

# The Structure of Zero-Divisor Graphs

Natalia I. Córdova, Clyde Gholston, Helen A. Hauser

July 25, 2005

## Abstract

Let  $\Gamma(\mathbb{Z}_n)$  be the zero-divisor graph whose vertices are the nonzero zero-divisors of  $\mathbb{Z}_n$ , and such that two vertices  $u, v$  are adjacent if  $n$  divides  $uv$ . Here, the authors investigate the size of the maximum clique in  $\Gamma(\mathbb{Z}_n)$ . This leads to results concerning a conjecture posed by S. Hedetniemi, the core of  $\Gamma(\mathbb{Z}_n)$ , vertex colorings of  $\Gamma(\mathbb{Z}_n)$  and  $\overline{\Gamma(\mathbb{Z}_n)}$ , and values of  $n$  for which  $\overline{\Gamma(\mathbb{Z}_n)}$  is Hamiltonian. Additional work is done to determine the cases in which  $\Gamma(\mathbb{Z}_n)$  is Eulerian.

## 1 Introduction

A *graph*  $G$  is a set of vertices  $V(G)$  and a set of edges  $E(G)$  consisting of unordered pairs of distinct vertices, such that there are no loops or multiple edges. We say that two vertices  $u, v \in V(G)$  are *adjacent* if there exists an edge  $\{u, v\} \in E(G)$ . Also, we say that the edge  $\{u, v\}$  is *incident* on the vertices  $u$  and  $v$ . Given a graph  $H = (V(H), E(H))$ ,  $H$  is a *subgraph* of  $G$  if  $V(H) \subseteq V(G)$  and  $E(H) \subseteq E(G)$ .

Our work here is on zero-divisor graphs, particularly, the zero-divisor graph of  $\mathbb{Z}_n$ , denoted  $\Gamma(\mathbb{Z}_n)$ . In order to understand the graph  $\Gamma(\mathbb{Z}_n)$ , we must first give a few definitions. The ring  $\mathbb{Z}_n = \{\overline{0}, \overline{1}, \dots, \overline{n-1}\}$  is the set of all residue classes modulo  $n$ . The operations are addition mod  $n$  and multiplication mod  $n$ . From now on we will not use the bars in order to simplify notation. An example of such a ring is

$$\mathbb{Z}_{12} = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11\}.$$

An element  $x$  of a ring  $R$  is called a *zero-divisor* if there exists some nonzero element  $y \in R$  such that  $xy = 0$ . The set of zero-divisors of  $\mathbb{Z}_{12}$  is

$$Z(\mathbb{Z}_{12}) = \{0, 2, 3, 4, 6, 8, 9, 10\}.$$

For example, since  $2 \times 6 = 12$  and  $12 | (2 \times 6)$ , then 2 and 6 are both zero-divisors.

Now we can consider  $\Gamma(\mathbb{Z}_n)$ , the *zero-divisor* graph. The set of vertices of this graph is the set  $V(\Gamma(\mathbb{Z}_n)) = Z(\mathbb{Z}_n) - \{0\}$ . Also,  $\{u, v\} \in E(\Gamma(\mathbb{Z}_n))$  if and only if  $uv = 0$  in  $\mathbb{Z}_n$  or  $n | uv$  as ordinary integers. Here we provide the graph  $\Gamma(\mathbb{Z}_{12})$  as an example in Figure 1. As we can see, there are only edges between those vertices whose product is zero.

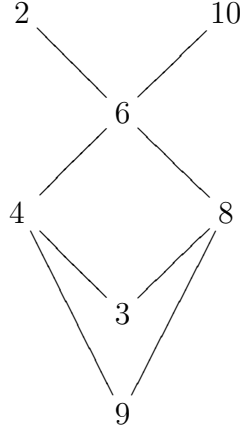


Figure 1:  $\Gamma(\mathbb{Z}_{12})$

In the following section, we provide more definitions and notation which will be used in our work. Also, for the convenience of the reader, we have included a glossary of definitions at the end of this paper.

## 2 Preliminaries

Let  $G$  be a graph. The *subgraph induced* by  $S$ , where  $S \subseteq V(G)$ , is the graph  $H$  with vertex set  $S = V(H)$  such that  $\{x, y\} \in E(H)$  if and only if  $\{x, y\} \in E(G)$  and  $x, y \in V(H)$ . A graph is called *complete* if for every  $u, v \in V(G)$  there exists an edge  $\{u, v\}$ . For example, the graph in Figure 2 is complete. When the subgraph induced by a set of vertices  $S \subseteq V(G)$  is complete, we call  $S$  a *clique*. The *clique number*  $\omega(G)$  is the size of the largest clique in a graph  $G$ .

The number of edges incident on  $v$ , or the number of vertices adjacent to  $v$ , is the *degree* of  $v$ , denoted  $\deg(v)$ . The set of vertices adjacent to  $v$  is called the *neighborhood* of  $v$ . Each member of the neighborhood is called a *neighbor* of  $v$ . In particular, in a complete graph, each vertex is a neighbor of all other vertices.

A *walk* is a sequence of vertices and edges  $v_0, e_0, v_1, e_1 \dots v_{k-1}, e_{k-1}, v_k$  where  $e_i = \{v_i, v_{i+1}\}$ . Here, edges and vertices may be repeated any number of times. A walk is

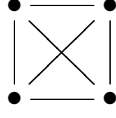


Figure 2: Complete Graph

called *closed* if the initial and terminal vertices coincide. A *path* is a walk in which no vertex is repeated, except possibly the initial and terminal vertices, and a closed path is called a *cycle*. A walk in which no edge is repeated is called a *trail*.

The *complement* of a graph  $G$  is the graph  $\overline{G}$  whose vertex set is  $V(G)$ , but  $\{x, y\} \in E(\overline{G})$  if and only if  $\{x, y\} \notin E(G)$ . For example, if the original graph  $G$  is complete, then  $\overline{G}$  is a set of vertices with no edges between them, which we call an *independent set*. In Figure 3 we show  $\overline{\Gamma(\mathbb{Z}_{12})}$  as an example.

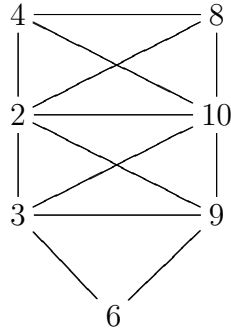


Figure 3:  $\overline{\Gamma(\mathbb{Z}_{12})}$

In this paper we deal with the vertex coloring of zero-divisor graphs. A coloring of a graph  $G$  has the property that no two adjacent vertices have the same color. Moreover, we say that  $G$  is *k-colorable* if there exists a coloring of  $G$  using only  $k$  colors. We denote  $\chi(G) = \min\{k : G \text{ is } k\text{-colorable}\}$ . For example,  $\chi(\mathbb{Z}_{12}) = 2$ .

Throughout this paper, we write  $n = p_1^{e_1} \dots p_r^{e_r}$ , to mean a factorization of  $n$ , as a product of distinct primes  $p_1 < \dots < p_r$  with positive exponents  $e_i$ .

## 2.1 Associate Classes

Let  $R$  be a ring and  $x$ ,  $y$ , and  $u$  be elements in  $R$ . An element  $u$  is called a *unit* if there exists an element  $y$  such that  $uy = 1$ . In this paper we only consider finite rings. Thus, every element is either a zero-divisor or a unit. If there exists a unit  $u$  such that  $x = uy$ , then we say that  $x$  and  $y$  are *associates*. As Lauve observed in [4],

if two elements of the vertex set  $V(\Gamma(\mathbb{Z}_n))$  are associates, then they have the same set of neighbors. Because of this, we can partition the elements of  $V(\Gamma(\mathbb{Z}_n))$  in such a way that two vertices are in the same class if and only if they are associates. This partition divides the elements of  $V(\Gamma(\mathbb{Z}_n))$  into *associate classes*. We denote by  $A_v$  the associate class represented by a vertex  $v$ . Note that the associate classes of  $\Gamma(\mathbb{Z}_n)$  and  $\overline{\Gamma(\mathbb{Z}_n)}$  are the same.

