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INTEL'S OPTIMIZING COMPILER

Software Solutions Group Intel® Corporation

Agenda

Intel's Optimizing Compiler: Brief Overview/Recap

Vector Programming

Intel® AVX-512

Consistency of Floating-Point Results

High Bandwidth Memory

Some Tips for Using OpenMP*

Intel® Compiler: Brief Recap

Supports standards

- Fortran77, Fortan90, Fortran95, Fortran2003, most Fortran 2008
- Up to C99 (most of C11); C++11 and much C++14
- Intel® Fortran (and C/C++) binary compatible with gcc, gdb, ...
 - But <u>not</u> binary compatible with gfortran

Supports all instruction sets via vectorization (auto- and explicit)

OpenMP* 4.0 support, some 4.5, no user-defined reductions

Optimized math libraries (including for KNL)

Many advanced optimizations

With detailed, structured optimization reports

Drivers: icc, icpc, ifort

To set the environment: source compilervars.sh intel64

Basic Optimizations with icc/ifort -O...

- -OO no optimization; sets -g for debugging
- -O1 scalar optimizations
 - Excludes optimizations tending to increase code size
- -O2 **default** (except with -g)
 - includes auto-vectorization; some loop transformations such as unrolling; inlining within source file;
 - Start with this (after initial debugging at -O0)
- -O3 more aggressive loop optimizations
 - Including cache blocking, loop fusion, loop interchange, ...
 - May not help all applications; need to test

-qopt-report

Generates compiler optimization reports in files *.optrpt

Inter Procedural Optimization (IPO)

icc -ipo or ifort -ipo

Analysis & Optimization across function and source file boundaries, e.g.

Function inlining; Interprocedural constant propagation; Alignment analysis;
 Disambiguation; Data & Function Layout; etc.

2-step process:

- Compile phase objects contain intermediate representation
- "Link" phase compile and optimize over all such objects
 - Fairly seamless: the linker automatically detects objects built with -ipo, and their compile options
 - · May increase build-time and binary size
 - But can be done in parallel with -ipo[n], -ipo-jobs<n>
 - Entire program need not be built with IPO/LTO, just hot modules

Particularly effective for C++ apps with many smaller functions

Get report on inlined functions with -qopt-report-phase=ipo

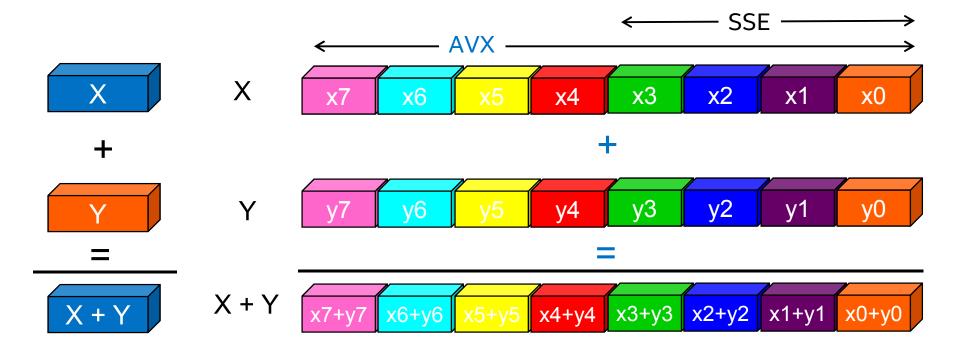


VECTOR PROGRAMMING

SIMD: Single Instruction, Multiple Data

- Scalar mode
 - one instruction produces one result
 - E.g. vaddss, (vaddsd)

- Vector (SIMD) mode
 - one instruction can produce multiple results
 - E.g. vaddps, (vaddpd)



Guidelines for Writing Vectorizable Code

Prefer simple "for" or "DO" loops

Write straight line code. Try to avoid:

- function calls (unless inlined or SIMD-enabled functions)
- branches that can't be treated as masked assignments

Avoid dependencies between loop iterations

Or at least, avoid read-after-write dependencies

Prefer arrays to the use of pointers

- Without help, the compiler often cannot tell whether it is safe to vectorize code containing pointers.
- Try to use the loop index directly in array subscripts, instead of incrementing a separate counter for use as an array address.
- Disambiguate function arguments, e.g. -fargument-noalias or "restrict"

