

# Asymptotics

Spencer Hirsch, Tyler Gutowski, Remington Greko

February 2, 2023

You have been randomly assigned to teams. Work together to write a report crossing this first bridge on algorithmic quest.

Submit the team's report on Canvas. Include a task matrix indicating who did what.

## Asymptotic Quest

After successful completion of these exercises you will understand the topic of *Asymptotics* and be able to explain and correctly answer questions about the topic.

### *The Pieces and their relationships*

The pieces are functions which we will call  $f$ ,  $g$ , and  $h$ , should we need others they can be named.

Standard relations include:

less than, equal, greater than, etc.

Relations can have properties such as:

Reflexing, Symmetric, Transitive

Quantifiers are also needed

For all, There exists...

Write precise (mathematical) definitions of the following relations:

1. Big-O:

Big-O is used as a general equation in order to find the time complexity of a function. It is hardly ever exact, however, it is unnecessary for it to

be exact as a general estimate is sufficient in determining the complexity of an algorithm based on  $n$ , the number of input.

$T(n)$  and  $f(n)$  are two positive functions. We can write  $T(n) \in O(f(n))$ , and say that  $T(n)$  has order of  $f(n)$ , if there are positive constants  $M$  and  $n_0$  such that  $T(n) \leq M * f(n)$  for all  $n \geq n_0$

(<https://yourbasic.org/algorithms/big-o-notation-explained/>)

## 2. Big- $\Omega$ :

Big  $\Omega$  is used to give a lower bound for the growth of a function. It is very similar to traditional Big-O, however the inequality is different.

$T(n)$  and  $f(n)$  are two positive functions. We write  $T(n) \in \Omega(f(n))$ , and say that  $T(n)$  is Big- $\Omega$  of  $f(n)$ , if there are positive constants  $m$  and  $n_0$  such that  $T(n) \geq m(f(n))$  for all  $n \geq n_0$

## 3. Big- $\Theta$ :

Now that general complexity and lower bound complexity have been defined, we can now look at Big- $\Theta$  which is used to determine both the upper and the lower bound of the time complexity function.

$T(n) \in \Theta(f(n))$  if  $T(n)$  is both  $O(f(n))$  and  $\Omega(f(n))$

Give examples of functions that satisfy these relations.

Explain how these relations describe bounds on running time (or other resources) expended when an algorithm is executed on input of size  $n$ .

<b>Name</b>	<b>Section</b>
Remington Greko	
Tyler Gutowski	
Spencer Hirsch	Mathematical definitions and explanations