**Lemma 2.1.**  $|A_m| = \phi(n/\gcd(n, m))$

**Proof.** Let  $m$  be a zero-divisor in  $\mathbb{Z}_n$  and  $m = uv$  where  $u$  is a unit and  $v$  divides  $n$ . Then, the ideal generated by  $m$ ,  $(m)$ , is a subset of  $\mathbb{Z}_n$ . Also, since  $v = \gcd(n, m)$ , the ideal generated by  $v$ ,  $(v) = (m)$ . Therefore,  $|(m)| = |(v)| = \frac{n}{\gcd(n, m)}$ .

To find the number of associates of  $m$ , note that  $k$  is associate to  $m \Leftrightarrow (k) = (m) \Leftrightarrow \gcd(k, |(m)|) = 1$  and  $k \neq 1$ . Hence, the number of generators of  $(m)$  is equal to  $\phi(|(m)|)$ , which is equal to  $\phi(n/\gcd(n, m))$ .

□

Phillips, et al. [5], also worked with associate classes and proved that every associate class of zero divisors of  $\mathbb{Z}_n$  is either a clique or an independent set in  $\Gamma(\mathbb{Z}_n)$ . We use this result in our paper.

Also, in order to facilitate our study of  $\Gamma(\mathbb{Z}_n)$ , we sometimes use the notion of structure graphs defined by Brickel in [1]. A *structure graph* is a graph whose vertices are associate classes of zero divisors in  $\mathbb{Z}_n$ . Here, two vertices  $X$  and  $Y$  are adjacent if and only if  $xy = 0$  for every  $x \in X$  and every  $y \in Y$ .

### 3 Eulerian Graphs

A graph is *Eulerian* if it contains a closed trail containing every edge. In this section, we determine for which values of  $n$  is  $\Gamma(\mathbb{Z}_n)$  Eulerian, using a result from Euler and Hierholzer that states that a graph is Eulerian if and only if every vertex in  $V(\Gamma(\mathbb{Z}_n))$  has even degree. An example of an Eulerian graph is shown below in Figure 4.

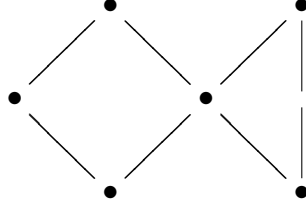


Figure 4: An Eulerian graph

**Theorem 3.1.** *Let  $n = p_1^{e_1} \dots p_r^{e_r}$ . The graph  $\Gamma(\mathbb{Z}_n)$  is Eulerian if and only if:*  
(i)  $n$  is odd and square-free or  
(ii)  $n = 4$ .

**Proof.** Take an arbitrary zero-divisor of  $n$ , say  $v = p_1^{f_1} p_2^{f_2} \dots p_r^{f_r}$  where  $0 \leq f_i \leq e_i$ . From [5], we know

$$\deg(v) = \begin{cases} \gcd(n, v) - 1, & \text{if } v^2 \neq 0; \\ \gcd(n, v) - 2, & \text{if } v^2 = 0. \end{cases} \quad (1)$$

To prove that a graph  $\Gamma(\mathbb{Z}_n)$  is Eulerian, we need to show that there exists no vertex  $v \in V(\mathbb{Z}_n)$  such that  $\deg(v)$  is odd. With the preceding formula we have the tools we need to prove the theorem.

**Case 1:** When  $n$  is odd

**Subcase 1a:** Let  $n = p_1^{e_1} \dots p_i^{e_i} \dots p_r^{e_r}$ , such that  $e_i \geq 2$ . In other words,  $n$  is not square-free. Let  $v = p_1^{e_1} \dots p_i^{e_i-1} \dots p_r^{e_r}$ . Then we know that  $v^2 = 0$ . This means that  $\deg(v) = \gcd(n, v) - 2$ . Since  $\gcd(n, v) = v$ , and  $v$  is odd, then it follows that  $\deg(v)$  is odd. Therefore,  $\Gamma(\mathbb{Z}_n)$  is not Eulerian.

**Subcase 1b:** Let  $n = p_1^{e_1} \dots p_i^{e_i} \dots p_r^{e_r}$ , such that  $e_1 = \dots = e_i = \dots = e_r = 1$ , and let  $v = p_1 \dots p_i^{e_i-1} \dots p_r^{e_r}$ . In this case,  $n$  is square-free, and  $v^2 \neq 0$ . This means  $\deg(v) = \gcd(n, v) - 1$ . Since  $\gcd(n, v) = v$ , and  $v$  is odd, then it follows that  $\deg(v)$  is even for every  $v$  of the form  $v = p_1 \dots p_i^{e_i-1} \dots p_r^{e_r}$ . However, every vertex of  $\Gamma(\mathbb{Z}_n)$  is associate to a vertex of this form. Hence, every vertex in  $\Gamma(\mathbb{Z}_n)$  has even degree. So, if  $n$  is odd and square-free, then  $\Gamma(\mathbb{Z}_n)$  is Eulerian.

**Case 2:** When  $n$  is even

**Subcase 2a:**  $n \neq 4$

If  $n = p_1^{e_1} \dots p_r^{e_r}$  then  $p_1 = 2$ . This means that 2 is a zero-divisor of  $\mathbb{Z}_n$ . If  $n \neq 4$ ,

$\Gamma(\mathbb{Z}_n)$  is not an Eulerian graph because, the only neighbor of 2 is  $n/2$ .

**Subcase 2b:**  $n = 4$

If  $n = 4$ ,  $\Gamma(\mathbb{Z}_n)$  consists of only one vertex. This is trivially Eulerian.

□

## 4 S. Hedetniemi's Conjecture

The *direct-product* of two graphs  $G$  and  $G'$  is the graph in which  $V(G \times G') = V(G) \times V(G')$ . That is, the vertex set is the set of ordered pairs of vertices from  $G$  and  $G'$ . The edge set of  $G \times G'$  consists of all  $\{(u, u'), (v, v')\}$ , such that,  $u, v \in V(G)$  and  $u', v' \in V(G')$ ,  $\{u, v\} \in E(G)$  and  $\{u', v'\} \in E(G')$ .

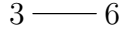


Figure 5:  $\Gamma(\mathbb{Z}_9)$

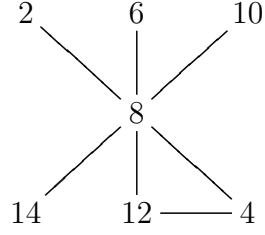


Figure 6:  $\Gamma(\mathbb{Z}_{16})$

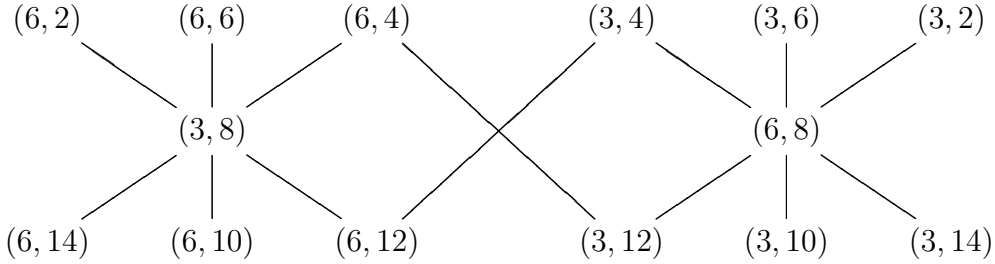


Figure 7:  $\Gamma(\mathbb{Z}_9) \times \Gamma(\mathbb{Z}_{16})$

S. Hedetniemi conjectured that for all graphs  $G$  and  $G'$ ,  $\chi(G \times G') = \min\{\chi(G), \chi(G')\}$  [3]. Here we prove Hedetniemi's conjecture for all zero-divisor graphs  $\Gamma(R)$  that satisfy  $\chi(\Gamma(R)) = \omega(\Gamma(R))$ . In particular, we prove the conjecture for  $\Gamma(\mathbb{Z}_n)$ .

