Lecture 3: Indexing for Sequence Alignment + Estimating Sequencing Coverage



ECE 365 - Data Science and Genomics

Announcements

□ Lab 2 (sequence alignment) released tomorrow

From last lecture: Best global alignment

We use an algorithm based on dynamic programming

match: +1mismatch: -1gap: -1

		G	С	A	Т	Т	С
	0	-1	-2	-3	-4	-5	-6
G	-1	1	0	-1	-2	-3	-4
A	-2	0	0	1	0	-1	-2
Т	-3	-1	-1	0	2	1	0
Т	-4	-2	-2	-1	1	3	2
A	-5	-3	-3	-1	0	2	2
С	-6	-4	-2	-2	-1	1	3

Needleman-Wunsch algorithm

$$H_{ij} = \max \begin{cases} H_{i-1,j-1} + m, \\ (\text{if } x_i = y_j) \\ H_{i-1,j-1} - s, \\ (\text{if } x_i \neq y_j) \\ H_{i-1,j} - d \\ H_{i,j-1} - d \end{cases}$$

GCATT-C G-ATTAC

From last lecture: Best local alignment

□ We can adapt the previous algorithm to find local alignments

match: +3mismatch: -3gap: -2

		Т	G	Т	Т	A	С	G	G
	0	0	0	0	0	0	0	0	0
G	0	0	3	1	0	0	0	3	3
G	0	0	3	1	0	0	0	3	6
Т	0	3	1	6	4	2	0	1	4
Т	0	3	1	4	9	7	5	3	2
G	0	1	6	4	7	6	4	8	6
A	0	0	4	3	5	10	8	6	5
С	0	0	2	1	3	8	13	11	9
Т	0	3	1	5	4	6	11	10	8
A	0	1	0	3	2	7	9	8	7

Smith-Waterman algorithm

$$H_{ij} = \max \begin{cases} H_{i-1,j-1} + m, \\ (\text{if } x_i = y_j) \\ H_{i-1,j-1} - s, \\ (\text{if } x_i \neq y_j) \\ H_{i-1,j} - d \\ H_{i,j-1} - d \\ 0 \end{cases}$$

GTT-AC
GTTGAC

From last lecture: Alignment problem variations

Global alignment

GCATT-C G-ATTAC

		G	С	A	Т	Т	С
	0	-1	-2	-3	-4	-5	-6
G	-1	1	0	-1	-2	-3	-4
Α	-2	0	0	1	0	-1	-2
Т	-3	-1	-1	0	2	1	0
Т	-4	-2	-2	-1	1	3	2
Α	-5	-3	-3	-1	0	2	2
U	-6	-4	-2	-2	-1	1	3

Local alignment

T GTT-AC GG G GTTGAC TA

		т	G	Т	Т	A	С	G	G
	0	0	0	0	0	0	0	0	0
G	0	0	3	1	0	0	0	3	3
G	0	0	3	1	0	0	0	3	6
Т	0	3	1	6	4	2	0	1	4
Т	0	3	1	4	9	7	5	3	2
G	0	1	6	4	7	6	4	8	6
А	0	0	4	3	5	10	8	6	5
С	0	0	2	1	3	8	13	11	9
Т	0	3	1	5	4	6	11	10	8
А	0	1	0	3	2	7	9	8	7

Alignment to reference

GCACT-C
TCT G-ACTAC GCGT

	Т	С	Т	G	Α	С	Т	Α	С	G	С	G	Т
G													
С													
А													
С													
Т													
С													

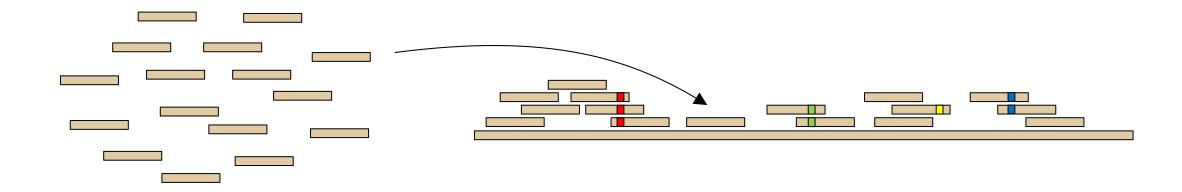
Overlap alignment

GCACT-C GCGTTCA
TTCT G-ACTAC

_	 			_	_	_	_	_	_	_	_		
	G	С	Α	С	Т	С	G	U	G	Т	Т	С	Α
Т													
Т													
С													
Т													
G													
Α													
С													
Т													
Α													
С													

Are these methods computationally efficient?

