CIS 5020 Analysis of Algorithms

 $Fall\ 2024$

Assignment 2

Due: September 27

Name: Yujie Xu, Richard Zhang

PennID: 51233809, 19331985

1 Question 1

1. We know that for any random variable X:

$$Var[X] = \mathbb{E}[(X - \mathbb{E}[X])^2]$$

We can deduce that $\mathbb{E}[X]$ is the minimum mean squared deviation from any constant. Then, suppose c is a constant.

$$Var[X] = \min_{c} \mathbb{E}[(X - c)^{2}]$$

This implies that for any constant c,

$$Var[X] \le \mathbb{E}[(X-c)^2]$$

Applying this equation to X = f(Z), we can get:

$$Var[f(Z)] \le \mathbb{E}[(f(Z) - c)^2]$$

If $c = f(\mathbb{E}[Z])$, then the equation becomes:

$$Var[f(Z)] \le \mathbb{E}[(f(Z) - f(\mathbb{E}[Z]))^2]$$

Because f is 1-Lipschitz, for any $x, y \in \mathbb{R}$, we have:

$$|f(x) - f(y)| \le |x - y|$$

If x = Z and $y = \mathbb{E}[Z]$, then we can get:

$$|f(Z) - f(\mathbb{E}[Z])| \le |Z - \mathbb{E}[Z]|$$

Squaring both sides and taking the expected values of both sides:

$$(f(Z) - f(\mathbb{E}[Z]))^2 \le (Z - \mathbb{E}[Z])^2.$$

$$\mathbb{E}\left[(f(Z) - f(\mathbb{E}[Z]))^2\right] \le \mathbb{E}[(Z - \mathbb{E}[Z])^2] = \operatorname{Var}[Z]$$

Because we know $Var[f(Z)] \leq \mathbb{E}[(f(Z) - f(\mathbb{E}[Z]))^2]$, the equation becomes:

$$\operatorname{Var}[f(Z)] \le \mathbb{E}[(f(Z) - f(\mathbb{E}[Z]))^2] \le \mathbb{E}[(Z - \mathbb{E}[Z])^2] = \operatorname{Var}[Z]$$

Thus, when f is a 1-Lipschitz function.

$$Var[f(Z)] \le Var[Z]$$

2. Suppose W is a random variable. Based on the definition of variance, we can write the variance of W as:

$$Var(W) = E[W^2] - (E[W])^2$$

Based on the question, we let $W = X - \beta Y$, where β is a constant.

$$W^{2} = (X - \beta Y)^{2} = X^{2} - 2\beta XY + \beta^{2}Y^{2}$$

Taking the expected values of both sides, we get:

$$E[W^{2}] = E[X^{2}] - 2\beta E[XY] + \beta^{2} E[Y^{2}]$$

Because variance of any random variable should be greater or equal to 0, we know $Var(W) \ge 0$. Thus, $E[W^2] - (E[W])^2 \ge 0$

Let us assume E[W] = 0 to simplify this question a bit. Then:

$$E[W^2] \ge 0$$

$$E[X^2] - 2\beta E[XY] + \beta^2 E[Y^2] \ge 0$$

Then we want to find a value of β such that $E[W^2]$ can be as small as possible (which is 0 in this case).

Taking the derivative with respect to β :

$$\frac{d}{d\beta} \left(E[X^2] - 2\beta \cdot E[XY] + \beta^2 \cdot E[Y^2] \right) = -2 \cdot E[XY] + 2\beta \cdot E[Y^2] = 0$$

$$\beta = \frac{E[XY]}{E[Y^2]}$$

Plugging in $\beta = \frac{E[XY]}{E[Y^2]}$ to the original inequality $E[X^2] - 2\beta E[XY] + \beta^2 E[Y^2] \ge 0$, we get:

$$E[X^2] - 2\left(\frac{E[XY]}{E[Y^2]}\right)E[XY] + \left(\frac{E[XY]}{E[Y^2]}\right)^2E[Y^2] \ge 0$$

$$E[X^2] - \frac{E[XY]^2}{E[Y^2]} \ge 0$$

$$E[XY]^2 \le E[X^2]E[Y^2]$$

Taking the square root of both sides

$$|E[XY]| \le \sqrt{E[X^2]E[Y^2]}$$

Thus,

$$E[|X\cdot Y|] \leq \sqrt{E[X^2]E[Y^2]}$$

2 Question 2

1. Let $v_1, v_2 \cdots, v_m, m = |E|$ denotes all edges in G = (V, E). Let w_i be the indicator of whether an edge is in cut(S). i.e.