**Notation.** For ease of notation, let  $\chi(\Gamma(\mathbb{Z}_n) \times \Gamma(\mathbb{Z}_m)) = \chi_{n,m}$ ,  $\chi(\Gamma(\mathbb{Z}_n)) = \chi_n$ , and  $\chi(\Gamma(\mathbb{Z}_m)) = \chi_m$ . Also, let  $\omega_n$  denote the clique number of  $\Gamma(\mathbb{Z}_n)$ . Here, the parity

of the exponents of the primes in the prime factorization of  $n$  is relevant. Therefore, we now let  $n = p_1^{e_1} \dots p_r^{e_r} q_1^{f_1} \dots q_s^{f_s}$ , where  $p_i$  and  $q_j$  are distinct primes, such that for every  $1 \leq i \leq r$ ,  $e_i$  is even, and for every  $1 \leq j \leq s$ ,  $f_j$  is odd.

**Theorem 4.1.**  $\chi_{n,m} = \min\{\chi_n, \chi_m\}$

In order to prove this theorem, we must prove  $\chi_{n,m} \leq \min\{\chi_n, \chi_m\}$ , and  $\chi_{n,m} \geq \min\{\chi_n, \chi_m\}$ . However, to prove  $\chi_{n,m} \geq \min\{\chi_n, \chi_m\}$  we need Lemma 4.2.

**Lemma 4.2.** *Let  $n = p_1^{e_1} \dots p_r^{e_r} q_1^{f_1} \dots q_s^{f_s}$ , where  $p_i$  and  $q_j$  are distinct primes, such that for every  $1 \leq i \leq r$ ,  $e_i$  is even, and for every  $1 \leq j \leq s$ ,  $f_j$  is odd. In  $\Gamma(\mathbb{Z}_n)$  there exists a clique of size*

$$\left| \bigcup_{v|n, n|v^2} A_v \right| + s = p_1^{\frac{e_1}{2}} \dots p_r^{\frac{e_r}{2}} q_1^{\frac{f_1-1}{2}} \dots q_s^{\frac{f_s-1}{2}} + s - 1.$$

**Proof.** Let  $v_0 = p_1^{\frac{e_1}{2}} \dots p_r^{\frac{e_r}{2}} q_1^{\frac{f_1+1}{2}} \dots q_s^{\frac{f_s+1}{2}}$ . Note that  $v_0^2 = 0$  and  $v_0 \in V(\Gamma(\mathbb{Z}_n))$ . Then we have,

$$\bigcup_{v|n, n|v^2} A_v = \{m p_1^{\frac{e_1}{2}} \dots p_r^{\frac{e_r}{2}} q_1^{\frac{f_1+1}{2}} \dots q_s^{\frac{f_s+1}{2}} : 1 \leq m \leq p_1^{\frac{e_1}{2}} \dots p_r^{\frac{e_r}{2}} q_1^{\frac{f_1-1}{2}} \dots q_s^{\frac{f_s-1}{2}} - 1\}.$$

In other words, the above union is the set of all the nonzero multiples of  $v_0$  that exist in  $\mathbb{Z}_n$ . Thus, we know,

$$\sum_{v|n, n|v^2} |A_v| = \left| \bigcup_{v|n, n|v^2} A_v \right| = p_1^{\frac{e_1}{2}} \dots p_r^{\frac{e_r}{2}} q_1^{\frac{f_1-1}{2}} \dots q_s^{\frac{f_s-1}{2}} - 1.$$

Hence, for each vertex  $u, w$  in the union of associate classes,  $v_0$  divides both  $u$  and  $w$ , and  $v_0^2$  divides  $uw$ . Because  $v_0^2 = 0$ , we know  $n$  also divides  $uw$ . Therefore,  $uw = 0$ .

We claim that there exists  $s$  more multiples  $x_i \in V(\Gamma(\mathbb{Z}_n))$  of  $v_0$  such that  $x_i v = 0$ , for all

$$v \in \bigcup_{v|n, n|v^2} A_v$$

where  $s$  is as in the statement of Lemma 4.2. To see this, recall  $v_0 = p_1^{\frac{e_1}{2}} \dots p_r^{\frac{e_r}{2}} q_1^{\frac{f_1+1}{2}} \dots q_i^{\frac{f_i+1}{2}} \dots q_s^{\frac{f_s+1}{2}}$ . For  $1 \leq i \leq s$ , let  $x_i = p_1^{\frac{e_1}{2}} \dots p_r^{\frac{e_r}{2}} q_1^{\frac{f_1+1}{2}} \dots q_i^{\frac{f_i-1}{2}} \dots q_s^{\frac{f_s+1}{2}}$ , so that  $x_i = v_0/q_i$ . Clearly,  $x_i$  is not in the above union. Then  $v_0 x_i = p_1^{e_1} \dots p_r^{e_r} q_1^{f_1+1} \dots q_i^{f_i} \dots q_s^{f_s+1}$ , which is a

multiple of  $n$ . Therefore, we also know that  $v_0x_i = 0$ . Hence,  $mv_0x_i = 0$  for any  $m$ . So  $x_i$  is adjacent to all elements in

$$\bigcup_{v|n, n|v_2} A_v.$$

Also, choose  $j$  such that  $i \neq j$ . Then  $x_ix_j = p_1^{e_1} \dots p_r^{e_r} q_1^{f_1+1} \dots q_i^{f_i} \dots q_j^{f_j} \dots q_s^{f_s+1}$ . Since  $n$  clearly divides  $x_ix_j$ , then  $x_ix_j = 0$ . So every  $x_i$  is adjacent to every  $x_j$ , as well as all elements in

$$\bigcup_{v|n, n|v^2} A_v.$$

The total number of  $x_i$ 's is  $s$ . Therefore, in  $\Gamma(\mathbb{Z}_n)$ , there exists a clique of size

$$\left| \bigcup_{v|n, n|v^2} A_v \right| + s = p_1^{\frac{e_1}{2}} \dots p_r^{\frac{e_r}{2}} q_1^{\frac{f_1-1}{2}} \dots q_s^{\frac{f_s-1}{2}} + s - 1.$$

□

From Beck's results (cited in [4]) we know  $\chi(\Gamma_0) = N' + s$ , where  $N' = p_1^{\frac{e_1}{2}} \dots p_r^{\frac{e_r}{2}} q_1^{\frac{f_1-1}{2}} \dots q_s^{\frac{f_s-1}{2}}$ , and  $s$  denotes the number of primes raised to an odd power in the prime factorization of  $n$ . We also know that  $\chi(G - \{v\}) = \chi(G) - 1$  if  $v$  is adjacent to every other vertex  $w \in G$ . Thus, it follows that

$$\chi_n = N' + s - 1.$$

Therefore,  $\chi_n = p_1^{\frac{e_1}{2}} \dots p_r^{\frac{e_r}{2}} q_1^{\frac{f_1-1}{2}} \dots q_s^{\frac{f_s-1}{2}} + s - 1$ , which, from Lemma 4.2, is the size of a clique in  $\Gamma(\mathbb{Z}_n)$ .

**Corollary 4.3.**  $\chi_n = \omega_n$

We now return to the proof of the theorem.

**Proof.** Assume without loss of generality that  $\chi_n \leq \chi_m$ .

I.  $\chi_{n,m} \leq \min\{\chi_n, \chi_m\}$

Fix  $f : V(\Gamma(\mathbb{Z}_n)) \rightarrow \{1, 2, \dots, \chi_n\}$ . Define  $\hat{f} : V(\Gamma(\mathbb{Z}_n) \times \Gamma(\mathbb{Z}_m)) \rightarrow \{1, 2, \dots, \chi_n\}$  by  $\hat{f}(v, u) = f(v)$ . Claim:  $\hat{f}$  is a coloring, i.e., if  $(v, u)$  and  $(\hat{v}, \hat{u})$  are adjacent in  $\Gamma(\mathbb{Z}_n) \times \Gamma(\mathbb{Z}_m)$ , then  $\hat{f}(v, u) \neq \hat{f}(\hat{v}, \hat{u})$ .

