哈尔滨工业大学

<<计算机网络>> 实验报告

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姓名:	张茗帅
学号:	1140310606
学院:	计算机科学与技术学院
教师:	聂兰顺

实验四 利用 Wireshark 进行协议分析

一、实验目的

熟悉并掌握 Wireshark 的基本操作,了解网络协议实体间进行交互以及报文交换的情况

实验环境:

- 我的个人 PC 机
- Windows10 操作系统
- Wireshark

二、实验内容

必做内容:

- (1) 学习 Wireshark 的使用
- (2) 利用 Wireshark 分析 HTTP 协议
- (3) 利用 Wireshark 分析 TCP 协议
- (4) 利用 Wireshark 分析 IP 协议
- (5) 利用 Wireshark 分析 Ethernet 数据帧

选做内容:

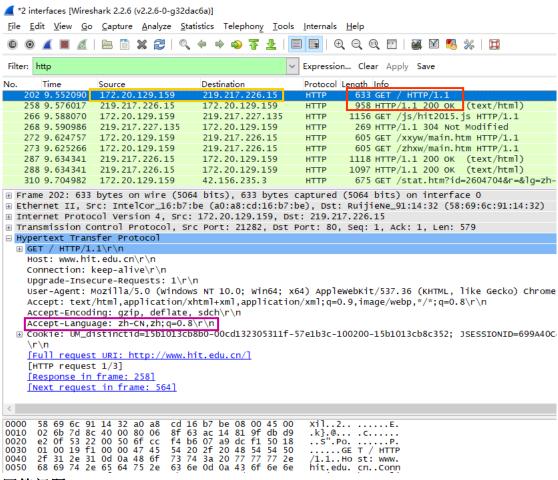
- (a) 利用 Wireshark 分析 DNS 协议
- (b) 利用 Wireshark 分析 UDP 协议
- (c) 利用 Wireshark 分析 ARP 协议

三、实验过程及结果

1 HTTP 分析

(1) HTTP GET/response 交互

打开 Wireshark 进行 capture 配置,准备进行分组俘获,然后打开浏览器,输入地址http://www.hit.edu.cn/(指导书中给的地址已经不存在了),页面加载完毕后,停止分组俘获,得到的俘获窗口内容截图如下:



回答问题:

a) 你的浏览器运行的是 HTTP1.0,还是 HTTP1.1?你所访问的服务器所运行 HTTP 协议的版本号是多少?

从上面的红框内可以看到,我的浏览器和服务器所运行的 HTTP 协议的版本号都是 HTTP 1.1

b) 你的浏览器向服务器指出它能接收何种语言版本的对象?

在紫框中可以看到,他接受的语言为 zh-CN,即简体中文(中国)

c) 你的计算机的 IP 地址是什么?服务器 http://www.hit.edu.cn 的 IP 地址是多 少?

在橙颜色的框中可以看到源 IP 地址和目的 IP 地址

我计算机的 IP 地址为:: 172.20.129.159

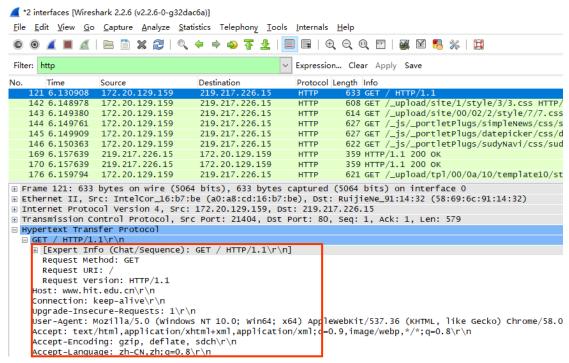
服务器 IP: 219.217.226.15

d) 从服务器向你的浏览器返回的状态代码是多少?

200

(2) HTTP 条件 GET/response 交互

截图如下:



回答问题:

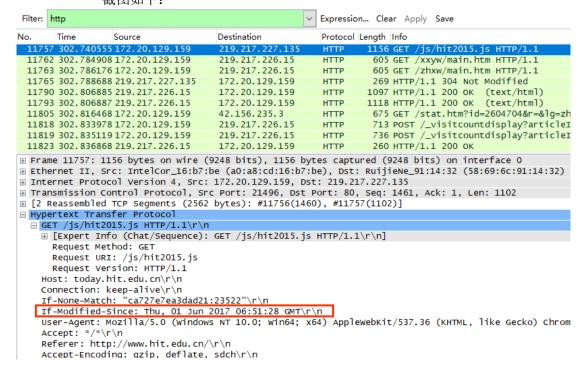
a) 分析你的浏览器向服务器发出的第一个 HTTP GET 请求的内容, 在该请求报文中,是否有一行是:IF-MODIFIED-SINCE?

没有 (在红框内没有发现)

b) 分析服务器响应报文的内容,服务器是否明确返回了文件的内容?如何获知? 服务器明确返回了内容

HTTP Status Code (状态代码) 为 304 时不明确返回文件 HTTP Status Code (状态代码) 为 200 时明确返回文件

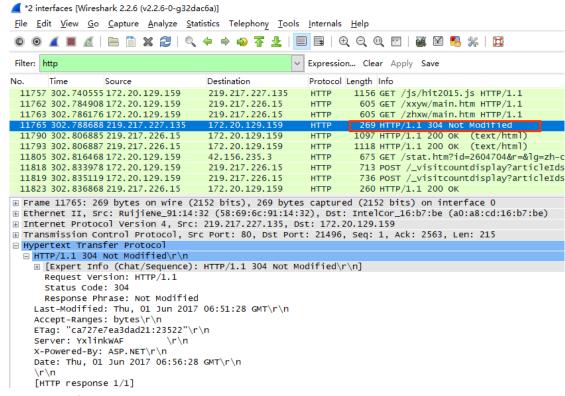
c) 分析你的浏览器向服务器发出的较晚的"HTTP GET"请求,在该请 求报文中是否有一行是:IF-MODIFIED-SINCE?如果有,在该首部行后面跟着的信息是什么? 截图如下:



如图在红框内,发现了这个字段。

这个字段后面代表的是时间,即咨询服务器在这个时候之后是否有更新

d) 服务器对较晚的 HTTP GET 请求的响应中的 HTTP 状态代码是多少?服务器是 否明确返回了文件的内容?请解释。

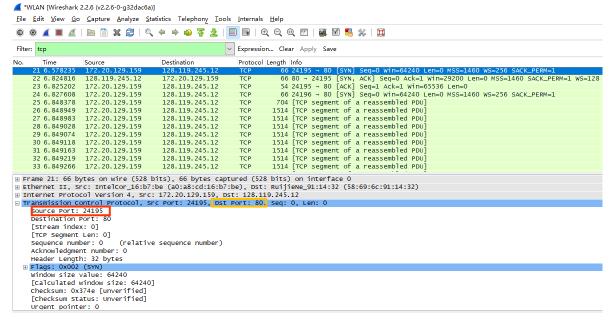


请求响应中的 HTTP 状态代码为 304。(见红色框内)

不会明确返回文件,因为根据之前 HTTP 的 GET 请求中 IF-MODIFIED-SINCE 字段内的时间服务器判断结果为 Not Modified, 于是客户端可以使用本地这个没有过期的缓存文件。

2 TCP 分析

获取的 TCP 报文整体情况截图如下



回答问题:

1) 向gaia.cs.umass.edu服务器传送文件的客户端主机的IP地址和TCP端口号是多少?

IP地址: 172.20.129.159 端口号: 24195 (见红色框内)

2) Gaia.cs.umass.edu服务器的IP地址是多少?对这一连接,它用来发送和接收TCP报文的端口号是多少?