Use efficient memory accesses

- Favor inner loops with unit stride
- Minimize indirect addressing a[i] = b[ind[i]]
- Align your data consistently where possible (to 16, 32 or 64 byte boundaries)

Example for Vectorizable Code: Search Loop Pattern

```
void mysearch(float *a,float *b, int n)
{
    int i;
    for (i = 0; i < n; i++)
    {
        if (a [i] < 0.) break;
        c[i] = sqrt(a[i]) * b[i];
    }
    return;
}</pre>
```

```
void mysearch(float *a,float *b, int n)
{
    int i, ns;
    for (i = 0; i < n; i++)
    {
        if (a [i] < 0. ) break;
    }
    ns = i;
    for (i = 0; i < ns; i++)
    {
        c[i] = sqrt(a[i]) * b[i];
    }
    return;
}</pre>
```

Left-hand side will not vectorize:

remark #15520: loop was not vectorized: loop with multiple exits cannot be vectorized unless it meets search loop idiom criteria

Right-hand side: Both loops will vectorize

- Simple search pattern of the first loop is recognized by the compiler
- · Computationally intensive second loop will be vectorized

Explicit SIMD (Vector) Programming

Modeled on OpenMP* for threading (explicit parallel programming)

```
#pragma omp simd <clauses> (for loops)
#pragma omp declare simd <clauses> (for functions and subroutines)
```

Enables reliable vectorization of complex loops that compiler can't autovectorize

E.g. outer loops

Directives are commands to the compiler, not hints

- Programmer is responsible for correctness (like OpenMP threading)
 - E.g. PRIVATE and REDUCTION clauses
- Overrides all dependency and cost-benefit analysis

Incorporated in OpenMP 4.0 ⇒ portable

-qopenmp or -qopenmp-simd to enable

```
void addit(double* a, double * b, int m, int n, int x) {
#pragma omp simd // If you know x<0
for (int i = m; i < m+n; i++) a[i] = b[i] + a[i-x];
}</pre>
```

Loops Containing Function Calls

- Function calls can have side effects that introduce a loop-carried dependency, preventing vectorization
- Possible remedies:
 - Inlining (-ip or -ipo)
 - best for small functions
 - !\$OMP SIMD directive (last resort)
 - Calls to the routine will be serialized
 - Remaining parts of the loop body will be vectorize
 - SIMD-enabled functions
 - Good for large, complex functions and in contexts where inlining is difficult
 - Two ways to vectorize:
 - o Call from regular DO loop
 - Adding "ELEMENTAL" keyword allows SIMD-enabled function to be called with array section argument

SIMD-enabled Subroutine

Compiler generates SIMD-enabled (vector) version of a scalar subroutine that can be called from a vectorized loop:

```
subroutine test_linear(x, y)

!$omp declare simd (test_linear) linear(ref(x, y))
    real(8),intent(in) :: x
    real(8),intent(out) :: y
    y = 1. + sin(x)**3
end subroutine test_linear
...
interface
...
do j = 1,n
    call test_linear(a(j), b(j))
enddo

| Important because arguments
passed by reference in Fortran

| remark #15301: FUNCTION WAS VECTORIZED.
```

SIMD-enabled routine must have <u>explicit interface</u>!\$omp simd not needed in simple cases like this

SIMD-enabled Subroutine

The LINEAR(REF) clause is very important

- In C, compiler places consecutive argument values in a vector register
- But Fortran passes arguments by reference
 - By default compiler places consecutive addresses in a vector register
 - Leads to a gather of the 4 addresses (slow)
 - LINEAR(REF(X)) tells the compiler that the addresses are consecutive;
 only need to dereference once and copy consecutive values to vector register
 - New in compiler version 16.0.1
- Same method could be used for C arguments passed by reference



How to Align Data (Fortran)

Align array on an "n"-byte boundary (n must be a power of 2)

```
!dir$ attributes align:n :: array
```

Works dynamically allocated, automatic and static arrays (not in common)

For a 2D array, choose column length to be a multiple of n, so that consecutive columns have the same alignment (pad if necessary)

