Lab: Abstract Interpretation

(Week 8)

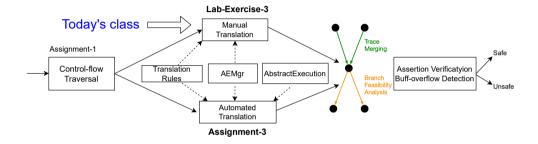
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Lab-2 Marks and Lab-3 Code Template

- Lab-2 marks are out and let us go through Quiz-2 and Exercise-2!
- Remember to git pull or docker pull to get the code template for Lab-Exercise-3

Today's class



Quiz-3 + Lab-Exercise-3 + Assignment-3

- Quiz-3 (5 points) (due date: 23:59, Wednesday, Week 10)
 - Abstract domain and soundness
 - Handling loops with widening and narrowing
- Lab-Exercise-3 (5 points) (due date: 23:59, Wednesday, Week 10)
 - Goal: Coding exercise to manually update abstract trace based on abstract execution rules and verify the assertions embedded in the code.
 - Specification: https://github.com/SVF-tools/ Software-Security-Analysis/wiki/Lab-Exercise-3

Quiz-3 + Lab-Exercise-3 + Assignment-3

- Quiz-3 (5 points) (due date: 23:59, Wednesday, Week 10)
 - Abstract domain and soundness
 - Handling loops with widening and narrowing
- Lab-Exercise-3 (5 points) (due date: 23:59, Wednesday, Week 10)
 - **Goal:** Coding exercise to manually update abstract trace based on abstract execution rules and verify the assertions embedded in the code.
 - Specification: https://github.com/SVF-tools/ Software-Security-Analysis/wiki/Lab-Exercise-3
- Assignment-3 (25 points) (due date: 23:59, Wednesday, Week 11)
 - Goal: Perform automated abstract trace update on ICFG for assertion checking and buffer overflow detection
 - Specification: https:

//github.com/SVF-tools/Software-Security-Analysis/wiki/Assignment-3

• SVF AE APIs: https:

//github.com/SVF-tools/Software-Security-Analysis/wiki/AE-APIs

Lab-3 Exercise: Manual Translation to Compute Abstract States

- Let us look at how to write abstract execution code to analyze examples of a loop-free and a loop C-like code by manually collecting abstract states at each program statement to form the abstract trace
- You will need to finish all the coding tests in AEMgr.cpp under Lab-Exercise-3

```
1 struct A{int f0;};
2 void main() {
     struct A * p :
     int * q :
    int x
    p = malloc;
     q = \&(p \rightarrow f0);
     *a = 10:
     x = *a:
10
    svf_assert(x == 10);
11 }
```

```
NodeID p = getNodeID("p", 1);
NodeID q = getNodeID("q");
NodeID x = getNodeID("x");
...
```

```
-----Var and Value-----
```

AEState:printAbstractState()

Source code

Translation for Abstract execution

```
1 struct A{int f0;};
2 void main() {
     struct A * p;
     int*q:
     int x:
    p = malloc;
     q = \&(p \rightarrow f0);
     *a = 10:
     x = *a:
10
    svf_assert(x == 10);
11 }
```

```
NodeID p = getNodeID("p", 1);
NodeID q = getNodeID("q");
NodeID x = getNodeID("x");
as[p] = AddressValue(getMemObjAddress("malloc"));
...
```

```
------Var and Value-------
Var4 (malloc) Value: 0x7f000004
Var1 (p) Value: 0x7f000004
```

0x7f000004 (or 2130706436 in decimal) represents the virtual memory address of this object Each SVF object starts with 0x7f + its ID.

Source code

Translation for Abstract execution

```
1 struct A{int f0;};
2 void main() {
     struct A * p;
     int*q:
     int x:
     p = malloc;
     q = \&(p \rightarrow f0);
     *a = 10:
     x = *a:
10
    svf_assert(x == 10);
11 }
```

```
NodeID p = getNodeID("p", 1);
NodeID q = getNodeID("q");
NodeID x = getNodeID("x");
NodeID x = getNodeID("x");
as[p] = AddressValue(getMemObjAddress("malloc"));
as[q] = AddressValue(getGepObjAddress("p", 0));
...
```

```
Var4 (malloc) Value: 0x7f000004
Var1 (p) Value: 0x7f000004
Var2 (q) Value: 0x7f000005
```

getGepObjAddress returns the field address of the aggregate object p The virual address also in the form of 0x7f.. + VaxID