IP 地址: 128.119.245.12 端口号: 80 (见橙色框)

3) 客户服务器之间用于初始化TCP连接的TCP SYN报文段的序号(sequence number)是多少?在该报文段中,是用什么来标示该报文段是SYN报文段的?

```
[TCP Segment Len: 0]
 Sequence number: 0
                      (relative sequence number)
 Acknowledgment number: 0
 Header Length: 32 bytes
Flags: 0x002 (SYN)
   000. .... = Reserved: Not set
   ...0 .... = Nonce: Not set
   .... 0... = Congestion Window Reduced (CWR): Not set
   .... .0.. .... = ECN-Echo: Not set
   .... ..0. .... = Urgent: Not set
   .... ...0 .... = Acknowledgment: Not set
   .... 0... = Push: Not set
   .... .... .O.. = Reset: Not set
 ± .... ...1. = Syn: Set
   .... .... 0 = Fin: Not set
   [TCP Flags: .....S.]
```

如图,初始化TCP连接的TCP SYN报文段的序号是0;

通过Flags标志位,将其中的SYN位置为1,表示该报文段是SYN报文段

4) 服务器向客户端发送的SYNACK报文段序号是多少?该报文段中,Acknowledgement 字段的值是多少?Gaia.cs.umass.edu服务器是如何决定此值的?在该报文段中,是用什么来标示该报文段是SYNACK报文段的?

```
[TCP Segment Len: 0]
 Sequence number: 0
                      (relative sequence number)
 Acknowledgment number: 1
                            (relative ack number)
 Header Length: 32 bytes

☐ Flags: 0x012 (SYN, ACK)

   000. .... = Reserved: Not set
   ...0 .... = Nonce: Not set
   .... O... = Congestion Window Reduced (CWR): Not set
   .... .0.. .... = ECN-Echo: Not set
   .... ..0. .... = Urgent: Not set
   .... ...1 .... = Acknowledgment: Set
   .... .... 0... = Push: Not set
   .... .... .0.. = Reset: Not set
 .... .... ...0 = Fin: Not set
   [TCP Flags: .....A..S.]
 Window size value: 29200
```

如上图,服务器端向客户端发送的报文段序号为0;

服务器发的acknowledgment number字段是根据上一次客户端发给服务器的seq+1得到的;

通过Flags标志位中的SYN位和ACK位都是1来确定该报文段是一个SYN ACK报文段的。

5) 你能从捕获的数据包中分析出tcp三次握手过程吗?

```
TCP 66 24195 → 80 [SYN] Seq=0 win=64240 Len=0 MSS=1460 wS=256 SACK_PERM=1
TCP 66 80 → 24195 [SYN, ACK] Seq=0 Ack=1 win=29200 Len=0 MSS=1460 SACK_PERM=1 wS=128
TCP 54 24195 → 80 [ACK] Seq=1 Ack=1 win=65536 Len=0
```

首先客户端向服务器发送 seq=0 的建立连接的请求

然后服务器向客户端返回 seq=0, ack=0+1=1 的响应

客户端收到响应,返回 seq=1,ack=0+1=1 的确认报文,连接建立

6) 包含HTTP POST命令的TCP报文段的序号是多少?

No.	Time	Source	Destination	Protocol Le	ength	Info		
18	7 7.602480	172.20.129.159	128.119.245.12			_	segment of	a reassembled PDU]
18	8 7.602505	172.20.129.159	128.119.245.12	TCP	1514	[TCP	segment of	a reassembled PDU]
18	9 7.602534	172.20.129.159	128.119.245.12	HTTP	539	POST	/wireshark-	labs/lab3-1-reply.htm HTTP/1.1 (text/plain)
19	0 7.681965	SamsungE_2e:3a:13	Broadcast	ARP	56	Who I	nas 172.20.1	28.1? Tell 172.20.153.128
	1 7.852503	128.119.245.12	172.20.129.159	TCP				Seq=1 Ack=98471 Win=183296 Len=0
19	2 7.852505	128.119.245.12	172.20.129.159	TCP				Seq=1 Ack=99931 win=186112 Len=0
	3 7.852506	128.119.245.12	172.20.129.159	TCP				Seq=1 Ack=101391 Win=189056 Len=0
		128.119.245.12	172.20.129.159	TCP				Seq=1 Ack=102851 Win=192000 Len=0
		128.119.245.12	172.20.129.159	TCP				Seq=1 Ack=104311 Win=194944 Len=0
		128.119.245.12	172.20.129.159	TCP				Seq=1 Ack=105771 Win=197888 Len=0
	7 7.852508	128.119.245.12	172.20.129.159	TCP				Seq=1 Ack=107231 Win=200704 Len=0
		128.119.245.12	172.20.129.159	TCP				Seq=1 Ack=108691 Win=203648 Len=0
19	199 7.852510 128.119.245.12 172.20.129.159 TCP 60 80 - 24195 [ACK] Seq=1 Ack=110151 Win=206592 Len=0							
⊕ Fran	me 189: 539	bytes on wire (4312	bits), 539 bytes c	aptured (4312	bits)	on interfa	ce 0
			oe (a0:a8:cd:16:b7:b				91:14:32 (58	:69:6c:91:14:32)
			172.20.129.159, Dst					
			Port: 24195, Dst P	ort: 80,	seq:	15248	37, Ack: 1,	Len: 485
	ource Port:							
	estination P							
	Stream index							
	TCP Segment							
	Sequence number: 152487 (relative sequence number)							
	[Next sequence number: 152972 (relative sequence number)]							
	Acknowledgment number: 1 (relative ack number)							
	Header Length: 20 bytes ⊕ Flags: 0x018 (PSH, ACK)							
	indow size v							
[c	Calculated w	/indow size: 65536]						

如图,序号为152487(在红色的框内)

7) 如果将包含HTTP POST命令的TCP报文段看作是TCP连接上的第一个报文段,那么该TCP连接上的第六个报文段的序号是多少?是何时发送的?该报文段所对应的ACK是何时接收的?

```
shark-labs/lab3-1-reply.htm HTTP/1.1 (1
[ACK] Seq=1 Ack=98471 win=183296 Len=0
[ACK] Seq=1 Ack=99931 win=186112 Len=0
      192 7.852505
                               128.119.245.12
                                                                  172.20.129.159
                                                                                                                       60 80 → 24195
                                                                                                     TCP
                                                                                                                       00 80 - 24195 [ACK] Seq=1 ACK=101391 Win=180112 Len=0
60 80 - 24195 [ACK] Seq=1 ACK=101391 Win=189056 Len=0
60 80 - 24195 [ACK] Seq=1 ACK=102851 Win=192000 Len=0
60 80 - 24195 [ACK] Seq=1 ACK=104311 Win=194944 Len=0
60 80 - 24195 [ACK] Seq=1 ACK=105771 Win=197888 Len=0
60 80 - 24195 [ACK] Seq=1 ACK=107231 Win=200704 Len=0
                                                                                                     TCP
TCP
TCP
      193 7.852506
                               128, 119, 245, 12
                                                                  172, 20, 129, 159
      194 7.852507
195 7.852507
196 7.852508
197 7.852508
                               128.119.245.12
128.119.245.12
128.119.245.12
                                                                  172.20.129.159
172.20.129.159
172.20.129.159
                               128.119.245.12
                                                                  172.20.129.159
                                                                                                     TCP
 Frame 189: 539 bytes on wire (4312 bits), 539 bytes captured (4312 bits) on interface 0
[Frame: 29, payload: 5030-6489 (1460 bytes)]
      [Frame: 30, pavload: 6490-7949 (1460 bytes)]

[Frame: 31, pavload: 7950-9409 (1460 bytes)]

[Frame: 32, pavload: 9410-10869 (1460 bytes)]

[Frame: 33, pavload: 10870-12329 (1460 bytes)]
```

由上图可知,该HTTP协议被108个TCP报文段所封装,双击下面蓝色的第六个(红框中)。 得到下图:

No.	Time	Source	Destination	Protocol	Length	Info			
	23 6.825202	172.20.129.159	128.119.245.12	TCP	54	24195 → 80 [ACK] Seq=1 Ack=1 Win=65536 Len=0			
	25 6.848378	172.20.129.159	128.119.245.12	TCP	704	[TCP segment of a reassembled PDU]			
	26 6.848949	172.20.129.159	128.119.245.12	TCP		[TCP segment of a reassembled PDU]			
	27 6.848983	172.20.129.159	128.119.245.12	TCP		[TCP segment of a reassembled PDU]			
	28 6.849028	172.20.129.159	128.119.245.12	TCP		[TCP segment of a reassembled PDU]			
	29 6.849074	172.20.129.159	128.119.245.12	TCP		[TCP segment of a reassembled PDU]			
		172.20.129.159	128.119.245.12	TCP		[TCP segment of a reassembled PDU]			
		172.20.129.159	128.119.245.12	TCP		[TCP segment of a reassembled PDU]			
	32 6.849219	172.20.129.159	128.119.245.12	TCP		[TCP segment of a reassembled PDU]			
	33 6.849266	172.20.129.159	128.119.245.12	TCP		[TCP segment of a reassembled PDU]			
	34 6.849310	172.20.129.159	128.119.245.12	TCP		[TCP segment of a reassembled PDU]			
	38 7.093676	128.119.245.12	172.20.129.159	TCP		80 - 24195 [ACK] Seq=1 Ack=651 Win=30592 Len=0			
	39 7.093678	128.119.245.12	172.20.129.159	TCP	60	80 → 24195 [ACK] Seq=1 Ack=2111 Win=33536 Len=0			
+ E	⊕ Frame 30: 1514 bytes on wire (12112 bits), 1514 bytes captured (12112 bits) on interface 0 ⊕ Ethernet II, Src: IntelCor_16:b7:be (a0:a8:cd:16:b7:be), Dst: RuijieNe_91:14:32 (58:69:6c:91:14:32) ⊕ Internet Protocol Version 4, Src: 172.20.129.159, Dst: 128.119.245.12 ⊟ Transmission Control Protocol, Src Port: 24195, Dst Port: 80, Seq: 6491, Ack: 1, Len: 1460 Source Port: 24195								
	Destination F								
	[Stream index	-							
	[TCP Segment								
	Sequence number: 6491 (relative sequence number)								
	[Next sequence number: 7951 (relative sequence number)]								
	Acknowledgmer		elative ack number)						
	Header Length								
+	Flags: 0x010								
	Window size								
	Calculated	[Calculated window size: 65536]							

即这个TCP段就是我们要找的第6个TCP段,序号为6491,发送时间为在http post发送之前,tcp连接建立之后发送的(发送了前五个TCP段后),该报文所对应的ACK则我们需要找到6491+1460=7951,即我们要找到Acknowledgment number=7951的TCP段(由服务器发回的),仔细寻找后,成功找到如下图:

```
52 7.101079 128.119.245.12
53 7.101081 128.119.245.12
                                                172.20.129.159
172.20.129.159
                                                                                       60 80 - 24195 [ACK] Seq=1 Ack=7951 Win=45184 Len=0 60 80 - 24195 [ACK] Seq=1 Ack=9411 Win=48128 Len=0
                                                                          TCP
      54 7.101082 128.119.245.12
55 7 101100 172 20 120 150
                                                172.20.129.159
                                                                         TCP
                                                                                       60 80 → 24195 [ACK] Seq=1 Ack=10871 Win=50944 Len=0
                                                128 119 245 12
                                                                          TCD
                                                                                     1514 TTCD SAG
                                                                                                                a reassembled poul
⊕ Frame 52: 60 bytes on wire (480 bits), 60 bytes captured (480 bits) on interface 0
⊞ Ethernet II, Src: RuijieNe_91:14:32 (58:69:66:91:14:32), Dst: IntelCor_16:b7:be (a0:a8:cd:16:b7:be)

⊞ Internet Protocol Version 4, Src: 128.119.245.12, Dst: 172.20.129.159
☐ Transmission Control Protocol, Src Port: 80, Dst Port: 24195, Seq: 1, Ack: 7951, Len: 0
     Source Port: 80
     Destination Port: 24195
     [Stream index: 0]
     [TCP Segment Len: 0]
      Sequence number: 1
                                 (relative sequence number)
   Acknowledgment number: 7951
Header Length: 20 bytes
                                             (relative ack number)
```

8) 前六个TCP报文段的长度各是多少?

```
☐ [108 Reassembled TCP Segments (152971 bytes): #25(650),
        [Frame: 25, payload: 0-649 (650 bytes)]
        [Frame: 26, payload: 650-2109 (1460 bytes)]
        [Frame: 27, payload: 2110-3569 (1460 bytes)]
        [Frame: 28, payload: 3570-5029 (1460 bytes)]
        [Frame: 29, payload: 5030-6489 (1460 bytes)]
        [Frame: 30, payload: 6490-7949 (1460 bytes)]
        [Frame: 31, payload: 7950-9409 (1460 bytes)]
```

如图,前六个TCP报文段的长度为650bytes、1460bytes、1460bytes、1460bytes、1460bytes、1460bytes、1460bytes。(不包含TCP头部字段的20字节)

9) 在整个跟踪过程中,接收端公示的最小的可用缓存空间是多少?限制发送端的传输以后,接收端的缓存是否仍然不够用?

```
Protocol Length Info
No.
      Time
                 Source
                                     Destination
                 172.20.129.159
    21 6.578235
                                     128.119.245.12
                                                                   66 24195 - 80 [SYN] Seq=0 Win=64240 I
                                                         TCP
    22 6.824816
                 128.119.245.12
                                     172.20.129.159
                                                         TCP
                                                                   66 80 - 24195 [SYN, ACK] Seq=0 Ack=1
                                                         TCP
    23 6.825202
                 172.20.129.159
                                     128.119.245.12
                                                                   54 24195 - 80 [ACK] Seq=1 Ack=1 Win=6
    25 6.848378
                 172.20.129.159
                                     128.119.245.12
                                                         TCP
                                                                  704 [TCP segment of a reassembled PDU]
                                                                 1514 [TCP segment of a reassembled PDU]
    26 6.848949
                 172.20.129.159
                                     128.119.245.12
                                                         TCP
    27 6.848983
                 172.20.129.159
                                     128.119.245.12
                                                         TCP
                                                                 1514 [TCP segment of a reassembled PDU]
    28 6.849028 172.20.129.159
                                     128, 119, 245, 12
                                                         TCP
                                                                 1514 [TCP segment of a reassembled PDU]
                                                         TCP
    29 6.849074
                 172.20.129.159
                                     128.119.245.12
                                                                 1514 [TCP segment of a reassembled PDU]
    30 6.849118
                 172.20.129.159
                                     128.119.245.12
                                                         TCP
                                                                 1514 [TCP segment of a reassembled PDU]
    31 6.849163
                 172.20.129.159
                                     128.119.245.12
                                                         TCP
                                                                 1514 [TCP segment of a reassembled PDU]
    32 6.849219
                 172.20.129.159
                                     128.119.245.12
                                                         TCP
                                                                 1514 [TCP segment of a reassembled PDU]
    33 6.849266 172.20.129.159
                                     128.119.245.12
                                                         TCP
                                                                 1514 [TCP segment of a reassembled PDU]
    34 6.849310 172.20.129.159
                                     128.119.245.12
                                                        TCP
                                                                 1514 [TCP segment of a reassembled PDU]
⊕ Frame 22: 66 bytes on wire (528 bits), 66 bytes captured (528 bits) on interface 0
⊕ Ethernet II, Src: RuijieNe_91:14:32 (58:69:6c:91:14:32), Dst: IntelCor_16:b7:be (a0:a8:cd:16:b7:be)
⊕ Internet Protocol Version 4, Src: 128.119.245.12, Dst: 172.20.129.159
☐ Transmission Control Protocol, Src Port: 80, Dst Port: 24195, Seq: 0, Ack: 1, Len: 0
    Source Port: 80
    Destination Port: 24195
    [Stream index: 0]
    [TCP Segment Len: 0]
    Sequence number: 0
                          (relative sequence number)
    Acknowledgment number: 1
                                (relative ack number)
    Header Length: 32 bytes

⊕ Flags: 0x012 (SYN, ACK)

    Window size value: 29200
```