$$w_i = \begin{cases} 1, & \text{if the i-th edge is in cut(S)} \\ 0, & \text{otherwise} \end{cases}$$

Let random variable X = |cut(S)| denotes the size of the cut.

Then the size of the size equals to the sum of indicators.

$$X = |\mathrm{cut}(S)| = \sum_{i=1}^{m} w_i$$

When w = 1, the 2 vertices of the edge belongs to different vertex set S and \overline{S} . Also note that the chance of 1 vertex to be chosen in S is random and have a possibility of 1/2. Thus,

$$\mathbb{E}(w_i) = \sum Pr(w_i = 1) = \frac{1}{2}$$

$$Var(w_i) = \mathbb{E}(w_i^2) - \mathbb{E}(w_i)^2 = \sum Pr(w_i^2 = 1) - \frac{1}{4} = \sum Pr(w_i = 1) - \frac{1}{4} = \frac{1}{4}$$

Note that the indicator w is independent from each other.

$$\mathbb{E}(X) = \mathbb{E}(\sum_{i=1}^{m} w_i) = \sum_{i=1}^{m} \mathbb{E}(w_i) = \sum_{i=1}^{m} \frac{1}{2} = \frac{m}{2}$$

$$Var(X) = Var(\sum_{i=1}^{m} w_i) = \sum_{i=1}^{m} Var(w_i) = \sum_{i=1}^{m} \frac{1}{4} = \frac{m}{4}$$

Conclusion: $\mathbb{E}(|\operatorname{cut}(S)|) = \frac{|E|}{2}, \operatorname{Var}(|\operatorname{cut}(S)|) = \frac{|E|}{4}$

2. Let $p = Pr(|\text{cut}(S)| \ge \frac{|E|}{2}) = Pr(X \ge \frac{m}{2}).$

Recall the expectation of X:

$$\mathbb{E}(X) = \sum_{x} x \cdot Pr(X = x)$$

$$= \sum_{x < \frac{m}{2}} x \cdot Pr(X = x) + \sum_{x \ge \frac{m}{2}} x \cdot Pr(X = x)$$

Note that the probability in single point is less equal than cumulative probability.

$$\max_{x < \frac{m}{2}} Pr(X = x) \le Pr(x < \frac{m}{2}) = 1 - p$$

$$\max_{x \ge \frac{m}{2}} Pr(X = x) \le Pr(x \ge \frac{m}{2}) = p$$

Also note that:

$$\max_{x < \frac{m}{2}} x = \frac{m}{2} - 1$$

$$\max_{x \ge \frac{m}{2}} x = m$$

Lemma. $\sum_i a_i \cdot b_i \leq \max_i(a_i) \cdot \max_i(b_i)$ **Proof.** $\sum_i a_i \cdot b_i \leq \sum_i a_i \max_i(b_i) \leq \max_i(a_i) \cdot \max_i(b_i)$ Thus,

$$\begin{split} \mathbb{E}(X) &= \\ \frac{m}{2} &= \sum_{x < \frac{m}{2}} x \cdot Pr(X = x) + \sum_{x \geq \frac{m}{2}} x \cdot Pr(X = x) \\ &\leq \max_{x < \frac{m}{2}} Pr(X = x) \max_{x < \frac{m}{2}} x + \max_{x \geq \frac{m}{2}} Pr(X = x) \max_{x \geq \frac{m}{2}} x \\ &\leq (1 - p)(\frac{m}{2} - 1) + p \cdot m \\ &\leq \frac{m}{2} + \frac{mp}{2} + p - 1 \end{split}$$

Therefore,

$$p \ge \frac{1}{\frac{m}{2} + 1}$$

Assume existing a positive constant c,

$$p \ge \frac{1}{\frac{m}{2} + 1} \ge \frac{c}{m}$$
$$0 < c \le \frac{2}{3} \le \frac{2}{1 + \frac{2}{m}}, \ m = |E| \ge 1$$

Therefore, there must exists a positive constant $0 < c \le \frac{2}{3}$,

s.t.
$$p = Pr(|\text{cut}(S)| \ge \frac{|E|}{2}) \ge \frac{c}{|E|}$$

3 Question 3

1. If we assign M items to N bins uniformly and independently, we can easily get the following conclusion:

Suppose there are two items x and y \in M. The possibility that x would collide with y is 1/N.

