

Handbook

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1 Grafos

1.1 Dinic

Insertar utilidad del algoritmo:

C++:

```
1 struct edge{
2     int x, y, flow;
3 };
4
5 int ans;
6 vector<edge> edges;
7 vector<vector<int> > grafo;
8 vector<int> sn;
9
10 void addEdge(int x, int y, int flow){
11     grafo[x].PB(edges.size());
12     edges.PB({x, y, flow});
13
14     grafo[y].PB(edges.size());
15     edges.PB({y, x, 0});
16 }
17
18 int bfs(int &ori, int &target){
19     int x = ori, y, flow;
20
21     FOR(i, 0, target + 1) sn[i] = INF;
22
23     sn[x] = 0;
24     deque<int> q(1, x);
25
26     while(!q.empty()){
27         x = q.F(); q.P_F();
28
29         for(auto &e: grafo[x]){
30             auto &edge = edges[e];
31             y = edge.y;
32             flow = edge.flow;
33
34             if(flow <= 0) continue;
35             if(sn[y] != INF) continue;
36             sn[y] = sn[x] + 1;
37             q.PB(y);
38         }
39     }
40
41     return sn[target];
42 }
43
44 int dfs(int ori, int &target, int min_flow){
45     int flow = INF, y;
46
47     for(auto &e_id: grafo[ori]){
48         auto &e = edges[e_id];
49         y = e.y;
```

```

50
51     if(sn[y] != 1 + sn[ori]) continue;
52     if(e.flow <= 0) continue;
53
54     if(y == target){
55         flow = min(min_flow, e.flow);
56         ans += flow;
57         edges[e_id].flow -= flow;
58         edges[e_id^1].flow += flow;
59         return flow;
60     }
61
62     flow = dfs(y, target, min(min_flow, e.flow));
63
64     if(flow != INF){
65         edges[e_id].flow -= flow;
66         edges[e_id^1].flow += flow;
67         return flow;
68     }
69 }
70
71 return flow;
72 }

```

2 Strings

2.1 Función Z

Insertar utilidad del algoritmo:

C++:

```
1 vector<int> z_function(string s) {
2     int n = (int) s.length();
3     vector<int> z(n);
4     for (int i = 1, l = 0, r = 0; i < n; ++i) {
5         if (i <= r)
6             z[i] = min (r - i + 1, z[i - l]);
7         while (i + z[i] < n && s[z[i]] == s[i + z[i]])
8             ++z[i];
9         if (i + z[i] - 1 > r)
10            l = i, r = i + z[i] - 1;
11     }
12     return z;
13 }
```

2.2 KMP

Insertar utilidad del algoritmo:

C++:

```
1 vector<int> z;
2
3 void kmp(string &s){
4     int j, n = s.size();
5     z.resize(n);
6
7     FOR(i, 1, n){
8         j = z[i - 1];
9         while(j > 0 and s[i] != s[j]) j = z[j - 1];
10        if(s[i] == s[j]) j++;
11        z[i] = j;
12    }
13 }
```

2.3 Suffix array

Devuelve un arreglo con el orden lexicográfico de los sufijos de un string S

C++:

```
1 vector<int> p, c;
2 void count_sort(vector<int> &p, vector<int> &c){
3     int n = p.size();
4     vector<int> cnt(n), p_new(n), pos(n);
5     pos[0] = 0;
6     for(auto x : c) cnt[x]++;
```

```

7   for(int i = 1 ; i < n ; ++i) pos[i] = pos[i - 1] + cnt[i - 1];
8   for(auto x : p){
9       int i = c[x];
10      p_new[pos[i]] = x;
11      pos[i]++;
12  }
13  p = p_new;
14 }
15 vector<int> suffix_array(string &s){
16     s+=" ";
17     int n = s.size();
18     p.resize(n);
19     c.resize(n);
20     vector<pair<char, int>> a(n);
21     for(int i = 0 ; i < n ; ++i) a[i] = {s[i], i};
22     sort(a.begin(), a.end());
23     for(int i = 0 ; i < n ; ++i) p[i] = a[i].second;
24     c[p[0]] = 0;
25     for(int i = 1 ; i < n ; ++i) c[p[i]] = a[i].first == a[i - 1].first ? c[p[i - 1]] : c[p[i]
26     int k = 0, shift;
27     while( (1<<k) < n ){
28         shift = 1<<k;
29         for(int i = 0 ; i<n ; ++i)
30             p[i] = (p[i] - (1<<k) + n) % n;
31         count_sort(p,c);
32         vector<int> c_new(n);
33         c_new[p[0]] = 0;
34         for(int i = 1 ; i < n ; ++i){
35             pair<int, int> prev = {c[p[i - 1]], c[ (p[i - 1] + shift) % n]};
36             pair<int, int> now = {c[p[i]], c[(p[i] + shift) % n]};
37             if(prev == now) c_new[p[i]] = c_new[p[i - 1]];
38             else c_new[p[i]] = c_new[p[i - 1]] + 1;
39         }
40         c = c_new;
41         k++;
42     }
43     return p;
44 }