-align array64byte compiler tries to align all array types

And tell the compiler...

```
!dir$ vector aligned OR
!$omp simd aligned( var [,var...]:<n>)
```

- Asks compiler to vectorize, assuming all array data accessed in loop are aligned for targeted processor
 - May cause fault if data are not aligned

```
!dir$ assume aligned array:n [,array2:n2, ...]
```

- Compiler may assume array is aligned to n byte boundary
- Typical use is for dummy arguments
- Extension for allocatable arrays v17

```
!DIR$ ASSUME_ALIGNED a(1):64, b(1):64, c(1):64, d(1):64
```

n=16 for Intel® SSE, n=32 for Intel® AVX, n=64 for Intel® AVX-512

How to Align Data (C/C++)

Allocate memory on heap aligned to n byte boundary:

```
void* _mm_malloc(int size, int n)
int posix_memalign(void **p, size_t n, size_t size)
void* aligned_alloc(size_t alignment, size_t size) (C11)
#include <aligned new> (C++11)
```

Alignment for variable declarations:

```
__attribute__((aligned(n))) var_name or
__declspec(align(n)) var_name
```

And tell the compiler...

```
#pragma vector aligned
#pragma omp simd aligned( var [,var...]:<n>)
```

- Asks compiler to vectorize, overriding cost model, and assuming all array data accessed in loop are aligned for targeted processor
- May cause fault if data are not aligned

```
__assume_aligned(array, n)
```

Compiler may assume array is aligned to n byte boundary

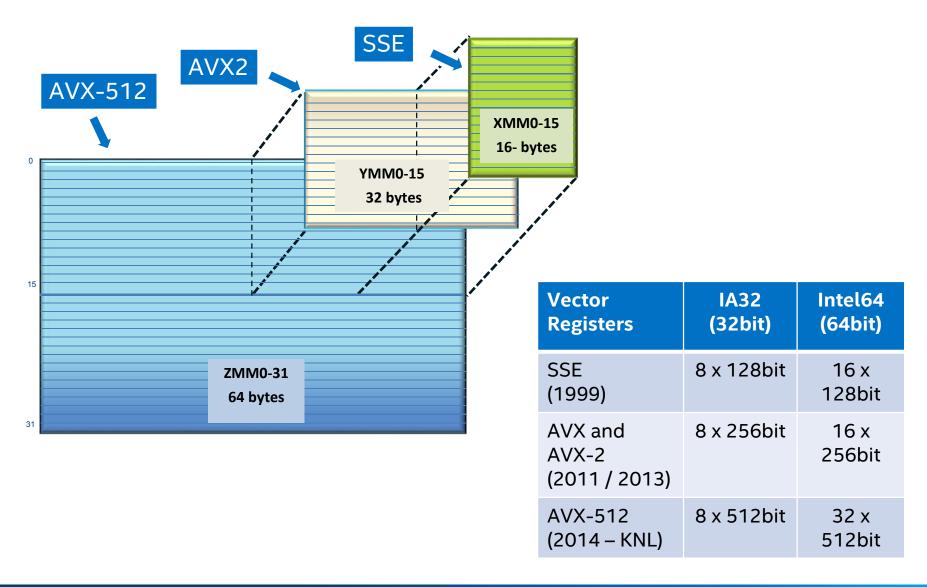


n=64 for Intel® Xeon Phi™ coprocessors, n=32 for Intel® AVX, n=16 for Intel® SSE

INTEL AVX-512 INSTRUCTION SET ARCHITECTURE

Changes for SW development resulting from new Intel® AVX-512 ISA

AVX-512 - Greatly increased Register File



AVX-512 - KNL and future Xeon

- KNL and future Intel[®] Xeon architecture share a large set of instructions
 - but sets are not identical
- Subsets are represented by individual feature flags (CPUID)