If we select two items from M, then the number of possible combinations is:

$$\binom{M}{2} = \frac{M(M-1)}{2}$$

Thus,

$$E[C] = \frac{1}{N} \cdot \frac{M(M-1)}{2} = \frac{M(M-1)}{2N}$$

When $M \leq 0.01\sqrt{N}$, we can rewrite E[C] as follows:

$$E[C] \le \frac{(0.01\sqrt{N})(0.01\sqrt{N} - 1)}{2N} \approx \frac{0.0001N}{2N} = \frac{0.0001}{2}$$

$$E[C] \le \frac{0.0001}{2} = 0.00005$$

Therefore, the probability that every item gets its own bin is at least (using Markov's inequality):

$$P(C=0) = 1 - P(C \ge 1) \ge 1 - 0.00005 = 0.99995 \ge \frac{9}{10}$$

Thus, when $M \leq 0.01\sqrt{N}$, every item gets its own bin with probability at least $\frac{9}{10}$.

2. To compute $E[C^2]$, we start by expressing C in terms of indicator random variables. Let X_{ij} be the indicator that two items i and j are assigned to the same bin. Then,

$$C = \sum_{1 \le i < j \le M} X_{ij}.$$

$$E[C] = \sum_{i \le j} E[X_{ij}] = {M \choose 2} \cdot \frac{1}{N} = \frac{M(M-1)}{2N}$$

Now, we need to compute $E[C^2]$. Let X_{kl} be another indicator that two items k and l are assigned to the same bin:

$$C^{2} = \left(\sum_{i < j} X_{ij}\right)^{2} = \sum_{i < j} X_{ij}^{2} + \sum_{i < j, k < l} X_{ij} X_{kl}$$

We know: $1 \le i < j \le M$ and $1 \le k < l \le M$.

We know $X_{ij}^2 = X_{ij}$ since it's a indicator variable (it can either be 1 or 0). Then, we get the following equation:

$$E[C^2] = \sum_{i < j} E[X_{ij}] + \sum_{i < j, k < l} E[X_{ij} X_{kl}] = E[C] + \sum_{i < j, k < l} E[X_{ij} X_{kl}]$$

$$E[C^{2}] = \frac{M(M-1)}{2N} + \sum_{i < j, k < l} E[X_{ij}X_{kl}]$$

Then, we need to consider computing $E[X_{ij}X_{kl}]$. We will discuss it under three different scenarios.

(1)Assuming $(i, j) \neq (k, l)$: In other words, i, j, k, l are four different items. Because the two event are independent, we have:

$$E[X_{ij}X_{kl}] = E[X_{ij}] \cdot E[X_{kl}] = (\frac{1}{N})^2 = \frac{1}{N^2}$$

The number of possible terms under this scenario is $\binom{M}{2}\binom{M-2}{2}$ as we want to select 4 different points from M.

Thus,

$$\sum_{i < j, k < l} E[X_{ij} X_{kl}] = \frac{M(M-1)}{2} \cdot \frac{(M-2)(M-3)}{2} \cdot \frac{1}{N^2}$$
$$= \frac{M(M-1)(M-2)(M-3)}{4N^2}$$

(2) Assuming there is one overlapping item:

When there is only one overlapping item, the probability that both pairs collide is the probability that the three items are all in the same bin (suppose i is the overlapping item):

$$E[X_{ij}X_{il}] = E[X_{ij}] \cdot E[X_{il}] = \frac{1}{N^2}$$

The number of possible terms is $M \cdot \binom{M-1}{2} \cdot 2$ as we want to find one overlapping item from all items first and then find the other two items from the remaining M-1 items. Thus,

$$\sum_{i < j, k < l} E[X_{ij} X_{il}] = M \cdot {M - 1 \choose 2} \cdot 2 \cdot \frac{1}{N^2}$$
$$= \frac{M(M - 1)(M - 2)}{N^2}$$

(3) Assuming (i, j) is the same with (k, l):

This situation is the same with $\sum_{i < j} E[X_{ij}]$, so we won't consider it here.

