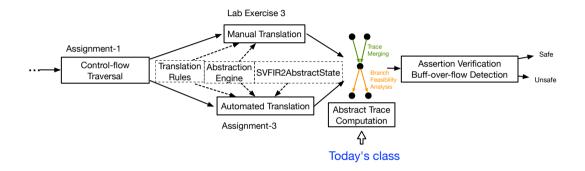
# **Lab: Abstract Interpretation**

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## Today's class



# Quiz-3 + Lab-Exercise-3 + Assignment-3

- Quiz-3 (5 points)
  - Abstract domain and soundness
  - Handling loops with widening and narrowing
- Lab-Exercise-3 (5 points)
  - Goal: Coding exercise to manually update abstract trace based on abstract execution rules and verify the assertions embedded in the code.
  - Specification: https://github.com/SVF-tools/ Software-Security-Analysis/wiki/Lab-Exercise-3

# Quiz-3 + Lab-Exercise-3 + Assignment-3

- Quiz-3 (5 points)
  - Abstract domain and soundness.
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- Lab-Exercise-3 (5 points)
  - Goal: Coding exercise to manually update abstract trace based on abstract execution rules and verify the assertions embedded in the code.
  - **Specification:** https://github.com/SVF-tools/ Software-Security-Analysis/wiki/Lab-Exercise-3
- Assignment-3 (25 points)
  - Goal: Perform automated abstract trace update on ICFG for assertion checking and buffer overflow detection
  - Specification: https:

//github.com/SVF-tools/Software-Security-Analysis/wiki/Assignment-3

SVF AE APIs: https:

//github.com/SVF-tools/Software-Security-Analysis/wiki/AE-APIs

# Lab 3 Coding Exercise: Abstract States and Abstract Traces

- Let us look at how to write abstract execution code to analyze examples of a loop-free and a loop C-like code by manually collecting abstract states at each program statement to form the abstract trace
- You will need to finish all the coding tests in AEMgr.cpp under Lab-Exercise-3

```
1 struct A{int f0;};
2 void main() {
    struct A * p :
    int*q:
    int x
   p = malloc;
    q = \&(p \rightarrow f0);
    *a = 10:
    x = *a:
    if(x == 10)
    x + +:
    assert(x == 11);
13 }
```

```
NodeID p = getNodeID("p", 1);
NodeID q = getNodeID("q");
NodeID x = getNodeID("x");
...
```

```
-----Var and Value-----
```

```
-----Loc and Value-----
```

Source code

Abstract execution

```
1 struct A{int f0;};
2 void main() {
     struct A * p;
     int*q:
     int x:
    p = malloc;
     q = \&(p \rightarrow f0);
     *a = 10:
     x = *a:
     if(x == 10)
11
     x + +:
     assert(x == 11):
13 }
```

```
1 NodeID p = getNodeID("p", 1);
2 NodeID q = getNodeID("q");
3 NodeID x = getNodeID("x");
4 es[p] = AddressValue(getMemObjAddress("malloc"));
5 ...
```

```
Var4 (malloc) Value: 0x7f000004
Var1 (p) Value: 0x7f000004
```

```
-----Loc and Value-----
```

0x7f000004 (or 2130706436 in decimal) represents the virtual memory address of this object Each SVF object starts with 0x7f + its ID.

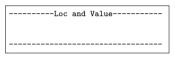
Source code

Abstract execution

```
1 struct A{int f0;};
2 void main() {
     struct A * p:
     int*q;
     int x:
     p = malloc:
     q = \&(p \rightarrow f0):
     *q = 10:
     x = *a:
     if(x == 10)
11
     x + +:
     assert(x == 11);
13 }
```

```
NodeID p = getNodeID("p", 1);
NodeID q = getNodeID("q");
NodeID x = getNodeID("x");
Stock = getNodeID("x");
ses[p] = AddressValue(getMemObjAddress("malloc"));
ses[q] = AddressValue(getGepObjAddress("p", 0));
...
```

```
------Var and Value-------
Var4 (malloc) Value: 0x7f000004
Var1 (p) Value: 0x7f000005
Var2 (q) Value: 0x7f000005
```



getGepObjAddress returns the field address of the aggregate object p The virual address also in the form of 0x7f.. + VaxID

