

TOWARDS WIDE-COVERAGE SEMANTICS

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QUESTION:

How can we parse arbitrary text?

QUESTION:

How can we parse arbitrary text *and*
assign it a representation of its meaning?

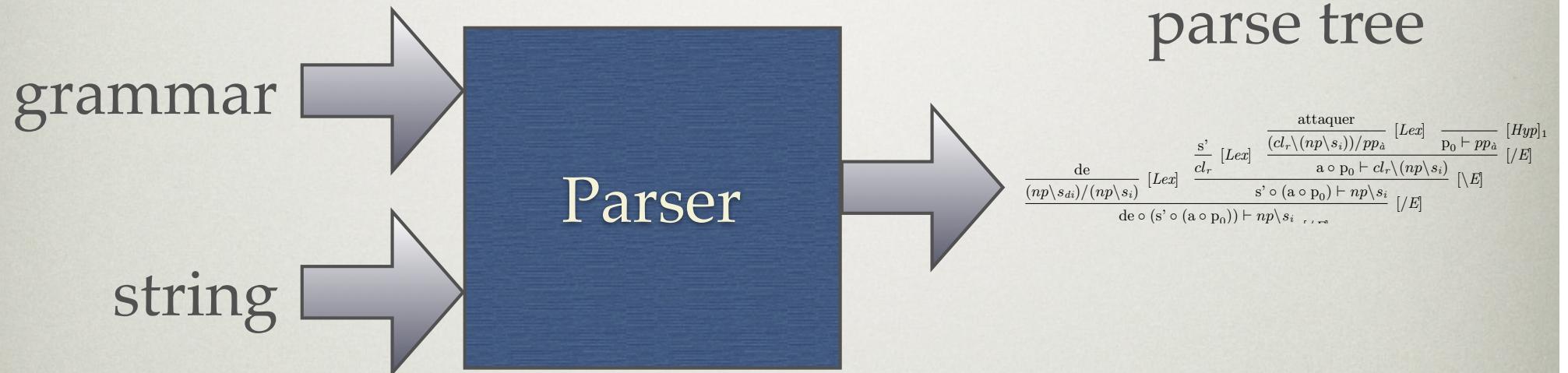
Two Solutions

1. Program a grammar by hand
2. Machine Learning

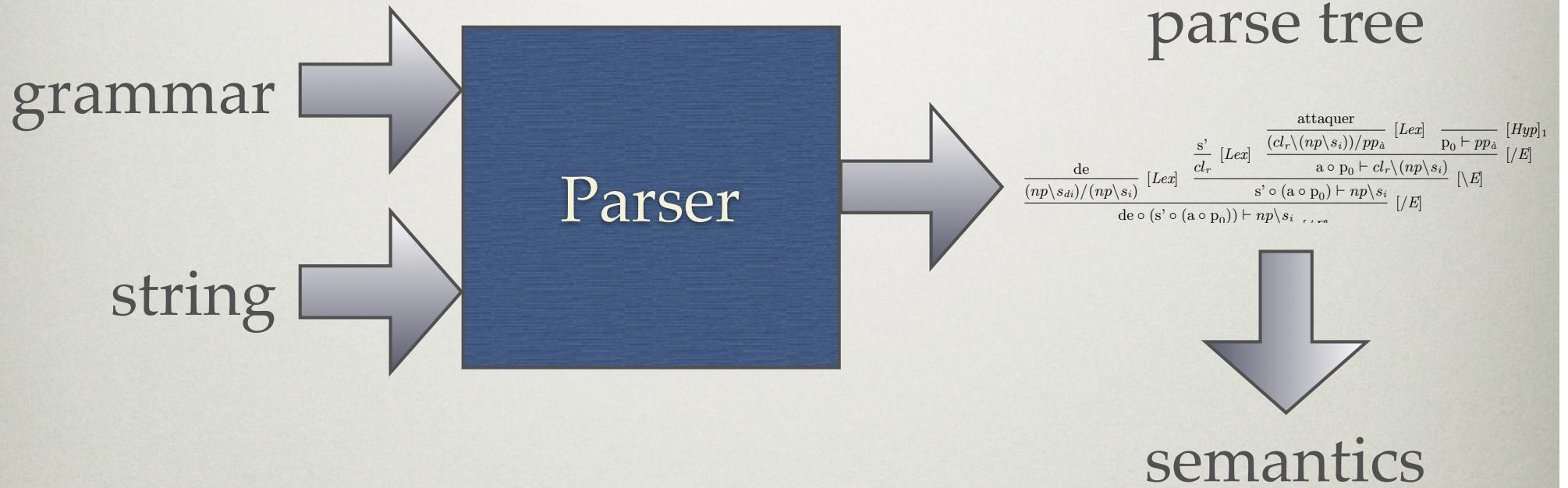
Writing grammars by hand is very useful and very important, but not well-suited to wide-coverage parsing

Treebanks contain a lot of information we can exploit. So we use treebank transformations to make them as close as possible to what we really want.

PARSING AND SEMANTICS



PARSING AND SEMANTICS



$(\lambda x. \exists y. (x y) \lambda z. ((\lambda v. \lambda w. \text{shot}(v, w) z) \text{ vito}))$

e_1	y_1					
$y_1 :$	<table border="1"> <tr> <td>$e_2 \ e_3 \ x_3$</td> </tr> <tr> <td>$x_3 = ?$</td> </tr> <tr> <td>aider.à(e_2, x_0, x_3, e_3)</td> </tr> <tr> <td>partir(e_3, x_3)</td> </tr> <tr> <td>demander(e_1, y_0, x_0, y_1)</td> </tr> </table>	$e_2 \ e_3 \ x_3$	$x_3 = ?$	aider.à(e_2, x_0, x_3, e_3)	partir(e_3, x_3)	demander(e_1, y_0, x_0, y_1)
$e_2 \ e_3 \ x_3$						
$x_3 = ?$						
aider.à(e_2, x_0, x_3, e_3)						
partir(e_3, x_3)						
demander(e_1, y_0, x_0, y_1)						

MANY APPLICATIONS BENEFIT FROM MEANING

- Textual entailment
- Question answering
- Anaphora resolution

ENTAILMENT EXAMPLE (FROM JOHNSON-LAIRD)

None of the artists is a beekeeper.

T All the beekeepers are chemists.

(There is at least one beekeeper.)

H ... artists ... chemists.

ENTAILMENT EXAMPLE (FROM JOHNSON-LAIRD)

None of the artists is a beekeeper.

T All the beekeepers are chemists.

(There is at least one beekeeper.)

H Some of the chemists are not artists

ENTAILMENT EXAMPLE (FROM RTE)

T Eating lots of foods that are a good source of fiber may keep your blood glucose from rising fast after you eat.

H Fiber improves blood sugar control.

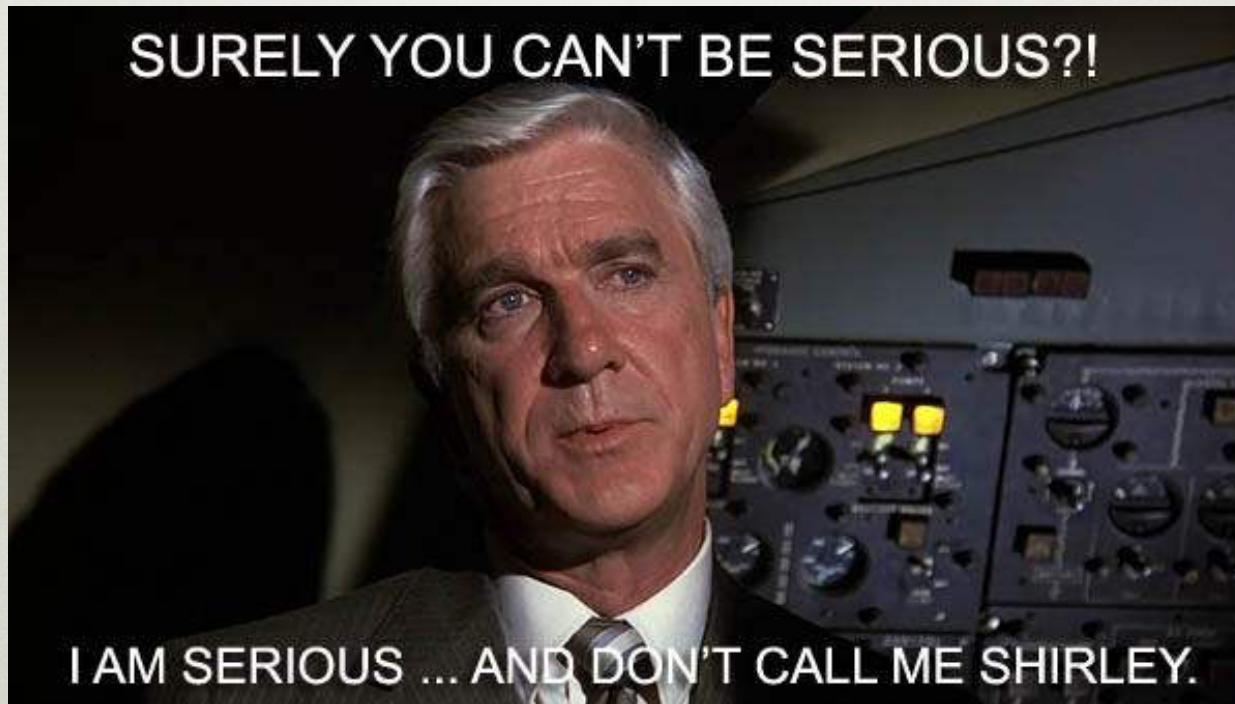
QUESTION ANSWERING NATURAL LANGUAGE INTERFACE



QUESTION ANSWERING NATURAL LANGUAGE INTERFACE



SURELY, YOU CAN'T BE SERIOUS



SURELY YOU CAN'T BE SERIOUS?!

I AM SERIOUS ... AND DON'T CALL ME SHIRLEY.

QUESTION ANSWERING NATURAL LANGUAGE INTERFACE

Wanted for general evilness; last seen at the tower of Barad-Dur; it's a giant eye, folks, kinda hard to miss



Watson does use deep parsing
and rather successfully

QUESTION ANSWERING NATURAL LANGUAGE INTERFACE

U.S. Cities

Its largest airport is named for a World War II hero; its second largest, for a World War II battle

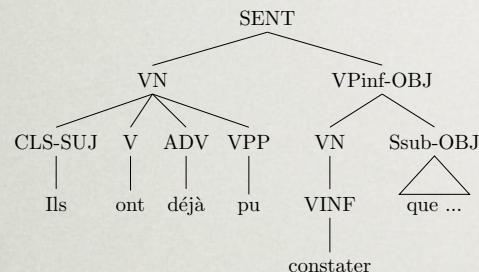


Still, there seems to be room for improvement in the parser component: here the second sentence should be interpreted as
“its second largest [airport] [is named] for a World War II battle.

Note: this is very difficult!
 $np\ pp \rightarrow np$
 $s ; np \rightarrow s$

OUTLINE

Treebank



Grammar Extraction

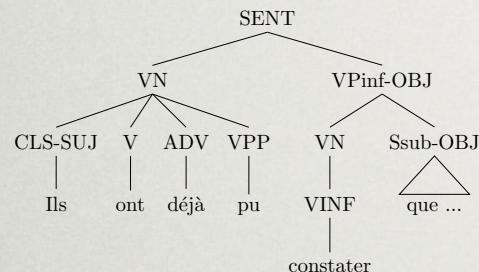
$$\frac{\frac{de}{(np \setminus s_{di}) / (np \setminus s_i)} [Lex]}{de \circ (s' \circ (a \circ p_0)) \vdash np \setminus s_i} \quad \frac{\frac{s'}{cl_r} [Lex] \quad \frac{\frac{attaquer}{(cl_r \setminus (np \setminus s_i)) / pp_a} [Lex] \quad \frac{p_0 \vdash pp_a}{[]} [Hyp]_1}{a \circ p_0 \vdash cl_r \setminus (np \setminus s_i)} [/ E]}{s' \circ (a \circ p_0) \vdash np \setminus s_i} [/ E]$$

Applications

e_1	y_1
$y_1 :$	$e_2 \ e_3 \ x_3$ $x_3 = ?$ $\text{aider_à}(e_2, x_0, x_3, e_3)$ $\text{partir}(e_3, x_3)$ $\text{demander}(e_1, y_0, x_0, y_1)$

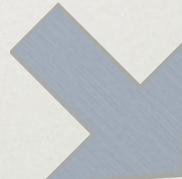
OUTLINE

Treebank



Grammar Extraction

$$\frac{\frac{\frac{de}{(np \setminus s_{di}) / (np \setminus s_i)} [Lex] \quad \frac{s'}{cl_r} [Lex] \quad \frac{\text{attaquer}}{(cl_r \setminus (np \setminus s_i)) / pp_a} [Lex] \quad \frac{p_0 \vdash pp_a}{a \circ p_0 \vdash cl_r \setminus (np \setminus s_i)} [Hyp]_1}{a \circ p_0 \vdash np \setminus s_i} [/E]}{s' \circ (a \circ p_0) \vdash np \setminus s_i} [/E]$$

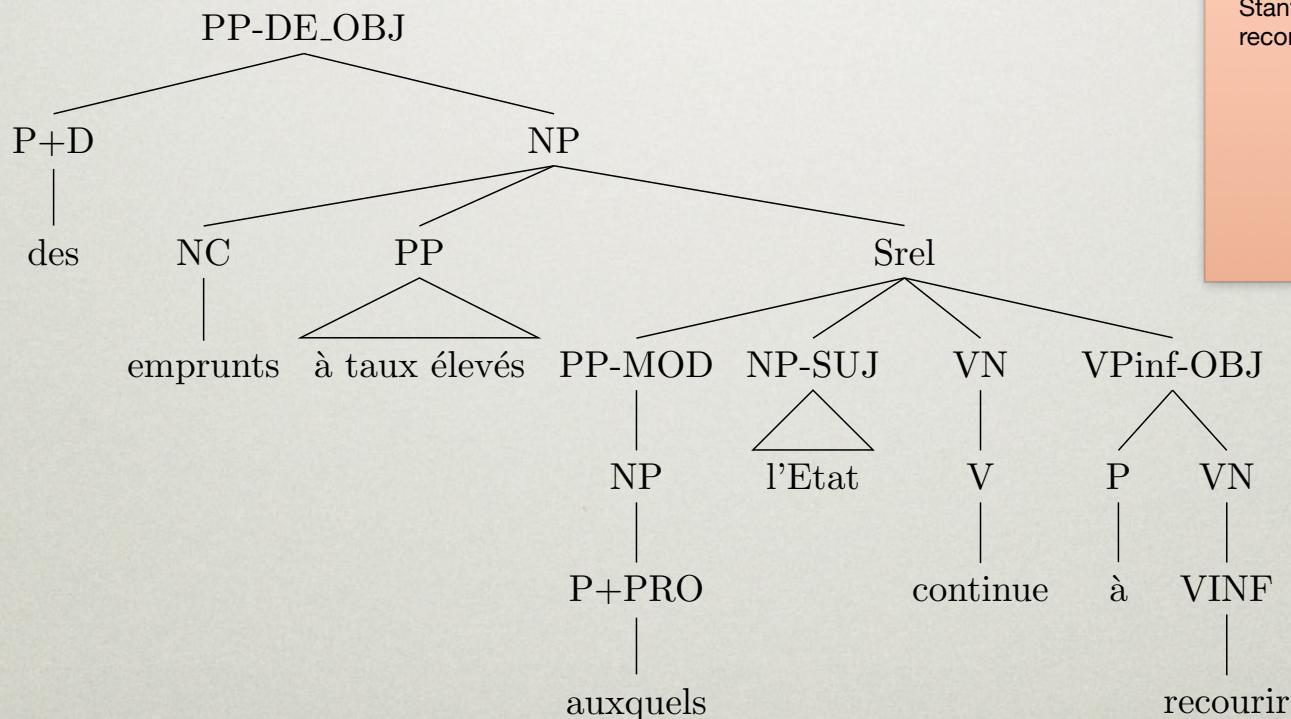


Applications

e_1	y_1
$y_1 :$	$e_2 \ e_3 \ x_3$ $x_3 = ?$ $\text{aider_à}(e_2, x_0, x_3, e_3)$ $\text{partir}(e_3, x_3)$ $\text{demander}(e_1, y_0, x_0, y_1)$

THE FRENCH TREEBANK

Example PP from the corpus (slightly simplified)

