

# The Structure of Zero Divisor Sum Graphs

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## Abstract

Let  $\Sigma_n$  be the graph whose vertex set is the set of non-zero zero divisors of  $\mathbb{Z}_n$  where  $vw$  is an edge if  $v + w$  is a non-zero zero divisor. We study various graph-theoretic properties of  $\Sigma_n$  including vertex degree, connectivity and cycles. Further investigation is also made into planar graphs and automorphisms of these kinds of graphs.

## 1 Preliminaries

### 1.1 Ring Structure

In order to fully understand the topic at hand it is important to review the properties of rings. A *ring* is a non-empty set  $R$  that has two binary operations, addition and multiplication, satisfying the following:

- $(\mathbf{R}, +)$  is an abelian group.
- Multiplication is associative and commutative.
- For all  $a, b, c \in \mathbf{R}$ , the distributive law,  $a(b + c) = (ab) + (ac)$  holds.

We use the ring  $\mathbb{Z}_n$ , this is the set of integers  $0, 1, \dots, (n - 1)$  under addition and multiplication modulo  $n$ . One of the most important concepts is that of a zero divisor.

**Definition 1.1.** *In a ring  $\mathbf{R}$ , a zero divisor is an element  $z \in \mathbf{R}$  such that there exists  $x \in \mathbf{R}$ ,  $x \neq 0$  where  $zx = 0$ .*

**Theorem 1.2.** *[4] In a ring  $\mathbb{Z}_n$ , the zero divisors are precisely those non-zero elements that are not relatively prime to  $n$ .*

In other words,  $k$  is a zero divisor in  $\mathbb{Z}_n$  if and only if  $\gcd(k, n) > 1$ , a fact which we will use often throughout the paper. The other important class of elements are the units.

**Definition 1.3.** A unit is an element  $u \in R$  that has a multiplicative inverse in  $R$ , i.e. there exists a  $v \in R$  such that  $uv = 1$ .

**Theorem 1.4.** [4] Let  $n > 0$ . Every element of  $\mathbb{Z}_n$  is either a zero divisor or a unit.

## 1.2 Graph Structure

We use a graph theoretic approach to study the non-zero zero divisors of the ring  $\mathbb{Z}_n$ .

**Definition 1.5.** A graph  $G$  consists of a vertex set,  $V(G)$ , an edge set  $E(G)$ , and an association to each edge,  $e \in E(G)$  of two vertices called the endpoints of  $e$ .

Two vertices are *adjacent* if they share a common edge. If  $x$  and  $y$  are adjacent in a graph we denote this  $x \sim y$ .

Two adjacent vertices are referred to as *neighbors* of each other. The set of neighbors of a vertex  $v$ , denoted  $N(v)$  is called the *neighborhood* of  $v$ . An *edge* is incident at a vertex if that vertex is one of its endpoints. The *degree* of a vertex  $v$  is the number of edges incident at it, denoted by  $\deg(v)$ .

A *uv-walk* in a graph  $G$  is a sequence of vertices in  $G$  beginning with  $u$  and ending at  $v$  such that consecutive vertices in the sequence are adjacent. A *trail* in a graph  $G$  is a walk in which no edge is traversed more than once. A *uv-path* is a walk in a graph starting at a vertex  $u$  and ending at a vertex  $v$  in which no vertex is repeated except for maybe the first and the last. A *cycle* is a closed path.

**Definition 1.6.** Let  $G$  be a graph. A subgraph of  $G$  is a graph  $H$  such that  $V(H) \subseteq V(G)$  and  $E(H) \subseteq E(G)$ .

## 1.3 Sum Graphs

**Definition 1.7.** A sum graph is a graph whose vertices are labeled with integers where vertices  $i$  and  $j$  are joined by a line if and only if the vertex  $i + j$  is in the graph.

We investigate the graph  $\Sigma_n$ . The vertex set  $V(\Sigma_n)$  consists of the non-zero zero divisors of  $\mathbb{Z}_n$ . Distinct vertices  $x, y \in V(\Sigma_n)$  are adjacent if  $x + y \in V(\Sigma_n)$ .

It is important to notice that no vertex  $v$  is adjacent to its inverse  $-v$  because this would imply that  $0$  is in the graph.

Note that when  $n$  is prime,  $V(\Sigma_n) = \emptyset$ . For our study we will never consider this case.

**Theorem 1.8.** *If  $h \mid n$ , then  $\Sigma_h$  is a subgraph of  $\Sigma_n$ .*

**Proof.**

For any  $k \in V(\Sigma_h)$ , the  $\gcd(h, k) > 1$ . Since  $h \mid n$ , then  $\gcd(n, k) > 1$ . Thus, every vertex in  $\Sigma_h$  is a vertex in  $\Sigma_n$ .

For all edges  $xy$  in  $\Sigma_h$ , we know that  $x, y, x + y \in V(\Sigma_h)$ . From the statement above, we know that  $x, y, x + y \in V(\Sigma_n)$ . Therefore,  $\Sigma_h$  is a subgraph of  $\Sigma_n$ .  $\blacktriangle$

## 2 Vertex Degrees

**Definition 2.1.** *A dominating vertex is a vertex in a graph that is adjacent to every other vertex of  $G$ .*

**Proposition 2.2.** *Let  $n = 2x$  for some integer  $x$ . Then  $x$  is a dominating vertex in  $\Sigma_n$  if and only if  $x$  is even.*

**Proof.**

( $\Leftarrow$ ) Let  $x \in V(\Sigma_n)$  be even. We want to show for all  $v \in V(\Sigma_n)$ , with  $v \neq x$ , we have  $v + x \in V(\Sigma_n)$ .

Case I: Suppose  $v$  is even, or that  $2 \mid v$ .

Since  $2 \mid x$ , then  $2 \mid v + x$ . This implies that  $\gcd(v + x, n) \geq 2 > 1$ , therefore  $v + x \in V(\Sigma_n)$ .

Case II: Suppose  $v$  is odd, or that  $2 \nmid v$ .

Since  $v$  is a zero-divisor, there exists  $w$  such that  $vw = 0 \in \mathbb{Z}_n$ , or, viewing  $v$  and  $w$  as integers,  $n \mid vw$ . Since  $2 \mid n$  and  $2 \nmid v$ , we know  $2 \mid w$ . Then

$$\begin{aligned} w(v + x) &= wv + wx \\ &= wv + w(n/2) \\ &= 0 + (w/2)n \\ &= 0 \end{aligned}$$

as we are in  $\mathbb{Z}_n$  and  $w/2$  is an integer.

