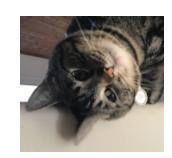
## Intro to a blockchain - the Teddy Token

A blockchain is a secure record of transactions



Transactions may consist of

- registrations or exchanges of "digital assets" such as coins and tokens\*
- registrations of and transactions with "smart contracts"

Transactions in a blockchain are secure: they cannot be reversed or revised

If an attempt is made to reverse or revise a transaction, the attempt is obvious and can be stopped

Transactions are grouped into blocks and the blocks are linked into a chain by a secure method involving cryptographic hashes

\*Technically, Teddy Token is a digital coin, but "token" rhymes with Teddy!

# Technology involved in blockchains

- cryptographic hashes
- hexadecimal and binary numbers
- consensus by proof-of-work
- modular arithmetic
- public key encryption
- peer-to-peer networks over the Internet P2P

All discussed here except P2P

## Transactions are anonymous and irreversible

From: bf61f561918f37cf441fb8a6a1094b7a

To: 9412daed1e0bd204f652677a80192ea9

Amount: 2

Date: Mon Feb 22 2021 13:44:10 GMT-0800 (PST)

Hash: b819e116c0aca3dd431c411b31040160

Addresses of sender and recipient are unique and the people's identity are known only to themselves... or their wallet company, that is

A hash is a cryptographic digest of the transaction that can be used to verify that the transaction hasn't be revised after it was entered into the blockchain

Any tampering? Re-hash and see if the new hash matches that in the chain!

Teddy Token transactions have only one sender and one recipient. Other blockchains allow multiple senders and recipients in one transaction.

Other blockchains may charge the sender a transaction fee

## Cryptographic hashes are key to blockchains

A cryptographic hash algorithm processes a message to produce a hash value that is:

- of fixed length, regardless of the length of the message > modular arithmetic
- that cannot be decrypted to determine the original message > "confusion"
- that changes dramatically when only a trivial change is made in the message > "diffusion"

The terms "confusion" and "diffusion" in cryptography are due to Claude Shannon

Teddy Token uses the MD2 hash - out of date but relatively easy to understand Bitcoin uses SHA-256 - in current use, e.g., by Bitcoin

From: bf61f561918f37cf441fb8a6a1094b7a

To: 9412daed1e0bd204f652677a80192ea9

Amount: 2

Date: Mon Feb 22 2021 13:44:10 GMT-0800 (PST)

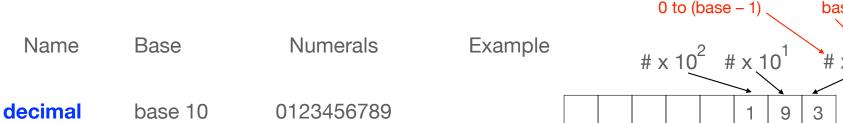
Hash: b819e116c0aca3dd431c411b31040160

Addresses are hashes of a user's public cryptographic key

This hash is a cryptographic digest of the transaction that can be used to verify that the transaction hasn't be revised after it was entered into the blockchain What are those strings of numbers and letters? 000869e6e37d5f11f5b71d2641c139d4 Those are **hexidecimal** numbers

## Some Number Systems

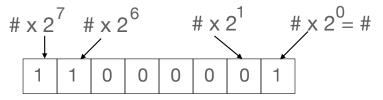
power, counting from 0 to the left



Why? Easy to compute in your head & by hand

 $1 \times 10^{2} + 9 \times 10^{1} + 3 = 100 + 90 + 3 = 193$  decimal

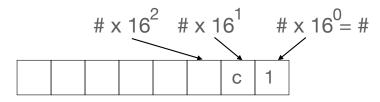
binary base 2 01



Why? Digital storage is by bits: 0/1, on/off, high/low... This is how numbers are stored physically by computers

$$2^{7} + 2^{6} + 1 = 128 + 64 + 1 = 193$$
 decimal

hexidecimal base 16 0123456789abcdef



Why? Easy conversion between hex & binary, 1 hex to 4 bits  $12 \times 16^{1} + 1 = 192 + 1 = 193$  decimal

$$12 \times 16^{1} + 1 = 192 + 1 = 193$$
 decima

## The start of the Teddy Token blockchain - a secure record of anonymous transactions

A CRYPTOGRAPHIC HASH of a message - in this case a block header - will change dramatically if any single character in the message is changed. This and the presence of the previous block hash in each block header ensure the integrity of the blockchain

