

# A Guide to the agop 0.2-0 Package for R

## *Aggregation Operators and Preordered Sets in R*

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*Any suggestions and comments are welcome!*

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## 1 Getting Started

*“The process of combining several numerical values into a single representative one is called **aggregation**, and the numerical function performing this process is called **aggregation function**. This simple definition demonstrates the size of the field of application of aggregation: applied mathematics (e.g. probability, statistics, decision theory), computer science (e.g. artificial intelligence, operation research), as well as many applied fields (economics and finance, pattern recognition and image processing, data fusion, multicriteria decision making, automated reasoning etc.). Although history of aggregation is probably as old as mathematics (think of the arithmetic mean), its existence has reminded underground till only recent (...)” [30, p. xiii]*

R [44] is a free, open source software environment for statistical computing and graphics, which includes an implementation of a very powerful and quite popular high-level language called S. It runs on all major operating systems, i.e. Windows, Linux, and MacOS X. To install R and/or find some information on the S language please visit R Project’s Homepage at [www.R-project.org](http://www.R-project.org). Perhaps you may also wish to install RStudio, a convenient development environment for R. It is available at [www.rstudio.org](http://www.rstudio.org).

*agop* is an open source (licensed under GNU LGPL 3) package for  $R \geq 2.12$  to which anyone can contribute. It started as a fork of the CITAN (*Citation Analysis Toolpack*, [19]) package.

To install latest “official” release of the package available on *CRAN* we type<sup>1</sup>:

```
install.packages('agop')
```

Alternatively, we may fetch its current development snapshot from *GitHub*:

```
install.packages('devtools')  
devtools::install_github('agop', 'Rexamine')
```

Note that in this case you will need a working C/C++ compiler<sup>2</sup>.

Each session with *agop* should be preceded by a call to:

```
library('agop') # Load the package
```

To view the main page of the manual we type:

```
library(help='agop')
```

For more information please visit the package’s homepage [24]. In case of any problems, comments, or suggestions feel free to contact the authors. Good luck!

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<sup>1</sup>You are viewing the **development** version of the tutorial. Some of the features presented in this document may be missing in the CRAN release. Please, upgrade to the **latest** development version from *GitHub* if you need the new functionality. Note that you will need a working C/C++ compiler.

<sup>2</sup>Windows users should have Rtools installed, see [cran.r-project.org/bin/windows/Rtools/](http://cran.r-project.org/bin/windows/Rtools/).

## 2 Theoretical Background

Let us establish some basic notation convention used throughout this tutorial. From now on let  $\mathbb{I} = [a, b]$ , possibly with  $a = -\infty$  or  $b = \infty$ . Note that in many practical situations we commonly choose  $\mathbb{I} = [-1, 1]$ ,  $\mathbb{I} = [0, 1]$  or  $\mathbb{I} = [0, \infty]$ . A set of all vectors of arbitrary length with elements in  $\mathbb{I}$  is denoted by  $\mathbb{I}^{1,2,\dots} = \bigcup_{n=1}^{\infty} \mathbb{I}^n$ .

For two equal-length vectors  $\mathbf{x}, \mathbf{y} \in \mathbb{I}^n$  we write  $\mathbf{x} \leq_n \mathbf{y}$  if and only if for all  $i = 1, \dots, n$  it holds  $x_i \leq y_i$ . Moreover, all binary arithmetic operations on vectors  $\mathbf{x}, \mathbf{y} \in \mathbb{I}^n$  will be performed element-wise, e.g.  $\mathbf{x} + \mathbf{y} = (x_1 + y_1, \dots, x_n + y_n) \in \mathbb{I}^n$ . Similar behavior is assumed for  $-$ ,  $\cdot$ ,  $/$ ,  $\wedge$  (min),  $\vee$  (max), etc. Additionally, each function of one variable  $f : \mathbb{I} \rightarrow \mathbb{I}$  can be extended to the vector space: we write  $f(\mathbf{x})$  to denote  $(f(x_1), \dots, f(x_n))$ .

Let  $x_{(i)}$  denote the  $i$ th order statistic, i.e. the  $i$ th smallest value in  $\mathbf{x}$ . Moreover, for convenience, let  $x_{\{i\}} = x_{|\mathbf{x}|-i+1}$  denote the  $i$ th greatest value in  $\mathbf{x}$ .

For any  $n \in \mathbb{N}$  and  $c \in \mathbb{I}$ , we set  $(n * c) = (c, \dots, c) \in \mathbb{I}^n$ . Also,  $[n] := \{1, 2, \dots, n\}$  with  $[0] = \emptyset$ .

### 2.1 A Note on Representing Numeric Data and Applying Operations in R

Recall how we create numeric vectors in R:

```
(x1 <- c(5, 2, 3, 1, 0, 0))
## [1] 5 2 3 1 0 0
class(x1)
## [1] "numeric"
(x2 <- 10:1) # the same as seq(10, 1)
## [1] 10 9 8 7 6 5 4 3 2 1
(x3 <- seq(1, 5, length.out=6))
## [1] 1.0 1.8 2.6 3.4 4.2 5.0
(x4 <- seq(1, 5, by=1.25))
## [1] 1.00 2.25 3.50 4.75
```

To obtain  $(n * c)$ , e.g. for  $n = 10$  and  $c = 3$ , we call:

```
rep(10, 3)
## [1] 10 10 10
```

Note that in R all the arithmetic operations on vectors are performed element-wise, i.e. in a manner indicated above. This is called **vectorization**. The same holds for mathematical functions: they are extended to the vector space.

```
x <- c(1, 3, 3, 2)
y <- c(2, 3, -1, 0)
x+y
## [1] 3 6 2 2
x*y
## [1] 2 9 -3 0
pmin(x,y) # parallel minimum
```

```
## [1] 1 3 -1 0
pmax(x,y) # parallel maximum
## [1] 2 3 3 2
abs(y)
## [1] 2 3 1 0
```

Thus, we calculated  $\mathbf{x} + \mathbf{y}$ ,  $\mathbf{x} \cdot \mathbf{y}$ ,  $\mathbf{x} \wedge \mathbf{y}$ ,  $\mathbf{x} \vee \mathbf{y}$ , and  $|\mathbf{x}|$  (try to determine yourself what happens if we deal with two vectors of unequal length in R).

Moreover, given two equal-length vectors, for the  $\leq_n$  relation we write:

```
all(x <= y)
## [1] FALSE
```

To get  $x_{\{i\}}$  we have to sort the given vector nonincreasingly:

```
(xg <- sort(x, decreasing=TRUE)) # `decreasing' may be misleading
## [1] 3 3 2 1
xg[3] # the third greatest value in x
## [1] 2
```

and for  $x_{(i)}$  we type:

```
(xs <- sort(x)) # sorted nondecreasingly
## [1] 1 2 3 3
xs[3] # the third smallest value in x
## [1] 3
```

## 2.2 A Note on Storing Multiple Numeric Vectors in R

Vectors of the same length can be conveniently stored in a matrices. Keep in mind that elements are stored in a columnwise order, so for performance reasons please do store each vector in a separate matrix's column (not: row). Please note that the `dimnames` attribute of a matrix may define its row and column labels. Its value may be set to `NULL` (no names given) or to a list with two character vectors (rows and columns names, respectively). Another simple way to set the labels is by using the `rownames()` and `colnames()` functions.

The `apply()` function may be called to evaluate a given method on each matrix row or column (parameter `MARGIN` set to 1 and 2, respectively).

```
expertopinions <- matrix(c(
  6,7,2,3,1, # this will be the first COLUMN
  8,3,2,1,9, # 2nd
  4,2,4,1,6 # 3rd
),
ncol=3,
dimnames=list(NULL, c("A", "B", "C"))) # only column names set
class(expertopinions)
## [1] "matrix"
```

```
print(expertopinions)  # or print(authors)

##      A B C
## [1,] 6 8 4
## [2,] 7 3 2
## [3,] 2 2 4
## [4,] 3 1 1
## [5,] 1 9 6

apply(expertopinions, 2, mean) # apply the mean() function on each COLUMN

##      A      B      C
## 3.8 4.6 3.4
```

Vectors that are not of the same length may be store in a list (with possibly named elements). In that case, the functionality of `apply()` is provided by `lapply()` or `sapply()` functions.

```
authors <- list(
  "John S." = c(7,6,2,1,0),
  "Kate F." = c(9,8,7,6,4,1,1,0)
)
class(authors)

## [1] "list"

str(authors)  # or print(authors)

## List of 2
## $ John S.: num [1:5] 7 6 2 1 0
## $ Kate F.: num [1:8] 9 8 7 6 4 1 1 0

index_h(authors[[1]]) # the h-index /see below/ for 1st author
## [1] 2

sapply(authors, index_h) # calculate the h-index for all vectors in a list

## John S. Kate F.
##      2      4
```

## 2.3 Aggregation Operators and Their Basic Properties

Dealing with huge amounts of data faces us with the problem of constructing their synthetic descriptions. The aggregation theory, a relatively new research domain at the border of mathematics and computer science, is interested in the analysis of functions that may be used in this task. Thus, we should start with the formal definition of objects of our interest. Here is the most general setting:

**Definition 1.** A function  $F : \mathbb{I}^{1,2,\dots} \rightarrow \mathbb{I}$  is called an **(extended<sup>3</sup>) aggregation operator** if it is at least **nondecreasing** in each variable, i.e. for all  $n$  and  $\mathbf{x}, \mathbf{y} \in \mathbb{I}^n$  if  $\mathbf{x} \leq_n \mathbf{y}$ , then  $F(\mathbf{x}) \leq F(\mathbf{y})$ .