Therefore, for all  $v \in V(\Sigma_n)$ , where  $v \neq x$ , it follows that  $x \sim v$ . Hence,  $x$  is a dominating vertex in  $\Sigma_n$ .

( $\Rightarrow$ ) If  $x$  is a dominating vertex, then  $x$  is even.

Suppose towards a contradiction that  $x$  is odd. Because  $x$  is a dominating vertex,  $x \sim 2$ , so  $x + 2 \in V(\Sigma_n)$ . Thus  $x + 2$  is odd and  $\gcd(x + 2, n) > 1$ . This forces  $x + 2 = x$ , since  $x$  is the only odd non-zero zero divisor. This is a contradiction because  $x + 2 \neq x$ , or  $2 \neq 0$ . Therefore  $x$  is even. Hence,  $x$  is a dominating vertex if and only if  $x$  is even. ▲

**Definition 2.3.** Let  $R$  be a ring, and  $a, b \in R$ . We say that  $a$  is associate to  $b$  if  $a = ub$  for some unit  $u$ .

**Proposition 2.4.** If  $g, g' \in V(\Sigma_n)$  and  $g$  is associate to  $g'$ , then

$$\deg(g) = \deg(g').$$

**Proof.**

We wish to show that if  $g \sim x$  then there exists  $y \in V(\Sigma_n)$  such that  $g' \sim y$ .

Suppose  $g \sim x$ . This implies that  $g + x$  is a non-zero zero divisor. Since  $g$  is associate to  $g'$ , we know  $gk = g'$  for some unit  $k$ .

Let  $y = xk$ . Observe that  $x$  is associate to  $y$ . Since  $x$  is a zero divisor, there exists non-zero  $s \in \mathbb{Z}_n$  such that  $xs = 0$ . Also observe:

$$\begin{aligned} ys &= (xk)s \\ &= (xs)k \\ &= 0 \end{aligned}$$

So  $y$  is a zero divisor. We now show that  $y$  is a non-zero zero divisor.

Assume that  $y = 0$ . Since  $y = xk$ , this implies that  $xk = 0$ . Since  $k$  is a unit, we can multiply by its inverse and end up with  $x = 0$ . This is a contradiction since  $x$  is a non-zero zero divisor. Therefore  $y \neq 0$ , and  $y \in V(\Sigma_n)$ .

Now, we show that  $g' \sim y$ , or that  $g' + y$  is a non-zero zero divisor. Observe:

$$\begin{aligned} &g' + y \\ &= gk + xk \\ &= (g + x)k, \end{aligned}$$

and since  $g + x$  is a zero divisor, there exists a non-zero  $c \in \mathbb{Z}_n$  such that

$$(g + x)c = 0.$$

and so

$$[(g+x)k]c = 0$$

Since  $kc \neq 0$ , this implies that  $g' + y$  is a non-zero zero divisor and therefore that  $g' \sim y$ .

Define the map  $\Phi: N(g) \mapsto N(g')$  by

$$\Phi(x) = xk.$$

This map is injective because if  $xk = yk$ , multiplying on the right side by  $k^{-1}$  gives

$$x = y.$$

Thus, for each neighbor of  $g$  we are able to find a unique neighbor of  $g'$ . Therefore,  $\deg(g') \geq \deg(g)$ . By symmetry,  $\deg(g) \geq \deg(g')$ . Therefore, if  $g$  is associate to  $g'$ ,

$$\deg(g) = \deg(g').$$

▲

**Definition 2.5.** *The minimum degree of a graph  $G$  is the smallest degree of all the vertices in a graph, denoted  $\delta(G)$ .*

**Proposition 2.6.**  $\Sigma_n$ ,  $\delta(\Sigma_n) = 1$  if and only if  $n = 4q$ , where  $q$  is prime.

**Proof.**

( $\Leftarrow$ ) Notice  $V(\Sigma_n) = A \cup B$  where  $A = \{2, 4, 6, \dots, 4q - 2\}$  and  $B = \{q, 2q, 3q\}$ . Each element in  $A$  is adjacent to every other element in  $A$  except its inverse. Thus  $\deg(v) > 1$  for all  $v \in A$ . Also  $q \sim 2q$ , but  $q \not\sim 3q$  (because they are inverses) and  $q \not\sim v$  for all  $v \in A \setminus \{2q\}$  because the sum of an even number and an odd number is odd. Therefore  $q$  will have degree 1.

( $\Rightarrow$ ) Let  $v \in V(\Sigma_n)$  such that  $\deg(v) = 1$ .

Case I:  $n = tv$  where  $t \geq 5$ .

Then,  $v, 2v, 3v, 4v \in V(\Sigma_n)$ . As a result,  $v \sim 2v$  and  $v \sim 3v$ . So,  $\deg(v) > 1$ , which is a contradiction.

Case II:  $n = 2v$

This implies  $v = -v$ . Since  $\deg(v) = 1$ ,  $v$  is adjacent to some vertex, say  $x$ . This means

$$\begin{aligned} v + x &\in V(\Sigma_n) \\ \Rightarrow -v - x &\in V(\Sigma_n) \\ \Rightarrow v - x &\in V(\Sigma_n) \text{ since } v = -v \\ \Rightarrow v &\sim -x \end{aligned}$$

Suppose  $x = -x$ . Then, since  $x < n$ , we have  $n = 2x$ . Thus  $v = \frac{n}{2} = x$ . This is a contradiction because  $v$  is not adjacent to itself. So  $x \neq -x$  which implies  $\deg(v) > 1$ , a contradiction. Therefore  $n \neq 2v$ .

Case III:  $n = 3v$

Subcase A: If  $2 \mid n$ , and since  $3 \mid n$ , this implies that  $6 \mid n$ . Since  $6 \mid n$ , this subcase has been reduced to Case 1.

Subcase B: If  $2 \nmid n$ , then 2 is a unit in  $\mathbb{Z}_n$ . Now  $-2v = v$  (as  $3v = 0$ ). Since  $\deg(v) = 1$ , suppose  $v \sim x$ . This implies

$$\begin{aligned} v + x &\in V(\Sigma_n) \\ \Rightarrow (-2)(v + x) &\in V(\Sigma_n) \text{ [ as } -1 \text{ and } 2 \text{ are units.]} \\ \Rightarrow v - 2x &\in V(\Sigma_n) \\ \Rightarrow v &\sim -2x \end{aligned}$$

Suppose  $x = -2x$ . Then, since  $x < n$ , we have  $x = \frac{n}{3}$  or  $x = \frac{2n}{3}$ . Thus either  $v = x$  or  $v = -x$ . This is a contradiction because  $v$  is not adjacent to itself or to its inverse. So  $x \neq -2x$  which implies  $\deg(v) > 1$ , a contradiction. Therefore  $n \neq 3v$ .

