

Exercise 1 (value 12)

Due to the **Energy Crisis** Europe is looking for a plan to gradually replace the supplies of gas from Russia with alternative sources. Today Europe buys, each month, Q billions of cubic meters from Russia. In the next T months this quantity must be reduced to q billions with a linear decrease. At the same time the quantity not bought from Russia must be bought from other suppliers. There are n possible suppliers. In month $t = 1, \dots, T$ each supplier i can sell v_{it} billions of cubic meters of gas at a price of c_{it} for billion. Moreover, if the overall gas bought in a month is more than the planned reduction, there is the possibility to store the exceeding cubic meters at cost of s_t for billion/month. The maximum storage for a month is S billions. At current month (0) no storage exists, but at the end of month T a storage of $\bar{S} \leq S$ billions must be guaranteed.

Write a linear mathematical model to help Europe to define the cheaper way to buy gas in the next T months, while reducing the quantities bought from Russia.

Modify the above model by adding the following constraint: to avoid dependence from a single supplier, Europe decided to buy, in the T months, from at least 3 suppliers.

Answer

Variables

x_{it} = billions of cubic meters of gas bought from supplier i in month t

I_t = billions of cubic meters of gas stored at the end of month t

y_i = 1 if supplier i is used

Constants

Δ = $\frac{1}{T}(Q - q)$ monthly reduction of gas from Russia

M a big number

$$\begin{aligned}
 \min \quad & \sum_{i=1}^n \sum_{t=1}^T c_{it} x_{it} + \sum_{t=1}^T s_t I_t \\
 & I_{t-1} + \sum_{i=1}^n x_{it} = t\Delta + I_t \quad t = 1, \dots, T \\
 & I_0 = 0 \\
 & I_T \geq \bar{S} \\
 & \sum_{i=1}^n y_i \geq 3 \\
 & y_i \leq M \sum_{t=1}^T x_{it} \quad i = 1, \dots, n \\
 & 0 \leq I_t \leq S \quad t = 1, \dots, T \\
 & 0 \leq x_{it} \leq v_{it} \quad i = 1, \dots, n, t = 1, \dots, T
 \end{aligned}$$

Exercise 2 (value 8)

Consider the following LP problem.

$$\begin{aligned}
 \min \quad & 4x_1 + 5x_2 + 3x_3 \\
 & 2x_1 + 3x_2 + x_3 \geq 3 \\
 & 3x_1 + 5x_2 + 4x_3 \geq 1 \\
 & x_1, x_2, x_3 \geq 0
 \end{aligned}$$

- Write the corresponding tableau.
- Manipulate the tableau, if necessary, and solve the problem with the most convenient (for this case) simplex method, following the Bland's rule.

- Expose your motivations for the chosen method.
- Write the optimal solution and its cost.

Answer

x_1	x_2	x_3	y_1	y_2	$-z$	
4	5	3	0	0	0	
2	3	1	-1	0	3	y_1
3	5	4	0	-1	1	y_2

Since all the reduced costs are positive, it is worth to multiply both constraints by -1 and apply the Dual simplex method.

x_1	x_2	x_3	y_1	y_2	$-z$	
4	5	3	0	0	0	
-2	-3	-1	1	0	-3	y_1
-3	-5	-4	0	1	-1	y_2

x_1	x_2	x_3	y_1	y_2	$-z$	
$\frac{2}{3}$	0	$\frac{4}{3}$	$\frac{5}{3}$	0	-5	
$\frac{2}{3}$	1	$\frac{1}{3}$	$-\frac{1}{3}$	0	1	x_2
$\frac{1}{3}$	0	$-\frac{7}{3}$	$-\frac{5}{3}$	1	4	y_2

The optimal solution is $x = (0, 1, 0)$ with cost 5.

Exercise 3 (value 7)

Find the optimal solution of the following knapsack problem using a Dynamic Programming method, without reordering the items: $n = 4, c = 32, (p_j, w_j) = [(1, 18), (2, 25), (3, 14), (2, 3)]$

Answer

$$P = 1 + 2 + 3 + 2 = 8$$

	0	1	2	3	4	5	6	7	8
f_0	0	∞	∞	∞	∞	∞	∞	∞	∞
f_1	0	18	∞	∞	∞	∞	∞	∞	∞
f_2	0	18	25	43	∞	∞	∞	∞	∞
f_3	0	18	25	14	32	39	57	∞	∞
f_4	0	18	3	14	28	17	35	42	60

	0	1	2	3	4	5	6	7	8
J_0	\emptyset	\emptyset	\emptyset	\emptyset	\emptyset	\emptyset	\emptyset	\emptyset	\emptyset
J_1	\emptyset	{1}	\emptyset	\emptyset	\emptyset	\emptyset	\emptyset	\emptyset	\emptyset
J_2	\emptyset	{1}	{2}	{1, 2}	\emptyset	\emptyset	\emptyset	\emptyset	\emptyset
J_3	\emptyset	{1}	{2}	{3}	{1, 3}	{2, 3}	{1, 2, 3}	\emptyset	\emptyset
J_4	\emptyset	{1}	{4}	{3}	{2, 4}	{3, 4}	{1, 3, 4}	{2, 3, 4}	{1, 2, 3, 4}

The optimal solution $x = \{3, 4\}$ has profit 5 and weight 17.