Random Graphs

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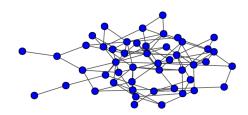
August 2019

Outline

- Erdos-Renyi models (ER)
- ② Chung-Lu model (CL)
- Configuration model

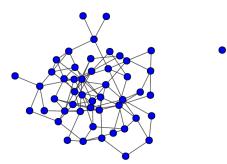
The G(n, m) model:

- n nodes
- m edges chosen at random from $N = \binom{n}{2}$ pairs
- average degree k = 2m/n



The G(n, p) model:

- n nodes
- each possible $N = \binom{n}{2}$ pair of nodes is connected with probability p
- expected number of edges is Np
- expected average degree k = p(n-1)



The G(n, p) model:

- with p = m/N, the expected average degree is k = 2m/n
- this is the model typically used in practice
- allows for easy calculation of graph statistics

Number of edges:

Let P_m the probability of getting m edges with the G(n, p) model, and let $N = \binom{n}{2}$.

$$P_m = \binom{N}{m} p^m (1-p)^{N-m}$$

Given that N is large and p is (typically) small, we can use the Poisson approximation to the binomial with $\lambda = Np$:

$$P_m \approx \frac{e^{-\lambda}\lambda^m}{m!}$$

Degree distribution:

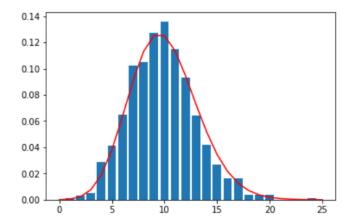
Let p_k , the probability that some node has degree k in G(n, p).

$$p_k = \binom{n-1}{k} p^k (1-p)^{n-1-k} \approx \frac{e^{-\lambda} \lambda^k}{k!}$$

with $\lambda = (n-1)p$, the expected average degree.

In view of the Poisson distribution, such models do not generate *hubs*, the high-degree nodes typically seen in real networks.

Degree distribution (n = 1000, p = .01):

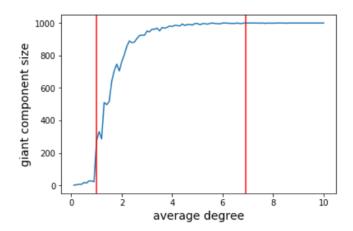


Giant connected component

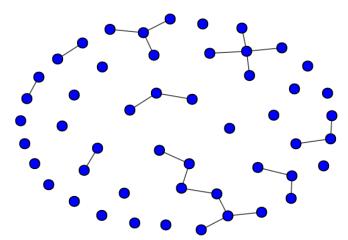
For G(n, p) with expected average degree k = p(n - 1), we distinguish 3 regimes:

- k < 1: subcritical regime; no giant component, clusters are mostly trees.
- k > 1: supercritical regime; single giant component, small clusters are mostly trees.

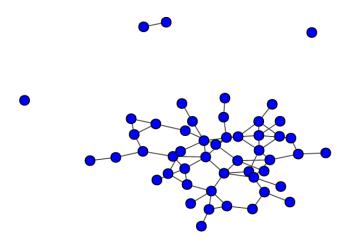
Giant connected component with n = 1000:



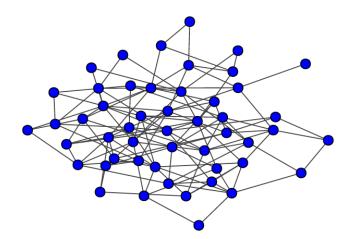
Subcritical regime:



Supercritical regime:



Connected regime:



Next algorithms taken from [WDS]:



LLNL-TR-678729

An In-Depth Analysis of the Chung-Lu Model

M. Winlaw, H. DeSterck, G. Sanders

October 28, 2015

Let *G* be a graph with vertices $V = \{v_1, \dots, v_n\}$.

Assume also the degree sequence: $k_i = deg_G(v_i)$.

Model I: probability of an edge between vertices v_i and v_j is given by:

$$p_{ij} = \frac{deg_G(v_i)deg_G(v_j)}{vol(V)}, \ i \neq j \text{ and } p_{ii} = \frac{(deg_G(v_i))^2}{2vol(V)}.$$

This is also known as the Bernoulli Chung-Lu model.

Model I:

Algorithm 1: Bernoulli Chung-Lu Algorithm

If we define $\mathcal{CL}_1(G)$ to be the distribution of graphs obtained with Model I, then $\mathbb{E}_{G'\sim\mathcal{CL}_1(G)}(deg_{G'}(v_i))=deg_G(v_i), 1\leq i\leq n$.

