CS-430: Problem Definition for the Reactive Agent Assignment

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1 Problem definition

We are trying to implement an agent that functions according to the following basic procedure.

In this case optimality is defined as *minimum cost function*. And the cost function is computed for any road taken as such :

$$c(\text{road}) = l_{road} \cdot c_{kil}$$

Where the value l_{road} is the length of a given road and the value c_{kil} is the cost per kilometer for a given agent.

Computing an optimal plan can be though of as finding a plan p_{opt} such that:

$$p_{opt} = argmin_p\{cost(p) \mid p \text{ reaches a goal}\}$$

Where the cost of a plan is simply given by :

$$cost(p) = \sum_{r \in \text{ roads in the plan}} cost(r)$$

The *goal* of our agent is to deliver all tasks on the map, all of these tasks as well as the length of every road and the cost per kilometer are known prior to planning.

Our agent has the ability to perform multiple tasks at the same time but has a maximum weight that it can carry, it has to verify:

$$\sum$$
 carried tasks \leq agent capacity

We will use BFS and ASTAR to find an optimal solution.

2 States, Transitions and Goals

Since we are working in a deliberative agent paradigm, we will not generate every possible state prior to planning, instead we will define states as a set of variables that are defined by our algorithm as it searches the solution space.

2.1 States

The states the agent can be in are defined by the following variables, for a given agent agent:

- 1. **agent.position**: the city a given agent is in
- 2. agent.ctasks: the list of all tasks carried by the agent in a given state
- 3. **agent.cweight**: the total weight carried by the agent in a given state
- 4. **agent.free_tasks**: the complete list of all undelivered, not yet picked-up tasks at a given state

2.2 Transitions

Our agent can perform three basic actions

- 1. It can drive from one city to another
- 2. It can decide to pickup or not to pickup a task when it is in a given city
- 3. It can decide to deliver a task when it is in the city it should be delivered to

We therefore construct our state representation around those three fundamental transitions:

1. drive(start, destination): the movement action between two adjacent cities.

Cost function defined as: $cost_{drive}(road_{start,destination}) = l_{road} \cdot c_{kil}$

 $Possible \iff$ the agent is in the city start and the city destination is adjacent to start.

Consequences: agent is now in the city destination

2. pickup(task): the action of adding a task to be carried by the agent

Cost function defined as: $cost_{pickup} = 0$

 $Possible \iff \text{the agent is in the city where } task \text{ is and if}$

 $agent.cweight + w_{task} \leq capacity$

Consequences: agent appends task to agent.ctasks and adds the weight w_{task} of the new task to agent.cweight, task is removed from $agent.free_tasks$

3. **deliver(task)**: the action of delivering a task carried by the agent to it's destination

```
Cost function defined as: cost_{deliver} = 0
```

 $Possible \iff \text{the agent is in the city where } task \text{ should be delivered}$

Consequences: agent removes task from agent.ctasks and substracts the weight w_{task} from the task to agent.cweight

2.3 Goal

The goal state is defined as any state where:

```
agent.free tasks = \emptyset
```

3 Implementation of the re-planning algorithm

Problem: if every agent has it's own *agent.free_tasks* data-structure, it will not be able to cope with the fact that some other agent may take a task and remove it from the actual list of free tasks.

Solution: check at each state wether or not the free tasks have been modified, if it is the case, replan accordingly. This is essentially what the *Basic agent logic* algorithm suggests but this is a somewhat more detailed implementation:

Algorithm 2: Replanning Implementation

```
Initialization: p_{opt} \leftarrow \text{compute\_optimal\_plan}(global.free\_tasks)
while goal not reached do

| if agent.free_task \neq global.free_tasks then
| p_{opt} \leftarrow \text{compute\_optimal\_plan}(global.free\_tasks)
end
| execute next action in p_{opt}
end
```

4 Tree Search algorithms

4.1 Breadth First Search

From the course:

Algorithm 3: Breadth-first(InitialNode)

```
Q \leftarrow (\operatorname{InitialNode}) \ \mathbf{repeat}
n \rightarrow \operatorname{first}(Q), \ Q \leftarrow \operatorname{rest}(Q)
\mathbf{if} \ n \ is \ a \ goal \ state \ \mathbf{then}
+ \ \mathbf{return} \ n
\mathbf{end}
S \leftarrow \operatorname{succ}(n)
Q \leftarrow \operatorname{append}(Q,S)
\mathbf{until} \ Q \ is \ empty;
\mathbf{return} \ FAIL
```

4.2 ASTAR

 $From\ the\ course$:

Algorithm 4: ASTAR(InitialNode)

```
Q \leftarrow (\operatorname{InitialNode}) \operatorname{\mathbf{repeat}}
n \rightarrow \operatorname{first}(Q), \ Q \leftarrow \operatorname{rest}(Q)
\operatorname{\mathbf{if}} n \ is \ a \ goal \ state \ \mathbf{then}
+ \operatorname{\mathbf{return}} n
\operatorname{\mathbf{end}}
\operatorname{\mathbf{if}} n \not\in C \ has \ a \ lower \ cost \ than \ its \ copy \ in \ C \ \mathbf{then}
- C \leftarrow \operatorname{append}(n,C)
- S \leftarrow \operatorname{succ}(n)
- S \leftarrow \operatorname{succ}(n)
- S \leftarrow \operatorname{sort}(S,f)
- Q \leftarrow \operatorname{merge}(Q,S,f) \ (Q \ is \ ordered \ by \ f(n) = g(n) + h(n))
- \operatorname{\mathbf{end}}
- \operatorname{\mathbf{until}} Q \ is \ empty;
- \operatorname{\mathbf{return}} FAIL
```

4.2.1 ASTAR Heuristics Definition

To be defined

4.3 Proving Optimality of Solution

To be done