ESE5320 Homework 2

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September 12, 2025

Answers

1. Identify

1) **Scale**: this function downscales an image(stored in a 1d array "input") by only outputing every two pixel.

Filter: applies a convolution filter to the image; first horizontally, then vertically, using hard-coded coefficients. It takes the input image to produce a smoothed output image.

Differentiate: for each pixel, use the value of its left pixel and its up pixel to calculate an average, and output will be the difference between that pixel's value and the average.

Compress: This compresses the input array using Huffman coding, packing bits into the output array one byte at a time. It returns the total number of bytes written to the output.

2. Measure

Functions	$T_{\text{measured_avg}}(\text{ns})$	% of Total Latency	$T_{\text{measured_avg}}(\text{cycles})$
Scale	2.09032e+07	3.64	7.10708 + e07
Filter horizontal	1.03589e + 08	14.57	3.52202e+08
Filter vertical	1.04882e + 08	16.39	3.56599e + 08
Differentiate	2.59741e+07	3.64	8.83120e+07
Compress	3.22591e + 08	61.94	1.09681e+09

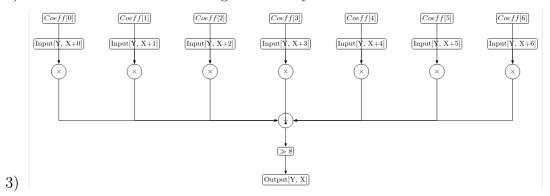
1) Done

- 2) Done
- 3) Done

3. Analyze

	Functions	$T_{\text{measured_avg}}$ (ns)	$T_{\text{measured_avg}}$ (cycles @ 1.2 GHz)
	Scale	8.426840e + 08	1.011221e+09
1)	Filter_horizontal	5.241310e + 09	6.289572e + 09
1)	Filter_vertical	5.299610e + 09	6.359532e + 09
	Differentiate	1.632450e + 09	1.958940e + 09
	Compress	5.363130e+09	6.435756e + 09

2) The Scale funtion has the highest latency.



- 4) total operation count is 3.35664e+08, and average latency is 3.35664e+08 cycles
- 5) I would choose **Compress** because it has the highest latency and therefore will benifit the most from speedups.
- 6) function Compress takes about 29% of total latency, so the maximum possible speedup is 1/(1-0.29) = 1.4
- 7) Done
- 8) critical path length is **5 operations**
- 9) need two cycles for mult, 3 for add and 1 for shift, therefore lower bound is $\bf 3$ cycles

4. Refine

1) 6.289572e + 09 cycles

	Assembly	Annotation	Number of
			Executions N
	add x6, x7, x12	LOOP1 index++	4000
	mov x4, x9	store input size	4000
	mov x0, 0	LOOP2 index++	23976000
	mov w1, 0	sum = 0	23976000
	ldrb w3, [x4, x0]	load Input byte	167832000
	ldr w2, [x5, x0, lsl 2]	load Coefficients[i]	167832000
	add x0, x0, 1	i++	167832000
	madd w1, w3, w2, w1	fused mult and add	167832000
2)	bne .L4	branch back LOOP3	167832000
	lsr w1, w1, 8	shift by 8	23976000
	strb w1, [x6], 1	store to output, then ++ output	23976000
		pointer for next one	
	add x4, x4, 1	X++	23976000
	cmp x7, x6	LOOP2 condition check	23976000
	bne .L7	branch back LOOP2	23976000
	add x7, x7, x10	++ to output pointer	4000
	add x9, x9, x11	++ to input pointer	4000
	cmp x7, x8	LOOP1 condition check	4000
	bne .L3	branch back LOOP1	4000

- 3) madd w1, w3, w2, w1 is the computation command
- 4) Done
- 5) $N_{\text{instr}} = 1174848000$

6)
$$N_{\rm mem} = 3.59640 \times 10^8$$

 $N_{\rm non_mem} = 8.15208 \times 10^8$
 $T_{\rm filter_h_measured} = 6.289572 \times 10^9$ cycles
 $T_{\rm cycle_mem} \approx 1.63552 \times 10^1$ cycles

7)
$$N_{\text{slow_mem}} = 47976007$$

 $N_{\text{fast_mem}} = 311663993$

8)
$$T_{\rm filter_h_measured} = \frac{N_{\rm non_mem}}{2} + N_{\rm fast_mem} + N_{\rm slow_mem} T_{\rm cycle_slow_mem}$$

 $T_{\rm cycle_slow_mem} \approx 1.16106 \times 10^2 \text{ cycles}$

5. Coding

1) Done

- 2) Done
- 3) cdc naive: 1711894ns, cdc rolling: 39424ns. The new function takes 97% less time