Thus we have shown that the minimum degree of  $\Sigma_n$  is 1 when  $n = 4v$ . However, we must still show  $v$  is prime. Suppose towards contradiction  $v = st$  where  $s, t > 1$ . Then  $n = 4st$ . We see

$$\begin{aligned} v &\sim s \text{ as } s(t+1) \text{ is a zero divisor and } st + s \neq n \text{ and} \\ v &\sim 2s \text{ as } s(t+2) \text{ is a zero divisor and } st + 2s \neq n. \end{aligned}$$

Thus  $\deg(v) > 1$ , which is a contradiction. So  $v$  is prime.

▲

**Proposition 2.7.** *There exists an isolated vertex,  $\delta(\Sigma_n) = 0$ , if and only if  $n = 3p$  or  $n = 2p$ , where  $p$  is prime.*

**Proof.** We will consider the two cases,  $n = 2p$  or  $n = 3p$ .

( $\Leftarrow$ ) Assume  $n = 3p$ .

The vertex set of  $\Sigma_{3p}$  is of the form:  $A \cup B$  where

$$A = \{3, 6, \dots, 3p-3\}$$

and

$$B = \{p, 2p\}$$

Notice,  $p \not\sim 2p$  because  $p + 2p \notin V(\Sigma_{3p})$ .

Let  $a \in A$ . For all these vertices,  $p \not\sim a$  because the vertex set  $A$  contains only multiples of 3. And  $p$  is not a multiple of 3. So  $p + a \notin A$  and  $2p + a \notin A$ . Also, for every  $b \in B$  we have  $p + b \neq 2p$  and  $2p + b \neq p$ , because  $b$  is not a multiple of  $p$ . Lastly,  $p \not\sim 2p$  because  $p + 2p = 3p$ , and  $3p \notin V(\Sigma_{3p})$ .

Therefore  $p$  and  $2p$  are isolated vertices in  $\Sigma_{3p}$ .

Note: If  $p=3$ , then  $n=9$ , which consists of two vertices, which are isolated.

( $\Leftarrow$ ) Assume  $n = 2p$ .

Consider  $\Sigma_n$  where  $n = 2p$  and  $p$  is prime. In this case, the vertices of  $\Sigma_n$  are of the form:

$$V = \{p, 2, 4, 6, \dots, 2p - 2\}.$$

Suppose towards contradiction that  $x \sim p$ ; therefore  $x + p \in V(\Sigma_n)$ . The vertex  $x$  corresponds to an even integer.

Because the sum of an even and an odd number is odd, and  $p$  is the only odd vertex,

$$x + p \equiv p \pmod{2p}$$

This implies that  $x = 0$ . However, the element 0 is not in the graph and therefore we have a contradiction. Hence,  $p$  is an isolated vertex of the graph.

Note: If  $p=2$ , then  $n=4$ , which is an isolated vertex.

( $\Rightarrow$ ) Case I: Suppose  $n = tp$ , where  $t \geq 4$  and  $t$  is not prime.

Let  $A = \{p, 2p, \dots, (t-1)p\}$ . Since  $t \geq 4$ , there are at least three elements in  $A$ . Let  $r$  be such that  $r \mid t$  and let  $B = \{r, 2r, \dots, (p-1)r\}$ . Since there are at least three elements in  $r$ ,  $\deg(b) > 0$  for  $b \in B$ . There are at least  $p-1$  elements in  $B$ . If  $p \geq 5$ , there are at least 3 elements in the set, and thus we satisfy our conditions. If  $p < 5$ , consider the following two cases.

Case II: Let  $p = 2$  so that  $n = 2t$ , where  $t$  is not prime.

Since  $t$  is not prime, it can be said that  $t = mp$ , where  $p$  is prime and  $m \geq 2$ . Then

by substitution,  $n = 2mp$ . Since  $m \geq 2$ , this implies that  $2m \geq 4$ , which was treated in Case 1.

Case III: Let  $p = 3$  so that  $n = 3t$ , where  $t$  is not prime.

Since  $t$  is not prime, it can be said that  $t = mp$ , where  $p$  is prime and  $m > 1$ . Then by substitution,  $n = 3mp$ . Since  $m > 1$ , this implies that  $3m \geq 3$ , which was treated in Case 1.

Case IV: Suppose  $n = tp$  where  $t$  is prime, where  $t, p > 3$ .

Notice  $V(\sigma(tp))$  can be partitioned into two sets:

$$A = \{p, 2p, \dots, (q-1)p\}$$

$$B = \{q, 2q, \dots, (p-1)q\}$$

All the elements of  $A$  are adjacent to everything except for itself and its inverse. All the elements of  $B$  are adjacent to everything except for itself and its inverse. Since  $t$  and  $p$  are primes greater than three, the degree of any element in  $A$  and  $B$  will be greater than zero.

▲

**Proposition 2.8.** *Let  $p > 3$  be prime. Then the graph  $\Sigma_{p^k}$  for  $k \in \mathbb{Z}$ , has  $p^{k-1} - 1$  vertices, each of degree  $p^{k-1} - 3$ . In particular, all vertices have even degree.*

**Proof.**

Observe that

$$V(\Sigma_{p^k}) = \{p, 2p, 3p, \dots, (p-1)p, p^2, \dots, (p^{k-1} - 1)p\}.$$

In particular there are  $p^{k-1} - 1$  vertices in the graph. Each vertex is adjacent to all others except itself and its inverse in  $\mathbb{Z}_{p^k}$ .

We show that no vertex can be its own inverse. If  $x \in V(\Sigma_{p^k})$  and  $x = -x$ , then  $2x \equiv 0 \pmod{p^k}$ . Because  $p > 3$  is prime,  $\gcd(2, p^k) = 1$ . Since 2 is a unit in  $\mathbb{Z}_{p^k}$ , we can multiply by the inverse of 2, which yields  $x \equiv 0 \pmod{p^k}$ , a contradiction.

Since there are  $p^{k-1} - 3$  neighbors of each vertex, the degree of each vertex in  $\Sigma_{p^k}$  is  $p^{k-1} - 3$ .

▲



### 3 Connectedness

Recall that  $uv$ -path is a walk in a graph from a vertex  $u$  to a vertex  $v$  in which no vertices are repeated.

**Definition 3.1.** A graph  $G$  is connected if there exists a  $uv$ -path between all pairs of distinct vertices  $u$  and  $v$  of  $G$ .

