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# 移动机器人运动规划 ——第五章作业讲解



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- Homework 1.1: In matlab, use the quadprog QP solver, write down a minimum snap trajectory generator
- Homework 1.2: In matlab, generate minimum snap trajectory based on the closed form solution
- Homework 2.1: In C++/ROS, use the OOQP solver, write down a minimum snap trajectory generator
- Homework 2.2: In C++/ROS, use Eigen, generate minimum snap trajectory based on the closed form solution

## Minimum Snap Trajectory Generation

- 多项式轨迹描述

本次作业采用many relative timeline, 分段轨迹表达式如下:

$$f_k(t) = \sum_{i=0}^N p_{k,i} t^i \quad 0 \leq t \leq \Delta T_k$$

- 轨迹约束条件

Derivative constraints: 规定起始状态和终止状态 $p, v, a, j$ 、中间点 $p$

Continuity constraints: 规定前后两端轨迹交点处 $p, v, a, j$ 连续

- 典型的QP问题

本次作业独立求解 $x$ 和 $y$ 轴的多项式系数, 该功能封装为函数  $\min$  MinimumSnapQP-Solver()和函数MinimumSnapCloseformSolver(), 分别表示QP解法和闭式解法

$$\begin{aligned} \min \quad & \begin{bmatrix} p_1 \\ \vdots \\ p_M \end{bmatrix}^T \begin{bmatrix} Q_1 & 0 & 0 \\ 0 & \ddots & 0 \\ 0 & 0 & Q_M \end{bmatrix} \begin{bmatrix} p_1 \\ \vdots \\ p_M \end{bmatrix} \\ \text{s.t.} \quad & A_{eq} \begin{bmatrix} p_1 \\ \vdots \\ p_M \end{bmatrix} = d_{eq} \end{aligned}$$

## Minimum Snap ---- getQ()分析

$$\begin{aligned}
 f(t) &= \sum_i p_i t^i \\
 \Rightarrow f^{(4)}(t) &= \sum_{i \geq 4} i(i-1)(i-2)(i-3)t^{i-4} p_i \\
 \Rightarrow (f^{(4)}(t))^2 &= \sum_{i \geq 4, l \geq 4} i(i-1)(i-2)(i-3)l(l-1)(l-2)(l-3)t^{i+l-8} p_i p_l \\
 \Rightarrow J(T) &= \int_{T_{j-1}}^{T_j} (f^{(4)}(t))^2 dt = \sum_{i \geq 4, l \geq 4} \frac{i(i-1)(i-2)(i-3)l(l-1)(l-2)(l-3)}{i+l-7} (T_j^{i+l-7} - T_{j-1}^{i+l-7}) p_i p_l \\
 \Rightarrow J(T) &= \int_{T_{j-1}}^{T_j} (f^{(4)}(t))^2 dt \\
 &= \begin{bmatrix} \vdots \\ p_i \\ \vdots \end{bmatrix}^T \left[ \dots \frac{i(i-1)(i-2)(i-3)l(l-1)(l-2)(l-3)}{i+l-7} T^{i+l-7} \dots \right] \begin{bmatrix} \vdots \\ p_l \\ \vdots \end{bmatrix}
 \end{aligned}$$

$$\Rightarrow J_j(T) = \mathbf{p}_j^T \mathbf{Q}_j \mathbf{p}_j \quad \text{Minimize this!}$$

```

for i=4:n_order
    for j=4:n_order
        Q_k(i+1, j+1)=factorial(i)/factorial(i-4)*factorial(j)/factorial(j-4)/(i+j-7)*ts(k)^(i+j-7);
    end
end
end
Q = blkdiag(Q, Q_k);
    
```

$$\begin{aligned}
 f(t) &= [p_0, p_1, \dots, p_7] \cdot [t^0, t^1, \dots, t^7]^T \\
 \Rightarrow \begin{cases} f(t) = P \cdot [t^0, t^1, t^2, t^3, t^4, t^5, t^6, t^7]^T \\ f^{(1)}(t) = P \cdot [0, 1, 2t, 3t^2, 4t^3, 5t^4, 6t^5, 7t^6]^T \\ f^{(2)}(t) = P \cdot [0, 0, 2, 6t, 12t^2, 20t^3, 30t^4, 42t^5]^T \\ f^{(3)}(t) = P \cdot [0, 0, 0, 6, 24t, 60t^2, 120t^3, 210t^4]^T \\ f^{(4)}(t) = P \cdot [0, 0, 0, 0, 24, 120t, 360t^2, 820t^3]^T \end{cases}
 \end{aligned}$$