#### Genesis Block - Catoshi's gift to his friends

Version: TedTok 01 Block number: 0

Merkle root: 5f5de6dab6f9f1a636690177cbe629dc Date: Mon Feb 22 2021 13:43:47 GMT-0800 (PST)

Target: 3 Nonce: 11911

Hash: 0001b04ff27cb9ab55b86477f95815de

Number transactions: 5

To: 8fca7b9ab6fac411b2443c1303d3bc56

Amount: 10

Hash: 425a3eba0375c91ff94af7ff9fd29209

\_\_\_\_\_\_

To: f157bab61715d2bca9d81ada45bad57c Amount: 10

Hash: 402ea2348edb95de31565ed152729aaf

To: 263f99c540f4168132cc5fee5726b5ce

Amount: 10

Hash: ebc383bbf91325b0997c9147591c35e5

To: bf61f561918f37cf441fb8a6a1094b7a

Amount: 10

Hash: dba01bfd264db14cb52123c6dbeb9f51

\_\_\_\_\_\_

To: 9412daed1e0bd204f652677a80192ea9

Amount: 10

Hash: d4704122a45d67f174fa36d53dd6b951

\_\_\_\_\_\_

#### Block 1

#### block header

Version: TedTok 01 Block number: 1

Previous hash: 0001b04ff27cb9ab55b86477f95815de Merkle root: 21e3776b697c3d22922876e7e8ae36c2

Date: Mon Feb 22 2021 13:44:14 GMT-0800 (PST)

Target: 3 Nonce: 4328

Hash: 000f5b9d7ca44d91ba33e9934e82cf27

Number transactions: 3

Miner reward to: f157bab61715d2bca9d81ada45bad57c

Amount: 0.99

Hash: ad5ce451141e8c30baf4f55e6b4ddb14

-----

From: bf61f561918f37cf441fb8a6a1094b7a To: 9412daed1e0bd204f652677a80192ea9

Amount: 2

Date: Mon Feb 22 2021 13:44:10 GMT-0800 (PST)

Hash: b819e116c0aca3dd431c411b31040160

From: f157bab61715d2bca9d81ada45bad57c To: 8fca7b9ab6fac411b2443c1303d3bc56

Amount: 1

Date: Mon Feb 22 2021 13:44:03 GMT-0800 (PST)

Hash: 5c0a2c83ae1366d3074e923198fec388

\_\_\_\_\_\_

#### Block 2

Version: TedTok 01 Block number: 2

Previous hash: 000f5b9d7ca44d91ba33e9934e82cf27

Merkle root: 2569e9f51001fdb295b836228bbfcc84 Date: Mon Feb 22 2021 13:44:40 GMT-0800 (PST)

Target: 3 Nonce: 856

Hash: 000869e6e37d5f11f5b71d2641c139d4

Number transactions: 3

Miner reward to: 263f99c540f4168132cc5fee5726b5ce

Amount: 0.99

Hash: 3490eb317353dcc32715788714a13769

From: 9412daed1e0bd204f652677a80192ea9 To: 8fca7b9ab6fac411b2443c1303d3bc56

Amount: 3

Date: Mon Feb 22 2021 13:44:36 GMT-0800 (PST)

Hash: bb281c938be6c2bee7ffc05bdb4f1e23

From: bf61f561918f37cf441fb8a6a1094b7a To: f157bab61715d2bca9d81ada45bad57c

Amount: 1

Date: Mon Feb 22 2021 13:44:29 GMT-0800 (PST)

Hash: 26a537138d436bbd55a7287edf152951

Any tampering? Re-hash and see if the new hash matches that in the chain!