如图,接收端公示的最小的可用缓存空间是29200,由于该窗口大小会一直增加,所以 不会出现接收端的缓存是否仍然不够用的情况,而且并没有限制发送端的传输。

10) 在跟踪文件中是否有重传的报文段?进行判断的依据是什么?

```
[Frame: 88, pay10a0: 03430-04889 (1400 DyteS)]
[Frame: 89, payload: 64890-66349 (1460 bytes)]
       90, payload: 66350-67809 (1460 bytes)]
        101, payload: 67810-69269 (1460 bytes)]
[Frame: 102, payload: 69270-70729 (1460 bytes)]
[Frame: 103, payload: 70730-72189 (1460 bytes)]
[Frame: 104, payload: 72190-73649 (1460 bytes)]
[Frame: 105, payload: 73650-75109 (1460 bytes)]
[Frame: 106, payload: 75110-76569 (1460 bytes)]
[Frame: 107, payload: 76570-78029 (1460 bytes)]
        108, payload: 78030-79489 (1460 bytes)]
       109, payload: 79490-80949 (1460 bytes)]
[Frame:
[Frame: 110, payload: 80950-82409 (1460 bytes)]
[Frame: 111, payload: 82410-83869 (1460 bytes)]
[Frame: 112, payload: 83870-85329 (1460 bytes)]
[Frame: 113. pavload: 85330-86789 (1460 bytes)]
```

没有出现重传的报文段,判断的依据是客户端没有发送序列号相同的报文段,或者也可以在上面的绿色区域观察,一是通过序列号有没有重复的,二是如果重发会被标记为红色,并有提示文字 "Resend TCP Segment",大概是长这个样子,直接某次抓包遇到过重传的,不过这次我抓取的这些包都没有重复的。

11) TCP连接的throughput (bytes transferred per unit time)是多少?请写出你的计算过程

☐ [108 Reassembled TCP Segments (152971 bytes):

首先总共发出的有效字节数(不包含头部字段)为152971 bytes 发出这些TCP

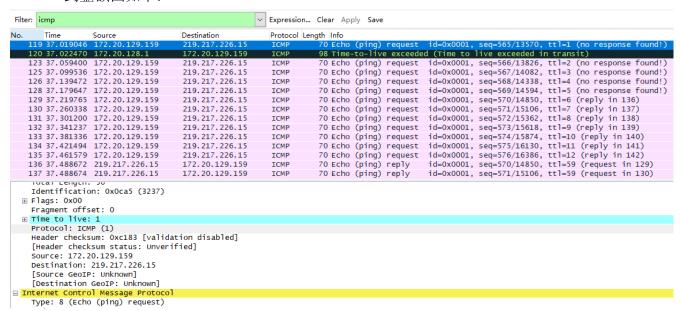
报文段所经历的总时间为(发送第一个TCP报文段的时间到最后一个TCP报文段被ACK确认的时间)7.856015-6.848378=1.007637 s,因此按照下面这个公式进行计算

TCP的最大吞吐率 = 一个RTT传输的有效数据 / RTT

得到 throughput = 152971 b/1.007637 s = 151811.615 bps

3 IP 分析

(1) 在你的捕获窗口中,应该能看到由你的主机发出的一系列ICMP Echo Request 包和中间路由器返回的一系列ICMP TTL-exceeded消息。选择第一个你的主机发出的 ICMP Echo Request消息,在packet details窗口展开数据包的Internet Protocol部分,得到金额图如下:



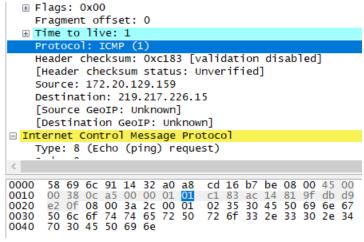
回答问题:

1) 你主机的IP地址是什么?

172. 20. 129. 159

2) 在IP数据包头中,上层协议(upper layer)字段的值是什么?

如下图:上层协议字段的值为01,代表为ICMP报文



3) IP头有多少字节?该IP数据包的净载为多少字节?并解释你是怎样确定该IP数据包的 净载大小的?

```
0100 .... = Version: 4
.... 0101 = Header Length: 20 bytes (5)

→ Differentiated Services Field: 0x00 (DSCP: CS0, ECN: Not-ECT)
Total Length: 56
```

如上图,头部为20字节,IP报文总长度为56字节,所以该IP数据包的净载为56-20=36字节。

4) 该IP数据包分片了吗?解释你是如何确定该P数据包是否进行了分片

```
    Flags: 0x00
    0..... = Reserved bit: Not set
    .0.... = Don't fragment: Not set
    .0.... = More fragments: Not set
    Fragment offset: 0
```

如上图,可以知道该 IP 数据包未进行分片。

首先看到 DF=0,则证明是允许分片的,那到底分片了没有,就接着看 MF 和 Fragment offset, MF 标志位是 0,代表这或者未分片,或者这是众多分片的最后一片,再看段的偏移 Fragment offset,由于其为 0,可以知道这是该 IP 数据包未分片。

(2) 单击 Source 列按钮,这样将对捕获的数据包按源 IP 地址排序。选择第一个你的主机发出的 ICMP Echo Request 消息,在 packet details 窗口展开数据包的 Internet Protocol 部分。得到截图如下:

				_	
Filter:	icmp			Expression	Clear Apply Save
No.	Time	Source	▲ Destination	Protocol Lei	ngth Info
12	0 37.022470	172.20.128.1	172.20.129.159	ICMP	98 Time-to-live exceeded (Time to live exceeded in transit)
15	7 40.095483	172.20.128.1	172.20.129.159	ICMP	120 Destination unreachable (Port unreachable)
		172.20.128.1	172.20.129.159	ICMP	120 Destination unreachable (Port unreachable)
		172.20.128.1	172.20.129.159	ICMP	120 Destination unreachable (Port unreachable)
		172.20.128.1	172.20.129.159	ICMP	98 Time-to-live exceeded (Time to live exceeded in transit)
		172.20.128.1	172.20.129.159	ICMP	98 Time-to-live exceeded (Time to live exceeded in transit)
		172.20.128.1	172.20.129.159	ICMP	590 Time-to-live exceeded (Time to live exceeded in transit)
		172.20.128.1	172.20.129.159	ICMP	590 Time-to-live exceeded (Time to live exceeded in transit)
		172.20.129.159	219.217.226.15	ICMP	70 Echo (ping) request id=0x0001, seq=565/13570, ttl=1 (no
		172.20.129.159	219.217.226.15	ICMP	70 Echo (ping) request id=0x0001, seq=566/13826, ttl=2 (no
		172.20.129.159	219.217.226.15	ICMP	70 Echo (ping) request id=0x0001, seq=567/14082, ttl=3 (no
		172.20.129.159	219.217.226.15	ICMP	70 Echo (ping) request id=0x0001, seq=568/14338, ttl=4 (no
		172.20.129.159	219.217.226.15	ICMP	70 Echo (ping) request id=0x0001, seq=569/14594, ttl=5 (no
		172.20.129.159	219.217.226.15	ICMP	70 Echo (ping) request id=0x0001, seq=570/14850, ttl=6 (rep
		172.20.129.159	219.217.226.15	ICMP	70 Echo (ping) request id=0x0001, seq=571/15106, ttl=7 (rep
		172.20.129.159	219.217.226.15	ICMP	70 Echo (ping) request id=0x0001, seq=572/15362, ttl=8 (rep
		172.20.129.159	219.217.226.15	ICMP	70 Echo (ping) request id=0x0001, seq=573/15618, ttl=9 (rep
		172.20.129.159	219.217.226.15	ICMP	70 Echo (ping) request id=0x0001, seq=574/15874, ttl=10 (re
		172.20.129.159	219.217.226.15	ICMP	70 Echo (ping) request id=0x0001, seq=575/16130, ttl=11 (re
13	5 37.461579	172.20.129.159	219.217.226.15	ICMP	70 Echo (ping) request id=0x0001, seq=576/16386, ttl=12 (re
					bits) on interface 0
⊕ Eth	ernet II, Sr	c: IntelCor_16:b7	':be (a0:a8:cd:16:b7	:be), Dst: R	uijieNe_91:14:32 (58:69:6c:91:14:32)
			: 172.20.129.159, D	st: 219.217.	226.15
	100 = V				
		Header Length: 20			
			0x00 (DSCP: CS0, E	CN: Not-ECT)	
T	otal Length:	56			
I	dentificatio	n: 0x0ca5 (3237)			
⊟ F	lags: 0x00				
		Reserved bit: No			
		- Don't framment.			
0000	58 69 6c 91	. 14 32 a0 a8 cd	16 b7 be 08 00 45 0		E.
0010			83 ac 14 81 9f db d 35 30 45 50 69 6e 6		FORDing
0020			6f 33 2e 33 30 2e 3		.50EPing ro3.30.4
	70 30 45 50		0. 33 20 33 30 20 3	p0EPin	10313014
_				•	

