27.0 INSTRUCTION SET SUMMARY

PIC18(L)F2X/45K50 devices incorporate the standard set of 75 PIC18 core instructions, as well as an extended set of eight new instructions, for the optimization of code that is recursive or that utilizes a software stack. The extended set is discussed later in this section.

27.1 Standard Instruction Set

The standard PIC18 instruction set adds many enhancements to the previous PIC® MCU instruction sets, while maintaining an easy migration from these PIC® MCU instruction sets. Most instructions are a single program memory word (16 bits), but there are four instructions that require two program memory locations.

Each single-word instruction is a 16-bit word divided into an opcode, which specifies the instruction type and one or more operands, which further specify the operation of the instruction.

The instruction set is highly orthogonal and is grouped into four basic categories:

- **Byte-oriented** operations
- **Bit-oriented** operations
- **Literal** operations
- **Control** operations

The PIC18 instruction set summary in Table 27-2 lists **byte-oriented**, **bit-oriented**, **literal** and **control** operations. Table 27-1 shows the opcode field descriptions.

Most **byte-oriented** instructions have three operands:

1. The file register (specified by ‘f’)
2. The destination of the result (specified by ‘d’)
3. The accessed memory (specified by ‘a’)

The file register designator ‘f’ specifies which file register is to be used by the instruction. The destination designator ‘d’ specifies where the result of the operation is to be placed. If ‘d’ is zero, the result is placed in the WREG register. If ‘d’ is one, the result is placed in the file register specified in the instruction.

All **bit-oriented** instructions have three operands:

1. The file register (specified by ‘f’)
2. The bit in the file register (specified by ‘b’)
3. The accessed memory (specified by ‘a’)

The bit field designator ‘b’ selects the number of the bit affected by the operation, while the file register designator ‘f’ represents the number of the file in which the bit is located.

The **literal** instructions may use some of the following operands:

- A literal value to be loaded into a file register (specified by ‘k’)
- The desired FSR register to load the literal value into (specified by ‘f’)
- No operand required (specified by ‘—’)

The **control** instructions may use some of the following operands:

- A program memory address (specified by ‘n’)
- The mode of the `CALL` or `RETURN` instructions (specified by ‘s’)
- The mode of the table read and table write instructions (specified by ‘m’)
- No operand required (specified by ‘—’)

All instructions are a single word, except for four double-word instructions. These instructions were made double-word to contain the required information in 32 bits. In the second word, the 4 MSbs are ‘1’s. If this second word is executed as an instruction (by itself), it will execute as a `NOP`.

All single-word instructions are executed in a single instruction cycle, unless a conditional test is true or the program counter is changed as a result of the instruction. In these cases, the execution takes two instruction cycles, with the additional instruction cycle(s) executed as a `NOP`.

The double-word instructions execute in two instruction cycles.

One instruction cycle consists of four oscillator periods. Thus, for an oscillator frequency of 4 MHz, the normal instruction execution time is 1 μs. If a conditional test is true, or the program counter is changed as a result of an instruction, the instruction execution time is 2 μs. Two-word branch instructions (if true) would take 3 μs.

Figure 27-1 shows the general formats that the instructions can have. All examples use the convention ‘nnh’ to represent a hexadecimal number.

The Instruction Set Summary, shown in Table 27-2, lists the standard instructions recognized by the Microchip Assembler (MPASM™).

Section 27.1.1 “Standard Instruction Set” provides a description of each instruction.

TABLE 27-1: OPCODE FIELD DESCRIPTIONS

Field	Description
a	RAM access bit a = 0: RAM location in Access RAM (BSR register is ignored) a = 1: RAM bank is specified by BSR register
bbb	Bit address within an 8-bit file register (0 to 7).
BSR	Bank Select Register. Used to select the current RAM bank.
C, DC, Z, OV, N	ALU Status bits: **C**arry, **D**igit **C**arry, **Z**ero, **Ov**erflow, **N**egative.
d	Destination select bit d = 0: store result in WREG d = 1: store result in file register f
dest	Destination: either the WREG register or the specified register file location.
f	8-bit Register file address (00h to FFh) or 2-bit FSR designator (0h to 3h).
f_s	12-bit Register file address (000h to FFFh). This is the source address.
f_d	12-bit Register file address (000h to FFFh). This is the destination address.
GIE	Global Interrupt Enable bit.
k	Literal field, constant data or label (may be either an 8-bit, 12-bit or a 20-bit value).
label	Label name.
mm	The mode of the TBLPTR register for the table read and table write instructions. Only used with table read and table write instructions:
*	No change to register (such as TBLPTR with table reads and writes)
*+	Post-Increment register (such as TBLPTR with table reads and writes)
*-	Post-Decrement register (such as TBLPTR with table reads and writes)
+*	Pre-Increment register (such as TBLPTR with table reads and writes)
n	The relative address (2's complement number) for relative branch instructions or the direct address for CALL/BRANCH and RETURN instructions.
PC	Program Counter.
PCL	Program Counter Low Byte.
PCH	Program Counter High Byte.
PCLATH	Program Counter High Byte Latch.
PCLATU	Program Counter Upper Byte Latch.
$\overline{\mathrm{PD}}$	Power-down bit.
PRODH	Product of Multiply High Byte.
PRODL	Product of Multiply Low Byte.
s	Fast Call/Return mode select bit s = 0: do not update into/from shadow registers s = 1: certain registers loaded into/from shadow registers (Fast mode)
TBLPTR	21-bit Table Pointer (points to a Program Memory location).
TABLAT	8-bit Table Latch.
$\overline{\mathrm{TO}}$	Time-out bit.
TOS	Top-of-Stack.
u	Unused or unchanged.
WDT	Watchdog Timer.
WREG	Working register (accumulator).
x	Don't care ('0' or '1'). The assembler will generate code with x = 0. It is the recommended form of use for compatibility with all Microchip software tools.
z_s	7-bit offset value for indirect addressing of register files (source).
z_d	7-bit offset value for indirect addressing of register files (destination).
{ }	Optional argument.
[text]	Indicates an indexed address.
(text)	The contents of text.
[expr]<n>	Specifies bit n of the register indicated by the pointer expr.
→	Assigned to.
< >	Register bit field.
∈	In the set of.
italics	User defined term (font is Courier).