## The genesis block of all blockchains - Satoshi's gift to Hal Finney of 50 bitcoins created out of thin air!

Genesis block - Bitcoin Wiki

#### Main network genesis block

Here is a representation of the genesis block<sup>[1]</sup> as it appeared in a comment in an old version of Bitcoin (line 1613 ©). The first section defines exactly all of the variables necessary to recreate the block. The second section is the block in standard printblock format, which contains shortened versions of the data in the first section.

```
GetHash()
           = 0x00000000019d6689c085ae165831e934ff763ae46a2a6c172b3f1b60a8ce26f
hashMerkleRoot = 0x4a5e1e4baab89f3a32518a88c31bc87f618f76673e2cc77ab2127b7afdeda33b
txNew.vin[0].scriptSig
                     = 486604799 4
txNew.vout[0].nValue
                     = 5000000000
txNew.vout[0].scriptPubKey =
0x5F1DF16B2B704C8A578D0BBAF74D385CDE12C11EE50455F3C438EF4C3FBCF649B6DE611FEAE06279A60939E028A8D65C10B7307
OP CHECKSIG
block.nVersion = 1
block.nTime
          = 1231006505
block.nBits
           = 0x1d00ffff
block.nNonce = 2083236893
CBlock(hash=00000000019d6, ver=1, hashPrevBlock=000000000000, hashMerkleRoot=4a5e1e,
nTime=1231006505, nBits=1d00ffff, nNonce=2083236893, vtx=1)
 CTransaction(hash=4a5e1e, ver=1, vin.size=1, vout.size=1, nLockTime=0)
   CTxIn(COutPoint(000000, -1), coinbase
CTxOut(nValue=50.00000000, scriptPubKey=0x5F1DF16B2B704C8A578D0B)
 vMerkleTree: 4a5e1e
```

#### Hash

The hash of the genesis block, 00000000019d6689c085ae165831e934ff763ae46a2a6c172b3f1b60a8ce26f, [1] has two more leading hex zeroes than were required for an early block.

#### Coinbase

The coinbase parameter (seen above in hex) contains, along with the normal data, the following text:<sup>[2]</sup>

The Times 03/Jan/2009 Chancellor on brink of second bailout for banks<sup>[1]</sup>



Bitcoin uses the SHA-256 hash, which is more secure than Teddy Token's MD2 hash and twice as long

The very first 50 bitcoin

a comment in the block included a reference to bank bailouts published in The Times, Jan 3, 2009

# The <u>blockchain header</u> lists all addresses and their current balances

## blockchain header

Address: 8fca7b9ab6fac411b2443c1303d3bc56

Balance: 14

-----

Address: f157bab61715d2bca9d81ada45bad57c

Balance: 10.99

-----

Address: 263f99c540f4168132cc5fee5726b5ce

Balance: 10.99

-----

Address: bf61f561918f37cf441fb8a6a1094b7a

Balance: 7

\_\_\_\_\_

Address: 9412daed1e0bd204f652677a80192ea9

Balance: 9

An anonymous user's ADDRESS is the hash of their PUBLIC CRYPTOGRAPHIC KEY

## The **MERKLE ROOT** is a hash produced from the hashes of the transactions in the block.

Version: TedTok\_01
Block number: 2

Previous hash: 000f5b9d7ca44d91ba33e9934e82cf27
Merkle root: 2569e9f51001fdb295b836228bbfcc84
Date: Mon Feb 22 2021 13:44:40 GMT-0800 (PST)

Target: 3
Nonce: 856

Hash: 000869e6e37d5f11f5b71d2641c139d4

Number transactions: 3

Miner reward to: 263f99c540f4168132cc5fee5726b5ce

Amount: 0.99

Hash: 3490eb317353dcc32715788714a13769

-----

From: 9412daed1e0bd204f652677a80192ea9 To: 8fca7b9ab6fac411b2443c1303d3bc56

Amount: 3

Date: Mon Feb 22 2021 13:44:36 GMT-0800 (PST)

Hash: bb281c938be6c2bee7ffc05bdb4f1e23

-----

From: bf61f561918f37cf441fb8a6a1094b7a To: f157bab61715d2bca9d81ada45bad57c

Amount: 1

Date: Mon Feb 22 2021 13:44:29 GMT-0800 (PST)

Hash: 26a537138d436bbd55a7287edf152951

\_\_\_\_\_

The Merkle root is placed in the block header.