回答问题:

1) 你主机发出的一系列ICMP消息中IP数据报中哪些字段总是发生改变?

我通过查看多个 ICMP 消息,发现 ID(标识符字段)、TTL、Header checksum 这 三个字段总在变化。

2) 哪些字段必须保持常量?哪些字段必须改变?为什么?

必须改变:

ID 鉴别码,用于区分不同的数据包;

TTL 来自于 traceroute 的要求,用来测试路径上的路由信息;

Header Checksum 首部校验和,前面的字段改变,该值也必须跟着改变;必须保持常量:

除以上(ID, TTL, Header Checksum)外的字段保持常量。

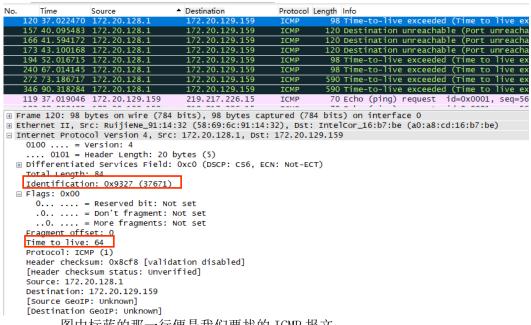
3) 描述你看到的IP数据包Identification字段值的形式。

标识符是 16 位的, 在某一范围内是+1 递增的, 如下两幅图,

	CHINE I WHI	17) >0	1•		
126 37.139472 172.20.129.159	219.217.226.15	ICMP			
128 37.179647 172.20.129.159	219.217.226.15	ICMP			
129 37.219765 172.20.129.159	125 37.099536 1	72. 20. 129. 159	219, 217, 226, 15	ICMP	70 Echo
130 37.260338 172.20.129.159	126 37.139472 1		219, 217, 226, 15	ICMP	70 Echo
131 37.301200 172.20.129.159	128 37.179647 1	.72.20.129.159	219.217.226.15	ICMP	70 Echo
132 37.341237 172.20.129.159	129 37.219765 1	.72.20.129.159	219.217.226.15	ICMP	70 Echo
133 37.381336 172.20.129.159	130 37.260338 1	72.20.129.159	219.217.226.15	ICMP	70 Echo
134 37.421494 172.20.129.159	131 37.301200 1	.72.20.129.159	219.217.226.15	ICMP	70 Echo
135 37.461579 172.20.129.159	132 37.341237 1		219.217.226.15	ICMP	70 Echo
⊕ Frame 126: 70 bytes on wire (560 l	133 37.381336 1		219.217.226.15	ICMP	70 Echo
	134 3/.421494 1	.72.20.129.159	219.217.226.15	ICMP	70 Echo
⊕ Ethernet II, Src: IntelCor_16:b7:l	135 37.461579 1	.72.20.129.159	219.217.226.15	ICMP	70 Echo
□ Internet Protocol Version 4, Src:	⊕ Frame 128: 70 byt	tes on wire (560 b	oits). 70 bytes cap	tured (560	bits) on
0100 = Version: 4	Ethornot II Src	· IntelCor 16:h7:l	on (a0.a8.cd.16.b7.	ha) Det : I	DudiioNo O
0101 = Header Length: 20 by	□ Internet Protoco	l Version 4. Src:	172.20.129.159. DS	t: 219.217.	226.15
⊕ Differentiated Services Field: (0100 = Ver	rsion: 4			
Total Length: 56	0101 = Hea	ader Length: 20 by	/tes (5)		
Identification: 0x0ca8 (3240)			0x00 (DSCP: CSO, EC	N: Not-ECT))
	Total Length:	56	•	Ĩ	
	Identification	: 0x0ca9 (3241)			
	⊟ Flags: 0x00				

通过这两幅图可以看出,相邻的两个 IP 数据包第一个的 ID 为 0x0ca8,而第二个 IP 数据包的 ID 为 0xca9。

(3) 找到由最近的路由器(第一跳)返回给你主机的 ICMP Time-to-live exceeded 消息。得到的截图如下



图中标蓝的那一行便是我们要找的 ICMP 报文

回答问题:

1) Identification字段和TTL字段的值是什么?

如图ID为0x9327, TTL为64(见图中红框所圈中的)

2) 最近的路由器(第一跳)返回给你主机的ICMP Time-to-live exceeded消息中这些值是 否保持不变?为什么?

都保持不变, IP 是无连接服务,相同的标识是为了分段后组装成同一段,给同一个主机返回的 ICMP,标识不代表序号,因此 Identification 不变,TTL 消息是相同的,都是 64。

(4) 单击Time列按钮,这样将对捕获的数据包按时间排序。找到在将包大小改为2000字节后你的主机发送的第一个ICMP Echo Request消息。 得到下图

Filter:				Expression.	Clear Apply Save			
No.	Time	Source	Destination	Protocol L	ength Info			
27	70 73.179466	172.20.129.159	219.217.226.15	IPV4	1514 Fragmented IP protoco			
27	71 73.179484	172.20.129.159	219.217.226.15	ICMP	534 Echo (ping) request			
27	72 73.186717	172.20.128.1	172.20.129.159	ICMP	590 Time-to-live exceeded			
27	73 73.219980	172.20.129.159	219.217.226.15	IPV4	1514 Fragmented IP protoco			
27	74 73.219998	172.20.129.159	219.217.226.15	ICMP	534 Echo (ping) request			
27	75 73.223634	192.168.80.1	172.20.129.159	ICMP	70 Time-to-live exceeded			
27	76 73.260126	172.20.129.159	219.217.226.15	IPV4	1514 Fragmented IP protoco			
27	77 73.260161	172.20.129.159	219.217.226.15	ICMP	534 Echo (ping) request			
27	78 73.262922	202.118.168.86	172.20.129.159	ICMP	70 Time-to-live exceeded			
27	79 73.301173	172.20.129.159	219.217.226.15	IPV4	1514 Fragmented IP protoco			
28	30 73.301214	172.20.129.159	219.217.226.15	ICMP	534 Echo (ping) request			
20	14 77 700747	173 30 150 300	173 30 150 355	HENC	AS Name allegations desired to			
					(4272 bits) on interface 0			
			•		RuijieNe_91:14:32 (58:69:60			
			: 172.20.129.159, D	st: 219.217	.226.15			
	100 = \							
		Header Length: 20						
			0x00 (DSCP: CS0, E	CN: Not-ECT	r)			
	otal Length							
		on: 0x0cbd (3261)						
	lags: 0x00							
F	ragment offs	set: 1480						
± T	ime to live	: 1						
P	rotocol: ICM	MP (1)						
Н	Header checksum: Oxbee2 [validation disabled]							
[Header check	ksum status: Unver	ified]					
S	ource: 172.2	20.129.159						
D	estination:	219.217.226.15						
_ [Source Geoi	P: Unknown]						
[Destination	GeoIP: Unknown]						
⊕ [2 IPv4 Fragr	ments (1980 bytes)	: #270(1480), #271(500)]				

回答问题:

1) 该消息是否被分解成不止一个IP数据报?

