

# Recursive Models of Theories with Few Models

Bakhadyr Khoussainov\*

Mathematics Department, Cornell University, Ithaca NY 14850 USA

Andre Nies†

Mathematics Department, The University of Chicago, IL 60637 USA

Richard A. Shore §

Mathematics Department, Cornell University, Ithaca NY 14853 USA

## 1 Introduction

We begin by presenting some basic definitions from effective model theory. A **recursive structure** is one with a recursive domain and uniformly recursive atomic relations. Without loss of generality, we can always suppose that the domain of every recursive structure is the set of all natural numbers  $\omega$  and that its language does not contain function symbols. If a structure  $\mathcal{A}$  is isomorphic to a recursive structure  $\mathcal{B}$ , then  $\mathcal{A}$  is **recursively presentable** and  $\mathcal{B}$  is a **recursive presentation of  $\mathcal{A}$** . Let  $\sigma$  be an effective signature. Let  $\sigma_0 \subset \sigma_1 \subset \sigma_2 \subset \dots$  be an effective sequence of finite signatures such that  $\sigma = \bigcup_t \sigma_t$ . It is clear that a structure  $\mathcal{A}$  of signature  $\sigma$  is recursive if and only if there exists an effective sequence  $\mathcal{A}_0 \subset \mathcal{A}_1 \subset \mathcal{A}_2 \subset \dots$  of finite structures such that for each  $i$  the domain of  $\mathcal{A}_i$  is  $\{0, \dots, t_i\}$ , the function  $i \rightarrow t_i$  is recursive,  $\mathcal{A}_i$  is a structure of signature  $\sigma_i$ ,  $\mathcal{A}_{i+1}$  is an expansion and

---

\*Partially supported by ARO through MSI, Cornell University, DAAL03-91-C0027.

†Partially supported by MSI, Cornell University and the NSF grant DMS 9500983

§Partially supported by NSF Grant DMS-9204308, DMS-9503503 and ARO through MSI, Cornell University, DAAL-03-C-0027.

extension of  $\mathcal{A}_i$ , and the structure  $\mathcal{A}$  is the union  $\bigcup_i \mathcal{A}_i$ . The domain of  $\mathcal{A}$  is denoted by  $A$ . For a structure  $\mathcal{A}$  of signature  $\sigma$  we write  $P^{\mathcal{A}}$  to denote the interpretation of the predicate symbol  $P \in \sigma$  in  $\mathcal{A}$ . When it does not cause confusion, we write  $P$  instead of  $P^{\mathcal{A}}$ . In this paper we only deal with finite or countable structures.

A basic question in recursive model theory is whether a given first order theory  $T$  has a recursive model. A standard Henkin type construction shows that each decidable theory has a recursive model. Moreover the satisfaction predicate for this model is recursive. Such recursive models are called **decidable**. Constructing recursive (decidable) presentations for specific models of  $T$  has been an intensive area of research in effective model theory [2], [9], [4]. For example, the recursiveness of homogeneous models, in particular of prime and saturated models has been well studied. In [2], [9] it is proved that the saturated model of  $T$  has a decidable presentation if and only if there exists a procedure which uniformly computes the set of all types of  $T$ . Goncharov [4] and Harrington [8] gave criteria for prime models to have decidable presentations. It is also known that the decidability of the saturated model of  $T$  implies the existence of a decidable presentation of the prime model of  $T$  [2], [12]. Thus, a general question arises as to how recursive models of undecidable theories behave in comparison to recursive models of decidable theories. In this paper we investigate recursive models of complete theories with “few countable models” (M. Morley [12]). Examples of such theories are theories with countably many countable models such as  $\omega_1$ -categorical theories and theories with finitely many countable models (**Ehrenfeucht theories**).

In [1] Baldwin and Lachlan developed the theory of  $\omega_1$ -categoricity in terms of strongly minimal sets. They settled affirmatively Vaught’s conjecture for  $\omega_1$ -categorical complete theories by proving that each complete  $\omega_1$ -categorical theory has either exactly one or  $\omega$  many countable models up to isomorphisms. Their paper also shows that all the countable models of any  $\omega_1$ -categorical theory  $T$  can be listed in an  $\omega + 1$  chain:

$$chain(T) : \quad \mathcal{A}_0 \preceq \mathcal{A}_1 \preceq \dots \preceq \mathcal{A}_n \preceq \dots \mathcal{A}_\omega$$

of elementary embeddings with  $\mathcal{A}_0$  and  $\mathcal{A}_\omega$  being the prime and saturated models of  $T$ , respectively [1]. The results of Baldwin and Lachlan lead one to investiagte the effective content of  $\omega_1$ -categorical theories and their models. Based on the theory developed by Baldwin and Lachlan, Harrington

[8] and Khissamiev [6] proved that every countable model of each decidable  $\omega_1$ -categorical theory  $T$  has a decidable presentation.

This result of Harrington and Khissamiev motivated the study of recursive models of  $\omega_1$ -categorical undecidable theories. In 1972, S. Goncharov [3] constructed an example of an  $\omega_1$ -categorical but not  $\omega$ -categorical theory  $T$  for which the only model with a recursive presentation is the prime model, that is the first element of  $\text{chain}(T)$ . Later in 1980, K.Kudeiberganov [7] modified Goncharov's construction to provide an example of an  $\omega_1$ -categorical but not  $\omega$ -categorical theory  $T$  with exactly  $n$  **recursive** models. These models are the first  $n$  elements of  $\text{chain}(T)$ . These results lead to the following two questions which have remained open:

**Question 1.1** (*S.Goncharov [5]*) *If an  $\omega_1$ -categorical but not  $\omega$ -categorical theory  $T$  has a recursive model, is the prime model of  $T$  recursively presentable?*

**Question 1.2** *If all models  $\mathcal{A}_0, \mathcal{A}_1, \dots, \mathcal{A}_i, \dots, i \in \omega$ , in  $\text{chain}(T)$  of an  $\omega_1$ -categorical but not  $\omega$ -categorical theory  $T$ , have recursive presentations, is the saturated model  $\mathcal{A}_\omega$  of  $T$  recursively presentable?*

