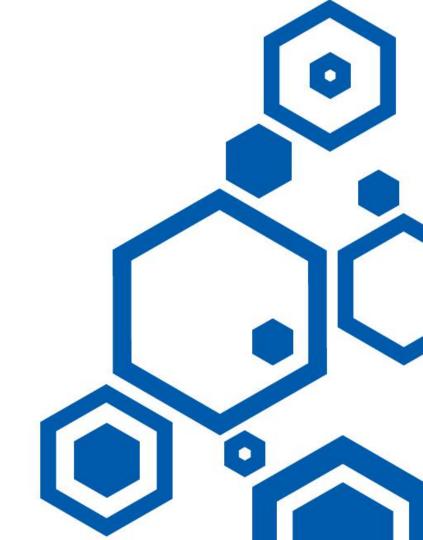


#### 第六章作业思路分享

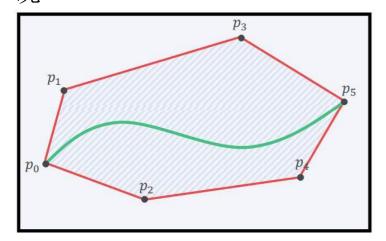
主讲人 夏韵凯

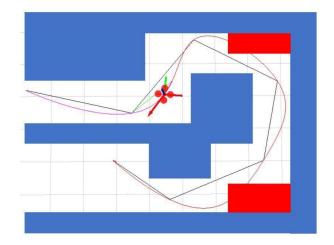


# 作业内容简介



●本章作业的核心目标是采用走廊法硬约束+贝塞尔曲线的建模方式,使用二次规划作为求解方法,解决最小snap路径生成问题。利用贝塞尔曲线轨迹在控制点内部的特性,保证轨迹在安全区内,不会出现右图的情况







●基本解题思路与第五章类似,核心是贝塞尔与多项式曲线的相互转换。

$$P_{j}(t) = p_{j}^{0} + p_{j}^{1}t + p_{j}^{2}t^{2} + \dots + p_{j}^{n}t^{n}$$

$$b_{n}^{i}(t) = c_{j}^{0}b_{n}^{0}(t) + c_{j}^{1}b_{n}^{1}(t) + \dots + c_{j}^{n}b_{n}^{n}(t) = \sum_{i=0}^{n}c_{j}^{i}b_{n}^{i}(t)$$

$$b_{n}^{i}(t) = {n \choose i} \cdot t^{i} \cdot (1-t)^{n-i}$$

●两者的系数可以通过M矩阵进行转换,有p=Mc

$$\mathbf{M} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ -6 & 6 & 0 & 0 & 0 & 0 & 0 \\ 15 & -30 & 15 & 0 & 0 & 0 & 0 \\ -20 & 60 & -60 & 20 & 0 & 0 & 0 \\ 15 & -60 & 90 & -60 & 15 & 0 & 0 \\ -6 & 30 & -60 & 60 & -30 & 6 & 0 \\ 1 & -6 & 15 & -20 & 15 & -6 & 1 \end{bmatrix}$$



●对于贝塞尔曲线,参数t的取值是[0,1]内,所以对于时间t需要进行归一化,多项式可以写成

$$f_{\mu}(t) = egin{cases} s_1 \cdot \sum_{i=0}^n c_{\mu 1}^i b_n^i \left(rac{t}{s_1}
ight), & t \in [0,T_1] \ s_2 \cdot \sum_{i=0}^n c_{\mu 2}^i b_n^i \left(rac{t}{s_2}
ight), & t \in [0,T_2] \ dots & dots \ s_m \cdot \sum_{i=0}^n c_{\mu m}^i b_n^i \left(rac{t}{s_m}
ight), & t \in [0,T_m] \end{cases}$$