FIGURE 27-1: GENERAL FORMAT FOR INSTRUCTIONS

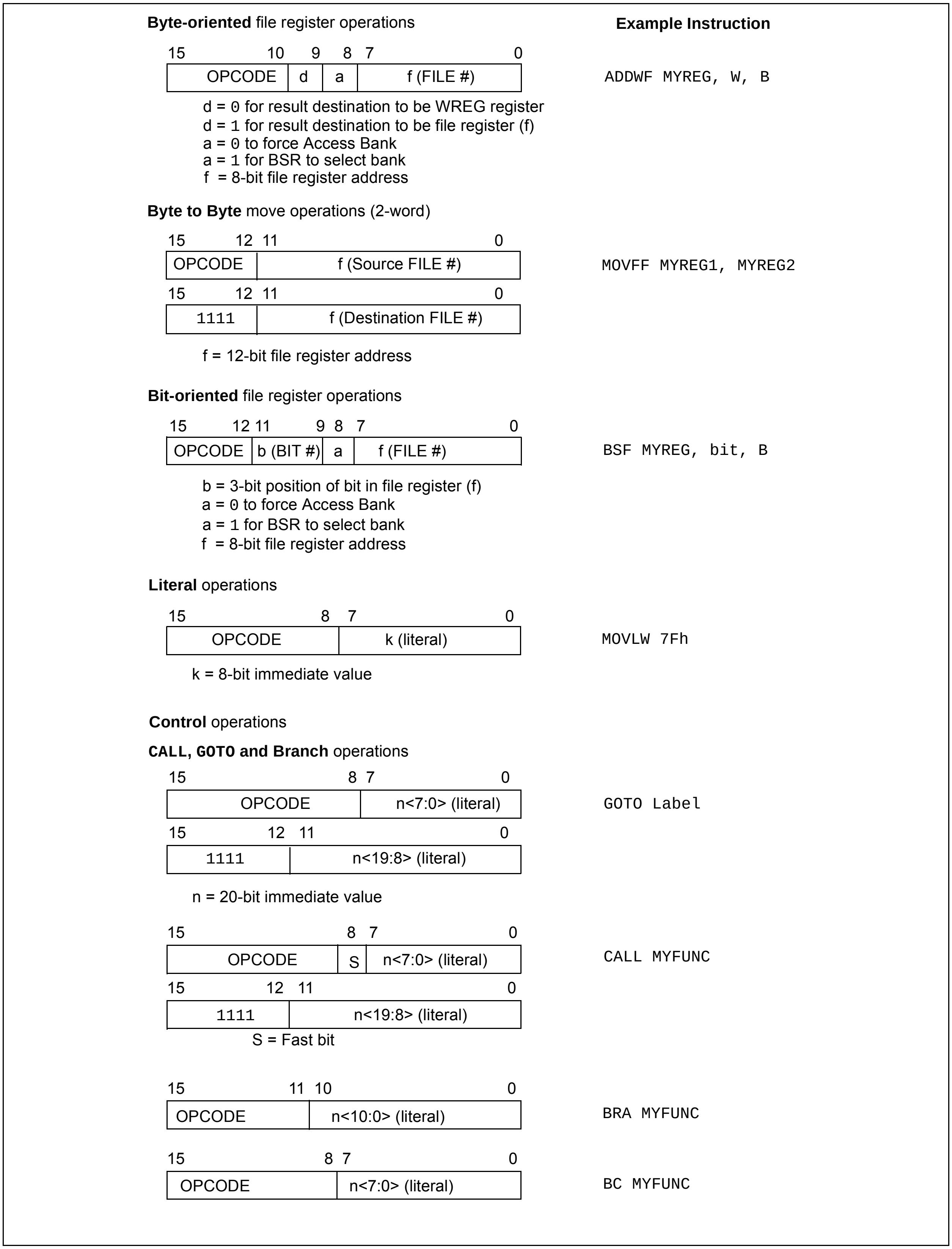

TABLE 27-2: PIC18 INSTRUCTION SET

Mnemonic, Operands		Description	Cycles	16-Bit Instruction Word				Status Affected	Notes
				MSb			LSb		
BYTE-ORIENTED OPERATIONS									
ADDWF	f, d, a	Add WREG and f	1	0010	01da	ffff	ffff	C, DC, Z, OV, N	**1, 2**
ADDWFC	f, d, a	Add WREG and CARRY bit to f	1	0010	00da	ffff	ffff	C, DC, Z, OV, N	**1, 2**
ANDWF	f, d, a	AND WREG with f	1	0001	01da	ffff	ffff	Z, N	**1,2**
CLRF	f, a	Clear f	1	0110	101a	ffff	ffff	Z	**2**
COMF	f, d, a	Complement f	1	0001	11da	ffff	ffff	Z, N	**1, 2**
CPFSEQ	f, a	Compare f with WREG, skip =	1 (2 or 3)	0110	001a	ffff	ffff	None	**4**
CPFSGT	f, a	Compare f with WREG, skip >	1 (2 or 3)	0110	010a	ffff	ffff	None	**4**
CPFSLT	f, a	Compare f with WREG, skip <	1 (2 or 3)	0110	000a	ffff	ffff	None	**1, 2**
DECF	f, d, a	Decrement f	1	0000	01da	ffff	ffff	C, DC, Z, OV, N	**1, 2, 3, 4**
DECFSZ	f, d, a	Decrement f, Skip if 0	1 (2 or 3)	0010	11da	ffff	ffff	None	**1, 2, 3, 4**
DCFSNZ	f, d, a	Decrement f, Skip if Not 0	1 (2 or 3)	0100	11da	ffff	ffff	None	**1, 2**
INCF	f, d, a	Increment f	1	0010	10da	ffff	ffff	C, DC, Z, OV, N	**1, 2, 3, 4**
INCFSZ	f, d, a	Increment f, Skip if 0	1 (2 or 3)	0011	11da	ffff	ffff	None	**4**
INFSNZ	f, d, a	Increment f, Skip if Not 0	1 (2 or 3)	0100	10da	ffff	ffff	None	**1, 2**
IORWF	f, d, a	Inclusive OR WREG with f	1	0001	00da	ffff	ffff	Z, N	**1, 2**
MOVF	f, d, a	Move f	1	0101	00da	ffff	ffff	Z, N	**1**
MOVFF	f_s, f_d	Move f_s (source) to 1st word f_d (destination) 2nd word	2	1100 1111	ffff ffff	ffff ffff	ffff ffff	None	
MOVWF	f, a	Move WREG to f	1	0110	111a	ffff	ffff	None	
MULWF	f, a	Multiply WREG with f	1	0000	001a	ffff	ffff	None	**1, 2**
NEGF	f, a	Negate f	1	0110	110a	ffff	ffff	C, DC, Z, OV, N	
RLCF	f, d, a	Rotate Left f through Carry	1	0011	01da	ffff	ffff	C, Z, N	**1, 2**
RLNCF	f, d, a	Rotate Left f (No Carry)	1	0100	01da	ffff	ffff	Z, N	
RRCF	f, d, a	Rotate Right f through Carry	1	0011	00da	ffff	ffff	C, Z, N	
RRNCF	f, d, a	Rotate Right f (No Carry)	1	0100	00da	ffff	ffff	Z, N	
SETF	f, a	Set f	1	0110	100a	ffff	ffff	None	**1, 2**
SUBFWB	f, d, a	Subtract f from WREG with borrow	1	0101	01da	ffff	ffff	C, DC, Z, OV, N	
SUBWF	f, d, a	Subtract WREG from f	1	0101	11da	ffff	ffff	C, DC, Z, OV, N	**1, 2**
SUBWFB	f, d, a	Subtract WREG from f with borrow	1	0101	10da	ffff	ffff	C, DC, Z, OV, N	
SWAPF	f, d, a	Swap nibbles in f	1	0011	10da	ffff	ffff	None	**4**
TSTFSZ	f, a	Test f, skip if 0	1 (2 or 3)	0110	011a	ffff	ffff	None	**1, 2**
XORWF	f, d, a	Exclusive OR WREG with f	1	0001	10da	ffff	ffff	Z, N	

Note 1: When a PORT register is modified as a function of itself (e.g., `MOVF PORTB, 1, 0`), the value used will be that value present on the pins themselves. For example, if the data latch is '1' for a pin configured as input and is driven low by an external device, the data will be written back with a '0'.

2: If this instruction is executed on the TMR0 register (and where applicable, 'd' = 1), the prescaler will be cleared if assigned.

3: If Program Counter (PC) is modified or a conditional test is true, the instruction requires two cycles. The second cycle is executed as a `NOP`.

4: Some instructions are two-word instructions. The second word of these instructions will be executed as a `NOP` unless the first word of the instruction retrieves the information embedded in these 16 bits. This ensures that all program memory locations have a valid instruction.

TABLE 27-2: PIC18 INSTRUCTION SET (CONTINUED)

Mnemonic, Operands		Description	Cycles	16-Bit Instruction Word				Status Affected	Notes
				MSb			LSb		
BIT-ORIENTED OPERATIONS									
BCF	f, b, a	Bit Clear f	1	1001	bbba	ffff	ffff	None	**1, 2**
BSF	f, b, a	Bit Set f	1	1000	bbba	ffff	ffff	None	**1, 2**
BTFSC	f, b, a	Bit Test f, Skip if Clear	1 (2 or 3)	1011	bbba	ffff	ffff	None	**3, 4**
BTFSS	f, b, a	Bit Test f, Skip if Set	1 (2 or 3)	1010	bbba	ffff	ffff	None	**3, 4**
BTG	f, b, a	Bit Toggle f	1	0111	bbba	ffff	ffff	None	**1, 2**
CONTROL OPERATIONS									
BC	n	Branch if Carry	1 (2)	1110	0010	nnnn	nnnn	None	
BN	n	Branch if Negative	1 (2)	1110	0110	nnnn	nnnn	None	
BNC	n	Branch if Not Carry	1 (2)	1110	0011	nnnn	nnnn	None	
BNN	n	Branch if Not Negative	1 (2)	1110	0111	nnnn	nnnn	None	
BNOV	n	Branch if Not Overflow	1 (2)	1110	0101	nnnn	nnnn	None	
BNZ	n	Branch if Not Zero	1 (2)	1110	0001	nnnn	nnnn	None	
BOV	n	Branch if Overflow	1 (2)	1110	0100	nnnn	nnnn	None	
BRA	n	Branch Unconditionally	2	1101	0nnn	nnnn	nnnn	None	
BZ	n	Branch if Zero	1 (2)	1110	0000	nnnn	nnnn	None	
CALL	k, s	Call subroutine 1st word	2	1110	110s	kkkk	kkkk	None	
		2nd word		1111	kkkk	kkkk	kkkk		
CLRWDT	—	Clear Watchdog Timer	1	0000	0000	0000	0100	$\overline{TO}$, $\overline{PD}$	
DAW	—	Decimal Adjust WREG	1	0000	0000	0000	0111	C	
GOTO	k	Go to address 1st word	2	1110	1111	kkkk	kkkk	None	
		2nd word		1111	kkkk	kkkk	kkkk		
NOP	—	No Operation	1	0000	0000	0000	0000	None	
NOP	—	No Operation	1	1111	xxxx	xxxx	xxxx	None	**4**
POP	—	Pop top of return stack (TOS)	1	0000	0000	0000	0110	None	
PUSH	—	Push top of return stack (TOS)	1	0000	0000	0000	0101	None	
RCALL	n	Relative Call	2	1101	1nnn	nnnn	nnnn	None	
RESET		Software device Reset	1	0000	0000	1111	1111	All	
RETFIE	s	Return from interrupt enable	2	0000	0000	0001	000s	GIE/GIEH, PEIE/GIEL	
RETLW	k	Return with literal in WREG	2	0000	1100	kkkk	kkkk	None	
RETURN	s	Return from Subroutine	2	0000	0000	0001	001s	None	
SLEEP	—	Go into Standby mode	1	0000	0000	0000	0011	$\overline{TO}$, $\overline{PD}$	

Note 1: When a PORT register is modified as a function of itself (e.g., `MOVF PORTB, 1, 0`), the value used will be that value present on the pins themselves. For example, if the data latch is '1' for a pin configured as input and is driven low by an external device, the data will be written back with a '0'.

2: If this instruction is executed on the TMR0 register (and where applicable, 'd' = 1), the prescaler will be cleared if assigned.

3: If Program Counter (PC) is modified or a conditional test is true, the instruction requires two cycles. The second cycle is executed as a `NOP`.

4: Some instructions are two-word instructions. The second word of these instructions will be executed as a `NOP` unless the first word of the instruction retrieves the information embedded in these 16 bits. This ensures that all program memory locations have a valid instruction.

TABLE 27-2: PIC18 INSTRUCTION SET (CONTINUED)

Mnemonic, Operands		Description	Cycles	16-Bit Instruction Word				Status Affected	Notes
				MSb			LSb		
LITERAL OPERATIONS									
ADDLW	k	Add literal and WREG	1	0000	1111	kkkk	kkkk	C, DC, Z, OV, N	
ANDLW	k	AND literal with WREG	1	0000	1011	kkkk	kkkk	Z, N	
IORLW	k	Inclusive OR literal with WREG	1	0000	1001	kkkk	kkkk	Z, N	
LFSR	f, k	Move literal (12-bit) 2nd word	2	1110	1110	00ff	kkkk	None	
		to FSR(f) 1st word		1111	0000	kkkk	kkkk		
MOVLB	k	Move literal to BSR<3:0>	1	0000	0001	0000	kkkk	None	
MOVLW	k	Move literal to WREG	1	0000	1110	kkkk	kkkk	None	
MULLW	k	Multiply literal with WREG	1	0000	1101	kkkk	kkkk	None	
RETLW	k	Return with literal in WREG	2	0000	1100	kkkk	kkkk	None	
SUBLW	k	Subtract WREG from literal	1	0000	1000	kkkk	kkkk	C, DC, Z, OV, N	
XORLW	k	Exclusive OR literal with WREG	1	0000	1010	kkkk	kkkk	Z, N	
DATA MEMORY ↔ PROGRAM MEMORY OPERATIONS									
TBLRD*		Table Read	2	0000	0000	0000	1000	None	
TBLRD*+		Table Read with post-increment		0000	0000	0000	1001	None	
TBLRD*-		Table Read with post-decrement		0000	0000	0000	1010	None	
TBLRD+*		Table Read with pre-increment		0000	0000	0000	1011	None	
TBLWT*		Table Write	2	0000	0000	0000	1100	None	
TBLWT*+		Table Write with post-increment		0000	0000	0000	1101	None	
TBLWT*-		Table Write with post-decrement		0000	0000	0000	1110	None	
TBLWT+*		Table Write with pre-increment		0000	0000	0000	1111	None	

Note 1: When a PORT register is modified as a function of itself (e.g., `MOVF PORTB, 1, 0`), the value used will be that value present on the pins themselves. For example, if the data latch is '1' for a pin configured as input and is driven low by an external device, the data will be written back with a '0'.

