

## Proof for LP strong duality:

Here, consider  $\min c^T x$ , s.t.  $Ax = b$ ,  $x \geq 0$ ,  $x \in \mathbb{R}^n$ ,  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$

i.e. we only consider the standard form LP since any kind of LP can be transformed to it.

(Assume  $C, D \subseteq \mathbb{R}^n$ ).

Lemma 1: (Separation thm): For two closed and convex set, denoted by  $C$  and  $D$ , if one is bounded, then we have:  
 $\exists d \neq 0 \in \mathbb{R}^n$  and  $r \in \mathbb{R}$ , s.t.  $d^T c < r < d^T d$  for  $\forall c \in C, d \in D$ .

if  $C \cap D = \emptyset$ .

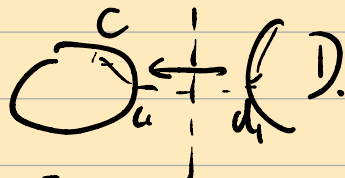
Lemma's pf: w.l.o.g. assume  $C$  is bounded.

It's easy to verify that:  $\exists c \in C, d_1 \in D$ , s.t.  $d(C, D) = \|c_1 - d_1\|_2$ .

$$\text{Here, } d(C, D) = \inf_{\substack{c \in C \\ d \in D}} \|c - d\|_2 = \inf_{c \in C} p(c, D)$$

(Can be done by choosing  $\{d_n\} \subset D$  within  $\frac{1}{n}$  respectively  
Or just leverage  $p(c, D)$  is cts and  $C$  is cpt,  $D$  is cls).

$$\text{Then, let } d = c_1 - d_1, r = \frac{c_1^T - d_1^T}{2}$$



$$\text{Since } \begin{cases} d_1 = \arg\min_D \frac{1}{2} \|c_1 - d\|^2 \\ c_1 = \arg\min_C \frac{1}{2} \|c - d_1\|^2 \end{cases}, \text{ we get } \begin{cases} (c_1 - d_1)^T (d - d_1) \geq 0, \text{ any } d \in D \\ (c_1 - d_1)^T (c - c_1) \geq 0, \text{ any } c \in C \end{cases}$$

$$\text{Hence, } \begin{cases} (c_1 - d_1)^T d_1 \geq (c_1 - d_1)^T d, \text{ any } d \in D \\ (c_1 - d_1)^T c \geq (c_1 - d_1)^T c_1, \text{ any } c \in C. \end{cases}$$

$$\text{Hence, } (c_1 - d_1)^T c \geq (c_1 - d_1)^T c_1 > (c_1 - d_1)^T d_1 \geq (c_1 - d_1)^T d, \text{ any } c, d \in C, D.$$

$$\text{It's easily to verify that: } (c_1 - d_1)^T c_1 > (c_1 - d_1)^T \frac{c_1 + d_1}{2} > (c_1 - d_1)^T d_1.$$

// In  $\mathbb{R}^n$ .

(~~||~~ implies without bdd, even if  $d(C, D) > 0$ , we can't get the strict result).

• Corollary: For any closed convex set  $C$ , any point  $d \notin C$ , we have:  
 $\exists e \neq 0$ , s.t.  $e^T d > e^T c$  for any  $c \in C$ .

• Corollary: For any cvx set  $C \neq \emptyset \subseteq \mathbb{R}^n$ ,  $z \in \text{bdry}(C)$ ,  $\exists a \neq 0 \in \mathbb{R}^n$ ,  
s.t.  $a^T z \geq a^T x$ ,  $\forall x \in C$ .

(Just using  $\text{cl}(C)$  is closed cvx set. Since  $z \in \text{bdry}(C)$ ,  $\exists z_i \in (\text{cl}(C))^c$   
with  $\|a_i\| = 1$ , s.t.  $a_i^T z_i > a_i^T x$ ,  $\forall x \in \text{cl}(C)$  and  $z_i \rightarrow z$ .

