

Gramatica de precedenta simpla

- mai exista (si nu vom face): gramatici de precedenta slaba, gramatici de precedent a operatorilor
- se va face doar la seminar

- Analiza ascendenta
- Depisteaza limita dreapta si a celei stanga pentru a face o reducere
Se folosesc relatiile $<\bullet$, $\bullet>$, $=\bullet$ (relatii de precedent)

Relatii de precedenta Wirth-Weber

$$R_{<\bullet} \subset (N \cup \Sigma \cup \{\$\}) \times (N \cup \Sigma \cup \{\$\})$$

$$R_{=\bullet} \subset (N \cup \Sigma) \times (N \cup \Sigma)$$

$$R_{\bullet>} \subset (N \cup \Sigma \cup \{\$\}) \times (\Sigma \cup \{\$\})$$

$$X = \bullet Y : A \rightarrow \alpha XY \gamma \in P$$

$$X < \bullet Y : A \rightarrow \alpha XB \gamma \in P, B \Rightarrow^+ Y \gamma$$

$$X \bullet > a : A \rightarrow \alpha BY \gamma \in P, B \Rightarrow^+ \gamma X, Y \Rightarrow^* a \delta$$

$$\$ < \bullet X : S \Rightarrow^+ X \alpha$$

$$X \bullet > \$: S \Rightarrow^+ \alpha X$$

Definitie:

gramatica de precedenta simpla

este o gramatica indep. de context proprie

(inclusiv ϵ -independenta)

- unic invertibila:
nu exista 2 reguli de productie cu acelasi membru drept
- intre oricare 2 simboluri exista cel mult o relatie de precedenta

Analizorul de precedenta simpla

- construiesc tabelul de precedenta a operatorilor
- analizeaza o secventa de terminale

modelul stivei $\sim LR$

$<\bullet$ si $=\bullet$ - deplasare

$\bullet>$ - reducere $Y <\bullet X_1 =\bullet \dots =\bullet X_i \bullet> Z$
 $A \rightarrow X_1 \dots X_i$

Gramatici de precedenta simpla

$$X = Y \quad \exists A \rightarrow \alpha XY \beta$$

$$X < Y \quad A \rightarrow \alpha XB \beta, B \Rightarrow^+ Y \gamma$$

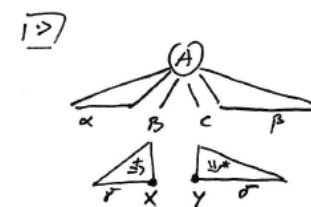
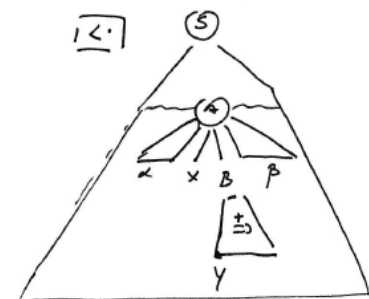
$$X \bullet Y \quad A \rightarrow \alpha BC \beta$$

$$B \Rightarrow^+ \gamma X$$

$$C \Rightarrow^+ Y \delta, Y \in \Sigma$$

$$\$ < X \quad S \Rightarrow^+ X \beta$$

$$Y \bullet \$ \quad S \Rightarrow^+ \alpha Y$$



Exemplu:

$$S \rightarrow aSSb$$

$$S \rightarrow c$$

	S	a	b	c	\$
S	= \bullet	< \bullet	= \bullet	< \bullet	
a	= \bullet	< \bullet		< \bullet	
b		> \bullet	> \bullet	> \bullet	> \bullet
c		> \bullet	> \bullet	> \bullet	> \bullet
\$		< \bullet		< \bullet	

Cuvantul: accb $? \in L(G)$

curăntul accl

\$ a c c l \$
 $\underbrace{c \ c}_{S}$

\$ a s c l \$
 $\underbrace{c \ = \ c}_{S}$

\$ a s s l \$
 $\underbrace{c \ \cdot \cdot \cdot \cdot}_{S}$
\$ S \$

Modelul otivei
 (metode deplasare-reducere)

\$ a c c c l \$	
\$ a c c c l \$	
\$ a c c l \$	<u>rp. 1</u>
\$ a s c l \$	
\$ a s c l \$	<u>rp. 2</u>
\$ a s s l \$	
\$ a s s l \$	
\$ a s s l \$	<u>rp. 1</u>
\$ S \$	

Gramatica de precedenta a operatorilot

Notatie: $\mu \in \mathbf{N} \cup \{\varepsilon\}$

exista o reg.productie de forma:

$$a = \bullet b : \quad A \rightarrow \alpha a \mu b \beta \in P$$

$$a < \bullet b : \quad A \rightarrow \alpha a B \beta \in P, B \Rightarrow^+ \mu b \delta$$

$$a \bullet > b : \quad A \rightarrow \alpha B b \beta \in P, B \Rightarrow^+ \gamma a \mu$$

$$\$ < \bullet b : \quad S \Rightarrow^+ \mu b \alpha$$

$$X \bullet > \$: \quad S \Rightarrow^+ \alpha a \mu$$

$$X \bullet > a : \quad A \rightarrow \alpha B Y \gamma \in P, B \Rightarrow^+ \gamma X, Y \Rightarrow^* a \delta$$

$$\$ < \bullet X : \quad S \Rightarrow^+ X \alpha$$

$$X \bullet > \$: \quad S \Rightarrow^+ \alpha X$$

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