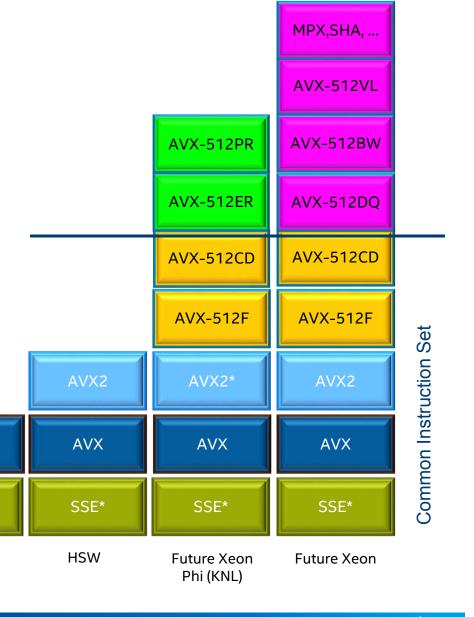
SSE*

NHM

AVX

SSE*

SNB



Intel® Compiler Switches Targeting Intel® AVX-512

Switch	Description
-xmic-avx512	KNL only
-xcore-avx512	Future Xeon only
-xcommon-avx512	AVX-512 subset common to both. Not a fat binary.
-m, -march, /arch	Not yet!
-axmic-avx512 etc.	Fat binaries. Allows to target KNL and other Intel® Xeon® processors
-qoffload-arch=mic-avx512	Offload to KNL coprocessor
-mmic	No – not for KNL

All supported in 16.0 and forthcoming 17.0 compilers

Binaries built for earlier Intel® Xeon® processors will run unchanged on KNL Binaries built for Intel® Xeon Phi™ coprocessors will not.

Optimization Improvements

Vectorization works as for other targets

- 512, 256 and 128 bit instructions available
- 64 byte alignment is best, like for KNC
- New instructions can help

Vectorization of compress/expand loops:

Uses vcompress/vexpand on KNL

```
for (int i; i <N; i++) {
    if (a[i] > 0) {
        b[j++] = a[i]; //
    compress
        c[i] = a[k++]; //
    expand
      }
}
• Cross-iteration
    dependencies
    by i and k
```

Convert certain gathers to vector loads

Can auto-generate Conflict Detection instructions (AVX512CD)

Compress Loop with Intel® AVX2

```
nb = 0
  do ia=1, na     ! line 11
  if(a(ia) > 0.) then
    nb = nb + 1
    b(nb) = a(ia)
  endif
enddo
```

With Intel® AVX2, does not auto-vectorize

And vectorizing with SIMD would be too inefficient

```
ifort -c -xcore-avx2 -qopt-report-file=stderr -qopt-report=3 -qopt-report-phase=vec compress.f90
...

LOOP BEGIN at compress.f90(11,3)

remark #15344: loop was not vectorized: vector dependence prevents vectorization.

First dependence is shown below. Use level 5 report for details

remark #15346: vector dependence: assumed ANTI dependence between line 13 and line 13
LOOP END
```

C code behaves the same

Compress Loop with Intel® AVX-512

Compile for KNL:

```
ifort -c -qopt-report-file=stderr -qopt-report=3 -qopt-report-phase=vec -xmic-avx512 compress.f90
LOOP BEGIN at compress.f90(11,3)
 remark #15300: LOOP WAS VECTORIZED
 remark #15450: unmasked unaligned unit stride loads: 1
 remark #15457: masked unaligned unit stride stores: 1
 remark #15478: estimated potential speedup: 14.040
 remark #15497: vector compress: 1
LOOP END
grep vcompress compress.s
   vcompressps %zmm4, -4(%rsi,%rdx,4){%k1}
                                                     #14.7 c7 stall 1
   vcompressps %zmm1, -4(%rsi,%r12,4){%k1}
                                                      #14.7 c5
   vcompressps %zmm1, -4(%rsi,%r12,4){%k1}
                                                      #14.7 c5
                                                     #14.7 c7 stall 1
   vcompressps %zmm4, -4(%rsi,%rdi,4){%k1}
```

Observed speed-up is substantial but depends on problem size, data layout, etc.