2: If this instruction is executed on the TMR0 register (and where applicable, 'd' = 1), the prescaler will be cleared if assigned.

3: If Program Counter (PC) is modified or a conditional test is true, the instruction requires two cycles. The second cycle is executed as a `NOP`.

4: Some instructions are two-word instructions. The second word of these instructions will be executed as a `NOP` unless the first word of the instruction retrieves the information embedded in these 16 bits. This ensures that all program memory locations have a valid instruction.

27.1.1 STANDARD INSTRUCTION SET

ADDLW **ADD literal to W**

Syntax:	ADDLW k
Operands:	$0 \leq k \leq 255$
Operation:	(W) + k → W
Status Affected:	N, OV, C, DC, Z
Encoding:	`0000` `1111` `kkkk` `kkkk`
Description:	The contents of W are added to the 8-bit literal 'k' and the result is placed in W.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	Write to W

Example: `ADDLW 15h`

Before Instruction
W = 10h
After Instruction
W = 25h

ADDWF **ADD W to f**

Syntax:	ADDWF f {,d {,a}}
Operands:	$0 \leq f \leq 255$ d ∈ [0,1] a ∈ [0,1]
Operation:	(W) + (f) → dest
Status Affected:	N, OV, C, DC, Z
Encoding:	`0010` `01da` `ffff` `ffff`
Description:	Add W to register 'f'. If 'd' is '0', the result is stored in W. If 'd' is '1', the result is stored back in register 'f'. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank. If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example: `ADDWF REG, 0, 0`

Before Instruction
W = 17h
REG = 0C2h
After Instruction
W = 0D9h
REG = 0C2h

Note: All PIC18 instructions may take an optional label argument preceding the instruction mnemonic for use in symbolic addressing. If a label is used, the instruction format then becomes: {label} instruction argument(s).

ADDWFC — ADD W and CARRY bit to f

Syntax:	ADDWFC f {,d {,a}}
Operands:	$0 \leq f \leq 255$ $d \in [0,1]$ $a \in [0,1]$
Operation:	$(W) + (f) + (C) \rightarrow$ dest
Status Affected:	N,OV, C, DC, Z
Encoding:	`0010` `00da` `ffff` `ffff`
Description:	Add W, the CARRY flag and data memory location 'f'. If 'd' is '0', the result is placed in W. If 'd' is '1', the result is placed in data memory location 'f'. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank. If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example: `ADDWFC REG, 0, 1`

Before Instruction
- CARRY bit = 1
- REG = 02h
- W = 4Dh

After Instruction
- CARRY bit = 0
- REG = 02h
- W = 50h

ANDLW — AND literal with W

Syntax:	ANDLW k
Operands:	$0 \leq k \leq 255$
Operation:	(W) .AND. $k \rightarrow W$
Status Affected:	N, Z
Encoding:	`0000` `1011` `kkkk` `kkkk`
Description:	The contents of W are AND'ed with the 8-bit literal 'k'. The result is placed in W.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	Write to W

Example: `ANDLW 05Fh`

Before Instruction
- W = A3h

After Instruction
- W = 03h

ANDWF — AND W with f

Syntax:	ANDWF f {,d {,a}}
Operands:	$0 \leq f \leq 255$ $d \in [0,1]$ $a \in [0,1]$
Operation:	(W) .AND. (f) $\rightarrow$ dest
Status Affected:	N, Z
Encoding:	`0001` `01da` `ffff` `ffff`
Description:	The contents of W are AND'ed with register 'f'. If 'd' is '0', the result is stored in W. If 'd' is '1', the result is stored back in register 'f'. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank. If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example: `ANDWF REG, 0, 0`

Before Instruction
- W = 17h
- REG = C2h

After Instruction
- W = 02h
- REG = C2h

BC — Branch if Carry

Syntax:	BC n
Operands:	$-128 \leq n \leq 127$
Operation:	if CARRY bit is '1' (PC) + 2 + 2n $\rightarrow$ PC
Status Affected:	None
Encoding:	`1110` `0010` `nnnn` `nnnn`
Description:	If the CARRY bit is '1', then the program will branch. The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC + 2 + 2n. This instruction is then a two-cycle instruction.
Words:	1
Cycles:	1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example: `HERE BC 5`

Before Instruction
- PC = address (`HERE`)

After Instruction
- If CARRY = 1;
- PC = address (`HERE + 12`)
- If CARRY = 0;
- PC = address (`HERE + 2`)

BCF — Bit Clear f

Syntax:	BCF f, b {,a}
Operands:	$0 \le f \le 255$ $0 \le b \le 7$ $a \in [0,1]$
Operation:	$0 \rightarrow$ f<b>
Status Affected:	None
Encoding:	`1001` `bbba` `ffff` `ffff`
Description:	Bit 'b' in register 'f' is cleared. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank. If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write register 'f'

Example: `BCF FLAG_REG, 7, 0`

Before Instruction
FLAG_REG = C7h
After Instruction
FLAG_REG = 47h

BN — Branch if Negative

Syntax:	BN n
Operands:	$-128 \le n \le 127$
Operation:	if NEGATIVE bit is '1' $(PC) + 2 + 2n \rightarrow PC$
Status Affected:	None
Encoding:	`1110` `0110` `nnnn` `nnnn`
Description:	If the NEGATIVE bit is '1', then the program will branch. The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC + 2 + 2n. This instruction is then a two-cycle instruction.
Words:	1
Cycles:	1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example: `HERE BN Jump`

Before Instruction
PC = address (HERE)
After Instruction
If NEGATIVE = 1;
PC = address (Jump)
If NEGATIVE = 0;
PC = address (HERE + 2)

BNC — Branch if Not Carry

Syntax:	BNC n
Operands:	$-128 \leq n \leq 127$
Operation:	if CARRY bit is '0' (PC) + 2 + 2n → PC
Status Affected:	None
Encoding:	`1110` `0011` `nnnn` `nnnn`
Description:	If the CARRY bit is '0', then the program will branch. The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC + 2 + 2n. This instruction is then a two-cycle instruction.
Words:	1
Cycles:	1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example:
```
HERE      BNC   Jump
```

Before Instruction
PC = address (HERE)
After Instruction
If CARRY = 0;
PC = address (Jump)
If CARRY = 1;
PC = address (HERE + 2)

BNN — Branch if Not Negative

Syntax:	BNN n
Operands:	$-128 \leq n \leq 127$
Operation:	if NEGATIVE bit is '0' (PC) + 2 + 2n → PC
Status Affected:	None
Encoding:	`1110` `0111` `nnnn` `nnnn`
Description:	If the NEGATIVE bit is '0', then the program will branch. The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC + 2 + 2n. This instruction is then a two-cycle instruction.
Words:	1
Cycles:	1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example:
```
HERE      BNN   Jump
```

Before Instruction
PC = address (HERE)
After Instruction
If NEGATIVE = 0;
PC = address (Jump)
If NEGATIVE = 1;
PC = address (HERE + 2)

BNOV — Branch if Not Overflow

Syntax: BNOV n

Operands: $-128 \le n \le 127$

Operation: if OVERFLOW bit is '0'
(PC) + 2 + 2n → PC

Status Affected: None

Encoding:

1110	0101	nnnn	nnnn

Description: If the OVERFLOW bit is '0', then the program will branch.
The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC + 2 + 2n. This instruction is then a two-cycle instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example: `HERE BNOV Jump`

Before Instruction
PC = address (`HERE`)
After Instruction
If OVERFLOW = 0;
PC = address (`Jump`)
If OVERFLOW = 1;
PC = address (`HERE + 2`)

BNZ — Branch if Not Zero

Syntax: BNZ n

Operands: $-128 \le n \le 127$

Operation: if ZERO bit is '0'
(PC) + 2 + 2n → PC

Status Affected: None

Encoding:

1110	0001	nnnn	nnnn

Description: If the ZERO bit is '0', then the program will branch.
The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC + 2 + 2n. This instruction is then a two-cycle instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example: `HERE BNZ Jump`

Before Instruction
PC = address (`HERE`)
After Instruction
If ZERO = 0;
PC = address (`Jump`)
If ZERO = 1;
PC = address (`HERE + 2`)

BRA — Unconditional Branch

Syntax:	BRA n
Operands:	-1024 ≤ n ≤ 1023
Operation:	(PC) + 2 + 2n → PC
Status Affected:	None
Encoding:	`1101` `0nnn` `nnnn` `nnnn`
Description:	Add the 2’s complement number ‘2n’ to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC + 2 + 2n. This instruction is a two-cycle instruction.
Words:	1
Cycles:	2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal ‘n’	Process Data	Write to PC
No operation	No operation	No operation	No operation

Example: `HERE BRA Jump`

Before Instruction
PC = address (`HERE`)
After Instruction
PC = address (`Jump`)

BSF — Bit Set f

Syntax:	BSF f, b {,a}
Operands:	0 ≤ f ≤ 255 0 ≤ b ≤ 7 a ∈ [0,1]
Operation:	1 → f<b>
Status Affected:	None
Encoding:	`1000` `bbba` `ffff` `ffff`
Description:	Bit ‘b’ in register ‘f’ is set. If ‘a’ is ‘`0`’, the Access Bank is selected. If ‘a’ is ‘`1`’, the BSR is used to select the GPR bank. If ‘a’ is ‘`0`’ and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever f ≤ 95 (5Fh). See **Section 27.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode”** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register ‘f’	Process Data	Write register ‘f’

Example: `BSF FLAG_REG, 7, 1`

Before Instruction
FLAG_REG = 0Ah
After Instruction
FLAG_REG = 8Ah

BTFSC — Bit Test File, Skip if Clear

Syntax: BTFSC f, b {,a}

Operands: $0 \leq f \leq 255$
$0 \leq b \leq 7$
$a \in [0,1]$

Operation: skip if (f<b>) = `0`

Status Affected: None

Encoding:

`1011`	`bbba`	`ffff`	`ffff`

Description: If bit ‘b’ in register ‘f’ is ‘`0`’, then the next instruction is skipped. If bit ‘b’ is ‘`0`’, then the next instruction fetched during the current instruction execution is discarded and a `NOP` is executed instead, making this a two-cycle instruction.
If ‘a’ is ‘`0`’, the Access Bank is selected. If ‘a’ is ‘`1`’, the BSR is used to select the GPR bank.
If ‘a’ is ‘`0`’ and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode”** for details.