For
$$G' = (V, E') \sim \mathcal{CL}_1(G)$$
:

- $\mathbb{E}(|E'|) = |E|$, but we may have $|E'| \neq |E|$,
- there are no multi-edges, and
- there can be self-edges.

This algorithm is not useful in practice, as it requires $O(n^2)$ multinomial experiments.

Model II is more practical algorithm with O(|E|) steps only.

Generate a graph over vertices V by selecting |E| edges $e = (u_1, u_2)$ where each u_i is independently sampled from V according to the multinomial distribution where $p(v_i) = deg_G(v_i)/vol(V)$.

Edges can be repeated, so what we have are expected number of edges instead of probabilities.

Model II:

```
Algorithm 3: \mathcal{O}(|\text{Edges}|) Chung-Lu Algorithm
```

```
for k=1 to m do

Draw node i with probability \frac{k_i}{2m};

Draw node j with probability \frac{k_j}{2m};

*/* Add edge (i,j) to the graph */

if i \neq j then

a_{ij} = a_{ji} = 1;

else

a_{ii} = 2;

end

end
```

If we define $\mathcal{CL}_2(G)$ to be the distribution of graphs obtained with Model II allowing for multi-edges, then

$$\mathbb{E}_{G' \sim \mathcal{CL}_2(G)}(deg_{G'}(v_i)) = deg_G(v_i), 1 \leq i \leq n.$$

For
$$G' = (V, E') \sim \mathcal{CL}_2(G)$$
:

- we always have |E'| = |E|,
- there can be multi-edges, and
- there can be self-edges.

For Model II, we can ignore multi-edges, which reduces the overall expected degree (volume) of the graph.

For both models, we can ignore self-edges, again reducing the overall volume.

For model II:

Algorithm 4: $\mathcal{O}(|Edges|)$ Chung-Lu Algorithm without Self-Edges.

```
for k=1 to m do

Draw node i with probability \frac{k_i}{2m};

Draw node j with probability \frac{k_j}{2m};

/* Add edge (i,j) to the graph */

if i \neq j then

a_{ij} = a_{ji} = 1;

end

end
```

In summary:

CHAPTER 1. AN IN-DEPTH ANALYSIS OF THE CHUNG-LU MODEL

Model IV

Probability Probability Expected Degree of Edge (i, j)of Edge (i, i) $E(\mathbf{D}_i)$ Self-Edges: $\frac{k_i k_j}{2m}$ k_i Model I Bernoulli Chung-Lu No Self-Edges: $\frac{k_i k_j}{2m}$ $k_i - \frac{k_i^2}{2m}$ 0 Model III $\left| 1 - (1 - 2\frac{k_i k_j}{4m^2})^m \right| 1 - (1 - \frac{k_i^2}{4m^2})^m \left| \sum_{j \neq i} 1 - (1 - 2\frac{k_i k_j}{4m^2})^m \right|$ Self-Edges: $<\frac{k_i k_j}{2m}$ $<\frac{k_i^2}{4m}$ $+1-(1-\frac{k_i^2}{4m^2})^m < k_i$ Model II $\mathcal{O}(m)$ Chung-Lu $\sum_{j\neq i} 1 - (1 - 2\frac{k_i k_j}{4m^2})^m$ No Self-Edges: $1 - (1 - 2\frac{k_i k_j}{4m^2})^m$ $< k_i - \frac{k_i^2}{2}$ $<\frac{k_ik_j}{2m}$

Table 1.2: Model Summary



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The Chung-Lu models are probabilistic models for edge generation, with *expected* degree sequence.

For the *configuration model*, we specify a precise degree sequence for the nodes $d_1, ..., d_n$.

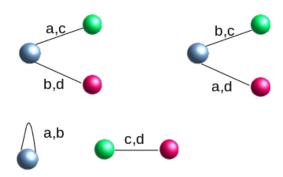
Each graph with *n* nodes and this exact degree sequence is assigned the same probability.

For each node i, we assign d_i stubs (half-edges);

Stubs are connected at random.



This can generate self-edges and multi-edges



For this model, the expected number of edges between nodes $i \neq j$ is:

$$e_{ij} = \frac{d_i d_j}{2m-1}$$

and:

$$e_{ii}=\frac{d_i(d_i-1)}{2(2m-1)}$$

The overall expected number of edges is $m = \frac{1}{2} \sum_{i} d_{i}$.

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Notebook #2