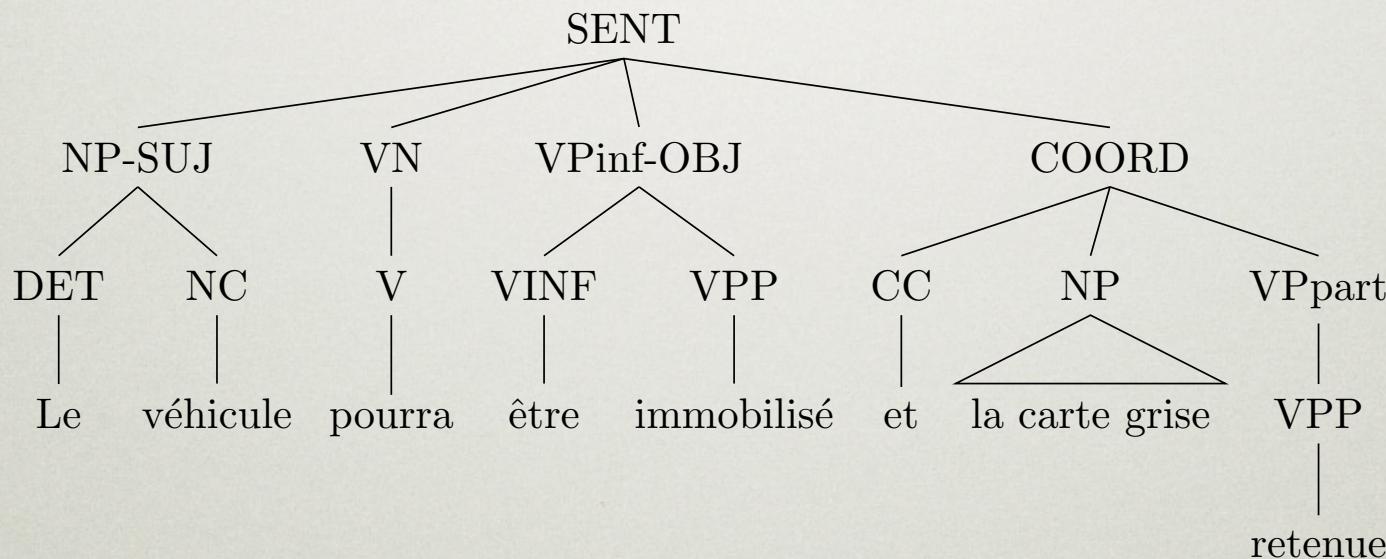


This is the FTB converted to Stanford TregEx (a tool I highly recommend)

≈ continue_to_resort_to(state,loans)

THE FRENCH TREEBANK

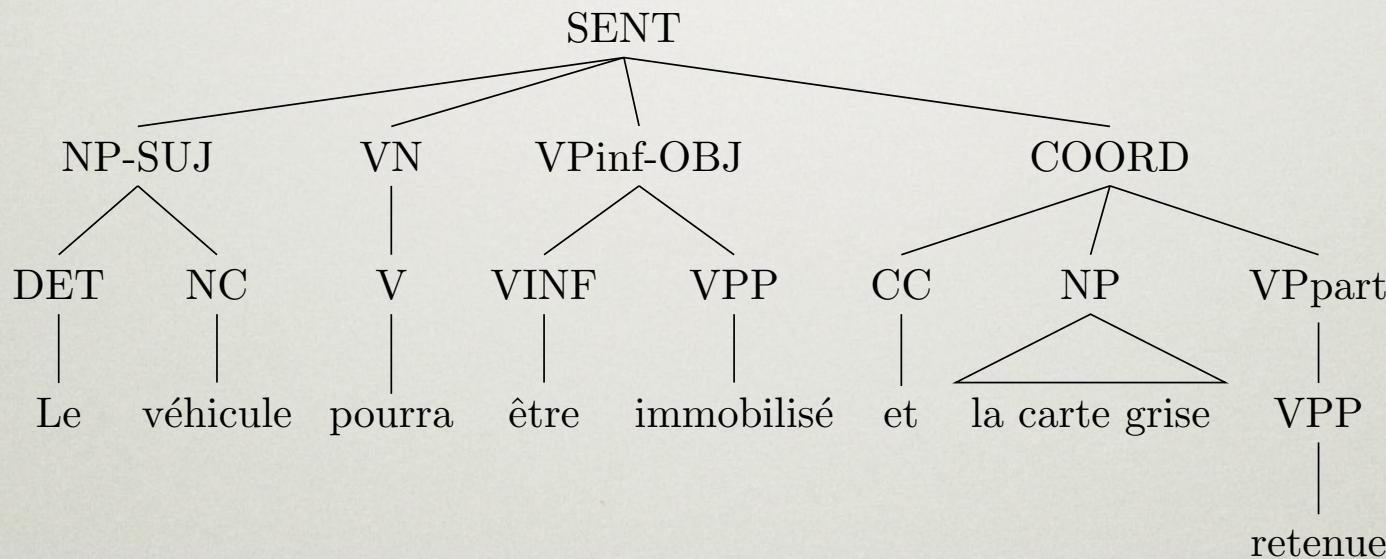
Example from the corpus (slightly simplified)



“The car could be immobilized and
the car registration (could be) taken”

THE FRENCH TREEBANK

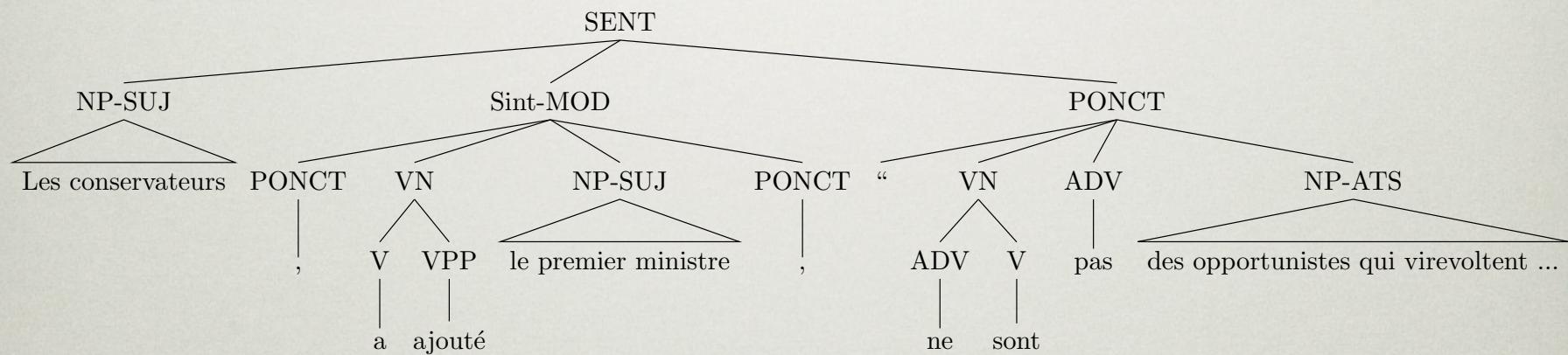
Example from the corpus (slightly simplified)



$\approx \exists x. \text{pouvoir}(\text{immobiliser}(x, \text{véhicule})) \wedge$
 $\exists y. \text{pouvoir}(\text{retenir}(y, \text{cart_grise}))$

THE FRENCH TREEBANK

Example from the corpus (slightly simplified)



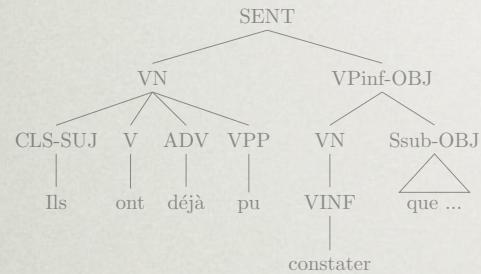
The conservative party, added the prime minister, “are not
opportunists who flip-flop...”

FTB: THE GOOD, THE BAD AND THE UGLY

- Context-free grammar
- Verb group as a constituent
- Treatment of coordination
- No annotation of passives, long-distance dependencies, ...

OUTLINE

French Treebank



Grammar Extraction

$$\frac{\frac{\frac{de}{(np \setminus s_{di}) / (np \setminus s_i)} [Lex]}{cl_r} \frac{s'}{cl_r} [Lex] \frac{\frac{attaque}{(cl_r \setminus (np \setminus s_i)) / pp_a} [Lex]}{a \circ p_0 \vdash cl_r \setminus (np \setminus s_i)} [Hyp]_1}{s' \circ (a \circ p_0) \vdash np \setminus s_i} [/E]$$



Applications

e_1	y_1
$y_1 :$	$e_2 \ e_3 \ x_3$ $x_3 = ?$ $aider_à(e_2, x_0, x_3, e_3)$ $partir(e_3, x_3)$ $demander(e_1, y_0, x_0, y_1)$

CATEGORIAL GRAMMARS

Formulas and corresponding expressions

- np
 - Jean, l'étudiant, ...
 - n
 - étudiant, économie, ...
 - s
 - Jean dort, Jean aime Marie
 - np \ s
 - dort, aime Marie
 - np / n
 - un, chaque, l'
 - (np \ s) / np
 - aime, étudie

CATEGORIAL GRAMMARS

Rules

Lambek categorial grammars have only four rules: an elimination and an introduction rule for both “\” and “/”

$$\frac{A/B \quad B}{A} [/ E] \quad \frac{B \quad B \setminus A}{A} [\setminus E]$$
$$\dots \quad [B]^i \quad [B]^i \quad \dots$$
$$\vdots \quad \vdots$$
$$\frac{A}{A/B} [/ I]^i \quad \frac{A}{B \setminus A} [\setminus I]^i$$

CATEGORIAL GRAMMARS

Example

un étudiant dort
np/n n np\s

$$\frac{A/B \quad B}{A} [/E] \quad \frac{B \quad B \setminus A}{A} [\setminus E]$$

$$\begin{matrix} & [B]^i & & [B]^i & \\ \dots & & & & \dots \\ & \vdots & & & \vdots \\ & & & & \end{matrix}$$

$$\frac{A}{A/B} [/I]^i \quad \frac{A}{B \setminus A} [\setminus I]^i$$

CATEGORIAL GRAMMARS

Example

un étudiant dort
 $\frac{\text{np/n} \quad \text{n}}{\text{np}}[/\text{E}]$ $\text{np} \setminus \text{s}$

$$\frac{\begin{array}{c} \text{A/B} \\ \text{B} \end{array} [/E]}{\text{A}} \quad \frac{\text{B}}{\text{B} \setminus \text{A}} [\backslash E]$$

$$\dots \quad [B]^i \quad [B]^i \quad \dots$$
$$\vdots \qquad \qquad \vdots$$

$$\frac{\text{A}}{\text{A/B}} [/I]^i \quad \frac{\text{A}}{\text{B} \setminus \text{A}} [\backslash I]^i$$

CATEGORIAL GRAMMARS

Example

un étudiant
 $\frac{\text{np/n} \quad n}{\text{np}} \quad \frac{\text{n}}{\text{np}} \quad \frac{\text{dort}}{\text{np}\backslash s}$

$\frac{\text{np}}{\text{s}}$

$$\frac{A/B \quad B}{A} [/E] \quad \frac{B \quad B/A}{A} [\backslash E]$$

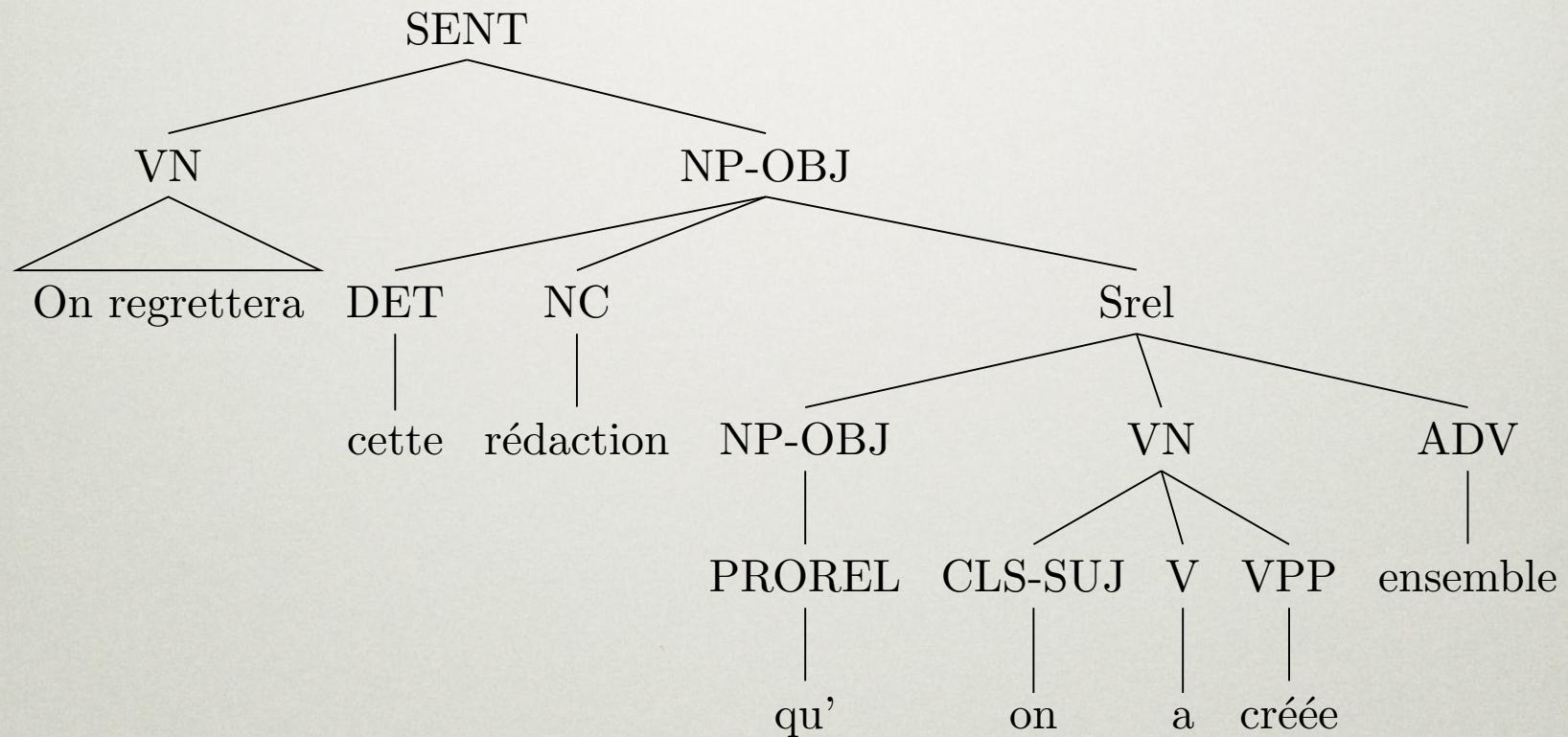
$$\dots \quad [B]^i \quad [B]^i \quad \dots$$
$$\vdots \qquad \qquad \vdots$$

$$\frac{A}{A/B} [/I]^i \quad \frac{A}{B/A} [\backslash I]^i$$

LAMBEK GRAMMARS AND BEYOND

- getting the semantics right requires a somewhat richer system than AB grammars
- introduction rules (“traces” or the original “slash categories” and their semantics)
- structural rules (“movement” or “head wrap”, essentially restricted tree rewrite operations)

INTRODUCTION RULES: EXAMPLE



INTRODUCTION RULES: EXAMPLE

redaction qu' on a créée
n $(n \setminus n) / (s / np)$ np $(np \setminus s) / (np \setminus s_{ppart})$ $(np \setminus s_{ppart}) / np$

INTRODUCTION RULES: EXAMPLE

redaction qu' on a créée
n (n\n)/(s\np) np (np\s)/(np\s_{ppart}) (np\s_{ppart})/np np

INTRODUCTION RULES: EXAMPLE

redaction qu' on a créée
n $(n \setminus n) / (s / np)$ np $(np \setminus s) / (np \setminus s_{ppart})$ $\frac{(np \setminus s_{ppart}) / np}{np \setminus s_{ppart}}$ np [/E]

INTRODUCTION RULES: EXAMPLE

redaction	qu'	on	créée
n	(n\ n)/(s/np)	np	a
			$\frac{(np \setminus s_{ppart})/np}{np \setminus s_{ppart}} [/ E]$
			$\frac{(np \setminus s)/(np \setminus s_{ppart})}{np \setminus s_{ppart}} [/ E]$

INTRODUCTION RULES: EXAMPLE

redaction qu'

$$n \quad (n \setminus n) / (s / np)$$

$$\frac{\frac{a}{(np \setminus s_{ppart}) / np} \quad np}{\frac{(np \setminus s) / (np \setminus s_{ppart}) \quad np \setminus s_{ppart}}{np \setminus s} \quad [\setminus E]} \quad [E]$$

INTRODUCTION RULES: EXAMPLE

redaction qu'

$$\frac{\frac{\frac{n \quad (n \setminus n) / (s / np)}{on \quad \frac{(np \setminus s) / (np \setminus s_{ppart})}{np}} a \quad \frac{(np \setminus s_{ppart}) / np}{np \setminus s_{ppart}} [np]^1_{[/E]} }{np \quad \frac{np \setminus s}{[\setminus E]}} }{s \quad \frac{s}{s / np} [/I]^1}$$

INTRODUCTION RULES: EXAMPLE

redaction

n

a

créée

$$\frac{(np \setminus s_{\text{ppart}}) / np}{np \setminus s_{\text{ppart}}} [np]^1 [/E]$$

np

$$\frac{np}{s} \backslash [E]$$

[/E]

qu'