Note that each aggregation operator is a mapping into  $\mathbb{I}$ , thus for all  $n$  we have  $\inf_{\mathbf{x} \in \mathbb{I}^n} F(\mathbf{x}) \geq a$  and  $\sup_{\mathbf{x} \in \mathbb{I}^n} F(\mathbf{x}) \leq b$ . By nondecreasingness, however, these conditions reduce to  $F(n * a) \geq a$  and  $F(n * b) \leq b$ .

---

<sup>3</sup>Extended to the space of vectors of arbitrary length, cf. e.g. [6, 30]; Classical approach considers only fixed-length vectors. In `agop` we are as much general as possible.

Also keep in mind that some authors assume (cf. [30]) that aggregation operators must fulfill the two following **strong boundary conditions**: for all  $n$  we have  $\inf_{\mathbf{x} \in \mathbb{I}^n} F(\mathbf{x}) = a$  and  $\sup_{\mathbf{x} \in \mathbb{I}^n} F(\mathbf{x}) = b$ . Such aggregation operators are sometimes called **averaging functions**. In our case, this does not necessarily hold – we want to be more general.

Here are some interesting properties of averaging functions. Later on we will characterize some of the classes of functions that fulfill them.

**Definition 2.** We call  $F : \mathbb{I}^{1,2,\dots} \rightarrow \mathbb{I}$  **symmetric** if:

$$(\forall n \in \mathbb{N}) (\forall \mathbf{x}, \mathbf{y} \in \mathbb{I}^n) \mathbf{x} \cong \mathbf{y} \implies F(\mathbf{x}) = F(\mathbf{y}),$$

where  $\mathbf{x} \cong \mathbf{y}$  if and only if there exists a permutation  $\sigma$  of  $[n]$  such that  $\mathbf{x} = (y_{\sigma(1)}, \dots, y_{\sigma(n)})$ .

It may be shown, see [30, Thm. 2.34], that  $F : \mathbb{I}^n \rightarrow \mathbb{I}$  is symmetric if and only if there exists a function  $G' : \mathbb{I}^{1,2,\dots} \rightarrow \mathbb{I}$  such that  $F(x_1, \dots, x_n) = G'(x_{(1)}, \dots, x_{(n)})$ , or, equivalently, a function  $G'' : \mathbb{I}^{1,2,\dots} \rightarrow \mathbb{I}$ , for which we have  $F(x_1, \dots, x_n) = G''(x_{\{1\}}, \dots, x_{\{n\}})$ . In other words,  $F$  may be defined solely using order statistics: its value is independent of the aggregated vector’s elements presentation.

By the way:

```
x <- c(0.5, 0.4, 0.1, 0.3, 0.2) # an exemplary vector
sigma1 <- c(1, 3, 5, 2, 4) # an exemplary permutation
x[sigma1]
## [1] 0.5 0.1 0.2 0.4 0.3

(sigma2 <- order(x)) # ordering permutation of x
## [1] 3 5 4 2 1

x[sigma2]
## [1] 0.1 0.2 0.3 0.4 0.5
```

Idempotence is well-known from algebra, where we say that element  $x$  is idempotent with respect to binary operator  $*$  if we have  $x * x = x$ . The following definition extends this property to  $n$ -ary aggregation functions, cf. [30].

**Definition 3.** We call  $F : \mathbb{I}^{1,2,\dots} \rightarrow \mathbb{I}$  **idempotent** if:

$$(\forall n \in \mathbb{N}) (\forall x \in \mathbb{I}) F(n * x) = x.$$

Idempotent aggregation operators fulfilling the strong boundary conditions (see p. 6) are sometimes called **averaging functions**, cf. [30].

An example of such object is the arithmetic mean or median.

**Definition 4.** We call  $F : \mathbb{I}^{1,2,\dots} \rightarrow \mathbb{I}$  **additive** if:

$$F(\mathbf{x} + \mathbf{y}) = F(\mathbf{x}) + F(\mathbf{y}),$$

for all  $n \in \mathbb{N}, \mathbf{x}, \mathbf{y} \in \mathbb{I}^n$  such that  $\mathbf{x} + \mathbf{y} \in \mathbb{I}^n$ .

Please note that for  $a \leq 0$ , if  $F$  is additive, then necessarily it holds  $F(\mathbf{0}) = 0$ .

**Definition 5.** We call  $F$  **minitive** if:

$$(\forall n \in \mathbb{N}) (\forall \mathbf{x}, \mathbf{y} \in \mathbb{I}^n) F(\mathbf{x} \wedge \mathbf{y}) = F(\mathbf{x}) \wedge F(\mathbf{y}).$$

**Definition 6.** We call  $F$  *maxitive* if:

$$(\forall n \in \mathbb{N}) (\forall \mathbf{x}, \mathbf{y} \in \mathbb{I}^n) F(\mathbf{x} \vee \mathbf{y}) = F(\mathbf{x}) \vee F(\mathbf{y}).$$

**Definition 7.** We call  $F$  *modular* (cf. [5, 30, 36]) if:

$$(\forall n \in \mathbb{N}) (\forall \mathbf{x}, \mathbf{y} \in \mathbb{I}^n) F(\mathbf{x} \vee \mathbf{y}) + F(\mathbf{x} \wedge \mathbf{y}) = F(\mathbf{x}) + F(\mathbf{y})$$

It may easily be seen that each additive operator is also modular (i.e. modularity is more general than additivity), because for any additive aggregation operator  $F$ , since  $(\mathbf{x} \vee \mathbf{y}) + (\mathbf{x} \wedge \mathbf{y}) = \mathbf{x} + \mathbf{y}$ , we have  $F(\mathbf{x}) + F(\mathbf{y}) = F(\mathbf{x} + \mathbf{y}) = F((\mathbf{x} \vee \mathbf{y}) + (\mathbf{x} \wedge \mathbf{y})) = F(\mathbf{x} \vee \mathbf{y}) + F(\mathbf{x} \wedge \mathbf{y})$ .

Apart from the “ordinary” minitivity, maxitivity, and modularity we may introduce their symmetrized versions, using  $\mathbf{x} \overset{S}{+} \mathbf{y} = (x_{(1)} + y_{(1)}, \dots, x_{(n)} + y_{(n)})$ ,  $\mathbf{x} \overset{S}{\vee} \mathbf{y} = (x_{(1)} \vee y_{(1)}, \dots, x_{(n)} \vee y_{(n)})$  and  $\mathbf{x} \overset{S}{\wedge} \mathbf{y} = (x_{(1)} \wedge y_{(1)}, \dots, x_{(n)} \wedge y_{(n)})$ .

## 2.4 Impact Functions and The Producers Assessment Problem

We already noticed the important class of aggregation operators: the averaging functions. They may be used to represent the most “typical” value of a numeric vector. Here is another interesting class that represents solutions to some very interesting practical issue.

The **Producers Assessment Problem** (PAP, [28]) concerns evaluation of a set of **producers** (e.g. scientists, artists, writers, craftsman) according to some quality or popularity **ratings** of **products** (e.g. scientific articles, works, books, artifacts) that were outputted by an entity.

**Tab. 1.** The Producer Assessment Problem – typical instances

	Producer	Products	Rating method	Discipline
A	Scientist	Scientific articles	Number of citations	Scientometrics
B	Scientific institute	Scientists	The $h$ -index	Scientometrics
C	Web server	Web pages	Number of in-links	Webometrics
D	Artist	Paintings	Auction price	Auctions
E	Billboard company	Advertisements	Sale results	Marketing
F	R package author	Packages	PageRank values w.r.t. the dependency graph	Software Engineering

PAP instances may be found in many real-life situations, like those encountered for example in scientometrics, webometrics, marketing, manufacturing, or quality engineering, see Table 1 and e.g. [16]. Our main interest here is focused on constructing and analyzing aggregation operators which may be used in the producers’ rating task. Such functions should take into account the two following aspects of a producer’s quality:

- his/her ability to output highly-rated products,
- his/her overall productivity.

For the sake of illustration, we will consider PAP in the scientometric context, where scientists “produce” papers that are cited by peers.

Let  $\mathbb{I} = [0, \infty]$  represent the set of values that some a priori chosen paper quality measure may take. These may of course be non-integers, for example when we consider citations normalized with respect to the number of papers’ authors.

It is widely accepted, see e.g. [49, 48, 50, 41, 39, 40, 28, 18, 17], that each aggregation operator  $F : \mathbb{I}^{1,2,\dots} \rightarrow \mathbb{I}$  to be applied in PAP should at least be:

- (a) nondecreasing in each variable (additional citations received by a paper or an improvement of its quality measure does not result in a decrease of the authors’ overall evaluation),
- (b) arity-monotonic (by publishing a new paper we never decrease the overall valuation of the entity),
- (c) symmetric (independent of the order of elements’ presentation, i.e. we may always assume that we aggregate vectors that are already sorted).

More formally, axiom (b) is fulfilled iff for any  $\mathbf{x} \in \mathbb{I}^{1,2,\dots}$  and  $y \in \mathbb{I}$  it holds  $F(\mathbf{x}) \leq F(x_1, \dots, x_n, y)$ . It may be seen that this property is **arity-dependent**, i.e. it takes into account the number of elements to be aggregated.

Moreover, (a) and (c) were defined in the previous section.

Here is a bunch of arity-dependent properties that can be useful while aggregating vectors of varying lengths, cf. also [9].

**Definition 8.** We call  $F \in \mathcal{E}(\mathbb{I})$  a **zero-insensitive** aggregation operator if for each  $\mathbf{x} \in \mathbb{I}^{1,2,\dots}$  it holds  $F(\mathbf{x}, 0) = F(\mathbf{x})$ .

It may be seen that, under nondecreasingness, zero-insensitivity implies arity-monotonicity, see [26]. What is interesting, each zero-insensitive impact function  $F$  may be defined by means of  $G : \mathbb{I}^\infty \rightarrow \mathbb{I}$  such that  $F(\mathbf{x}) = G(\mathbf{x}, 0, 0, \dots)$ , i.e. of function which domain is the space of vectors of infinite length.

