]$

2 $[1, 3, 5]$

3 $[2, 3, 5]$

4 $[1, 5]$

5 $[1, 2, 3, 4]$

d) Because it's undirected graph, we can just count ones in node's row (for node n we look at n -th row) and if there is one in G_{nn} field, we count it twice.

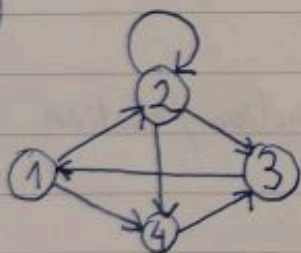
e) We count how many nodes are there in a list ^{of specific node} and if there is the same node number as the one we are looking for, we count it twice (self loop is counted twice)

1.3) $G=(V,E)$

$$V = \{1, 2, 3, 4\}$$

$$E = \{(1,2), (1,4), (2,2), (2,3), (2,4), (3,1), (4,3)\}$$

a)

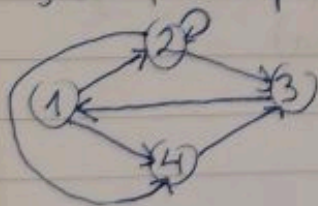


b) indegree of vertex 2 : 2 ; $(1,2), (2,2)$
out degree of vertex 2 : 3 ; $(2,4), (2,3), (2,2)$

c) acyclic?

No it is not because it contains cycle
 $(1 \rightarrow 4 \rightarrow 3 \rightarrow 1)$

c) yes, it is planar:



4-1 ✓
4-2 ✓
4-3 ✓

1-2 ✓ 2-1 ✓ 3-1 ✓
1-3 ✓ 2-3 ✓ 3-2 ✓
1-4 ✓ 2-4 ✓ 3-4 ✓

d) graph is strongly [↑] connected
because there is a path from every
node to every other node

1.5

a) example: $h-g-f-a-g-d-h-e-b-a-c-d-b$

My logic: if a node has odd degree, the path HAS to either start or finish at that node. In this graph, we have 2 nodes with odd degree so it's possible if we start at one and finish at other

b) No, it doesn't have to exist Euler tour/path because more than 2 nodes might have odd degree and still be complete.