

JUST ANOTHER IMPLEMENTATION OF THE FAST MULTIPOLE METHOD FOR THE ELECTROSTATIC POTENTIAL IN 2D

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Abstract.

1. Introduction.

2. The multipole expansion for well separated clusters of points. Consider the electrostatic potential $\{u_i\}_{i=1}^M \subset \mathbb{R}$ at some target points $\{\mathbf{x}\}_{i=1}^M \subset \mathbb{R}^2$ generated by a set of point charges $\{q_j\}_{j=1}^N \subset \mathbb{R}$ located at source points $\{\mathbf{y}_j\}_{j=1}^N$. Mathematically, this is expressed as

$$(2.1) \quad u(\mathbf{x}_i) := u_i = \sum_{j=1}^N G(\mathbf{x}_i, \mathbf{y}_j) q_j, \quad i = 1, 2, \dots, M,$$

where the kernel G is taken as a scaled version of the fundamental solution of the Laplace equation in two dimensions,

$$(2.2) \quad G(\mathbf{x}, \mathbf{y}) = \begin{cases} \log |\mathbf{x} - \mathbf{y}|, & \text{if } \mathbf{x} \neq \mathbf{y}, \\ 0, & \text{if } \mathbf{x} = \mathbf{y}. \end{cases}$$

Clearly, a direct evaluation of (2.1) would take $\mathcal{O}(NM)$ operations, which is prohibitively expensive when we have a large number of source and target points $N, M \gg 1$. In what follows, we will describe how to compute an arbitrarily close approximation to this sum in a much faster way, subject to the constraint that source points are far away from target points.

Suppose that the source points are clustered in a box Ω_σ , with center \mathbf{c}_σ , whereas the target points are clustered in a disjoint box Ω_τ of the same size, with center \mathbf{c}_τ , as shown in Fig. 1. The separation between boxes is $d := |\mathbf{c}_\tau - \mathbf{c}_\sigma|$.

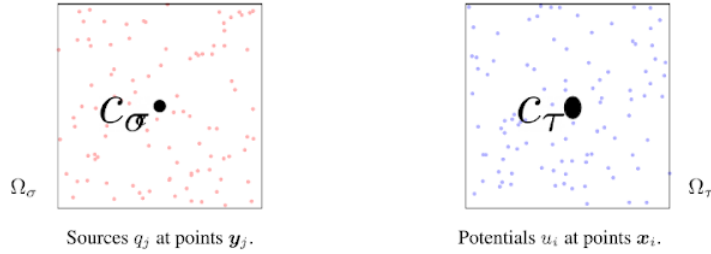


FIGURE 1. The problem of computing the electrostatic potential between well-separated source and target boxes.

For simplicity, let us switch to complex notation, where each point $\mathbf{x}, \mathbf{y} \in \mathbb{R}^2$ is reinterpreted as a point in the complex plane \mathbb{C} , and with abuse of notation we will write the complex kernel $G(\mathbf{x}, \mathbf{y}) = \log(\mathbf{x} - \mathbf{y})$, where \log is the principal branch of the complex logarithm. The real kernel in (2.2) is retrieved by taking $\text{Re}\{G(\mathbf{x}, \mathbf{y})\}$.

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We can now formally expand the potential u_i in a Taylor expansion,

$$\begin{aligned}
(2.3) \quad u_i &= \sum_{j=1}^N \log(\mathbf{x}_i - \mathbf{y}_j) q_j \\
&= \sum_{j=1}^N \log((\mathbf{x}_i - \mathbf{c}_\sigma) - (\mathbf{y}_j - \mathbf{c}_\sigma)) q_j \\
&= \sum_{j=1}^N \left[\log(\mathbf{x}_i - \mathbf{c}_\sigma) + \log \left(1 - \frac{\mathbf{y}_j - \mathbf{c}_\sigma}{\mathbf{x}_i - \mathbf{c}_\sigma} \right) \right] q_j \\
&= \sum_{j=1}^N \left[\log(\mathbf{x}_i - \mathbf{c}_\sigma) - \sum_{p=1}^{\infty} \frac{1}{p} \frac{(\mathbf{y}_j - \mathbf{c}_\sigma)^p}{(\mathbf{x}_i - \mathbf{c}_\sigma)^p} \right] q_j \\
&\approx \log(\mathbf{x}_i - \mathbf{c}_\sigma) \hat{q}_0^\sigma + \sum_{p=1}^P \frac{1}{p} \frac{1}{(\mathbf{x}_i - \mathbf{c}_\sigma)^p} \hat{q}_p^\sigma,
\end{aligned}$$