Source code

Translation for Abstract execution

```
1 struct A{int f0;};
2 void main() {
     struct A * p;
     int*q:
    int x:
     p = malloc;
     q = \&(p \rightarrow f0);
     *a = 10:
     x = *a:
10
    svf_assert(x == 10);
11 }
```

```
1 NodeID p = getNodeID("p", 1);
2 NodeID q = getNodeID("q");
3 NodeID x = getNodeID("x");
4 as[p] = AddressValue(getMemObjAddress("malloc"));
5 as[q] = AddressValue(getGepObjAddress("p", 0));
6 as.storeValue(q, IntervalValue(10, 10));
7 as[x] = as.loadValue(q);
8 ...
```

```
Var4 (malloc) Value: 0x7f000004
Var1 (p) Value: 0x7f000004
Var2 (q) Value: 0x7f000005
Var3 (x) Value: [10, 10]
Var5 (0x7f000005) Value: [10, 10]
```

store value of 5 to address ox7f000005

load the value from ox7f000005 to x

Source code

Translation for Abstract execution

```
1 struct A{int f0;};
2 void main() {
     struct A * p;
     int*q:
     int x:
     p = malloc;
     q = \&(p \rightarrow f0);
     *a = 10:
     x = *a:
10
    svf_assert(x == 10);
11 }
```

```
1 NodeID p = getNodeID("p", 1);
2 NodeID q = getNodeID("q");
3 NodeID x = getNodeID("x");
4 as[p] = AddressValue(getMemObjAddress("malloc"));
5 as[q] = AddressValue(getGepObjAddress("p", 0));
6 as.storeValue(q, IntervalValue(10, 10));
7 as[x] = as.loadValue(q);
```

 ${\tt svf_assert}$ checking is done in test.cpp.

```
Var4 (malloc) Value: 0x7f000004
Var1 (p) Value: 0x7f000004
Var2 (q) Value: 0x7f000005
Var3 (x) Value: [10, 10]
Var5 (0x7f000005) Value: [10, 10]
```

assertion checking

Source code

Translation for Abstract execution

```
1 NodeID argy = getNodeID("argy");
2 as[argy] = IntervalValue(5, 15)
3 ...
```

```
-----Var and Value------
Var1 (argv) Value: [5, 15]
```

assume $5 \le \text{argv} \le 15$

Source code

Translation for Abstract execution

```
1 NodeID argy = getNodeID("argy");
2 as[argy] = IntervalValue(5, 15)
3 NodeID x = getNodeID("x");
4 ...
```

as:

as_true:

```
------Var and Value---------
Var1 (argv) Value: [5, 15]
Var2 (x) Value: [11, 11]
```

Source code

Translation for Abstract execution

```
int main(int argv) {
  int x = 10;
  if(argv > 10)
     x + +;
    else
     x + = 2;
    svf_assert(x <= 12);
  }
</pre>
```

as:

as_true:

Source code

Translation for Abstract execution

```
1 int main(int argv) {
    int x = 10:
   if(argv > 10)
    x + +:
    else
                            12
     x + = 2:
                            13
   svf_assert(x \le 12);
```

```
1 NodeID argv = getNodeID("argv"):
 2 as[argv] = IntervalValue(5, 15)
 3 NodeID x = getNodeID("x"):
 5 AbstractState as after if:
6 AbstractValue cmp_true = as[argv] > IntervalValue(10, 10);
7 // feasibility checking
8 cmp true.meet with(IntervalValue(1, 1)):
9 if (!cmp_true.getInterval().isBottom()) {
      AbstractState as true = as:
      as_true[x] = as_true[x] + IntervalValue(1, 1);
      //Join the states at the control-flow joint point
      as_after_if.joinWith(as_true);
14 }
16 AbstractValue cmp false = as[argv] > IntervalValue(10, 10):
  cmp_false.meet_with(IntervalValue(0, 0));
18 if (!cmp_false.getInterval().isBottom()){
      AbstractState as_false = as:
      as false[x] = as false[x] + IntervalValue(2, 2):
      as after if.joinWith(as false):
22 }
23 . . .
```

as:

```
Var1 (argv) Value: [5, 15]
Var2 (x) Value: [10, 10]
```

as_true:

as_false:

```
------Var and Value---------
Var1 (argv) Value: [5, 15]
Var2 (x) Value: [12, 12]
```

Source code

Translation for Abstract execution

```
1 int main(int argv) {
    int x = 10:
   if(argv > 10)
     x + +:
    else
     x + = 2:
7
   svf_assert(x \le 12);
8
```

```
NodeID argy = getNodeID("argy"):
2 as[argv] = IntervalValue(5, 15)
3 NodeID x = getNodeID("x"):
 5 AbstractState as after if:
 6 AbstractValue cmp true = as[argv] > IntervalValue(10, 10);
7 // feasibility checking
8 cmp_true.meet_with(IntervalValue(1, 1));
9 if (!cmp true.getInterval().isBottom()) {
      AbstractState as_true = as;
      as true[x] = as true[x] + IntervalValue(1, 1):
      //Join the states at the control-flow joint point
      as_after_if.joinWith(as_true);
13
16 AbstractValue cmp false = as[argv] > IntervalValue(10, 10):
17 // feasibility checking
18 cmp_false.meet_with(IntervalValue(0, 0));
19 if (!cmp_false.getInterval().isBottom()){
      AbstractState as_false = as:
21
      as_false[x] = as_false[x] + IntervalValue(2, 2);
      //Join the states at the control-flow joint point
      as_after_if.joinWith(as_false);
24 }
25 as = as after if:
```

svf_assert checking is done in test.cpp.