Zero-sensitivity may be strengthened as follows, cf. [26] and [49, Axiom A1].

**Definition 9.**  $F \in \mathcal{E}(\mathbb{I})$  is **F-insensitive** if

$$(\forall \mathbf{x} \in \mathbb{I}^{1,2,\dots}) (\forall y \in \mathbb{I}) y \leq F(\mathbf{x}) \implies F(\mathbf{x}, y) = F(\mathbf{x}).$$

Note that the above property was called R-stability in [4].

**Definition 10.**  $F \in \mathcal{E}(\mathbb{I})$  is **F+sensitive** if

$$(\forall \mathbf{x} \in \mathbb{I}^{1,2,\dots}) (\forall y \in \mathbb{I}) y > F(\mathbf{x}) \implies F(\mathbf{x}, y) > F(\mathbf{x}).$$

## 2.5 Fuzzy Logic Connectives

Another set of tools in which the theory of aggregation is interested consist of fuzzy logic connectives, cf. e.g. [33, 3]. Most of them are binary operations and assume that  $\mathbb{I} = [0, 1]$ .

**Definition 11.** A function  $T : [0, 1] \times [0, 1] \rightarrow [0, 1]$  is a **t-norm** if for all  $x, y, z \in [0, 1]$  it holds:

1.  $T(x, y) = T(y, x)$  (symmetry/commutativity),
2. if  $y \leq z$ , then  $T(x, y) \leq T(x, z)$  (nondecreasingness),
3.  $T(x, T(y, z)) = T(T(x, y), z)$  (associativity),
4.  $T(x, 1) = x$  (neutral element).

Thus, a t-norm is a special kind of symmetric averaging function on  $[0, 1]^2$ . Moreover, each t-norm has 0 as its annihilator element, i.e.  $T(x, 0) = T(0, x) = 0$  for all  $x$ . It is easily seen that the restriction of any t-norm to  $\{0, 1\}$  gives us the conjunction operation known from the classical logic.



**Definition 12.** A function  $S : [0, 1] \times [0, 1] \rightarrow [0, 1]$  is a **t-conorm** if for all  $x, y, z \in [0, 1]$  it holds:

1.  $S(x, y) = S(y, x)$  (symmetry/commutativity),
2. if  $y \leq z$ , then  $S(x, y) \leq S(x, z)$  (nondecreasingness),
3.  $S(x, S(y, z)) = S(S(x, y), z)$  (associativity),
4.  $S(x, 0) = x$  (neutral element).

It is easily seen that the restriction of any t-conorm to  $\{0, 1\}$  gives us the classical logical alternative.

**Definition 13.** A function  $N : [0, 1] \rightarrow [0, 1]$  is a **fuzzy negation** if for all  $x, y \in [0, 1]$  it holds:

1. if  $x \leq y$ , then  $N(x) \geq N(y)$  (nonincreasingness),
2.  $N(0) = 1$ ,
3.  $N(1) = 0$ .

**Definition 14.** A function  $I : [0, 1] \times [0, 1] \rightarrow [0, 1]$  is a **fuzzy implication** if for all  $x, y, x', y' \in [0, 1]$  it holds:

1. if  $x \leq x'$ , then  $I(x, y) \geq I(x', y)$  (nonincreasingness w.r.t.  $x$ ),
2. if  $y \leq y'$ , then  $I(x, y) \leq I(x, y')$  (nondecreasingness w.r.t.  $y$ ),
3.  $I(1, 1) = 1$ ,
4.  $I(0, 0) = 1$ ,
5.  $I(1, 0) = 0$ .

It is easily seen that  $I(x, 1) = 1$  and  $I(0, y) = 1$  for all  $x, y$ .

Note that fuzzy negations and implications are not averaging functions in the above-mentioned sense. It is because they do not fulfill the nondecreasingness condition.

## 2.6 Copulas

Copulas form another group of interesting and useful aggregation operators. They may be used in probability and statistics to model the kind of dependency between random variables, see e.g. [38].

For given  $n$ , each  $n$ -copula  $C : [0, 1]^n \rightarrow [0, 1]$  is a cumulative distribution function of a  $n$ -dimensional random variable having uniform margins. In particular, for  $n = 2$  we have what follows.

**Definition 15.** A function  $C : [0, 1] \times [0, 1] \rightarrow [0, 1]$  is a **2-copula** if for all  $x, y, x', y' \in [0, 1]$  it holds:

1. if  $x \leq x'$  and  $y \leq y'$ , then  $C(x, y) + C(x', y') - C(x, y') - C(x', y) \geq 0$  (2-increasingness),
2.  $C(x, 0) = C(0, x) = 0$  (annihilator element),
3.  $C(x, 1) = x$  (neutral element).

Note that each t-norm fulfills (2) and (3). Moreover, each copula is nondecreasing. However, there are 2-copulas that are not t-norms and conversely, see e.g. [32].

## 2.7 Spread Measures

Classically, in the theory of aggregation we consider the broadly-conceived averaging functions and logical connectives. However, one often needs a very different kind of a proper synthesis of multidimensional numeric data into a single number. In [23] an axiomatization of **spread measures**, which may be used to measure data variability, spread, or scatter, was proposed.

.....TO BE DONE.....

Given  $\mathbf{x}, \mathbf{x}' \in \mathbb{R}^n$ , we write  $\mathbf{x} \preccurlyeq \mathbf{x}'$  and say that  $\mathbf{x}$  **has not greater spread than**  $\mathbf{x}'$ , if and only if for all  $i, j = 1, \dots, n$  it holds

$$x_i \geq x_j \implies x_i - x_j \leq x'_i - x'_j. \quad (1)$$

It is easily seen that  $\preccurlyeq$  is a preorder on  $\mathbb{R}^n$ , i.e. a relation that is reflexive and transitive, and that it is not necessarily total, i.e. not all vectors are comparable with each other, see also Sec. 4.3.

**Definition 16.** A **spread measure** is an aggregation operator  $V : \mathbb{R}^n \rightarrow [0, \infty]$  such that:

(d1) For any  $c \in \mathbb{R}$  it holds  $V(n * c) = V(c, c, \dots, c) = 0$ ,

(d2) for each  $\mathbf{x} \preccurlyeq \mathbf{x}'$  it holds  $V(\mathbf{x}) \leq V(\mathbf{x}')$ .

This class includes e.g. the sample variance, standard deviation, range, interquartile range, median absolute deviation etc., that is functions widely used in exploratory data analysis.

## 3 Visualization

### 3.1 Depicting Producers

The `plot_producer()` function may be used to draw a graphical representation of a given numeric vector, i.e. what is sometimes called a **citation function** in scientometrics.

As in the PAP we are interested in symmetric agops, a given vector  $\mathbf{x} = (x_1, \dots, x_n)$  may be represented by a step function defined for  $0 \leq y < n$  and given by:

$$\pi(y) = x_{\lfloor y+1 \rfloor}.$$

This function may be obtained by setting `type='right.continuous'` argument in `plot_producer()`. Recall that  $x_{\{i\}}$  denotes  $i$ -th greatest value in  $\mathbf{x}$ .

On the other hand, for `type='left.continuous'` (the default), we get

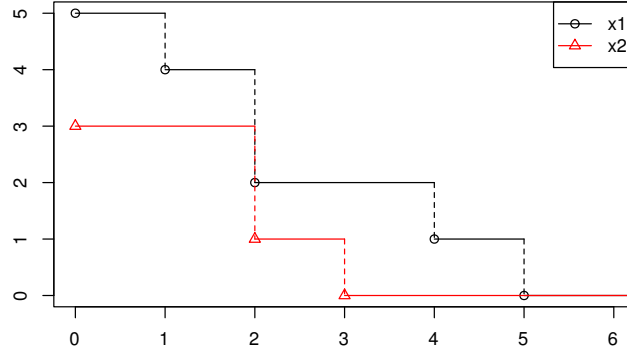
$$\pi(y) = x_{\lfloor y \rfloor}$$

for  $0 < y \leq n$ .

Note that this function depicts the curve that joins the sequence of points  $(0, x_{\{1\}}), (1, x_{\{1\}}), (1, x_{\{2\}}), (2, x_{\{2\}}), \dots, (n, x_{\{n\}})$ .

The `plot_producer()` function behaves much like the well-known R's `plot.default()` and allows for passing all its graphical parameters. For example, let us depict the state of two given producers,  $\mathbf{x}^{(1)}$  and  $\mathbf{x}^{(2)}$ .

```
x1 <- c(5, 4, 2, 2, 1)
x2 <- c(3, 3, 1, 0, 0, 0, 0)
plot_producer(x1, extend=TRUE)
plot_producer(x2, add=TRUE, col=2, pch=2, extend=TRUE)
legend('topright', c('x1', 'x2'), col=c(1, 2), lty=1, pch=c(1, 2))
```



## 4 Binary Relations

The `agop` package includes a few functions aiming to deal with binary relations<sup>4</sup> defined on finite sets consisting of distinctive elements  $V = \{v_1, \dots, v_n\}$ . Each binary relation  $R \subseteq V \times V$  may be represented by a square 0-1<sup>5</sup> matrix  $A = (a_{i,j})_{i,j=1,\dots,n}$  such that  $a_{i,j} = 1$  if and only if  $v_i R v_j$ .

Note that we write  $R \subseteq R'$  if  $v_i R v_j \implies v_i R' v_j$ , and that if  $R = V \times V$  then for all  $i, j$  it holds  $v_i R v_j$ .

Table 2 contains an overview of the most often considered properties of binary relations. Moreover, Table 3 lists popular classes of relations.