If  $(v, u)$  and  $(\hat{v}, \hat{u})$  are adjacent in  $\Gamma(\mathbb{Z}_n) \times \Gamma(\mathbb{Z}_m)$ , then  $v$  and  $\hat{v}$  are adjacent in  $\Gamma(\mathbb{Z}_n)$ . It follows that, since  $f$  is a coloring,  $f(v) \neq f(\hat{v})$ . Hence,  $\hat{f}(v, u) \neq \hat{f}(\hat{v}, \hat{u})$ .

II.  $\chi_{n,m} \geq \min\{\chi_n, \chi_m\}$

We know by Corollary 4.3 that in  $\mathbb{Z}_n$  there exists a clique  $C_n$  of size  $k = \chi_n$ , and in



$\mathbb{Z}_m$  there exists a clique  $C_m$  of size  $j = \chi_m$ . Let  $V(C_n) = \{v_1, \dots, v_k\}$ , and  $V(C_m) = \{u_1, \dots, u_j\}$ . Consider the subgraph of  $\Gamma(\mathbb{Z}_n) \times \Gamma(\mathbb{Z}_m)$  induced by  $\{(v_i, u_i) : 1 \leq i \leq \chi_n\}$ . Since every  $v_i$  is adjacent to every other vertex in  $V(C_n)$  and every  $u_i$  is adjacent to every other vertex in  $V(C_m)$ , the subgraph induced by  $\{(v_i, u_i) : 1 \leq i \leq \chi_n\}$  forms a clique in  $\Gamma(\mathbb{Z}_n) \times \Gamma(\mathbb{Z}_m)$  of size  $\chi_n$ . This means that we need at least  $\chi_n$  colors to color  $\Gamma(\mathbb{Z}_n) \times \Gamma(\mathbb{Z}_m)$ .

By I. and II.,  $\chi_{n,m} = \min\{\chi_n, \chi_m\}$ .

□

**Theorem 4.4.** *Let  $R$  and  $R'$  be any rings such that  $\chi(\Gamma(R)) = \omega(\Gamma(R))$  and  $\chi(\Gamma(R')) = \omega(\Gamma(R'))$ . Then,  $\chi(\Gamma(R) \times \Gamma(R')) = \min\{\chi(\Gamma(R)), \chi(\Gamma(R'))\}$ .*

**Proof.** We prove this theorem using the same reasoning as in the proof of Theorem 4.1.

As in Theorem 4.1, we assume without loss of generality that  $\chi(\Gamma(R)) \leq \chi(\Gamma(R'))$ . Then, we prove that  $\chi(\Gamma(R) \times \Gamma(R')) \leq \min\{\chi(\Gamma(R)), \chi(\Gamma(R'))\}$  by defining a coloring for  $\Gamma(R) \times \Gamma(R')$  using only the colors necessary to color  $\Gamma(R)$ .

On the other hand, we prove that  $\chi(\Gamma(R) \times \Gamma(R')) \geq \min\{\chi(\Gamma(R)), \chi(\Gamma(R'))\}$  using the fact that  $\chi(\Gamma(R)) = \omega(\Gamma(R))$  and  $\chi(\Gamma(R')) = \omega(\Gamma(R'))$ . Let  $C_R$  and  $C_{R'}$  be the maximal cliques in  $\Gamma(R)$  and  $\Gamma(R')$  respectively. Also, let  $V(C_R) = \{v_1, \dots, v_u\}$  and  $V(C_{R'}) = \{w_1, \dots\}$ . Consider the subgraph of  $\chi(\Gamma(R) \times \Gamma(R'))$  induced by  $\{(v_i, u_i) : 1 \leq i \leq \chi(\Gamma(R))\}$ . By the same reasoning as in Theorem 4.2, the subgraph forms a clique in  $\Gamma(R) \times \Gamma(R')$  of size  $\chi(\Gamma(R))$ . This means that we need at least  $\chi(\Gamma(R))$  colors to color  $\Gamma(R) \times \Gamma(R')$ . Hence,  $\chi(\Gamma(R) \times \Gamma(R')) = \min\{\chi(\Gamma(R)), \chi(\Gamma(R'))\}$ .

□

## 5 The Core of a Graph

In this section, we determine the core of the graph  $\Gamma(\mathbb{Z}_n)$ . To understand what a core is, we must first define homomorphism and automorphism. A *homomorphism* is a mapping  $f$  from  $V(G)$  to  $V(H)$ , such that for any  $x, y \in V(G)$ , where  $\{x, y\} \in E(G)$ ,  $\{f(x), f(y)\} \in E(H)$ . An *automorphism* is a bijective homomorphism  $f$  from a graph  $G$  to itself. Two graphs  $G$  and  $H$  are called *isomorphic* if and only if there exists a one-to-one and onto mapping of vertices  $f : V(G) \longrightarrow V(H)$  which preserves adjacency.

By definition, a graph  $G$  is a *core* if any homomorphism from  $G$  to itself is an automorphism. Also, a subgraph  $H$  of  $G$  is called a core of  $G$  if  $H$  is a core itself, and there is a homomorphism from  $G$  to  $H$ . Every graph has at least one core, and in cases where a graph has more than one core, all of the cores are isomorphic [3].

**Theorem 5.1.** *The core of the graph  $\Gamma(\mathbb{Z}_n)$  is the maximal clique in  $\Gamma(\mathbb{Z}_n)$ .*

**Proof.** By Corollary 4.3 we know that  $\chi(\Gamma(\mathbb{Z}_n)) = \omega(\Gamma(\mathbb{Z}_n))$ . Let  $\chi(\Gamma(\mathbb{Z}_n)) = m$ , and  $C_m$  be a clique of size  $m$ . Also, let  $V(C_m) = \{v_1, \dots, v_m\}$  where  $v_i$  is colored  $i$ .

Let  $f : V(\Gamma(\mathbb{Z}_n)) \rightarrow V(C_m)$  be given by  $f(w) = v_i$ , where  $\chi(w) = \chi(v_i)$ . Note that  $f(v_i) = v_i$  for every  $i = 1, \dots, m$ , and hence,  $f$  is onto.

Let  $\{w, w'\} \in E(\Gamma(\mathbb{Z}_n))$ . It follows that the color of  $w$  is not equal to the color of  $w'$ , which implies that  $f(w) \neq f(w')$ . Hence,  $\{f(w), f(w')\} \in E(\Gamma(\mathbb{Z}_n))$ . Therefore,  $f$  is a homomorphism. By definition,  $C_m$  is the core of  $\Gamma(\mathbb{Z}_n)$ .

In particular, since the maximal clique of a complete graph is the graph itself, then the core of a complete graph is itself a core.

□

**Corollary 5.2.**  *$\Gamma(\mathbb{Z}_n)$  is a core if and only if  $n = p^2$ .*

**Proof.** The proof of this is rather simple. We know that  $\Gamma(\mathbb{Z}_n)$  is complete if and only if  $n = p^2$ . By Theorem 5.1, the core of  $\Gamma(\mathbb{Z}_n)$  is the largest clique. Since  $\Gamma(\mathbb{Z}_{p^2})$  is complete, it is its own largest clique, and therefore also the core. Hence,  $\Gamma(\mathbb{Z}_n)$  is a core if and only if  $n = p^2$ .

□

**Theorem 5.3.** *Let  $R$  be a ring such that  $\chi(\Gamma(R)) = \omega(\Gamma(R))$ . Then, the core of  $\Gamma(R)$  is the maximal clique in  $\Gamma(R)$ .*

**Proof.** We prove this theorem using the same reasoning we used for the proof of Theorem 5.1.

Since,  $\chi(\Gamma(R)) = \omega(\Gamma(R))$  we can let  $\chi(\Gamma(R)) = m$ , and  $C_m$  be a clique of size  $m$ . Also let,  $V(C_m) = \{v_1, \dots, v_m\}$ , and  $\chi(v_i) = i$ .

Let  $f : V(\Gamma(R)) \rightarrow V(C_m)$  be given by  $f(w) = v_i$ , where  $\chi(w) = \chi(v_i)$ . Note that  $f(v_i) = v_i$  for every  $i = 1, \dots, m$ , and hence,  $f$  is onto.