Source code

Translation for Abstract execution

Abstract trace

as_after_if, as:

as_true:

as_false:

Before entering loop

```
1 int main() {
2    int a = 0;
3    while(a < 10) {
4        a + +;
5    }
6    svf_assert(a == 10);
7    return 0;
8 }</pre>
```

as, entry_as, pre_as and post_as:

The initialization of a.

Source code

Translation for Abstract execution

Widening delay stage

```
1 int main() {
2   int a = 0;
3   while(a < 10) {
4      a + +;
5   }
5   }
6   svf_assert(a == 10);
7   return 0;
8   }
1 int main() {
3      AbstractState to tmp_as.joinWith tmp_as.joinWit
```

pre₋as after Line 8:

```
Var1 (a) Value: [0, 0]
```

as after Line 15:

Widening delay with i=0.

Source code

Translation for Abstract execution

Widening delay stage

pre_as after Line 8:

```
Var1 (a) Value: [0, 1]
```

as after Line 15:

Widening delay with i=1.

Source code

Translation for Abstract execution

Widening delay stage

pre₋as after Line 8:

as after Line 15:

Widening delay with i=2.

Source code

Translation for Abstract execution

Widening stage

```
1 int main() {
2   int a = 0;
3   while(a < 10) {
4     a + +;
5   }
6   svf_assert(a == 10);
7   return 0;
8 }</pre>
```

```
2 for (int i = 0: : ++i) {
      if (i < 3) {
           pre as = as:
      } else {
           // widen and widen fixpoint
           if (increasing) {
               as = pre as.widening(as):
               if (pre as >= as) {
                   pre as = as:
                   increasing = false;
                   continue:
               pre as = as:
           } else {
               // narrow
      as[a].meet with(IntervalValue(
20
           IntervalValue::minus_infinity(), 9));
      as[a] = as[a] + IntervalValue(1, 1);
22
      post as = as:
24
```

pre_as before Line 9:

```
Var1 (a) Value: [0, 2]
```

as before Line 9:

as after Line 9:

```
------Var and Value----------
Var1 (a) Value: [0, +∞]
```

Widening stage where i=3.

Source code

Translation for Abstract execution

Widening stage

```
int main() {
   int a = 0;
   while(a < 10) {
        a + +;
    }
   svf_assert(a == 10);
   return 0;
   }
}</pre>
```

```
2 for (int i = 0: : ++i) {
      if (i < 3) {
           pre as = as:
      } else {
           // widen and widen fixpoint
           if (increasing) {
               as = pre as.widening(as):
               if (pre as >= as) {
                   pre as = as:
                   increasing = false;
                   continue:
               pre as = as:
           } else {
               // narrow
      as[a].meet with(IntervalValue(
20
           IntervalValue::minus_infinity(), 9));
      as[a] = as[a] + IntervalValue(1, 1);
22
      post as = as:
24
```

pre_as before Line 9:

as before Line 9:

```
Var1 (a) Value: [0, 9]
```

as after Line 9:

Widening stage where i=4.

Source code

Translation for Abstract execution

Narrowing stage

```
int main() {
  int a = 0;
  while(a < 10) {
    a + +;
  }
}

svf_assert(a == 10);

return 0;
}</pre>
```

```
2 for (int i = 0; ; ++i) {
       if (i < 3) {
           pre as = as:
           // widen and widen fixpoint
           if (increasing) {
           } else {
               as = pre as.narrowing(as):
               if (as >= pre_as) {
                   break:
               pre as = as:
       as[a].meet with(IntervalValue(
           IntervalValue::minus_infinity(), 9));
       as[a] = as[a] + IntervalValue(1, 1):
       post as = as:
22 }
23 . . .
```

pre_as before Line 11:

```
Var1 (a) Value: [0, +∞]
```

as before Line 11:

as after Line 11:

Narrowing stage where i=5.

Source code

Translation for Abstract execution

Narrowing stage

```
int main() {
   int a = 0;
   while(a < 10) {
        a + +;
      }
   svf_assert(a == 10);
   return 0;
   }
}</pre>
```

```
2 for (int i = 0; ; ++i) {
       if (i < 3) {
           pre as = as:
           // widen and widen fixpoint
           if (increasing) {
           } else {
               as = pre as.narrowing(as):
               if (as >= pre_as) {
                   break:
               pre as = as:
       as[a].meet with(IntervalValue(
           IntervalValue::minus_infinity(), 9));
       as[a] = as[a] + IntervalValue(1, 1):
       post as = as:
22 }
23 . . .
```

pre_as before Line 11:

as before Line 11:

as after Line 11:

Narrowing stage where i=6.

Source code

Translation for Abstract execution

After exiting loop

```
1 int main() {
2    int a = 0;
3    while(a < 10) {
4        a + +;
5    }
6    svf_assert(a == 10);
7    return 0;
8   }</pre>
```

```
as:
```

After analyzing loop.

Source code Translation for Abstract execution