**Tab. 2.** Properties of binary relations.

Property	Definition	agop implementation
Reflexivity	$(\forall i) v_i R v_i$	<code>rel_is_reflexive()</code>
Irreflexivity	$(\forall i) \neg v_i R v_i$	<code>rel_is_irreflexive()</code>
Symmetry	$(\forall i, j) v_i R v_j \Rightarrow v_j R v_i$	<code>rel_is_symmetric()</code>
Antisymmetry	$(\forall i, j) v_i R v_j \text{ and } v_j R v_i \Rightarrow i = j$	<code>rel_is_antisymmetric()</code>
Asymmetry	$(\forall i, j) v_i R v_j \Rightarrow \neg v_j R v_i$	<code>rel_is_asymmetric()</code>
Totality	$(\forall i, j) v_i R v_j \text{ or } v_j R v_i$	<code>rel_is_total()</code>
Transitivity	$(\forall i, j, k) v_i R v_j \text{ and } v_j R v_k \Rightarrow v_i R v_k$	<code>rel_is_transitive()</code>
Cyclicity	transitive closure of $R$ is not antisymmetric	<code>rel_is_cyclic()</code>

**Reductions and closures** To determine the **reflexive closure**, i.e. the minimal reflexive  $R' \supseteq R$  call `rel_closure_reflexive()`. On the other hand, with a call to `rel_reduction_reflexive()` we get the **reflexive reduction** of  $R$ , i.e. the minimal  $R' \subseteq R$  such that the reflexive closures of  $R'$  and  $R$  are equal. In other words,  $R'$  is the largest irreflexive relation contained in  $R$ .

<sup>4</sup>On CRAN there is also a package `relations`, see [37], that provides data structures and algorithms for  $k$ -ary relations with arbitrary domains, featuring relational algebra, predicate functions, and fitters for consensus relations. For some time, `agop`’s functionality will be a subset of `relations`’s (yet faster). In future versions, however, we’d like to add fuzzy relations handling.

<sup>5</sup>Or, equivalently, logical; we have `as.logical(0) == FALSE` and `as.logical(x) == TRUE` if  $x \neq 0$ , and, on the other hand, `as.integer(FALSE) == 0` `as.integer(TRUE) == 1`.

**Tab. 3.** Types of binary relations.

Class	Properties
preorder (quasiorder)	reflexive, transitive
total preorder (weak order, preference)	total ( $\Rightarrow$ reflexive), transitive
partial order	reflexive, transitive, antisymmetric
linear order	total ( $\Rightarrow$ reflexive), transitive, antisymmetric
equivalence relation	symmetric, reflexive, transitive

To find the **transitive closure**, cf. [47], of a given binary relation  $R$ , i.e. the minimal transitive  $R' \supseteq R$ , we call *agop*’s `rel_closure_transitive()` function. On the other hand, the **transitive reduction** of acyclic  $R$ , see [1] and the `rel_reduction_transitive()` function, is the minimal  $R' \subseteq R$  such that the transitive closures of  $R$  and  $R'$  are equal.

A mixture of the reflexive reduction and some kind of transitive reduction, particularly useful when drawing Hasse diagrams of preordered sets (which may not necessarily be represented by an acyclic relation  $R$ ) may be determined with `rel_reduction_hasse()`.

In general, a total closure and a total reduction are not well-defined. However, when dealing with preorders, the following notion may be useful, see [21]. To determine the so-called **fair total closure** i.e. minimal total  $R' \supseteq R$  such that if  $\neg xRy$  and  $\neg xR'y$  then  $xR'y$  and  $yR'x$ , we call `rel_closure_total_fair()`.

The **symmetric closure**, the smallest symmetric binary relation that contains a given one, is available via a call to `rel_closure_symmetric()`.

#### 4.1 Weak Dominance Relation (for PAP)

Let us consider the following relation on  $\mathbb{I}^{1,2,\dots}$ . For any  $\mathbf{x} \in \mathbb{I}^n$  and  $\mathbf{y} \in \mathbb{I}^m$  we write  $\mathbf{x} \preceq \mathbf{y}$  if and only if  $n \leq m$  and  $x_{\{i\}} \leq y_{\{i\}}$  for all  $i = 1, \dots, n \wedge m$ .

Of course,  $\preceq$  is symmetric and transitive, i.e. it is a preorder. Moreover, it would have been a partial order (in general it’s not), if we had defined it on the set of *sorted* vectors.

Intuitively, we say that an author (scientometric context again)  $X$  is (weakly) dominated by an author  $Y$ , if  $X$  has no more papers than  $Y$  and each the  $i$ th most cited paper of  $X$  has no more citations than the  $i$ th most cited paper of  $Y$ . Not that the  $(m - n)$  least cited  $Y$ ’s papers are not taken into account here.

Most importantly, however, there exist pairs of vectors that are *incomparable* with respect to  $\preceq$  (see the illustration below). In other words, this dominance relation is not total.

Whether this relationship between a pair of vectors holds may be determined using *agop*’s `pord_weakdom()` function.

```
c(pord_weakdom(5:1, 10:1), pord_weakdom(10:1, 5:1)) # 5:1 <= 10:1
## [1] TRUE FALSE

c(pord_weakdom(3:1, 5:4), pord_weakdom(5:4, 3:1)) # 3:1 ?? 5:4
## [1] FALSE FALSE
```

The following result was shown in [28]. Let  $F \in \mathcal{E}(\mathbb{I})$ . Then  $F$  is symmetric, nondecreasing in each variable and arity-monotonic if and only if for any  $\mathbf{x}, \mathbf{y}$  if  $\mathbf{x} \preceq \mathbf{y}$ , then  $F(\mathbf{x}) \leq F(\mathbf{y})$ . Therefore, the class of impact functions may be equivalently defined as all the aggregation operators that are nondecreasing with respect to this preorder.

Additionally, we will write  $\mathbf{x} \triangleleft \mathbf{y}$  if  $\mathbf{x} \leq \mathbf{y}$  and  $\mathbf{x} \neq \mathbf{y}$  (strict dominance).

## 4.2 Weak Dominance Relation (for fixed arity)

Consider the following relation on  $\mathbb{I}^{1,2,\dots}$ . For any  $\mathbf{x} \in \mathbb{I}^n$  and  $\mathbf{y} \in \mathbb{I}^m$  we write  $\mathbf{x} \leq \mathbf{y}$  if and only if  $n = m$  and  $x_i \leq y_i$  for all  $i = 1, \dots, n$ .

The `all(x <= y)` expression may not work well for vectors of unequal lengths.

...TO DO...

Moreover, a version that is independent of the vectors’ elements ordering (i.e. defined using  $x_{(i)} \leq y_{(i)}$ ) is available:

...TO DO...

## 4.3 Vector Spread

The function `pord_spread()` can be used to compare spread, scater, or variability of two numeric vectors, see [23] with the  $\preceq$  relation introduced in Sec. 2.7 and [23].

Moreover, for `pord_spreadsym()` we get the relation which acts on sorted input vectors (which leads to symmetric dispersion functions).

```
pord_spreadsym(c(1,5,2), c(1,7,3))
```

```
## [1] TRUE
```

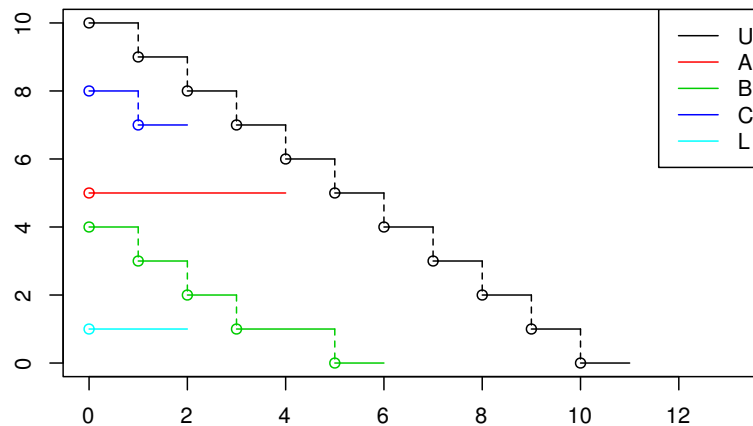
## 4.4 Operations on Preorders and Other Binary Relations

**Example.** Let us consider the 5 following vectors.

```
ex1 <- list(
  U = 10:0,           # some upper bound
  A = c(5,5,5,5),     # moderate productivity & quality
  B = c(4,3,2,1,1,0), # high productivity
  C = c(8,7),         # high quality
  L = c(1,1)          # some lower bound
)
```

Plot of “citation” curves:

```
for (i in seq_along(ex1))
  plot_producer(ex1[[i]], add=(i>1), col=i)
legend("topright", legend=names(ex1), col=1:length(ex1), lty=1)
```



Here is the adjacency matrix for the preordered set  $(\{A, B, C, L, U\}, \preceq)$  created with the `rel_graph()` function, that takes each pair of elements from the list passed as its first argument and compares them with a given function passed as its second argument.

```
ord <- rel_graph(ex1, pord_weakdom) # compare each (ex1[[i]], ex1[[j]]) with pord_weakdom
print(ord)

##      U      A      B      C      L
## U TRUE FALSE FALSE FALSE FALSE
## A TRUE  TRUE FALSE FALSE FALSE
## B TRUE FALSE  TRUE FALSE FALSE
## C TRUE FALSE FALSE  TRUE FALSE
## L TRUE  TRUE  TRUE  TRUE  TRUE

rel_is_reflexive(ord) # is reflexive
## [1] TRUE

rel_is_transitive(ord) # is transitive
## [1] TRUE

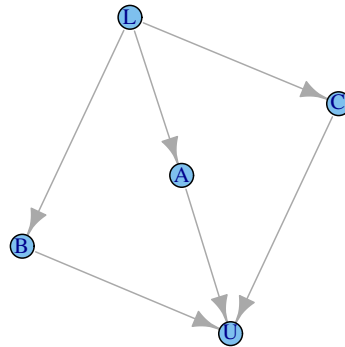
rel_is_total(ord)      # not a total preorder...
## [1] FALSE
```