●论文连接:

https://ieeexplore.ieee.org/document/8462878



●使用Bezier曲线,其Q函数会有些变化,每一项需要乘上s<sup>^</sup>-(2\*4-3), 积分上限是1,其余部分和第五章的一致。s直接取该时间段的时间值

```
0 = [1:
M = []:
M k = getM(n order);
for k = 1:n seg
    % STEP 2.1 calculate Q_k of the k-th segment, minimize snap
    0 k = [];
    t k = 1;
    s_k = ts(k);
    for i = 4:n_order
        for l = 4:n order
            den = i + l - 7:
            0_k(i+1,l+1) = i*(i-1)*(i-2)*(i-3)*l*(l-1)*(l-2)*(l-3)/den*(t_k^den)/s_k^(2*4-3);
        end
    end
    Q = blkdiag(Q, Q_k);
    M = blkdiag(M, M_k);
end
```



Stack

- ●等式约束: Boundary constraints (1) 开始点的P,V,A(2) 终点的P,V,A
- ●在使用控制点时,需要乘上归一化系数

```
n all poly = n seg*(n order+1);
% STEP 2.1 p,v,a constraint in start
Aeq_start = zeros(3, n_all_poly);
beg_start = start_cond';
S = ts(1):
k = 0; Aeq_start(k+1, 1:3) = [1,0,0]*S^(1-k);
k = 1; Aeq_start(k+1, 1:3) = [-1,1,0]*n_order*S^(1-k);
k = 2; Aeq_start(k+1, 1:3) = [1,-2,1]*n_order*(n_order-1)*S^(1-k);
% STEP 2.2 p,v,a constraint in end
Aeq_end = zeros(3, n_all_poly);
beg end = end cond';
S = ts(end);
idx = (n_{seg-1})*(n_{order+1}) + n_{order+1} - 3;
k = 0; Aeq_end(k+1, idx+(1:3)) = [0,0,1]*S^(1-k);
k = 1; Aeq_end(k+1, idx+(1:3)) = [0,-1,1]*n_order*S^(1-k);
k = 2; Aeq_{end}(k+1, idx+(1:3)) = [1,-2,1]*n_{order}(n_{order}-1)*S^(1-k);
```

Boundary Constraints:

$$a_{\mu j}^{l,0} \cdot s_j^{(1-l)} = d_{\mu j}^{(l)}$$

Continuity Constraints:

$$a_{\mu j}^{\phi,n} \cdot s_j^{(1-\phi)} = a_{\mu,j+1}^{\phi,0} \cdot s_{j+1}^{(1-\phi)}, \ a_{\mu j}^{0,i} = c_{\mu j}^i.$$

Safety Constraints:

$$\beta_{\mu j}^- \leq c_{\mu j}^i \mathbf{S} \leq \beta_{\mu j}^+, \ \mu \in \{x,y,z\}, \ i=0,1,2,...,n,$$

Dynamical Feasibility Constraints:

$$v_m^- \le n \cdot (c_{\mu j}^i - c_{\mu j}^{i-1}) \le v_m^+, a_m^- \le n \cdot (n-1) \cdot (c_{\mu j}^i - 2c_{\mu j}^{i-1} + c_{\mu j}^{i-2})/s_j \le a_m^+$$



#### ●等式约束: Continuity constraints (3,4,5) 段的P,V,A连续

```
% STEP 2.3 position continuity constrain between 2 segments
Aeq con p = zeros(n seg - 1, n all poly);
beg con p = zeros(n seg - 1, 1):
for k = 1:n \text{ seg} - 1
   s1 = ts(k);
    s2 = ts(k+1):
    Aeq con p(k, 8*k:8*k+1) = [1*s1, -1*s2]:
end
% STEP 2.4 velocity continuity constrain between 2 segments
Aeq\_con\_v = zeros(n\_seg - 1, n\_all\_poly);
beg con v = zeros(n seg - 1, 1);
for k = 1:n seg - 1
    Aeq_{con_v(k, 8*k-1:8*k+2)} = n_{order*[-1, 1, 1, -1]};
end
```