Adjacent Gather Optimizations

Or "Neighborhood Gather Optimizations"

```
do j=1,n

y(j) = x(1,j) + x(2,j) + x(3,j) + x(4,j) ....
```

- Elements of x are adjacent in memory, but vector index is in other dimension
- Compiler generates simd loads and shuffles for x instead of gathers
 - Before AVX-512: gather of x(1,1), x(1,2), x(1,3), x(1,4)
 - With AVX-512: SIMD loads of x(1,1), x(2,1), x(3,1), x(4,1) etc., followed by permutes to get back to x(1,1), x(1,2), x(1,3), x(1,4) etc.
 - Message in optimization report: remark #34029: adjacent sparse (indexed) loads optimized for speed
- Arrays of short vectors or structs are very common

Histogramming with Intel® AVX2

```
! Accumulate histogram of sin(x) in h
do i=1,n
    y = sin(x(i)*twopi)
    ih = ceiling((y-bot)*invbinw)
    ih = min(nbin,max(1,ih))
    h(ih) = h(ih) + 1     ! line 15
enddo
```

With Intel® AVX2, this does not vectorize

- Store to h is a scatter
- ih can have the same value for different values of i
- Vectorization with a SIMD directive would cause incorrect results

```
ifort -c -xcore-avx2 histo2.f90 -qopt-report-file=stderr -qopt-report-phase=vec LOOP BEGIN at histo2.f90(11,4)
```

remark #15344: loop was not vectorized: vector dependence prevents vectorization.

First dependence is shown below. Use level 5 report for details

remark #15346: vector dependence: assumed FLOW dependence between line 15 and line 15 LOOP END

Histogramming with Intel® AVX-512

Compile for KNL:

```
ifort -c -xmic-avx512 histo2.f90 -qopt-report-file=stderr -qopt-report=3 –S
...

LOOP BEGIN at histo2.f90(11,4)
remark #15300: LOOP WAS VECTORIZED
remark #15458: masked indexed (or gather) loads: 1
remark #15459: masked indexed (or scatter) stores: 1
remark #15478: estimated potential speedup: 13.930
```

LOOP END

vpconflictd vpgatherdd vptestmd kmovw vpaddd	%zmm5, %zmm21, %zmm3 %zmm3, %zmm1 -4(%rsi,%zmm3,4), %zmm6{%k1} %zmm18, %zmm1, %k0 %k0, %r10d %zmm19, %zmm6, %zmm2	#14.7 c19 #15.7 c21 #15.15 c21 #15.7 c23 #15.7 c27 stall 1 #15.7 c27
testl vpscatterdd	%r10d, %r10d %zmm2, -4(%rsi,%zmm3,4){%k1}	#15.7 c3

remark #15499: histogram: 2

Histogramming with Intel® AVX-512 (cont'd)

Observed speed-up between AVX2 (non-vectorized) and AVX512 (vectorized) can be large, but depends on problem details

- Comes mostly from vectorization of other heavy computation in the loop
 - Not from the scatter itself
- Speed-up may be (much) less if there are many conflicts
 - E.g. histograms with a singularity or narrow spike

Other problems map to this

E.g. energy deposition in cells in particle transport Monte Carlos



Intel® Software Development Emulator (SDE)

Intel Compilers Already Recognize Intel® AVX-512 and Will Generate KNL Code Use Intel® Software Development Emulator (SDE) to Test Code

- Will test instruction mix, not performance
- Does not emulate hardware (e.g. memory hierarchy) only ISA

Use the SDE to Answer

- Is my compiler generating Intel® AVX-512/KNL-ready code?
- How do I restructure my code so that Intel® AVX-512 code is generated?

Visit <u>Intel Xeon Phi Coprocessor code named "Knights Landing" - Application</u>
<u>Readiness</u>

FLOATING POINT CONSISTENCY

Dealing with FP consistency when moving from Intel® MIC or Intel Xeon to KNL

Floating-Point Reproducibility

- -fp-model precise disables most value-unsafe optimizations (especially re-associations)
- The primary way to get consistency between different platforms or different optimization levels

Example of **disabled** optimizations:

- Vectorization of loops containing transcendental functions
- Fast, approximate division and square roots
- Flush-to-zero of denormals
- Vectorization of reduction loops
- Evaluation of constant expressions at compile time
- ...