Words: 1

Cycles: 1(2)
Note: Three cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register ‘f’	Process Data	No operation

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

Example:

```
HERE    BTFSC   FLAG, 1, 0
FALSE   :
TRUE    :
```

Before Instruction
PC = address (`HERE`)
After Instruction
If FLAG<1> = 0;
PC = address (`TRUE`)
If FLAG<1> = 1;
PC = address (`FALSE`)

BTFSS — Bit Test File, Skip if Set

Syntax: BTFSS f, b {,a}

Operands: $0 \leq f \leq 255$
$0 \leq b < 7$
$a \in [0,1]$

Operation: skip if (f<b>) = `1`

Status Affected: None

Encoding:

`1010`	`bbba`	`ffff`	`ffff`

Description: If bit ‘b’ in register ‘f’ is ‘`1`’, then the next instruction is skipped. If bit ‘b’ is ‘`1`’, then the next instruction fetched during the current instruction execution is discarded and a `NOP` is executed instead, making this a two-cycle instruction.
If ‘a’ is ‘`0`’, the Access Bank is selected. If ‘a’ is ‘`1`’, the BSR is used to select the GPR bank.
If ‘a’ is ‘`0`’ and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode”** for details.

Words: 1

Cycles: 1(2)
Note: Three cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register ‘f’	Process Data	No operation

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

Example:

```
HERE    BTFSS   FLAG, 1, 0
FALSE   :
TRUE    :
```

Before Instruction
PC = address (`HERE`)
After Instruction
If FLAG<1> = 0;
PC = address (`FALSE`)
If FLAG<1> = 1;
PC = address (`TRUE`)

BTG — Bit Toggle f

Syntax: BTG f, b {,a}

Operands: $0 \le f \le 255$
$0 \le b < 7$
$a \in [0,1]$

Operation: $(\overline{f<b>}) \rightarrow f<b>$

Status Affected: None

Encoding:

`0111`	`bbba`	`ffff`	`ffff`

Description: Bit ‘b’ in data memory location ‘f’ is inverted.
If ‘a’ is ‘`0`’, the Access Bank is selected. If ‘a’ is ‘`1`’, the BSR is used to select the GPR bank.
If ‘a’ is ‘`0`’ and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh). See **Section 27.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode”** for details.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register ‘f’	Process Data	Write register ‘f’

Example: `BTG PORTC, 4, 0`

Before Instruction:
PORTC = `0111 0101` [75h]
After Instruction:
PORTC = `0110 0101` [65h]

BOV — Branch if Overflow

Syntax: BOV n

Operands: $-128 \le n \le 127$

Operation: if OVERFLOW bit is ‘1’
$(PC) + 2 + 2n \rightarrow PC$

Status Affected: None

Encoding:

`1110`	`0100`	`nnnn`	`nnnn`

Description: If the OVERFLOW bit is ‘`1`’, then the program will branch.
The 2’s complement number ‘2n’ is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC + 2 + 2n. This instruction is then a two-cycle instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal ‘n’	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal ‘n’	Process Data	No operation

Example: `HERE BOV Jump`

Before Instruction
PC = address (`HERE`)
After Instruction
If OVERFLOW = 1;
PC = address (`Jump`)
If OVERFLOW = 0;
PC = address (`HERE + 2`)

BZ Branch if Zero

Syntax:	BZ n
Operands:	-128 ≤ n ≤ 127
Operation:	if ZERO bit is '1' (PC) + 2 + 2n → PC
Status Affected:	None
Encoding:	`1110` `0000` `nnnn` `nnnn`
Description:	If the ZERO bit is '1', then the program will branch. The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC + 2 + 2n. This instruction is then a two-cycle instruction.
Words:	1
Cycles:	1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example: `HERE BZ Jump`

Before Instruction
PC = address (HERE)
After Instruction
If ZERO = 1;
PC = address (Jump)
If ZERO = 0;
PC = address (HERE + 2)

CALL Subroutine Call

Syntax:	CALL k {,s}
Operands:	0 ≤ k ≤ 1048575 s ∈ [0,1]
Operation:	(PC) + 4 → TOS, k → PC<20:1>, if s = 1 (W) → WS, (Status) → STATUSS, (BSR) → BSRS
Status Affected:	None

Encoding:

1st word (k<7:0>)	1110	110s	$k_7$kkk	kkkk_0
2nd word(k<19:8>)	1111	k_{19}kkk	kkkk	kkkk_8

Description:	Subroutine call of entire 2-Mbyte memory range. First, return address (PC + 4) is pushed onto the return stack. If 's' = 1, the W, Status and BSR registers are also pushed into their respective shadow registers, WS, STATUSS and BSRS. If 's' = 0, no update occurs. Then, the 20-bit value 'k' is loaded into PC<20:1>. CALL is a two-cycle instruction.
Words:	2
Cycles:	2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'<7:0>,	PUSH PC to stack	Read literal 'k'<19:8>, Write to PC
No operation	No operation	No operation	No operation

Example: `HERE CALL THERE, 1`

Before Instruction
PC = address (HERE)
After Instruction
PC = address (THERE)
TOS = address (HERE + 4)
WS = W
BSRS = BSR
STATUSS = Status

CLRF — Clear f

Syntax:	CLRF f {,a}
Operands:	$0 \le f \le 255$ $a \in [0,1]$
Operation:	$000h \rightarrow f$ $1 \rightarrow Z$
Status Affected:	Z
Encoding:	`0110` `101a` `ffff` `ffff`
Description:	Clears the contents of the specified register. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank. If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write register 'f'

Example: `CLRF FLAG_REG, 1`

Before Instruction
FLAG_REG = 5Ah
After Instruction
FLAG_REG = 00h

CLRWDT — Clear Watchdog Timer

Syntax:	CLRWDT
Operands:	None
Operation:	$000h \rightarrow WDT$, $000h \rightarrow$ WDT postscaler, $1 \rightarrow \overline{TO}$, $1 \rightarrow \overline{PD}$
Status Affected:	$\overline{TO}$, $\overline{PD}$
Encoding:	`0000` `0000` `0000` `0100`
Description:	`CLRWDT` instruction resets the Watchdog Timer. It also resets the postscaler of the WDT. Status bits, $\overline{TO}$ and $\overline{PD}$, are set.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No operation	Process Data	No operation

Example: `CLRWDT`

Before Instruction
WDT Counter = ?
After Instruction
WDT Counter = 00h
WDT Postscaler = 0
$\overline{TO}$ = 1
$\overline{PD}$ = 1

COMF — Complement f

Syntax: COMF f {,d {,a}}

Operands: $0 \le f \le 255$
$d \in [0,1]$
$a \in [0,1]$

Operation: $(\bar{f}) \rightarrow$ dest

Status Affected: N, Z

Encoding:

0001	11da	ffff	ffff

Description: The contents of register ‘f’ are complemented. If ‘d’ is ‘0’, the result is stored in W. If ‘d’ is ‘1’, the result is stored back in register ‘f’.
If ‘a’ is ‘0’, the Access Bank is selected. If ‘a’ is ‘1’, the BSR is used to select the GPR bank.
If ‘a’ is ‘0’ and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh). See **Section 27.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode”** for details.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register ‘f’	Process Data	Write to destination

Example: `COMF REG, 0, 0`

Before Instruction
REG = 13h
After Instruction
REG = 13h
W = ECh

CPFSEQ — Compare f with W, skip if f = W

Syntax: CPFSEQ f {,a}

Operands: $0 \le f \le 255$
$a \in [0,1]$

Operation: (f) – (W),
skip if (f) = (W)
(unsigned comparison)

Status Affected: None

Encoding:

0110	001a	ffff	ffff

Description: Compares the contents of data memory location ‘f’ to the contents of W by performing an unsigned subtraction.
If ‘f’ = W, then the fetched instruction is discarded and a NOP is executed instead, making this a two-cycle instruction.
If ‘a’ is ‘0’, the Access Bank is selected. If ‘a’ is ‘1’, the BSR is used to select the GPR bank.
If ‘a’ is ‘0’ and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh). See **Section 27.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode”** for details.

Words: 1

Cycles: 1(2)
Note: Three cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register ‘f’	Process Data	No operation

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

Example:

```
HERE     CPFSEQ REG, 0
NEQUAL   :
EQUAL    :
```

Before Instruction
PC Address = HERE
W = ?
REG = ?
After Instruction
If REG = W;
PC = Address (EQUAL)
If REG ≠ W;
PC = Address (NEQUAL)

CPFSGT — Compare f with W, skip if f > W

Syntax:	CPFSGT f {,a}
Operands:	$0 \leq f \leq 255$ $a \in [0,1]$
Operation:	(f) – (W), skip if (f) > (W) (unsigned comparison)
Status Affected:	None
Encoding:	`0110` `010a` `ffff` `ffff`
Description:	Compares the contents of data memory location 'f' to the contents of the W by performing an unsigned subtraction. If the contents of 'f' are greater than the contents of WREG, then the fetched instruction is discarded and a `NOP` is executed instead, making this a two-cycle instruction. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank. If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1(2) **Note:** Three cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	No operation

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

<u>Example</u>:

```
HERE      CPFSGT REG, 0
NGREATER  :
GREATER   :
```

Before Instruction
- PC = Address (`HERE`)
- W = ?