S

$$\frac{s/np}{\sqrt{[E]}} \sim N(0, 1)$$

$$(n \setminus n) / (s / np)$$

n\n

INTRODUCTION RULES: EXAMPLE

			créée
	a	$(np \setminus s_{ppart}) / np$	$[np]^1$
on	$\frac{(np \setminus s) / (np \setminus s_{ppart})}{np \setminus s_{ppart}}$	$[/E]$	$[/E]$
np	$np \setminus s$	$[\setminus E]$	
qu'		s	$[/I]^1$
redaction	$\frac{(n \setminus n) / (s / np)}{n \setminus n}$	s / np	$[/E]$
n			
	n		

BEYOND AB GRAMMARS

INTRODUCTION RULES

qu'	on	a	créée	ensemble
(n\n)/(s\np)	np	(np\s)/(np\s _{ppart})	(np\s _{ppart})/np	(np\s)\(np\s)

LAMBDA CALCULUS AND PROOFS AS TERMS

- Proofs in categorial grammar correspond to lambda terms
- These lambda terms "forget" the directions of the implications.

$$\frac{t:A/B \quad u:B}{(t\ u):A} \qquad \frac{u:B \quad t:B\setminus A}{(t\ u):A}$$
$$\frac{[x:B] \quad [x:B]}{\vdots \quad \vdots}$$
$$\frac{t:A}{A/B:\lambda x.t} \qquad \frac{t:A}{B\setminus A:\lambda x.t}$$

COMPUTING SEMANTICS

... cet équilibre délicat que l'Iran et l'Arabie saoudite chercheront à sauvegarder à Genève.

$$\frac{\dots [x:B] \dots \\ \vdots \\ t:A}{A/\Diamond\square B:\lambda x.t}$$

l'Iran et
l'Arabie saoudite $w_1:(n\backslash n)/(s/\Diamond\square np)$

que
 $w_2:np$

chercheront à $w_3:(np\backslash s)/(np\backslash s_{inf})$

sauvegarder

$w_4:(np\backslash s_{inf})/np$

COMPUTING SEMANTICS

... cet équilibre délicat que l'Iran et l'Arabie saoudite chercheront à sauvegarder à Genève.

$$\frac{\vdots}{A/\Diamond\Box B:\lambda x.t}$$

l'Iran et l'Arabie saoudite w3:(np\s)/(np\sinf)
w2:np
que
w1:(n\n)/(s/◊□np)

COMPUTING SEMANTICS

... cet équilibre délicat que l'Iran et l'Arabie saoudite chercheront à sauvegarder ε à Genève.

$$\frac{\dots [x:B] \dots}{A/\Diamond\square B:\lambda x.t}$$

l'Iran et
l'Arabie saoudite

que
 $w_1:(n\backslash n)/(s/\Diamond\square np)$

$$\frac{\text{chercheront à } w_3:(np\backslash s)/(np\backslash s_{inf}) \quad (w_4 x):np\backslash s_{inf}}{\text{sauvegarder } \epsilon \quad w_4:(np\backslash s_{inf})/np \quad x:np}$$

COMPUTING SEMANTICS

... cet équilibre délicat que l'Iran et l'Arabie saoudite chercheront à sauvegarder ε à Genève.

$$\frac{\dots [x:B] \dots}{\frac{\vdots}{\frac{t:A}{A/\Diamond\square B:\lambda x.t}}}$$

$$\frac{\text{l'Iran et } \text{l'Arabie saoudite}}{\frac{\text{que} \quad w_1:(n\backslash n)/(s/\Diamond\square np)}{\frac{w_2:np}{\frac{chercheront \ à \quad w_3:(np\backslash s)/(np\backslash s_{inf})}{\frac{(w_4 x):np\backslash s_{inf}}{(w_3(w_4 x)):np\backslash s}}}}}$$

sauvegarder ε

w₄:(np\ s_{inf})/np x:np

chercheront à

w₃:(np\ s)/(np\ s_{inf}) (w₄ x):np\ s_{inf}

(w₃(w₄ x)):np\ s

que
w₁:(n\ n)/(s/\Diamond\square np)

COMPUTING SEMANTICS

... cet équilibre délicat que l'Iran et l'Arabie saoudite chercheront à sauvegarder ε à Genève.

$$\frac{\dots [x:B] \dots}{\frac{\vdots}{\frac{t:A}{A/\Diamond\square B:\lambda x.t}}}$$

$$\frac{\text{l'Iran et} \quad \text{sauvegarder} \quad \varepsilon}{\frac{\text{l'Arabie saoudite} \quad \frac{\text{chercheront à} \quad w_4:(np \setminus s_{inf})/np \quad x:np}{\frac{(w_4 x):np \setminus s_{inf}}{w_2:np \quad ((w_3(w_4 x)):np \setminus s)}}}{que \quad ((w_3(w_4 x)) \ w_2):s}}$$
$$w_1:(n \setminus n)/(s/\Diamond\square np)$$

COMPUTING SEMANTICS

... cet équilibre délicat que l'Iran et l'Arabie saoudite chercheront à sauvegarder ε à Genève.

$$\frac{\dots [x:B] \dots}{\frac{\vdots}{\frac{t:A}{A/\Diamond\square B:\lambda x.t}}}$$

$$\frac{\text{l'Iran et} \quad \text{sauvegarder} \quad \varepsilon \\ \text{l'Arabie saoudite} \quad \frac{\text{chercheront à} \quad w_4:(np \setminus s_{inf})/np \quad x:np}{\frac{w_3:(np \setminus s)/(np \setminus s_{inf}) \quad (w_4 x):np \setminus s_{inf}}{\frac{w_2:np \quad ((w_3(w_4 x)):np \setminus s)}{\frac{\text{que} \quad ((w_3(w_4 x)) w_2):s}{\frac{w_1:(n \setminus n)/(s/\Diamond\square np) \quad \lambda x.((w_3(w_4 x)) w_2):s/\Diamond\square np}}}}}}{}}$$

COMPUTING SEMANTICS

... cet équilibre délicat que l'Iran et l'Arabie saoudite chercheront à sauvegarder ε à Genève.

$$\frac{\dots [x:B] \dots}{\frac{\vdots}{\frac{t:A}{A/\Diamond\square B:\lambda x.t}}}$$

$$\frac{\frac{\frac{\frac{\frac{\text{l'Iran et}}{\text{l'Arabie saoudite}} \frac{\frac{w_3:(np\s)/((np\s_{inf})/np \ x:np)}{(w_4 x):np\s_{inf}}}{w_4:(np\s_{inf})/np \ x:np}}{(w_4 x):np\s_{inf}}}{(w_3(w_4 x)):np\s}}{((w_3(w_4 x)) \ w_2):s}}{\frac{w_1:(n\backslash n)/(s/\Diamond\square np)}{\frac{\lambda x.((w_3(w_4 x)) \ w_2):s/\Diamond\square np}{(w_1 \ \lambda x.((w_3(w_4 x)) \ w_2)):n\backslash n}}}}$$

COMPUTING SEMANTICS

w₁: $\lambda P \lambda Q \lambda x (Px) \wedge (Qx)$

w₂: $I \wp A s$

w₃:*chercher_à*

w₄:*sauvegarder*

sauvegarder ε

chercheront à $\frac{w_4:(np \setminus s_{inf})/np \quad x:np}{(w_4 x):np \setminus s_{inf}}$

l'Iran et
l'Arabie saoudite $\frac{w_3:(np \setminus s)/(np \setminus s_{inf}) \quad (w_4 x):np \setminus s_{inf}}{w_2:np \quad (w_3(w_4 x)):np \setminus s}$

que $\frac{((w_3(w_4 x)) \ w_2):s}{(w_3(w_4 x)) \ w_2):s}$

$w_1:(n \setminus n)/(s/\diamond \square np) \quad \lambda x.((w_3(w_4 x)) \ w_2):s/\diamond \square np$

$(w_1 \ \lambda x.((w_3(w_4 x)) \ w_2)):n \setminus n$

COMPUTING SEMANTICS

w₁: $\lambda P \lambda Q \lambda y (P_y) \wedge (Q_y)$

w₂: $I \wp A s$

w₃:*chercher_à*

w₄:*sauvegarder*

(w₁ $\lambda x.((w_3(w_4 x)) w_2))$



(($\lambda P \lambda Q \lambda y (P_y) \wedge (Q_y)$)

$\lambda x.((chercher_à(sauvegarder x)) I \wp A s))$



($\lambda Q \lambda y (\lambda x.((chercher_à(sauvegarder x)) I \wp A s) y) \wedge (Q_y))$

COMPUTING SEMANTICS

w₁: $\lambda P \lambda Q \lambda y (P_y) \wedge (Q_y)$

w₂: $I \wp A s$

w₃:*chercher_à*

w₄:*sauvegarder*

(w₁ $\lambda x.((w_3(w_4 x)) w_2))$



(($\lambda P \lambda Q \lambda y (P_y) \wedge (Q_y)$)

$\lambda x.((chercher_à(sauvegarder x)) I \wp A s))$



($\lambda Q \lambda y (\lambda x.((chercher_à(sauvegarder x)) I \wp A s) y) \wedge (Q_y))$



($\lambda Q \lambda y ((chercher_à(sauvegarder_y)) I \wp A s) \wedge (Q_y))$



$\lambda x.équilibre(x) \wedge délicat(x) \wedge$
chercher_à(I $\wp A s, (sauvegarder x))$

A TALE OF TWO
CATEGORIAL
GRAMMARS

LOGIC AND COMBINATORS

combinatorial logic

combinators

logic

SKI, BCK, BCI



IIL, IAL, ILL

LOGICS AND COMBINATORS

combinatorial logic

combinators

logic

SKI, BCK, BCI



IIL, IAL, ILL

CCG

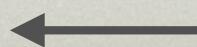
combinators (?)



logic???

Lambek / TLG

combinators???



logic

CCG AND TYPE-LOGICAL GRAMMAR

Both complicate locally and simplify globally
but with a different notion of locality

	CCG	TLG
Sequence	Local	Non-local
Connective	Non-local	Local

CCG: each category can be traced back to a continuous sequence of the input categories

TLG: a formula can correspond to a discontinuous sequence of lexical formulas

CCG: Rules can refer to multiple connectives
TLG: logical rules operate on one connective at a time (subformula property, independence of

CROSSED COMPOSITION IN CCG

qu'	on	a	créée	ensemble
$(n \setminus n) / (s / np)$	np	$(np \setminus s) / (np \setminus s_{ppart})$	$(np \setminus s_{ppart}) / np$	$(np \setminus s) \setminus (np \setminus s)$

CROSSED COMPOSITION IN CCG

$$\frac{A}{B/(A \setminus B)} T \quad \frac{A/B \quad B/C}{A/C} B \quad \frac{B/C \quad B \setminus A}{A/C} Bx$$

qu'

$$\begin{array}{c} a \\ \hline (np \setminus s)/(np \setminus s_{ppart}) & (np \setminus s_{ppart})/np \\ \hline np & (np \setminus s)/np \\ \hline s/(np \setminus s) & (np \setminus s)/np \\ \hline & s/np \end{array}$$

créée

ensemble

$$\frac{(np \setminus s)/np}{(np \setminus s) \setminus (np \setminus s)}$$

CROSSED COMPOSITION IN CCG

$$\frac{A}{B/(A \setminus B)} T \quad \frac{A/B \quad B/C}{A/C} B \quad \frac{B/C \quad B \setminus A}{A/C} Bx$$

qu'

$$\begin{array}{c} a \\ \hline (n \setminus n) / (s / np) \quad \text{on} \quad \frac{(np \setminus s) / (np \setminus s_{\text{ppart}}) \quad (np \setminus s_{\text{ppart}}) / np}{(np \setminus s) / np} \quad \text{créée ensemble} \\ \hline \frac{np}{s / (np \setminus s)} \quad \hline (np \setminus s) \setminus (np \setminus s) \\ \hline s / np \end{array}$$

CROSSED COMPOSITION IN CCG

$$\frac{A}{B/(A \setminus B)} T \quad \frac{A/B \quad B/C}{A/C} B \quad \frac{B/C \quad B \setminus A}{A/C} Bx$$

qu'
 $(n \setminus n)/(s/np)$ a créée
on $\frac{(np \setminus s)/(np \setminus s_{ppart}) \quad (np \setminus s_{ppart})/np}{(np \setminus s)/np}$ ensemble
np $\frac{(np \setminus s)/np}{(np \setminus s) \setminus (np \setminus s)}$
 $\underline{s/(np \setminus s)}$ $\underline{(np \setminus s)/np}$
 $\underline{\underline{s/np}}$

CROSSED COMPOSITION IN CCG

$$\frac{A}{B/(A \setminus B)} T \quad \frac{A/B \quad B/C}{A/C} B \quad \frac{B/C \quad B \setminus A}{A/C} Bx$$

qu'
 $(n \setminus n)/(s/np)$ a créée
on $\frac{(np \setminus s)/(np \setminus s_{ppart}) \quad (np \setminus s_{ppart})/np}{(np \setminus s)/np}$ ensemble
np $\frac{(np \setminus s)/np}{(np \setminus s) \setminus (np \setminus s)}$
 $\underline{s/(np \setminus s)}$ $\underline{(np \setminus s)/np}$
 $\underline{\underline{s/np}}$

CROSSED COMPOSITION IN CCG

$$\frac{A}{B/(A \setminus B)} T \quad \frac{A/B \quad B/C}{A/C} B \quad \frac{B/\textcolor{red}{C} \quad B \setminus \textcolor{blue}{A}}{A/\textcolor{green}{C}} Bx$$

qu'

$$\frac{(n \setminus n)/(s/np) \quad \text{on} \quad \frac{a}{\frac{(np \setminus s)/(np \setminus s_{ppart}) \quad (np \setminus s_{ppart})/np}{(np \setminus s)/np}} \quad \text{créée}}{\frac{np}{\frac{s/(np \setminus s)}{(np \setminus s) \setminus (np \setminus s)}}} \quad \text{ensemble}$$

s/np

CROSSED COMPOSITION IN CCG

$$\frac{A}{B/(A \setminus B)}^T$$

harmless

CROSSED COMPOSITION IN CCG

$$\frac{A}{B/(A \setminus B)} T \quad \frac{A/B \quad B/C}{A/C} B$$

harmless

mostly
harmless

CROSSED COMPOSITION IN CCG

$$\frac{A}{B/(A \setminus B)} T$$

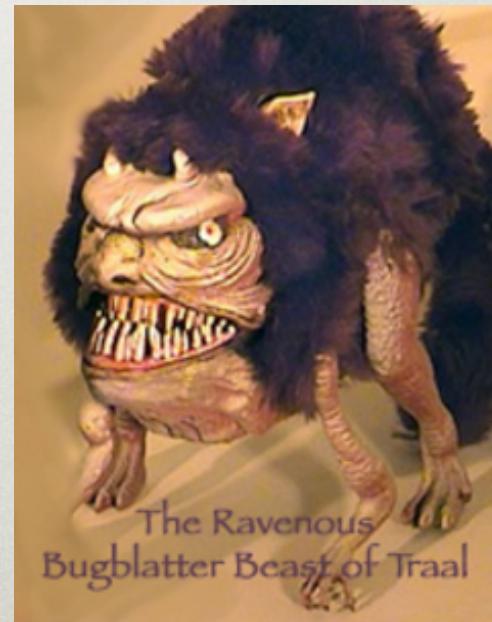
harmless

$$\frac{A/B \quad B/C}{A/C} B$$

mostly
harmless

$$\frac{B/C \quad B \setminus A}{A/C} Bx$$

danger!



CROSSED COMPOSITION IN CCG

$$\frac{A}{B/(A \setminus B)} T$$

harmless

$$\frac{A/B \quad B/C}{A/C} B$$

mostly
harmless

$$\frac{B/C \quad B \setminus A}{A/C} Bx$$

danger!

$$\frac{A \quad [A \setminus B]^i}{B} \frac{\backslash E}{/I_i}$$
$$\frac{}{B/(A \setminus B)}$$

CROSSED COMPOSITION IN CCG

$$\frac{A}{B/(A \setminus B)} T$$

$$\frac{A/B \quad B/C}{A/C} B$$

$$\frac{B/C \quad B \setminus A}{A/C} Bx$$

harmless

mostly
harmless

$$\frac{A \quad [A \setminus B]^i}{B} \frac{\backslash E}{/I_i} \frac{}{B/(A \setminus B)}$$

$$\frac{A/B}{\frac{A}{A/C}} \frac{\frac{B/C \quad [C]^i}{B} /E}{/E} /I^i$$

CROSSED COMPOSITION IN CCG

$$\frac{A}{B/(A \setminus B)} T$$

$$\frac{A/B \quad B/C}{A/C} B$$

$$\frac{B/C \quad B \setminus A}{A/C} Bx$$

harmless

mostly
harmless

$$\frac{x : A \quad [f : A \setminus B]^i}{(f\ x) : B} \setminus^E / I_i \quad \frac{\lambda f.(f\ x) : B/(A \setminus B)}{\lambda f.(f\ x) : B/(A \setminus B)}$$

$$\frac{f : A/B \quad \frac{g : B/C \quad [x : C]^i}{(g\ x) : B} / E}{(f\ (g\ x)) : A} / E \quad \frac{(f\ (g\ x)) : A}{\lambda x.(f\ (g\ x)) : A/C} / I^i$$

CROSSED COMPOSITION

$$\frac{\frac{B/C \quad [C]^i}{B} \ /E \quad B\backslash A}{\frac{A}{A/C} \ ?I_i} \ \backslash E$$

CROSSED COMPOSITION

$$\frac{\frac{B/C \quad [C]^i}{B} \ /E \quad B\backslash A}{\frac{A}{A/C} \ /I_i} \ \backslash E$$

CROSSED COMPOSITION

$$\frac{\frac{B/C \quad [C]^i}{B} \ /_E \ B \setminus A}{\frac{A}{A/\diamond \square C} \ /_{\diamond \square_i} \ \backslash E}$$

CROSSED COMPOSITION

$$\frac{\frac{g : B/C \quad [x : C]^i}{(g\ x) : B} \ /_E \ f : B \setminus A}{(f\ (g\ x)) : A} \ \backslash_E \ ?I_i}{\lambda x.(f\ (g\ x)) : A/C}$$

HOW TO EXTEND OR MODIFY A LOGIC

- Suppose we are unhappy with our logic and we want to change it, then what is our job?