The above result of Harrington and Khissamiev also inspired Nerode to ask whether the hypothesis of  $\omega_1$ -categoricity of  $T$  can be replaced by the hypothesis that  $T$  has only finitely many countable models, that is whether every countable model of a decidable Ehrenfeucht theory has a decidable presentation. Morley noted that if the countable saturated model of a such theory is decidable, then the theory has at least three recursive models [12]. Lachlan answered Nerode's question by giving an example of a decidable theory with exactly 6 models of which only the prime one has a recursive presentation. Later, for each natural number  $n > 3$ , Peretyatkin constructed an example of decidable theory with exactly  $n$  models such that the prime model of the theory is recursive and none of the other models of the theory has recursive presentations [13]. In [7] Kudeiberganov constructed an example of a theory with exactly 3 models such that the theory has only one recursive model and that model is prime. The saturated model of the theory can not be decidable since, otherwise, all 3 model of the theory would have recursive presentations. These results lead Morley to ask as whether any countable model of a decidable Ehrenfeucht theory  $T$  with a decidable saturated model

has a decidable presentation [12]. There is a natural analog of this question for recursive models:

**Question 1.3** *If the saturated model of an Ehrenfeucht theory is recursive, does there exist a nonsaturated, recursive model of the theory?*

In this paper we answer the above three questions by providing appropriate counterexamples. Our examples of models which answer the first two questions have infinite signatures. However these questions remain open for theories of finite signatures.

The general problem suggested by these results is to characterize **the spectrum of recursive models of  $\omega_1$ -categorical theories**: Let  $T$  be an  $\omega_1$ -categorical but not  $\omega$ -categorical complete theory. Consider  $\text{chain}(T)$ . **The spectrum of recursive models of  $T$** , denoted by  $\mathbf{SRM}(T)$ , is the set

$$\{i \leq \omega \mid \text{the model } \mathcal{A}_i \text{ in } \text{chain}(T) \text{ has a recursive presentation}\}.$$

**Problem.** *Describe all subsets of  $\omega$  which are of the form  $\mathbf{SRM}(T)$  for some  $\omega_1$ -categorical theory  $T$ .*

The result of Harrington and Khissamiev shows that if  $T$  is decidable, then  $\mathbf{SRM}(T) = \omega \cup \{\omega\}$ . The results of Goncharov and Kudeiberganov show that the sets  $\{1, \dots, n\}$ , where  $n \in \omega$ , are spectra of recursive models of  $\omega_1$ -categorical theories. In this paper we show that the sets  $\omega - \{0\} \cup \{\omega\}$  and  $\omega$  are also spectra of recursive models of  $\omega_1$ -categorical theories.

## 2 Main Results

The results of this paper are based on the idea of coding  $\Sigma_2^0$  or  $\Pi_2^0$  sets with certain recursion-theoretic properties into  $\omega_1$ -categorical theories. Our first result is the following theorem which answers Question 1.1.

**Theorem 2.1** *There exists an  $\omega_1$ -categorical but not  $\omega$  categorical theory  $T$  such that all the countable models of  $T$  except its prime model have recursive presentations (and so  $\mathbf{SRM}(T) = \omega - \{0\} \cup \{\omega\}$ ).*

Before proving this theorem we would like to give the basic idea of our proof. For an infinite subset  $S \subset \omega$  we construct a structure  $\mathcal{A}_S$  of infinite signature  $(P_0, P_1, P_2, \dots)$ , where each  $P_i$  is a binary predicate symbol. We will show that the theory  $T_S$  of the structure  $\mathcal{A}_S$  is  $\omega_1$ -categorical and  $\mathcal{A}_S$  is the prime model of  $T_S$ . The countable models of  $T_S$  will have the following property: Every non prime model  $\mathcal{A}$  of  $T_S$  has a recursive presentation if and only if the set  $S$  is a  $\Sigma_2^0$ -set. The existence of a recursive presentation of the prime model will imply that the set  $S$  has a certain recursion-theoretic property. Our recursion-theoretic lemma (Lemma 2.1.) will show that there exists a  $\Sigma_2^0$ -set  $S$  which does not have this properties.

**The Construction of Cubes.** Let  $n$  be a nonzero natural number. Let  $\sigma_n = (P_0, \dots, P_{n-1})$  be a signature such that each  $P_i$  is a binary predicate symbol. For each nonzero natural number  $n$  we define a finite structure of signature  $\sigma_n$ , called an  $n$ -cube, as follows.

A 1-cube  $\mathcal{C}_1$  is a structure  $(\{a, b\}, P_0)$  such that  $P_0(x, y)$  holds in  $\mathcal{C}_1$  if and only if  $x = a$  and  $y = b$  or  $y = a$  and  $x = b$ .

Suppose that  $n$ -cubes have been defined. Let  $\mathcal{A} = (A, P_0^{\mathcal{A}}, \dots, P_{n-1}^{\mathcal{A}})$  and  $\mathcal{B} = (B, P_0^{\mathcal{B}}, \dots, P_{n-1}^{\mathcal{B}})$  be  $n$ -cubes such that  $A \cap B = \emptyset$ . These two  $n$ -cubes are isomorphic. Let  $f$  be an isomorphism from  $\mathcal{A}$  to  $\mathcal{B}$ . Then a  $n+1$ -cube  $\mathcal{C}_{n+1}$  is

$$(A \cup B, P_0^{\mathcal{A}} \cup P_0^{\mathcal{B}}, \dots, P_{n-1}^{\mathcal{A}} \cup P_{n-1}^{\mathcal{B}}, P_n),$$

where  $P_n(x, y)$  holds if and only if  $f(x) = y$  or  $f^{-1}(x) = y$ . It follows that we can naturally define an  $\omega$ -cube  $\mathcal{C}_\omega = \bigcup_{i \in \omega} \mathcal{C}_i$  as an increasing union of  $n$ -cubes formed in this way.

An  $\omega$ -cube  $\mathcal{C}_\omega$  is a structure of the infinite signature  $\sigma = (P_0, P_2, \dots)$ . From these definitions of cubes it follows

**Claim 2.1** *For each  $n \leq \omega$  any two  $n$ -cubes are isomorphic.*  $\square$

Each binary predicate  $P_i$  in any cube  $\mathcal{A}$  is a partial function and sets up a one-to-one mapping from  $\text{dom}(P_i)$  onto  $\text{range}(P_i)$ . Therefore we can also write  $P_i(x) = y$  instead  $P_i(x, y)$ . Moreover by the definition of  $P_i$ ,  $\text{dom}(P_i) = \text{range}(P_i)$ .