是的,该消息被分解成了2片

2) 观察第一个IP分片,IP头部的哪些信息表明数据包被进行了分片?IP头部的哪些信息表明数据包是第一个而不是最后一个分片?该分片的长度是多少

观察下面左面的图,先看第一个分片,MF为1代表后面还有更多分片,同时Fragment

```
offset=0 说明这是第一个分片
□ Flags: 0x01 (More Fragments)
0... = Reserved bit: Not set
.0. ... = Don't fragment: Not set
.1. ... = More fragments: Set
Fragment offset: 0
□ Flags: 0x00
0... = Reserved bit: Not set
.0. ... = Don't fragment: Not set
.0. ... = More fragments: Not set
Fragment offset: 1480
```

再看第二个分片的 ICMP 报文,如上面右面的图,可以看到此时 MF=0,则证明这是最后一个分片了,即共有两个 IP 分片

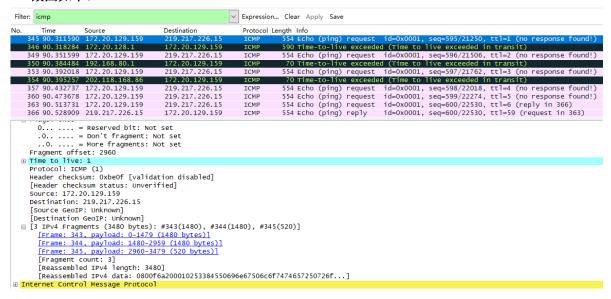
```
[2 IPv4 Fragments (1980 bytes): #270(1480), #271(500)]
```

[Frame: 270, payload: 0-1479 (1480 bytes)]
[Frame: 271, payload: 1480-1979 (500 bytes)]

[Fragment count: 2]

通过上面的图就可以知道每个分片的长度分别为 1480 bytes 和 500 bytes (当然通过 之前的每一个分片的偏移量也可以得到该结果)

C 找到在将包大小改为3500字节后你的主机发送的第一个ICMP Echo Request消息。 得到 截图如下:



如上,这里面直接显示了包大小为3500字节拆分后的最后一个分片,前两个分片见下图:

No.	Time	Source	Destination	Protocol	l Length Info					
		172.20.129.159	219.217.226.15	IPV4	1514 Fragmented IP protocol (proto=ICMP					
		172.20.129.159	219.217.226.15	IPV4	1514 Fragmented IP protocol (proto=ICMP					
		172.20.129.159	219.217.226.15	ICMP	554 Echo (ping) request id=0x0001, se					
		172.20.128.1	172.20.129.159	ICMP	590 Time-to-live exceeded (Time to liv					
_		172.20.129.159	219.217.226.15	IPV4	1514 Fragmented IP protocol (proto=ICMP					
		172.20.129.159	219.217.226.15	IPV4	1514 Fragmented IP protocol (proto=ICMP					
		172.20.129.159	219.217.226.15	ICMP	554 Echo (ping) request id=0x0001, se					
		192.168.80.1	172.20.129.159	ICMP	70 Time-to-live exceeded (Time to liv					
		172.20.129.159 172.20.129.159	219.217.226.15 219.217.226.15	IPV4	1514 Fragmented IP protocol (proto=ICMP					
	32 90.391991	1/2.20.129.139	219.217.226.13	IPV4	1514 Fragmented IP protocol (proto=ICMP	Ι,				
⊕ Fr	ame 343: 1514	4 bytes on wire (1	12112 bits), 1514 byt	tes captur	ured (12112 bits) on interface O					
# Et	hernet II, Sr	rc: IntelCor_16:b7	:be (a0:a8:cd:16:b7:	be), Dst:	:: RuijieNe_91:14:32 (58:69:6c:91:14:32)					
□ In	ternet Proto	col Version 4, Sro	: 172.20.129.159, Ds	st: 219.21	217.226.15					
	0100 = \									
		Header Length: 20								
			0x00 (DSCP: CS0, E0	CN: Not-EC	ECT)					
	Total Length									
		on: 0x0cc3 (3267)								
		(More Fragments)								
		= Reserved bit: No								
		= Don't fragment:								
		= More fragments:	set							
	Fragment offs									
_	Time to live: 1									
	Protocol: ICMP (1) Header checksum: 0x9bc1 [validation disabled]									
		ksum status: Unver								
	Source: 172.2		ITTedj							
		219.217.226.15								
	[Source GeoI									
1	Lacui de Geori	r. onknownj								

如上图其中, No。一列下343,344,345组成了这三个分片

回答问题:

1) 原始数据包被分成了多少片?

[3 IPv4 Fragments (3480 bytes): #343(1480), #344(1480), #345(520)]

[Frame: 343, payload: 0-1479 (1480 bytes)]
[Frame: 344, payload: 1480-2959 (1480 bytes)]
[Frame: 345, payload: 2960-3479 (520 bytes)]
[Fragment count: 3]

如图可知被分成了三片

2) 这些分片中IP数据报头部哪些字段发生了变化?

总长度字段、MF 标志位、偏移量(Fragment offset)、校验和 这四个字段发生了变化。 第一片与第二片的总长度字段均为 1500 bytes(因为包含 20 字节的头部),第三片 的总长度字段为 540 bytes;

第一片与第二片的 MF 标志位均为 1, 第三片的 MF 标志位为 0

第一片的偏移量为0,第二片的偏移量为1480,第三片的偏移量为2960;

由于每一个分片的报文不是完全相同,则每一个分片的校验和字段一定也不相同。

4 抓取 ARP 数据包

(1) 利用MS-DOS命令: arp或c:\windows\system32\arp查看主机上ARP缓存的内容。说明ARP缓存中每一列的含义是什么?

输入apr -a查看主机上ARP缓存的内容,结果如下图所示 ARP 缓存中的每一列分别表示 IP 地址所对应的物理地址和类型(动态配置或静态配置)

■ 管理员: Command Prompt

```
C:\WINDOWS\system32>arp -a
接□: 172.20.129.159 --- 0x3
  Internet 地址
                        58-69-6c-91-14-32
 172.20.128.1
 172.20.159.255
                        ff-ff-ff-ff-ff-ff
  224.0.0.22
                        01-00-5e-00-00-16
  224.0.0.251
                        01-00-5e-00-00-fb
 224.0.0.252
                        01-00-5e-00-00-fc
 239.255.255.250
                        01-00-5e-7f-ff-fa
 255.255.255.255
                        ff-ff-ff-ff-ff-ff
```

(2) 清除主机上ARP缓存的内容,抓取ping命令时的数据包。分析数据包,回答下面的问题: 通过 arp -d 命令来清楚本机的ARP的缓存内容,如下图

```
C:\WINDOWS\system32>arp -d

C:\WINDOWS\system32>arp -a

接口: 172.20.129.159 --- 0x3

Internet 地址 物理地址 类型

172.20.128.1 58-69-6c-91-14-32 动态
224.0.0.22 01-00-5e-00-00-16 静态
```