where we used the fact that $\log(1 - z) = -\sum_p z^p/p$ for $|z| < 1$, the distance d is chosen such that $\left| \frac{\mathbf{y}_j - \mathbf{c}_\sigma}{\mathbf{x}_i - \mathbf{c}_\sigma} \right| < 1$ for every pair of target/source points, and we have approximated the infinite sum by P terms, where P is known as the *interaction rank*. The quantities $\hat{\mathbf{q}}^\sigma := \{\hat{q}_p^\sigma\}_{p=0}^{P-1}$ given by

$$\begin{aligned}
(2.4) \quad \hat{q}_0^\sigma &:= \sum_{j=1}^N q_j, \\
\hat{q}_p^\sigma &:= \sum_{j=1}^N \frac{-1}{p} (\mathbf{y}_j - \mathbf{c}_\sigma)^p q_j, \quad p = 1, 2, \dots, P-1,
\end{aligned}$$

are the *outgoing expansions* of the source box σ . On the other hand, the potential $u(x)$ is analytic, thus it can be expanded as a (truncated) Taylor series around the target box center \mathbf{c}_τ ,

$$(2.5) \quad u(\mathbf{x}_i) = u_i \approx \sum_{p=0}^{P-1} (\mathbf{x}_i - \mathbf{c}_\tau)^p \hat{\mathbf{v}}_p^\tau,$$

where $\hat{\mathbf{v}}^\tau := \{\hat{v}_p^\tau\}_{p=0}^{P-1}$ are the *incoming expansions* of the target box τ . By expanding (2.3) in Taylor, and again choosing the distance d large enough, it can be deduced that

$$\begin{aligned}
(2.6) \quad \hat{v}_0^\tau &= \log(\mathbf{c}_\tau - \mathbf{c}_\sigma) \hat{q}_0^\sigma + \sum_{p=1}^{P-1} (-1)^p \frac{1}{(\mathbf{c}_\sigma - \mathbf{c}_\tau)^p} \hat{q}_p^\sigma, \\
\hat{v}_r^\tau &= -\frac{1}{r(\mathbf{c}_\sigma - \mathbf{c}_\tau)^r} \hat{q}_0^\sigma + \sum_{p=1}^{P-1} (-1)^p \binom{r+p-1}{p-1} \frac{1}{(\mathbf{c}_\sigma - \mathbf{c}_\tau)^{r+p}} \hat{q}_p^\sigma, \quad r = 1, 2, \dots, P-1.
\end{aligned}$$

Denote N_τ and N_σ the numbers of target and source points, respectively, $\mathbf{q}^\sigma :=$

$\{q_j\}_{j=1}^{N_\sigma}$ and $\mathbf{u}^\tau := \{u_i\}_{i=1}^{N_\tau}$. Note that we have implicitly defined the linear maps

$$(2.7) \quad \hat{\mathbf{q}}^\sigma = \mathbf{T}_\sigma^{\text{ofs}} \mathbf{q}^\sigma,$$

$$(2.8) \quad \hat{\mathbf{v}}^\tau = \mathbf{T}_{\tau,\sigma}^{\text{ifo}} \hat{\mathbf{q}}^\sigma \quad \text{and}$$

$$(2.9) \quad \mathbf{u}^\tau = \mathbf{T}_\tau^{\text{tfi}} \hat{\mathbf{v}}^\tau,$$

where $\mathbf{T}_\sigma^{\text{ofs}}$ is the $P \times N_\sigma$ *outgoing-from-sources* translation operator defined by (2.4), $\mathbf{T}_{\tau,\sigma}^{\text{ifo}}$ is the $P \times P$ *incoming-from-outgoing* translation operator defined by (2.6) and $\mathbf{T}_\tau^{\text{tfi}}$ is the $N_\tau \times P$ *target-from-incoming* translation operator defined by (2.5). With this, we have effectively computed a *rank-P* approximation of the potential

$$(2.10) \quad \mathbf{u}^\tau \approx \mathbf{T}_\tau^{\text{tfi}} \left(\mathbf{T}_{\tau,\sigma}^{\text{ifo}} \left(\mathbf{T}_\sigma^{\text{ofs}} \mathbf{q}^\sigma \right) \right),$$

which takes $\mathcal{O}(P(N_\sigma + N_\tau) + P^2)$ operations, much less than the $\mathcal{O}(N_\sigma N_\tau)$ operations of the naive method for (2.1).