After Instruction
- If REG > W;
  - PC = Address (`GREATER`)
- If REG ≤ W;
  - PC = Address (`NGREATER`)

CPFSLT — Compare f with W, skip if f < W

Syntax:	CPFSLT f {,a}
Operands:	$0 \leq f \leq 255$ $a \in [0,1]$
Operation:	(f) – (W), skip if (f) < (W) (unsigned comparison)
Status Affected:	None
Encoding:	`0110` `000a` `ffff` `ffff`
Description:	Compares the contents of data memory location 'f' to the contents of W by performing an unsigned subtraction. If the contents of 'f' are less than the contents of W, then the fetched instruction is discarded and a `NOP` is executed instead, making this a two-cycle instruction. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank.
Words:	1
Cycles:	1(2) **Note:** Three cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	No operation

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

<u>Example</u>:

```
HERE   CPFSLT REG, 1
NLESS  :
LESS   :
```

Before Instruction
- PC = Address (`HERE`)
- W = ?

After Instruction
- If REG < W;
- PC = Address (`LESS`)
- If REG ≥ W;
- PC = Address (`NLESS`)

DAW — Decimal Adjust W Register

Syntax:	DAW
Operands:	None

Operation:	If [W<3:0> > 9] or [DC = 1] then
	(W<3:0>) + 6 → W<3:0>;
	else
	(W<3:0>) → W<3:0>;
	
	If [W<7:4> + DC > 9] or [C = 1] then
	(W<7:4>) + 6 + DC → W<7:4> ;
	C = 1;
	else
	(W<7:4>) + DC → W<7:4>

Status Affected:	C
Encoding:	`0000` `0000` `0000` `0111`
Description:	`DAW` adjusts the eight-bit value in W, resulting from the earlier addition of two variables (each in packed BCD format) and produces a correct packed BCD result.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register W	Process Data	Write W

Example1:

`DAW`

Before Instruction

W = A5h
C = 0
DC = 0

After Instruction

W = 05h
C = 1
DC = 0

Example 2:

Before Instruction

W = CEh
C = 0
DC = 0

After Instruction

W = 34h
C = 1
DC = 0

DECF — Decrement f

Syntax:	DECF f {,d {,a}}
Operands:	$0 \leq f \leq 255$ $d \in [0,1]$ $a \in [0,1]$
Operation:	(f) – 1 → dest
Status Affected:	C, DC, N, OV, Z
Encoding:	`0000` `01da` `ffff` `ffff`
Description:	Decrement register ‘f’. If ‘d’ is ‘`0`’, the result is stored in W. If ‘d’ is ‘`1`’, the result is stored back in register ‘f’. If ‘a’ is ‘`0`’, the Access Bank is selected. If ‘a’ is ‘`1`’, the BSR is used to select the GPR bank. If ‘a’ is ‘`0`’ and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode”** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register ‘f’	Process Data	Write to destination

Example: `DECF CNT, 1, 0`

Before Instruction

CNT = 01h
Z = 0

After Instruction

CNT = 00h
Z = 1

DECFSZ — Decrement f, skip if 0

Syntax:	DECFSZ f {,d {,a}}
Operands:	0 ≤ f ≤ 255 d ∈ [0,1] a ∈ [0,1]
Operation:	(f) – 1 → dest, skip if result = `0`
Status Affected:	None
Encoding:	`0010` `11da` `ffff` `ffff`
Description:	The contents of register ‘f’ are decremented. If ‘d’ is ‘`0`’, the result is placed in W. If ‘d’ is ‘`1`’, the result is placed back in register ‘f’. If the result is ‘`0`’, the next instruction, which is already fetched, is discarded and a `NOP` is executed instead, making it a two-cycle instruction. If ‘a’ is ‘`0`’, the Access Bank is selected. If ‘a’ is ‘`1`’, the BSR is used to select the GPR bank. If ‘a’ is ‘`0`’ and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever f ≤ 95 (5Fh). See **Section 27.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode”** for details.
Words:	1
Cycles:	1(2) **Note:** Three cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register ‘f’	Process Data	Write to destination

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

Example:

```
HERE        DECFSZ   CNT, 1, 1
            GOTO     LOOP
CONTINUE
```

Before Instruction
PC = Address (`HERE`)
After Instruction
CNT = CNT - 1
If CNT = 0;
PC = Address (`CONTINUE`)
If CNT ≠ 0;
PC = Address (`HERE + 2`)

DCFSNZ — Decrement f, skip if not 0

Syntax:	DCFSNZ f {,d {,a}}
Operands:	0 ≤ f ≤ 255 d ∈ [0,1] a ∈ [0,1]
Operation:	(f) – 1 → dest, skip if result ≠ `0`
Status Affected:	None
Encoding:	`0100` `11da` `ffff` `ffff`
Description:	The contents of register ‘f’ are decremented. If ‘d’ is ‘`0`’, the result is placed in W. If ‘d’ is ‘`1`’, the result is placed back in register ‘f’. If the result is not ‘`0`’, the next instruction, which is already fetched, is discarded and a `NOP` is executed instead, making it a two-cycle instruction. If ‘a’ is ‘`0`’, the Access Bank is selected. If ‘a’ is ‘`1`’, the BSR is used to select the GPR bank. If ‘a’ is ‘`0`’ and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever f ≤ 95 (5Fh). See **Section 27.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode”** for details.
Words:	1
Cycles:	1(2) **Note:** Three cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register ‘f’	Process Data	Write to destination

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

Example:

```
HERE    DCFSNZ   TEMP, 1, 0
ZERO    :
NZERO   :
```

Before Instruction
TEMP = ?
After Instruction
TEMP = TEMP – 1,
If TEMP = 0;
PC = Address (`ZERO`)
If TEMP ≠ 0;
PC = Address (`NZERO`)

GOTO — Unconditional Branch

Syntax:	GOTO k
Operands:	$0 \le k \le 1048575$
Operation:	k → PC<20:1>
Status Affected:	None

Encoding:

1st word (k<7:0>)	1110	1111	$k_7$kkk	kkkk$_0$
2nd word(k<19:8>)	1111	k_{19}kkk	kkkk	kkkk$_8$

Description: GOTO allows an unconditional branch anywhere within entire 2-Mbyte memory range. The 20-bit value ‘k’ is loaded into PC<20:1>. GOTO is always a two-cycle instruction.

Words: 2

Cycles: 2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal ‘k’<7:0>,	No operation	Read literal ‘k’<19:8>, Write to PC
No operation	No operation	No operation	No operation

Example: `GOTO THERE`

After Instruction
PC = Address (THERE)

INCF — Increment f

Syntax:	INCF f {,d {,a}}
Operands:	$0 \le f \le 255$ $d \in [0,1]$ $a \in [0,1]$
Operation:	(f) + 1 → dest
Status Affected:	C, DC, N, OV, Z

Encoding:

0010	10da	ffff	ffff

Description: The contents of register ‘f’ are incremented. If ‘d’ is ‘0’, the result is placed in W. If ‘d’ is ‘1’, the result is placed back in register ‘f’.
If ‘a’ is ‘0’, the Access Bank is selected. If ‘a’ is ‘1’, the BSR is used to select the GPR bank.
If ‘a’ is ‘0’ and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh). See **Section 27.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode”** for details.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register ‘f’	Process Data	Write to destination

Example: `INCF CNT, 1, 0`

Before Instruction
CNT = FFh
Z = 0
C = ?
DC = ?

After Instruction
CNT = 00h
Z = 1
C = 1
DC = 1

INCFSZ — Increment f, skip if 0

Syntax:	INCFSZ f {,d {,a}}
Operands:	$0 \le f \le 255$ $d \in [0,1]$ $a \in [0,1]$
Operation:	(f) + 1 → dest, skip if result = `0`
Status Affected:	None
Encoding:	`0011` `11da` `ffff` `ffff`
Description:	The contents of register 'f' are incremented. If 'd' is '`0`', the result is placed in W. If 'd' is '`1`', the result is placed back in register 'f'. If the result is '`0`', the next instruction, which is already fetched, is discarded and a `NOP` is executed instead, making it a two-cycle instruction. If 'a' is '`0`', the Access Bank is selected. If 'a' is '`1`', the BSR is used to select the GPR bank. If 'a' is '`0`' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever f ≤ 95 (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1(2) **Note:** 3 cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

Example:

```
HERE    INCFSZ   CNT, 1, 0
NZERO   :
ZERO    :
```

Before Instruction
PC = Address (HERE)
After Instruction
CNT = CNT + 1
If CNT = 0;
PC = Address (ZERO)
If CNT ≠ 0;
PC = Address (NZERO)

INFSNZ — Increment f, skip if not 0

Syntax:	INFSNZ f {,d {,a}}
Operands:	$0 \le f \le 255$ $d \in [0,1]$ $a \in [0,1]$
Operation:	(f) + 1 → dest, skip if result ≠ `0`
Status Affected:	None
Encoding:	`0100` `10da` `ffff` `ffff`
Description:	The contents of register 'f' are incremented. If 'd' is '`0`', the result is placed in W. If 'd' is '`1`', the result is placed back in register 'f'. If the result is not '`0`', the next instruction, which is already fetched, is discarded and a `NOP` is executed instead, making it a two-cycle instruction. If 'a' is '`0`', the Access Bank is selected. If 'a' is '`1`', the BSR is used to select the GPR bank. If 'a' is '`0`' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever f ≤ 95 (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1(2) **Note:** 3 cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

Example:

```
HERE    INFSNZ   REG, 1, 0
ZERO
NZERO
```

Before Instruction
PC = Address (HERE)
After Instruction
REG = REG + 1
If REG ≠ 0;
PC = Address (NZERO)
If REG = 0;
PC = Address (ZERO)

