

# Intelligent Learning and Analysis Systems SS19, Exercise Sheet 2

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## Exercise 3

### Lemma 1

Each bucket has exactly size  $w = \lceil \frac{1}{\varepsilon} \rceil$ .  $\Rightarrow$  total number of buckets is

$$\frac{m}{w} = \frac{m}{\lceil 1/\varepsilon \rceil} \leq \varepsilon m$$

So the current bucket id is at most  $\varepsilon m$ .

### Lemma 2

Proof by induction.

Base Case: Let  $b_{\text{current}} = 1$ , let  $(e, f, \Delta) \in \mathcal{D}$ . Since there have been no deletions and no decrements of  $f$  until now, we know that  $f$  is the true frequency up until this point, i.e.  $f = f_e$ .

Also,  $(e, f, \Delta)$  is only deleted when  $f \leq 1$ , in other words  $(e, f, \Delta)$  is only deleted when  $f_e = f \leq 1 = b_{\text{current}}$  which proves the base case.

Step Case: Let  $k > 1$ . Suppose that for  $b_{\text{current}} < k$  we know that  $f_e \leq b_{\text{current}}$  when  $(e, f, \Delta)$  gets deleted.

Now let  $b_{\text{current}} = k$  and  $(e, f, \Delta) \in \mathcal{D}$ . The true frequency of  $e$  in the buckets with ids  $\Delta + 1, \dots, b_{\text{current}}$  is equal to  $f$ , because since then  $f$  hasn't been decreased. Let  $b'$  be the bucket id where  $e$  was deleted the last time, if it exists. Else set  $b' = 0$ . By the induction hypothesis, the true frequency of  $e$  in buckets  $1, \dots, b'$  is  $\leq b'$ , if  $b'$  exists. The true frequency in buckets  $b' + 1, \dots, \Delta$  is 0. So the true frequency  $f_e$  at the current time (bucket id  $b_{\text{current}}$ ) is at most  $f_e \leq f + b'$ .

If  $(e, f, \Delta)$  gets deleted, we have  $f \leq 1$ , so  $f_e \leq f + b' \leq b' + 1$  and  $b'$  is at most  $b_{\text{current}} - 1$ , so  $f_e \leq b_{\text{current}}$ .

### Lemma 3

If there is no entry for  $e$  in  $\mathcal{D}$ , then there are 2 cases:

Case 1: There has never been an entry for  $e$  in  $\mathcal{D}$ .

Then  $f_e = 0 \leq \varepsilon m$ .

Case 2: There has been an entry for  $e$  in  $\mathcal{D}$  before.

Let  $b$  be the bucket id when  $e$  was deleted the last time. Then  $f_e \leq b$  by Lemma 2, and  $f_e \leq b \leq \varepsilon m$  by Lemma 1.

### Lemma 4

$f \leq f_e$ , because  $f$  gets incremented by 1 at most  $f_e$  times.

If  $e$  was deleted after processing bucket  $b$ , then by Lemma 3 the true frequency of  $e$  in buckets  $1, \dots, b$  is at most  $\varepsilon m$ . The true frequency of  $e$  in buckets  $b + 1, \dots, \Delta$  is 0, and the true frequency in buckets  $\Delta + 1, \dots, b_{\text{current}}$  is  $f$ , because since  $e$  was added (in bucket  $\Delta + 1$ ), there have been no decreases of  $f$ , and for every occurrence of  $e$ ,  $f$  was incremented by 1. So putting it all together:

$$f_e \leq \varepsilon m + 0 + f = f + \varepsilon m$$

## Exercise 4

$d_i$  is the number of elements of  $\mathcal{D}$  that were last added to  $\mathcal{D}$  during processing of bucket  $B - i + 1$ . Call the  $i$ -th summand the contribution of bucket  $B - i + 1$  to the sum. We want that the contribution of all summands does not exceed the total size of buckets  $B - j + 1, \dots, B$  (which is  $jw$ ).

So if element  $e$  was created during processing of  $B - i + 1$ , then  $e$  survives the deletion/decrementing process  $i - 1$  times. This can only happen if  $e$  occurs at least  $i$  times in buckets  $B - i + 1, \dots, B$ . So  $e$  is allowed a contribution of  $i$ . So the contribution of all the elements last added to  $\mathcal{D}$  during processing of bucket  $B - i + 1$  is allowed to be  $id_i$ , because this way the contribution of all summands can't exceed the total size of  $B - j + 1, \dots, B$ , i.e.

$$\sum_{i=1}^B id_i \leq jw$$