We see that we have  $A \not\preceq B$ ,  $A \not\preceq C$ ,  $B \not\preceq C$  (no pair from  $\{A, B, C\}$  is comparable w.r.t.  $\preceq$ ):  
 ..TO DO..

```
#incomp <- get_incomparable_pairs(ord)
#incomp <- incomp[incomp[,1]<incomp[,2],] # remove permutations: ((1,2), (2,1)) -> (1,2)
#incomp[,] <- rownames(ord)[incomp]
#print(incomp) # all incomparable pairs
## the other way: generate maximal independent sets
#lapply(get_independent_sets(ord), function(set) rownames(ord)[set])
```

To draw the Hasse diagram, we base on the reflexive and a kind of transitive reduction of the graph, which is determined with `rel_reduction_hasse()`.

```
require(igraph)
hasse <- graph.adjacency(rel_reduction_hasse(ord))
set.seed(1234567) # igraph's drawing facilities are far from perfect
plot(hasse, layout=layout_fruchterman_reingold(hasse, dim=2))
```

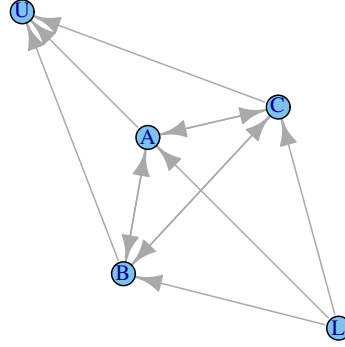


$(\{A, B, C, L, U\}, \preceq)$  is not totally ordered, let us determine the **fair total closure** of  $\preceq$  (set  $x \preceq'' y$  and  $y \preceq'' x$  whenever  $\neg(x \preceq y \text{ or } y \preceq x)$ , see [21] for discussion), and then calculate its transitive closure, as the resulting matrix may not necessarily be transitive.

```
ord_total <- rel_closure_transitive(rel_closure_total_fair(ord)) # a total preorder
print(ord_total)

##      U      A      B      C      L
## U TRUE FALSE FALSE FALSE FALSE
## A TRUE  TRUE  TRUE  TRUE  FALSE
## B TRUE  TRUE  TRUE  TRUE  FALSE
## C TRUE  TRUE  TRUE  TRUE  FALSE
## L TRUE  TRUE  TRUE  TRUE  TRUE

hasse <- graph.adjacency(rel_reduction_hasse(ord_total))
set.seed(123)
plot(hasse, layout=layout_fruchterman_reingold(hasse, dim=2))
```



Note that each total preorder  $\leq''$  induces an equivalence relation ( $x \simeq y$  iff  $x \leq'' y$  and  $y \leq'' x$ ; the equivalence classes may be ordered with  $\leq''$ ).

## 5 Predefined Classes of Aggregation Operators in *agop*

### 5.1 A Review of Notable Classes of Aggregation Operators

Here are some well-known classes of aggregation operators. Originally, they were defined for fixed-length vector and for  $\mathbb{I} = [0, 1]$ .

**Definition 17.** Let  $\mathbf{w} = (w_1, \dots, w_n) \in [0, 1]^n$  be a **weighting vector** such that  $\sum_{i=1}^n w_i = 1$ . Then, for any  $\mathbf{x} \in \mathbb{I}^n$ :

1. The **weighted arithmetic mean** associated with  $\mathbf{w}$ ,  $\text{WAM}_{\mathbf{w}} : \mathbb{I}^n \rightarrow \mathbb{I}$ , is defined as

$$\text{WAM}_{\mathbf{w}}(\mathbf{x}) = \sum_{i=1}^n w_i x_i.$$

2. The **ordered weighted averaging operator** (cf. [51]) associated with  $\mathbf{w}$ ,  $\text{OWA}_{\mathbf{w}} : \mathbb{I}^n \rightarrow \mathbb{I}$ , is defined as

$$\text{OWA}_{\mathbf{w}}(\mathbf{x}) = \sum_{i=1}^n w_i x_{(i)}.$$

We see that both functions are idempotent, additive, and that OWA is the symmetrized version of WAM. Moreover, for  $\mathbf{w} = (n * \frac{1}{n})$ ,  $\text{WAM}_{\mathbf{w}}$  defines the arithmetic mean (`mean()` in R). **Truncated mean** is an interesting example of an OWA operator (see `mean(x, trim=...)`).

In *agop* the WAM and OWA operators are available as `wam()` and `owa()`.

```
wam(c(1,2,2,2), c(0.1,0.4,0.4,0.1))
## [1] 1.9
owa(c(1,3,5,2), rep(1,4)) # should be normalized
## Warning: elements of 'w' does not sum up to 1. correcting.
## [1] 2.75
```



Note that there is a strong, well-known connection between the OWA operators and the Choquet integral [10] w.r.t. some monotone measure, see e.g. [30].

**Definition 18.** Let  $\mathbf{w} = (w_1, \dots, w_n) \in \mathbb{I}^n$  be a vector such that  $\bigvee_{i=1}^n w_i = b = \sup \mathbb{I}$ . Then, for any  $\mathbf{x} \in \mathbb{I}^n$ :

1. The **weighted maximum** associated with  $\mathbf{w}$ ,  $\text{WMax}_{\mathbf{w}} : \mathbb{I}^n \rightarrow \mathbb{I}$ , is defined as

$$\text{WMax}_{\mathbf{w}}(\mathbf{x}) = \bigvee_{i=1}^n (w_i \wedge x_i).$$

2. The **ordered weighted maximum** (cf. [14, 13]) associated with  $\mathbf{w}$ ,  $\text{OWMax}_{\mathbf{w}} : \mathbb{I}^n \rightarrow \mathbb{I}$ , is defined as

$$\text{OWMax}_{\mathbf{w}}(\mathbf{x}) = \bigvee_{i=1}^n (w_i \wedge x_{(i)}).$$

*agop* implementation: `wmax()` and `owmax()`.

```
wmax(c(1,3,5,2), rep(Inf, 4)) # greatest value /default behavior/
## [1] 5
owmax(1:10, 1:10)
## [1] 10
```

**Definition 19.** Let  $\mathbf{w} = c(w_1, \dots, w_n) \in [0, 1]^n$  be such that  $\bigwedge_{i=1}^n w_i = a = \inf \mathbb{I}$ . Then, for any  $\mathbf{x} \in \mathbb{I}^n$ :

1. The **weighted minimum**  $\text{WMin}_{\mathbf{w}} : \mathbb{I}^n \rightarrow \mathbb{I}$  associated with the weight vector  $\mathbf{w}$  is defined as

$$\text{WMin}_{\mathbf{w}}(\mathbf{x}) = \bigwedge_{i=1}^n (w_i \vee x_i).$$

2. The **ordered weighted minimum**  $\text{OWMin}_{\mathbf{w}} : \mathbb{I}^n \rightarrow \mathbb{I}$  associated with the weight vector  $\mathbf{w}$  is defined as

$$\text{OWMin}_{\mathbf{w}}(\mathbf{x}) = \bigwedge_{i=1}^n (w_i \vee x_{(i)}).$$

*agop* implementation: `wmin()` and `owmin()`.

It is clear to see that  $\text{OWMax}$  operators fulfill the maxitivity property and  $\text{OWMin}$  operators fulfill the minitivity property. Interestingly, it may be shown, cf. [30], that for each  $\text{OWMax}$  operator there exist an equivalent  $\text{OWMin}$  operator and inversely.

As stated above, “classical” aggregation operators were defined for vectors of fixed lengths. Let us present some notable generalizations of these operators.

Let  $\mathbb{I}^{\mathbb{I}}$  denote the set of functions from  $\mathbb{I}$  to  $\mathbb{I}$ . The following object will be needed for further considerations.

**Definition 20.** A **triangle of functions** is a sequence  $\Delta = (f_{i,n} \in \mathbb{I}^{\mathbb{I}} : i \in [n], n \in \mathbb{N})$ .

Here is a graphical interpretation of  $\Delta$ :

$$\begin{array}{cccc} f_{1,1} & & & \\ f_{1,2} & f_{2,2} & & \\ f_{1,3} & f_{2,3} & f_{3,3} & \\ \vdots & \vdots & \vdots & \ddots \end{array}$$

**Definition 21.** Let  $\Delta = (f_{i,n})_{i \in [n], n \in \mathbb{N}}$  be a triangle of functions such that  $(\forall n) \sum_{i=1}^n \inf f_{i,n} \geq a$  and  $(\forall n) \sum_{i=1}^n \sup f_{i,n} \leq b$ . Then the **quasi-L-statistic** generated by  $\Delta$  is a function  $\mathbf{qL}_\Delta : \mathbb{I}^{1,2,\dots} \rightarrow \mathbb{I}$  such that

$$\mathbf{qL}_\Delta(\mathbf{x}) = \sum_{i=1}^n f_{i,n}(x_{\{i\}}).$$

It is easily seen that quasi-L-statistics generalize OWA operators if we set  $f_{i,n}(x) = c_{n-i+1,n}x$ ,  $c_{i,n} \in [0, 1]$ , and  $(\forall n) \sum_{i=1}^n c_{i,n} = 1$ .