- Needleman-Wunsch and Smith-Waterman algorithms are slow in practice
- Impractical to align many short reads to a reference genome:



Is there a way to quickly figure out where a read belongs in the genome?

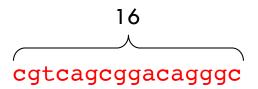
Does the sequence cgtcagcggacagggc appear in the genome below?

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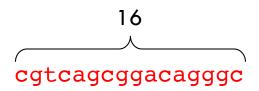
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Does the sequence cgtcagcggacagggc appear in the genome below?

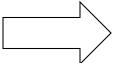
One idea: sort substrings like in a dictionary!



ggtttaatgtggttctgcttggcggtagtcattaagagccccgtggtggccaat caagaaaatgtcacgccgcttcccagcactttcagctgttttgtcgtagcccat caccaccgtaagccaagacccagcttcaggccaagtagccttccgccagcggtt ctgcgtcggcatggattctgcacggcaaagttcacgcgtcggtttgccataatt aaggacgcgcctggattcaccttgcgatcggcaatcgcaggaatgagagagcag ataatgaaagcgttgacgtaagaaagccatcgttttcccggtaccggttttgc gcctgcccggctacgtcagcgacctcgccagcgtcagcggacagggcgcaagtg ccgtgaatgggccgtacagttatgaaaccctttttttctaaggggcttctacaa cccttggatgcagggcgaagtcgggaaaacttctgttctgtttaaaatgtgttt tgctcatagtgtggtagatctcagcttactattggctttaacgaaagccgtatt ccggtgaaaataacagtcacgcttttagttgttaatgttacaccaacaacgaaa ccaacacgccaggcttaattcctgtggggttaaatcggatcc



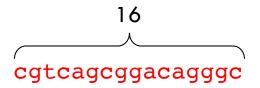
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all substrings of length 16

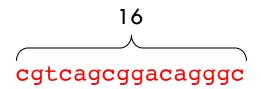
aaaacttctgttctgt	457
aaaataacagtcacgc	546
aaaatgtcacgccgct	58
aaaatgtgttttgctc	475
aaaccaacacgccagg	591
aaaccctttttttcta	402
aaacttctgttctgtt	458
aaagccatcgttttcc	292
aaagccgtattccggt	529
aaagcgttgacgtaag	276
aaagttcacgcgtcgg	188
aaataacagtcacgct	547
aaatgtcacgccgctt	59
aaatgtgttttgctca	476
aacaacgaaaccaaca	584
aacacgccaggcttaa	596
aacagtcacgctttta	551
aaccaacacgccaggc	592
aacccttggatgcagg	430
aaccctttttttctaa	403
aacgaaaccaacacgc	587
aacgaaagccgtattc	525
aacttctgttctgttt	459
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aagaaagccatcgttt	289
aagacccagcttcagg	122
aagagccccgtggtgg	33

16



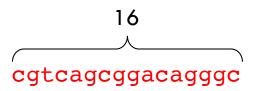
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Another way to do indexing: Python dictionary

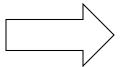


ggtttaatgtggttctgcttggcggtagtcattaagagccccgtggtggccaat caagaaatgtcacgccgttcccagcactttcagctgttttgtcgtagcccat caccaccgtaagccaagacccagcttcaggccaagtagccttccgccagcggtt ctgcgtcggcatggattctgcacggcaaagttcacgcgtcggtttgccataatt aaggacgcgcctggattcaccttgcgatcggcaatcgcaggaatgagagagcag ataatgaaagcgttgacgtaagaaagccatcgttttcccggtaccggttttgc gcctgcccggctacgtcagcgacctcgccagcgtcagcggacagggcgcaagtg ccgtgaatgggccgtacagttatgaaaccctttttttctaaggggcttctacaa cccttggatgcagggcgaagtcgggaaaacttctgttctgtttaaaatgtgttt tgctcatagtgtggtagatctcagcttactattggctttaacgaaagccgtatt ccggtgaaaataacagtcacgcttttagttgttaatgttacaccaacaacgaaa ccaacacgccaggcttaattcctgtggagttaattctggcttaatatgggctaaatcggatcc