Let  $\{w, w'\} \in E(\Gamma(R))$ . It follows that  $\chi(w) \neq \chi(w')$ , which implies that  $f(w) \neq f(w')$ . Hence,  $\{f(w), f(w')\} \in E(\Gamma(R))$ . Therefore,  $f$  is a homomorphism. By definition,  $C_m$  is the core of  $\Gamma(R)$ .

□

## 6 Coloring of $\overline{\Gamma(\mathbb{Z}_n)}$

Melody Brickel addressed the coloring of  $\overline{\Gamma(\mathbb{Z}_n)}$  in [1]. She did not determine a formula for the coloring of  $\overline{\Gamma(\mathbb{Z}_n)}$ ; however, she did determine that the chromatic number of  $\overline{\Gamma(\mathbb{Z}_n)}$  equals the clique number of  $\overline{\Gamma(\mathbb{Z}_n)}$  for the cases  $n = p^e$ , and  $n = p_1 p_2 p_3$ , where  $p_1 < p_2 < p_3$ . Here we include the proofs of these cases, and extend to proofs of some additional cases. From these results, we make the following conjecture:

**Conjecture 6.1.**  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$  for all  $n \geq 2$ .

**Lemma 6.2.**  $\chi(G) \geq \omega(G)$  for any graph  $G$ .

**Proof.**

Consider a clique in  $G$ . For any vertices  $v, w$  in this clique,  $v$  and  $w$  must be colored differently. Thus, we need at least  $\omega(G)$  to color  $G$ .

□

**Proposition 6.3.** Let  $n = p^e$ . Then  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ .

**Proof.** Suppose  $e$  is even and let  $A_{p^i}$  be the associate class of  $p^i$ , where  $1 \leq i \leq e$ . The largest clique in  $\overline{\Gamma(\mathbb{Z}_n)}$  is

$$\left(\bigcup_{i=1}^{\frac{e}{2}-1} A_{p^i}\right) \bigcup \{x\}$$

for some  $x \in A_{p^{\frac{e}{2}}}$ . All of these vertices have a different color. Any vertex  $v$  not in the clique is included in

$$\bigcup_{i=\frac{e}{2}}^e A_{p^i},$$

and thus  $xv = up^q$  where  $q \geq e$ , and  $u$  is a unit. Also  $v_i v_j$  is a multiple of  $p^e$  for any  $v_i, v_j$  not in the maximal clique. Therefore, all  $v$  can be colored with the color of  $x$ . Hence,  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ .

Now suppose  $e$  is odd. The largest clique is now

$$\bigcup_{i=1}^{\frac{e-1}{2}} A_{p^i}.$$

All of the vertices not in this clique form an independent set, since the product of these elements is divisible by  $p^e$ . Also, for the same reason any  $v$  not in the maximal clique cannot be adjacent to any  $y$  in  $A_{p^{\frac{e-1}{2}}}$ . Therefore, all  $v$  can be colored with the color of any  $y$ . Thus,  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ .

□

**Proposition 6.4.** *Let  $n = p_1 p_2 p_3$ . Then  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ .*

**Proof.** We know from Lemma 2.1 that  $|A_{p_1}| > |A_{p_2}| > |A_{p_3}|$ , and  $|A_{p_1 p_2}| > |A_{p_1 p_3}| > |A_{p_2 p_3}|$ .

**Case 1:** The largest clique is  $A_{p_1} \cup A_{p_2} \cup A_{p_3}$ . Since  $A_{p_1}$  is not adjacent to  $A_{p_2 p_3}$  in the structure graph, the colors of  $A_{p_1}$  can be used to color  $A_{p_2 p_3}$ . The same strategy can be used for the coloring of  $A_{p_1 p_2}$  and  $A_{p_1 p_3}$ . Thus,  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) \leq \omega(\overline{\Gamma(\mathbb{Z}_n)})$ . So,  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ .

**Case 2:** Suppose  $|A_{p_3}| < |A_{p_1 p_2}|$ . The largest clique is now  $A_{p_1} \cup A_{p_2} \cup A_{p_1 p_2}$ . Color each of the vertices of the clique. The vertices of  $A_{p_3}$  can be colored with the colors of  $A_{p_1 p_2}$ . Similarly,  $A_{p_2 p_3}$  can be colored with the colors of  $A_{p_1}$ , and  $A_{p_1 p_3}$  can be colored with the colors of  $A_{p_2}$ . So,  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) \leq \omega(\overline{\Gamma(\mathbb{Z}_n)})$ , and  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ . Thus, the proposition is proved.

□

**Proposition 6.5.** *Let  $n = p_1 p_2$ , where  $p_1 < p_2$ . Then  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ .*

**Proof.** Consider the associate classes of zero-divisors, and note that there are only two:  $A_{p_1}$  and  $A_{p_2}$ . In  $\overline{\Gamma(\mathbb{Z}_n)}$ , each class forms a disjoint clique. From Lemma 2.1, and since  $p_1 < p_2$ , we know  $p_2 - 1 = |A_{p_1}| > |A_{p_2}| = p_1 - 1$ . Hence  $A_{p_1}$  forms the larger clique. All of the vertices of  $A_{p_1}$  must have a different color. These colors can be used to color all of the vertices of  $A_{p_2}$ . Thus,  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) \leq \omega(\overline{\Gamma(\mathbb{Z}_n)})$ . So,  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ .

□

**Proposition 6.6.** *Let  $n = p_1^2 p_2$ , where  $p_1 < p_2$ . Then  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ .*

**Proof.** Again consider the associate classes of zero-divisors. In  $\overline{\Gamma(\mathbb{Z}_n)}$ ,  $A_{p_1^2}$  is adjacent to  $A_{p_1}$ , which is also adjacent to  $A_{p_2}$ , which is also adjacent to  $A_{p_1 p_2}$  in the structure graph. The sizes of the classes determined using Lemma 2.1 are as follows:

$$|A_{p_1}| = p_1 p_2 - p_1 - p_2 + 1$$

$$|A_{p_2}| = p_1^2 - p_1$$

$$|A_{p_1^2}| = p_2 - 1$$

$$|A_{p_1 p_2}| = p_1 - 1.$$

**Case 1:**  $|A_{p_1^2}| < |A_{p_2}|$  Then we have,  $|A_{p_1 p_2}| < |A_{p_1^2}| < |A_{p_2}| < |A_{p_1}|$ . Thus, the largest clique is  $A_{p_1} \cup A_{p_2}$ . Color each vertex of the largest clique a different color. Then color all of  $A_{p_1 p_2}$  with one color of  $A_{p_1}$ , and use colors of  $A_{p_2}$  to color  $A_{p_1^2}$ . Therefore, we have  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) \leq \omega(\overline{\Gamma(\mathbb{Z}_n)})$ . Hence,  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ .

**Case 2:**  $|A_{p_1^2}| > |A_{p_2}|$  Then we have,  $|A_{p_1 p_2}| < |A_{p_2}| < |A_{p_1^2}| < |A_{p_1}|$ . So the largest clique is now  $A_{p_1} \cup A_{p_1^2}$ . Color each vertex of the largest clique a different color. Then color all of  $A_{p_1 p_2}$  with one color of  $A_{p_1}$ , and use colors of  $A_{p_1^2}$  to color  $A_{p_2}$ . Therefore,  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ .

□

**Proposition 6.7.** *Let  $n = p_1 p_2^2$ , where  $p_1 < p_2$ . Then  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ .*

**Proof.** This proof follows the same logic as the previous one, so we will be brief. Again, the largest clique is formed by  $A_{p_1} \cup A_{p_2}$ . All of  $A_{p_1 p_2}$  can again be colored by one color, this time from  $A_{p_2}$ , and  $A_{p_2^2}$  can be colored with colors of  $A_{p_1}$ . Thus,  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ .

□

**Proposition 6.8.** *Let  $n = p_1^2 p_2^2$ . Then  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ .*

**Proof.** Consider the associate classes of zero-divisors. Using the formula in Lemma 2.1, the sizes of the classes are as follows:

$$|A_{p_1}| = p_1 p_2^2 - p_1 p_2 - p_2^2 + p_2$$

$$|A_{p_2}| = p_1^2 p_2 - p_1 p_2 - p_1^2 + p_1$$

$$|A_{p_1 p_2}| = p_1 p_2 - p_1 - p_2 + 1$$

$$|A_{p_1^2}| = p_2^2 - p_2$$

$$|A_{p_2^2}| = p_1^2 - p_1$$

$$|A_{p_1 p_2^2}| = p_1 - 1$$

$$|A_{p_1^2 p_2}| = p_2 - 1.$$

Thus, we have  $|A_{p_1 p_2^2}| < |A_{p_1^2 p_2}| < |A_{p_2^2}| < |A_{p_1 p_2}| < |A_{p_1^2}| < |A_{p_2}| < |A_{p_1}|$ . The classes that form cliques in  $\overline{\Gamma(\mathbb{Z}_n)}$  are  $A_{p_1}$ ,  $A_{p_2}$ ,  $A_{p_1^2}$ , and  $A_{p_2^2}$ , however  $A_{p_1^2}$  and  $A_{p_2^2}$  are not adjacent in the complement. Also, all of the members of  $A_{p_1 p_2}$  are adjacent to all of the classes that form cliques. Thus, the largest clique in  $\overline{\Gamma(\mathbb{Z}_n)}$  is  $A_{p_1} \cup A_{p_2} \cup A_{p_1^2} \cup \{v\}$ , where  $v \in A_{p_1 p_2}$ .

Color each member of the clique a different color. All of the members of  $A_{p_1 p_2}$  can be the same color since the members of  $A_{p_1 p_2}$  form an independent set. Since  $A_{p_1^2}$  and  $A_{p_2^2}$  are not adjacent in  $\overline{\Gamma(\mathbb{Z}_n)}$ ,  $A_{p_2^2}$  can be colored with colors of  $A_{p_1^2}$ . Also,  $A_{p_1 p_2^2}$ , which forms an independent set, is not adjacent to  $A_{p_1}$  and thus, can be colored with one color of  $A_{p_1}$ . By symmetry  $A_{p_1^2 p_2}$  can be colored with one color of  $A_{p_2}$ . Therefore,  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) \leq \omega(\overline{\Gamma(\mathbb{Z}_n)})$ . So,  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$

□

The proof of the following result is immediate:

**Corollary 6.9.** *If  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) = \omega(\overline{\Gamma(\mathbb{Z}_n)})$ , and  $\chi(\overline{\Gamma(\mathbb{Z}_m)}) = \omega(\overline{\Gamma(\mathbb{Z}_m)})$ , then  $\chi(\overline{\Gamma(\mathbb{Z}_n)}) \times \chi(\overline{\Gamma(\mathbb{Z}_m)}) = \min\{\chi(\overline{\Gamma(\mathbb{Z}_n)}), \chi(\overline{\Gamma(\mathbb{Z}_m)})\}$ .*

## 7 Hamiltonian Graphs

A graph is *Hamiltonian* if it has a cycle which contains every vertex. Figure 3 is an example of a Hamiltonian graph. In this section we ask the question: For which values of  $n$  is  $\overline{\Gamma(\mathbb{Z}_n)}$  Hamiltonian? M. Brickel [1] proved by induction that  $\overline{\Gamma(\mathbb{Z}_n)}$  is Hamiltonian when  $n = p_1 \dots p_r$ , where  $r \geq 3$ . In this section, we extend this result to the cases where  $n = p_1^{e_1} p_2^{e_2}$ , and  $n = p_1^2 p_2 \dots p_r$ .

A subset  $S \subseteq V(G)$  is an *independent set* in a graph  $G$  if the subgraph induced by  $S$  has no edges. Moreover, the *independence number*  $\alpha(G)$  is the number of vertices in the largest independent set of  $G$ . On the other hand, the *connectivity*  $\kappa(G)$  of a graph  $G$  is the size of the smallest subset  $S \subseteq V(G)$  such that  $G - S$  is disconnected.

In this section we also refer to the vertex of *minimum degree*  $\delta(\Gamma(\mathbb{Z}_n))$ . We know from Phillips, et al. [5] that in  $\overline{\Gamma(\mathbb{Z}_n)}$  the vertex of minimum degree has degree  $\delta(G) = p_1^{e_1-1}(p_1 - 1)p_2^{e_2} \dots p_r^{e_r} - p_1^{e_1-1}(p_1 - 1) \dots p_r^{e_r-1}(p_r - 1)$ . Also, since the largest independent set in  $\overline{\Gamma(\mathbb{Z}_n)}$  is equal to the largest clique in  $\Gamma(\mathbb{Z}_n)$ , we know from Lemma 4.2 that  $\alpha(\overline{\Gamma(\mathbb{Z}_n)}) = p_1^{\frac{e_1}{2}} \dots p_r^{\frac{e_r}{2}} q_1^{\frac{f_1-1}{2}} \dots q_s^{\frac{f_s-1}{2}} + s - 1$ .

**Theorem 7.1.** *If  $n = p_1^{e_1} p_2^{e_2}$ , then  $\overline{\Gamma(\mathbb{Z}_n)}$  is Hamiltonian.*

**Proof.** To prove this theorem we will use a result from Diestel [2] that states that, for any graph  $G$ , if  $\kappa(G) \geq \alpha(G)$ , then  $G$  is Hamiltonian, for any graph  $G$ .

**Case 1:**  $e_1 > 1$

We know from Phillips, et al. [5] that when  $n = p_1^{e_1} p_2^{e_2}$  and  $e_1 > 1$ , then  $\kappa(\Gamma(\mathbb{Z}_n)) = \delta(\Gamma(\mathbb{Z}_n))$ . Hence, for this case, we only need to show that  $\delta(\overline{\Gamma(\mathbb{Z}_n)}) \geq \alpha(\overline{\Gamma(\mathbb{Z}_n)})$ .

*Subcase 1:*  $e_1, e_2$  even

Here we need to show that

$$p_1^{e_1-1}(p_1-1)p_2^{e_2} - p_1^{e_1-1}(p_1-1)p_2^{e_2-1}(p_2-1) \geq p_1^{\frac{e_1}{2}} p_2^{\frac{e_2}{2}} - 1.$$

Simplifying the left hand side we obtain,

$$(p_1^{e_1-1})(p_2^{e_2-1})(p_1-1).$$

But,

$$(p_1^{e_1-1})(p_2^{e_2-1})(p_1-1) \geq (p_1^{e_1-1})(p_2^{e_2-1}) \geq p_1^{\frac{e_1}{2}} p_2^{\frac{e_2}{2}} \geq p_1^{\frac{e_1}{2}} p_2^{\frac{e_2}{2}} - 1.$$

*Subcase 2:*  $e_1, e_2$  odd

For this case we need to show that

$$p_1^{e_1-1}(p_1-1)p_2^{e_2} - p_1^{e_1-1}(p_1-1)p_2^{e_2-1}(p_2-1) \geq p_1^{\frac{e_1-1}{2}} p_2^{\frac{e_2-1}{2}} + 1.$$

As in Subcase 1, we have,

$$p_1^{e_1-1}(p_1-1)p_2^{e_2} - p_1^{e_1-1}(p_1-1)p_2^{e_2-1}(p_2-1) = (p_1^{e_1-1})(p_2^{e_2-1})(p_1-1) \geq (p_1^{e_1-1})(p_2^{e_2-1}).$$

$$\text{Claim: } (p_1^{e_1-1})(p_2^{e_2-1}) \geq p_1^{\frac{e_1-1}{2}} p_2^{\frac{e_2-1}{2}} + 1.$$

This is equivalent to,  $(p_1^{e_1-1})(p_2^{e_2-1}) - p_1^{\frac{e_1-1}{2}} p_2^{\frac{e_2-1}{2}} \geq 1$ .

Factoring we obtain,  $p_1^{\frac{e_1}{2}} p_2^{\frac{e_2}{2}} (p_1^{\frac{e_1}{2}} p_2^{\frac{e_2}{2}} - 1) \geq 1$ , which is a true statement since  $e_1 \geq 3$  and  $p_1 \geq 2$ .