Does not prevent differences due to:

- Different implementations of math functions:
 - For consistency of math functions between KNL and Intel® Xeon® processors use
 - -fimf-arch-consistency=true available for Xeon and KNL
- Use of fused multiply-add instructions (FMAs): see next slide

FMAs

Most common cause of differences between Intel® Xeon® processors and Intel® Xeon Phi™

- Not disabled by -fp-model precise
- Can disable for testing with -no-fma
- Or by function-wide pragma or directive:

```
#pragma float_control(fma,off)
```

!dir\$ nofma

- With some impact on performance
- -fp-model strict disables FMAs, amongst other things
 - But on KNC, results in non-vectorizable x87 code
- The fma() intrinsic in C should always give a result with a single rounding, even on processors with no FMA instruction

FMAs

Can cause issues even when both platforms support them (e.g. Haswell and KNL)

- Optimizer may not generate them in the same places
 - No language rules
- FMAs may break the symmetry of an expression:

```
c = a; d = -b;
result = a*b + c*d; ( = 0 if no FMAs )
```

If FMAs are supported, the compiler may convert to either

```
result = fma(c, d, (a*b)) or result = fma(a, b, (c*d))
```

Because of the different roundings, these may give results that are non-zero and/or different from each other.

Bottom Line for FP consistency

To get consistent results between KNL and Intel Xeon processors, use

-fp-model precise -fimf-arch-consistency=true -no-fma

(you could try omitting -no-fma for Xeon processors that support FMA, but FMA's could still possibly lead to differences)

To come: One-Stop-Shop" in the 17.0 compiler:

-fp-model consistent

Note: even with this there will be the usual caveats: threaded reductions, math libraries, different compilation targets, SIMD directives...)

To get consistent results that are as close as possible between KNC and Intel® Xeon® processors or KNL, try

-fp-model precise -no-fma on both.

Tips for Debugging OpenMP applications

Build with -qopenmp to generate threaded code

- but run a single thread, OMP_NUM_THREADS=1
- If it works, try Intel® Inspector XE to detect threading errors
- If still fails, excludes race conditions or other synchronization issues as cause

Build with -openmp-stubs -auto

- RTL calls are resolved; but no threaded code is generated
- allocates local arrays on the stack, as for OpenMP
- If works, check for missing FIRSTPRIVATE, LASTPRIVATE
- If still fails, eliminates threaded code generation as cause
- If works without -auto, implicates changed memory model
 - Perhaps insufficient stack size
 - Perhaps values not preserved between successive calls

Additional Resources (Optimization)

Webinars:

https://software.intel.com/videos/getting-the-most-out-of-the-intel-compiler-with-new-optimization-reports

https://software.intel.com/videos/new-vectorization-features-of-the-intel-compiler https://software.intel.com/articles/further-vectorization-features-of-the-intel-compiler-webinar-code-samples

https://software.intel.com/videos/from-serial-to-awesome-part-2-advanced-code-vectorization-and-optimization

https://software.intel.com/videos/data-alignment-padding-and-peel-remainder-loops

Vectorization Guide (C):

https://software.intel.com/articles/a-guide-to-auto-vectorization-with-intel-c-compilers/

Explicit Vector Programming in Fortran:

https://software.intel.com/articles/explicit-vector-programming-in-fortran

Initially written for Intel® Xeon Phi™ coprocessors, but also applicable elsewhere:

https://software.intel.com/articles/vectorization-essential

https://software.intel.com/articles/fortran-array-data-and-arguments-and-vectorization

Compiler User Forums at http://software.intel.com/forums

Additional Resources

https://software.intel.com/articles/intel-xeon-phi-coprocessor-code-named-knights-landing-application-readiness

https://software.intel.com/sites/default/files/managed/4c/1c/parallel_mag_issue20.pdf

https://software.intel.com/articles/what-disclosures-has-intel-made-about-knights-landing

https://software.intel.com/articles/intel-software-development-emulator

https://github.com/memkind

https://software.intel.com/articles/consistency-of-floating-point-results-using-the-intel-compiler