$\exists \{a_{i_k}\}_{k=1}^\infty \subset \{a_i\}_{i=1}^\infty$  converging to  $a$ , then we get:

$$a^T z \geq a^T x, \forall x \in C. \quad )$$

• Corollary: For any two cvx set  $C, D$ , if  $C \cap D = \emptyset$ , then  $\exists a \neq 0$ , s.t.  
 $a^T c \leq a^T d$ ,  $\forall c \in C, \forall d \in D$ .

(Considering  $C-D$ ,  $0 \notin C-D$ . consider  $\begin{cases} 0 \in \text{cl}(C-D) \\ 0 \notin \text{cl}(C-D) \end{cases}$  )

Now, we can propose:

Lemma 2: (Farkas Strongy Alternative).

$\begin{cases} \textcircled{1} \{x \mid Ax = b, x \geq 0\} \\ \textcircled{2} \{y \mid A^T y \leq 0, b^T y > 0\} \end{cases}$ , exactly one of  $\textcircled{1}, \textcircled{2}$  is empty

Proof: If  $\textcircled{1}$  is nonempty,  $\textcircled{2}$  must be empty.

If  $\textcircled{1}$  is empty, Let  $A = (a_1, \dots, a_n)$ , then  $b \notin \text{cone}(a_1, \dots, a_n)$  which is closed. then,  $\exists y \neq 0$ , s.t.  $y^T b > y^T c$ ,  $c \in \text{cone}(a_1, \dots, a_n)$ .

Since  $0 \in \text{cone}(a_1, \dots, a_n)$ ,  $y^T b > 0$ .

And  $y^T a_i \leq 0, i=1, \dots, n$ , (Otherwise  $y^T c \rightarrow \infty$  for some  $c \in \text{cone}(a_1, \dots, a_n)$ )  
 $\Rightarrow y^T A \leq 0, \dots$

Lemma 3:

- ①  $\exists x, Ax \leq b$   
②  $\exists y \geq 0, \text{ s.t. } y^T A = 0, y^T b < 0.$  } exactly one of them holds

Proof: If ① holds, then ② can't hold.

If ① doesn't hold. Rewrite it in standard form:

$$\{(x^+, x^-, s) \mid (A, -A, I) \begin{pmatrix} x^+ \\ x^- \\ s \end{pmatrix} = b, x^+, x^-, s \geq 0\} \text{ is empty.}$$

$$\text{Hence, } \exists y \neq 0 \in \mathbb{R}^n, \text{ s.t. } \begin{pmatrix} A^T \\ -A^T \\ I \end{pmatrix} y \geq 0, b^T y < 0 \Rightarrow A^T y = 0.$$

Now, we begin our prove

$$(P) \min c^T x \\ \text{s.t. } Ax = b, x \geq 0$$

$$(D) \max b^T \lambda \\ \text{s.t. } c \geq A^T \lambda$$

Proof: If (P) isn't bdd, i.e.  $p^* = -\infty$ , then  $d^* = -\infty$  as well.

If  $p^*$  is finite, we just need to prove  $\begin{cases} c \geq A^T \lambda \\ b^T \lambda \geq p^* \end{cases}$  is feasible, i.e.

$$\begin{pmatrix} A^T \\ -b^T \end{pmatrix} \lambda \leq \begin{pmatrix} c \\ -p^* \end{pmatrix} \text{ is feasible.}$$

If not.  $\exists (y_1, y_2) \geq 0, \text{ s.t. } Ay_1 - by_2 = 0, y_1^T c - y_2^T p^* < 0$

① If  $y_2 = 0$ , then.  $Ay_1 = 0$ ,  $y_1^T c < 0$ ,  $y_1 \geq 0 \Rightarrow$  Contradiction.

② If  $y_2 \neq 0$ , then  $A\left(\frac{y_1}{y_2}\right) = b$ ,  $c^T \frac{y_1}{y_2} < p^* \Rightarrow$  Contradiction. □