1. reprove formal results (normalisation, subformula property)
2. verify we keep what worked and can do some things which didn't work before (possibly by direct translation, but this is not always possible)

BEYOND LAMBEK GRAMMARS

NEW INTRODUCTION RULES

.... [B]ⁱ

$$\frac{A}{A/\diamond \Box B} [/I]^i$$

qu'	on	a	créée	ensemble
(n\n)/(s/\diamond \Box np)	np	(np\s)/(np\s _{part})	(np\s _{part})/np	(np\s)\(np\s)

BEYOND LAMBEK GRAMMARS

NEW INTRODUCTION RULES

qu'	on	a	créée	ensemble
$(n \setminus n) / (s / \diamond \Box np)$	np	$(np \setminus s) / (np \setminus s_{ppart})$	$(np \setminus s_{ppart}) / np$	$(np \setminus s) \setminus (np \setminus s)$

BEYOND LAMBEK GRAMMARS

NEW INTRODUCTION RULES

qu'	on	a	créée	ensemble
$(n \setminus n) / (s / \diamond \Box np)$	np	$(np \setminus s) / (np \setminus s_{ppart})$	$(np \setminus s_{ppart}) / np$	np $(np \setminus s) \setminus (np \setminus s)$

BEYOND LAMBEK GRAMMARS

NEW INTRODUCTION RULES

qu'	on	a	créée	ensemble
$(n \setminus n) / (s / \diamond \Box np)$	np	$(np \setminus s) / (np \setminus s_{ppart})$	$(np \setminus s_{ppart}) / np$	np $(np \setminus s) \setminus (np \setminus s)$
<hr/>				$np \setminus s_{ppart}$

BEYOND LAMBEK GRAMMARS

NEW INTRODUCTION RULES

$$\begin{array}{ccccc} \text{qu'} & \text{on} & & \text{créée} & \text{ensemble} \\ (\mathbf{n}\backslash\mathbf{n})/(\mathbf{s}/\diamond\Box\mathbf{np}) & \mathbf{np} & \text{a} & (\mathbf{np}\backslash\mathbf{s}_{\text{ppart}})/\mathbf{np} & \mathbf{np} \quad (\mathbf{np}\backslash\mathbf{s})\backslash(\mathbf{np}\backslash\mathbf{s}) \\ & & & \hline & \\ & & (\mathbf{np}\backslash\mathbf{s})/(\mathbf{np}\backslash\mathbf{s}_{\text{ppart}}) & \mathbf{np}\backslash\mathbf{s}_{\text{ppart}} & \\ & & & \hline & \\ & & & \mathbf{np}\backslash\mathbf{s} & \end{array}$$

BEYOND LAMBEK GRAMMARS

NEW INTRODUCTION RULES

$$\begin{array}{ccccc} \text{qu'} & & \text{on} & & \text{créée} \\ (\text{n}\backslash\text{n})/(\text{s}/\diamond\Box\text{np}) & \text{np} & a & (\text{np}\backslash\text{s}_{\text{ppart}})/\text{np} & \text{np} \\ & & & \hline & (\text{np}\backslash\text{s})/(\text{np}\backslash\text{s}_{\text{ppart}}) & \text{np}\backslash\text{s}_{\text{ppart}} & & \text{ensemble} \\ & & & \hline & \text{np}\backslash\text{s} & & (\text{np}\backslash\text{s})\backslash(\text{np}\backslash\text{s}) & \\ & & & \hline & & & \text{np}\backslash\text{s} & \end{array}$$

BEYOND LAMBEK GRAMMARS

NEW INTRODUCTION RULES

$$\frac{\frac{\frac{qu'}{(n \setminus n) / (s / \diamond \square np)} a \frac{(np \setminus s_{\text{ppart}}) / np \quad np}{(np \setminus s) / (np \setminus s_{\text{ppart}}) \quad np \setminus s_{\text{ppart}}} \text{ créée}}{(np \setminus s) \setminus (np \setminus s)} \text{ ensemble}}{np} \text{ on} \quad \frac{np \setminus s}{s}}$$

BEYOND LAMBEK GRAMMARS

NEW INTRODUCTION RULES

$$\frac{\frac{\frac{\frac{qu'}{(n \setminus n) / (s / \diamond \Box np)} a \frac{(np \setminus s_{\text{ppart}}) / (np \setminus s_{\text{ppart}})}{(np \setminus s) / (np \setminus s_{\text{ppart}})} crée}{(np \setminus s_{\text{ppart}}) / np} [np]^1}{np \setminus s_{\text{ppart}}} ensemble}{np \setminus s} (np \setminus s) \setminus (np \setminus s)}$$
$$\frac{on}{np} \frac{s}{s / \diamond \Box np} [/I]^1$$

BEYOND LAMBEK GRAMMARS

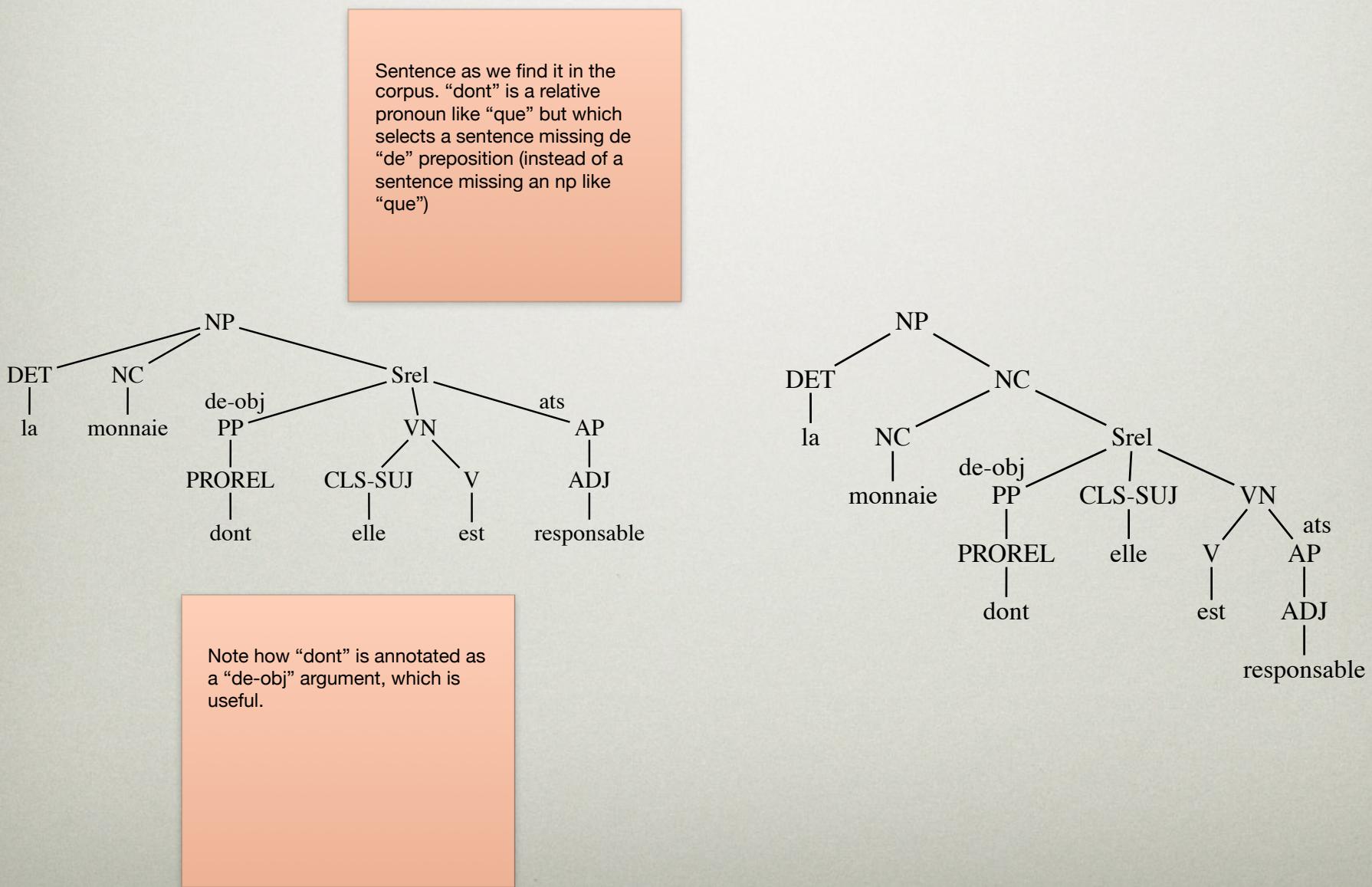
NEW INTRODUCTION RULES

$$\frac{\frac{\frac{\frac{\frac{\frac{créée}{(np \setminus s_{\text{ppart}}) / np} [np]^1}{(np \setminus s) / (np \setminus s_{\text{ppart}})} np \setminus s_{\text{ppart}} ensemble}{on}{np \setminus s}{(np \setminus s) \setminus (np \setminus s)}}{np}{qu'}}{(n \setminus n) / (s / \diamond \Box np)} s / \diamond \Box np [/I]^1}{n \setminus n}$$

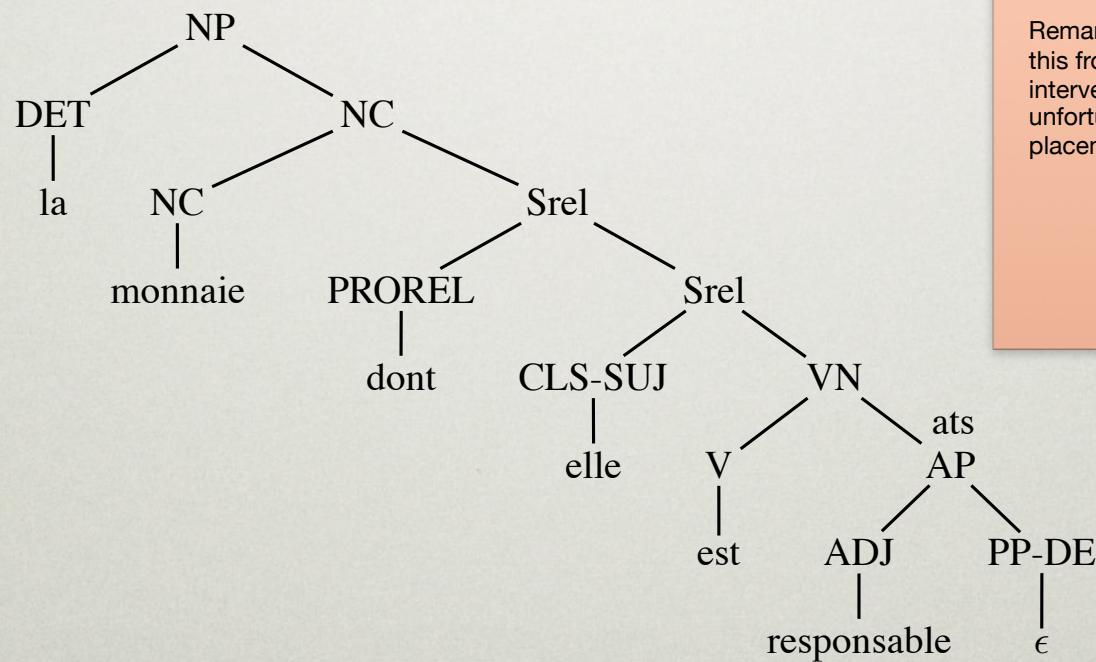
TREEBANK EXTRACTION

- A rather classical approach. The extraction algorithm is parametric wrt three functions:
 1. a function identifying heads / functors
 2. a function identifying modifiers
 3. a function given a formula for each syntactic category of the annotation

TREEBANK EXTRACTION



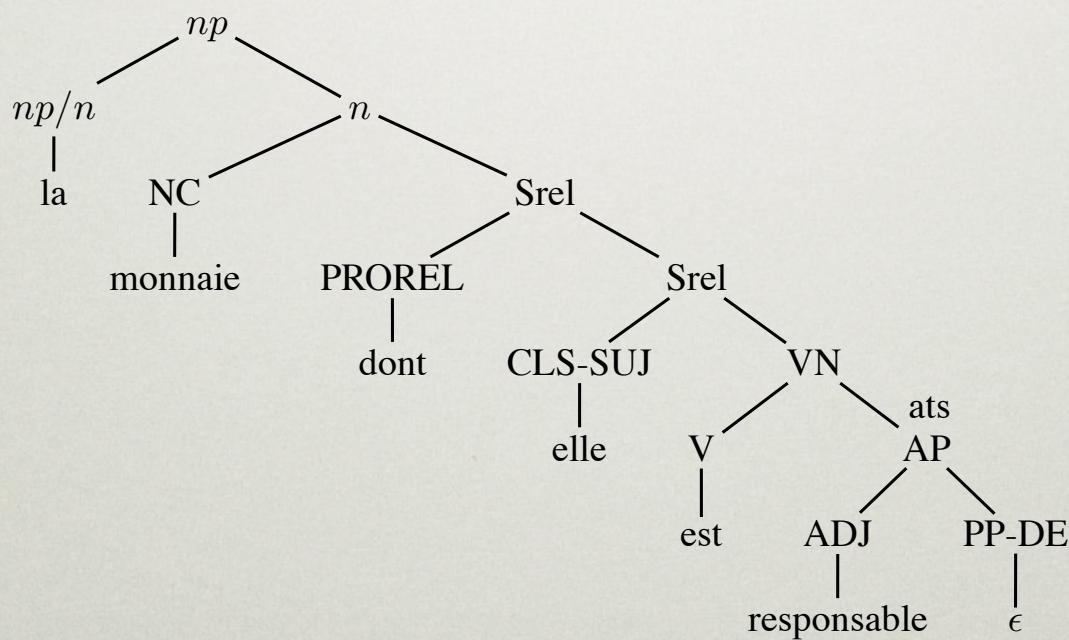
TREEBANK EXTRACTION



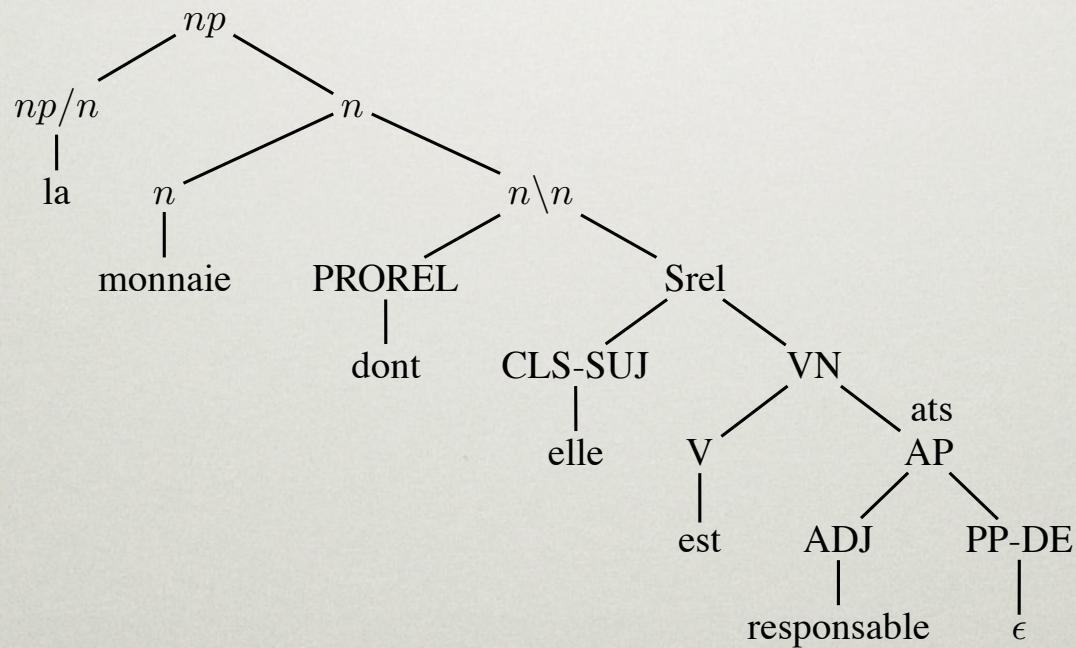
However, the “de” preposition belongs to “responsable” (some adjectives select for prepositions: “responsable de X” functions as an adjective just as “responsable”)

Remark however, that there is no way to derive this from the annotation as it is given. Manual intervention (or at least verification!) is unfortunately necessary to assure the correct placement of the hypothetical preposition.