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Notice revision #20110804



Optimization Report Example

\$ icc -c -qopt-report=5 -xavx -fargument-noalias foo.c Begin optimization report for: foo LOOP BEGIN at foo.c(4,7) remark #15389: vectorization support: reference theta has unaligned access [foo.c(5,20)] remark #15389: vectorization support: reference sth has unaligned access [foo.c(5,11)] remark #15381: vectorization support: unaligned access used inside loop body remark #15305: vectorization support: vector length 4 remark #15309: vectorization support: normalized vectorization overhead 0.043 remark #15417: vectorization support: number of **FP up converts**: single precision to double precision 1 [foo.c(5,20)] remark #15418: vectorization support: number of **FP down converts**: double precision to single precision 1 [foo.c(5,11)] remark #15300: LOOP WAS VECTORIZED remark #15450: unmasked unaligned unit stride loads: 1 remark #15451: unmasked unaligned unit stride stores: 1 remark #15475: --- begin vector loop cost summary ---#include <math.h> remark #15476: scalar loop cost: 114

#include <math.h>
void foo (float * theta, float * sth) {
 int i;
 for (i = 0; i < 128; i++)
 sth[i] = sin(theta[i]+3.1415927);
}</pre>

LOOP END

remark #15487: type converts: 2

remark #15477: vector loop cost: 40.750

remark #15478: estimated potential speedup: 2.790

remark #15488: --- end vector loop cost summary --- remark #25015: Estimate of max trip count of loop=32

remark #15482: vectorized math library calls: 1

Optimization Report Example

\$ icc -c -qopt-report=5 -fargument-noalias -xavx foo.c

```
Begin optimization report for: foo
LOOP BEGIN at foo.c(6,7)
 remark #15388: vectorization support: reference theta has aligned access [foo.c(7,20)]
 remark #15388: vectorization support: reference sth has aligned access [foo.c(7,11)]
 remark #15412: vectorization support: streaming store was generated for sth [foo.c(7,11)]
 remark #15305: vectorization support: vector length 8
 remark #15309: vectorization support: normalized vectorization overhead 0.013
 remark #15300: LOOP WAS VECTORIZED
 remark #15448: unmasked aligned unit stride loads: 1
 remark #15449: unmasked aligned unit stride stores: 1
                                                           #include <math.h>
 remark #15467: unmasked aligned streaming stores: 1
                                                           void foo (float * theta, float * sth) {
 remark #15475: --- begin vector loop cost summary ---
                                                            int i:
 remark #15476: scalar loop cost: 110
                                                             assume aligned(theta,32);
 remark #15477: vector loop cost: 9.870
 remark #15478: estimated potential speedup: 11.130
                                                               assume aligned(sth,32);
 remark #15482: vectorized math library calls: 1
                                                           For (I = 0; I < 2000000; i++)
 remark #15488: --- end vector loop cost summary ---
                                                              sth[i] = sinf(theta[i]+3.1415927f);
 remark #25015: Estimate of max trip count of loop=250000
LOOP END
```

Estimated performance speedup from type conversion: 9.7
Additional performance gain due to alignment and streaming stores

Optimization Report Example

\$ icc -c -qopt-report=5 -fargument-noalias -xmic-avx512 foo.c

```
Begin optimization report for: foo
LOOP BEGIN at foo.c(6,7)
 remark #15388: vectorization support: reference theta has aligned access [foo.c(7,20)]
 remark #15388: vectorization support: reference sth has aligned access [foo.c(7,11)]
 remark #15412: vectorization support: streaming store was generated for sth [foo.c(7,11)]
 remark #15305: vectorization support: vector length 16
 remark #15309: vectorization support: normalized vectorization overhead 0.013
 remark #15300: LOOP WAS VECTORIZED
 remark #15448: unmasked aligned unit stride loads: 1
 remark #15449: unmasked aligned unit stride stores: 1
 remark #15467: unmasked aligned streaming stores: 1
 remark #15475: --- begin vector loop cost summary ---
 remark #15476: scalar loop cost: 111
 remark #15477: vector loop cost: 4.930
 remark #15478: estimated potential speedup: 22.480
 remark #15482: vectorized math library calls: 1
 remark #15488: --- end vector loop cost summary ---
 remark #25015: Estimate of max trip count of loop=125000
LOOP END
```

```
#include <math.h>
void foo (float * theta, float * sth) {
 int i:
 assume aligned(theta,64);
   assume aligned(sth,64);
for (i = 0; i < 2000000; i++)
  sth[i] = sinf(theta[i]+3.1415927f);
```

