

**Laboratory**  
**8**

Expected delivery of **lab\_08.zip** must include:

- zipped project folders for Exercise1, Exercise2
- this lab track completed and converted to pdf format.

## Exercise 1) Experiment the SVC instruction.

- Download the **template project** for Keil  $\mu$ Vision “**01\_SVC**” from the course material.

Write, compile, and execute a code that invokes an SVC instruction in the main (in main.c) as following:

```
call_svc();
```

The `call_svc` is a [naked function \(LINK\)](#) written in assembly (i.e., no save and restore of registers) declared into a code, read only AREA called `svc_code`. It calls the `svc` with a fixed SVC number.

Q1: Where is the code of the function when the AREA is declared as READONLY?

*When the AREA is declared as READONLY first of all this is the size of my code:*

*Program Size: Code=468 RO-data=300 RW-data=0 ZI-data=1120*

*From the memory map (stored in Listings -> .SVC.map) it can be seen that the section is stored in the ER\_RO (readonly) region.*

Q2: Where is the code of the function when the AREA is declared as READWRITE?

*When the AREA is declared as READWRITE first of all this is the size of my code:*

*Program Size: Code=504 RO-data=276 RW-data=0 ZI-data=1120*

*From the memory map (stored in Listings -> .SVC.map) it can be seen that the section is stored in the ER\_RW (readonly) region.*

You must set the control to user mode(unprivileged) before leaving the Reset\_Handler function.

By means of invoking a SuperVisor Call, implement an error-correction method to enable a fault-tolerant data transmission by leveraging **Repetition Coding**. The same bit is transmitted multiple times to increase reliability in data transmission, particularly in noisy communication channels. Instead of sending each data bit once, the bit is repeated several times, allowing the receiver to determine the original bit based on majority voting among the received bits.

For example, suppose that we send the message «10101» with a **repetition code (3,1)**. It will be encoded in the following sequence of data «111000111000111», where each transmitted bit is repeated three times.

Correct Transmission	Received Message
111000111000111	10101
Faulty Transmission	Received Message
111010111000101000	10101

It can correct a message when one of the transmitted bits is flipped - a single bit error - meaning that the **correcting ability of this code is 1 bit**.

You are required to implement the following ASM code.

- Use a SuperVisor Call (SVC) to request encoding of a binary message.

- The message to be encoded is stored in the first 8 bits of the SVC instruction (bit 0 through 7).
- In SVC handler, implement the **repetition code (3,1)** on the input message.
- For each bit in the 8-bit message:
  - Repeat each bit **three times** to create an encoded output.
  - This repetition ensures each bit is represented as three identical bits (e.g., "1" becomes "111" and "0" becomes "000").
- Return the encoded message using the stack.

Example:

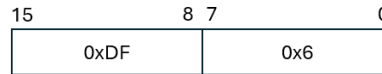


Figure 1: SVC format

SVC 0x5; 2\_0101 binary value of the SVC number

Returning from the SVC, the process ToS (Top of the stack) must contain the value "000111000111".

Q3: Describe how the stack structure is used by your project. Where is the return value? Please provide the addresses for the stack structures.

*At first, when the svc\_call assembly function is entered, the stack pointer is at address 0x10000560. Here I chose to push R3 in the stack, in order to have a place to store the return address from SVC\_Handler at the stack top. After this push, SP=0x1000055C.*

*When the SVC\_Handler is entered, following the AAPCS, registers r0-r3, r12, LR, PC are stored in the stack. Before doing so, a word of all 0s is inserted, because in ARM Cortex-M3, if bit 9 of the Configuration Control Register (located at address 0xE000ED14) is set at 1 (default value), then 8-byte stack alignment during exception entry is enforced.*

*After that, we have as said before registers r0-r3, r12, LR, PC, bringing SP to 0x10000538.*

*Then, with the instruction "STMFD SP!, {R0-R12, LR}", also these registers are stored in the stack, which address becomes 0x10000500.*