**Definition 3.2.** The distance between  $u, v \in V(G)$  is the smallest length of a  $uv$ -path in  $G$  and is denoted  $d(u, v)$ .

**Definition 3.3.** The diameter of a graph is the greatest distance between any two vertices of a connected graph  $G$  and is denoted  $\text{diam}(G)$ .

**Theorem 3.4.** The graph  $\Sigma_n$  is disconnected if and only if:

$$\begin{aligned} n &= 9 \text{ or} \\ n &= pq \text{ where } p \text{ and } q \text{ are distinct primes.} \end{aligned}$$

**Proof.**

( $\Leftarrow$ ) The graph  $\Sigma_9$  has two vertices and no edges, hence  $\Sigma_9$  is disconnected.

Now let  $n = pq$ , where  $p$  and  $q$  are distinct primes. We can partition the vertices into two sets:

$$\begin{aligned} A &= \{p, 2p, \dots, (q-1)p\} \\ B &= \{q, 2q, \dots, (p-1)q\} \end{aligned}$$

There are no edges between vertices in  $A$  and vertices in  $B$  because no vertex that is a multiple of  $q$  is adjacent to a vertex that is a multiple of  $p$ . Therefore  $\Sigma_n$  has at two connected components and so it is disconnected.

In the case where  $n = 6$ , there are three isolated vertices- and therefore disconnected.

( $\Rightarrow$ ) We now show that for every  $n$  not of this form,  $\Sigma_n$  is connected. So we considered all of the possible prime factorizations of  $n$  that are not of the form  $3^2$  or  $pq$ .

Case I: Let  $n = p^k$  where  $k > 1$  and  $n \neq 9$ . In order to show that  $\Sigma_n$  is connected, it suffices to show that there exists a  $uv$ -path between all vertices  $u$  and  $v$  in  $\Sigma_n$ .

Let  $u \in V(\Sigma_{p^k})$ . We know from Proposition 2.8 that the graph of  $\Sigma_{p^k}$  has  $p^{k-1} - 1$  vertices each of degree  $p^{k-1} - 3$ . For distinct  $u, v \in V(\Sigma_{p^k})$ , if  $u \not\sim v$ , then  $u$  and  $v$  are inverses. Therefore, there must be another vertex  $r$  such that  $u \sim r$  and  $v \sim r$ . Since inverses are unique,  $r$  cannot be the inverse of either  $u$  or  $v$ . Therefore there exists an  $uv$ -path through  $r$ . Hence,  $\Sigma_{p^k}$  is connected with diameter 2.

Case II: Let  $n = p^a q^b$  where  $p, q$  are distinct primes and  $a > 1$  or  $b > 1$ . Without loss of generality, assume that  $a > 1$ .

In this case,  $pq \not\sim pq$  and  $pq \not\sim (n - pq)$ . However,  $pq$  is adjacent to every other vertex in the graph; therefore  $pq \sim p$ . We know that  $p \sim (n - pq)$ , so  $pq, p, (n - pq)$  is a path. Therefore the graph is connected and has diameter 2.

Case III: Let  $n = p_1^{e_1} p_2^{e_2} \dots p_r^{e_r}$  for  $r \geq 3$  where the  $p_i$  are distinct primes. Let

$$a_i = n/p_i = p_1^{e_1} p_2^{e_2} \dots p_i^{e_i-1} \dots p_r^{e_r}$$

For any  $i, j \in \mathbb{Z}$ ,  $a_i \sim a_j$  because all  $a_i$  are multiples of  $p_k$ , and therefore  $a_i + a_j$  is a multiple of  $p_k$  where  $k \neq i, j$  and  $k \in \mathbb{N}$ . The vertex  $a_j$  is of the form:

$$a_j = p_1^{e_1} p_2^{e_2} \dots p_j^{e_j-1} \dots p_r^{e_r}.$$

The inverse of  $a_i$  is  $-a_i = p_1^{e_1} p_2^{e_2} \dots (p_i^{e_i} - p_i^{e_i-1}) \dots p_r^{e_r}$ , so  $a_j \neq -a_i$ , and since  $a_j$  is a multiple of  $p_k$ , it is true that  $a_i \sim a_j$  for all  $i, j$ . So,  $a_1, a_2, \dots, a_r$  form a clique and are therefore in the same component.

Now we want to show that for all  $v \in V(\Sigma_n)$ ,  $v \sim a_s$  for all  $s$  except maybe one. Suppose  $v = -a_i$  for some  $i$ . Then obviously,  $v + a_j \neq 0$  for all  $j \neq i$ . So we have shown that  $v + a_j$  is non-zero.

Now we need to show that  $\gcd(n, v) > 1$ . We know that  $v$  is a multiple of  $p_k$  for all  $k \neq i$ , and that  $a_j$  is a multiple of all  $p_k$  for  $k \neq j$ . Since  $r \geq 3$  we know we can pick  $k \neq i, j$ . Therefore  $p_k$  is a factor of  $a_j$  and  $v$  and since  $a_j$  and  $v$  are multiples of  $p_k$ , it is true that  $a_j + v$  is a multiple of  $p_k$  and hence, a zero divisor. So,  $v \sim a_k$  for all  $k \neq i$  if  $v = -a_i$ .

Now, suppose  $v$  is not of the form  $a_i$  for any  $i$ . Then there exists  $i$  such that  $\gcd(v, p_i) > 1$ , therefore  $v$  is a multiple of  $p_i$ . For all  $s \neq i$ , we know  $a_s$  is a multiple

of  $p_i$ , since  $v$  is a multiple of  $p_i$  and  $a_s$  is a multiple of  $p_i$ , it is true that  $v + a_s$  is a multiple of  $p_i$  and therefore is a zero divisor. Here, their sum is obviously not 0 because  $v$  is not the inverse of any  $a_i$ .

Therefore  $v \sim a_s$  for all  $s \neq i$ .

Finally, we need to show that for  $u, v \in V(\Sigma_n)$ , there exists a  $uv$ -path. Find  $i, j \in V(\Sigma_n)$  such that  $\gcd(p_i, u) > 1$  and  $\gcd(p_j, v) > 1$ . Since  $r > 3$ , we find  $k \neq i, j$ , therefore  $u \sim a_k$  and  $v \sim a_k$  and thus there exists a  $uv$ -path of distance 2, so  $\Sigma_n$  is connected.