Thus,

$$E[C^{2}] = E[C] + \frac{M(M-1)(M-2)(M-3)}{4N^{2}} + \frac{M(M-1)(M-2)}{N^{2}}$$

$$\approx \frac{M^{2}}{2N} + \frac{M^{4}}{4N^{2}} + \frac{M^{3}}{N^{2}}$$

Computing $(E[C])^2$:

$$(E[C])^2 \approx (\frac{M^2}{2N})^2$$
$$\approx \frac{M^4}{4N^2}$$

Calculating $\frac{(E[C])^2}{E[C^2]}$:

$$\frac{(E[C])^2}{E[C^2]} = \frac{(E[C])^2}{\frac{M^2}{2N} + \frac{M^4}{4N^2} + \frac{M^3}{N^2}} = \frac{\frac{M^4}{4N^2}}{\frac{M^2}{2N} + \frac{M^4}{4N^2} + \frac{M^3}{N^2}}$$

The Paley-Zygmund inequality states for a non-negative random variable $Z,\,0<\alpha<1$:

$$Pr(Z > \alpha E[Z]) \ge (1 - \alpha)^2 \frac{(E[Z])^2}{E[Z^2]}$$

Let $\alpha = \frac{1}{2}$, then:

$$Pr(C > \frac{E[C]}{2}) \ge \left(1 - \frac{1}{2}\right)^2 \frac{(E[C])^2}{E[C^2]}$$

$$\ge \left(\frac{1}{2}\right)^2 \times \frac{(E[C])^2}{E[C^2]}$$

$$\ge \frac{1}{4} \times \frac{(E[C])^2}{E[C^2]}$$

$$\ge \frac{1}{4} \times \frac{\frac{M^4}{4N^2}}{\frac{M^2}{2N} + \frac{M^4}{4N^2} + \frac{M^3}{N^2}}$$

We know C can only be non-negative integer, so $Pr(C \ge 1) = Pr(C > \frac{1}{2})$. Since $M \ge 0.01\sqrt{N}$, we can set $M = \sqrt{2N}$. We know N is a positive integer:

$$Pr(C>\frac{E[C]}{2}) = Pr(C>\frac{\frac{M^2}{2N}}{2}) = Pr(C>\frac{M^2}{4N}) = Pr(C>\frac{1}{2})$$

$$Pr(C > \frac{1}{2}) \ge \frac{1}{4} \times \frac{\frac{M^4}{4N^2}}{\frac{M^2}{2N} + \frac{M^4}{4N^2} + \frac{M^3}{N^2}}$$

$$\ge \frac{\frac{4N^2}{4N^2}}{\frac{2N}{2N} + \frac{4N^2}{4N^2} + \frac{2N \cdot \sqrt{2N}}{N^2}} \cdot \frac{1}{4}$$

$$\ge \frac{1}{1 + 1 + \frac{2N \cdot \sqrt{2N}}{N^2}} \cdot \frac{1}{4}$$

$$\ge \frac{1}{2 + \frac{2N \cdot \sqrt{2N}}{N^2}} \cdot \frac{1}{4}$$

$$\ge \frac{1}{8 + \frac{8N \cdot \sqrt{2N}}{N^2}}$$

Since N is a positive integer, $\frac{1}{8 + \frac{8N \cdot \sqrt{2N}}{N^2}}$ is a positive constant.

Thus,

$$Pr(C \ge 1) = Pr(C > \frac{1}{2}) \ge \frac{1}{8 + \frac{8N \cdot \sqrt{2N}}{N^2}} \ge 0$$

Therefore, if $M \ge 0.01N$, with some constant probability $\gamma > 0$, there is a collision.