IORLW — Inclusive OR literal with W

Syntax:	IORLW k
Operands:	$0 \leq k \leq 255$
Operation:	(W) .OR. k $\rightarrow$ W
Status Affected:	N, Z
Encoding:	`0000` `1001` `kkkk` `kkkk`
Description:	The contents of W are ORed with the eight-bit literal 'k'. The result is placed in W.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	Write to W

Example: `IORLW 35h`

Before Instruction
W = 9Ah
After Instruction
W = BFh

IORWF — Inclusive OR W with f

Syntax:	IORWF f {,d {,a}}
Operands:	$0 \leq f \leq 255$ $d \in [0,1]$ $a \in [0,1]$
Operation:	(W) .OR. (f) $\rightarrow$ dest
Status Affected:	N, Z
Encoding:	`0001` `00da` `ffff` `ffff`
Description:	Inclusive OR W with register 'f'. If 'd' is '`0`', the result is placed in W. If 'd' is '`1`', the result is placed back in register 'f'. If 'a' is '`0`', the Access Bank is selected. If 'a' is '`1`', the BSR is used to select the GPR bank. If 'a' is '`0`' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example: `IORWF RESULT, 0, 1`

Before Instruction
RESULT = 13h
W = 91h
After Instruction
RESULT = 13h
W = 93h

LFSR — Load FSR

Syntax: LFSR f, k

Operands: $0 \le f \le 2$
$0 \le k \le 4095$

Operation: $k \rightarrow FSRf$

Status Affected: None

Encoding:

1110	1110	00ff	k_{11}kkk
1111	0000	$k_7$kkk	kkkk

Description: The 12-bit literal 'k' is loaded into the File Select Register pointed to by 'f'.

Words: 2

Cycles: 2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k' MSB	Process Data	Write literal 'k' MSB to FSRfH
Decode	Read literal 'k' LSB	Process Data	Write literal 'k' to FSRfL

Example: `LFSR 2, 3ABh`

After Instruction
FSR2H = 03h
FSR2L = ABh

MOVF — Move f

Syntax: MOVF f {,d {,a}}

Operands: $0 \le f \le 255$
$d \in [0,1]$
$a \in [0,1]$

Operation: $f \rightarrow dest$

Status Affected: N, Z

Encoding:

0101	00da	ffff	ffff

Description: The contents of register 'f' are moved to a destination dependent upon the status of 'd'. If 'd' is '0', the result is placed in W. If 'd' is '1', the result is placed back in register 'f'. Location 'f' can be anywhere in the 256-byte bank.
If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank.
If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write W

Example: `MOVF REG, 0, 0`

Before Instruction
REG = 22h
W = FFh
After Instruction
REG = 22h
W = 22h

MOVFF	**Move f to f**
Syntax:	MOVFF f_s,f_d
Operands:	$0 \le f_s \le 4095$ $0 \le f_d \le 4095$
Operation:	$(f_s) \rightarrow f_d$
Status Affected:	None

Encoding:				
1st word (source)	1100	ffff	ffff	$ffff_s$
2nd word (destin.)	1111	ffff	ffff	$ffff_d$

Description: The contents of source register 'f_s' are moved to destination register 'f_d'. Location of source 'f_s' can be anywhere in the 4096-byte data space (000h to FFFh) and location of destination 'f_d' can also be anywhere from 000h to FFFh.
Either source or destination can be W (a useful special situation).
`MOVFF` is particularly useful for transferring a data memory location to a peripheral register (such as the transmit buffer or an I/O port).
The `MOVFF` instruction cannot use the PCL, TOSU, TOSH or TOSL as the destination register.

Words: 2

Cycles: 2 (3)

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f' (src)	Process Data	No operation
Decode	No operation No dummy read	No operation	Write register 'f' (dest)

Example: `MOVFF REG1, REG2`

Before Instruction
REG1 = 33h
REG2 = 11h
After Instruction
REG1 = 33h
REG2 = 33h

MOVLB	**Move literal to low nibble in BSR**
Syntax:	MOVLB k
Operands:	$0 \le k \le 255$
Operation:	$k \rightarrow$ BSR
Status Affected:	None

Encoding:				
	0000	0001	kkkk	kkkk

Description: The eight-bit literal 'k' is loaded into the Bank Select Register (BSR). The value of BSR<7:4> always remains '0', regardless of the value of k_7:k_4.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	Write literal 'k' to BSR

Example: `MOVLB 5`

Before Instruction
BSR Register = 02h
After Instruction
BSR Register = 05h

MOVLW — Move literal to W

Syntax:	MOVLW k
Operands:	$0 \leq k \leq 255$
Operation:	$k \rightarrow W$
Status Affected:	None
Encoding:	`0000` `1110` `kkkk` `kkkk`
Description:	The eight-bit literal 'k' is loaded into W.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	Write to W

Example: `MOVLW 5Ah`

After Instruction

W = 5Ah

MOVWF — Move W to f

Syntax:	MOVWF f {,a}
Operands:	$0 \leq f \leq 255$ $a \in [0,1]$
Operation:	$(W) \rightarrow f$
Status Affected:	None
Encoding:	`0110` `111a` `ffff` `ffff`
Description:	Move data from W to register 'f'. Location 'f' can be anywhere in the 256-byte bank. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank. If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write register 'f'

Example: `MOVWF REG, 0`

Before Instruction

W = 4Fh
REG = FFh

After Instruction

W = 4Fh
REG = 4Fh

MULLW — Multiply literal with W

Syntax:	MULLW k
Operands:	$0 \leq k \leq 255$
Operation:	(W) x k → PRODH:PRODL
Status Affected:	None
Encoding:	`0000` `1101` `kkkk` `kkkk`
Description:	An unsigned multiplication is carried out between the contents of W and the 8-bit literal 'k'. The 16-bit result is placed in the PRODH:PRODL register pair. PRODH contains the high byte. W is unchanged. None of the Status flags are affected. Note that neither overflow nor carry is possible in this operation. A zero result is possible but not detected.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	Write registers PRODH: PRODL

Example: `MULLW 0C4h`

Before Instruction

W	=	E2h
PRODH	=	?
PRODL	=	?

After Instruction

W	=	E2h
PRODH	=	ADh
PRODL	=	08h

MULWF — Multiply W with f

Syntax:	MULWF f {,a}
Operands:	$0 \leq f \leq 255$ $a \in [0,1]$
Operation:	(W) x (f) → PRODH:PRODL
Status Affected:	None
Encoding:	`0000` `001a` `ffff` `ffff`
Description:	An unsigned multiplication is carried out between the contents of W and the register file location 'f'. The 16-bit result is stored in the PRODH:PRODL register pair. PRODH contains the high byte. Both W and 'f' are unchanged. None of the Status flags are affected. Note that neither overflow nor carry is possible in this operation. A zero result is possible but not detected. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank. If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write registers PRODH: PRODL

Example: `MULWF REG, 1`

Before Instruction

W	=	C4h
REG	=	B5h
PRODH	=	?
PRODL	=	?

After Instruction

W	=	C4h
REG	=	B5h
PRODH	=	8Ah
PRODL	=	94h

NEGF — Negate f

Syntax:	NEGF f {,a}
Operands:	$0 \leq f \leq 255$ $a \in [0,1]$
Operation:	$(\overline{f}) + 1 \rightarrow f$
Status Affected:	N, OV, C, DC, Z
Encoding:	`0110` `110a` `ffff` `ffff`
Description:	Location ‘f’ is negated using two’s complement. The result is placed in the data memory location ‘f’. If ‘a’ is ‘0’, the Access Bank is selected. If ‘a’ is ‘1’, the BSR is used to select the GPR bank. If ‘a’ is ‘0’ and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode”** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register ‘f’	Process Data	Write register ‘f’

Example: `NEGF REG, 1`

Before Instruction
REG = `0011 1010` [3Ah]
After Instruction
REG = `1100 0110` [C6h]

NOP — No Operation

Syntax:	NOP
Operands:	None
Operation:	No operation
Status Affected:	None
Encoding:	`0000` `0000` `0000` `0000` `1111` `xxxx` `xxxx` `xxxx`
Description:	No operation.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No operation	No operation	No operation

Example:

None.

POP — Pop Top of Return Stack

Syntax:	POP
Operands:	None
Operation:	(TOS) → bit bucket
Status Affected:	None
Encoding:	`0000` `0000` `0000` `0110`
Description:	The TOS value is pulled off the return stack and is discarded. The TOS value then becomes the previous value that was pushed onto the return stack. This instruction is provided to enable the user to properly manage the return stack to incorporate a software stack.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No operation	POP TOS value	No operation

Example:

```
POP
GOTO    NEW
```

Before Instruction
- TOS = 0031A2h
- Stack (1 level down) = 014332h

After Instruction
- TOS = 014332h
- PC = NEW

PUSH — Push Top of Return Stack

Syntax:	PUSH
Operands:	None
Operation:	(PC + 2) → TOS
Status Affected:	None
Encoding:	`0000` `0000` `0000` `0101`
Description:	The PC + 2 is pushed onto the top of the return stack. The previous TOS value is pushed down on the stack. This instruction allows implementing a software stack by modifying TOS and then pushing it onto the return stack.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	PUSH PC + 2 onto return stack	No operation	No operation

Example:

```
PUSH
```

Before Instruction
- TOS = 345Ah
- PC = 0124h

After Instruction
- PC = 0126h
- TOS = 0126h
- Stack (1 level down) = 345Ah

RCALL Relative Call

Syntax:	RCALL n
Operands:	$-1024 \le n \le 1023$
Operation:	(PC) + 2 → TOS, (PC) + 2 + 2n → PC
Status Affected:	None
Encoding:	`1101` `1nnn` `nnnn` `nnnn`
Description:	Subroutine call with a jump up to 1K from the current location. First, return address (PC + 2) is pushed onto the stack. Then, add the 2's complement number '2n' to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC + 2 + 2n. This instruction is a two-cycle instruction.
Words:	1
Cycles:	2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'n' PUSH PC to stack	Process Data	Write to PC
No operation	No operation	No operation	No operation

Example: `HERE RCALL Jump`

Before Instruction
PC = Address (`HERE`)
After Instruction
PC = Address (`Jump`)
TOS = Address (`HERE + 2`)