Assume that  $\mathbb{I} = [0, b]$ . Interestingly, it has been shown ([36], cf. also [20]) that an aggregation operator  $F : \mathbb{I}^{1,2,\dots} \rightarrow \mathbb{I}$  fulfills the symmetric modularity property if and only if  $F$  is a nondecreasing quasi-L-statistic. What is more, in [20] we may find that  $\mathbf{qL}_\Delta$  is nondecreasing if and only if there exists  $\nabla = (g_{i,n})_{i \in [n], n \in \mathbb{N}}$  such that  $(\forall n) (\forall i \in [n]) g_{i,n}$  is nondecreasing,  $(\forall n) \sum_{i=1}^n g_{i,n} \leq b$ ,  $(\forall n) (\forall i > 1) g_{i,n}(0) = 0$  and  $\mathbf{qL}_\Delta = \mathbf{qL}_\nabla$ .

**Definition 22.** The **quasi-S-statistic** for a given triangle of functions  $\Delta = (f_{i,n})_{i \in [n], n \in \mathbb{N}}$  is a function  $\mathbf{qS}_\Delta : \mathbb{I}^{1,2,\dots} \rightarrow \mathbb{I}$  such that

$$\mathbf{qS}_\Delta(\mathbf{x}) = \bigvee_{i=1}^n f_{i,n}(x_{\{i\}}),$$

for any  $\mathbf{x} \in \mathbb{I}^{1,2,\dots}$ .

Quasi-S-statistic generalize the OWMax operators, if  $f_{i,n}(x) = x \wedge c_{n-i+1,n}$ ,  $c_{i,n} \in \mathbb{I}$  and  $(\forall n) \bigvee_{i=1}^n c_{i,n} = b$ .

There is an equivalence between symmetric maxitive aggregation operators and nondecreasing quasi-S-statistics. Moreover, without loss of generality we may assume that a nondecreasing quasi-S-statistic is always generated by triangle of functions in which  $(\forall n) (\forall i \in [n]) f_{i,n}$  is nondecreasing,  $(\forall n) (\forall i \in [n]) f_{i,n}(a) = f_{n,n}(a)$  and  $(\forall n) f_{1,n} \preceq \dots \preceq f_{n,n}$ , see [20].

**Definition 23.** The **quasi-I-statistic** generated by  $\Delta = (f_{i,n})_{i \in [n], n \in \mathbb{N}}$  is a function  $\mathbf{qI}_\Delta : \mathbb{I}^{1,2,\dots} \rightarrow \mathbb{I}$  such that

$$\mathbf{qI}_\Delta(\mathbf{x}) = \bigwedge_{i=1}^n f_{i,n}(x_{\{i\}}),$$

for any  $\mathbf{x} \in \mathbb{I}^{1,2,\dots}$ .

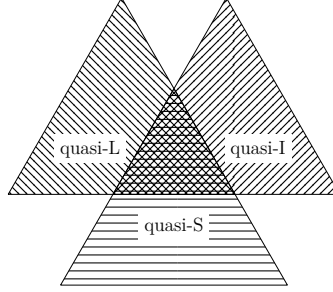
Quasi-I-statistics are generalizations of the OWMin operators, if  $f_{i,n}(x) = x \vee c_{n-i+1,n}$ ,  $c_{i,n} \in \mathbb{I}$  and  $(\forall n) \bigwedge_{i=1}^n c_{i,n} = a$ .

Like above, it has been shown that every symmetric minitive aggregation operator is a nondecreasing quasi-I-statistic, and conversely. Additionally, with no loss in generality we may assume that nondecreasing quasi-S-statistic is generated by triangle of functions in which  $(\forall n) (\forall i \in [n]) f_{i,n}$  is nondecreasing,  $(\forall n) (\forall i \in [n]) f_{i,n}(b) = f_{n,n}(b)$  and  $(\forall n) f_{1,n} \preceq \dots \preceq f_{n,n}$ , see [20].

Note: sometimes we also consider L-, S-, and I-statistics, i.e. special cases of the above-defined quasi--statistics, generated by triangles of coefficients (i.e. sequences  $\Delta = (c_{i,n} \in \mathbb{I} : i \in [n], n \in \mathbb{N})$ , cf. [6]). An **L-statistic** is a quasi-L-statistic for which we have  $f_{i,n}(x) = c_{i,n}x$ . Similarly, by setting  $f_{i,n}(x) = x \wedge c_{i,n}$  we obtain an **S-statistic** from the quasi-S-statistics class, and by setting  $f_{i,n}(x) = x \vee c_{i,n}$  we get an **I-statistic** from quasi-I-statistics.

Also note that L-statistics are known from the probability theory. However, sometimes under this name some authors understand sums of a function of order statistics.

Most interestingly, in [20] it has been shown that the intersection of any two of the three “quasi” classes is the same:



Basing on this result, the **OM3** class (symmetric maxitive, minitive, and also modular aggregation operators) was proposed in [7, 8].

**Definition 24.** A sequence of nondecreasing functions  $\mathbf{w} = (w_1, w_2, \dots)$ ,  $w_i : \mathbb{I} \rightarrow \mathbb{I}$ , and a triangle of coefficients  $\Delta = (c_{i,n})_{i \in [n], n \in \mathbb{N}}$ ,  $c_{i,n} \in \mathbb{I}$  such that  $(\forall n) c_{1,n} \leq c_{2,n} \leq \dots \leq c_{n,n}$ ,  $0 \leq w_n(0) \leq c_{1,n}$ , and  $w_n(b) = c_{n,n}$ , generates a nondecreasing **OM3 operator**  $M_{\Delta, \mathbf{w}} : \mathbb{I}^n \rightarrow \mathbb{I}$  such that for  $\mathbf{x} \in \mathbb{I}^n$  we have:

$$\begin{aligned} M_{\Delta, \mathbf{w}}(\mathbf{x}) &= \bigvee_{i=1}^n w_n(x_{(n-i+1)}) \wedge c_{i,n} = \bigwedge_{i=1}^n (w_n(x_{(n-i+1)}) \vee c_{i-1,n}) \wedge c_{n,n} \\ &= \sum_{i=1}^n \left( (w_n(x_{(n-i+1)}) \vee c_{i-1,n}) \wedge c_{i,n} - c_{i-1,n} \right). \end{aligned}$$

We see that the OM3 class contains i.a. all order statistics (whenever  $w_n(x) = x$ , and  $c_{i,n} = 0$ ,  $c_{j,n} = b$  for  $i < k$ ,  $j \geq k$ , and some  $k$ ), OWM<sub>ax</sub> operators (for  $w_n(x) = x$ ), and the famous Hirsch  $h$ -index (see below).

## 5.2 Interesting Impact Functions

Let us go back to the Producers Assessment Problem. Below we assume that  $\mathbb{I} = [0, \infty]$ .

**The  $h$ -index.** Given a sequence  $\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{I}^{1,2,\dots}$ , the **Hirsch index** [31] of  $\mathbf{x}$  is defined as  $H(\mathbf{x}) = \max\{i = 1, \dots, n : x_{\{i\}} \geq i\}$  if  $n \geq 1$  and  $x_{\{1\}} \geq 1$ , or  $H(\mathbf{x}) = 0$  otherwise. It may be shown that the  $h$ -index is a zero-insensitive OM3 aggeration operator, see [20], with:

$$H(\mathbf{x}) = \bigvee_{i=1, \dots, n} i \wedge \lfloor x_{\{i\}} \rfloor.$$

Interpretation: “an author has  $h$ -index of  $H$  if  $H$  of his/her  $n$  most cited papers have at least  $H$  citations each, and the other  $n - H$  papers are cited no more that  $H$  times each”. The  $h$ -index may also be expressed as a Sugeno integral [43] w.r.t. to a counting measure, cf. [45] and [29].

*agop* implementation: `index_h()`.

```
index_h(c(6,5,4,2,1,0,0,0,0,0))
```

```
## [1] 3
```

Moreover, we have  $H(\mathbf{x}) \leq \min\{n, x_1\}$ .

Note that the  $h$ -index was defined in original context (aggregation of citation counts) for integer vectors. More generally, it is better to use the OM3 operator with  $w_i(x) = x = \text{Id}(x)$  and  $c_{i,n} = i$  (two identity “objects” = one of the simplest setting). Interestingly, such aggregation operator is then asymptotically idempotent, i.e. for all  $x \in \mathbb{I}$  we have  $\lim_{n \rightarrow \infty} M_{\Delta, \mathbf{w}}(n * x) = x$ .

**The  $g$ -index.** Egghe’s  $g$ -index [15] is defined as  $G(\mathbf{x}) = \max\{g = 1, \dots, n : \sum_{i=1}^g x_{\{g\}} \geq g^2\}$ , and is available in *agop* by calling `index_g()`. We have  $G(\mathbf{x}) \geq H(\mathbf{x})$  with  $G(n*n) = H(n*n) = n$ .

Note that this aggregation operator is not zero-insensitive, for example  $G(9, 0) = 2$  and  $G(9, 0, 0) = 3$ . Thus, we also provide the `index_g_zi()` function, which treats  $\mathbf{x}$  as it would be padded with 0s.

```
index_g(9)
## [1] 1
index_g(c(9,0,0))
## [1] 3
index_g_zi(9)
## [1] 3
```

The index is interesting from the computational point of view – it may be calculated on the nondecreasing vector of cumulative sums, `cumsum(sort(x, decreasing=TRUE))`, however, it cannot directly be expressed as a symmetric maxitive aggregation operator.

However, it might be shown (see [29] for the proof) that if  $\mathbf{x}$  is sorted nondecreasingly, then:

$$G(\mathbf{x}) = H(\mathbf{x})(0 \vee \text{cummin}(\text{cumsum}(x) - (1:n)^2 + (1:n))),$$

where  $1:n = (1, 2, 3, \dots, n)$ .