**Construction of  $\mathcal{A}_S$ .** For each natural number  $n \in \omega$  consider an  $n$ -cube denoted by  $\mathcal{A}_n$ . Assume that  $\mathcal{A}_n \cap \mathcal{A}_t = \emptyset$  for all  $n \neq t$ . Let  $S$  be a subset of  $\omega$ . Define a structure  $\mathcal{A}_S$  by

$$\mathcal{A}_S = \bigcup_{n \in S} \mathcal{A}_n.$$

Thus the structure  $\mathcal{A}_S$  is the disjoint union of all cubes  $\mathcal{A}_n$ ,  $n \in S$ , with the natural interpretations of predicate symbols of signature  $\sigma$ . Let  $T_S$  be the theory of the structure  $\mathcal{A}_S$ .

**Claim 2.2** *If  $S$  is an infinite set, then the theory  $T_S$  is  $\omega_1$  categorical but not  $\omega$ -categorical.*

**Proof.** The model  $\mathcal{A}_S$  satisfies the following list of statements. It is easy to see that this list of statements can be written as an (infinite) set of statements in the first order logic.

1.  $\forall x \exists y P_0(x, y)$  and for each  $n$ ,  $P_n$  is a partial one to one function.
2. For all  $n \neq m$  and for all  $x$ ,  $P_n(x) \neq P_m(x)$ .
3. For each  $n$  and for all  $x$  if  $P_n(x)$  is defined, then  $P_0(x)$ ,  $P_1(x)$ ,  $\dots$ ,  $P_{n-1}(x)$  are also defined.
4. For all  $n, m$  and for all  $x$  if  $P_n(x)$  and  $P_m(P_n(x))$  are defined, then  $P_m(P_n(x)) = P_n(P_m(x))$ .
5. For all  $k$ ,  $n > n_1 \geq n_2 \geq \dots \geq n_{k-1} \geq n_k$ , for all elements  $x$ ,  $P_{n_1}(\dots(P_{n_k}(x)\dots) \neq P_n(x)$ .
6. For each  $n \in \omega$ ,  $n \in S$  if and only if there exists exactly one  $n$ -cube which is not contained in an  $n+1$ -cube.

Let  $\mathcal{M}$  be a model which satisfies all the above statements. Then for each  $n \in S$ ,  $\mathcal{M}$  must have an  $n$ -cube which is not contained in an  $n+1$ -cube. Moreover if an  $x \in M$  does not belong to any  $n$ -cube for  $n \in S$ , then  $x$  is in an  $\omega$ -cube. Note that each  $\omega$ -cube is countable. Using the previous claim it can be seen that any two models which satisfy the above list of axioms are isomorphic if and only if these two models have the same number of  $\omega$ -cubes.

Suppose that  $\mathcal{M}_1$  and  $\mathcal{M}_2$  are models of  $T_S$  and their cardinalities are  $\omega_1$ . Since each cube is a countable set it follows that the number of  $\omega$ -cubes in  $\mathcal{M}_1$  and  $\mathcal{M}_2$  is  $\omega_1$ . Therefore the models  $\mathcal{M}_1$  and  $\mathcal{M}_2$  are isomorphic. Hence  $T_S$  is an  $\omega_1$ -categorical but not  $\omega$ -categorical theory.  $\square$

**Claim 2.3** *The set  $S$  is in  $\Sigma_2^0$  if and only if every nonprime model of  $T_S$  possesses a recursive presentation.*

**Proof.** Each  $\omega$ -cube has a recursive presentation. Therefore it suffices to prove that  $S \in \Sigma_2^0$  if and only if the nonprime model  $\mathcal{M}$  of  $T_S$  with exactly one  $\omega$ -cube has a recursive presentation. If  $\mathcal{M}$  is recursive, then  $s \in S$  if and only if  $\exists x \exists y \forall z (P_s(x, y) \& \neg P_{s+1}(x, z))$ . Therefore  $S \in \Sigma_2^0$ .

Now suppose that  $S \in \Sigma_2^0$ . There exists a recursive function  $f$  such that for every  $n \in \omega$ ,  $n \in S$  if and only if  $W_{f(n)}$  is finite. We construct an effective sequence

$$\mathcal{M}_0 \subset \mathcal{M}_1 \subset \mathcal{M}_2 \subset \dots$$

of finite structures by stages such that

1. The model  $\mathcal{M}$  is isomorphic to  $\bigcup_n \mathcal{M}_n$ ,
2. Each  $\mathcal{M}_t$  has exactly  $t + 1$  cubes and the function  $t \rightarrow \text{card}(\mathcal{M}_t)$  is recursive,
3. Each  $\mathcal{M}_t$  is a structure of signature  $(P_0, \dots, P_{n_i})$ , where  $i \rightarrow n_i$  is a recursive function.

**Stage 0.** Construct a 1-cube  $\mathcal{M}_0$  and mark this structure with the symbol  $\square_\omega$ .

**Stage s+1.** Suppose that  $\mathcal{M}_s$  has been constructed as the disjoint union

$$\mathcal{M}_{s,0} \bigcup \mathcal{M}_{s,1} \bigcup \dots \bigcup \mathcal{M}_{s,s} \bigcup \mathcal{M}_{s,\omega},$$

where each  $\mathcal{M}_{s,i}$ ,  $i \leq s$  is a  $i$ -cube, and  $\mathcal{M}_{s,\omega}$  is the cube marked with  $\square_\omega$  at the previous stage. Compute  $W_{f(0),s+1}, \dots, W_{f(s),s+1}, W_{f(s+1),s+1}$ . For each  $i \leq s + 1$  define  $\mathcal{M}_{i,s+1}$  and  $\mathcal{M}_{s+1,\omega}$  as follows:

1. If  $W_{f(i),s+1} = W_{f(i),s}$ , then let  $\mathcal{M}_{i,s+1} = \mathcal{M}_{i,s}$ .

2. If  $W_{f(i),s+1} \neq W_{f(i),s}$ , then construct a new  $i$ -cube and let  $\mathcal{M}_{i,s+1}$  be this new cube.
3. Extend the cube  $\mathcal{M}_{s,\omega}$  to a finite cube denoted by  $\mathcal{M}_{s+1,\omega}$  such that for each  $i \leq s$  if  $W_{f(i),s+1} \neq W_{f(i),s}$ , then  $\mathcal{M}_{s+1,\omega}$  contains  $\mathcal{M}_{s,i}$ .