Let us introduce the concept of *well separated* boxes, which gives us an estimate of the error committed in (2.10). We will say that a box Ω_τ of center \mathbf{c}_τ is *well separated* from a box Ω_σ of center \mathbf{c}_σ and side length $2a$ if

$$(2.11) \quad \|\mathbf{c}_\tau - \mathbf{c}_\sigma\| \geq 4a.$$

Now, for *well separated* source and target boxes of the same size, Ω_σ and Ω_τ , the error committed in (2.10) using a *rank-P* approximate potential \mathbf{u}_P^τ is given by [1]

$$(2.12) \quad \|\mathbf{u}^\tau - \mathbf{u}_P^\tau\|_\infty \leq \frac{\eta^P}{P(1-\eta)} \|\mathbf{q}^\sigma\|_1,$$

where $\eta = \sqrt{2}/(4 - \sqrt{2}) \approx 0.547$. This bound can be derived by considering the remainders of the Taylor expansions of (2.5) and (2.3), and the definition of *well separated* boxes. It follows that the error decays exponentially with respect to the interaction rank P , and we can pick $P = \mathcal{O}(\log(1/\epsilon))$ to achieve an error of ϵ .

3. The Fast Multipole Method for an uniform cluster of charges. Consider an uniform cluster of N points \mathbf{x}_i with associated charges q_i , $i = 1, \dots, N$, as shown in Figure 2. We wish to compute the electrostatic potential at every point \mathbf{x}_i , this is

$$(3.1) \quad u(\mathbf{x}_i) := u_i = \sum_{j=1}^N G(\mathbf{x}_i, \mathbf{x}_j) q_j, \quad i = 1, 2, \dots, N.$$

As before, computing the potentials using the naive formula (3.1) takes $\mathcal{O}(N^2)$ operations. Here, we will describe the Fast Multipole Method (FMM), which allows us to compute all potentials with only $\mathcal{O}(N)$, using the multipole technique presented in section 2.

In the FMM we construct a *quadtrees* of boxes, a hierarchical tree data structure, where we start with a *root* box which is divided into four smaller and identical *children* boxes. The children constitute the first level of the tree. To construct the second level, each child is again subdivided into four smaller and identical children. This process gets repeated until we reach a prescribed number of levels $L + 1$, where level 0 corresponds

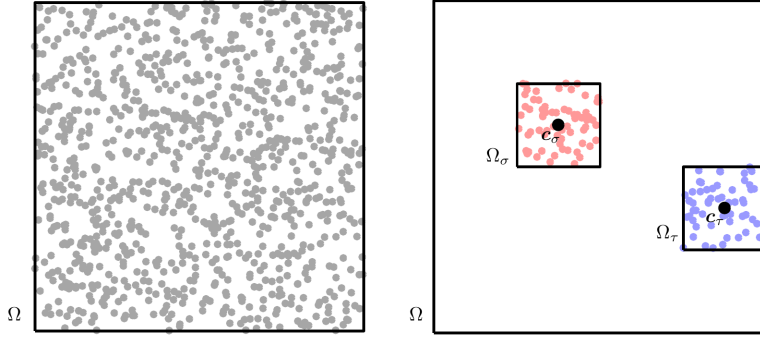


FIGURE 2. Left: uniform cluster of charges. Right: for well separated boxes inside the domain, we could resort to the fast multipole approach of [section 2](#).

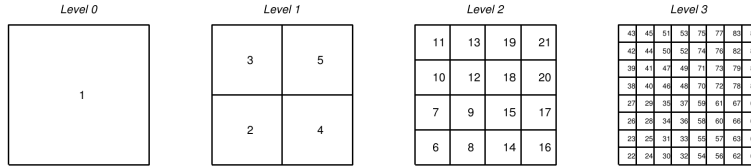


FIGURE 3. Quadtree of boxes with $L + 1$ levels, $L = 3$.

only to the root box, whereas level L contains the finer boxes, the *leaves*. An example of a quadtree structure with $L = 3$ is shown in [Figure 3](#).