**The  $w$ -index.** The  $w$ -index [49] is defined as

$$W(\mathbf{x}) = \max \left\{ w = 0, 1, 2, \dots : x_{\{i\}} \geq w - i + 1, i = 1, \dots, w \right\}$$

and is available in *agop* by calling `index_w()`.

Interestingly, we have shown in [29] that if  $\mathbf{x}$  is sorted nondecreasingly, then:

$$W(\mathbf{x}) = H(\mathbf{x})(\text{cummin}(\mathbf{x} + (1:n) - 1)).$$

Thus, it is easily seen that this is a zero-insensitive impact function. What is more we have  $H(\mathbf{x}) \leq W(\mathbf{x}) \leq 2H(\mathbf{x})$  and  $W(\mathbf{x}) \leq \min\{n, x_1\}$ .

**The  $r_p$ -indices.** The  $r_p$ -index, for  $p \geq 1$  is expressed as

$$r_p(\mathbf{x}) = \sup \{r > 0 : \mathbf{s}^{p,r} \preceq \mathbf{x}\},$$

where  $\mathbf{s}^{p,r} = (\sqrt[p]{r^p - 0^p}, \sqrt[p]{r^p - 1^p}, \dots, \sqrt[p]{r^p - \lfloor r \rfloor^p})$ . For more details see [18, 25].

Please note that for integer vectors we have  $r_1 = W$  and  $r_\infty = H$  (cf. [25]). Hence it easily seen that, this is a zero-insensitive impact function.

*agop* implementation: `index_rp()`.

**The  $l_p$ -indices.** The  $l_p$ -index (cf. [18, 25]) for  $p \in [1, \infty)$ ,  $u > 0$  and  $v > 0$  is a function  $l_p : \mathbb{I}^{1,2,\dots} \rightarrow \mathbb{I}^2$  given by the equation

$$l_p(\mathbf{x}) = \arg \sup_{(u,v)} \{uv : \mathbf{e}^{p,u,v} \leq \mathbf{x}\},$$

where  $\mathbf{e}^{p,u,v} = \left( \sqrt[p]{v^p - (\frac{v}{u}0)^p}, \sqrt[p]{v^p - (\frac{v}{u}1)^p}, \dots, \sqrt[p]{v^p - (\frac{v}{u}\lfloor u \rfloor)^p} \right)$ .

*agop* implementation: `index_lp()`.

**The MAXPROD-index.** The MAXPROD-index [35] is given by the equation

$$\text{MP}(\mathbf{x}) = \max \left\{ i \cdot x_{\{i\}} : i = 1, 2, \dots \right\}$$

is another example of zero-insensitive impact function. Interestingly, this index is a particular case of a projected  $l_\infty$ -index, see [25], and can be also expressed in terms of Shilkret integral [42], see [29] for discussion.

In *agop* the MAXPROD-index is implemented in the `index_maxprod()` function.

**Simple transformations of the  $h$ -index.** Bibliometricians in many papers considered very simple, direct modifications of the  $h$ -index. For example, the  $h(2)$ -index [34] is defined as:

$$\text{H2}(\mathbf{x}) = \max \left\{ h = 0, 1, 2, \dots : x_h \geq h^2 \right\}.$$

Some authors introduced other settings than “ $h^2$ ” on the right side of (5.2), e.g. “ $2h$ ”, “ $\alpha h$ ” for some  $\alpha > 0$ , or “ $h^\beta$ ”,  $\beta \geq 1$ , cf. [2].

It may easily be shown that these reduce to the  $h$ -index for properly transformed input vectors, e.g.  $\text{H2}(\mathbf{x}) = \text{H}(\sqrt{\mathbf{x}})$ .

### 5.3 Noteworthy Fuzzy Logic Connectives

All the predefined fuzzy logic connectives have been vectorized in *agop*. In other words, any e.g. binary operation  $B : [0, 1]^2 \rightarrow [0, 1]$  has been extended to act on vectors of arbitrary length. Given  $\mathbf{x}, \mathbf{y} \in [0, 1]^n$  we have  $B(\mathbf{x}, \mathbf{y}) = (B(x_1, y_1), \dots, B(x_n, y_n))$ . For instance:

```
x <- c(0, 0.5, 1)
y <- c(0.4, 0.6, 0.8)
tnorm_lukasiewicz(x, y)
## [1] 0.0 0.1 0.8
```

Note that many new logical connectives may be generated via existing ones. For example, given any fuzzy implication  $I$ ,  $N(x) = I(x, 0)$  is a fuzzy negation. Moreover, given any t-conorm  $S$  and any negation  $N$ ,  $I(x, y) = S(N(x), y)$  is a fuzzy implication (a so-called (S-N)-implication), see e.g. [3] for more details.

**t-norms** Table 4 lists all the t-norms predefined by the *agop* package. For any t-norm  $T$  and all  $x, y$  it holds  $T_D(x, y) \leq T(x, y) \leq T_M(x, y)$ . Moreover, we have  $T_L(x, y) \leq T_P(x, y)$ .

**t-conorms** Table 5 lists all the t-conorms in the *agop* package. For any t-conorm  $S$  and all  $x, y$  it holds  $S_M(x, y) \leq T(x, y) \leq S_D(x, y)$ . Moreover, we have  $S_P(x, y) \leq S_L(x, y)$ . Also note that  $S$  is a t-conorm if and only if there exists a t-norm  $t$  such that for all  $x, y$  it holds  $S(x, y) = 1 - T(1 - x, 1 - y)$ , see [32].

**Tab. 4.** Exemplary t-norms

Name	Function	Definition
Minimum	<code>tnorm_minimum()</code>	$T_M(x, y) = x \wedge y$
Product	<code>tnorm_product()</code>	$T_P(x, y) = xy$
Łukasiewicz	<code>tnorm_lukasiewicz()</code>	$T_L(x, y) = (x + y - 1) \vee 0$
Drastic	<code>tnorm_drastic()</code>	$T_D(x, y) = \begin{cases} 0 & \text{if } x, y \in [0, 1) \\ x \wedge y & \text{if } x = 1 \text{ or } y = 1 \end{cases}$
Fodor	<code>tnorm_fodor()</code>	$T_F(x, y) = \begin{cases} 0 & \text{if } x + y \leq 1 \\ x \wedge y & \text{if } x + y > 1 \end{cases}$

**Tab. 5.** Exemplary t-conorms

Name	Function	Definition
Maximum	<code>tconorm_minimum()</code>	$S_M(x, y) = x \vee y$
Product	<code>tconorm_product()</code>	$S_P(x, y) = x + y - xy$
Łukasiewicz	<code>tconorm_lukasiewicz()</code>	$S_L(x, y) = (x + y) \wedge 1$
Drastic	<code>tconorm_drastic()</code>	$S_D(x, y) = \begin{cases} 1 & \text{if } x, y \in (0, 1] \\ x \vee y & \text{if } x = 0 \text{ or } y = 0 \end{cases}$
Fodor	<code>tconorm_fodor()</code>	$S_F(x, y) = \begin{cases} 1 & \text{if } x + y \geq 1 \\ x \vee y & \text{if } x + y < 1 \end{cases}$

**Fuzzy negations** TO BE DONE

**Fuzzy implications** TO BE DONE

## 5.4 A Note on Copulas

Copulas are used in probability and statistics to model dependency between random variables (cf. the Sklar theorem). Many copulas are defined e.g. by the `copula` package – we decided not to duplicate its features.

```
library("copula")
cc <- frankCopula(1, dim=2)
pCopula(c(0.5, 0.8), cc) # 0.4197217
pCopula(c(0.9, 1.0), cc) # 0.9
```

Note that t-norms such as  $T_M$ ,  $T_P$ , and  $T_L$  are examples of 2-copulas. On the other hand,  $T_D$  and  $T_F$  are not 2-copulas.

Interestingly, by the Frechet-Hoeffding theorem, we have  $T_L(x, y) \leq C(x, y) \leq T_M(x, y)$  for any  $x, y$  and 2-copula  $C$ , see e.g. [32].

## 5.5 Interesting Spread Measures

..... (see `var()`), standard deviation (see `sd()`), range, interquartile range (IQR, see `IQR()`), median absolute deviation (MAD) etc., that is functions widely used in exploratory data analysis as descriptive statistics.

...TO DO...

D2OWA (`d2owa()`):

$$\text{D2OWA}_{\mathbf{w}}(\mathbf{x}) = \sqrt{\frac{1}{n} \sum_{i=1}^n (x_i - \text{OWA}_{\mathbf{w}}(\mathbf{x}))^2}$$

## 6 Aggregation Operators from the Probabilistic Perspective

By default, theory of aggregation looks at the aggregation operators mainly from the algebraic perspective. Of course, we may also be interested in their probabilistic properties, e.g. in i.i.d. RVs models (the simplest and the most “natural” ones in statistics), cf. [18] for discussion.

Intuitively, a random variable is a method for “producing” input data. An aggregation operator applied on a random variable (possibly multidimensional) is classically called a **statistic**.

### 6.1 Some Notable Probability Distributions

Let  $(X_1, \dots, X_n)$  i.i.d.  $F$ , where  $\text{supp } F = \mathbb{I}$ . In social phenomena modeling, if  $F$  is continuous, we often assume that the underlying density  $f$  is decreasing and convex on  $\mathbb{I}$ , possibly with heavy-tails. E.g. in the bibliometric impact assessment problem, this assumption reflects the fact that higher paper valuations are more difficult to obtain than the lower ones, most of the papers have very small valuation (near 0), and the probability of attaining a high note decreases in at least linear pace.

Let us make a review of some useful statistical distributions, that are not available through “base” R (for other ones, e.g. exponential, normal, uniform, Weibull, etc. refer to the widely-available literature).