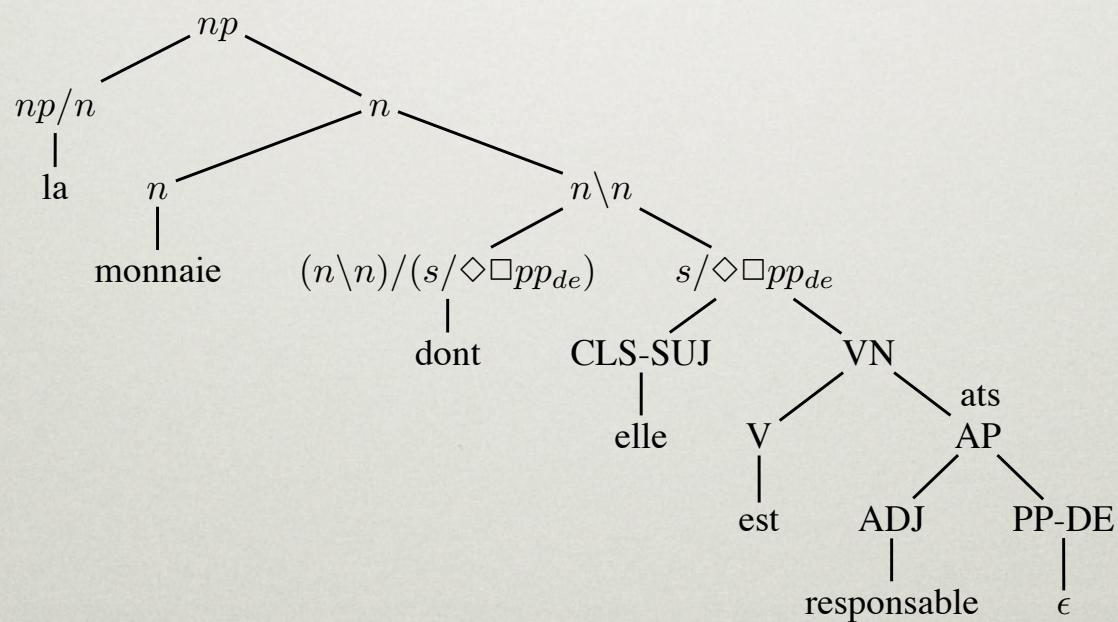
TREEBANK EXTRACTION



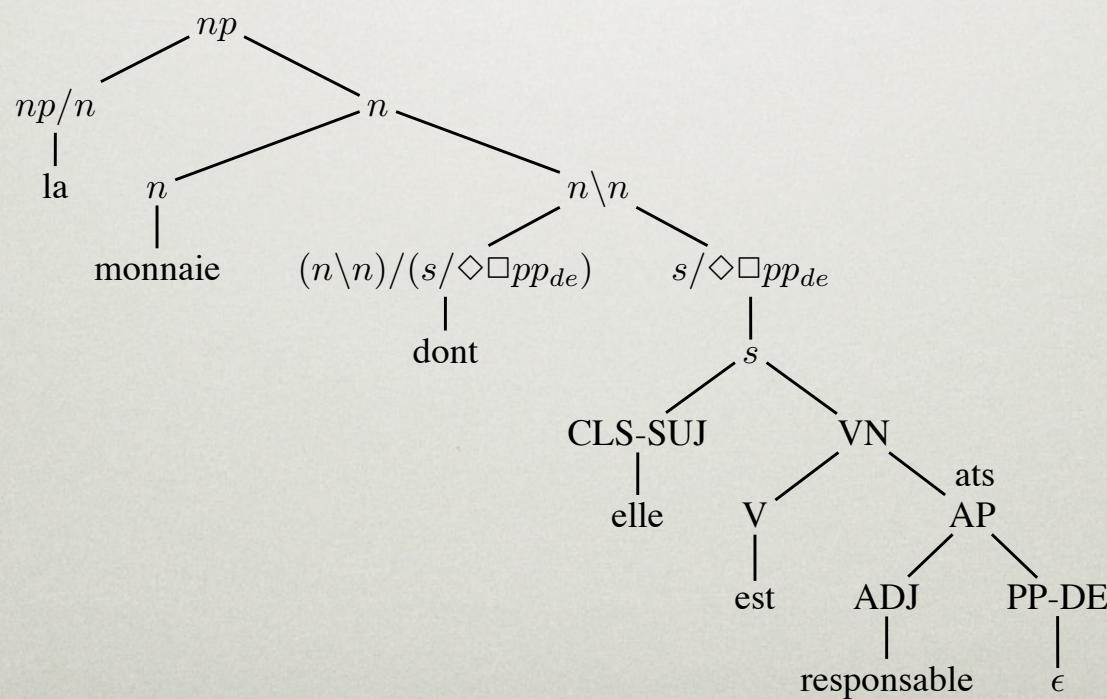
TREEBANK EXTRACTION



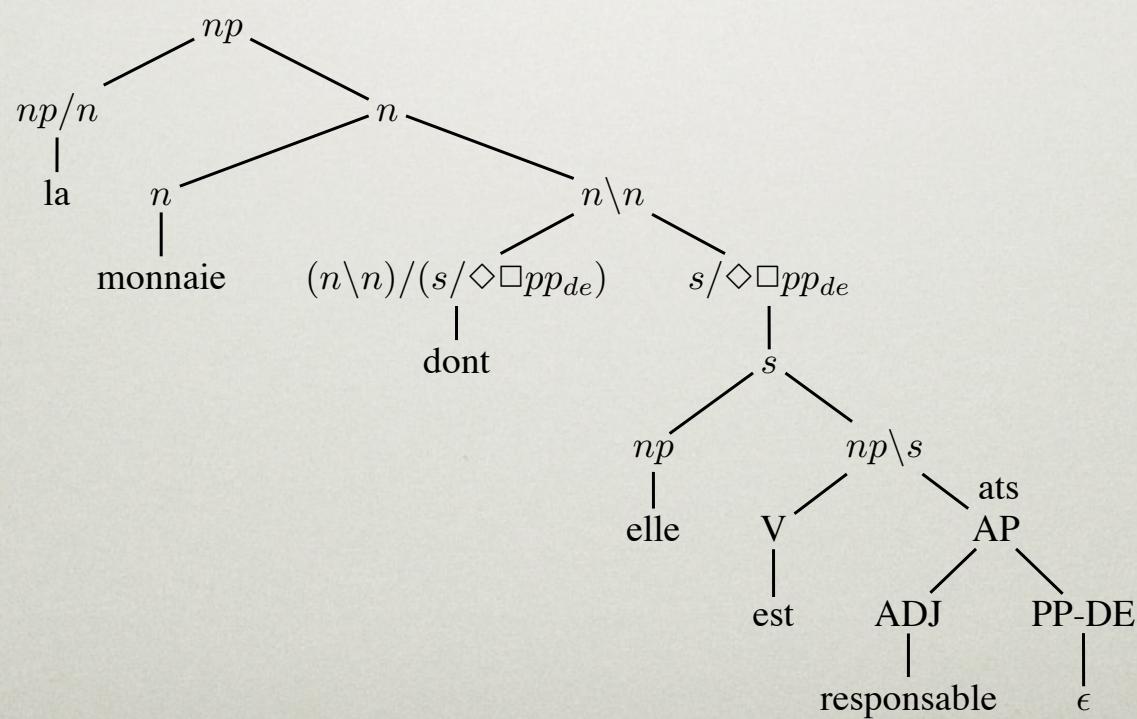
TREEBANK EXTRACTION



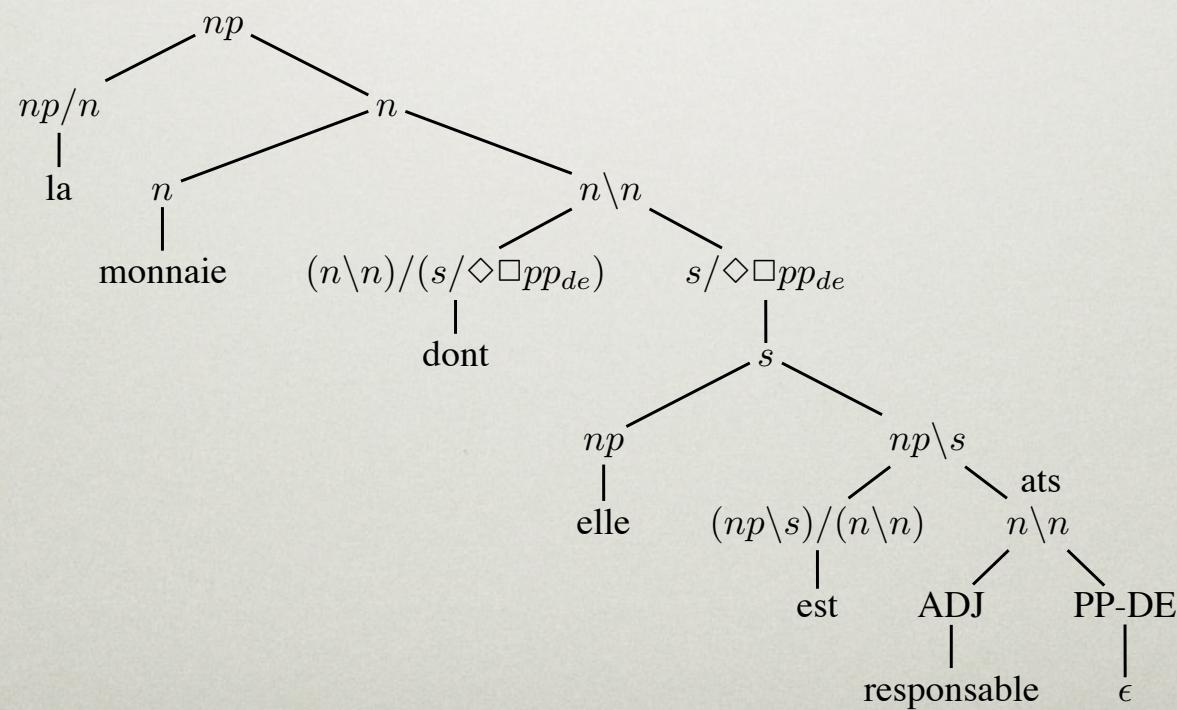
TREEBANK EXTRACTION



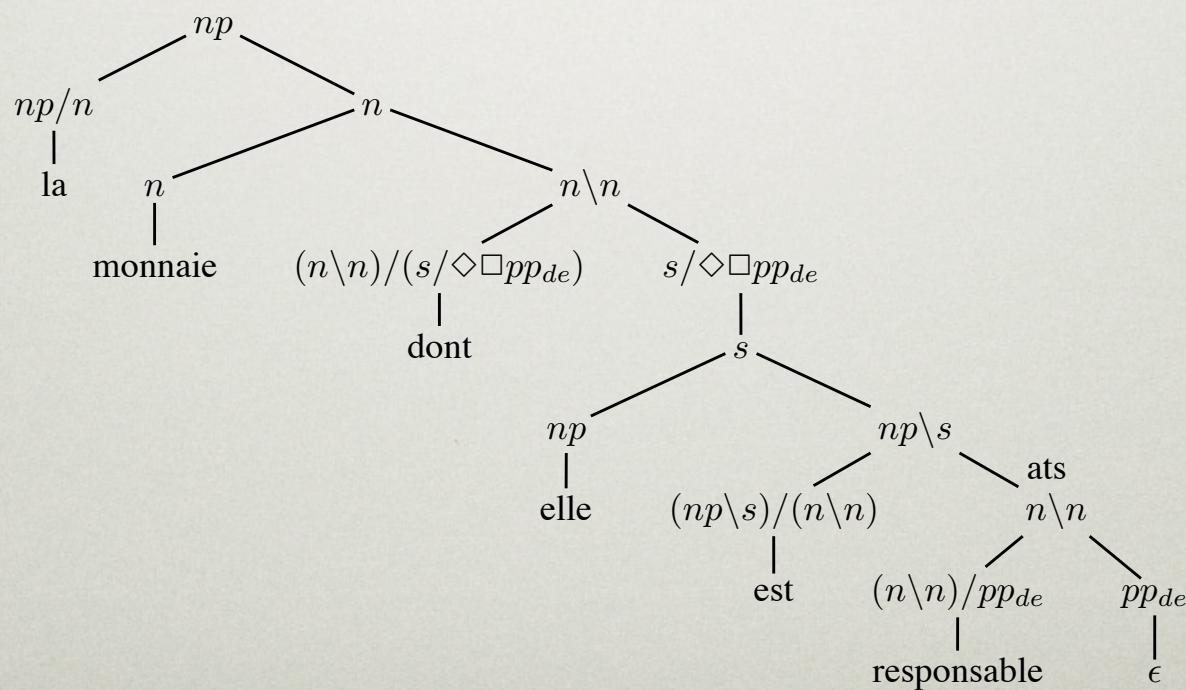
TREEBANK EXTRACTION



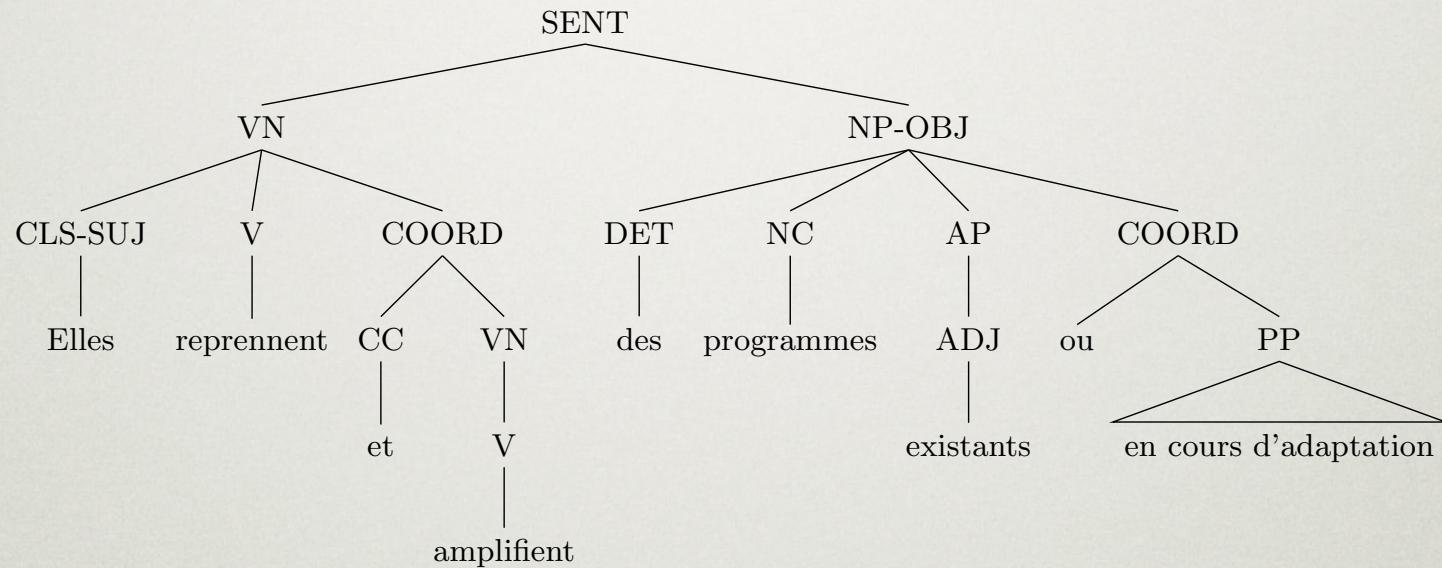
TREEBANK EXTRACTION



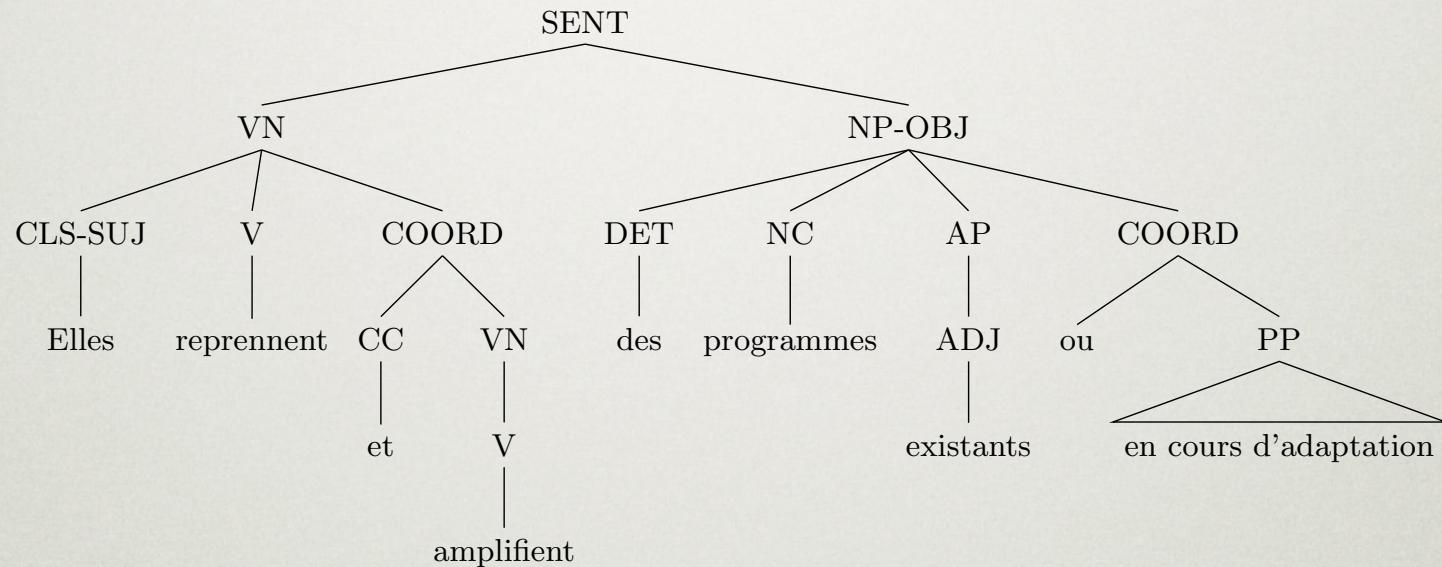
TREEBANK EXTRACTION



TREEBANK EXTRACTION: COORDINATION



TREEBANK EXTRACTION: COORDINATION



$$\frac{\frac{existant}{n \setminus n} \ Lex \quad \frac{ou}{((n \setminus n) \setminus (n \setminus n)) / (n \setminus n)} \ Lex \quad \frac{en \ cours...}{n \setminus n} \ Lex}{\frac{(n \setminus n) \setminus (n \setminus n)}{n \setminus n} \setminus E} \setminus E$$

STATISTICS ABOUT THE EXTRACTED FRENCH TREEBANK

- 15,590 sentences 445,918 words
- 43,098 distinct lexical entries
- 859 different formulas
- By comparison: 12,617 CFG rules

Note: these are the statistics after considerable cleanup: the first version had over 4,000 different formulas!

