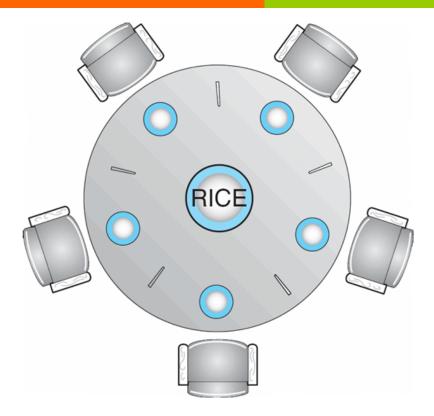
Dining-Philosophers Problem



- Shared data
 - Bowl of rice (data set)
 - Semaphore chopstick [5] initialized to 1

Dining-Philosophers Problem (Cont.)

☞ The structure of Philosopher *i*:

```
while (1) {
    P ( chopstick[i] );
    P ( chopStick[ (i + 1) % 5] );
    // eat
    V ( chopstick[i] );
    V (chopstick[ (i + 1) % 5] );
    // think
}
```

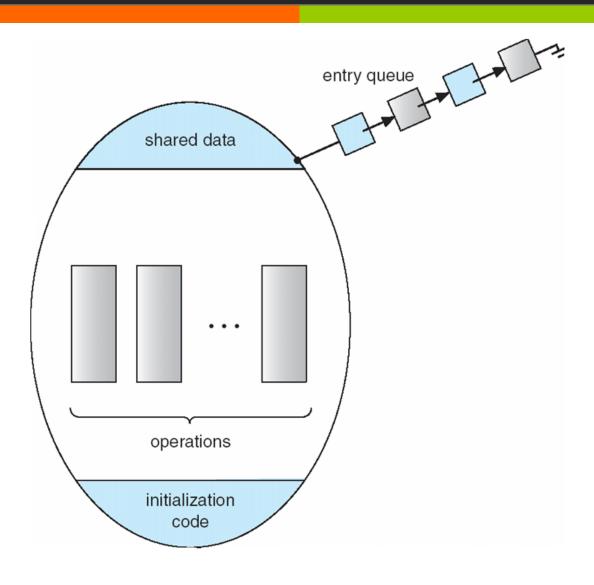
Monitors

- Higher-level, considered more intuitive than semaphores
- Justify states a language construct (i.e. C doesn't have it)
- Consists of shared data, subroutines, initialization code
- Only one process can be in any monitor subroutine at any time, all others are forced to wait

Monitors

- To use a monitor to ensure mutual exclusion for a critical section:
 - put the CS in a subroutine that can be called by all processes that want access to the CS
 - put that subroutine in a monitor (e.g. put it in the "procedures" section of the monitor)
- easy to extend to several related critical sections
 - e.g. list manipulation routines

Schematic View of a Monitor

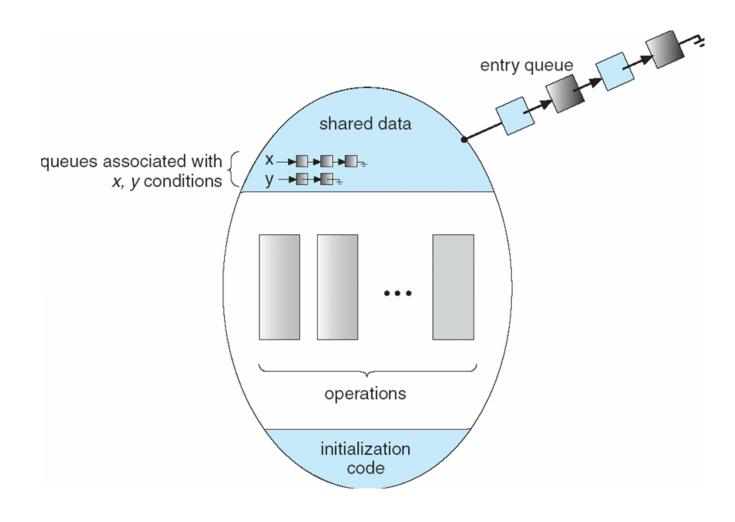


Monitors

Wait(C)

- suspends the current process until another process calls Signal(C), C is a *condition variable*
 - a suspended process is considered to be "out of the monitor" so other processes can access monitor routines
- **ℬ** Signal(C)
 - if there is a process suspended on C, resume it and wait until it has left the monitor before proceeding
 - otherwise, continue execution
- Notify(C)
 - like Signal, but resume the suspended process after the notifying process leaves the monitor

Monitor with Condition Variables



Binary Semaphore via a Monitor

```
MONITOR SEM
    -shared data:
        int busy;
         condition nonbusy;
    -Procedures: P() {
                          if (busy)
                             nonbusy.wait();
                          busy++;
                 V() {
                          busy = 0;
                          nonbusy.signal();
    -Init Code:
                  begin() {
                          busy = 0;
END SEM
```

Producer/Consumer via Monitor

MONITOR LISTMON

```
- shared data:
   LIST itemList;
   int maxListSize = N;
   condition bufavail, itemavail;
-Procedures:
   enqueueItem(item) {
      if(ListCount(itemList) == maxListSize)
         bufAvail.wait();
      ListPrepend(itemList, item);
      itemAvail.signal();
```

Producer/Consumer via Monitor

```
getItem() {
          if(ListCount(itemList) == 0) /* list is empty */
            itemAvail.wait();
         item = ListTrim(itemList);
          bufAvail.signal();
          return(item);
   -Init Code:
       itemList = ListCreate();
END LISTMON
```