New productions for Pat

```
Pat ::= '[' ']' | Pat '::' Pat | '[' Pat (';' Pat)* ']'
```

Remark

- all variables in a pattern must be distinct
 - this makes pattern matching more efficient
- patterns are built with constructors, but not with operators
 - \times x::y is a valid pattern, x@y or x+y is not
 - constructors guarantee unique decomposition of values

What is pattern matching?

Examples:

```
let add (x,y) = x+y;;
add (3,5);; (* does (3,5) match with pattern (x,y)? *)
```

- (3,5) matches with pattern (x,y) if and only if x=3 and y=5
- if we substitute x and y in (x,y) with 3 and 5, respectively, then we obtain the value (3,5)

```
let hd (h::t) = h;; (* returns the head of the list *)
hd [3;5];; (* does [3;5] match with pattern h::t? *)
```

- [3;5] matches with pattern (h::t) if and only if h=3 and t=[5]
- if we substitute h and t in (h::t) with 3 and [5], respectively, then we obtain the value (3::[5]) = [3;5]

Do we always need to use a single pattern to match all values?

```
let hd (h::t) = h;; (* returns the head of the list *)
hd [];; // does [] match with pattern h::t?
```

- [] does not match with (h::t) for any value associated with h and t
- indeed [] \neq (h::t) for all possible values associated with h and t
- this is correct, because the head of a list is undefined for the empty list

Functions that cannot be directly defined with a single pattern

- the length of a list
- the sum of all the elements of a list
- the list with the first two elements swapped

An expression to match values with multiple patterns

```
Exp ::= 'match' Exp 'with' Pat '->' Exp ('|' Pat '->' Exp) \star
```

Examples

```
(* defined with two patterns *)
let rec length l = match l with
   [] -> 0
   | hd::tl -> l+length tl;;

let rec sum l = match l with
   [] -> 0
   | hd::tl -> hd+sum tl;; (* hd and tl are local variables *)
(* defined with more than two patterns *)
let swap l = match l with
   [] -> []
   | [x] -> [x] (* x is a local variable *)
   | x::y::l -> y::x::l;; (* x, y and l are local variables *)
```

```
match e with p_1 \rightarrow e_1 \mid \ldots \mid p_n \rightarrow e_n
```

Static semantics

- the expression e and all patterns $p_1 \dots p_n$ must have the same type
- all expressions $e_1 \dots e_n$ must have the same type

Dynamic semantics

- e is evaluated
- all patterns $p_1 \dots p_n$ tried from left to right, top to bottom
- at the first match with p_i, the expression e_i is evaluated, with variables defined by the match with p_i
- if there is no match, then Match_failure is raised

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Static semantics: further checks

A warning is reported if:

- patterns are not exhaustive, that is, some case is missing
- a pattern is unused

Example

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Unique decomposition

Constructors ensure that if there is a match with p, then there exists a unique substitution for the variables in p

```
Counter-example
```

```
# fun ls -> match ls with l1@12 -> l1;; (* @ is not a constructor! *)
Error: Syntax error

If ls is [1;2;3] what would be the value of l1???
[]???
[1]???
[1;2]???
[1;2;3]???
```

Constructors for primitive types

All literals (=tokens that represent values) are constant constructors

Example of pattern matching with primitive types

```
let mynot b = match b with false -> true | true -> false;;
let iszero i = match i with 0 -> true | any -> false;;
(* a variable is needed for the second pattern *)
```

Remark

Pattern matching with primitive types is seldom used

Shorthand notation

- the wildcard _ is the pattern which matches all values when no variable is needed
- function $p_1 \rightarrow e_1 \mid \ldots \mid p_n \rightarrow e_n$ is a shorthand for fun $var \rightarrow match \ var \ with \ p_1 \rightarrow e_1 \mid \ldots \mid p_n \rightarrow e_n$
- p as id: a pattern (or sub-pattern) p can be associated with an id to refer to the matched value more directly

Examples