Another way to do indexing: Python dictionary (hash function)



ggtttaatgtggttctgcttggcggtagtcattaagagccccgtggtggccaat caagaaaatgtcacgccgcttcccagcactttcagctgttttgtcgtagcccat caccaccgtaagccaagacccagcttcaggccaagtagccttccgccagcggtt ctgcgtcggcatggattctgcacggcaaagttcacgcgtcggtttgccataatt aaggacgcgcctggattcaccttgcgatcggcaatcgcaggaatgagagagcag ataatgaaagcgttgacgtaagaaagccatcgttttcccggtaccggttttgc gcctgcccggctacgtcagcgacctcgccagcgtcagcggacagggcgcaagtg ccgtgaatgggccgtacagttatgaaaccctttttttctaaggggcttctacaa cccttggatgcagggcgaagtcgggaaaacttctgttctgtttaaaatgtgttt tgctcatagtgtggtagatctcagcttactattggctttaacgaaagccgtatt ccggtgaaaataacagtcacgcttttagttgttaatgttacaccaacaacgaaa ccaacacgccaggcttaattcctgtggagttatatatagagcgtaaatcggatcc



all substrings of length 16

Dict()

```
ggtttaatgtggttct
gtttaatgtggttctg
tttaatgtggttctgc
ttaatgtggttctgct
taatgtggttctgctt
aatgtggttctgcttg
atgtggttctgcttgg
tgtggttctgcttggc
gtggttctgcttggcg
tggttctgcttggcgg
ggttctgcttggcggt
gttctgcttggcggta
ttctgcttggcggtag
tctgcttggcggtagt
ctgcttggcggtagtc
tgcttggcggtagtca
gcttggcggtagtcat
cttggcggtagtcatt
ttggcggtagtcatta
```

- Another way to do indexing: Python dictionary (hash function)
- Let's look at this in a notebook



Y =

$$X = cgtaagcggacatggc$$

 $lue{}$ One idea: Consider indexing substrings of length $k \in \mathcal{C}$





all substrings of length k = 6

 $lue{}$ One idea: Consider indexing substrings of length k

```
Python Dict ()
ggttta --> [0]
gtttaa --> [1, 471]
tttaat --> [2]
ttaatg --> [3, 571]
taatgt --> [4, 572]
aatqtq --> [5, 477]
atgtgg --> [6]
tgtggt --> [7, 495]
qtqqtt --> [8]
tggttc --> [9]
ggttct --> [10, 158]
gttctg --> [11, 159, 466]
ttctqc --> [12, 160, 177]
tctqct --> [13]
ctgctt --> [14]
tqcttq --> [15]
gcttgg --> [16]
cttagc --> [17]
ttqqcq --> [18]
tagcag --> [19]
ggcggt --> [20]
```

Y =

```
X = cgtaagcggacatggc
```

take substrings of length k = 6

```
X = cgtaagcggacatggc
```

take substrings of length k = 6

```
cgtaag --> (114) 286]
gtaagc --> (115)
taagcg --> aagcgg --> agcgga --> (359)
gcggac --> (360)
cggaca --> gacatg --> acatgg --> catggc -->
```

```
gtaagc -->
                                          taagcg -->
                                          aagcgg -->
X = cgtaagcggacatggc
                                          agcgga -->
                                          gcggac --> (3601
                                          cggaca --> 3611
     take substrings
                                          ggacat -->
     of length k = 6
                                          gacatg -->
                                                                  Smith Washiman (X, Y)
                                          acatgg -->
                                          catggc -->
                      cgtaagcggacatggc
  ...tagcccatcaccaccgtaagccaagacccagcttcaggccaagtagccttccgccagcgg...
                     114
  ...ccgctacgtcagcgacctcgccagcgtcagcggacagggcgcaagtgccgtgaatgggc...
```

How do we choose k?

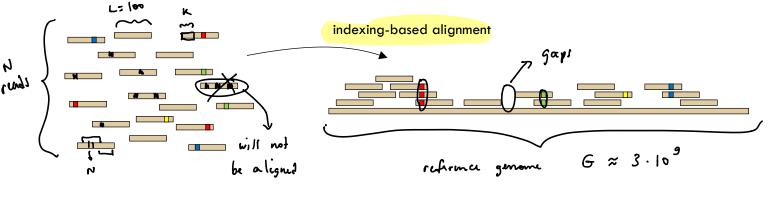
 \square Based on the sequencing technology (e.g., Illumina: L=100, error rate = 0.1%)

Claim: there is a regiment of length $\frac{100-1}{2}$ = 49.5 $\left[\frac{L-t}{t+1}\right]$ with no errors

the of errors in a read: $X \sim Binomial(100, 0.001)$ P(X>2) = 0.00015 fraction of reads with > t errors: P(X>t)

$$p(\chi > 2) = 0.000 15$$
 $K = \left(\frac{100 - 2}{3}\right) = 33$

Big picture:



How to pick N?