

Are there other (possibly better) solutions?

Observation

Although ambiguous, the infix notation is a more intuitive and practical solution!

Elimination of ambiguity of infix notation

- define associativity rules for binary operators
 - ▶ addition is left-associative: " $1+1+1$ " means " $(1+1)+1$ "
 - ▶ addition is right-associative: " $1+1+1$ " means " $1+(1+1)$ "
- define precedence rules for operators, use parentheses to override them
 - ▶ multiplication has higher precedence over addition: " $1+1*1$ " means " $1+(1*1)$ "
 - ▶ addition has higher precedence over multiplication: " $1*1+1$ " means " $1*(1+1)$ "

More on precedence and associativity rules

Operators with the same precedence

- binary operators can have the same precedence; in this case they also share the same associativity rule
 - ▶ addition and multiplication have the same precedence and are left-associative: " $1+1*1$ " means " $(1+1)*1$ " and " $1*1+1$ " means " $(1*1)+1$ "
 - ▶ addition and multiplication have the same precedence and are right-associative: " $1+1*1$ " means " $1+(1*1)$ " and " $1*1+1$ " means " $1*(1+1)$ "

Remark on associativity rules

Associativity rules resolve ambiguities between binary operators with the same precedence

Operators with different arities

Mixing together operators of different arities (typically 1, 2 and 3) makes elimination of ambiguity more complex!

Standard techniques for grammar disambiguation

Details

- an ambiguous grammar G is transformed into an *equivalent non-ambiguous* grammar G'
- *equivalent* means that for all non-terminal B of G , the languages generated by G and G' from B are *equal*
- the transformation is able to encode associativity and precedence rules in the non-ambiguous grammar G'

Example 1: + and * with the same precedence

Ambiguous grammar

```
Exp ::= Num | Exp '+' Exp | Exp '*' Exp | '(' Exp ')'  
Num ::= '0' | '1'
```

Non-ambiguous grammar, left associative operations

```
Exp ::= Atom | Exp '+' Atom | Exp '*' Atom  
Atom ::= Num | '(' Exp ')'  
Num ::= '0' | '1'
```

Non-ambiguous grammar, right associative operations

```
Exp ::= Atom | Atom '+' Exp | Atom '*' Exp  
Atom ::= Num | '(' Exp ')'  
Num ::= '0' | '1'
```

Example 1: + and * with the same precedence

Non-ambiguous grammar, left associative operations

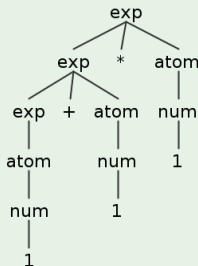
$\text{Exp} ::= \text{Atom} \mid \text{Exp} \text{ '+' } \text{Atom} \mid \text{Exp} \text{ '*' } \text{Atom}$

$\text{Atom} ::= \text{Num} \mid \text{'(' Exp ') '}$

$\text{Num} ::= \text{'0' } \mid \text{'1'}$

Remark: $\text{Exp} \text{ '+' } \text{Atom} (\text{Exp} \text{ '*' } \text{Atom})$ means that on the right side of + (*) additions (multiplications) are allowed only if surrounded by parentheses

Unique derivation tree for $1+1*1$



Example 1: + and * with the same precedence

Non-ambiguous grammar, right associative operations

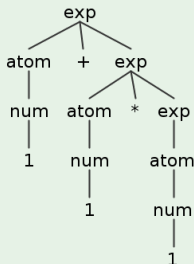
$\text{Exp} ::= \text{Atom} \mid \text{Atom} \text{ '+' } \text{Exp} \mid \text{Atom} \text{ '*' } \text{Exp}$

$\text{Atom} ::= \text{Num} \mid \text{'(' Exp ') '}$

$\text{Num} ::= \text{'0'} \mid \text{'1'}$

Remark: $\text{Atom} \text{ '+' } \text{Exp}$ ($\text{Atom} \text{ '*' } \text{Exp}$) means that on the left side of + (*) additions (multiplications) are allowed only if surrounded by parentheses

Unique derivation tree for $1+1*1$



Example 2: * with higher precedence

Ambiguous grammar

```
Exp ::= Num | Exp '+' Exp | Exp '*' Exp | '(' Exp ')'  
Num ::= '0' | '1'
```

Non-ambiguous grammar, left associative operations

```
Exp ::= Mul | Exp '+' Mul  
Mul ::= Atom | Mul '*' Atom  
Atom ::= Num | '(' Exp ')'  
Num ::= '0' | '1'
```

Non-ambiguous grammar, right associative operations

```
Exp ::= Mul | Mul '+' Exp  
Mul ::= Atom | Atom '*' Mul  
Atom ::= Num | '(' Exp ')'  
Num ::= '0' | '1'
```

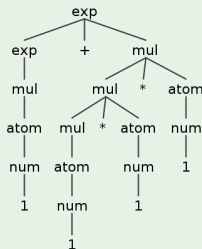
Example 2: $*$ with higher precedence

Non-ambiguous grammar, left associative operations

```
Exp ::= Mul | Exp '+' Mul
Mul ::= Atom | Mul '*' Atom
Atom ::= Num | '(' Exp ') '
Num ::= '0' | '1'
```

Remark: `Mul '*' Atom` means that on both sides of $*$ additions are allowed only if surrounded by parentheses

Unique derivation tree for $1+1*1*1$



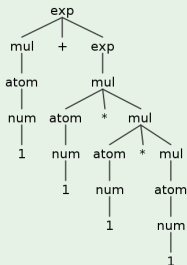
Example 2: * with higher precedence

Non-ambiguous grammar, right associative operations

```
Exp ::= Mul | Mul '+' Exp
Mul ::= Atom | Atom '*' Mul
Atom ::= Num | '(' Exp ') '
Num ::= '0' | '1'
```

Remark: Atom '*' Mul means that on both sides of * additions are allowed only if surrounded by parentheses

Unique derivation tree for $1+1*1*1$



Remaining examples

Solutions left for exercise

- $*$ higher precedence and left associative, $+$ right associative
- $*$ higher precedence and right associative, $+$ left associative
- $+$ higher precedence and left associative, $*$ right associative
- $+$ higher precedence and right associative, $*$ left associative

The non-ambiguous grammars can be easily defined by symmetry

Ambiguous syntax for statements

Ambiguity: not only expressions ...

$\text{Stmt} ::= \text{ID '=' Exp} \mid \text{'if' '(' Exp ')' Stmt} \mid \text{Stmt ';' Stmt}$
 $\mid \text{'{' Stmt '}'}$

$\text{Exp} ::= \text{ID} \mid \text{BOOL}$ // *ID and BOOL defined by regular expressions*

Two derivation trees for "if (x) x=false; y=true"