The idea of this hierarchical data structure is the following. Leaf boxes compute the potentials at their points using multipole expansions for *well separated* boxes and direct evaluations for neighboring boxes and self-interactions. For *closer* well separated boxes we use the multipole expansion shown in [section 2](#). This set of *closer* well separated boxes is known as the *interaction list* $\mathcal{L}_\sigma^{\text{int}}$ of the leaf box σ , we will provide a precise definition below. Instead, for farther away boxes, it is much more efficient that the leaves' parents compute the multipole interactions (*outgoing* or *incoming* expansions) for the parent's *interaction list*, and then they pass the interaction to their children. Again, if two *well separated* parent boxes are very far away (meaning that they do not belong to each others *interaction list*) then it is more efficient that the parents' parents compute the interaction and so on. This hierarchical transfer of information is the secret behind FMM which allows to achieve a $\mathcal{O}(N)$ complexity. In the end, boxes only care about boxes in their interaction list. If they are farther away, then their ancestry (parents, grandparents, etc.) compute the interaction. And for leaf boxes, they have to care about their interaction list, neighbors and self interactions to finally compute the potentials at their points.

Before describing the FMM algorithm, we will explain how children/parents boxes transfer their *outgoing/incoming expansions* to each other. First, a parent box τ needs to compute its *outgoing expansion* \hat{q}^τ . In fact, this can be done exactly (without any approximation) using the *outgoing expansions* of their children, as shown next. Fix a child box σ of τ . Denote $I_\sigma = \{i : \mathbf{x}_i \in \Omega_\sigma\}$. Its outgoing expansion of order p is

$$(3.2) \quad \hat{q}_p^\sigma = \sum_{j \in I_\sigma} \frac{-1}{p} (\mathbf{x}_j - \mathbf{c}_\sigma)^p q_j,$$

where $p = 1, 2, \dots, P-1$. On the other hand, we can write the order $k = 1, 2, \dots, P-1$ *outgoing expansion* of τ given by the charges at Ω_σ as

$$(3.3) \quad \hat{q}_k^{\tau, \sigma} = \sum_{j \in I_\sigma} \frac{-1}{k} (\mathbf{y}_j - \mathbf{c}_\tau)^k q_j$$

$$(3.4) \quad = \sum_{j \in I_\sigma} \frac{-1}{k} ((\mathbf{y}_j - \mathbf{c}_\sigma) - (\mathbf{c}_\sigma - \mathbf{c}_\tau))^k q_j$$

$$(3.5) \quad = \sum_{j \in I_\sigma} \frac{-1}{k} \sum_{p=0}^k \binom{k}{p} (\mathbf{c}_\sigma - \mathbf{c}_\tau)^{k-p} (\mathbf{y}_j - \mathbf{c}_\sigma)^p q_j$$

$$(3.6) \quad = \frac{-1}{k} (\mathbf{c}_\sigma - \mathbf{c}_\tau)^k \hat{q}_0^\sigma + \sum_{p=1}^k \frac{-1}{k} \binom{k}{p} (\mathbf{c}_\sigma - \mathbf{c}_\tau)^{k-p} \sum_{j \in I_\sigma} (\mathbf{y}_j - \mathbf{c}_\sigma)^p$$

$$(3.7) \quad = \frac{-1}{k} (\mathbf{c}_\sigma - \mathbf{c}_\tau)^k \hat{q}_0^\sigma + \sum_{p=1}^k \frac{p}{k} \binom{k}{p} (\mathbf{c}_\sigma - \mathbf{c}_\tau)^{k-p} \sum_{j \in I_\sigma} \frac{-1}{p} (\mathbf{y}_j - \mathbf{c}_\sigma)^p$$

$$(3.8) \quad = \frac{-1}{k} (\mathbf{c}_\sigma - \mathbf{c}_\tau)^k \hat{q}_0^\sigma + \sum_{p=1}^k \frac{p}{k} \binom{k}{p} (\mathbf{c}_\sigma - \mathbf{c}_\tau)^{k-p} \hat{q}_p^\sigma.$$