1) **ARP数据包的格式是怎样的?由几部分构成,各个部分所占的字节数是多少?** ARP数据包的格式如下图:



由 9 部分构成,分别是硬件类型(2 字节),协议类型(2 字节),硬件地址长度(1 字节),协议地址长度(1 字节),0P(2 字节),发送端 MAC 地址(6 字节),发送端 IP 地址(4 字节),目的 MAC 地址(6 字节),目的 IP 地址(4 字节)。

通过 Ping 命令首先给 219. 217. 226. 15 这个 IP 地址发送测试包,如下图

```
C:\WINDOWS\system32\ping 219.217.226.15

正在 Ping 219.217.226.15 具有 32 字节的数据:
来自 219.217.226.15 的回复: 字节=32 时间=65ms TTL=59
来自 219.217.226.15 的回复: 字节=32 时间=98ms TTL=59
来自 219.217.226.15 的回复: 字节=32 时间=98ms TTL=59
来自 219.217.226.15 的回复: 字节=32 时间=80ms TTL=59
来自 219.217.226.15 的回复: 字节=32 时间=80ms TTL=59

219.217.226.15 的 Ping 统计信息:
数据包: 已发送 = 4,已接收 = 4,丢失 = 0(8½ 丢失),
往返行程的估计时间(以毫秒为单位):
最短 = 65ms,最长 = 98ms,平均 = 80ms

C:\WINDOWS\system32\ping 219.217.226.15

正在 Ping 219.217.226.15 自回复: 字节=32 时间=46ms TTL=59
来自 219.217.226.15 的回复: 字节=32 时间=75ms TTL=59
来自 219.217.226.15 的回复: 字节=32 时间=75ms TTL=59
来自 219.217.226.15 的回复: 字节=32 时间=16ms TTL=59
请求超时。

219.217.226.15 的 Ping 统计信息:
数据包: 已发送 = 4,已接收 = 3,丢失 = 1(25½ 丢失),
往返行程的估计时间(以毫秒为单位):
最短 = 16ms,最长 = 75ms,平均 = 45ms
```

然后打开 Wireshark 开始俘获包,得到如下截图:

```
Filter: arp
                                                                   Expression... Clear Apply Save
        Time
                                              Destination
                                                                       Protocol Length Info
      35 8.023373
                      IntelCor_16:b7:be
                                              RuijieNe_91:14:32 ARP
                                                                                     42 Who has 172.20.128.1? Tell 172.20.129.159
                                                                                     42 Who has 172.20.128.1 rell 172.20.129.159
     36 8.056062
                      RuijieNe_91:14:32
                                               IntelCor_16:b7:be
                                                                       ARP
     85 57.023704
                      IntelCor 16:b7:be
                                               RuijieNe_91:14:32
                                                                       ARP
     91 57.085059 RuijieNe_91:14:32
                                                                                     60 172.20.128.1 is at 58:69:6c:91:14:32
                                               IntelCor_16:b7:be
                                                                       ARP
    102 67.023799 IntelCor_16:b7:be
103 67.036404 RuijieNe_91:14:32
                                                                                     42 Who has 172.20.128.1? Tell 172.20.129.159 60 172.20.128.1 is at 58:69:6c:91:14:32
                                               RuijieNe_91:14:32
                                                                       ARP
                                               IntelCor_16:b7:be
                                                                       ARP
     335 216.523418 Intelcor_16:b7:be
                                               RuijieNe_91:14:32
                                                                                     42 Who has 172.20.128.1? Tell 172.20.129.159
                                                                        ARP
                                                                                     60 172.20.128.1 is at 58.69.6c;91:14:32
42 who has 172.20.128.1? Tell 172.20.129.159
42 who has 172.20.128.1? Tell 172.20.129.159
    336 216.730480 RuijieNe_91:14:32
378 270.524103 IntelCor_16:b7:be
                                               RuijieNe_91:14:32
    379 271.523816 IntelCor_16:b7:be
380 271.592338 RuijieNe_91:14:32
                                               RuijieNe_91:14:32 ARP
                                              IntelCor_16:b7:be ARP
                                                                                     60 172.20.128.1 is at 58:69:6c:91:14:32
    423 314.023399 Intelcor_16:b7:be
424 314.038897 RuijieNe_91:14:32
                                               RuijieNe_91:14:32 ARP
                                                                                     42 Who has 172.20.128.1? Tell 172.20.129.159
                                                                                     60 172.20.128.1 is at 58:69:6c:91:14:32
                                               IntelCor 16:b7:be
                                                                       ARP
    482 372.024089 IntelCor_16:b7:be
                                                                                     42 Who has 172.20.128.1? Tell 172.20.129.159
                                               RuijieNe_91:14:32
                                                                      ARP
    483 372 196716 RuiiiieNe 91:14:32
                                               IntelCor 16:h7:he
                                                                       ARP
                                                                                     60 172 20 128 1 is at 58.69.6c.91.14.32
⊕ Frame 336: 60 bytes on wire (480 bits), 60 bytes captured (480 bits) on interface 0
⊕ Ethernet II, Src: RuijieNe_91:14:32 (58:69:6c:91:14:32), Dst: IntelCor_16:b7:be (a0:a8:cd:16:b7:be)

□ Address Resolution Protocol (reply)
     Hardware type: Ethernet (1)
Protocol type: IPv4 (0x0800)
     Hardware size: 6
     Protocol size: 4
     Opcode: reply (2)
     Sender MAC address: RuijieNe_91:14:32 (58:69:6c:91:14:32)
Sender IP address: 172.20.128.1
Target MAC address: IntelCor_16:b7:be (a0:a8:cd:16:b7:be)
     Target IP address: 172.20.129.159
```

2) 如何判断一个ARP数据是请求包还是应答包?

通过0P字段。当0P字段值为0x0001时是请求包,当0P字段值为0x0002时是应答包。如下图这是一个请求包:

```
    Frame 379: 42 bytes on wire (336 bits), 42 bytes captured (336
    Ethernet II, Src: IntelCor_16:b7:be (a0:a8:cd:16:b7:be), Dst:
    Address Resolution Protocol (request)
    Hardware type: Ethernet (1)
    Protocol type: IPv4 (0x0800)
    Hardware size: 6
    Protocol size: 4
    Opcode: request (1)
    Sender MAC address: IntelCor_16:b7:be (a0:a8:cd:16:b7:be)
    Sender IP address: RuijieNe_91:14:32 (58:69:6c:91:14:32)
    Target MAC address: 172.20.128.1
```

而下面这幅图就是一个响应包:

```
    Frame 91: 60 bytes on wire (480 bits), 60 bytes captured (480
    Ethernet II, Src: RuijieNe_91:14:32 (58:69:6c:91:14:32), Dst:
    Address Resolution Protocol (reply)
    Hardware type: Ethernet (1)
    Protocol type: IPv4 (0x0800)
    Hardware size: 6
    Protocol size: 4
    Opcode: reply (2)
    Sender MAC address: RuijieNe_91:14:32 (58:69:6c:91:14:32)
    Sender IP address: 172.20.128.1
    Target MAC address: IntelCor_16:b7:be (a0:a8:cd:16:b7:be)
    Target IP address: 172.20.129.159
```

3) 为什么ARP查询要在广播帧中传送,而ARP响应要在一个有着明确目的局域网地址的 帧中传送?

