

Exercises

from Riehl, *Category Theory in Context*

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I am self-studying this alongside Aluffi, *Algebra: Chapter 0*, so there is a bit of mixed notation.

Chapter 1

Categories, Functors, Natural Transformations

1.1 Abstract and concrete categories

Exercise 1.1.i.

- (i) Consider a morphism $f : x \rightarrow y$. Show that if there exists a pair of morphisms $g, h : y \rightrightarrows x$ so that $gf = 1_x$ and $fh = 1_y$, then $g = h$ and f is an isomorphism.
- (ii) Show that a morphism can have at most one inverse isomorphism.

- (i) *Proof.* We can compose together $gfh = (gf)h = 1_x h = h$. But $gfh = g(fh) = g1_y = g$ by associativity. Therefore, $g = h$. Since $gf = 1_x$ and $fh = fg = 1_y$, f is an isomorphism. \square
- (ii) *Proof.* Suppose $f : x \rightarrow y$ has two inverses g and h . Then, as above, $g = g1_y = g(fh) = (gf)h = 1_x f = h$ by associativity. \square

Exercise 1.1.ii. Let \mathbf{C} be a category. Show that the collection of isomorphisms in \mathbf{C} defines a subcategory, the maximal groupoid inside \mathbf{C} .

Proof. Let X, Y , and Z be objects in \mathbf{C} . We must show identity and composition, since we get associativity for free from \mathbf{C} .

The identity morphisms 1_X are isomorphisms, since they are their own inverses. Therefore, they are present in the maximal groupoid.

Now, if $f : X \rightarrow Y$ and $g : Y \rightarrow Z$ are isomorphisms, then the composition $gf : X \rightarrow Z$ is an isomorphism with inverse $f^{-1}g^{-1}$ since $gff^{-1}g^{-1} = gg^{-1} = 1_Z$ and $f^{-1}g^{-1}gf = f^{-1}f = 1_X$.

Therefore, the maximal groupoid is in fact a category. \square

Exercise 1.1.iii. For any category \mathbf{C} and any object $c \in \mathbf{C}$, show that:

- (i) There is a category c/\mathbf{C} whose objects are morphisms $f : c \rightarrow x$ with domain c and in which a morphism from $f : c \rightarrow x$ to $g : c \rightarrow y$ is a map $h : x \rightarrow y$ between the codomains so that the triangle

$$\begin{array}{ccc} & c & \\ f \swarrow & & \searrow g \\ x & \xrightarrow{h} & y \end{array}$$

commutes, i.e., so that $g = hf$.

- (ii) There is a category \mathbf{C}/c whose objects are morphisms $f : x \rightarrow c$ with codomain c and in which a morphism from $f : x \rightarrow c$ to $g : y \rightarrow c$ is a map $h : x \rightarrow y$ between the codomains so that the triangle

$$\begin{array}{ccc} x & \xrightarrow{h} & y \\ f \searrow & & \swarrow g \\ & c & \end{array}$$

commutes, i.e., so that $f = gh$.

The categories c/\mathbf{C} and \mathbf{C}/c are called slice categories of \mathbf{C} under and over c , respectively.

(i) See Exercise I.3.7, Aluffi.

(ii) See Example I.3.5, Aluffi.

1.2 Duality

Exercise 1.2.i. Defining \mathbf{C}/c to be $(c/(\mathbf{C}^{\text{op}}))^{\text{op}}$, deduce Exercise 1.1.iii(ii) from Exercise 1.1.iii(i).

Proof. We must establish that \mathbf{C}/c is in fact $(c/(\mathbf{C}^{\text{op}}))^{\text{op}}$. Then, everything (inverses, composition, associativity) follows immediately from duality.

First, notice that in $c/(\mathbf{C}^{\text{op}})$, we have objects that are morphisms in \mathbf{C}^{op} , i.e., the same that are in \mathbf{C} but backwards:

$$\begin{array}{ccc} & c & \\ f^{\text{op}} \swarrow & & \searrow g^{\text{op}} \\ x & \xrightarrow{h^{\text{op}}} & y \end{array}$$

To get to the desired commutative diagram, we have to apply the opposite operation once more. \square

Exercise 1.2.ii.

- (i) Show that a morphism $f : x \rightarrow y$ is a split epimorphism in a category \mathbf{C} if and only if for all $c \in \mathbf{C}$, the post-composition function $f_* : \mathbf{C}(c, x) \rightarrow \mathbf{C}(c, y)$ is surjective.
- (ii) Argue by duality that f is a split monomorphism if and only if for all $c \in \mathbf{C}$, the pre-composition function $f^* : \mathbf{C}(y, c) \rightarrow \mathbf{C}(x, c)$ is surjective.

(i) *Proof.* (\Rightarrow) Suppose f is a split epimorphism, i.e., it is a left inverse of some morphism $g : y \rightarrow x$. That is, $fg = 1_y$.

Let $h : c \rightarrow y$. We must show $f_*(j) = fj = h$ for some $j : c \rightarrow x$. Notice that $f(gh) = (fg)h = 1_y h = h$. Therefore, if we let $j := gh$, we are done.