This defines a linear operator $\mathbf{T}_{\tau, \sigma}^{\text{of}}$, known as the $P \times P$ *outgoing-from-outgoing* translation operator, which maps the *outgoing expansion* $\hat{\mathbf{q}}^\sigma$ of the child σ into the parent τ expansion $\hat{\mathbf{q}}^\tau$. This definition of $\mathbf{T}_{\tau, \sigma}^{\text{of}}$ is slightly different to the one shown in [2]. Finally, the *total outgoing expansion* of τ is given by the sum of the contributions of their children, this is

$$(3.9) \quad \hat{\mathbf{q}}^\tau = \sum_{\sigma \in \mathcal{L}_\tau^{\text{child}}} \mathbf{T}_{\tau, \sigma}^{\text{of}} \hat{\mathbf{q}}^\sigma,$$

where $\mathcal{L}_\tau^{\text{child}}$ is the set of children of τ .

Once a parent τ computes its *incoming expansion* $\hat{\mathbf{v}}^\tau$ using the *outgoing expansions* of the boxes in its interaction list, it needs to transfer this *incoming expansion* to each of its four children. Fix a children σ of τ . Its *incoming expansion* is given by

$$(3.10) \quad \hat{\mathbf{v}}^\sigma = \mathbf{T}_{\tau, \sigma}^{\text{if}} \hat{\mathbf{v}}^\tau + \sum_{\nu \in \mathcal{L}_\sigma^{\text{int}}} \mathbf{T}_{\sigma, \nu}^{\text{if}} \hat{\mathbf{q}}^\nu,$$

where the first term is the contribution from the parent (this includes the contribution of boxes farther away than the interaction list $\mathcal{L}_\sigma^{\text{int}}$ of σ), and the second term correspond to the contribution of the interaction list $\mathcal{L}_\sigma^{\text{int}}$ of σ . The linear operator $\mathbf{T}_{\sigma, \tau}^{\text{if}}$ is the $P \times P$ *incoming-from-incoming* translation operator, which maps the *incoming expansion* $\hat{\mathbf{v}}^\tau$ of the parent τ into the child σ expansion $\hat{\mathbf{v}}^\sigma$. Using the expression (2.5) and the same binomial expansion tricks as used in (3.8), it can be derived that

$$(3.11) \quad \mathbf{T}_{\sigma, \tau}^{\text{if}} = \begin{cases} \binom{p}{r} (\mathbf{c}_\sigma - \mathbf{c}_\tau)^{p-r}, & \text{for } r \leq p, \\ 0, & \text{for } r > p, \end{cases}$$

for $p, r = 0, 1, \dots, P-1$.

For a box τ let us define the following:

- The parent of τ is the box on the next coarsest level that contains τ .

- The children list of τ is the set $\mathcal{L}_\tau^{\text{child}}$ of boxes whose parent is τ .
- The neighbor list of τ is the set $\mathcal{L}_\tau^{\text{nei}}$ of boxes on the same level that directly touch τ .
- The interaction list of τ is the set $\mathcal{L}_\tau^{\text{int}}$ of all boxes σ such that (1) σ and τ are on the same level, (2) σ and τ do not touch, and (3) the parents of σ and τ do touch.
- Denote $A(I_\tau, I_\sigma)$ as the (dense) operator that maps the charges \mathbf{q}^σ at box σ into the potentials \mathbf{u}^τ at box τ , i.e., $\mathbf{u}^\tau = A(I_\tau, I_\sigma)\mathbf{q}^\sigma$.

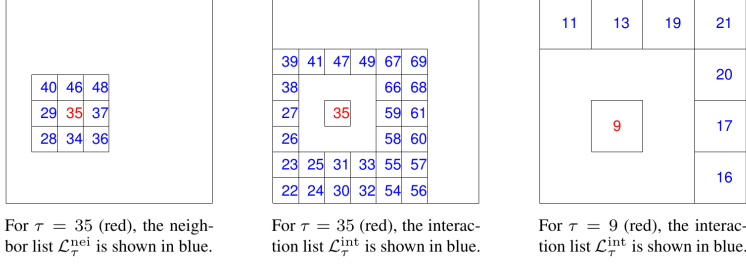


FIGURE 4. Examples of interaction and neighbor lists.