The block header is hashed and the hash passed to the next block.

Any change in any transaction will be obvious in changes in the hash of the transaction >>

>> the Merkle root >>

>> the hash of the block header >>

>> the hash of all later blocks

The Merkle root is obtained from the transactions by a method that provides fast verification of individual transactions in the blockchain

#### Proof-of-work consensus method

- Any miner on the peer network can build a block by combining pending transactions
- which block is added to the chain is determined by a consensus method
- Teddy Token and Bitcoin use the "proof-of-work" method
- the protocol specifies a target number of zeros with which the final block header hash must start
- a miner must find the number the nonce to add to the header that achieves this
- the first miner who finds the correct nonce can add its block to the chain
- that miner receives a reward, which increases the number of coins in circulation

To: 8fca7b9ab6fac411b2443c1303d3bc56

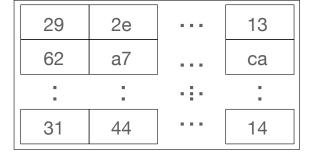
Hash
Address
Merkle root
Miner reward
Nonce
Target
Crypto key

nonce - number used once

## One step in the MD2 cryptographic hash (C array)

- "Confusion" by mixing message information with random numbers from the S-box
- "Diffusion" by value in current position in hash selecting S-box number for next message position such that confused information cascades and a change anywhere in message affects all locations by repeated passes through the message.

#### Substitution S-box (hexadecimal values)

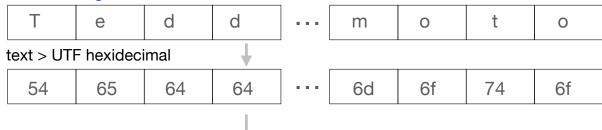


number hex c1 of S-box is hex ec

value from element

S-box is array of 256 elements, each containing a randomly placed, nonrepeating value in decimal range 0-255

#### The message

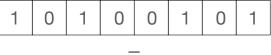


info for step 4

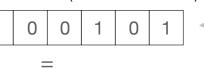
binary values



XOR (bitwise exclusive OR)



select value from element number hex c1 of S-box



0 0 0 0 0

ec

this message value is mixed with random value from S-box in hash from previous step, selected by operation with previous message value

info from step 3 for step 4

The hash being computed

info from step 4 for step 5

a5

MD2 hash length is 32 hex = 16 each, 8-bit bytes = 128 bits, regardless of length of message, and maintained at this length by MODULAR ARITHMETIC

#### Modular arithmetic

Messages are first converted to large integers using text-to-code tables, e.g., A = 65

All variables represent integers. The integers a and b are congruent modulo n.

$$a \equiv b \pmod{n}$$

where "mod" in parenthesis means that the modulus operation applies to all parts of the expression. This can also be expressed as

$$a = b + kn$$

That is, a is divisible k times by n, with remainder b.

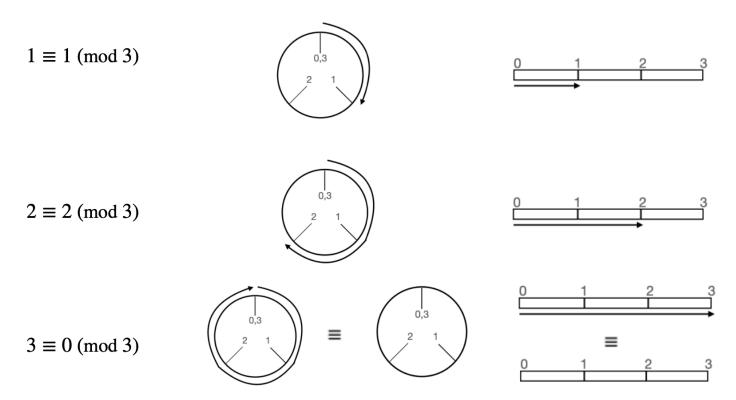
Those of us who use a 12 hour clock dial use modular arithmetic everyday. For example, the hour hand points to 6 on the dial at 6 am and 6 pm, where 6 pm is 18:00 in 24-hour time.

