

Part 30

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Transportation problem

Are given

- *n sources* with maximum capacity a_i , i = 1, ..., n
- ullet m sinks with a request $b_j, j=1,\ldots,m$
- c_{ij} = transport cost of a unit from i to j

Problem: Find the set of transports giving to each sink the required amount, without exceeding the sources' capacity and minimizing the total cost.

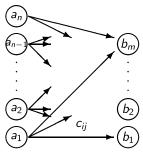


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 $x_{ij} =$ quantity transported from i to j

$$\min \sum_{i=1}^n \sum_{j=1}^n c_{ij} x_{ij}$$
 $\sum_{j=1}^n x_{ij} \leq a_i, \quad i=1,\ldots,n$ $\sum_{i=1}^n x_{ij} \geq b_j, \quad j=1,\ldots,n$ $x_{ij} \geq 0 \quad \text{integer} \quad i,j=1,\ldots,n$

Linear assignment

Given an $n \times n$ square matrix $C = [c_{ij}]$, assign each row to a column, minimizing the sum of the selected elements and such that each column is assigned to a unique row.

(given n workers, n works and the cost c_{ij} incurred if work j is done by worker i, assign each worker to one work, minimizing the cost.)



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(given n workers, n works and the cost c_{ij} incurred if work j is done by worker i, assign each worker to one work, minimizing the cost.)

$$x_{ij} = \left\{ egin{array}{ll} 1 & ext{if row } i ext{ is assigned to column } j \\ 0 & ext{otherwise} \end{array}
ight.$$

$$\min \sum_{i=1}^{n} \sum_{j=1}^{n} c_{ij} x_{ij}$$

$$\sum_{j=1}^{n} x_{ij} = 1, \quad i = 1, \dots, n$$

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$$x_{ij} \in \{0, 1\} \quad i, j = 1, \dots, n$$







Knapsack

- *n* items
- Each item j has a profit p_i and a size w_i
- The knapsack has capacity c

Problem: Find the set of items that can be inserted in the knapsack maximizing the profit



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$$\max \sum_{j=1}^n p_j x_j$$

$$\sum_{j=1}^n w_j x_j \leq c$$

$$x_j \in \{0,1\} \ j=1,\ldots,n$$



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Bin packing

- *n* items
- Each item j has a size w_j
- Infinite number of bins of capacity c

Problem: Pack all the items in the minimum number of bins



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$$x_{ij} = \left\{ egin{array}{ll} 1 & ext{if item j is packed in bin i} \\ 0 & ext{otherwise} \end{array} \right. \quad z_i = \left\{ egin{array}{ll} 1 & ext{if bin i is used} \\ 0 & ext{otherwise} \end{array} \right.$$

min
$$\sum_{i=1}^n z_i$$

$$\sum_{j=1}^n w_j x_{ij} \leq c z_i, \quad i=1,\ldots,n$$

$$\sum_{i=1}^n x_{ij} = 1, \quad j=1,\ldots,n$$

$$x_{ij} \in \{0,1\} \; i,j=1,\ldots,n$$

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Facility location

- $ightharpoonup N = \{1, ..., n\}$ potential facility locations
- ▶ $I = \{1, ..., m\}$ customers' locations
- ▶ Site $j \in N$ has capacity u_i and activation cost c_i
- ▶ b_i , $i \in I$ is the request of customer i
- \blacktriangleright h_{ij} cost for serving customer i from site j
- ► Each customer must be served by a unique facility.



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$$x_j = \left\{ \begin{array}{ll} 1 & \text{if a facility is set up in site } j \\ 0 & \text{otherwise} \end{array} \right. \\ y_{ij} = \left\{ \begin{array}{ll} 1 & \text{if customer } i \text{ is served from site } j \\ 0 & \text{otherwise} \end{array} \right.$$

$$\min \sum_{j \in N} c_j x_j + \sum_{i \in I} \sum_{j \in N} h_{ij} y_{ij}$$

$$\sum_{j \in N} y_{ij} = 1 \quad i \in I$$
(1)

$$\sum_{i\in I}b_iy_{ij} \leq u_jx_j \quad j\in N$$
 (2)

$$x_j \in \{0,1\} \ j \in N$$
 (3)
 $y_{ii} \in \{0,1\} \ i \in I, j \in N$

Fixed Charge

- n products
- $ightharpoonup k_i > 0$ Set up cost *charge*) for *j*-th product
- $ightharpoonup c_j$ unitary cost for making one product of type j
- \triangleright g_i revenue for one product unit of product j
- ▶ b_i quantity of resource i = 1, ..., m available
- $ightharpoonup a_{ij}$ usage of resource i for one unit of product j

Problem: Find the optimal production mix.



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Problem: Find the optimal production mix.



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$$y_j = \begin{cases} 1 & \text{if } x_j > 0 \\ 0 & \text{otherwise} \end{cases}$$



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$$y_j = \begin{cases} 1 & \text{if } x_j > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\max \sum_{j=1}^n (g_j-c_j)x_j-\sum_{j=1}^n k_jy_j$$

$$\sum_{j=1}^n a_{ij}x_j \leq b_i \qquad i=1,\ldots,m \qquad (1)$$

$$x_i \leq My_i, \quad j = 1, \dots, n$$
 (2)

$$x_j \le M y_j,$$
 $j = 1, ..., n$ (2)
 $y_j \in \{0, 1\},$ $j = 1, ..., n$ (3)
 $x_j \ge 0,$ $j = 1, ..., n$ (4)

$$x_i \geq 0, \qquad j = 1, \dots, n$$
 (4)