A diagram that shows an example of interaction and neighbor lists is presented in Figure 4. Now we have all the tools to describe the FMM algorithm. In the precomputation step we first assemble the quadtree, the boxes lists and precompute all of the translation operators. Now, to compute the potentials \mathbf{u} for a given set of charges \mathbf{q} we follow three steps:

1. Upward pass: compute the *outgoing expansions* of all boxes. Start at the finest level computing the expansions directly from the charges. For parent boxes, compute their expansions using their children's expansions.
2. Downward pass: compute the *incoming expansions* of all boxes. From level $l = 2$ to the finest, compute the *incoming expansions* using the parent's contribution and the interaction list's contribution.
3. Compute potentials: Each leaf computes the potentials at its points, using the incoming expansions and direct evaluations of the potentials for neighboring boxes and self-interactions.

The FMM algorithm is detailed in Figure 5, Figure 6 and Figure 7.

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loop over levels  $\ell$ , from the finest to the coarsest
  loop over all boxes  $\tau$  on level  $\ell$ 
    if ( $\tau$  is a leaf)
       $\hat{\mathbf{q}}_\tau = \mathbf{T}_\tau^{\text{ofs}} \mathbf{q}(I_\tau)$ 
    else
       $\hat{\mathbf{q}}_\tau = \sum_{\sigma \in \mathcal{L}_\tau^{\text{child}}} \mathbf{T}_{\tau,\sigma}^{\text{ofc}} \hat{\mathbf{q}}_\sigma$ 
    end if
  end loop
end loop

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FIGURE 5. FMM upward pass.

4. Error estimation. The error for FMM type of methods is usually estimated using a single translation operation, which includes the processes that maps charges

Set $\hat{\mathbf{u}}_\tau = \mathbf{0}$ for every box τ on level 1.
loop over levels ℓ , from level $\ell = 2$ to the finest
 loop over all boxes τ on level ℓ
 Let ν denote the parent of τ .
 $\hat{\mathbf{u}}_\tau = \mathbf{T}_{\tau,\nu}^{\text{ifi}} \hat{\mathbf{u}}_\nu + \sum_{\sigma \in \mathcal{L}_\tau^{\text{int}}} \mathbf{T}_{\tau,\sigma}^{\text{ifo}} \hat{\mathbf{q}}_\sigma$.
 end loop
end loop

FIGURE 6. FMM downward pass.

loop over every leaf box τ
 $\mathbf{u}(I_\tau) = \mathbf{A}_\tau^{\text{tff}} \hat{\mathbf{u}}_\tau + \mathbf{A}(I_\tau, I_\tau) \mathbf{q}(I_\tau) + \sum_{\sigma \in \mathcal{L}_\tau^{\text{nei}}} \mathbf{A}(I_\tau, I_\sigma) \mathbf{q}(I_\sigma)$
end loop

FIGURE 7. FMM compute potentials step.

into *outgoing expansions*, into *incoming expansions* and finally into potentials. As shown in [section 2](#), this error ϵ decays exponentially fast with respect to the interaction rank P , i.e., $P = \mathcal{O}(\log(1/\epsilon))$. A more careful analysis and a sharper bound is given in [\[3\]](#).

5. Cost estimation. Let N be the total number of charges, L the number of levels ($L + 1$ levels in total) and b the average number of charges in a leaf box ($b \approx N/4^L$, where 4^L is the number of leaf boxes, recall that the distribution of charges is uniform). Let us compute the cost of applying each of the translation operators:

- The *outgoing-from-sources* $\mathbf{T}_\sigma^{\text{ofs}}$ operator is of size $P \times b$ and is applied once for every leaf box. There are N/b leaf boxes, so $\mathbf{T}_\sigma^{\text{ofs}} \sim Pb \times N/b = PN$.
- The *outgoing-from-outgoing* $\mathbf{T}_{\tau,\sigma}^{\text{offo}}$ operator is of size $P \times P$ and is applied once for every box. The total number of boxes is bounded by twice the number of leaf boxes, hence $\mathbf{T}_{\tau,\sigma}^{\text{offo}} \sim P^2 \times 2N/b \sim P^2 N/b$.
- The *outgoing-from-incoming* $\mathbf{T}_{\tau,\sigma}^{\text{ofifi}}$ operator and the *incoming-from-incoming* $\mathbf{T}_{\sigma,\tau}^{\text{ififi}}$ operator follow a similar bound as the *outgoing-from-outgoing* $\mathbf{T}_{\tau,\sigma}^{\text{offo}}$ operator, therefore $\mathbf{T}_{\tau,\sigma}^{\text{ofifi}} \sim \mathbf{T}_{\sigma,\tau}^{\text{ififi}} \sim \mathbf{T}_{\tau,\sigma}^{\text{offo}} \sim P^2 N/b$.
- The cost of the *target-from-incoming* $\mathbf{T}_\tau^{\text{tffi}}$ operator is the same as the *outgoing-from-outgoing* $\mathbf{T}_{\tau,\sigma}^{\text{offo}}$ operator, so $\mathbf{T}_\tau^{\text{tffi}} \sim \mathbf{T}_{\tau,\sigma}^{\text{offo}} \sim P^2 N/b$.
- For the near interactions, i.e., neighbor and self-interactions, each particle has to compute at most $9b$ interactions, since there are at most 9 neighboring boxes. Thus, $\mathbf{T}^{\text{near}} \sim Nb$.

Adding up all costs, we end up with a total cost of $\mathbf{T}^{\text{FMM}} \sim P^2 N/b + PN + Nb$. Since we are free to choose b by tuning the numbers of levels L , we can choose $b \sim P$ so that the cost is minimized,

$$(5.1) \quad \mathbf{T}^{\text{FMM}} \sim PN \sim N \log(1/\epsilon),$$

where we used the fact $P \sim \log(1/\epsilon)$, with ϵ being the error.

6. Numerical results.

REFERENCES

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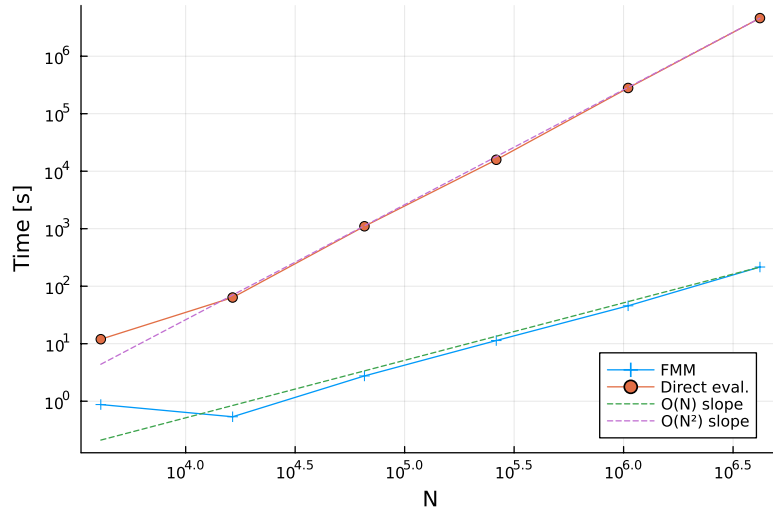


FIGURE 8.

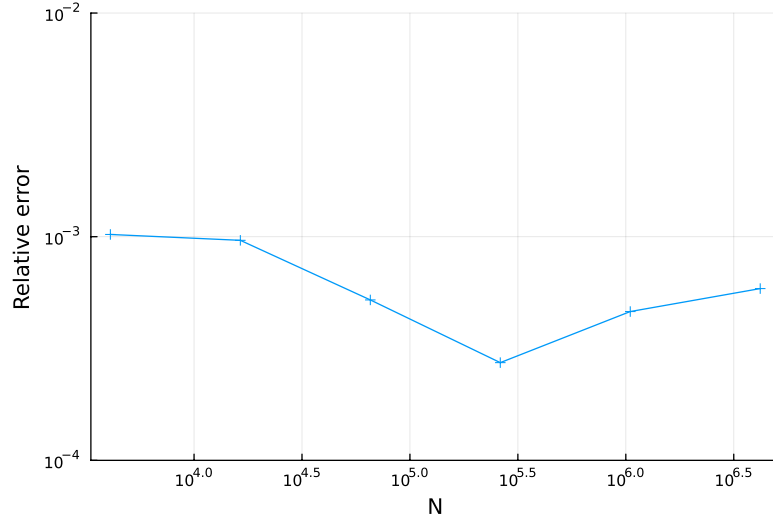


FIGURE 9.

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