*Subcase 3:*  $e_1$  odd, and  $e_2$  even

Here we use the same reasoning as in Subcase 1 to show that

$$p_1^{e_1-1}(p_1-1)p_2^{e_2} - p_1^{e_1-1}(p_1-1)p_2^{e_2-1}(p_2-1) \geq p_1^{\frac{e_1}{2}} p_2^{\frac{e_2-1}{2}}.$$

Note that if  $e_1$  is even, and  $e_2$  is odd we can also apply the same reasoning.

**Case 2:**  $e_1 = 1, e_2 \geq 1$

From Theorem 7.3 [5], we know that  $\kappa(G) = \delta(G) - |A_{v_1}| + 1$ , where  $v_1$  is a vertex of minimum degree. Here,  $v_1 = p_2^{e_2}$ . Therefore,  $|A_{v_1}| = |A_{p_2^{e_2}}| = p_1 - 1$ .

*Subcase 1:*  $e_2$  even

In this case,  $\alpha(\overline{\Gamma(\mathbb{Z}_n)}) = p_2^{\frac{e_2}{2}}$ . Therefore, we need to show that

$$\kappa(\overline{\Gamma(\mathbb{Z}_n)}) = (p_1-1)p_2^{e_2} - (p_1-1)p_2^{e_2-1}(p_2-1) - (p_1-1) + 1 \geq p_2^{\frac{e_2}{2}}.$$

Factoring the left hand side we obtain,

$$(p_1-1)[p_2^{e_2} - p_2^{e_2-1}(p_2-1) - 1] + 1 = (p_1-1)(p_2^{e_2-1} - 1) + 1 \geq p_2^{e_2-1} \geq p_2^{\frac{e_2}{2}}.$$

*Subcase 2:*  $e_2$  odd

Here,  $\alpha(\overline{\Gamma(\mathbb{Z}_n)}) = p_2^{\frac{e_2-1}{2}} + 1$ . So, we have to show that  
 $p_1^{e_1}(p_1 - 1)p_2^{e_2} - p_1^{e_1-1}(p_1 - 1)p_2^{e_2-1}(p_2 - 1) + 1 \geq p_2^{\frac{e_2-1}{2}} + 1$ .  
As in Subcase 1, we manipulate the left hand side and obtain,  
 $(p_1 - 1)[p_2^{e_2} - (p_1 - 1)] + 1 = (p_1 - 1)[p_2^{e_2-1}(p_1 - 1)] + 1 = (p_1 - 1)(p_2^{e_2-1} - 1) \geq$   
 $p_2^{e_2-1} - 1 = (p_2^{\frac{e_2-1}{2}} - 1)(p_2^{\frac{e_2-1}{2}} + 1) \geq p_2^{\frac{e_2-1}{2}} + 1 \geq p_2^{\frac{e_2-1}{2}}$ .

□

**Theorem 7.2.** *If  $n = p_1^2 p_2 \dots p_r$ , then  $\overline{\Gamma(\mathbb{Z}_n)}$  is Hamiltonian.*

*Basecase :* Let  $n = p_1^2 p_2$ . This is a special case of  $n = p_1^{e_1} p_2^{e_2}$  which is Hamiltonian in the previous theorem.

**Proof.** We will use induction on  $r$ . For the induction hypothesis, assume  $\overline{\Gamma(\mathbb{Z}_n)}$  is hamiltonian when  $n = p_1^2 p_2 \dots p_r$  and  $r \leq k$ . Now suppose  $n = p_1^2 p_2 \dots p_{k+1}$ . Then we can partition the vertices of  $\overline{\Gamma(\mathbb{Z}_n)}$  into two subsets  $V_1 = \{v \in \overline{\Gamma(\mathbb{Z}_n)} : p_{k+1} \mid v\}$  and  $V_2 = \{v \in \overline{\Gamma(\mathbb{Z}_n)} : p_{k+1} \nmid v\}$ . Let  $\mathbf{R}$  be the subgraph induced by  $V_1$  and  $\mathbf{H}$  be the subgraph induced by  $V_2$ . Note that  $V(\mathbf{H})$  is a clique, and thus  $\mathbf{H}$  is Hamiltonian. Similarly we will show that  $\mathbf{R}$  is Hamiltonian.

Let  $K = \overline{\Gamma(\mathbb{Z}_{p_1^2 p_2 \dots p_k})}$ . To show  $\mathbf{R}$  to be Hamiltonian we will first show that  $\mathbf{R}$  has a subgraph isomorphic to  $K$ .

Let  $\mathbf{B} = \mathbf{R} - A_{p_{k+1}}$  and let  $v \in V(\mathbf{B})$ . Then  $v = m \cdot p_{k+1}$  where  $m$  is a zero-divisor. Let  $A_m'$  be the associate class of  $m$  in  $K$ . Using the Euler  $\phi$  function we have  $|A_v| = \phi(\frac{n}{v}) = \phi(\frac{n}{mp_{k+1}}) = \phi(\frac{p_1^2 \dots p_{k+1}}{mp_{k+1}}) = \phi(\frac{p_1^2 \dots p_k}{m})$ . In the subgraph  $K$ ,  $|A_m'| = \phi(\frac{p_1^2 \dots p_k}{m})$ . Therefore  $|A_v| = |A_m'|$ .

We claim that  $A_m'$  is a clique *iff*  $A_v$  is a clique and  $A_m'$  is an independent set if  $A_v$  is an independent set. If  $A_v$  is a clique, then  $v^2 \neq 0$  in  $\mathbb{Z}_{p_1^2 p_2 \dots p_k}$ . Therefore, there exists a  $p_i, 1 \leq i \leq k+1$  such that  $p_i$  does not divide  $v^2$ . Thus there exists a  $p_i, 1 \leq i \leq k$  such that  $p_i$  does not divide  $A_m'$ . From this we know that  $A_m'$  is a clique. Similarly if  $A_v$  is an independent set, then  $v^2 = 0$  in  $\mathbb{Z}_{p_1^2 p_2 \dots p_k}$ . Thus  $p_1^2 \dots p_{k+1} \mid v^2$  and so every  $p_i, 1 \leq i \leq k+1$  divides  $v^2$ . Therefore every  $p_i$  divides  $v/p_{k+1}$ , hence  $(v/p_{k+1})^2 = 0$  in  $\mathbb{Z}_{p_1^2 p_2 \dots p_k}$ . This provides sufficient evidence that  $A_m'$  is an independent set.

From what we have just written we can conclude that the subgraph  $\mathbf{B} = \mathbf{R} - A_{p_{k+1}}$  is isomorphic to  $\overline{\Gamma(\mathbb{Z}_{p_1^2 p_2 \dots p_k})}$ . Therefore, by induction  $\mathbf{B}$  contains a Hamiltonian path  $P$ .



Let  $u$  and  $x$  be the terminal vertices of  $P$  respectively. Without loss of generality, assume  $u \notin A_{p_1 p_2 \dots p_{k+1}}$ . Also,  $A_{p_{k+1}}$  is adjacent to every vertex in  $B$  and is a clique. So  $\mathbf{R}$  has a Hamiltonian path  $P_r$  given as follows: Follow  $P$  from  $u$  to  $x$ ; then follow an edge from  $x$  to  $w$  where  $w$  is in  $A_{p_{k+1}}$ ; then follow a Hamiltonian path of  $A_{p_{k+1}}$  from  $w$  to  $y$ .

Finally construct a Hamiltonian cycle in  $\overline{\Gamma\mathbb{Z}_n}$  as follows: Recall  $u \notin A_{p_1 p_2 \dots p_{k+1}}$ . Thus  $\exists 1 \leq i \leq k$  such that  $p_i \nmid u$ . Choose  $a \in V(H)$  such that  $p_i \nmid a$  and let  $b = p_i$ . Then  $a \neq b$ ,  $a$  is adjacent to  $u$ , and  $b$  is adjacent to  $y$ . Let  $P_H$  be a Hamiltonian path in  $H$  from  $a$  to  $b$ . We take a Hamiltonian cycle by first taking  $P_H$  from  $a$  to  $b$ . Next we take an edge from  $b$  to  $y$ . Then we can follow  $P_R$  from  $y$  to  $u$ . Finally we take an edge from  $u$  to  $a$ . Thus the claim holds by induction.