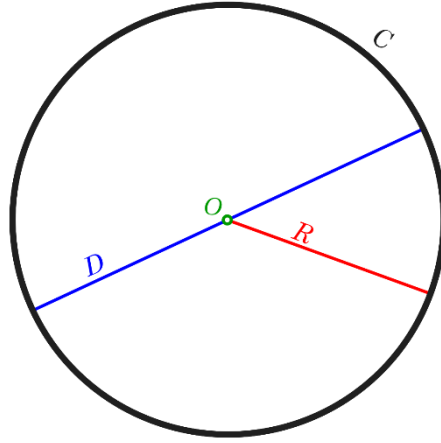
*At the end of the repetition code implementation, when we have the final value, I decided to manually address the stack address where I pushed R3 at the beginning of my program to store the return value of my function, using "STREQ R3, [R1, #0x5C]". In this way the SP is not modified, and the following instruction "LDMFD SP!, {R0-R12, LR}" is allowed to be executed, since I maintain a symmetric number of pushes and pops. After this last instruction, SP goes back again to 0x10000538, and exiting SVC\_Handler automatically registers r0-r3, r12, LR, PC are restored, thus making SP=0x1000055C.*

*In this way, now, going back to the main.c file, we have SP=0x1000055C, and at this address is stored the SVC\_Handler return value.*

## Exercise 2) – Compute the value of $\pi$ using the circle method.

- Download the template project for Keil  $\mu$ Vision "ABP" from the course material.

Pythagoreans called a set of points equally spaced from a given origin *Monad* (from Ancient Greek μονάς (*monas*) 'unity', and μόνος (*monos*) 'alone'). As originally conceived by the Pythagoreans, the *Monad* is the Supreme Being, divinity, the totality of all things, or the unreachable perfection by



human beings (but we can at least try 😊). The *Monad* (aka circle) in geometry is strongly intertwined with  $\pi$  (a mathematical transcendental irrational constant), commonly defined as the ratio between a circle's circumference and diameter.

$$\pi = \frac{C}{D}$$

An irrational number cannot be expressed exactly as a ratio of two integers. Consequently, its decimal representation never ends!

**3.14159265358979323846264338327950288419716939937510...**

However, **mathematical operations in computers are finite!** In the literature, some interesting methods and algorithms to compute an approximated value of  $\pi$  have been developed.

One of the most intuitive methods is the circle method (not the best in terms of performance). It is based on the following observations.

- The area of the circle is:

$$A = \pi r^2 \text{ (Eq. 1)}$$

- Where  $\pi$  can be computed as:

$$\pi = \frac{A}{r^2} \text{ (Eq. 2)}$$

- The assumption is to have a circle of radius **r** **centered in the origin (0,0)**.

Therefore, the Pythagoras theorem states that the distance from the origin is:

$$d^2 = x^2 + y^2 \text{ (Eq. 3)}$$

The cartesian plane can be built thinking of unitary squares centered in every  $(x, y)$  point, where  $x$  and  $y$  are the integers between  $-r$  and  $r$ .

Squares whose center belongs within or on the circumference contribute to the final area. They must satisfy the following:

$$x^2 + y^2 \leq r^2 \text{ (Eq. 4)}$$

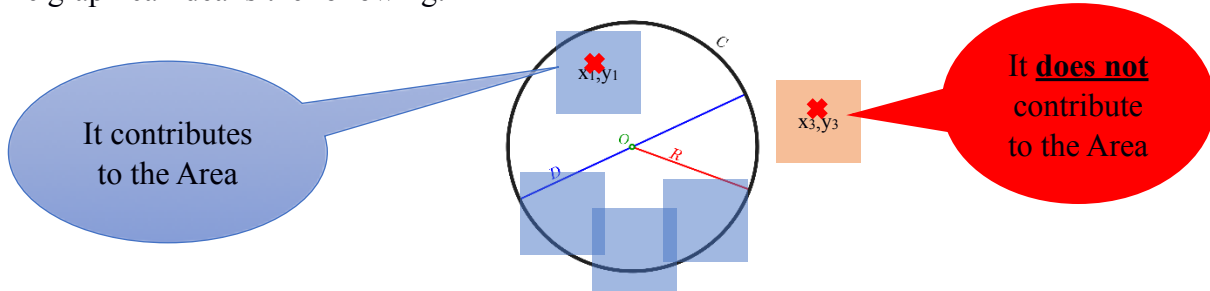
The number of points that satisfy the above condition (Eq. 4) approximates the area of the circle.