(\Leftarrow) Suppose f_* is surjective, i.e., for all $g : c \rightarrow y$ there exists an $h : c \rightarrow x$ such that $f_*(h) = fh = g$. In particular, for $g' = 1_y$, there exists $h' : y \rightarrow x$ such that $fh' = 1_y$. That is, f is a split epimorphism, the retraction of h' . \square

(ii) *Proof.* Apply part (i) to the category \mathbf{C}^{op} :

$f^{\text{op}} \in \mathbf{C}^{\text{op}}(x, y)$ is a split epimorphism if and only if for all $c \in \mathbf{C}^{\text{op}}$, the post-composition function $f_*^{\text{op}} : \mathbf{C}^{\text{op}}(c, x) \rightarrow \mathbf{C}^{\text{op}}(c, y)$ is surjective.

But this is exactly the same as saying

$f \in \mathbf{C}(y, x)$ is a split monomorphism if and only if for all $c \in \mathbf{C}$, the pre-composition function $f^* : \mathbf{C}^{\text{op}}(x, c) \rightarrow \mathbf{C}^{\text{op}}(y, c)$ is surjective.

because $\mathbf{C}^{\text{op}}(x, y) = \mathbf{C}(y, x)$ and a split epimorphism $f^{\text{op}}g^{\text{op}} = 1_y$ becomes a split monomorphism $gf = 1_y$. \square

Exercise 1.2.iii. Prove Lemma 1.2.11 by proving either (i) and (ii), then arguing by duality. Conclude that the monomorphisms in any category define a subcategory of that category and dually that the epimorphisms also define a subcategory.

Proof. (i) Let $f : x \rightarrow y$ and $g : y \rightarrow z$. We must show that gf is monic, i.e., $gfh = gfk \implies h = k$ for all $h, k : c \rightrightarrows x$. Suppose $gfh = gfk$. Since g is monic, we know that $g(fh) = g(fk) \implies fh = fk$. But then, since f is monic, $fh = fk \implies h = k$.

(ii) Let $f : x \rightarrow y$ and $g : y \rightarrow z$ such that $gf : x \rightarrow z$. We must show f is monic, i.e., $fh = fk \implies h = k$ for all $h, k : c \rightrightarrows x$. Suppose $fh = fk$. Then, by pre-composing g , $(gf)h = (gf)k$. Since gf is monic, $h = k$.

Now, we can show that the subcategory $\mathbf{C}_{\text{Monic}}$ of the same objects as \mathbf{C} and with only its monomorphisms is a category. First, since for all $h, k \in \mathbf{C}(c, x)$, $1_x h = 1_x k \implies h = k$, we have the identities in $\mathbf{C}_{\text{Monic}}$. Then, due to (i), the compositions of all morphisms are in the subcategory. Since the objects are unchanged, we can conclude $\mathbf{C}_{\text{Monic}}$ is a subcategory.

Dually, the opposite monomorphism f^{op} in \mathbf{C}^{op} such that $f^{\text{op}}h^{\text{op}} = f^{\text{op}}k^{\text{op}} \implies h^{\text{op}} = k^{\text{op}}$ is the epimorphism f in \mathbf{C} such that $hf = kf \implies h = k$. Therefore, the dual of (i) is (i') and the dual of (ii) is (ii'). Finally, the dual of the existence of the subcategory $\mathbf{C}_{\text{Monic}}$ is a subcategory \mathbf{C}_{Epic} of epimorphisms. \square

Exercise 1.2.vi. Prove that a morphism that is both a monomorphism and a split epimorphism is necessarily an isomorphism. Argue by duality that a split monomorphism that is an epimorphism is also an isomorphism.

Proof. Suppose $f : x \rightarrow y$ is also a split epimorphism such that $fg = 1_y$ for some $g : y \rightarrow x$. Then, since f is monic, $f(gf) = (fg)f = (1_y)f = f = f(1_x) \implies gf = 1_x$. That is, f is an isomorphism.

By duality, a split epimorphism $f^{\text{op}}g^{\text{op}} = 1_y$ is a split homomorphism $gf = 1_y$. Therefore, a split homomorphism that is an epimorphism is an isomorphism too. \square

Exercise 1.2.vii. Regarding a poset (P, \leq) as a category, define the supremum of a subcollection of objects $a \in P$ in such a way that the dual statement defines the infimum. Prove that the supremum of a subset of objects is unique, whenever it exists, in such a way that the dual proof demonstrates the uniqueness of the infimum.

Proof. Let a be the object such that $f_x : x \rightarrow a$ exists for all other x in the subcollection. That is, for every object x , $y \leq a$. Then, the dual definition is the object a^{op} in P^{op} such that $f_X^{\text{op}} : x \rightarrow a$ always exists, i.e., the object a' in P such that $f_x : a \rightarrow x$ always exists.

Suppose there are two suprema a and a' . Then, by the definition for a , we just have a morphism $a \rightarrow a'$. But if a' is a supremum, there must also be a morphism $a' \rightarrow a$. Therefore, $a \leq a' \leq a$, which means $a = a'$ since P is a poset (and not a pre-ordered set). \square