Source code

Abstract execution

```
1 struct A{int f0;};
2 void main() {
     struct A * p;
     int*q:
     int x:
     p = malloc;
     q = \&(p \rightarrow f0);
     *a = 10:
     x = *a:
     if(x == 10)
11
     x + +:
     assert(x == 11):
13 }
```

```
Var4 (malloc) Value: 0x7f000004
Var1 (p) Value: 0x7f000004
Var2 (q) Value: 0x7f000005
Var3 (x) Value: [10, 10]
```

store value of 5 to address ox7f000005

Source code

Abstract execution

```
1 struct A{int f0;};
2 void main() {
     struct A * p;
    int*q:
    int x:
    p = malloc;
    q = \&(p \rightarrow f0);
     *a = 10:
    x = *a:
    if(x == 10)
     x + +:
     assert(x == 11):
13
```

```
NodeID p = getNodeID("p", 1):
  NodeID g = getNodeID("g"):
  NodeID x = getNodeID("x");
  es[p] = AddressValue(getMemObiAddress("malloc")):
 5 es[q] = AddressValue(getGepObjAddress("p", 0));
 6 for (auto addr : es[q].getAddrs()) {
      es.store(addr. IntervalValue(10, 10)):
 9 for (const auto &addr: es[q].getAddrs()) {
      es[x].join with(es.load(addr)):
  AbstractState es after if:
13 AbstractValue cmp_true = es[x] == IntervalValue(10, 10);
14 cmp_true.meet_with(IntervalValue(1, 1));
15 if (!cmp true.isBottom()) {
      es[x] = es[x] + IntervalValue(1, 1);
17 3
18 . . .
```

```
Var4 (malloc) Value: 0x7f000004
Var1 (p) Value: 0x7f000004
Var2 (q) Value: 0x7f000005
Var3 (x) Value: [11, 11]
```

handle branch

Source code

Abstract execution

```
1 struct A{int f0;};
2 void main() {
     struct A * p;
     int*q:
     int x:
     p = malloc;
     q = \&(p \rightarrow f0);
     *a = 10:
     x = *a:
     if(x == 10)
11
     x + +:
     assert(x == 11):
12
13
```

```
Var4 (malloc) Value: 0x7f000004
Var1 (p) Value: 0x7f000004
Var2 (q) Value: 0x7f000005
Var3 (x) Value: [11, 11]
```

assertion checking

Source code

Abstract execution

## Before entering loop

```
1 int main() {
2    int a = 0;
3    while(a < 10) {
4       a + +;
5    }
6    assert(a == 10);
7    return 0;
8 }</pre>
```

es, entry\_es, pre\_es and post\_es:

The initialization of a.

Source code

Abstract execution

## Widening delay stage

```
2 for (int i = 0; AbstractStat tmp.es.joinW tmp.es.joinW tmp.es.joinW es = tmp.es;
3 while(a < 10) { 8 pre.es = 6 }
4 a + +; 9 else { // widen }
5 } 11 csture 0; 12 }
6 assert(a == 10); 13 es[a].meet_w return 0; 14 Interest es = 6 }
7 return 0; 14 es[a] = es[a]
```

## pre₋es after Line 8:

```
Var1 (a) Value: [0, 0]
```

## es after Line 15:

Widening delay with i=0.

Source code

Abstract execution

## Widening delay stage

```
| 2 | for (int i | Abstract | Abs
```

```
1 ...

2 for (int i = 0; ; ++i) {
3    AbstractState tmp_es;
4 tmp_es.joinWith(entry_es);
5 tmp_es.joinWith(entry_es);
6 es = tmp_es;
7 if (i < 3) {
9    pre_es = AbstractState(es);
9 } else {
10    // widen and widen fixpoint
11    ...
12 }
13 es[a].meet_with(IntervalValue(
14    IntervalValue::minus_infinity(), 9));
15 es[a] = es[a] + IntervalValue(1, 1);
16 post_es = es;
17 }
18 ...
```

## pre₋es after Line 8:

```
Var1 (a) Value: [0, 1]
```

### es after Line 15:

Widening delay with i=1.

Source code

Abstract execution

## Widening delay stage

```
2 for (int i = 0: : ++i) {
                                        AbstractState tmp_es;
                                        tmp es.joinWith(post es):
1 int main() {
                                        tmp es.joinWith(entry es):
     int a = 0:
                                        es = tmp_es;
                                        if (i < 3) {
    while(a < 10) {
                                            pre es = AbstractState(es):
                                       } else {
       a + +:
                                            // widen and widen fixpoint
     assert(a == 10):
                                        es[a] meet with(IntervalValue(
     return 0:
                                            IntervalValue::minus infinity(). 9)):
                                 15
                                        es[a] = es[a] + IntervalValue(1, 1):
8
                                        post_es = es;
```

### pre\_es after Line 8:

## es after Line 15:

Widening delay with i=2.