Let  $\mathcal{M}_{s+1}$  be  $\mathcal{M}_{s+1,0} \cup \mathcal{M}_{s+1,1} \cup \dots \cup \mathcal{M}_{s+1,s+1} \cup \mathcal{M}_{s+1,\omega}$ . Define

$$\mathcal{M}_\omega = \bigcup_s \mathcal{M}_s.$$

By the construction, the structure  $\mathcal{M}_\omega$  is recursive. The construction of  $\mathcal{M}_\omega$  guarantees that the structure  $\mathcal{M}_\omega$  is isomorphic to the model  $\mathcal{M}$ .  $\square$

Now we need the following definition and recursion theoretic lemma. We will prove the lemma at the end of this section.

**Definition 1** A function  $f$  is **limitwise monotonic** if there exists a recursive function  $\phi(x, t)$  such that  $\phi(x, t) \leq \phi(x, t+1)$  for all  $x, t \in \omega$ ,  $\lim_t \phi(x, t)$  exists for every  $x \in \omega$  and  $f(x) = \lim_t \phi(x, t)$ .

**Lemma 2.1 (Recursion Theoretic Lemma)** There exists a  $\Delta_2^0$  set  $A$  which is not the range of any limitwise monotonic function.  $\square$

**Proof of Theorem 2.1.** We need the following

**Lemma 2.2** If the prime model  $\mathcal{A}_S$  is recursive, then the set  $S$  is the range of a limitwise monotonic function.

**Proof.** Let  $x \in \mathcal{A}_S$ . Note that each cube in  $\mathcal{A}_S$  is finite. Define  $\phi(x)$  to be an  $s$  such that  $x$  is in an  $s$ -cube and this cube is not contained in a  $s+1$ -cube. It is clear that  $\phi$  witnesses that  $S$  is the range of a limitwise monotonic function.  $\square$

By the Recursion Theoretic Lemma there exists an  $S \in \Delta_2^0$  which is not the range of any limitwise monotonic function. Consider the structure  $\mathcal{A}_S$  and its theory  $T_S$ . The claims above and Lemma 2.2 show that  $T_S$  is the required theory and so prove Theorem 1.1.  $\square$

Now we give an answer to Question 1.2. The idea of our proof is the following. We take a  $\Pi_2^0$  but not  $\Sigma_2^0$  set  $S$  and code this set into a theory  $T_S$ .

The language of  $T_S$  will contain infinitely many unary predicates  $P_0, P_1, \dots$ , and infinitely many predicates of arity  $n$  for each  $n \in \omega$ . We will prove that  $T_S$  is an  $\omega_1$ -categorical but not  $\omega$ -categorical theory. Our construction of  $T_S$  guarantees that all the countable models of  $T_S$ , except the saturated model, have recursive presentations. The existence of a recursive presentation for the saturated model will imply that the set  $S$  is a  $\Sigma_2^0$  set. This will contradict with the choice of  $S$ .

**Theorem 2.2** *There exist an  $\omega_1$ -categorical but not  $\omega$ -categorical theory  $T$  such that all the countable models of  $T$  except the saturated model, have recursive presentations.*

**Proof.** We construct a structure of the infinite signature

$$(P_0, P_1, \dots, R_{1,0}, R_{1,1}, R_{1,2}, \dots, R_{k,0}, R_{k,1}, R_{k,2}, \dots),$$

where each  $P_i$  is a unary predicate and each  $R_{k,s}$  is a predicate of arity  $k$ .

Let  $S$  be a  $(\Pi_2^0 \setminus \Sigma_2^0)$  set. There exists a recursive predicate  $H$  such that  $n \in S$  if and only if  $\forall x \exists y H(x, y, n)$  holds. Below we present a step by step construction of a recursive structure denoted by  $\mathcal{A}_S$  and prove that the theory  $T_S$  of this structure satisfies the requirements of the theorem.

**Stage 0.** Let  $\mathcal{A}_0 = (\{0\}, P_0)$ , where  $P_0(0)$  holds.

**Stage  $t+1$ .** The domain  $A_{t+1}$  of  $\mathcal{A}_{t+1}$  is  $\{0, \dots, t+1\}$ . The signature of  $\mathcal{A}_{t+1}$  is

$$\sigma_{t+1} = (P_0, \dots, P_{t+1}, R_{1,0}, \dots, R_{1,t+1}, \dots, R_{t+1,0}, \dots, R_{t+1,t+1}).$$

For each  $i \leq t+1$  let  $P_i(x)$  hold if and only if  $x \geq i$ . For  $k, s \leq t+1$ , let  $R_{k,s}(x_1, \dots, x_k)$  hold if and only if  $x_1, \dots, x_k$  are pairwise different and for the maximal number  $j \leq t+1$  such that all  $P_j(x_1), \dots, P_j(x_k)$  hold we have  $\forall n \leq s \exists m \leq j H(n, m, k)$ . We have defined the model  $\mathcal{A}_{t+1}$ .

Thus we have an effective sequence  $\mathcal{A}_0, \mathcal{A}_1, \mathcal{A}_2, \dots$  of finite structures such that each  $\mathcal{A}_{i+1}$  is an extension and expansion of  $\mathcal{A}_i$ . Therefore we can define  $\mathcal{A}_S$  by

$$\mathcal{A}_S = \bigcup_i \mathcal{A}_i.$$

It is clear that the model  $\mathcal{A}_S$  is recursive.

**Claim 2.4** *The theory  $T_S$  of the model  $\mathcal{A}_S$  is a  $\omega_1$ -categorical but not  $\omega$ -categorical.*

**Proof.** The model  $\mathcal{A}_S$  satisfies the following list of properties which can be written as an infinite set of statements in the language of the first order logic.