$$18 \equiv 6 \pmod{12}$$

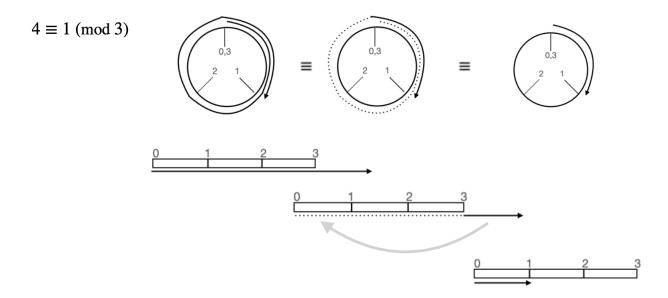
$$18 = 6 + 1(12)$$

## Modular arithmetic

Cryptography and hashing use very large integers. For illustration here, we use small integers. For some operations below, we will use a dial or linear scale with three divisions to simplify things. The relationships shown work for numbers of all magnitudes.



## Modular arithmetic



One way to think of this is that numbers greater than some multiple k of the modulus n "wraps around" so that the maximum value in the system is n. In the examples above k = 1 but k may have any integer value.

When we discuss hash algorithms, you will see that this "wrap around" feature is what keeps hash values a constant length, regardless of the length of the message being hashed.

The MD2 hash "wraps around" at 16 bytes = 32 hex numerals

# Addresses in blockchains are hashes of owners' public encryption keys

- one class of encryption is symmetric use same key for both encryption and decryption
- disadvantage is that anyone who discovers the key can produce fake messages and secretly decrypt private messages
- the class that blockchains use is <u>asymmetric</u> encryption
- public key encryption is asymmetric
- a message is encrypted with the public key, which may be known to the public
- an encrypted message cannot be decrypted with the public key
- the encrypted message can only be decrypted with the owner's private key

# Bitcoin uses Elliptic Curve Cryptography - Teddy Token uses RSA Cryptography

- both are asymmetric methods
- we discuss RSA cryptography in the following slides

The primes p and q are the basis numbers of the RSA public-key encryption method. Their choice results in the generation of the public and private keys. They are kept secret.

The public key consists of two integers, n and e. The integer e is chosen to be coprime with n = pq. That is, e < n and e will not evenly divide n.

The private key consists of two integers, n and d. The integer d is the "modular inverse" of e, such that

$$ed \equiv 1 \pmod{\phi(n)}$$

This can also be expressed as

$$ed = 1 + k\phi(n)$$

and  $\phi$  is Euler's phi function

$$\phi(n) = \phi(pq) = \phi(p)\phi(q) = (p-1)(q-1)$$

Message m < n is encrypted to the encoded message c by this equation:

$$m^e \equiv c \pmod{n}$$

The holder of the private key can decrypt c in order to recover m by this equation:

$$c^d \equiv m \pmod{n}$$

In actual RSA encryption, the values are large such that the computation of  $m^e$  cannot be done directly on a computer before the modulo operation, as can be done with small values such as  $2^3$ . This is because the value of  $m^e$  that would be obtained with the large values of m and e that are used would be so large that it could not be contained by a computer's method of storing numbers.

The solution of this problem is to compute the results using the modular exponentiation algorithm. In that algorithm, the largest value that must be stored during computation is  $m^2$ .

# Why is the private key secure & can't be determined from the public key?

- from the public key, n and e are known
- d has to be determined in order to get the private key
- d is the modular inverse of e:  $e^*d = 1 \pmod{\phi(n)}$
- to get d knowing e, one needs to get φ(n)
- $\phi(n) = (p-1)^*(q-1)$ , where  $n = p^*q$
- this would require finding which two prime numbers p and q when multiplied together equals the known value n
- this is called "prime factorization"
- with the large values of p and q used in RSA, the prime factorization of n would take a ridiculously long time

# See Web Labs at ReactorLab.net for interactive simulations of

- Teddy Token blockchain
- Cryptographic hash
- Public key encryption