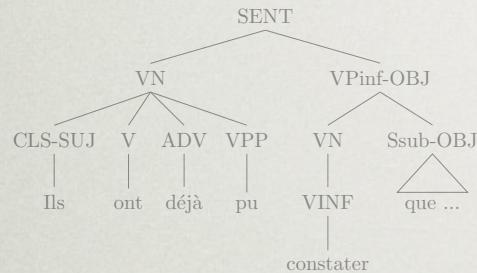
HOW TO IMPROVE AN EXTRACTED TREEBANK

Identify problem and look for clues to the underlying cause:

- annotation error
- need more sophisticated extraction
- need to add information by hand

OUTLINE

French Treebank



Grammar Extraction

$$\frac{\frac{de}{(np \setminus s_{di}) / (np \setminus s_i)} [Lex] \quad \frac{\frac{s'}{cl_r} [Lex] \quad \frac{\frac{attaquer}{(cl_r \setminus (np \setminus s_i)) / pp_a} [Lex] \quad \frac{p_0 \vdash pp_a}{a \circ p_0 \vdash cl_r \setminus (np \setminus s_i)} [Hyp]_1}{a \circ p_0 \vdash np \setminus s_i} [/E]}{s' \circ (a \circ p_0) \vdash np \setminus s_i} [/E]$$



Applications

e_1	y_1
$y_1 :$	$\begin{array}{ c }\hline e_2 & e_3 & x_3 \\ \hline x_3 = ? \\ \hline \text{aider_à}(e_2, x_0, x_3, e_3) \\ \hline \text{partir}(e_3, x_3) \\ \hline \text{demander}(e_1, y_0, x_0, y_1) \\ \hline \end{array}$

APPLICATIONS

- Wide-coverage parsing for French
- Wide-coverage semantics for French

WIDE-COVERAGE PARSING

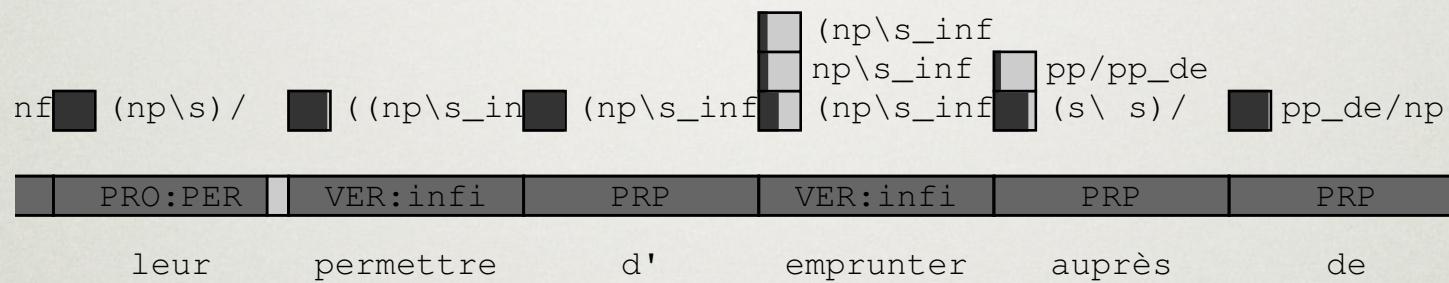
- How can we parse very big grammars efficiently?
- Bottlenecks: lexicon size, grammatical combinatorics

LEXICON SIZE

- Many frequent words occur with very many different formulas
- Classic solution: supertagging

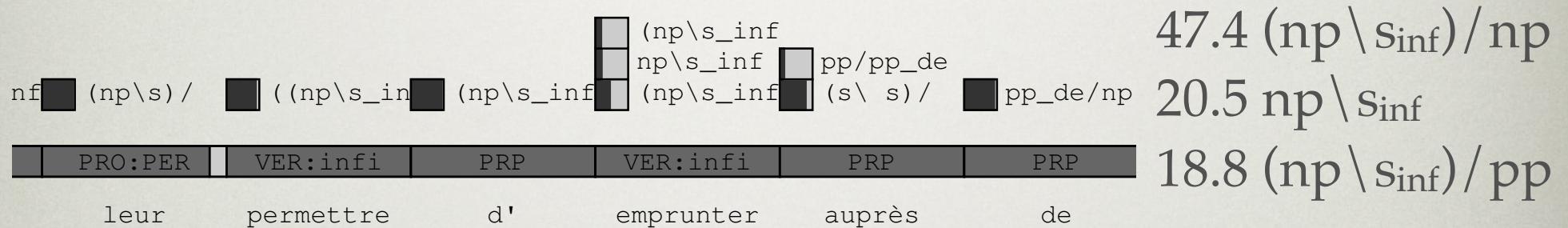
est - “is”	
(np \ s) / np	23,2 %
(np \ s) / (n \ n)	20,6 %
(np \ s) / (np \ s _{pass})	16,8 %
(cl _r \ (np \ s)) / (cl _r \ (np \ s _{ppart}))	10,8 %
(np \ s) / pp	8,1 %
(np \ s) / (np \ s _{ppart})	6,3 %
(np \ s) / (np \ s _{infX})	2,8 %
((np \ s) / s _q) / (n \ n)	2,2 %

WHAT SUPERTAGGING DOES



- Supertagging \cong statistical approximation of lexical lookup
- Assigns each word the contextually most likely (set of) formulas

WHAT SUPERTAGGING DOES



- Supertagging \approx statistical approximation of lexical lookup
- Assigns each word the contextually most likely (set of) formulas

HOW TO POSTAG/ SUPERTAG?

- Hidden Markov Model
- Maximum Entropy
- Deep learning (LSTM)

HOW TO POSTAG/ SUPERTAG?

HMM



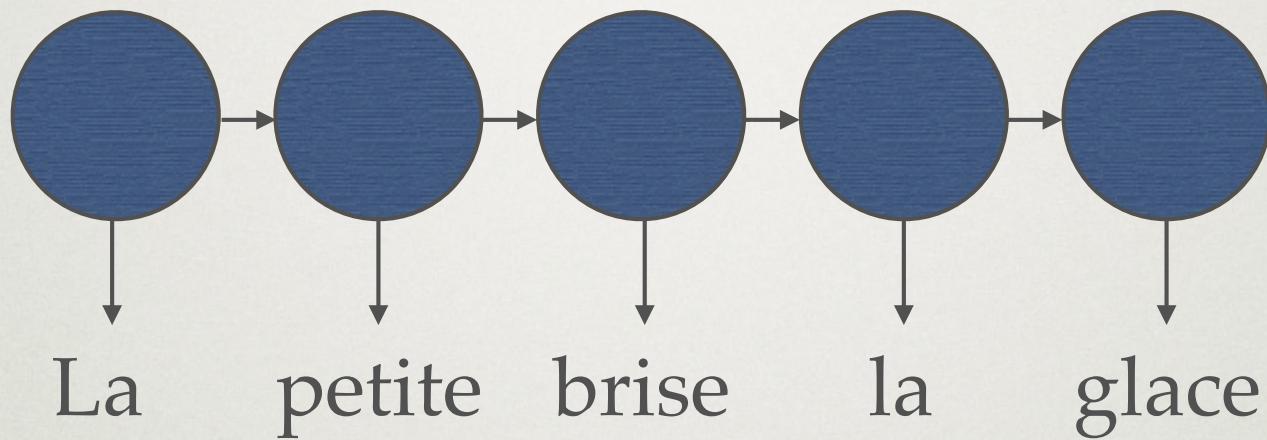
MaxEnt



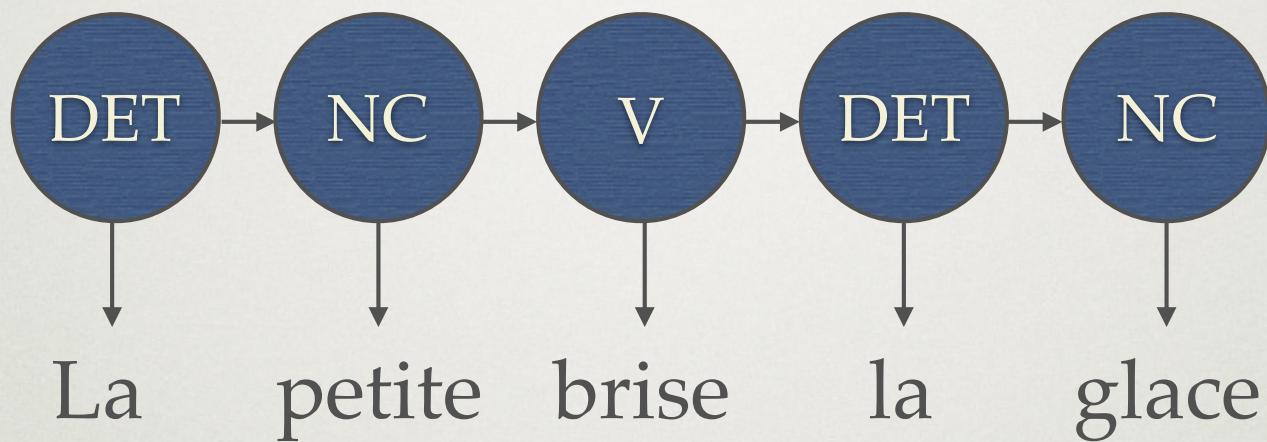
LSTM



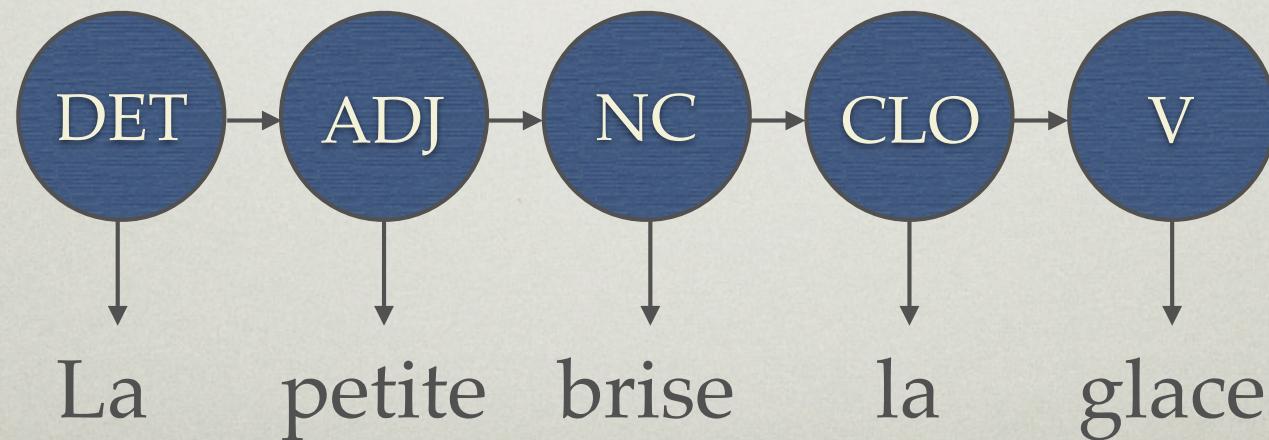
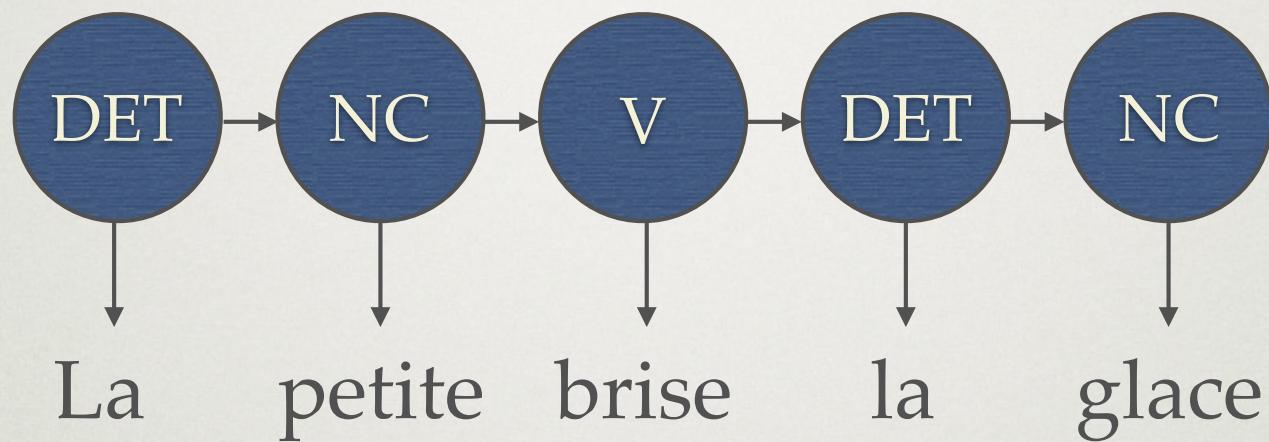
HIDDEN MARKOV MODELS



HIDDEN MARKOV MODELS



HIDDEN MARKOV MODELS

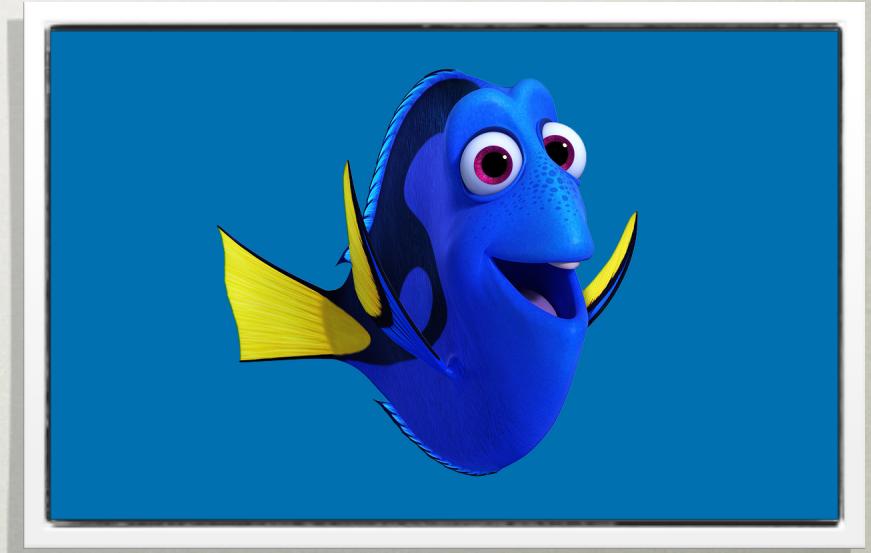


WHAT IS MAXIMUM ENTROPY?

- General-purpose statistical modelling tool
- Use any “features” we like with no assumptions of statistical independence
- Slogan: assume the least possible about data we haven’t seen

RECURSIVE NEURAL NETWORKS

- Recursive neural networks can handle long-distance dependencies in theory.
- However, in practice they do not.