1. For all  $x$  if  $P_{i+1}(x)$  holds, then  $P_i(x)$  also holds. Moreover  $\forall x P_0(x)$  is true.
2. For each  $i \in \omega$  there exists a unique  $x$  such that  $P_i(x) \& \neg P_{i+1}(x)$ ,  $i \in \omega$ .
3. For all  $k, s \in \omega$ , if  $R_{k,s}(x_1, \dots, x_k)$  holds, then  $x_1, \dots, x_k$  are pairwise distinct.
4. Let  $k \in S$ . For every  $s \in \omega$  there exists a  $j \in \omega$  such that  $\forall n \leq s \exists m < j H(n, m, s)$ . Let  $j_s$  be the minimal number which has this property. Then for all pairwise distinct  $x_1 \dots x_k$  if  $P_{j_s}(x_1) \& \dots P_{j_s}(x_k)$  holds, then  $R_{k,s}(x_1, \dots, x_k)$  holds.
5. Let  $k \notin S$ . There exists an  $s_0$  such that for all  $s \geq s_0$  and for all  $x_1, \dots, x_k$ ,  $R_{k,s}(x_1, \dots, x_k)$  does not hold.

Let  $\mathcal{A}$  be a model of  $T_S$ . Consider the set  $\bigcap_i P_i^{\mathcal{A}}$ . For any two elements  $a, b \in \bigcap_i P_i^{\mathcal{A}}$  there exists an automorphism  $\alpha$  of the model  $\mathcal{A}$  such that  $\alpha(a) = b$ . Thus a proof of  $\omega_1$ -categoricity can be based on the following observation. Two models  $\mathcal{B}$  and  $\mathcal{C}$  of the theory  $T_S$  are isomorphic if and only if the cardinalities of the sets  $\bigcap_i P_i^{\mathcal{B}}$  and  $\bigcap_i P_i^{\mathcal{C}}$  are equal. Hence if  $\mathcal{B}$  and  $\mathcal{C}$  are models of cardinality  $\omega_1$ , then both  $\bigcap_i P_i^{\mathcal{B}}$  and  $\bigcap_i P_i^{\mathcal{C}}$  have exactly  $\omega_1$  elements. It follows that  $\mathcal{B}$  and  $\mathcal{C}$  are isomorphic.  $\square$

From the proof of Claim 2.4, it follows that if  $\mathcal{B}$  is a countable unsaturated model of the theory  $T_S$ , then  $\bigcap_i P_i^{\mathcal{B}}$  has a finite number of elements.

**Claim 2.5** *If  $\mathcal{C}$  is a countable and unsaturated model of  $T_S$ , then  $\mathcal{C}$  has a recursive presentation.*

**Proof.** Let  $\mathcal{C}$  be a countable, unsaturated model of  $T_S$ . The set  $\bigcap_i P_i^{\mathcal{C}}$  has a finite number of elements, say  $n$ . We construct a recursive presentation of  $\mathcal{C}$  by stages.

Let  $a_1, \dots, a_n$  be new symbols. In our construction of a recursive presentation  $\mathcal{A}$  of  $\mathcal{C}$  we put the elements  $a_1, \dots, a_n$  into  $\bigcap_i P_i^{\mathcal{A}}$ . Let  $p_1, \dots, p_n$  be the all elements of  $S \cap \{0, 1, \dots, n\}$ .

**Stage 0.** Define  $\mathcal{A}_0 = (\{0, a_1, \dots, a_n\}, P_0)$ , letting  $P_0(0), P_0(a_1), \dots, P_0(a_n)$  hold.

**Stage t+1.** The domain  $A_{t+1}$  of  $\mathcal{A}_{t+1}$  is  $\{0, \dots, t+1, a_1, \dots, a_n\}$ . The signature of the  $\mathcal{A}_{t+1}$  is

$$\sigma_{t+1} = (P_0, \dots, P_{t+1}, R_{1,0}, \dots, R_{1,t+1}, \dots, R_{t+1,0}, \dots, R_{t+1,t+1}).$$

For each  $i \leq t+1$  let  $P_i(x)$  hold if and only if  $x \geq i$  or  $x \in \{a_1, \dots, a_n\}$ . For  $k, s \leq t+1$ , let  $R_{k,s}(x_1, \dots, x_s)$  hold if and only if one of the followings holds:

1.  $k \in \{p_1, \dots, p_n\}$ ,  $(x_1, \dots, x_k) \in \{a_1, \dots, a_n\}^n$ , and  $x_1, \dots, x_k$  are pairwise distinct.
2.  $\{x_1, \dots, x_k\} \setminus \{a_1, \dots, a_n\} \neq \emptyset$ , the elements  $x_1, \dots, x_k$  are pairwise different, and for the maximal number  $j \leq t+1$  such that all  $P_j(x_1), \dots, P_j(x_k)$  hold we have  $\forall n \leq s \exists m \leq j H(n, m, k)$ .

Thus this stage defines the structure  $\mathcal{A}_{t+1}$ . For each  $i \in \omega$   $\mathcal{A}_{i+1}$  is an extension and expansion of  $\mathcal{A}_i$ . Define  $\mathcal{A}$  by  $\mathcal{A} = \bigcup_i \mathcal{A}_i$ . It is clear that the structure  $\mathcal{A}$  is recursive and isomorphic to the model  $\mathcal{C}$ .  $\square$

**Claim 2.6** *The countable saturated model  $\mathcal{B}$  of  $T$  does not have a recursive presentation.*

**Proof.** Suppose that  $\mathcal{B}$  is recursive. Since  $\mathcal{B}$  is saturated the number of elements in  $\bigcap_i P_i^{\mathcal{B}}$  is infinite. It can be checked that for each  $k \in \omega$ ,  $k \in S$  if and only if there exist different elements  $y_1, \dots, y_k$  from  $\bigcap_i P_i^{\mathcal{B}}$  such that for all  $s \geq 1$ ,  $R_{k,s}(y_1, \dots, y_k)$  holds. The set  $S$  would then be a  $\Sigma_2^0$ -set. This contradicts with our assumption that  $S \in \Pi_2^0 \setminus \Sigma_2^0$ .  $\square$

These claims prove Theorem 3.  $\square$

Thus the above theorems prove the following corollary about **spectra of recursive models (SRM)** of  $\omega_1$  categorical theories.