Therefore, a  $\Sigma_n$  is disconnected if and only if  $n = 9$  or  $n = pq$ . ▲

**Corollary 3.5.** *If  $\Sigma_n$  is connected, its diameter is less than or equal to 2.*

**Proof.** From the proof of Proposition 3.4 we can see that  $\text{diam}(\Sigma_n) \leq 2$  when  $\Sigma_n$  is connected. ▲

**Definition 3.6.** *A  $uv$ -trail in a graph is a  $uv$ -walk in which no edge is traversed more than once.*

**Definition 3.7.** *An Eulerian graph is a graph which contains a closed trail containing every edge.*

**Theorem 3.8.** [3] *A nontrivial connected graph is Eulerian if and only if every vertex of  $G$  has even degree.*

Combining Propositions 2.8 and Case I of 3.4, we develop the following corollary.

**Corollary 3.9.** *Let  $p > 3$  be prime, then  $\Sigma_{p^k}$  is Eulerian.*

**Proof.**

From Proposition 2.8 we see that the vertices of the graph of  $\Sigma_{p^k}$  are all of even degree, and from Proposition 3.4 we see that  $\Sigma_{p^k}$  is connected. Therefore by Theorem 3.8,  $\Sigma_{p^k}$  is Eulerian. ▲

## 4 Cycles

**Theorem 4.1.** *Let  $n$  be composite, then  $\Sigma_n$  contains a cycle if and only if  $n > 9$ .*

**Proof.**

( $\Leftarrow$ ) Let  $n > 9$  be such that  $n = pm$  where  $p$  is the smallest prime that divides  $n$  and  $m \in \mathbb{N}$ . Note that this implies  $m \geq 4$ .

If  $n = 4p$ , then  $p \neq 2$  because  $n > 9$ . However, if  $p \geq 3$ , then  $2, 4, 6$  form a cycle.

If  $n = 5p$  then  $2p \not\sim 3p$ , but  $p, 2p, 4p, 3p$  form a cycle.

If  $n$  is not of the form  $4p$  or  $5p$ , then  $p, 2p, 3p$  form a cycle.

( $\Rightarrow$ ) Now we show by brute force that for  $n \leq 9$  the graph of  $\Sigma_n$  does not contain a cycle.

For  $n = 1, 2, 3, 5, 7$ ,  $n$  is prime, so  $\Sigma_n$  obviously contains no cycle because it is the empty set.

For  $n = 4$ , the graph contains only the vertex 2, so  $\Sigma_4$  does not contain a cycle either.

For  $n = 6$ , the graph contains 3 vertices, but 3 is an isolated vertex, so there is no cycle.

For  $n = 8$ , the graph contains 3 vertices, but 2 and 6 are not adjacent, so there is no cycle.

Therefore, for composite  $n$ , the graph  $\Sigma_n$  contains a cycle if and only if  $n > 9$ . ▲

**Definition 4.2.** *If a graph  $G$  has a cycle, then the girth of  $G$  is the length of the shortest cycle in  $G$ . The girth of a graph  $G$  is denoted  $gr(G)$ .*

**Corollary 4.3.** *If  $n \neq 5p$  and  $n > 9$  is composite, then  $gr(\Sigma_n) = 3$ .*

**Proof.** This is clear from the proof of Proposition 4.1. ▲

**Corollary 4.4.** *If  $n = 5p$  for  $p = 2$  or  $3$ , then  $gr(\Sigma_n) = 4$ .*

**Proof.**

Recall that  $p$  is the smallest prime that divides  $n$  where  $n > 9$ . Since  $5 \mid n$ , and  $p \leq 5$ , there are only 3 cases to consider,  $n = 10$ ,  $n = 15$ , and  $n = 25$ . In each case, inspection shows that  $\Sigma_n$  contains no 3-cycles.

If  $n = 10$  a cycle is formed by 2,4,8,6 and  $gr(\Sigma_{10}) = 4$ .

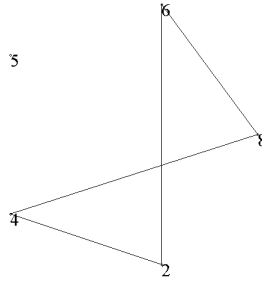


Figure 1: This is a picture of  $\Sigma_{10}$ ; note that there are no three cycles.

If  $n = 15$  a cycle is formed by 3,6,12,9 and  $gr(\Sigma_{15}) = 4$ .

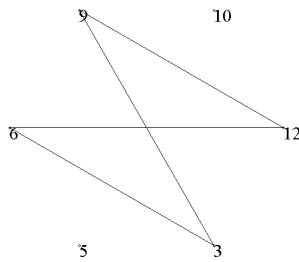


Figure 2: This is a picture of  $\Sigma_{15}$ ; note that there are no three cycles.

If  $n = 25$  a cycle is formed by 5, 10, 20, 15 and again  $gr(\Sigma_{25}) = 4$ .

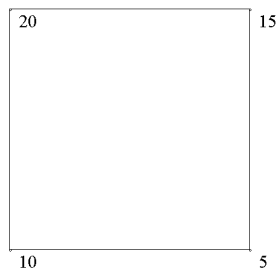


Figure 3: This is a picture of  $\Sigma_{25}$ ; note that there are no three cycles.

Thus,  $gr(\Sigma_n) = 4$  in all 3 cases.



## 5 Planarity

**Definition 5.1.** *A graph is planar if it can be drawn in the plane with no edge crossings.*

To prove our next lemma, we use the Handshaking Theorem[5]:

**Theorem 5.2.** *The sum of the degrees of the vertices of a graph is equal to twice the number of edges.*

We denote the number of vertices as  $m$  and the number of edges as  $e$  to obtain the equation

$$\sum_{v \in V} \deg(v) = 2e.$$

**Lemma 5.3.** *For all graphs,  $\delta m/2 \leq e$ .*

**Proof.**

For every  $v \in V(\Sigma_n)$ ,  $\delta \leq \deg(v)$ . Therefore,  $\delta m \leq \sum_{v \in V} \deg(v)$ . By the handshaking theorem,  $\delta m \leq 2e$ . Dividing by 2,  $\delta m/2 \leq e$ . ▲

**Definition 5.4.** *A face is a portion of a plane drawing of a planar graph which is bounded by edges and has no edge running through the interior.*

**Theorem 5.5.** [5] (Euler's Formula) *For any connected plane graph  $G$ ,  $m - e + f = 2$ , where  $f$  denotes the number of faces in  $G$ .*

**Lemma 5.6.** *If  $\Sigma_n$  is planar, then  $e \leq 3m - 6$ .*

**Proof.**

Every face in  $\Sigma_n$  has length  $\geq 3$ . Let  $F_1, F_2, \dots, F_k$  be the faces in  $\Sigma_n$  and  $l(F_i)$  be the length of the boundary of the face  $F_i$ . Because every edge  $e$  is part of the boundary of two faces,  $2e = \sum_1^k l(F_i) \geq \sum_1^k 3 = 3f$ . So  $f \leq 2e/3$ . Using Euler's formula, we know

$$\begin{aligned} f &= 2 - m + e \leq 2e/3 \\ e/3 &\leq m - 2 \\ e &\leq 3m - 6 \end{aligned}$$