LSTM

- LSTM (long short-term memory) have an explicit memory mechanism
- There are separate network parameters for “remembering/storing” and “forgetting”
- One of the most widely used modern technologies in natural language processing

TRANSLATING WORDS INTO MODEL INPUTS

first letter is "r"

last two letters
are "ue"

embedding
similar to
“boulevard” and
“avenue”

0	1	0	1	...
---	---	---	---	-----

0	0	0	1	...
---	---	---	---	-----

0.01	0.09	-0.31	0.06	...
------	------	-------	------	-----

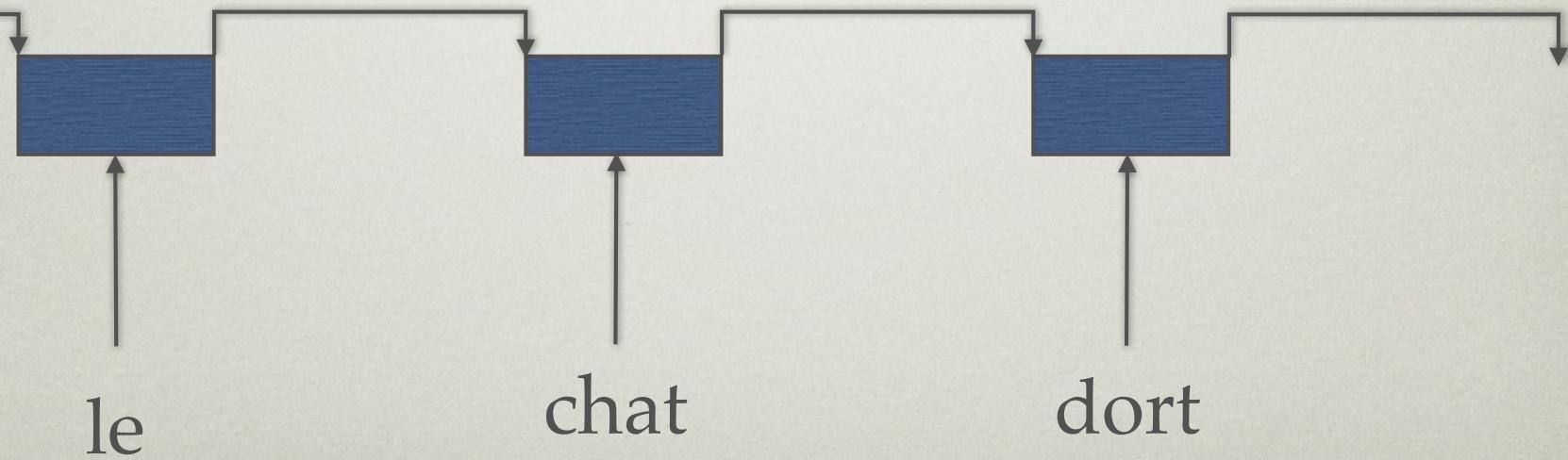
prefix

suffix

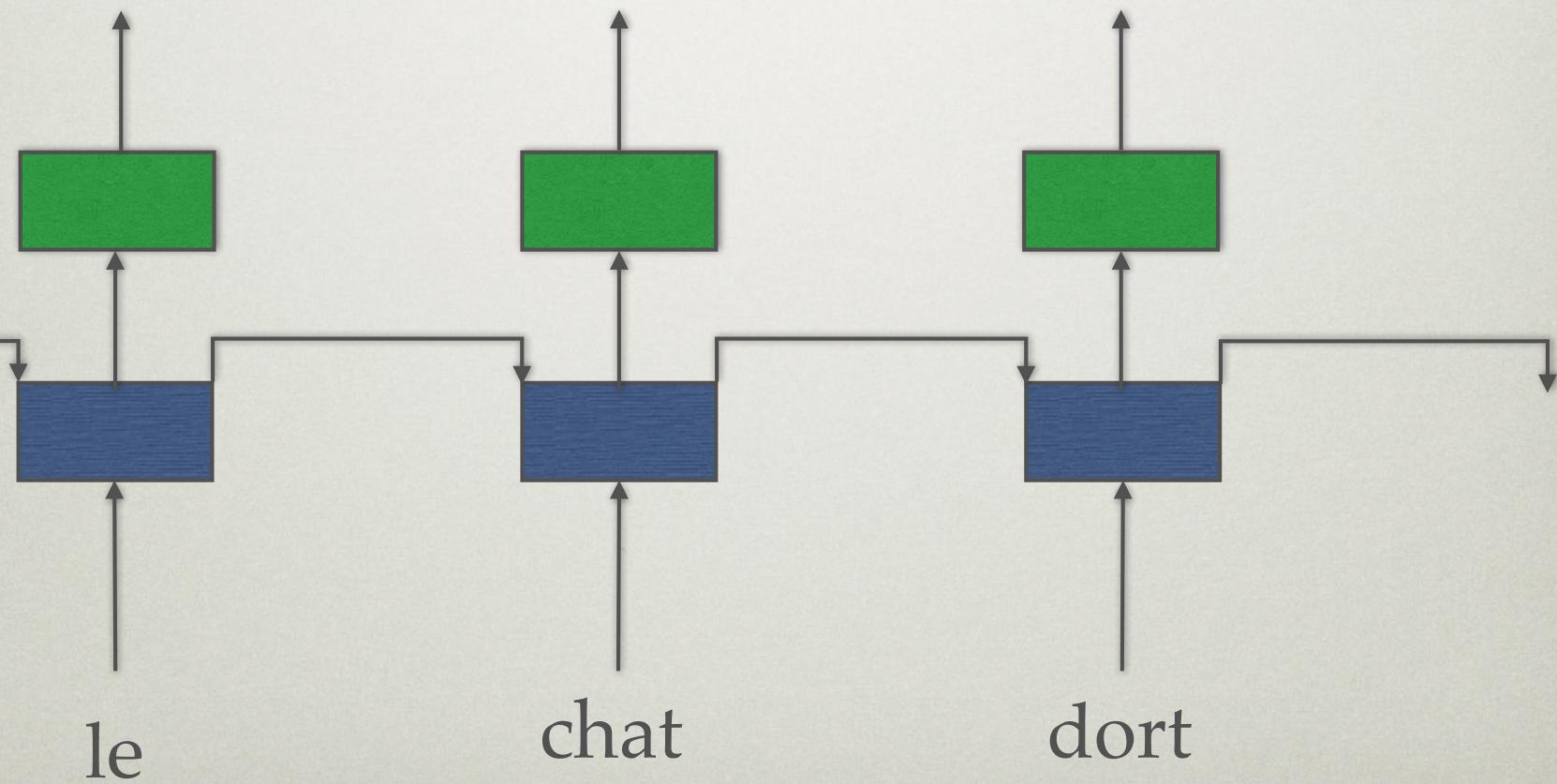
“rue”