Boundary Constraints:

$$a_{\mu j}^{l,0} \cdot s_j^{(1-l)} = d_{\mu j}^{(l)}$$

Continuity Constraints:

$$a_{\mu j}^{\phi,n} \cdot s_{j}^{(1-\phi)} = a_{\mu,j+1}^{\phi,0} \cdot s_{j+1}^{(1-\phi)}, \quad a_{\mu j}^{0,i} = c_{\mu j}^{i}.$$

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% STEP 2.5 acceleration continuity constrain between 2 segments
Aeq\_con\_a = zeros(n\_seg - 1, n\_all\_poly);
beq\_con\_a = zeros(n\_seg - 1, 1);

for k = 1:n\_seg - 1
 s1 = ts(k);
 s2 = ts(k+1);
 Aeq\_con\_a(k, 8\*k-2:8\*k+3) = n\_order\*(n\_order-1)\*[1/s1, -2/s1, 1/s1, -1/s2, 2/s2, -1/s2];
end



●不等式约束: (1)Safety constraints:控制点被corridor\_range包围

```
% STEP 3.2.1 p constraint
Aieq_p = zeros(n_seg*(n_order+1)*2, n_all_poly);
bieq p = zeros(n seg*(n order+1)*2, 1);
for k = 1:n seg
   s = ts(k);
   for i = 1: (n order+1)
       Aieq p(i+16*(k-1), i+8*(k-1)) = 1*s;
       Aieq p(8+i+16*(k-1), i+8*(k-1)) = -1*s;
       bieq_p(i+16*(k-1)) = corridor_range(k, 2);
       bieq p(8+i+16*(k-1)) = -corridor range(k, 1);
   end
end
```

· Safety Constraints:

$$\beta_{\mu j}^- \le c_{\mu j}^i \mathbf{S} \le \beta_{\mu j}^+, \ \mu \in \{x, y, z\}, \ i = 0, 1, 2, ..., n,$$

大于端的处理方法是两边同时乘-1变成小于位置要乘s



●不等式约束条件: (2,3) Dynamical Feasibility Constraints: 以控制点形式

的v,a受到限制

```
% STEP 3.2.2 v constraint
  Aieq v = zeros(n seg*(n order)*2, n all polv);
  bieq_v = zeros(n_seg*(n_order)*2, 1);
  for k = 1:n seg
     for i = 1: (n order)
         Aieq v(i+14*(k-1), i+8*(k-1):i+8*(k-1)+1) = n \text{ order}*[-1, 1]:
         Aieg v(7+i+14*(k-1), i+8*(k-1):i+8*(k-1)+1) = -n \text{ order}*[-1, 1]:
         bieq v(i+14*(k-1)) = v \max;
         bieq v(7+i+14*(k-1)) = v \max;
% STEP 3.2.3 a constraint
Aieq a = zeros(n seg*(n order-1)*2, n all poly);
bieq a = zeros(n seg*(n order-1)*2, 1);
for k = 1:n seg
   s = ts(k):
   for i = 1: (n order-1)
        Aieq a(i+12*(k-1), i+8*(k-1):i+8*(k-1)+2) = n \text{ order}*(n \text{ order}-1)/s*[1,-2,1]:
        Aieq_a(6+i+12*(k-1), i+8*(k-1):i+8*(k-1)+2) = n_order*(n_order-1)/s*[-1, 2, -1];
       bieq a(i+12*(k-1)) = a max:
       bieq a(6+i+12*(k-1)) = a \max;
   end
end
```

Dynamical Feasibility Constraints:

$$v_m^- \le n \cdot (c_{\mu j}^i - c_{\mu j}^{i-1}) \le v_m^+,$$
  
 $a_m^- \le n \cdot (n-1) \cdot (c_{\mu j}^i - 2c_{\mu j}^{i-1} + c_{\mu j}^{i-2})/s_j \le a_m^+$ 

# 其他



●更多Bezier曲线的性质可以参考

https://pomax.github.io/bezierinfo/zh-CN/index.html#extended

●关于本章的课程可以参考

https://ieeexplore.ieee.org/document/8462878

https://github.com/HKUST-Aerial-Robotics/Btraj



#### 感谢各位聆听 Thanks for Listening