**Corollary 2.1** 1. *There exists an  $\omega_1$ -categorical but not  $\omega$  categorical theory  $T$  such that  $\text{SRM}(T) = \omega - \{0\} \cup \{\omega\}$ .*

2. There exists an  $\omega_1$ -categorical but not  $\omega$  categorical theory  $T$  such that  $\mathbf{SRM}(T) = \omega$ .  $\square$

In the next theorem, which answers Question 1.3, we provide an example of a theory  $T_S$  with exactly 3 countable models of which only the saturated model is recursively presentable. To prove that  $T_S$  has exactly 3 countable models, we use the known ideas which show that the theory of the model  $(Q, \leq, c_0, c_1, \dots)$ , where  $\leq$  is the linear ordering of rationals, and the constants are such that  $c_0 > c_1 > c_2 > \dots$ , has exactly 3 countable models [14].

**Theorem 2.3** *There exists a theory  $T$  with exactly 3 countable models such that the only model of  $T$  which has a recursive presentation is the saturated model.*

**Proof.** Let  $Q$  be the set of all rational numbers. For each cardinal number  $m \in \omega \cup \{\omega\}$  define a structure  $Q_0(m)$  as follows. The domain of the structure is

$$\{q \in Q \mid 1 \leq q\} \cup \{c_{q,1}, \dots, c_{q,m} \mid q \in Q\},$$

where  $\{c_{q,i} \mid q \in Q, 1 \leq i \leq m\}$  is a set of new elements. The signature of the model is  $(\leq, f)$ , where  $\leq$  is a binary predicate and  $f$  is a unary function symbol. The predicate  $\leq$  and the function  $f$  are defined as follows. For all  $x, y$  we have  $x \leq y$  if and only if  $x, y \in Q$  and  $x$  is less or equal to  $y$  as rational numbers. For all  $z, y$  define  $f(z) = y$  if and only if for some rational number  $q$ ,  $y = q$  and  $z \in \{c_{q,1}, \dots, c_{q,m}\}$  or  $y = z = q$ . Let  $Q(m)$  be the structure obtained from  $Q_0(m)$  by removing the elements  $1, c_{1,1}, \dots, c_{1,m}$  from the domain of  $Q_0(m)$ .

If  $\mathcal{A}$  and  $\mathcal{B}$  are isomorphic copies of the structures  $Q_0(n)$  and  $Q_0(m)$ , respectively, and  $A \cap B = \emptyset$ , then one can naturally define the isomorphism type of the structure  $Q_0(n) + Q_0(m)$  as follows. The domain of the new structure is  $A \cup B$ . The predicate  $\leq$  in the new structure is the least partial ordering which contains the partial orderings of  $\mathcal{A}$ , the partial ordering of  $\mathcal{B}$ , and the relation  $\{(x, y) \mid x \in A \& f^{\mathcal{A}}(x) = x \& y \in B \& f^{\mathcal{B}}(y) = y\}$ . The unary function  $f$  in the new structure is the union of the unary operations of the first and the second structures.

If  $n_0, n_1, n_2, \dots, n_i, \dots$ ,  $i < \omega$  is a sequence of natural numbers, then as above we can define the structure

$$Q_0(n_0) + Q_0(n_1) + Q_0(n_2) + \dots .$$

Let  $S$  be a set in  $\Delta_2^0$  which is not the range of a limitwise monotonic function. There exists a recursive function  $g$  such that, for all  $n$   $h(n) = \lim_s g(n, s)$  exists and  $\text{range}(h) = S$ . Consider the model  $Q_0(S)$  defined by

$$Q_0(h(0)) + Q_0(h(1)) + Q_0(h(2)) + \dots .$$

Define the theory  $T_S$  to be the theory of the structure  $Q_0(S)$ .

**Claim 2.7.** *The theory  $T_S$  has exactly three countable models.*

**Proof.** The first model of  $T_S$  is  $Q_0(S)$ . This model is the prime model of the theory  $T_S$ . The second model of  $T_S$  is

$$Q'(S) = Q_0(h(0)) + Q_0(h(1)) + Q_0(h(2)) + \dots + Q_0(\omega).$$

The third model  $\mathcal{M}$  of  $T_S$  is

$$Q(h(0)) + Q_0(h(1)) + Q_0(h(2)) + \dots + Q(\omega).$$

These structures are indeed models of  $T_S$ . To see this, note that  $Q_0(S)$  is a submodel of  $Q'(S)$ , and  $Q'(S)$  is a submodel of  $\mathcal{M}$ . It can be checked that for any formula  $\exists x\phi(x, a_1, \dots, a_n)$  and all  $a_1, \dots, a_n \in Q_0(S)$  ( $a_1, \dots, a_n \in Q'(S)$ ) if the formula  $\exists x\phi(x, a_1, \dots, a_n)$  is true in  $Q'(S)$  (in  $\mathcal{M}$ ) then there exists a  $b \in Q_0(S)$  ( $b \in Q'(S)$ ) such that  $\phi(b, a_1, \dots, a_n)$  is true in  $Q_0(S)$  (in  $Q'(S)$ ). Therefore the embeddings are elementary.

We have to prove that any countable model of  $T_S$  is isomorphic to one of the three models described above. Let  $\mathcal{A}$  be a model of  $T_S$ . For each  $i \in \omega$  we define by induction an element  $a_i \in A$  as follows.

The element  $a_0$  is the minimal element with respect to the partial ordering in  $\mathcal{A}$ . Note that the set  $\{b | b \neq a_0 \& f(b) = a_0\}$  has exactly  $h(0)$  elements. Also put  $k_0 = 0$ .

Suppose that the elements  $a_0, \dots, a_{i-1} \in A$  and the numbers  $k_0, \dots, k_{i-1}$  have been defined. Let  $k_i$  be the least element such that  $h(k_i) \neq h(k_j)$  for  $j = 1, \dots, i-1$ . The element  $a_i$  is the one such that the following properties hold:

1. The set  $\{b|b \neq a_i \& f(b) = a_i\}$  has exactly  $h(k_i)$  elements,
2. For each  $x < a_i$  the cardinality of the set  $\{b|b \neq x \& f(b) = x\}$  is in  $\{h(k_0), \dots, h(k_{i-1})\}$ .