Source code

Abstract execution

## Widening stage

```
int main() {
  int a = 0;
  while(a < 10) {
    a + +;
  }
  assert(a == 10);
  return 0;
}</pre>
```

```
2 for (int i = 0: : ++i) {
       if (i < 3) {
           pre es = AbstractState(es):
       } else {
           // widen and widen fixpoint
           if (increasing) {
               es = pre es.widening(es):
               if (pre es >= es) {
                   pre es = es:
                   increasing = false;
                   continue:
               pre es = es:
           } else {
               // narrow
       es[a].meet with(IntervalValue(
20
           IntervalValue::minus_infinity(), 9));
       es[a] = es[a] + IntervalValue(1, 1);
       post es = es:
24 7
25
```

#### pre\_es before Line 9:

```
-----Var and Value------Var 1 (a) Value: [0, 2]
```

## es before Line 9:

## es after Line 9:

Widening stage where i=3.

Source code

Abstract execution

## Widening stage

```
int main() {
  int a = 0;
  while(a < 10) {
    a + +;
  }
  assert(a == 10);
  return 0;
}</pre>
```

```
2 for (int i = 0: : ++i) {
      if (i < 3) {
           pre es = AbstractState(es):
      } else {
           // widen and widen fixpoint
           if (increasing) {
               es = pre es.widening(es):
               if (pre es >= es) {
                   pre es = es:
                   increasing = false;
                   continue:
               pre es = es:
           } else {
               // narrow
      es[a].meet with(IntervalValue(
20
           IntervalValue::minus_infinity(), 9));
      es[a] = es[a] + IntervalValue(1, 1);
      post es = es:
24 }
25
```

#### pre\_es before Line 9:

## es before Line 9:

```
Var1 (a) Value: [0, 9]
```

## es after Line 9:

Widening stage where i=4.

## Source code

## Abstract execution

## **Narrowing stage**

```
int main() {
  int a = 0;
  while(a < 10) {
    a + +;
  }
  assert(a == 10);
  return 0;
}</pre>
```

```
2 for (int i = 0; ; ++i) {
       if (i < 3) {
           pre es = AbstractState(es):
           // widen and widen fixpoint
           if (increasing) {
           } else {
               es = pre es.narrowing(es):
               if (es >= pre_es) {
                   break:
               pre es = es:
       es[a].meet with(IntervalValue(
           IntervalValue::minus_infinity(), 9));
20
       es[a] = es[a] + IntervalValue(1, 1):
       post_es = es;
21
22 }
23 . . .
```

## pre\_es before Line 11:

### es before Line 11:

### es after Line 11:

```
-----Var and Value------
Var1 (a) Value: [0, 9]
```

Narrowing stage where i=5.

Source code

Abstract execution

## **Narrowing stage**

```
1 int main() {
2   int a = 0;
3   while(a < 10) {
4     a + +;
5   }
6   assert(a == 10);
7   return 0;
8 }</pre>
```

```
2 for (int i = 0; ; ++i) {
       if (i < 3) {
           pre es = AbstractState(es):
           // widen and widen fixpoint
           if (increasing) {
           } else {
               es = pre es.narrowing(es):
               if (es >= pre_es) {
                   break:
               pre es = es:
       es[a].meet with(IntervalValue(
           IntervalValue::minus_infinity(), 9));
20
       es[a] = es[a] + IntervalValue(1, 1):
       post_es = es;
21
22 }
23 . . .
```

### pre\_es before Line 11:

```
-----Var and Value------Var1 (a) Value: [0, 9]
```

### es before Line 11:

```
Var1 (a) Value: [0, 9]
```

### es after Line 11:

```
-----Var and Value------
Var1 (a) Value: [0, 9]
```

Narrowing stage where i=6.

Source code

Abstract execution

## After exiting loop

```
1 int main() {
2    int a = 0;
3    while(a < 10) {
4        a + +;
5    }
6    assert(a == 10);
7    return 0;
8 }</pre>
```

#### es:

```
Var1 (a) Value: [10, 10]
```

After analyzing loop.

Source code

Abstract execution