▲

We use Lemma 5.3 and Lemma 5.6 to obtain the following inequalities:

$$\begin{aligned}
\delta m/2 &\leq 3m - 6 \\
\delta m &\leq 6m - 12 \\
\delta &\leq 6 - 12/m
\end{aligned}$$

**Theorem 5.7.** *Let  $p$  be prime and  $k > 1$ . If  $\Sigma_{p^k}$  is planar, then  $k \leq 3$ . Further,  $p \leq 7$  for  $k \leq 2$  and  $p = 2$  for  $k = 3$ .*

**Proof.**

By the result above,  $\delta \leq 6 - 12/m$ . Let  $i$  be any integer such that  $1 < i \leq k$  and let  $c$  be any integer such that  $1 \leq c \leq p - 1$ . Since  $p$  is prime,  $cp^i \in V(\Sigma_{p^k})$  is adjacent to every vertex in  $V(\Sigma_{p^k})$  except for itself and  $p^k - cp^i$ . Therefore,  $\delta = m - 2$ , and we have

$$\begin{aligned}
m - 2 &\leq 6 - 12/m \\
m - 8 + 12/m &\leq 0 \\
m^2 - 8m + 12 &\leq 0 \\
(m - 6)(m - 2) &\leq 0 \\
m &\leq 6
\end{aligned}$$

But we know that  $m = p^{k-1} - 1$ , so

$$\begin{aligned}
m &= p^{k-1} - 1 \leq 6 \\
p^{k-1} &\leq 7
\end{aligned}$$

If  $k = 2$ , then  $p \leq 7$ .

If  $k = 3$ , then  $p^2 \leq 7$ , and  $p = 2$ .

If  $k \geq 4$ , then  $p^{k-1} \leq 7$ , and  $p < 2$ , which is impossible. ▲

**Definition 5.8.** *A subdivision of an edge  $e = uv$  occurs when a new vertex  $w$  is placed along  $e$  and the edge  $uv$  is replaced by the path  $uwv$  of length 2.*

**Definition 5.9.** *Let  $G$  and  $G'$  be graphs.  $G$  is homeomorphic to  $G'$  if there exists a graph  $H$  such that  $G$  and  $G'$  both result from subdivisions of  $E(H)$ .*

**Definition 5.10.** *The graph  $K_{3,3}$  has six vertices. Three of the vertices  $a$ ,  $b$ , and  $c$  are all adjacent to the other three vertices  $d$ ,  $e$ , and  $f$ . Also,  $a \approx b, c$ ;  $b \approx c$ ;  $d \approx e, f$ ;  $e \approx f$ .*

**Definition 5.11.** *The graph  $K_5$  has five vertices, each of which are adjacent to the other four.*

**Theorem 5.12.** [5] (*Kuratowski's Theorem*) A graph is planar if and only if it has no subgraph homeomorphic to  $K_5$  or  $K_{3,3}$ .

**Lemma 5.13.** Every subgraph of a planar graph is planar.

**Proof.**

Suppose not. Let  $G$  be a planar graph and  $H$  be a non-planar subgraph of  $G$ . Then  $H$  contains a subgraph homeomorphic to  $K_5$  or  $K_{3,3}$ . So  $G$  itself must contain a subgraph that is homeomorphic to  $K_5$  or  $K_{3,3}$ .  $\blacktriangle$

**Definition 5.14.** Let  $G$  be a graph and  $S \subseteq V(G)$ . The subgraph induced by  $S$ , denoted  $G[S]$  is the graph defined by  $V(G[S]) = S$  and  $e \in E(G[S])$  if  $e \in E(G)$  and both endpoints of  $e$  are elements of  $S$ .

**Theorem 5.15.** Let  $p_1^{e_1}, p_2^{e_2}, \dots, p_r^{e_r}$  be primes. If  $\Sigma_{p_1^{e_1} p_2^{e_2} \dots p_r^{e_r}}$  is planar, then  $p_i \leq 7$  for every  $i$ .

**Proof.**

Let  $I$  be the subgraph of  $\Sigma_{p_1^{e_1} p_2^{e_2} \dots p_r^{e_r}}$  induced by  $p_i, 2p_i, \dots, (p_j - 1)p_i$  for some  $1 \leq i \leq r$  and  $1 \leq j \leq r, i \neq j$ . By Lemma 5.13, if  $\Sigma_{p_1^{e_1} p_2^{e_2} \dots p_r^{e_r}}$  is planar, then  $I$  must be planar. Since the vertices of  $\Sigma_{p_j^2}$  are  $\{p_j, 2p_j, \dots, (p_j - 1)p_j\}$ , the mapping  $qp_i \mapsto qp_j$  defines an isomorphism from  $I$  to  $\Sigma_{p_j^2}$ . Since  $I$  is planar, then  $\Sigma_{p_j^2}$  must be planar. Therefore, by Proposition 5.7,  $p_j \leq 7$  for all  $1 \leq j \leq r$ .  $\blacktriangle$

**Corollary 5.16.**  $\Sigma_n$  is only planar for

$$n = 4, 6, 8, 9, 10, 12, 14, 15, 21, 25, 35, 49.$$

For any planar graph, each component is isomorphic to one of the following:

2

Figure 4:  $\Sigma_4$



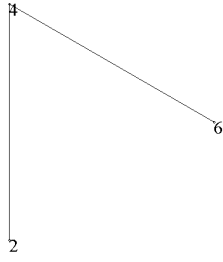


Figure 5:  $\Sigma_8$

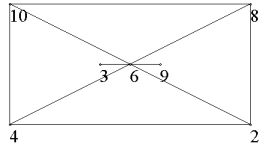


Figure 6:  $\Sigma_{12}$

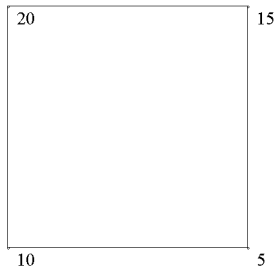


Figure 7:  $\Sigma_{25}$

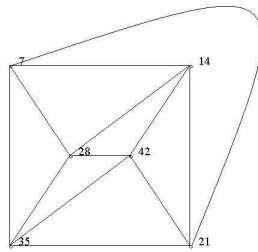


Figure 8:  $\Sigma_{49}$

**Definition 5.17.** A planar graph is outerplanar if every vertex of the graph lies on the boundary of the exterior region.