Consider the sequence  $a_0, a_1, a_2, \dots$ . Clearly  $a_0 < a_1 < a_2 < \dots$ . Thus we have three cases:

**Case 1.**  $\lim_i a_i$  does not exists and for any  $x \in A$  such that  $f(x) = x$  there exists an  $i$  such that  $a_i \geq x$ ,

**Case 2.**  $\lim_i a_i$  exists,

**Case 3.**  $\lim_i a_i$  does not exists and there exists an  $x$  such that  $f(x) = x$  and  $x \geq a_i$  for all  $a_i$ .

In the first case  $\mathcal{A}$  is isomorphic to  $Q_0(S)$ . In the second case  $\mathcal{A}$  is isomorphic to  $Q'(S)$ . In the third case  $\mathcal{A}$  is isomorphic to  $\mathcal{M}$ . Note that  $Q_0(S)$  is the prime model. The model  $Q'(S)$  is not saturated since it does not realize the type containing  $\{x > a_i \& c > x | i \in \omega\}$ , where  $c = \lim_i a_i$ . Hence  $\mathcal{M}$  is the saturated model of  $T_S$ .  $\square$

**Claim 2.8** *The unsaturated models of the theory  $T_S$  do not have recursive presentations.*

**Proof.** Consider the prime model  $Q_0(S)$ . Suppose  $Q_0(S)$  is a recursive model. Then it can be easily checked that the set  $S$  is the range of a limitwise monotonic function. This contradicts the assumption on  $S$ . If the other unsaturated model

$$Q'(S) = Q_0(h(0)) + Q_0(h(1)) + Q_0(h(2)) + \dots + Q_0(\omega)$$

were recursive, then  $Q_0(S)$  would be a recursively enumerable submodel of the model  $Q'(S)$ . Hence  $Q_0(S)$  would have a recursive presentation. This is again a contradiction.  $\square$

**Claim 2.9** *The saturated model  $\mathcal{M}$  of the theory  $T$  has a recursive presentation.*

**Proof.** We present a construction of the saturated model  $\mathcal{M}$  by stages. The construction will clearly show that the saturated model has a recursive presentation.

**Stage 0.** Consider the structure  $Q_0(g(0,0)) + Q(\omega)$ . Denote this model by  $\mathcal{A}_0$ .

**Stage n+1.** Suppose that  $\mathcal{A}_n$  has been defined and is isomorphic to

$$Q_0(g(0,n)) + \dots + Q_0(g(n,n)) + Q(\omega).$$

Compute  $g(0,n+1), \dots, g(n+1,n+1)$ . Let  $i \leq n$  be the minimal number such that  $g(i,n) \neq g(i,n+1)$ .  $\mathcal{A}_n$  can be extended to a structure  $\mathcal{A}_{n+1}$  isomorphic to

$$Q_0(g(0,n+1)) + \dots + Q_0(g(i-1,n+1)) + Q_0(g(i,n+1)) + \dots + Q_0(g(n+1,n+1)) + Q(\omega).$$

To see this, take the substructure

$$Q_0(g(i,n)) + \dots + Q_0(g(n,n)) + Q(\omega)$$

of  $\mathcal{A}_n$ ; extend this substructure to  $Q(\omega)$ ; insert the new structure

$$Q_0(g(i,n+1)) + \dots + Q_0(g(n+1,n+1))$$

between the structures  $Q_0(g(0,n+1)) + \dots + Q_0(g(i-1,n+1))$  and the extended structure  $Q(\omega)$ . The structure obtained in this way is  $\mathcal{A}_{n+1}$ .

Thus we have the sequence

$$\mathcal{A}_0 \subset \mathcal{A}_1 \subset \mathcal{A}_2 \subset \dots$$

Define  $\mathcal{A}_\omega = \bigcup_i \mathcal{A}_i$ . It is easy to see that the model  $\mathcal{A}_\omega$  is isomorphic to

$$Q_0(h(0)) + Q_0(h(1)) + \dots + Q_0(h(n)) + \dots + Q(\omega).$$

Now it is clear that the above description can be effectivized.  $\square$

These claims prove the theorem.  $\square$

Finally we have to prove the promised recursion theoretic lemma.

**Proof of the Recursion Theoretic Lemma.** Let  $\phi_e(x,t)$ ,  $e \in \omega$ , be a uniform enumeration of all partial recursive functions  $\phi$  such that for all  $t' \geq t$  if  $\phi(x,t')$  is defined, then  $\phi(x,t)$  is defined and  $\phi(x,t) \leq \phi(x,t')$ . At stage  $s$  of our construction we define a finite set  $A_s$  in such a way that  $A(y) = \lim_s A_s(y)$  exists for all  $y$ . We satisfy the requirement  $R_e$  asserting that, if  $f_e(x) = \lim_t \phi_e(x,t) < \omega$  for all  $x$ , then  $\text{range}(f_e) \neq A$ .

The strategy for a single  $R_e$  is as follows: At stage  $s$  pick a witness  $m_e$ , enumerate  $m_e$  into  $A$  (i.e.  $A_s(m_e) = 1$ ). Now  $R_e$  is satisfied (since  $m_e$  remains in  $A$ ) unless at some later stage  $t_0$  we find an  $x$  such that  $\phi_e(x, t_0) = m_e$ . If so,  $R_e$  ensures that  $A(\phi_e(x, t)) = 0$  for all  $t \geq t_0$ . Thus, either  $f_e(x) \uparrow$  or  $f_e(x) \downarrow$  and  $f_e(x) \notin A$ .

Keeping  $\phi_e(x, t)$  out of  $A$  for all  $t \geq t_0$  can conflict with a lower priority ( $i > e$ ) requirement  $R_i$  since it maybe the case that  $m_i = \phi_e(x, t')$  for some  $t' > t_0$ . However, if  $f_e(x) \downarrow$ , then this holds permanently for just one number, and if  $f_e(x) \uparrow$ , then the restriction is transitory for each number. So each lower priority  $R_i$  will be able to choose a stable witness at some stage.