#### 6.1.1 Pareto-Type II Distribution

Many generalizations of the Pareto distribution have been proposed (GPD, *Generalized Pareto Distributions*, cf. e.g. [46, 52]). Here we will introduce the so-called Pareto-Type II (Lomax) distribution, which has support  $\mathbb{I} = [0, \infty]$  and is defined with two parameters.

Formally,  $X$  follows the Pareto-II distribution with shape parameter  $k > 0$  and scale parameter  $s > 0$ , denoted  $X \sim \text{P2}(k, s)$ , if its density is of the form

$$f(x) = \frac{ks^k}{(s+x)^{k+1}} \quad (x \geq 0). \quad (2)$$

The cumulative distribution function of  $X$  is then:

$$F(x) = 1 - \frac{s^k}{(s+x)^k} \quad (x \geq 0). \quad (3)$$

The Pareto-Type II distribution is implemented in `agop`: `dpareto2()` gives the p.d.f. (2), `ppareto2()` gives the c.d.f. (3), `qpareto2()` calculates the quantile function,  $F^{-1}$ , and `rpareto2()` generates random deviates.

**Properties.** The expected value of  $X \sim P2(k, s)$  exists for  $k > 1$  and is equal to  $\mathbb{E}X = \frac{s}{k-1}$ . Variance exists for  $k > 2$  and is equal to  $\text{Var } X = \frac{ks^2}{(k-2)(k-1)^2}$ . More generally, the  $i$ -th raw moment for  $k > i$  is given by:  $\mathbb{E}X^i = \frac{\Gamma(i+1)\Gamma(k-i)}{\Gamma(k+1)} ks^i$ .

For a fixed  $s$ , if  $X \sim P2(k_x, s)$  and  $Y \sim P2(k_y, s)$ ,  $k_x < k_y$ , then  $X$  stochastically dominates  $Y$ , denoted  $X \succ Y$ . On the other hand, for a fixed  $k$ , if  $X \sim P2(k, s_x)$  and  $Y \sim P2(k, s_y)$ , then  $s_x > s_y$  implies  $X \succ Y$ .

Most importantly, if  $X \sim P2(k, s)$ , then the conditional distribution of  $X - t$  given  $X > t$ , is  $P2(k, s + t)$   $t \geq 0$ .

Additionally, it might be shown that if  $X \sim P2(k, s)$ , then  $\ln(s + X)$  has c.d.f.  $F(x) = 1 - s^k e^{-kx}$  and density  $f(x) = ks^k e^{-kx}$  for  $x \geq \ln s$ , i.e. has the same distribution as  $Z + \ln s$ , where  $Z \sim \text{Exp}(k) \equiv \Gamma(1, 1/k)$  (exponential distribution).

**Parameter estimation.** Let  $\mathbf{x} = (x_1, \dots, x_n)$  be a realization of the Pareto-Type II i.i.d. sample with known  $s > 0$ . The unbiased (corrected) maximum likelihood estimator for  $k$ :

$$\hat{k}(\mathbf{x}) = \frac{n-1}{\sum_{i=1}^n \ln\left(1 + \frac{1}{s}x_i\right)}.$$

It may be shown that for  $n > 2$  it holds  $\text{Var } \hat{k}(\mathbf{x}) = k^2 \frac{1}{n-2}$ .

`agop` implementation: `pareto2_estimate_mle()` with explicitly set argument `s`.

```
rowMeans(replicate(1000, {
  pareto2_estimate_mle(rpareto2(50, 2, 1.5), s=1.5)
}))
##          k          s
## 1.988461 1.500000
```

For both unknown  $k$  and  $s$  we have:

$$\begin{cases} \hat{k} = \frac{n}{\sum_{i=1}^n \ln(1+x_i/\hat{s})}, \\ 1 + \frac{1}{n} \sum_{i=1}^n \ln(1+x_i/\hat{s}) - \frac{n}{\sum_{i=1}^n (1+x_i/\hat{s})^{-1}} = 0. \end{cases}$$

Unfortunately, the second equation must be solved numerically. It is worth noting that the above system of equations may sometimes have no solution (as the local minimum of the likelihood function may not exist, see [12] for discussion). This estimator may be heavily biased and have a large mean squared error (of course, it is only asymptotically unbiased).

`agop` implementation: `pareto2_estimate_mle()` with explicitly set argument `s`.

```
rowMeans(replicate(1000, {
  pareto2_estimate_mle(rpareto2(50, 2, 1.5))
}), na.rm=TRUE)
##          k          s
## 2.876271 2.418572
```



We see that the estimator’s performance is weak.

A better (in general) estimation procedure was proposed in [53]. The Zhang-Stevens MMS (*minimum mean square error*) (Bayesian) estimator has relatively small bias (often positive) and mean squared error. In agop it is available as: `pareto2_estimate_mmse`.

```
suppressWarnings(rowMeans(replicate(1000, {
  pareto2_estimate_mmse(rpairoto2(50, 2, 1.5))
})))

##          k          s
## 4.037108 3.291486
```

**Goodness-of-fit tests.** `pareto2_test_ad()` ..... – Anderson-Darling goodness-of-fit test (approximate p-value)..... (TO DO: describe) for known  $s$  by means of the `exp_test_ad()` function and the above-mentioned relationship between Pareto-Type II distributions and Exponential ones.

```
x <- rpareto2(100, k=1, s=2)
pareto2_test_ad(x, s=2)

##
## Anderson-Darling goodness-of-fit test for Pareto Type-II
## distribution
##
## data: x
## W = 0.3169, p-value = 0.797
```

**Two-sample  $F$ -test.** The following simple test was introduced in [18]. Let  $(X_1, X_2, \dots, X_{n_1})$  *i.i.d.*  $P2(k_1, s)$  and  $(Y_1, Y_2, \dots, Y_{n_2})$  *i.i.d.*  $P2(k_2, s)$ , where  $s$  is an a-priori known scale parameter. We are going to verify the null hypothesis  $H_0 : k_1 = k_2$  against the two-sided alternative hypothesis  $K : k_1 \neq k_2$ .

It might be shown that  $\sum_{i=1}^n \ln(s + X_i) - n \ln s \sim \Gamma(n, 1/k)$ . This implies that under  $H_0$ , the following test statistic follows the Snedecor  $F$  distribution:

$$R(\mathbf{X}, \mathbf{Y}) = \frac{n_1 \sum_{i=1}^{n_2} \ln \left( 1 + \frac{Y_i}{s} \right)}{n_2 \sum_{i=1}^{n_1} \ln \left( 1 + \frac{X_i}{s} \right)} \stackrel{H_0}{\sim} F^{[2n_2, 2n_1]}. \quad (4)$$

The null hypothesis is accepted iff

$$R(\mathbf{x}, \mathbf{y}) \in [\mathbf{qf}(\frac{\alpha}{2}, 2n_2, 2n_1), \mathbf{qf}(1 - \frac{\alpha}{2}, 2n_2, 2n_1)],$$

where  $\mathbf{qf}(q, d_1, d_2)$  denotes the  $q$ -quantile of  $F^{[d_1, d_2]}$

The  $p$ -value may be determined as follows:

$$p = 2 \left( \frac{1}{2} - \left| \mathbf{pf}(R(\mathbf{x}, \mathbf{y}), 2n_2, 2n_1) - \frac{1}{2} \right| \right), \quad (5)$$

where  $\mathbf{pf}(x, d_1, d_2)$  is the c.d.f. of  $F^{[d_1, d_2]}$ .

agop implementation: `pareto2_test_f()`.

```
x <- rpareto2(35, 1.2, 1)
y <- rpareto2(25, 2.1, 1)
pareto2_test_f(x, y, s=1)
```

```
##
## Two-sample F-test for equality of shape parameters for Type
## II-Pareto distributions with known common scale parameter
##
## data: x and y
## F = 0.3858, p-value = 0.000547
## alternative hypothesis: two-sided
```

### 6.1.2 Discretized Pareto-Type II Distribution

We would say that  $X \sim \text{DP2}(k, s)$ , i.e. it follows the **discretized Pareto-Type II distribution** with shape parameter  $k > 0$  and scale parameter  $s > 0$ , if  $X = \lfloor Y \rfloor$ , where  $Y \sim \text{P2}(k, s)$ .

.....TO BE DONE.....

The Discretized Pareto-Type II distribution is implemented in *agop*: **ddpareto2()** gives the p.m.f., **pdpareto2()** gives the c.d.f., **qdpareto2()** calculates the quantile function, and **rdpareto2()** generates random deviates.

## 6.2 Stochastic Properties of Aggregation Operators

Given  $(X_1, X_2, \dots)$  i.i.d. following a continuous c.d.f.  $F$  it is well-known, see [11], that L-statistics with weights  $c_{i,n} = w(i/n)$ , for  $w : [0, 1] \rightarrow \mathbb{I}$ , are asymptotically normally distributed. A similar result for the same weight setting has been shown for S-statistics, see [27].

For i.i.d. samples of finite length we have e.g. the following result [22]:

**Theorem 25.** *Let  $\mathbf{X} = (X_1, \dots, X_n)$  be a sequence of i.i.d. random variables with continuous c.d.f.  $F$  defined on  $\mathbb{R}_{0+}$ . Then the c.d.f. of  $H(\mathbf{X})$  for  $x \in [0, n)$  is given by*

$$D_n(x) = \mathcal{I}(F(\lfloor x + 1 \rfloor^{-0}); n - \lfloor x \rfloor, \lfloor x \rfloor + 1),$$

where  $\mathcal{I}(p; a, b)$  is the regularized incomplete beta function (**pbeta()** in *R*).

More generally, the c.d.f. of some quasi-S-statistics may be expressed as an incomplete beta function, see [27]. Note that, unlike in the case of the distribution of “ordinary” order statistics (see [11]), the parameters  $a, b$  of  $\mathcal{I}$  are functions of  $x$  here.

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