Therefore, the final formula is (where the double sums are the Area):

$$\pi \approx \frac{1}{r^2} \sum_{x=-r}^r \sum_{y=-r}^r \text{check\_square}(x, y, r) \text{ (Eq. 5)}$$

$$\text{where } \text{check\_square}(x,y,r) = \begin{cases} 1, & x^2 + y^2 \leq r^2 \\ 0, & x^2 + y^2 > r^2 \end{cases}$$

The graphical idea is the following:



Declare the coordinates as couples of  $(x, y)$  (signed integer) into a read-only memory region 2-byte aligned (into an assembly file) as follows:

```
Matrix_Coordinates DCD -5,5,-4,5,-3,5,-2,5,-1,5,0,5,1,5,2,5,3,5,4,5,5,5
DCD -5,4,-4,4,-3,4,-2,4,-1,4,0,4,1,4,2,4,3,4,4,4,4,5,4
DCD -5,3,-4,3,-3,3,-2,3,-1,3,0,3,1,3,2,3,3,3,4,3,5,3
DCD -5,2,-4,2,-3,2,-2,2,-1,2,0,2,1,2,2,2,3,2,4,2,5,2
DCD -5,1,-4,1,-3,1,-2,1,-1,1,0,1,1,1,2,1,3,1,4,1,5,1
DCD -5,0,-4,0,-3,0,-2,0,-1,0,0,0,1,0,2,0,3,0,4,0,5,0
DCD -5,-1,-4,-1,-3,-1,-2,-1,-1,-1,0,-1,1,-1,2,-1,3,-1,4,-1,5,-1
DCD -5,-2,-4,-2,-3,-2,-2,-2,-1,-2,0,-2,1,-2,2,-2,3,-2,4,-2,5,-2
DCD -5,-3,-4,-3,-3,-3,-2,-3,-1,-3,0,-3,1,-3,2,-3,3,-3,4,-3,5,-3
DCD -5,-4,-4,-4,-3,-4,-2,-4,-1,-4,0,-4,1,-4,2,-4,3,-4,4,-4,5,-4
DCD -5,-5,-4,-5,-3,-5,-2,-5,-1,-5,0,-5,1,-5,2,-5,3,-5,4,-5,5,-5
ROWS DCB 11
COLUMNS DCB 22
```

The parsing of Matrix\_Coordinates must be done in C. Remember that the `extern` keyword must be used for referencing assembly data structures:

```
extern <datatype> _Matrix_Coordinates;
extern <datatype> _ROWS;
extern <datatype> _COLUMNS;
```

Also remember that any assembly data structure that you want to use in C must be `EXPORTED`!

In the loop body of Eq. 4, `check_square(x, y, r)` is called using an assembly function with the following prototype:

```
int check_square(int x, int y, int r)
```

which implements:

$$\text{check\_square}(x,y,r) = \begin{cases} 1, & x^2 + y^2 \leq r^2 \\ 0, & x^2 + y^2 > r^2 \end{cases}$$

1) Moreover, the Arm Cortex-M3 does not provide hardware floating point support. Therefore, we can resort to software emulated floating-point algorithms, including the type `float`. As an example, Arm FPlib has the following EABI-compliant function for float division:

```
float __aeabi_fdiv (float a ,float b) /* return a/b */
```

<https://developer.arm.com/documentation/dui0475/m/floating-point-support/the-software-floating-point-library--fplib/calling-fplib-routines?lang=en>

You are required to compute the division with  $r^2$  in (Eq. 5) by calling a second assembly function with the following prototype:

```
float my_division(float* a, float* b)
```

The function body to implement is:

```
my_division:
    /*save R4,R5,R6,R7,LR,PC*/
    /*obtain value of a and b and prepare for next function
    call*/
    /*call __eabi_fdiv*/
    /*results has to be returned!*/
    /*restore R4,R5,R6,R7,LR,PC*/
```

2) Compute the value of  $\pi$  using a radius of 2,3,5 and store it into a variable.

Radius (r)	Area	Approximated value of $\pi$ (3 decimal units)	Clock Cycle (xtal=18 MHz)
2	13	3.250	8932
3	29	3.222	8932
5	81	3.240	8932

Converter from hex to FP (and viceversa): <https://gregstoll.com/~gregstoll/floattohex/>