Semaphore Implementation

- Typically, semaphore operations are implemented within the operating system itself.
- Instead of a busy wait (while S <= 0;), the operating system can block the process
 - Saves the CPU for the use of other processes
- A process executing a V may cause another previously blocked process to be put on the ready queue.

Semaphore Implementation

- How to ensure P and V are atomic?
 - uni-processor: disable interrupts (P & V are fairly quick)
 - multi-processor: software or hardware solutions
 - busy wait is negligible for P & V

Readers-Writers Problem

- A data set is shared among a number of concurrent processes
 - Readers only read the data set; they do **not** perform any updates
 - → Writers can both read and write
- Problem allow multiple readers to read at the same time. Only one single writer can access the shared data at the same time

Readers-Writers Problem (Cont.)

- Shared Data
 - Data set
 - **尽** Semaphore mutex − to control access to critical section
 - **尽** Semaphore wrt − to ensure mutual exclusivity when writing
 - Integer readcount − to count the readers

Readers-Writers Problem (Cont.)

■ The structure of a writer process:

Readers-Writers Problem (Cont.)

```
The structure of a reader process:
                                SEM mutex = 1;
while (1) {
                               SEM wrt = 1;
           P (mutex);
                                int readcount = 0;
           readcount++;
           if (readcount == 1)
                  P (wrt);
           V (mutex)
                // reading is performed
            P (mutex);
            readcount--;
            if (readcount == 0)
                 V (wrt);
            V (mutex);
```