RESET Reset

Syntax:	RESET
Operands:	None
Operation:	Reset all registers and flags that are affected by a $\overline{\text{MCLR}}$ Reset.
Status Affected:	All
Encoding:	`0000` `0000` `1111` `1111`
Description:	This instruction provides a way to execute a $\overline{\text{MCLR}}$ Reset by software.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Start Reset	No operation	No operation

Example: `RESET`

After Instruction
Registers = Reset Value
Flags* = Reset Value

RETFIE — Return from Interrupt

Syntax:	RETFIE {s}
Operands:	s ∈ [0,1]
Operation:	(TOS) → PC, 1 → GIE/GIEH or PEIE/GIEL, if s = 1 (WS) → W, (STATUSS) → Status, (BSRS) → BSR, PCLATU, PCLATH are unchanged.
Status Affected:	GIE/GIEH, PEIE/GIEL.
Encoding:	`0000` `0000` `0001` `000s`
Description:	Return from interrupt. Stack is popped and Top-of-Stack (TOS) is loaded into the PC. Interrupts are enabled by setting either the high or low priority global interrupt enable bit. If 's' = 1, the contents of the shadow registers, WS, STATUSS and BSRS, are loaded into their corresponding registers, W, STATUS and BSR. If 's' = 0, no update of these registers occurs.
Words:	1
Cycles:	2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No operation	No operation	POP PC from stack Set GIEH or GIEL
No operation	No operation	No operation	No operation

Example: `RETFIE 1`

After Interrupt

PC	=	TOS
W	=	WS
BSR	=	BSRS
Status	=	STATUSS
GIE/GIEH, PEIE/GIEL	=	1

RETLW — Return literal to W

Syntax:	RETLW k
Operands:	0 ≤ k ≤ 255
Operation:	k → W, (TOS) → PC, PCLATU, PCLATH are unchanged
Status Affected:	None
Encoding:	`0000` `1100` `kkkk` `kkkk`
Description:	W is loaded with the eight-bit literal 'k'. The program counter is loaded from the top of the stack (the return address). The high address latch (PCLATH) remains unchanged.
Words:	1
Cycles:	2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	POP PC from stack, Write to W
No operation	No operation	No operation	No operation

Example:

```
    CALL TABLE ; W contains table
               ; offset value
               ; W now has
               ; table value
   :
TABLE
    ADDWF PCL  ; W = offset
    RETLW k0   ; Begin table
    RETLW k1   ;
   :
   :
    RETLW kn   ; End of table
```

Before Instruction

W = 07h

After Instruction

W = value of kn

RETURN — Return from Subroutine

Syntax:	RETURN {s}
Operands:	s ∈ [0,1]
Operation:	(TOS) → PC, if s = 1 (WS) → W, (STATUSS) → Status, (BSRS) → BSR, PCLATU, PCLATH are unchanged
Status Affected:	None
Encoding:	`0000 0000 0001 001s`
Description:	Return from subroutine. The stack is popped and the top of the stack (TOS) is loaded into the program counter. If 's'= 1, the contents of the shadow registers, WS, STATUSS and BSRS, are loaded into their corresponding registers, W, STATUS and BSR. If 's' = 0, no update of these registers occurs.
Words:	1
Cycles:	2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No operation	Process Data	POP PC from stack
No operation	No operation	No operation	No operation

Example: `RETURN`

After Instruction:
PC = TOS

RLCF — Rotate Left f through Carry

Syntax:	RLCF f {,d {,a}}
Operands:	$0 \le f \le 255$ d ∈ [0,1] a ∈ [0,1]
Operation:	(f<n>) → dest<n + 1>, (f<7>) → C, (C) → dest<0>
Status Affected:	C, N, Z
Encoding:	`0011 01da ffff ffff`
Description:	The contents of register 'f' are rotated one bit to the left through the CARRY flag. If 'd' is '0', the result is placed in W. If 'd' is '1', the result is stored back in register 'f'. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank. If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.

Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example: `RLCF REG, 0, 0`

Before Instruction
REG = `1110 0110`
C = `0`
After Instruction
REG = `1110 0110`
W = `1100 1100`
C = `1`

RLNCF — Rotate Left f (No Carry)

Syntax: RLNCF f {,d {,a}}

Operands: $0 \leq f \leq 255$
$d \in [0,1]$
$a \in [0,1]$

Operation: (f<n>) → dest<n + 1>,
(f<7>) → dest<0>

Status Affected: N, Z

Encoding:

`0100`	`01da`	`ffff`	`ffff`

Description: The contents of register 'f' are rotated one bit to the left. If 'd' is '0', the result is placed in W. If 'd' is '1', the result is stored back in register 'f'.
If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank.
If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example: `RLNCF REG, 1, 0`

Before Instruction
REG = `1010 1011`
After Instruction
REG = `0101 0111`

RRCF — Rotate Right f through Carry

Syntax: RRCF f {,d {,a}}

Operands: $0 \leq f \leq 255$
$d \in [0,1]$
$a \in [0,1]$

Operation: (f<n>) → dest<n – 1>,
(f<0>) → C,
(C) → dest<7>

Status Affected: C, N, Z

Encoding:

`0011`	`00da`	`ffff`	`ffff`

Description: The contents of register 'f' are rotated one bit to the right through the CARRY flag. If 'd' is '0', the result is placed in W. If 'd' is '1', the result is placed back in register 'f'.
If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank.
If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example: `RRCF REG, 0, 0`

Before Instruction
REG = `1110 0110`
C = `0`
After Instruction
REG = `1110 0110`
W = `0111 0011`
C = `0`

RRNCF — Rotate Right f (No Carry)

Syntax: RRNCF f {,d {,a}}

Operands: $0 \leq f \leq 255$
$d \in [0,1]$
$a \in [0,1]$

Operation: (f<n>) → dest<n – 1>,
(f<0>) → dest<7>

Status Affected: N, Z

Encoding:

`0100`	`00da`	`ffff`	`ffff`

Description: The contents of register 'f' are rotated one bit to the right. If 'd' is '`0`', the result is placed in W. If 'd' is '`1`', the result is placed back in register 'f'.
If 'a' is '`0`', the Access Bank will be selected, overriding the BSR value. If 'a' is '`1`', then the bank will be selected as per the BSR value.
If 'a' is '`0`' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example 1: `RRNCF REG, 1, 0`

Before Instruction
REG = `1101 0111`
After Instruction
REG = `1110 1011`

Example 2: `RRNCF REG, 0, 0`

Before Instruction
W = ?
REG = `1101 0111`
After Instruction
W = `1110 1011`
REG = `1101 0111`

SETF — Set f

Syntax: SETF f {,a}

Operands: $0 \leq f \leq 255$
$a \in [0,1]$

Operation: FFh → f

Status Affected: None

Encoding:

`0110`	`100a`	`ffff`	`ffff`

Description: The contents of the specified register are set to FFh.
If 'a' is '`0`', the Access Bank is selected. If 'a' is '`1`', the BSR is used to select the GPR bank.
If 'a' is '`0`' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write register 'f'

Example: `SETF REG, 1`

Before Instruction
REG = 5Ah
After Instruction
REG = FFh

SLEEP — Enter Sleep mode

Syntax:	SLEEP
Operands:	None
Operation:	00h → WDT, 0 → WDT postscaler, 1 → $\overline{TO}$, 0 → $\overline{PD}$
Status Affected:	$\overline{TO}$, $\overline{PD}$
Encoding:	0000 0000 0000 0011
Description:	The Power-down Status bit ($\overline{PD}$) is cleared. The Time-out Status bit ($\overline{TO}$) is set. The Watchdog Timer and its postscaler are cleared. The processor is put into Sleep mode with the oscillator stopped.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No operation	Process Data	Go to Sleep

Example: SLEEP

Before Instruction
- $\overline{TO}$ = ?
- $\overline{PD}$ = ?

After Instruction
- $\overline{TO}$ = 1 †
- $\overline{PD}$ = 0

† If WDT causes wake-up, this bit is cleared.

SUBFWB — Subtract f from W with borrow

Syntax:	SUBFWB f {,d {,a}}
Operands:	$0 \le f \le 255$ $d \in [0,1]$ $a \in [0,1]$
Operation:	$(W) - (f) - (\overline{C}) \rightarrow$ dest
Status Affected:	N, OV, C, DC, Z
Encoding:	0101 01da ffff ffff
Description:	Subtract register 'f' and CARRY flag (borrow) from W (2's complement method). If 'd' is '0', the result is stored in W. If 'd' is '1', the result is stored in register 'f'. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank. If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example 1: SUBFWB REG, 1, 0

Before Instruction
- REG = 3
- W = 2
- C = 1

After Instruction
- REG = FF
- W = 2
- C = 0
- Z = 0
- N = 1 ; result is negative

Example 2: SUBFWB REG, 0, 0

Before Instruction
- REG = 2
- W = 5
- C = 1

After Instruction
- REG = 2
- W = 3
- C = 1
- Z = 0
- N = 0 ; result is positive

Example 3: SUBFWB REG, 1, 0

Before Instruction
- REG = 1
- W = 2
- C = 0

After Instruction
- REG = 0
- W = 2
- C = 1
- Z = 1 ; result is zero
- N = 0

SUBLW — Subtract W from literal

Syntax:	SUBLW k
Operands:	$0 \leq k \leq 255$
Operation:	$k - (W) \rightarrow W$
Status Affected:	N, OV, C, DC, Z
Encoding:	`0000` `1000` `kkkk` `kkkk`
Description	W is subtracted from the eight-bit literal 'k'. The result is placed in W.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	Write to W

Example 1: `SUBLW 02h`

```
Before Instruction
    W    =  01h
    C    =  ?
After Instruction
    W    =  01h
    C    =  1     ; result is positive
    Z    =  0
    N    =  0
```

Example 2: `SUBLW 02h`

```
Before Instruction
    W    =  02h
    C    =  ?
After Instruction
    W    =  00h
    C    =  1     ; result is zero
    Z    =  1
    N    =  0
```

Example 3: `SUBLW 02h`

```
Before Instruction
    W    =  03h
    C    =  ?
After Instruction
    W    =  FFh   ; (2's complement)
    C    =  0     ; result is negative
    Z    =  0
    N    =  1
```