**Definition 5.18.** The graph  $K_{2,3}$  has five vertices. Two of the vertices  $a$  and  $b$  are adjacent to the three other vertices in the graph,  $c$ ,  $d$ , and  $e$ . Also,  $a \asymp b$ ,  $c \asymp d$ ,  $c \asymp e$ , and  $d \asymp e$ .

**Theorem 5.19.** Let  $G$  be a graph. If  $G$  has a subgraph homeomorphic to  $K_{2,3}$ , then  $G$  is not outerplanar.

**Proof.**

For a contradiction, assume  $G$  has a subgraph homeomorphic to  $K_{2,3}$  and  $G$  is outerplanar. Then it is possible to introduce a new vertex  $v$  inside the exterior region such that  $v$  is adjacent to  $c$ ,  $d$ , and  $e$  and the resulting graph is planar. This would imply that  $K_{3,3}$  is planar, which is not true by Kuratowski's Theorem.  $\blacktriangle$

**Proposition 5.20.** The graph  $\Sigma_n$  is outerplanar if and only if

$$n = 4, 6, 8, 9, 15, 25.$$

**Proof.**

If  $\Sigma_n$  is outerplanar, it must be planar. So we have a limited number of values for  $n$  to investigate. The graphs of  $\Sigma_4$ ,  $\Sigma_8$ , and  $\Sigma_{25}$  are clearly outerplanar. The graph of  $\Sigma_{12}$  is not outerplanar because the vertices 10 and 2 are each adjacent to 8, 4, and 6. Thus it has a subgraph isomorphic to  $K_{2,3}$ . The graph of  $\Sigma_{49}$  is not outerplanar because the vertices 7 and 42 are each adjacent to 14, 21, and 28. Thus it has a subgraph isomorphic to  $K_{2,3}$ .  $\blacktriangle$

## 6 Automorphisms

**Proposition 6.1.** Let  $s \in \mathbb{Z}_n$ . Then the map  $M_s : V(\Sigma_n) \rightarrow V(\Sigma_n)$  defined by  $M_s(x) = sx$  induces an automorphism of  $\Sigma_n$  if and only if  $s \in \mathbb{Z}_n^*$ .

**Proof.**

( $\Leftarrow$ ) Suppose  $u \in \mathbb{Z}_n$  is a unit and  $x, y$  are adjacent vertices. Then  $x + y$  is a vertex, so  $x + y \neq 0$ . Since  $u$  is a unit,

$$\begin{aligned} & ux + uy \\ &= u(x + y) \\ &\neq 0 \end{aligned}$$

Thus,  $ux \sim uy$ . This shows that the map  $M_u$  preserves edges.

We must show  $M_u$  is one-to-one. Suppose that there exists  $x, y \in V(\Sigma_n)$  such that  $ux = uy$ . Since  $u \in \mathbb{Z}_n^*$ , we have  $u^{-1}(ux) = u^{-1}(uy)$ . Therefore,  $M_u$  is one-to-one, since  $x = y$ .

We do not need to show  $M_u$  is onto because any one-to-one map from a finite set to itself is automatically onto.

Therefore  $M_u$  is a one-to-one correspondence between the vertices which preserves edges.

( $\Rightarrow$ ) To prove the contrapositive, assume  $z$  is not a unit. Thus  $z$  is a zero divisor. Now suppose  $z \in \mathbb{Z}_n$  is a zero divisor. Then there exists  $t \neq 0$ , such that  $zt = 0$ . Since  $t \in V(\Sigma_n)$ ,  $zt = 0$ , which is not a vertex. Therefore  $M_z$  is not an automorphism.

▲

The following table shows all 50 automorphisms of  $\Sigma_{16}$ . Keep in mind that 8 is a dominating vertex and is adjacent to every other vertex in the graph. Since no other vertex is adjacent to everything else, then 8 will always map to itself. For that reason we will not include in the table  $8 \mapsto 8$ .