```

Java:

```

1  static int[] p, c;
2  public static class Suffix implements Comparable<Suffix> {
3      int index, r, next;
4      public Suffix(int index, int rank, int next){
5          this.index = index; this.r = rank; this.next = next;
6      }
7      public int compareTo(Suffix s){
8          return r != s.r ? r - s.r : (next != s.next ? next - s.next : index - s.index);
9      }
10 }
11 public static int[] sort(int[] p, int[] c){
12     int N = p.length;
13     int[] cnt = new int[N], pos = new int[N], p_new = new int[N];
14     for(int e : c) cnt[e]++;
15     for(int i = 1 ; i < N ; ++i) pos[i] = pos[i - 1] + cnt[i - 1];
16     for(int x : p){
17         p_new[pos[c[x]]] = x; pos[c[x]]++;

```

```

18     }
19     p = p_new;
20     return p;
21 }
22 public static int[] suffixArray(String s) {
23     s+="$";
24     int n = s.length();
25     c = new int[n];
26     p = new int[n];
27     Suffix[] su = new Suffix[n];
28     for (int i = 0; i < n; ++i) su[i] = new Suffix(i, s.charAt(i), 0);
29     Arrays.sort(su);
30     for(int i = 0 ; i < n ; ++i) p[i] = su[i].index;
31     c[p[0]] = 0;
32     for(int i = 1 ; i < n ; ++i) c[p[i]] = su[i].r == su[i - 1].r ? c[p[i-1]] : c[p[i-1]] + 1;
33     int k = 0, shift;
34     while((1<<k) < n){
35         shift = (1<<k);
36         for(int i = 0 ; i < n ; ++i) p[i] = (p[i] - shift + n) % n;
37         p = sort(p, c);
38         int[] c_new = new int[n];
39         c_new[p[0]] = 0;
40         for(int i = 1 ; i < n ; ++i)
41             c_new[p[i]] = (c[p[i]] == c[p[i-1]] && c[(p[i]+shift) % n] == c[(p[i - 1] + shift)
42                 ? c_new[p[i - 1]] : c_new[p[i - 1]] + 1;
43         c = c_new;
44         ++k;
45     }
46     return p;
47 }

```

2.4 Longest Common Prefix on Suffixs

Devuelve un arreglo que contiene el largo del prefijo común máximo entre 2 sufijos i e $i+1$

C++:

```

1 vector<int> lcp(vector<int> &p, vector<int> &c, string &s){
2     int n = p.size();
3     vector<int> lcp(n);
4     int k = 0;
5     for(int i = 0 ; i < n - 1 ; ++i){
6         int pi = c[i];
7         int j = p[pi - 1];
8         while(s[i + k] == s[j + k]) k++;
9         lcp[pi] = k;
10        k = max(k - 1, 0);
11    }
12    return lcp;
13 }

```

Java:

```

1 static int[] p, c, LCP;
2

```

```

3 static int[] lcp(int[] p, int[] c, String s){
4     int n = p.length;
5     LCP = new int[n];
6     int k = 0;
7     for(int i = 0 ; i < n - 1 ; ++i){
8         int pi = c[i];
9         int j = p[pi - 1];
10        while(s.charAt(i + k) == s.charAt(j + k)) k++;
11        LCP[pi] = k;
12        k = Math.max(k - 1, 0);
13    }
14    return LCP;
15 }

```

3 Búsqueda

3.1 Ternary Search

Insertar utilidad del algoritmo:

C++:

```
1 #define ld long double
2
3 ld ternary_search(ld l, ld r) {
4     ld eps = 1e-9;
5     ld m1, m2, f1, f2;
6     while (r - l > eps) {
7         m1 = l + (r - l) / 3;
8         m2 = r - (r - l) / 3;
9         f1 = f(m1);           //evaluates the function at m1
10        f2 = f(m2);           //evaluates the function at m2
11        if (f1 < f2) l = m1;
12        else r = m2;
13    }
14
15    //return the maximum of f(x) in [l, r]
16    return f(l);
17 }
```


4 Geometría

4.1 Convex Hull

Insertar utilidad del algoritmo:

C++:

```
1 struct pt {
2     double x, y;
3 };
4
5 int orientation(pt a, pt b, pt c) {
6     double v = a.x*(b.y-c.y)+b.x*(c.y-a.y)+c.x*(a.y-b.y);
7     if (v < 0) return -1; // clockwise
8     if (v > 0) return +1; // counter-clockwise
9     return 0;
10 }
11
12 bool cw(pt a, pt b, pt c, bool include_collinear) {
13     int o = orientation(a, b, c);
14     return o < 0 || (include_collinear && o == 0);
15 }
16 bool ccw(pt a, pt b, pt c, bool include_collinear) {
17     int o = orientation(a, b, c);
18     return o > 0 || (include_collinear && o == 0);
19 }
20
21 void convex_hull(vector<pt>& a, bool include_collinear = false) {
22     if (a.size() == 1)
23         return;
24
25     sort(a.begin(), a.end(), [](pt a, pt b) {
26         return make_pair(a.x, a.y) < make_pair(b.x, b.y);
27     });
28     pt p1 = a[0], p2 = a.back();
29     vector<pt> up, down;
30     up.push_back(p1);
31     down.push_back(p1);
32     for (int i = 1; i < (int)a.size(); i++) {
33         if (i == a.size() - 1 || cw(p1, a[i], p2, include_collinear)) {
34             while (up.size() >= 2){
35                 if(cw(up[up.size()-2], up[up.size()-1], a[i], include_collinear)) break;
36                 up.pop_back();
37             }
38             up.push_back(a[i]);
39         }
40         if (i == a.size() - 1 || ccw(p1, a[i], p2, include_collinear)) {
41             while (down.size() >= 2){
42                 if(ccw(down[down.size()-2], down[down.size()-1], a[i], include_collinear)) break;
43                 down.pop_back();
44             }
45             down.push_back(a[i]);
46         }
47     }
48
49     if (include_collinear && up.size() == a.size()) {
```

```
50     reverse(a.begin(), a.end());
51     return;
52 }
53 a.clear();
54 for (int i = 0; i < (int)up.size(); i++)
55     a.push_back(up[i]);
56 for (int i = down.size() - 2; i > 0; i--)
57     a.push_back(down[i]);
58 }
```

5 Matemáticas

5.1 Inverso Modular

Insertar utilidad del algoritmo:

C++:

```
1 ll inv(int a){
2     int n = mod - 2;
3     ll dp[32], ans = 1;
4     dp[0] = a;
5     FOR(i, 1, 32) dp[i] = (dp[i - 1]*dp[i - 1])%mod;
6
7     FOR(i, 0, 32){
8         if(n & (1 << i)) ans = (ans*dp[i])%mod;
9     }
10
11     return ans;
12 }
```

5.2 Miller-Rabin

Insertar utilidad del algoritmo:

C++:

```
1 vector<int> a{2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37};
2
3 ll mult(ll a, ll b, ll mod) {
4     return ((__int128)a * b) % mod;
5 }
6
7 /*
8  a es la base.
9  d es la potencia.
10 n es el modulo.
11 */
12 ll pw(ll a, ll d, ll n){ // pow in log(n)
13     vector<ll> dp(63);
14     dp[0] = a;
15     ll res;
16
17     FOR(i, 1, 63) dp[i] = mult(dp[i - 1], dp[i - 1], n);
18
19     deque<int> bits;
20
21     FOR(i, 0, 63) if(d & (1ll)1 << i) bits.PB(i);
22
23     res = dp[bits.F()]%n;
24     bits.P_F();
25
26     while(!bits.empty()){
27         res = (mult(res, dp[bits.F()], n))%n;
28         bits.P_F();
29     }
```

```

29     }
30
31     return res;
32 }
33
34 bool prime(ll n){ // test de primalidad
35     ll r, x, m, d;
36     bool out;
37     r = 0;
38     m = n - 1;
39
40     while(m%2 == 0){
41         m /= 2;
42         r++;
43     }
44     d = m;
45
46     FOR(i, 0, a.size()){
47         x = pw(a[i], d, n);
48         out = false;
49         if(x == 1 or x == n - 1) continue;
50         else{
51             FOR(j, 0, r - 1){
52                 x = mult(x, x, n);
53                 if(x == n - 1){
54                     out = true;
55                     break;
56                 }
57             }
58         }
59
60         if(out) continue;
61         return false;
62     }
63     return true;
64 }

```

5.3 Pollard Rho

Insertar utilidad del algoritmo:

C++:

```

1 ll mult(ll a, ll b, ll mod) {
2     return ((__int128)a * b) % mod;
3 }
4
5 ll f(ll x, ll c, ll mod) {
6     return (mult(x, x, mod) + c) % mod;
7 }
8
9 ll rho(ll n) {
10     ll c = 1, x, y, g;
11     y = x = 2;
12     g = c;
13     while(g == 1){

```

```
14     x = f(x, c, n);
15     y = f(y, c, n);
16     y = f(y, c, n);
17     g = __gcd(abs(x - y), n);
18 }
19 return g;
20 }
```