## 8 Acknowledgements

For his intellectual guidance, expertise, and daily jokes, we extend our deepest thanks to Dr. Reza Akhtar of Miami University. We would also like to thank Leigh Cobbs for her enthusiasm for mathematics, willingness to help, everlasting patience, and friendship. We thank Dr. Thomas Farmer for assisting us in the editing and revision of our paper. We especially thank Dr. Vasant Waikar, Dr. Dennis Davenport, for believing in the capabilities of undergraduate students, and all others at Miami University who contributed to our learning experience this summer. Finally, we thank the National Science Foundation and the National Security Agency for their contributions which have made the Summer Undergraduate Mathematical Research Institute possible.

## A Appendix: Glossary of Terms

In all the following definitions, let  $R$  be a ring, and  $G$  be a graph with vertex set  $V(G)$  and edge set  $E(G)$ .

**adjacent:** Two vertices  $u, v \in V(G)$  are said to be *adjacent* if  $\{u, v\} \in E(G)$ .

**associates:** Two ring elements  $x$  and  $y$  are *associates* if there exists a unit  $u$  such that  $x = uy$ .

**associate class:** The *associate class* of a vertex  $v$ , denoted  $A_v$ , consists of all the associates of  $v$ .

**automorphism:** An *automorphism* of a graph is a bijective homomorphism from the graph to itself.

**clique:** A subset  $C \subseteq V(G)$  is called a *clique* if the subgraph of  $G$  induced by  $C$  is a complete graph.

**clique number:** The *clique number* of  $G$ , denoted  $\omega(G)$ , is the size of the maximal clique.

**closed walk:** A walk is called *closed* if the initial and terminal vertices coincide, and the closed path is called a *cycle*.

**complete graph:** A graph  $G$  is called *complete* if for every  $u, v \in V(G)$ ,  $u \neq v$ , there exists an edge  $\{u, v\}$ .

**complement:** The *complement* of a graph  $G$  is the graph  $\overline{G}$  whose vertex set is  $V(G)$ , but  $\{x, y\} \in E(\overline{G})$  if and only if  $\{x, y\} \notin E(G)$ .

**connectivity:** The *connectivity*  $\kappa(G)$  of a graph  $G$  is the least number  $k$  of vertices  $\{v_1, v_2, \dots, v_k\}$  such that  $G - \{v_1, v_2, \dots, v_k\}$  is disconnected.

**core:** A graph  $G$  is a *core* if any homomorphism from  $G$  to itself is an automorphism.

**degree:** For each  $v \in V(G)$ , the *degree* of  $v$  is the number of edges incident on  $v$ .

**direct-product graph:** A *direct-product graph* of two graphs  $G$  and  $G'$  is the graph where  $V(G \times G') = V(G) \times V(G')$ . Also, for any  $u, v \in V(G)$  and  $u', v' \in V(G')$ ,  $\{(u, u'), (v, v')\}$  is an edge if and only if  $\{u, v\} \in E(G)$  and  $\{u', v'\} \in E(G')$ .

**Eulerian:** A graph is called *Eulerian* if there exists a closed trail containing every edge.

**Euler phi function:** The *Euler phi function*  $\phi(n)$ , defined for  $n \geq 1$ , counts the number of  $a$ 's,  $1 \leq a \leq n - 1$  satisfying  $\gcd(a, n) = 1$ .

**graph:** A *graph* is a pair  $G = (V(G), E(G))$ , where  $V(G)$  is the vertex set and  $E(G)$  is the edge set. The edge set consists of unordered pairs of distinct elements of  $V(G)$ .

**Hamiltonian:** A graph is called *Hamiltonian* if there exists a cycle containing every vertex.

**homomorphism:** Let  $G$  and  $H$  be graphs. A mapping  $f$  from  $V(G)$  to  $V(H)$  is a *homomorphism* if  $f(x)$  and  $f(y)$  are adjacent in  $H$  whenever  $x$  and  $y$  are adjacent in  $G$ .

**incident:** An edge  $\{u, v\} \in E(G)$  is said to be *incident* to the vertices  $u, v \in V(G)$ .

**independent set:** A subset  $S \subseteq V(G)$  is called an *independent set* if the subgraph induced by  $S$  has no edges. The *independence number* of a graph  $G$  is the number of elements in the largest independent set of  $G$ . The independence number is denoted  $\alpha(G)$ .

**isomorphic** Two graphs  $G$  and  $H$  are called *isomorphic* if there exists a one-to-one and onto mapping  $f : V(G) \longrightarrow V(H)$  which preserves adjacency. If  $G$  and  $H$  are isomorphic, we write  $G \cong H$  and say there is an isomorphism between  $G$  and  $H$ .

**k-colorable:** A graph  $G$  is called *k-colorable* if there exists a vertex coloring of  $G$  using only  $k$  colors. Note:  $\chi(G) = \min\{k : G \text{ is } k\text{-colorable}\}$ .

**minimal degree:** The *minimal degree*  $\delta(G)$  of  $G$ , is defined as  $\min\{\deg(v) : v \in V\}$ .

**neighbors:** We say two vertices  $u, v \in V(G)$  are *neighbors* in a graph  $G$  if there exists an edge  $\{u, v\} \in E(G)$  that connects  $u$  and  $v$ .

**path:** A *path* is a walk in which no vertex is repeated, except possibly the initial and terminal vertices.

**subgraph:** A *subgraph* of  $G$  is a graph  $H = (V(H), E(H))$ , where  $V(H) \subseteq V(G)$  and  $E(H) \subseteq E(G)$ .

**subgraph induced:** For a graph  $G$  and a subset  $S \subseteq V(G)$ , the *subgraph induced* by  $S$  is the graph  $H$  with vertex set  $S = V(H)$  such that  $\{x, y\} \in E(H)$  if and only if  $\{x, y\} \in E(G)$  and  $x, y \in V(H)$ .

**structure graph:** A *structure graph* is a graph whose vertices are associate classes of zero divisors in  $\mathbb{Z}_n$  and where two vertices  $X$  and  $Y$  are adjacent if and only if  $xy = 0$  for some  $x \in X$  and some  $y \in Y$ .

**trail:** A *trail* is a walk in which no edge is repeated.

**unit:** An element  $x \in R$  is a *unit* if there exists  $y \in R$  such that  $x \cdot y = 1$ .

**vertex coloring:** A *vertex coloring* of  $G$  is a map  $f : V(G) \longrightarrow \mathbb{N}$  such that  $f(v) \neq f(w)$  if  $\{v, w\} \in E$ .

**walk:** A *walk* is a sequence of vertices and edges  $v_0, e_0, v_1, e_1 \dots v_{k-1}, e_{k-1}, v_k$  where  $e_i = \{v_i, v_{i+1}\}$ .

**zero-divisor:** An element  $x \in R$  is a *zero-divisor* if there exists a nonzero  $y \in R$  such that  $xy = 0$ .

**zero-divisor graph:** The *zero-divisor graph* is the graph denoted  $\Gamma(R)$ , where the vertex set consists of the zero-divisors of a ring  $R$ , and  $\{x, y\}$  is an edge if and only if  $xy = 0$  in  $R$ .

## References

- [1] M. Brickel. *Complements of Zero Divisor Graphs*. Plan B Project, Miami University, 2004.
- [2] R. Diestel. *Graph Theory*. Springer, New York, 1997.
- [3] C. Godsil and G. Royle. *Algebraic Graph Theory*. Springer, New York, 2001.
- [4] A. Lauve. *Zero-divisor graphs of finite commutative rings*. Senior Thesis, University of Oklahoma, 1999.
- [5] A. Phillips, J. Rogers, K. Tolliver, F. Worek. *Uncharted Territory of Zero Divisor Graphs and Their Complements*. Miami University, SUMSRI, 2004.

Natalia I. Córdova  
University of Puerto Rico  
San Juan, PR  
cordova.natalia@gmail.com

Clyde Gholston  
Shaw University  
Raleigh, NC  
cgholston@hotmail.com

Helen A. Hauser  
Ohio University  
Athens, OH  
hh328403@ohio.edu