$2 \mapsto 2$	$14 \mapsto 14$	$4 \mapsto 10$	$12 \mapsto 6$	$10 \mapsto 4$	$6 \mapsto 12$
	$14 \mapsto 14$	$4 \mapsto 10$	$12 \mapsto 6$	$10 \mapsto 12$	$6 \mapsto 4$
$2 \mapsto 4$	$14 \mapsto 12$	$4 \mapsto 2$	$12 \mapsto 14$	$10 \mapsto 10$	$6 \mapsto 6$
	$14 \mapsto 12$	$4 \mapsto 2$	$12 \mapsto 14$	$10 \mapsto 6$	$6 \mapsto 10$
	$14 \mapsto 12$	$4 \mapsto 6$	$12 \mapsto 10$	$10 \mapsto 2$	$6 \mapsto 14$
	$14 \mapsto 12$	$4 \mapsto 6$	$12 \mapsto 10$	$10 \mapsto 14$	$6 \mapsto 2$
	$14 \mapsto 12$	$4 \mapsto 10$	$12 \mapsto 6$	$10 \mapsto 2$	$6 \mapsto 14$
	$14 \mapsto 12$	$4 \mapsto 10$	$12 \mapsto 6$	$10 \mapsto 14$	$6 \mapsto 2$
	$14 \mapsto 12$	$4 \mapsto 14$	$12 \mapsto 2$	$10 \mapsto 10$	$6 \mapsto 6$
	$14 \mapsto 12$	$4 \mapsto 14$	$12 \mapsto 2$	$10 \mapsto 6$	$6 \mapsto 10$
$2 \mapsto 6$	$14 \mapsto 10$	$4 \mapsto 12$	$12 \mapsto 4$	$10 \mapsto 14$	$6 \mapsto 2$
	$14 \mapsto 10$	$4 \mapsto 12$	$12 \mapsto 4$	$10 \mapsto 2$	$6 \mapsto 14$
	$14 \mapsto 10$	$4 \mapsto 2$	$12 \mapsto 14$	$10 \mapsto 4$	$6 \mapsto 12$
	$14 \mapsto 10$	$4 \mapsto 2$	$12 \mapsto 14$	$10 \mapsto 12$	$6 \mapsto 4$
	$14 \mapsto 10$	$4 \mapsto 14$	$12 \mapsto 2$	$10 \mapsto 4$	$6 \mapsto 12$
	$14 \mapsto 10$	$4 \mapsto 14$	$12 \mapsto 2$	$10 \mapsto 12$	$6 \mapsto 4$
	$14 \mapsto 10$	$4 \mapsto 4$	$12 \mapsto 12$	$10 \mapsto 14$	$6 \mapsto 2$
	$14 \mapsto 10$	$4 \mapsto 4$	$12 \mapsto 12$	$10 \mapsto 2$	$6 \mapsto 14$
	$14 \mapsto 6$	$4 \mapsto 4$	$12 \mapsto 12$	$10 \mapsto 2$	$6 \mapsto 14$
	$14 \mapsto 6$	$4 \mapsto 4$	$12 \mapsto 12$	$10 \mapsto 14$	$6 \mapsto 2$
$2 \mapsto 10$	$14 \mapsto 6$	$4 \mapsto 2$	$12 \mapsto 14$	$10 \mapsto 12$	$6 \mapsto 4$
	$14 \mapsto 6$	$4 \mapsto 2$	$12 \mapsto 14$	$10 \mapsto 4$	$6 \mapsto 12$
	$14 \mapsto 6$	$4 \mapsto 12$	$12 \mapsto 4$	$10 \mapsto 2$	$6 \mapsto 14$
	$14 \mapsto 6$	$4 \mapsto 12$	$12 \mapsto 4$	$10 \mapsto 14$	$6 \mapsto 2$
	$14 \mapsto 6$	$4 \mapsto 14$	$12 \mapsto 2$	$10 \mapsto 4$	$6 \mapsto 12$
	$14 \mapsto 6$	$4 \mapsto 14$	$12 \mapsto 2$	$10 \mapsto 12$	$6 \mapsto 4$
	$14 \mapsto 4$	$4 \mapsto 2$	$12 \mapsto 14$	$10 \mapsto 10$	$6 \mapsto 6$
	$14 \mapsto 4$	$4 \mapsto 6$	$12 \mapsto 10$	$10 \mapsto 14$	$6 \mapsto 2$
	$14 \mapsto 4$	$4 \mapsto 6$	$12 \mapsto 10$	$10 \mapsto 2$	$6 \mapsto 14$
	$14 \mapsto 4$	$4 \mapsto 10$	$12 \mapsto 6$	$10 \mapsto 2$	$6 \mapsto 14$
$2 \mapsto 12$	$14 \mapsto 4$	$4 \mapsto 10$	$12 \mapsto 6$	$10 \mapsto 14$	$6 \mapsto 2$
	$14 \mapsto 4$	$4 \mapsto 14$	$12 \mapsto 2$	$10 \mapsto 10$	$6 \mapsto 6$
	$14 \mapsto 2$	$4 \mapsto 4$	$12 \mapsto 12$	$10 \mapsto 10$	$6 \mapsto 6$
	$14 \mapsto 2$	$4 \mapsto 4$	$12 \mapsto 12$	$10 \mapsto 6$	$6 \mapsto 10$
	$14 \mapsto 2$	$4 \mapsto 12$	$12 \mapsto 4$	$10 \mapsto 6$	$6 \mapsto 10$
	$14 \mapsto 2$	$4 \mapsto 12$	$12 \mapsto 4$	$10 \mapsto 10$	$6 \mapsto 6$
	$14 \mapsto 2$	$4 \mapsto 6$	$12 \mapsto 10$	$10 \mapsto 12$	$6 \mapsto 4$
	$14 \mapsto 2$	$4 \mapsto 6$	$12 \mapsto 10$	$10 \mapsto 4$	$6 \mapsto 12$
	$14 \mapsto 2$	$4 \mapsto 10$	$12 \mapsto 6$	$10 \mapsto 4$	$6 \mapsto 12$
	$14 \mapsto 2$	$4 \mapsto 10$	$12 \mapsto 6$	$10 \mapsto 12$	$6 \mapsto 4$

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## 8 A Appendix: Glossary of Terms

In the following definitions, let  $G = \Sigma_n$  be the graph with vertex set  $V(\Sigma_n)$  and edge set  $E(\Sigma_n)$ .

**Adjacent:** Two vertices  $u, v \in V(G)$  are said to be *adjacent* if  $uv \in E(G)$ .

**Associate:** Let  $R$  be a ring, and  $a, b \in R$ . We say that  $a$  is associate to  $b$  if  $a = ub$  for some unit  $u$ .

**Connected:** A graph  $G$  is connected if there exists a  $uv$ -path between all pairs of distinct vertices  $u$  and  $v$  of  $G$ .

**Degree:** The degree of a vertex  $v$  is the number of edges incident at it.

**Diameter:** The diameter of a graph is the greatest distance between any two vertices of a connected graph  $G$  and is denoted  $\text{diam}(G)$ .

**Distance:** The distance between  $u, v \in V(G)$  is the shortest length of a  $uv$ -path in  $G$ .

**Dominating vertex:** Any vertex,  $v \in V(\Sigma_n)$ , that is adjacent to every other vertex in  $\Sigma_n$ .

**Eulerian Graph:** A graph which has a closed trail containing every edge.

**Face:** A face is a portion of the graph which is bounded by edges and such that there is no edge running through the interior.

**Girth:** If a graph  $G$  has a cycle, then the girth of  $G$  is the length of the shortest cycle in  $G$ . The girth of a graph  $G$  is denoted  $gr(G)$ .

**Graph:** A graph  $G$  consists of a vertex set,  $V(G)$ , an edge set  $E(G)$ , and an association to each edge,  $e \in E(G)$  of two vertices, called the endpoints of  $e$ .

**Incident:** An edge is incident at a vertex if that vertex is one of its endpoints.

**Minimum degree:** The minimum degree of a graph  $G$  is the smallest degree of all

the vertices in a graph.

**Outerplanar:** A planar graph is outerplanar if every vertex of the graph lies on the boundary of the exterior region.

**Path::** A  $uv$ -path is a walk in a graph starting at a vertex  $u$  and ending at a vertex  $v$  in which no vertex is repeated

**Planar Graph:** A graph is planar if it can be drawn in the plane with no edge crossings.

**Subgraph:** Let  $\Sigma_n$  be a graph. A subgraph of  $\Sigma_n$  is a graph  $H$  such that  $V(H) \subseteq V(\Sigma_n)$  and  $E(H) \subseteq E(\Sigma_n)$ .

**Unit:** A unit is an element  $u \in R$  that has a multiplicative inverse in  $R$ , i.e. there exists a  $v \in R$  such that  $uv = 1$  and  $vu = 1$ .

**Walk:** A walk in a graph  $G$  is a sequence of vertices such that consecutive vertices in the sequence are adjacent.

**Zero divisor:** In a ring  $R$ , a zero divisor is an element  $z \in R$  for which there exists  $x \in R$ ,  $x \neq 0$  such that  $zx = 0$ .

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