**Construction.** At stage  $s$  we try to determine the values of parameters  $m_e$ ,  $x_e$ , and  $n_e = \phi_e(x_e, s)$  for  $R_e$ . Each parameter may remain undefined. Moreover we define the approximation  $A_s$  to  $A$  at stage  $s$ .

**State 0.** Let  $A_0 = \emptyset$ , and declare all parameters to be undefined.

**Stage s.** For each  $e = 0, \dots, s-1$  in turn go through substages  $e$  by performing the following actions.

1. If  $m_e$  is undefined, let  $m_e$  be the least number in  $\omega^{[e]}$  greater than all  $m_i$  ( $i < e$ ) which is not equal to any  $n_i$ . Let  $A_s(m_e) = 1$  and proceed to the next substages, or to stage  $s+1$  if  $e = s-1$ .
2. If  $x_e$  is undefined and  $\phi_e(x, s) = m_e$  for some  $x$ , let  $x_e = x$ ,  $n_e = m_e$ , and  $A_s(n_e) = 0$ , and proceed to the next stage  $s+1$  if  $e = s-1$ .
3. Let  $n_e = \phi_e(x_e, s)$  and  $A_s(n_e) = 0$ . If  $n_e = m_i$  for some  $i > e$ , declare all the parameters of the  $R_j$ ,  $j \geq i$ , to be undefined.

For each  $y$ , if  $A_s(y)$  is not determined by the end of stage  $s$ , then assign to  $A_s(y)$  its previous value  $A_{s-1}(y)$ . The stage is now completed.

Now we will verify that the construction succeeds.

**Claim 2.10** *Each  $m_e$  is defined and is constant from some stage on.*

**Proof.** Suppose inductively that the claim holds for each  $i < e$ . Let  $s_0$  be a stage such that each  $m_i$  has reached its limit for  $i < e$ , and if  $x_i$  ever

becomes defined after  $s_0$ , and  $\lim_s n_{i,s} < \infty$ , then the limit has been reached at  $s_0$ . Moreover, let  $k \geq e$  be the least number which does not equal any of these limits and is greater than all  $m_i$  for  $i < e$ . Also suppose that  $n_{i,s_0} > k$  if  $\lim_s n_{j,s} = \infty$ , ( $j < e$ ). If  $m_e$  is cancelled after stage  $s_0$ , then  $m_e = k$  is permanent from the next stage on. This proves the claim.

**Claim 2.11** *For each  $y$ ,  $\lim_s A_s(y)$  exists. Therefore the set  $A = \lim_s A_s$  is a  $\Delta_2^0$ -set.*

**Proof.** Suppose that  $y \in \omega^{[e]}$ , and let  $s_0$  be a stage at which  $m_e$  has reached its limit. Since  $y$  can only be enumerated into  $A$  if  $y = m_e$ , after stage  $s_0$   $A(y)$  can change at most once. This proves the claim.

**Claim 2.12** *Suppose that  $f_e(x) = \lim_t \phi_e(x, t)$  exists for each  $x$ . Then  $A \neq \text{range}(f_e)$ .*

**Proof.** Suppose that  $A = \text{range}(f_e)$ . Let  $s_0$  be the stage at which  $m_e$  reaches its limit. Then at some stage  $s > s_0$  we must reach the second instruction of the construction, otherwise  $A(m_e) = 1$  but  $m_e \notin \text{range}(f_e)$ . Suppose that  $\phi_e(x, s) = m_e$  for the minimal  $s \geq s_0$  at which we reach the second instruction of the construction. It follows that for  $t \geq s$ ,  $n_e = \phi_e(x, t)$  and  $A_t(n_e) = 0$ . So  $A(f_e(x)) = 0$ . This contradiction proves the claim and hence the lemma.  $\square$

**Remark.** It is possible to make  $A$  d.r.e, i.e.  $A = B - C$  for some r.e. sets  $B, C$ . To do so, we have to set aside an interval  $I_e$ , roughly of size  $2^e$ , for  $R_e$ ,  $I_0 < I_1 < \dots$ . As a first choice for  $m_e$ , we take the maximal element of  $I_e$ , and then we proceed downward. The point is that, if  $R_e$  is injured by  $R_i$ ,  $i < e$ , via  $n_i = m_e$ , then all further values of  $n_i$  are above the next values of  $m_e$  (unless  $R_i$  injured itself later). Obviously  $A$  can be neither r.e. nor co-r.e.

## References

- [1] J. Baldwin, A. Lachlan, *On Strongly Minimal Sets*, Journal of Symbolic Logic, v.36, no 1, 1971.

- [2] Yu. Ershov, *Constructive Models and Problems of Decidability*, Moskow, Nauka, 1980.
- [3] S.Goncharov, *Constructive Models of  $\omega_1$ -categorical Theories*, Matematicheskie Zametki, v.23, no 6, 1978.
- [4] S. Goncharov, *Strong Constructivability of Homogeneous Models*, Algebra and Logic, v.17, no 4, 1978.
- [5] *Logic Notebook*, Novosibirsk University, Editors Yu.Ershov, S.Goncharov, 1986.
- [6] N.Khissamiev, *On Strongly Constructive Models of Decidable Theories*, Izvestiya AN Kaz. SSR, no 1, 1974.
- [7] K.Kudeiberganov, *On Constructive Models of Undecidable Theories*, Siberian Mathematical Journal, v.21, no 5, 1980.
- [8] L.Harrington, *Recursively Presentable Prime Models*, Journal of Symbolic Logic, v.39, no 2, 1973.
- [9] T. Millar, *Ph.D. Dissertation*, Cornell University, 1976.
- [10] T. Millar, *Omitting Types, Type Spectrums, and Decidability*, Journal of Symbolic Logic, v.48, no 1, 1983.
- [11] T. Millar, *Foundations of Recursive Model Theory*, Annals of Mathematical Logic, v. 13, 1978.
- [12] M.Morley, *Decidable Models*, Israel Journal of Mathematics, v.25, 1976.
- [13] M.Peret'ytkin, *On Complete Theories with Finite Number of Countable Models*, Algebra and Logic, v.12, no 5, 1973.
- [14] J. Rosenstein, *Linear Orderings*, New York, Academic Press, 1982.
- [15] G. Sacks, *Saturated Model Theory*, Reading, Mass., W. A. Benjamin, 1972.