



Exercise 1 (value 9)

A telecommunication company needs to install antennae for the new 6G technology. A set L of *locations* where antennae can be installed (maximum one antenna per location) has been identified. A set A of *areas* that the company aims to cover with the signal is also given. The signal intensity s_{ij} , namely the strength of the signal from an antenna installed in location $i \in L$ received in area $j \in A$, is known for each pair (i, j) . An area j is classified as *covered* only when the sum of the strength of the signals received is at least T . Furthermore, due to interference issues, an area $j \in A$ is declared as *non-covered* if there is more than one signal with level greater than or equal to U (i.e., $s_{ij} \geq U$) in j . Finally, for technical reasons, an antenna can be placed in location α only if an antenna is installed in location β . Write an integer linear programming model to help the company to decide where to install antennae in order to maximize the number of *covered* areas.

Exercise 2 (value 9)

Solve the following LP problem with a simplex method of your choice:

$$\begin{aligned} \min \quad & 4x_1 + x_2 + x_3 \\ & x_1 + 4x_2 - 10x_3 \leq 4 \\ & -x_1 - 2x_2 + 5x_3 = 5 \\ & x_1, x_2, x_3 \geq 0. \end{aligned}$$

State the method you choose, show all the steps, and write the optimal solution together with its cost.

Exercise 3 (value 9)

Consider the graph $G = (V, A)$ described by the following adjacency matrix with costs:

	1	2	3	4	5	6
1	.	2	3	5	.	8
2	5	.	.	2	7	.
3	5	-2
4	5	1
5	6	3	4	.	.	1
6	1	12	14	.	1	.

Find the shortest path from node 1 to node 5 by using the Bellman-Ford algorithm. Show the procedure (show the tables of labels and predecessors) and write the optimal path with the associated cost.



Solution sketch

Exercise 1

$x_i = 1$ if an antenna is installed in location $i \in L$; 0 otherwise

$y_j = 1$ if the area $j \in A$ is covered; 0 otherwise

$$\begin{aligned}
 \max \quad & \sum_{j \in A} y_j \\
 \sum_{i \in L} s_{ij} x_i & \geq T y_j & j \in A \\
 \sum_{i \in L: s_{ij} \geq U} x_i & \leq 1 + |L|(1 - y_j) & j \in A \\
 x_\alpha & \leq x_\beta \\
 x_i & \in \{0, 1\} & i \in L \\
 y_j & \in \{0, 1\} & j \in A
 \end{aligned}$$

Exercise 2

Dual simplex method:

$$\begin{aligned}
 \min \quad & 4x_1 + x_2 + x_3 \\
 & x_1 + 4x_2 - 10x_3 + y_1 = 4 \\
 & -x_1 - 2x_2 + 5x_3 + y_2 = 5 \\
 & x_1 + 2x_2 - 5x_3 + y_3 = -5 \\
 & x_1, x_2, x_3, y_1, y_2, y_3 \geq 0.
 \end{aligned}$$

x_1	x_2	x_3	y_1	y_2	y_3	$-z$	
4	1	1	0	0	0	0	
1	4	-10	1	0	0	4	y_1
-1	-2	5	0	1	0	5	y_2
1	2	-5	0	0	1	-5	y_3
x_1	x_2	x_3	y_1	y_2	y_3	$-z$	
21/5	7/5	0	0	0	1/5	-1	
-1	0	0	1	0	-2	14	y_1
0	0	0	0	1	1	0	y_2
-1/5	-2/5	1	0	0	-1/5	1	x_3

Alternatively, two phases method:

x_1	x_2	x_3	y_1	a_1	$-z$	
0	0	0	0	1	0	
1	4	-10	1	0	4	y_1
-1	-2	5	0	1	5	a_1
x_1	x_2	x_3	y_1	a_1	$-z$	
1	2	-5	0	0	-5	
1	4	-10	1	0	4	y_1
-1	-2	5	0	1	5	a_1

x_1	x_2	x_3	y_1	a_1	$-z$	
0	0	0	0	1	0	
-1	0	0	1	2	14	y_1
-1/5	-2/5	1	0	1/5	1	x_3

x_1	x_2	x_3	y_1	z	
1/5	2/5	0	0	-1	
-1	0	0	1	14	y_1
-1/5	-2/5	1	0	1	x_3

The optimal solution is $x = \{0, 0, 1\}$ with cost 1.

Exercise 3

k	$f^k(j)$						$Pred_j$					
	1	2	3	4	5	6	1	2	3	4	5	6
0	0	—	—	—	—	—	1	—	—	—	—	—
1	0	2	3	5	—	8	1	1	1	1	—	1
2	0	1	3	4	9	6	1	3	1	2	2	4
3	0	1	3	3	7	5	1	3	1	2	6	4
4	0	1	3	3	6	4	1	3	1	2	6	4
5	0	1	3	3	5	4	1	3	1	2	6	4

The optimal path is $\{1, 3, 2, 4, 6, 5\}$ and has cost 5.