word embedding

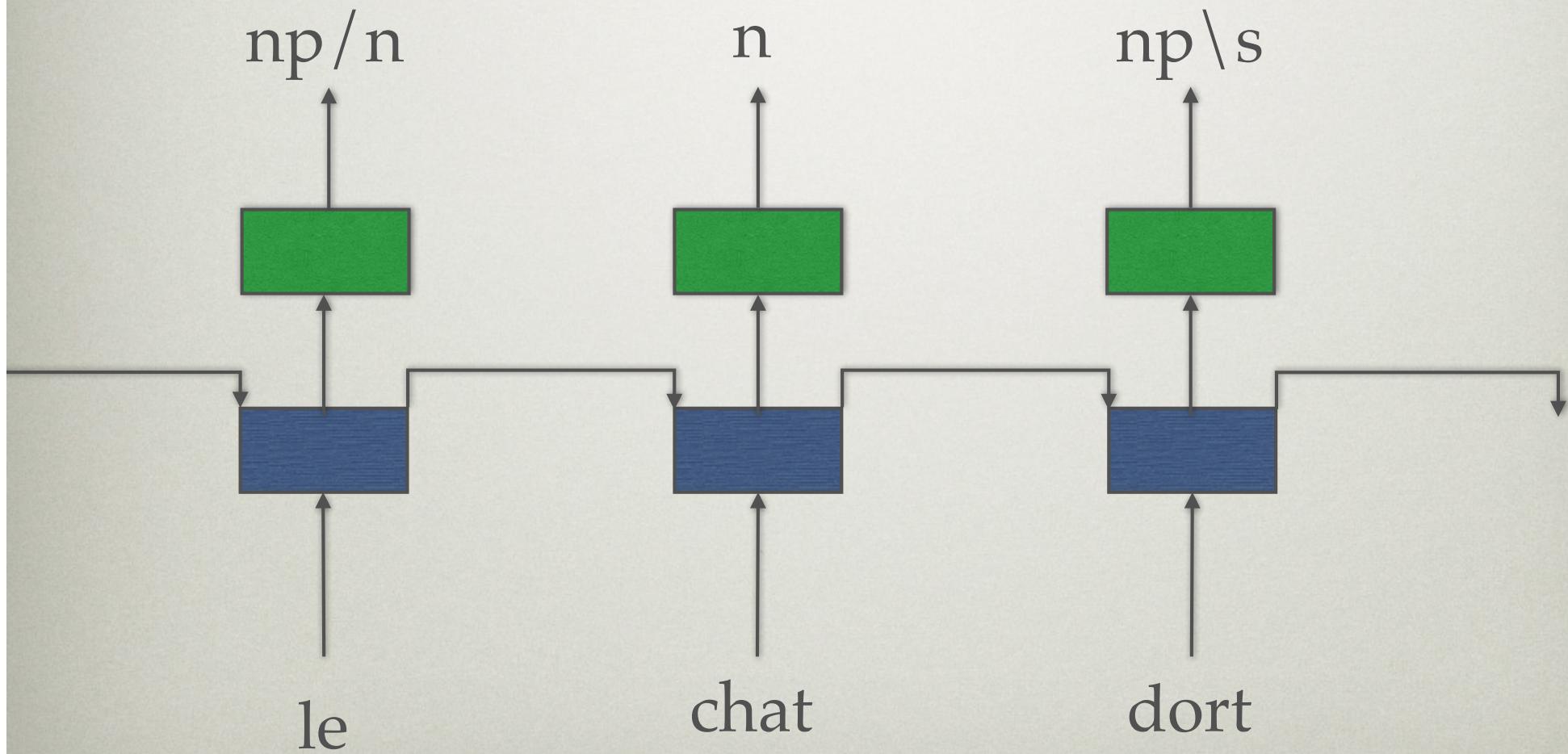
LSTM TAGGING/ SUPERTAGGING



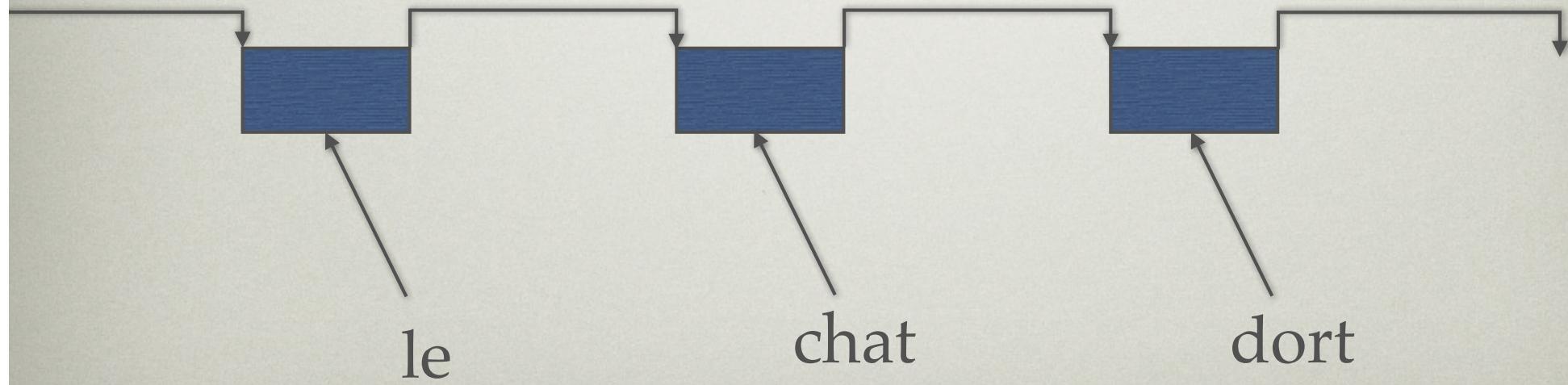
LSTM TAGGING/ SUPERTAGGING



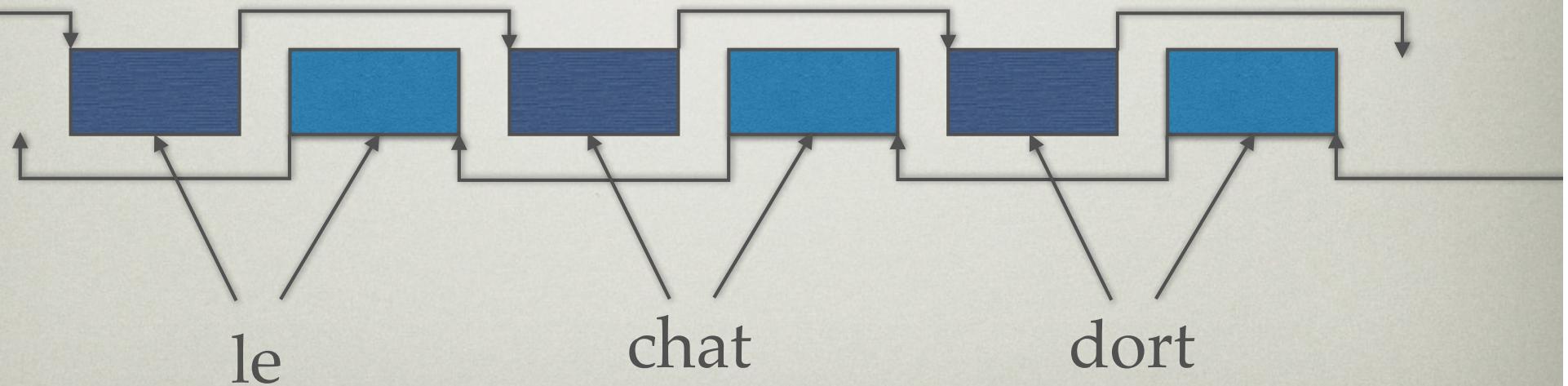
LSTM TAGGING/ SUPERTAGGING



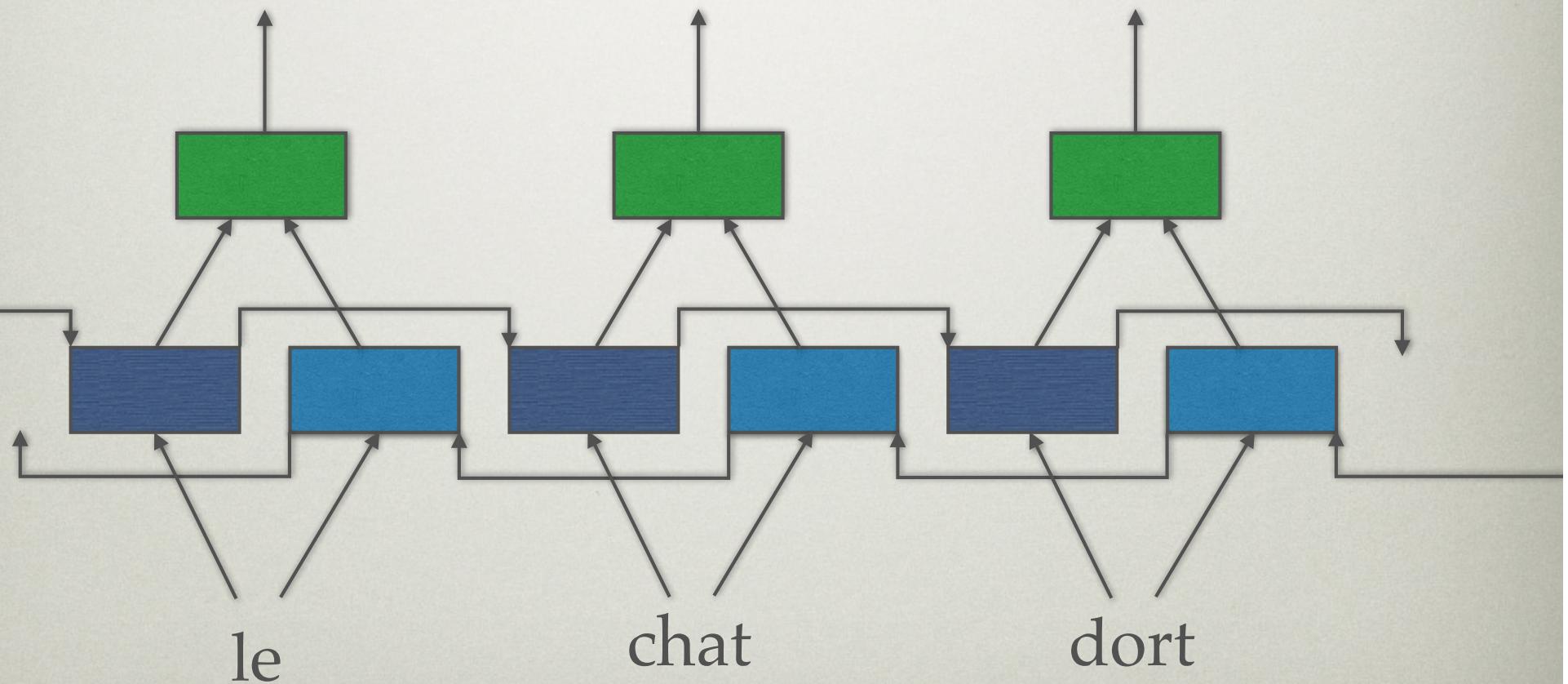
LSTM TAGGING/ SUPERTAGGING



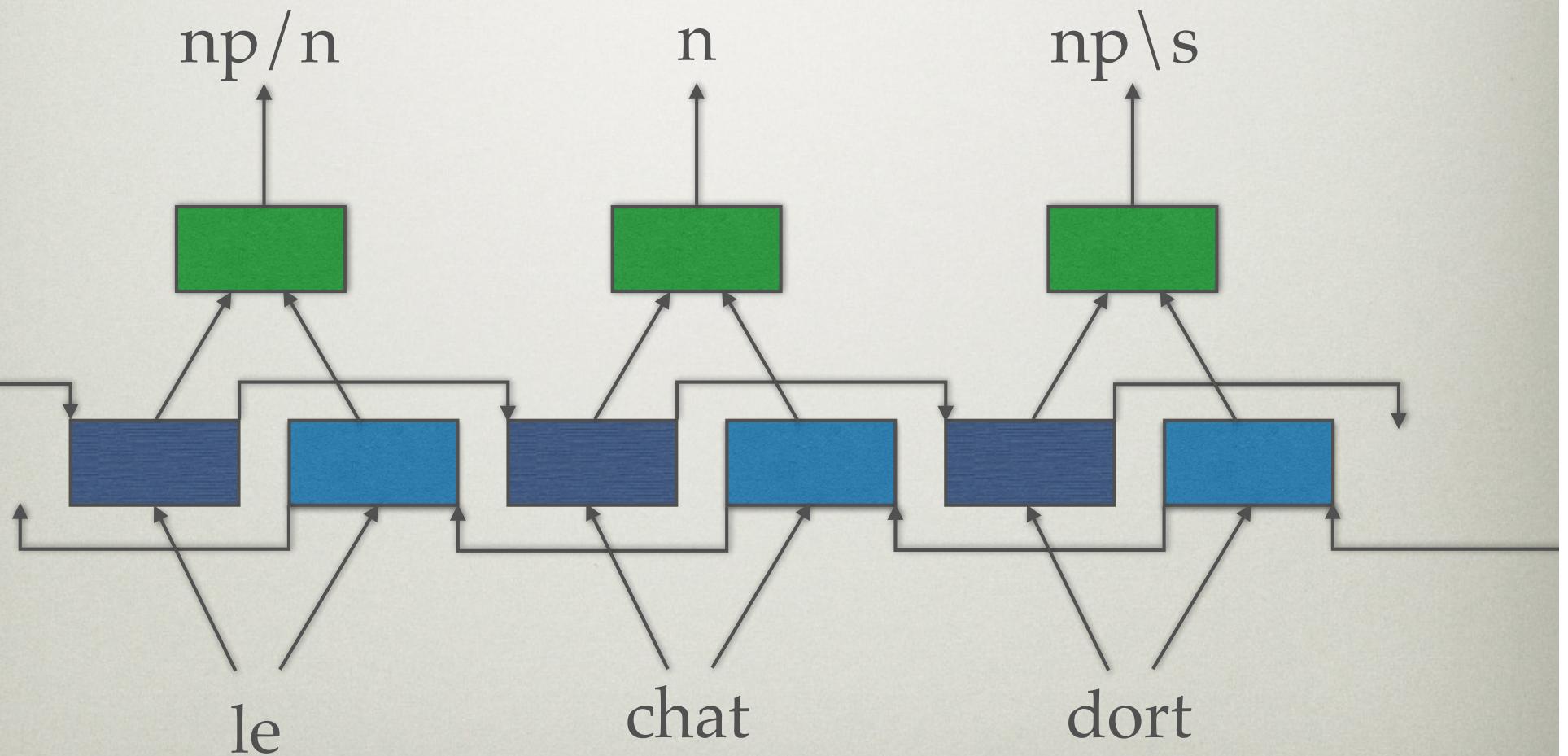
LSTM TAGGING/ SUPERTAGGING



LSTM TAGGING/ SUPERTAGGING



LSTM TAGGING/ SUPERTAGGING



SUPERTAGGER PERFORMANCE (MAXENT)

Corpus	POS	Super	0,1	0,01	F / w
FTB	97,8 %	90,6 %	96,4 %	98,4 %	2,3

SUPERTAGGER PERFORMANCE (MAXENT)

Corpus	POS	Super	0,1	0,01	F / w
FTB	97,8 %	90,6 %	96,4 %	98,4 %	2,3
Le Monde 2010	97,3 %	89,9 %	95,8 %	97,9 %	2,2
Sequoia / Annodis	97,3 %	88,1 %	94,8 %	97,6 %	2,4
Itipy / Forbes	95,7 %	86,7 %	93,8 %	97,1 %	2,6

HOW GOOD IS THIS?

- 90.6% accuracy for the best supertag sounds good, but this is given the correct part-of-speech tag
- When combining POS-tagger with supertagger, accuracy drops to 88.7% (without POS-tagger, we end up at 86.7%, so POS-tagging helps)

SUPERTAGGER PERFORMANCE

Corpus	POS	Super	0,1	0,01	0,001
MaxEnt	97,8	90,6	96,4 (1,4)	98,4 (2,3)	98,8 (4,7)
LSTM	98,4	92,2	95,8 (1,2)	97,9 (1,5)	99,0 (2,4)

with $\beta=0,0001$, we have 4,7 formulas per word
(same as ME with $\beta=0,001$) but accuracy of 99.5%

CHART-PARSING TYPE-LOGICAL GRAMMARS

- Even when using a supertagger, we there are still > 2 formulas per word
- Simple enumeration is not efficient enough
- Sharing subcomputations for proof nets requires more complex proof-theoretic machinery.

CHART-PARSING TYPE- LOGICAL GRAMMARS

- Original intention: preprocessing step for proof net parser
- At some point my preprocessor single-handedly parsed 97-98% of my extracted French treebank
- It seemed reasonable to add some extra chart rules to handle the rest.

CHART PARSING

- Categorial grammar parsing in the 90s has studied chart parsing for the Lambek calculus (and some variants)
- Chart parsing for the Lambek calculus turned out to be fairly complicated.
- An important source of complexity is the interaction of hypothetical formulas. Tiede (2001) showed that Lambek calculus proofs are not regular tree languages.

CHART-PARSING TYPE-LOGICAL GRAMMARS

- Prototype exhaustive chart parser for Type-Logical Grammars, strongly inspired by Shieber e.a. (1995)
- Fine-tuned for grammars in the restricted form as produced by the extraction algorithm
- LaTeX output

CHART-PARSING TYPE- LOGICAL GRAMMARS

$\frac{1'}{L' \vdash n / [Lex]}$ $\frac{L' \vdash n / [Lex]}{L \vdash \text{Etat} \vdash np / [Lex]}$	$\frac{\text{Etat} \vdash n / [Lex]}{n / [Lex]}$ $\frac{n / [Lex]}{(np / s_{\text{main}}) / (np / s_{\text{main}}) / [Lex]}$	$\frac{va / [Lex]}{(np / s_{\text{main}}) / (np / s_{\text{main}}) / [Lex]}$	$\frac{va / [Lex]}{(np / s_{\text{main}}) / (np / s_{\text{main}}) / [Lex]}$	$\frac{va / [Lex]}{(np / s_{\text{main}}) / (np / s_{\text{main}}) / [Lex]}$	
					$\frac{\text{Etat} \vdash np / [Lex]}{\text{Etat} \vdash np / [Lex]}$
					$\frac{\text{Etat} \vdash np / [Lex]}{\text{Etat} \vdash np / [Lex]}$
					$\frac{\text{Etat} \vdash np / [Lex]}{\text{Etat} \vdash np / [Lex]}$
					$\frac{\text{Etat} \vdash np / [Lex]}{\text{Etat} \vdash np / [Lex]}$

CHART-PARSING TYPE-LOGICAL GRAMMARS

$$\begin{array}{c}
 \frac{\text{leur}}{(np \setminus s_{inf}) / ((np \setminus s_{inf}) / \diamond_1 \square_1^\perp pp_a) \quad [Lex]} \\
 \frac{\text{leur} \circ (\text{permettre} \circ (d' \circ (\text{emprunter} \circ_1 (\text{au près} \circ (\text{de} \circ (\text{leur} \circ \text{banques})))))) \vdash np \setminus s_{inf} \quad r, r^*}{(np \setminus s_{inf}) / ((np \setminus s_{inf}) / \diamond_1 \square_1^\perp pp_a) \quad [Lex]}
 \end{array}$$

$$\frac{\text{permettre} \quad \frac{[Lex] \quad [p_0 \vdash pp_a]^0 \quad [E]}{((np \setminus s_{inf}) / (np \setminus s_{deinf})) / pp_a \quad [Lex]} \quad \frac{d'}{(np \setminus s_{deinf}) / (np \setminus s_{inf}) \quad [Lex]} \quad \frac{[Lex] \quad [pp_{de}/np] \quad [E]}{\text{de} \quad \frac{[Lex] \quad [np/n] \quad [E]}{\text{leur} \circ \text{banques} \vdash np \quad [E]}}}{\text{permettre} \circ p_0 \vdash (np \setminus s_{inf}) / (np \setminus s_{deinf}) \quad [E]}$$

$$\frac{\frac{[q_0 \vdash np]^1 \quad \frac{[Lex]}{np \setminus s_{inf} \quad [E]} \quad \frac{\text{emprunter} \quad [Lex] \quad \frac{[Lex]}{\text{au près} \quad [Lex]} \quad \frac{[Lex] \quad [pp_{de}/np] \quad [E]}{\text{de} \circ (\text{leur} \circ \text{banques}) \vdash pp_{de} \quad [E]}}}{q_0 \circ \text{emprunter} \vdash s_{inf} \quad [E]} \quad \frac{[q_0 \circ (\text{emprunter} \circ_1 (\text{au près} \circ (\text{de} \circ (\text{leur} \circ \text{banques})))) \vdash s_{inf} \quad [\backslash_1 E]}{\text{au près} \circ (\text{de} \circ (\text{leur} \circ \text{banques})) \vdash s_{inf} \setminus_1 s_{inf} \quad [\backslash_1 E]} \quad \frac{[Lex] \quad [np/n] \quad [E]}{\text{leur} \circ \text{banques} \vdash np \quad [E]}}}{q_0 \circ (\text{emprunter} \circ_1 (\text{au près} \circ (\text{de} \circ (\text{leur} \circ \text{banques})))) \vdash s_{inf} \quad [\backslash_1 E]} \quad \frac{[Lex] \quad [n] \quad [E]}{n \quad [E]}}$$

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$$\begin{array}{c}
 \frac{\text{emprunter}}{[q_0 \vdash np]^1} \quad \frac{\text{np}\backslash s_{inf}}{[\backslash E]} \quad [Lex] \quad []
 \\ \hline
 q_0 \circ \text{emprunter} \vdash s_{inf}
 \\ \hline
 \frac{\text{permettre}}{((np\backslash s_{inf}) / (np\backslash s_{deinf})) / pp_a} \quad [Lex] \quad \frac{[p_0 \vdash pp_a]^0}{[/ E]} \quad \frac{d'}{(np\backslash s_{deinf}) / (np\backslash s_{inf})} \quad [Lex] \quad \frac{q_0 \circ (\text{emprunter} \circ_1 (\varepsilon))}{\text{emprunter} \circ_1 (\text{aupr})}
 \\ \hline
 \frac{\text{permettre} \circ p_0 \vdash (np\backslash s_{inf}) / (np\backslash s_{deinf})}{(permettre \circ p_0) \circ (d' \circ (\text{emprunter} \circ_1 (\text{aupr} \circ (\text{de} \circ (\text{leurs} \circ \text{banques})))))) \vdash np\backslash s_{inf}}
 \\ \hline
 \frac{\text{leur}}{(np\backslash s_{inf}) / ((np\backslash s_{inf}) / \diamond_1 \square_1^\downarrow pp_a)} \quad [Lex] \quad \frac{(permettre \circ (d' \circ (\text{emprunter} \circ_1 (\text{aupr} \circ (\text{de} \circ (\text{leurs} \circ \text{banques})))))) \vdash (np\backslash s_{inf}) / \diamond_1 \square_1^\downarrow pp_a}{\text{leur} \circ (\text{permettre} \circ (d' \circ (\text{emprunter} \circ_1 (\text{aupr} \circ (\text{de} \circ (\text{leurs} \circ \text{banques})))))) \vdash np\backslash s_{inf}}
 \end{array}$$

SEMANTICS

- Type-logical grammar proofs are a subset of intuitionistic proofs, which correspond to lambda-terms.
- Lexical substitution follow by beta normalization gives us the full sentence meaning

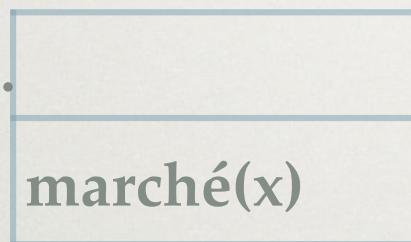
SEMANTICS

- Montague-style semantics, where the meaning of *love* is **love'**
- But which has the advantage of being scaleable, since many lexical entries follow a specific pattern
- Uses DRT

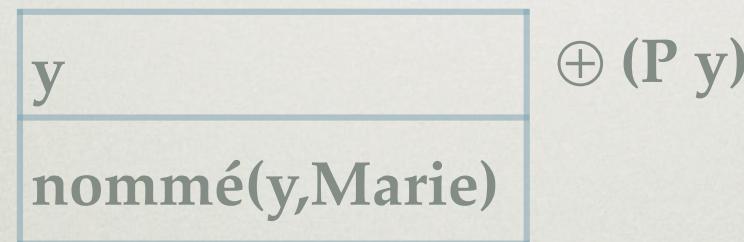
SEMANTICS

Example entries (slightly simplified)

marché: $\lambda x.$



Marie: $\lambda P.$



chaque: $\lambda P \lambda Q.$ (

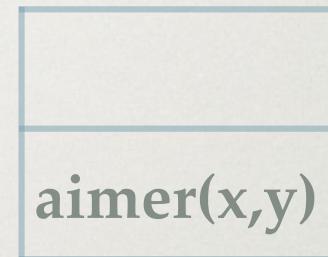


$\oplus (P\ z)) \rightarrow (Q\ z)$

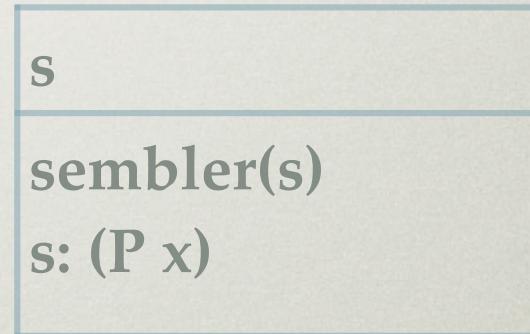
SEMANTICS

Example entries (slightly simplified)

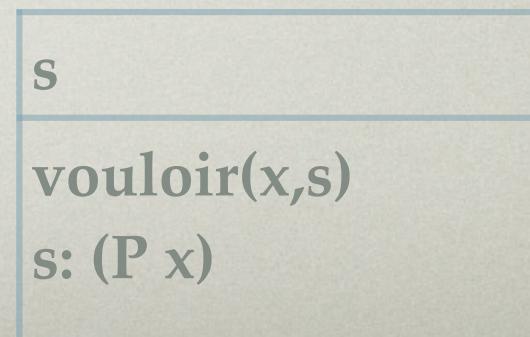
aime: $(np \setminus s)/np$ $\lambda y. \lambda x.$



semble: $(np \setminus s)/(np \setminus s_{inf})$ $\lambda P. \lambda x.$



veut: $(np \setminus s)/(np \setminus s_{inf})$ $\lambda P. \lambda x.$



SEMANTICS

- 692 words in the lexicon, with idiosyncratic properties
- 337 lexical schemata, eg. $(np \setminus s)/np$ for word w means semantics w' ($\equiv \lambda y. \lambda x. w'(x, y)$)

FUTURE WORK

- Connect word meaning to nodes in JeuxDeMots
- Add logical entailment component, possibly à la Natural Logic
- Add multiword units component (eg. “être en train de”, but also “plus X que Y”)

CONCLUSIONS

- Type-Logical Grammars factorize grammars in a way permitting an easy connection to formal semantics,
- this ensures a direct link between wide-coverage parsing and wide-coverage *semantics*.