SUBWF — Subtract W from f

Syntax:	SUBWF f {,d {,a}}
Operands:	$0 \leq f \leq 255$ $d \in [0,1]$ $a \in [0,1]$
Operation:	$(f) - (W) \rightarrow dest$
Status Affected:	N, OV, C, DC, Z
Encoding:	`0101` `11da` `ffff` `ffff`
Description:	Subtract W from register 'f' (2's complement method). If 'd' is '0', the result is stored in W. If 'd' is '1', the result is stored back in register 'f'. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank. If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example 1: `SUBWF REG, 1, 0`

```
Before Instruction
    REG  =  3
    W    =  2
    C    =  ?
After Instruction
    REG  =  1
    W    =  2
    C    =  1     ; result is positive
    Z    =  0
    N    =  0
```

Example 2: `SUBWF REG, 0, 0`

```
Before Instruction
    REG  =  2
    W    =  2
    C    =  ?
After Instruction
    REG  =  2
    W    =  0
    C    =  1     ; result is zero
    Z    =  1
    N    =  0
```

Example 3: `SUBWF REG, 1, 0`

```
Before Instruction
    REG  =  1
    W    =  2
    C    =  ?
After Instruction
    REG  =  FFh   ;(2's complement)
    W    =  2
    C    =  0     ; result is negative
    Z    =  0
    N    =  1
```


SUBWFB — Subtract W from f with Borrow

Syntax:	SUBWFB f {,d {,a}}
Operands:	$0 \le f \le 255$ $d \in [0,1]$ $a \in [0,1]$
Operation:	(f) – (W) – ($\overline{C}$) → dest
Status Affected:	N, OV, C, DC, Z
Encoding:	`0101` `10da` `ffff` `ffff`
Description:	Subtract W and the CARRY flag (borrow) from register 'f' (2's complement method). If 'd' is '0', the result is stored in W. If 'd' is '1', the result is stored back in register 'f'. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank. If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example 1: `SUBWFB REG, 1, 0`

```
Before Instruction
    REG   =   19h    (0001 1001)
    W     =   0Dh    (0000 1101)
    C     =   1
After Instruction
    REG   =   0Ch    (0000 1100)
    W     =   0Dh    (0000 1101)
    C     =   1
    Z     =   0
    N     =   0      ; result is positive
```

Example 2: `SUBWFB REG, 0, 0`

```
Before Instruction
    REG   =   1Bh    (0001 1011)
    W     =   1Ah    (0001 1010)
    C     =   0
After Instruction
    REG   =   1Bh    (0001 1011)
    W     =   00h
    C     =   1
    Z     =   1      ; result is zero
    N     =   0
```

Example 3: `SUBWFB REG, 1, 0`

```
Before Instruction
    REG   =   03h    (0000 0011)
    W     =   0Eh    (0000 1110)
    C     =   1
After Instruction
    REG   =   F5h    (1111 0101)
                     ; [2's comp]
    W     =   0Eh    (0000 1110)
    C     =   0
    Z     =   0
    N     =   1      ; result is negative
```

SWAPF — Swap f

Syntax:	SWAPF f {,d {,a}}
Operands:	$0 \le f \le 255$ $d \in [0,1]$ $a \in [0,1]$
Operation:	(f<3:0>) → dest<7:4>, (f<7:4>) → dest<3:0>
Status Affected:	None
Encoding:	`0011` `10da` `ffff` `ffff`
Description:	The upper and lower nibbles of register 'f' are exchanged. If 'd' is '0', the result is placed in W. If 'd' is '1', the result is placed in register 'f'. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank. If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example: `SWAPF REG, 1, 0`

```
Before Instruction
    REG   =   53h
After Instruction
    REG   =   35h
```

TBLRD	Table Read
Syntax:	TBLRD (*; *+; *-; +*)
Operands:	None
Operation:	if TBLRD *, (Prog Mem (TBLPTR)) → TABLAT; TBLPTR – No Change; if TBLRD *+, (Prog Mem (TBLPTR)) → TABLAT; (TBLPTR) + 1 → TBLPTR; if TBLRD *-, (Prog Mem (TBLPTR)) → TABLAT; (TBLPTR) – 1 → TBLPTR; if TBLRD +*, (TBLPTR) + 1 → TBLPTR; (Prog Mem (TBLPTR)) → TABLAT;
Status Affected:	None

Encoding:

`0000`	`0000`	`0000`	`10nn`
			`nn=0 *`
			`=1 *+`
			`=2 *-`
			`=3 +*`

Description: This instruction is used to read the contents of Program Memory (P.M.). To address the program memory, a pointer called Table Pointer (TBLPTR) is used.

The TBLPTR (a 21-bit pointer) points to each byte in the program memory. TBLPTR has a 2-Mbyte address range.

TBLPTR[0] = `0`:	Least Significant Byte of Program Memory Word
TBLPTR[0] = `1`:	Most Significant Byte of Program Memory Word

The `TBLRD` instruction can modify the value of TBLPTR as follows:

- no change
- post-increment
- post-decrement
- pre-increment

Words: 1

Cycles: 2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No operation	No operation	No operation
No operation	No operation (Read Program Memory)	No operation	No operation (Write TABLAT)

TBLRD Table Read (Continued)

Example1: `TBLRD *+ ;`

Before Instruction

TABLAT	=	55h
TBLPTR	=	00A356h
MEMORY (00A356h)	=	34h

After Instruction

TABLAT	=	34h
TBLPTR	=	00A357h

Example2: `TBLRD +* ;`

Before Instruction

TABLAT	=	AAh
TBLPTR	=	01A357h
MEMORY (01A357h)	=	12h
MEMORY (01A358h)	=	34h

After Instruction

TABLAT	=	34h
TBLPTR	=	01A358h

TBLWT — Table Write

Syntax:	TBLWT (*; *+; *-; +*)
Operands:	None
Operation:	if TBLWT*, (TABLAT) → Holding Register; TBLPTR – No Change; if TBLWT*+, (TABLAT) → Holding Register; (TBLPTR) + 1 → TBLPTR; if TBLWT*-, (TABLAT) → Holding Register; (TBLPTR) – 1 → TBLPTR; if TBLWT+*, (TBLPTR) + 1 → TBLPTR; (TABLAT) → Holding Register;
Status Affected:	None

Encoding:

0000	0000	0000	11nn nn=0 * =1 *+ =2 *- =3 +*

Description: This instruction uses the three LSBs of TBLPTR to determine which of the eight holding registers the TABLAT is written to. The holding registers are used to program the contents of Program Memory (P.M.). (Refer to **Section 7.0 “Flash Program Memory”** for additional details on programming Flash memory.)

The TBLPTR (a 21-bit pointer) points to each byte in the program memory. TBLPTR has a 2-MByte address range. The LSb of the TBLPTR selects which byte of the program memory location to access.

TBLPTR[0] = 0: Least Significant Byte of Program Memory Word
TBLPTR[0] = 1: Most Significant Byte of Program Memory Word

The `TBLWT` instruction can modify the value of TBLPTR as follows:

- no change
- post-increment
- post-decrement
- pre-increment

Words: 1

Cycles: 2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No operation	No operation	No operation
No operation	No operation (Read TABLAT)	No operation	No operation (Write to Holding Register)

TBLWT — Table Write (Continued)

Example1: `TBLWT *+;`

Before Instruction

TABLAT	=	55h
TBLPTR	=	00A356h
HOLDING REGISTER (00A356h)	=	FFh

After Instructions (table write completion)

TABLAT	=	55h
TBLPTR	=	00A357h
HOLDING REGISTER (00A356h)	=	55h

Example 2: `TBLWT +*;`

Before Instruction

TABLAT	=	34h
TBLPTR	=	01389Ah
HOLDING REGISTER (01389Ah)	=	FFh
HOLDING REGISTER (01389Bh)	=	FFh

After Instruction (table write completion)

TABLAT	=	34h
TBLPTR	=	01389Bh
HOLDING REGISTER (01389Ah)	=	FFh
HOLDING REGISTER (01389Bh)	=	34h

TSTFSZ — Test f, skip if 0

Syntax:	TSTFSZ f {,a}
Operands:	0 ≤ f ≤ 255 a ∈ [0,1]
Operation:	skip if f = `0`
Status Affected:	None
Encoding:	`0110` `011a` `ffff` `ffff`
Description:	If 'f' = `0`, the next instruction fetched during the current instruction execution is discarded and a `NOP` is executed, making this a two-cycle instruction. If 'a' is '`0`', the Access Bank is selected. If 'a' is '`1`', the BSR is used to select the GPR bank. If 'a' is '`0`' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever f ≤ 95 (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1(2) **Note:** Three cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	No operation

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

Example:

```
HERE    TSTFSZ  CNT, 1
NZERO   :
ZERO    :
```

Before Instruction
- PC = Address (`HERE`)

After Instruction
- If CNT = 00h,
- PC = Address (`ZERO`)
- If CNT ≠ 00h,
- PC = Address (`NZERO`)

XORLW — Exclusive OR literal with W

Syntax:	XORLW k
Operands:	0 ≤ k ≤ 255
Operation:	(W) .XOR. k → W
Status Affected:	N, Z
Encoding:	`0000` `1010` `kkkk` `kkkk`
Description:	The contents of W are XORed with the 8-bit literal 'k'. The result is placed in W.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	Write to W

Example:

```
XORLW   0AFh
```

Before Instruction
- W = B5h

After Instruction
- W = 1Ah

XORWF — Exclusive OR W with f

Syntax:	XORWF f {,d {,a}}
Operands:	$0 \leq f \leq 255$ $d \in [0,1]$ $a \in [0,1]$
Operation:	(W) .XOR. (f) $\rightarrow$ dest
Status Affected:	N, Z
Encoding:	`0001` `10da` `ffff` `ffff`
Description:	Exclusive OR the contents of W with register 'f'. If 'd' is '0', the result is stored in W. If 'd' is '1', the result is stored back in the register 'f'. If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank. If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 27.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example: `XORWF REG, 1, 0`

Before Instruction
REG = AFh
W = B5h

After Instruction
REG = 1Ah
W = B5h