

# Curried/uncurried functions

## Definition

- Curried function (from Haskell Curry): a higher-order function with a single argument returning a chain of functions with a single argument
- Uncurried function: a function with multiple arguments

## Facts

- an uncurried function can be transformed in the equivalent curried version
- a curried function can be transformed in the equivalent uncurried version

## Examples

```
(* addition of two integers *)  
fun x y->x+y;;           (* curried version int->int->int *)  
fun (x,y)->x+y;;         (* uncurried version int*int->int *)  
  
(* multiplication of three integers *)  
fun x y z->x*y*z;;       (* curried version int->int->int->int *)  
fun (x,y,z)->x*y*z;;     (* uncurried version int*int*int->int *)
```

# Curried/uncurried functions

## Partial application

- curried functions allow *partial* application: arguments can be passed once at time
- uncurried functions do not allow partial application: all arguments must be passed altogether

## Example

```
let curried_add x y=x+y;;  
let uncurried_add(x,y)=x+y;;  
(* computes 1+2 with the uncurried version *)  
uncurried_add(1,2);;  
(* computes 1+2 by partial application *)  
let inc=curried_add 1;; (* passes argument 1 and saves the result *)  
inc 2;; (* passes argument 2 and computes the final result *)
```

# Curried/uncurried functions

## Partial application promotes *generic programming*

Partial application allows function specialization: from a generic function it is possible to generate more specific ones with *no code duplication*.

- *software reuse* and *maintenance* are favored
- interesting examples will be shown later on

# Boolean values

## Syntax

```
Exp ::= BOOL | 'not' Exp | Exp '&&' Exp | Exp '||' Exp  
Type ::= 'bool'
```

BOOL **boolean values** **false** | **true**

## Standard syntactic rules

- **&&** and **||** are left-associative
- **not** higher precedence than **&&**
- **&&** higher precedence than **||**

# Boolean values

## Static semantics

- **false** and **true** have type `bool`
- **not**  $e$  has type `bool` if and only if  $e$  has type `bool`
- **not**  $e$  is **not** type correct if either  $e$  has type  $\neq \text{bool}$  or  $e$  is not type correct
- $e_1 \&\& e_2$  and  $e_1 || e_2$  have type `bool` if and only if  $e_1$  and  $e_2$  have type `bool`
- $e_1 \&\& e_2$  and  $e_1 || e_2$  are **not** type correct if either  $e_1$  or  $e_2$  has type  $\neq \text{bool}$  or  $e_1$  or  $e_2$  is **not** type correct

# Boolean values

## Standard semantics

- operands of `&&` and `||` evaluated left-to-right with “short circuit”
- if  $e_1$  evaluates to `false` then  $e_1 \&\& e_2$  evaluates to `false`, else it evaluates to the value of  $e_2$
- if  $e_1$  evaluates to `true` then  $e_1 || e_2$  evaluates to `true`, else it evaluates to the value of  $e_2$

# Boolean values

## Conditional expression

$\text{Exp} ::= \text{'if' Exp 'then' Exp 'else' Exp}$

Conditional expression has precedence lower than all other operators

## Static semantics

- **if**  $e$  **then**  $e_1$  **else**  $e_2$  has type  $t$  if and only if  
   $e$  has type `bool` and  $e_1$  and  $e_2$  have type  $t$
- **if**  $e$  **then**  $e_1$  **else**  $e_2$  is **not** type correct if
  - ▶  $e$  has type  $\neq \text{bool}$
  - ▶ or there is no type  $t$  such that  $e_1$  and  $e_2$  have type  $t$
  - ▶ or  $e$  or  $e_1$  or  $e_2$  is **not** type correct

# More on declarations of global “variables”

## Grammar

```
Dec ::= 'let' Def ('and' Def)*  
      | 'let' 'rec' FunDef ('and' FunDef)*  
Def  ::= Pat '=' Exp | FunDef  
FunDef ::= ID Pat* '=' Exp
```

## Remark

- recursive declarations allowed only for function types and other types
- for simplicity we consider only recursive declarations of functions

## Example

```
let rec sumsquare n = (*sumsquare can be used on the right-hand side*)  
  if n<=0 then 0 else n*n+sumsquare(n-1);;
```



# Curried functions and generic programming

## Example 1: addition of square numbers

```
let rec sumsquare n = (*sumsquare can be used on the right-hand side*)  
  if n<=0 then 0 else n*n+sumsquare(n-1);;
```

## Example 2: addition of cube numbers

```
let rec sumcube n = (*sumcube can be used on the right-hand side*)  
  if n<=0 then 0 else n*n*n+sumcube(n-1);;
```

## Remarks

- the two examples are almost identical!
- can we improve code reuse and maintenance?

**Solution: use a curried function with an argument of type function**

# Curried functions and generic programming

## Solution

```
let rec gen_sum f n = (* (int -> int) -> int -> int *)
  if n<=0 then 0 else f n+gen_sum f (n-1);;

let der_sumsquare = gen_sum (fun x->x*x);; (* int -> int *)
let der_sumcube = gen_sum (fun x->x*x*x);; (* int -> int *)
```

## Remarks

`gen_sum` can be specialized because

- it is curried
- the “first” argument is `f` rather than `n`

# Declarations of local “variables”

## Syntax

```
Dec ::= 'let' Def ('and' Def)* 'in' Exp
      | 'let' 'rec' FunDef ('and' FunDef)* 'in' Exp
Def  ::= Pat '=' Exp | FunDef
FunDef ::= ID Pat* '=' Exp
```

## Example

```
# let f x=x+1 and v=41 in f v;; (* f and v can only be used here *)
- : int = 42
# let x=1 in let x=x*2 in x*x (* nested declarations *)
- : int = 4
```

## Remark

Nested declarations overrides outer declarations with the same ID

# Static scope of declarations

## Example

```
let v=40;;
```

```
let f x = x*v;; (* v refers to the declaration above *)
```

```
f 3;; (* evaluates to 120 *)
```

```
let v=4;; (* declaration of v overridden *)
```

```
f 3;; (* evaluates to 120 *)
```

# Curried functions and generic programming (revisited)

## A slightly better solution

```
let gen_sum f = (* (int -> int) -> int -> int *)  
  let rec aux n = if n<=0 then 0 else f n+aux (n-1) (* int -> int *)  
  in aux;;
```

## Remarks

We do not have to pass argument  $f$  to the recursive function `aux`