CSI 113 Assignment 6 1. SIt instruction is done by performing A-B and taking the resulting sign bit as the result. If A × B then sign bit will be I and if A × B then sign bit will be 0. The control signal is OIII because the control signal is a combination of bit imposts to the ALU. The signal is represented by A inv B inv, SO and SI respectively. SO + SI are the input bits that control the Mux that selects the curtant. For SIt A inv is set to 0 and B inv is set to 2 to allow for A-B to be done. The sign bit from A-B is chosen by setting SO + SI to III.
resulting sign bit as the result. If A = B then sign bit will be I and if A = B then sign bit will be 0. The control signal is Oll because the control signal is a combination of bit inputs to the ALU. The signal is represented by A inv B inv, SO and SI respectively. SO + SI one the input bits that control the Mux that selects the output. For SIT A inv is set to 0 and B inv is set to 1 to allow for A-B to be done. The sign bit from A-B is chosen by setting SO + SI to I, That's why the control signal for SIT is OIII.
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done. The sign bit from A-B is choson by setting SO & 82 to 1. That's why the control signal for SIt is OIII.
1. That's why the control signal for sit is Olll.
2.0101 * 1110 5 * 14 > 70
MD AC MQ Step iterations
0101 0000 1110
0000 1110 ACHO ACHO
9000 0111 >>
0101 0111 AC+MD 2
000 1011 >> 2
CIII IOII ACTMD 3
0011 1101 >> 3
1000 1101 Actmb 4
0100 0110 27 4
result: 0100 0110 = 70 V

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	They 1	both achie	p.10 the	some perf	armanee in	the b	est case	, but
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