因为进行 ARP 查询时并不知道目的 IP 地址对应的 MAC 地址,所以需要广播查询;而 ARP 响应报文知道查询主机的 MAC 地址(通过查询主机发出的查询报文获得),且局域 网中的其他主机不需要此次查询的结果,因此 ARP 响应要在一个有着明确目的局域网地址的帧中传送。

5 抓取 UDP 数据包

利用 OO 发送消息, 抓包截图如下:

				_				
Filter:	udp		•	Expression	Clear Apply Sav	re		
No.	Time	Source	Destination	Protocol Le	ngth Info			
	3 1.879689	172.20.128.1	255.255.255.255	DHCP	370 DHCP ACK	- Transa		
	11 5.467809	123.151.13.169	172.20.129.159	OICQ	121 OICQ Proto	col		
	22 10.336361	172.20.129.159	123.151.13.169	UDP	193 4003 - 800	0 Len=151		
	23 10.463452	172.20.129.159	202.118.224.101	DNS	78 Standard q	uery 0xd150		
	27 10.527052	123.151.13.169	172.20.129.159	UDP	73 8000 → 400	3 Len=31		
	28 10.550387	172.20.129.159	202.118.224.100	DNS	78 Standard q	uery 0xd150		
	29 10.559664	172.20.129.159	202.118.224.101	DNS	75 Standard q	uery 0x7643		
⊕ Fra	ame 22: 193	bytes on wire (1544	bits), 193 bytes o	captured (1	544 bits) on in	terface 0		
# Et	hernet II, Si	rc: IntelCor_16:b7:	be (a0:a8:cd:16:b7:	:be), Dst: F	RuijieNe_91:14:	32 (58:69:6c		
			172.20.129.159, Ds		13.169			
□ Use	□ User Datagram Protocol, Src Port: 4003, Dst Port: 8000							
- :	Source Port:	4003						
1	Destination I	Port: 8000						
1	Length: 159							
(Checksum: 0x0bd0 [unverified]							
	[Checksum Sta	atus: Unverified]						
	[Stream index	x: 1]						
⊕ Dat	ta (151 byte:	5)						

回答问题:

1) 消息是基于UDP的还是TCP的?

消息是基于UDP的,见上面截图中的红框

2) 你的主机ip地址是什么?目的主机ip地址是什么?

打开IP报文段,可以得到下图: Source: 172.20.129.159 Destination: 123.151.13.169

我主机的IP地址: 172.20.129.159目的IP地址为: 123.151.13.169

3) 你的主机发送QQ消息的端口号和QQ服务器的端口号分别是多少?

Source Port: 4003 Destination Port: 8000

如图,发送QQ消息的端口号为: 4003

QQ服务器的端口号为: 8000

4) 数据报的格式是什么样的?都包含哪些字段,分别占多少字节?

UDP 数据报格式如下图:

来源端口	目的端口	长度域	校验和
------	------	-----	-----

UDP数据报格式有首部和数据两个部分。首部很简单,共8字节。包括:

源端口号: 2字节

目的端口号: 2字节

长度: 2字节, UDP用户数据报的总长度, 以字节为单位。

校验和: 2字节,用于校验UDP数据报的数字段和包含UDP数据报首部的"伪首部"。其校验方法同 IP分组首部中的首部校验和。

抓包得到的UDP报文段如下:

⊡ User Datagram Protocol, Src Port: 4003, Dst Port: 8000

Source Port: 4003 Destination Port: 8000

Length: 159

Checksum: 0x0bd0 [unverified] [Checksum Status: Unverified]

[Stream index: 1]

5) 为什么你发送一个ICQ数据包后,服务器又返回给你的主机一个ICQ数据包?这UDP的不可靠数据传输有什么联系?对比前面的TCP协议分析,你能看出UDP是无连接的吗?

因为服务器需返回接收的结果给客户端。

因为服务器只提供了一次返回的 ACK, 所以不保证数据一定送达。

可以看出。UDP 数据包没有序列号,因此不能像 TCP 协议那样先握手再发送数据,因为每次只发送一个数据报,然后等待服务器响应。

6 DNS 协议分析

打开浏览器,输入<u>http://www.hit.edu.cn/</u>观察 Wireshark 中抓到的包如下:

Filter:	dns		V	Expressio	n Clear Apply Save
No.	Time	Source	Destination	Protocol	Length Info
	1 0.000000	172.20.129.159	202.118.224.101	DNS	74 Standard query 0x61a9 AAAA www.hit.edu.cn
	2 0.006937	202.118.224.101	172.20.129.159	DNS	126 Standard query response Ox61a9 AAAA www.hit.edu.cn SOA hit01.hit.edu.cn
2	36 3.505297	202.118.224.101	172.20.129.159	DNS	73 Standard query response Oxe3af Server failure AAAA www.hitce.net
2	37 3.513504	202.118.224.100	172.20.129.159	DNS	73 Standard query response Oxe3af Server failure AAAA www.hitce.net
2	7.306036	172.20.129.159	202.118.224.101	DNS	75 Standard query Ox6ace A cdn.onenote.net
2	52 7.307625	172.20.129.159	202.118.224.101	DNS	75 Standard query Oxefb2 AAAA cdn.onenote.net
2	3 7.307824	172.20.129.159	202.118.224.101	DNS	88 Standard query 0x5480 A cdn.content.prod.cms.msn.com
2	54 7.307893	172.20.129.159	202.118.224.101	DNS	88 Standard query 0x551f AAAA cdn.content.prod.cms.msn.com
2	55 7.332249	172.20.129.159	202.118.224.100	DNS	75 Standard query Oxefb2 AAAA cdn.onenote.net
2	6 7.332249	172.20.129.159	202.118.224.100	DNS	88 Standard query 0x551f AAAA cdn.content.prod.cms.msn.com

选取一个 dns 的 query 包,查看其 IP 源地址和目的地址如下:

Source: 172.20.129.159

Destination: 202.118.224.101

即我主机的地址为: 172.20.129.159

本地域名服务器的 IP 地址为: 202.118.224.101

再看一下 UDP 对应的端口号:

□ User Datagram Protocol, Src Port: 57590, Dst Port: 53

Source Port: 57590 Destination Port: 53

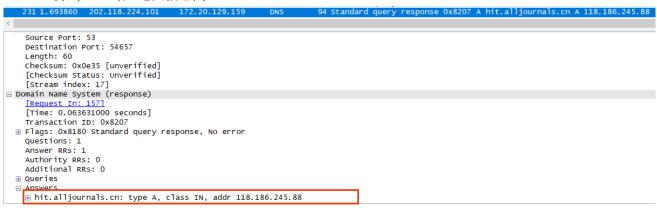
源端口号为: 57590 目的端口号为: 53

选取其中一个子查询报文

78 Standard query Oxclb3 AAAA hit.alljournals.cn 158 1.630749 172.20.129.159 Domain Name System (query) [Response In: 162] Transaction ID: 0xc1b3 Questions: 1 Answer RRs: 0 Authority RRs: 0 Additional RRs: 0 ■ Queries ⊟ hit.alljournals.cn: type AAAA, class IN Name: hit.alljournals.cn [Name Length: 18] [Label Count: 3] Type: AAAA (IPv6 Address) (28) class: IN (0x0001)

可以看到,他要查询主机域名为 hit. all journals. cn 的主机 IP 地址

得到 DNS 对应返回信息为



可以看到, 主机域名为 hit. all journals. cn 对应的主机 IP 地址为 118. 186. 245. 88

四、实验心得

通过本次实验,我掌握了使用 Wireshark 工具进行抓包的方法,并且我能够熟练地使用它。刚开始有些生疏,做前面实验速度非常的慢,但做的多了以后也就熟悉了,后面的实验做起来也相对快了许多。在对 IP 进行分析时,需要下载 pingplotter,刚开始在网上下载的版本可能太新了,根本用不了,后来换了低版本才可以使用。与此同时,我还对于之前学习的HTTP 协议、TCP 协议、UDP 协议、IP 协议、ARP 协议以及 DNS 域名搜索的过程进行了全面的复习,因为若想对 Wireshark 俘获的包进行正确的分析,首先必须要知道每个包的详细内容,例如头部字段都含有哪些部分,每个部分都代表着什么等等。总而言之,经过大量的时间打磨后,我收获颇丰。