

Algorithms Midterm Exam (Spring 2023)

Total : 110%

Name: _____

Student ID #: _____

Question	Score
1 (35%)	
2 (15%)	
3 (15%)	
4 (10%)	
5 (15%)	
6 (20%)	
Total	

1. **[Simple questions: 35%]** Complete the following questions with simple answers.
- (a) (5%) Suggest a best sorting algorithm when *in-place* as well as the *best* and the *average*-case performance are the top concerns.

- (b) (5%) Compare the two sorting algorithms merge sort and heapsort where both perform the best in the worst-case scenarios by giving pros and cons of the two algorithms.

- (c) (5%) Is $\Omega(f(n)) + O(f(n)) = \Theta(f(n))$ correct? Explain your answer.

(d) (5%) Is $O(n^2) + \Omega(n^3 \ln n) = \Omega(n^3 \ln n)$ correct? Explain your answer.

(e) (5%) Describe how the quadratic probing can be better than linear probing.

(f) (5%) A sequence has been stored into a hash table following the order k_1, k_2, \dots, k_n , now searching for the first element k_1 or the last element k_n takes the fastest time. Discuss the answer when the chaining strategy is used.

(g) (5%) Following the previous question, but for open addressing strategy.

2. **[Sorting: 15%]** Answer two questions regarding to the *stable* property of sorting.
- (a) (8%) Show how to detect a sorting algorithm is stable or not. That is, given an unknown sorting algorithm, can you help to write a short script to detect whether the output of the sorting algorithm produces stable outcomes.

- (b) (7%) Below is an implementation of the heapsort algorithm. Can you identify the line(s) where executing the line(s) can lead to non-stable sorting result, and why?

MAX-HEAPIFY(A, i)

```
1   $l \rightarrow \text{LEFT}(i)$ 
2   $r \rightarrow \text{RIGHT}(i)$ 
3  if  $l \leq \text{heap-size}[A]$  and  $A[l] > A[i]$ 
4  then  $\text{largest} \leftarrow l$ 
5  else  $\text{largest} \leftarrow i$ 
6  if  $r \leq \text{heap-size}[A]$  and  $A[r] > A[\text{largest}]$ 
7  then  $\text{largest} \leftarrow r$ 
8  if  $\text{largest} \neq i$ 
9    then  $\text{SWAP}(A[i], A[\text{largest}])$ 
10    $\text{MAX-HEAPIFY}(A, \text{largest})$ 
```

BUIDE-MAX-HEAP(A)

```
1   $\text{heap-size}[A] \leftarrow \text{length}[A]$ 
2  for  $i \leftarrow \lfloor \text{length}[A]/2 \rfloor$  downto 1
3    do  $\text{MAX-HEAPIFY}(A, i)$ 
```

HEAPSORT(A)

```
1   $\text{BUIDE-MAX-HEAP}(A)$ 
2  for  $i \leftarrow \text{length}[A]$  downto 2
3    do  $\text{SWAP}(A[1], A[i])$ 
4       $\text{heap-size}[A] \leftarrow \text{heap-size}[A] - 1$ 
5       $\text{MAX-HEAPIFY}(A, 1)$ 
```

3. **[Sorting: 15%]** Finish the following two tasks in $O(n)$ time.
- (a) (7%) Select the five smallest elements from an unknown sequence.

- (b) (8%) Sort n elements where the elements are in the range of 0 to $n^3 - 1$.

4. [Quicksort: 10%] The following is the quicksort implementation that we discussed in class.

QUICKSORT(A, p, r)

```
1  if  $p < r$ 
2      then  $q \leftarrow \text{PARTITION}(A, p, r)$ 
3          QUICKSORT( $A, p, q-1$ )
4          QUICKSORT( $A, q+1, r$ )
```

PARTITION(A, p, r)

```
1   $x \leftarrow A[r]$ 
2   $i \leftarrow p - 1$ 
3  for  $j \leftarrow p$  to  $r - 1$ 
4      do if  $A[j] \leq x$ 
5          then  $i \leftarrow i + 1$ 
6              SWAP( $A[i], A[j]$ )
7  SWAP( $A[i + 1], A[r]$ )
8  return  $i + 1$ 
```

Different from the previous implementation, we have now QUICKSORT(A, q, r) rather than QUICKSORT($A, q+1, r$) on line 4.

QUICKSORT(A, p, r)

```
1  if  $p < r$ 
2      then  $q \leftarrow \text{PARTITION}(A, p, r)$ 
3          QUICKSORT( $A, p, q-1$ )
4*         QUICKSORT( $A, q, r$ )
```

Can you comment on this modification from the correctness and efficiency point of view?

5. **[Recurrence equations: 15%]** Solve the following recurrence equations.

(a) (5%) $T(n) = T(n - 1) + \lg(n/2)$

(b) (5%) $T(n) = 4T\left(\frac{n}{5}\right) + n \lg n + n^2$

(c) (5%) $T(n) = 3T\left(\frac{n}{2} + 5\right) + n^2$

6. **[Hashing : 20%]** Answer the following questions related to hashing.

- (a) (7%) Demonstrate what happens when we insert the keys 17, 16, 10, 9, 8, 1, 2, 3, 15, 14, 13, 12 into a hash table of size 13 with collision resolved by double hashing. The hash functions are defined by $h_1(k) = k \bmod 13$ and $h_2(k) = 1 + (k \bmod 7)$. You do not need to finish all but up to the moment when you just finish the insertion with the second collision.

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- (b) (7%) If the function $h_1(k)$ is substituted by $h_1(k) = k \bmod 14$, what will happen?
Try to discuss as many problems as possible.

- (c) (6%) Describe any defect for the quadratic probing designed by $h(k, i) = (h'(k) + c_1 i + c_2 i^2) \bmod 24$, where $c_1 = 3$ and $c_2 = 6$.