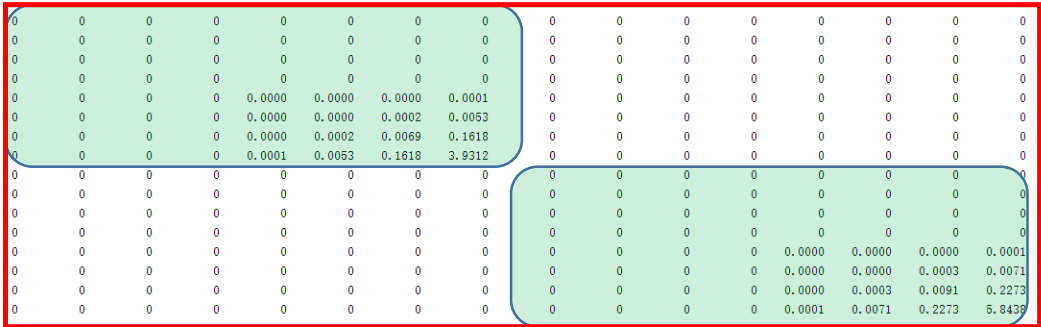
$$\begin{aligned}
 \min \quad & \begin{bmatrix} \mathbf{p}_1 \\ \vdots \\ \mathbf{p}_M \end{bmatrix}^T \begin{bmatrix} \mathbf{Q}_1 & \mathbf{0} & \mathbf{0} \\ \mathbf{0} & \ddots & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \mathbf{Q}_M \end{bmatrix} \begin{bmatrix} \mathbf{p}_1 \\ \vdots \\ \mathbf{p}_M \end{bmatrix} \\
 \text{s.t.} \quad & \mathbf{A}_{eq} \begin{bmatrix} \mathbf{p}_1 \\ \vdots \\ \mathbf{p}_M \end{bmatrix} = \mathbf{d}_{eq}
 \end{aligned}$$

$$f^{(4)}(t) = P \cdot [0, 0, 0, 0, 24, 120t, 360t^2, 820t^3]^T$$

$$\Rightarrow (f^{(4)}(t))^2 = P \cdot \begin{bmatrix} 0 \\ 24 \\ 120t \\ 360t^2 \\ 820t^3 \end{bmatrix} \cdot P^T$$

两段轨迹 $Q = \begin{bmatrix} Q_1 & 0 \\ 0 & Q_2 \end{bmatrix}$ 矩阵示意图:

$$\begin{aligned} \min \quad & \begin{bmatrix} \mathbf{p}_1 \\ \vdots \\ \mathbf{p}_M \end{bmatrix}^T \begin{bmatrix} \mathbf{Q}_1 & \mathbf{0} & \mathbf{0} \\ \mathbf{0} & \ddots & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \mathbf{Q}_M \end{bmatrix} \begin{bmatrix} \mathbf{p}_1 \\ \vdots \\ \mathbf{p}_M \end{bmatrix} \\ \text{s. t. } & \mathbf{A}_{\text{eq}} \begin{bmatrix} \mathbf{p}_1 \\ \vdots \\ \mathbf{p}_M \end{bmatrix} = \mathbf{d}_{eq} \end{aligned}$$



## Minimum Snap ---- getAbeq()分析

- Derivative constraint for one polynomial segment

- Also models waypoint constraint ( $0^{th}$  order derivative)

$$\begin{aligned}
 f_j^{(k)}(T_j) &= x_j^{(k)} \\
 \Rightarrow \sum_{i \geq k} \frac{i!}{(i-k)!} T_j^{i-k} p_{j,i} &= x_{T,j}^{(k)} \\
 \Rightarrow \begin{bmatrix} \dots & \frac{i!}{(i-k)!} T_j^{i-k} & \dots \end{bmatrix} \begin{bmatrix} \vdots \\ p_{j,i} \\ \vdots \end{bmatrix} &= x_{T,j}^{(k)} \\
 \Rightarrow \begin{bmatrix} \dots & \frac{i!}{(i-k)!} T_{j-1}^{i-k} & \dots \\ \dots & \frac{i!}{(i-k)!} T_j^{i-k} & \dots \end{bmatrix} \begin{bmatrix} \vdots \\ p_{j,i} \\ \vdots \end{bmatrix} &= \begin{bmatrix} x_{0,j}^{(k)} \\ x_{T,j}^{(k)} \end{bmatrix} \\
 \Rightarrow \mathbf{A}_j \mathbf{p}_j &= \mathbf{d}_j
 \end{aligned}$$

- Continuity constraint between two segments

- Ensures continuity between trajectory segments when no specific derivatives are given

$$\begin{aligned}
 f_j^{(k)}(T_j) &= f_{j+1}^{(k)}(T_j) \\
 \Rightarrow \sum_{i \geq k} \frac{i!}{(i-k)!} T_j^{i-k} p_{j,i} - \sum_{l \geq k} \frac{l!}{(l-k)!} T_j^{l-k} p_{j+1,l} &= 0 \\
 \Rightarrow \begin{bmatrix} \dots & \frac{i!}{(i-k)!} T_j^{i-k} & \dots & -\frac{l!}{(l-k)!} T_j^{l-k} & \dots \end{bmatrix} \begin{bmatrix} \vdots \\ p_{j,i} \\ \vdots \\ p_{j+1,l} \\ \vdots \end{bmatrix} &= 0 \\
 \Rightarrow [\mathbf{A}_j & -\mathbf{A}_{j+1}] \begin{bmatrix} \mathbf{p}_j \\ \mathbf{p}_{j+1} \end{bmatrix} &= 0
 \end{aligned}$$

这里主要根据导数约束与连续性约束构建方程，最终整理成 $\mathbf{A}_{eq}$ 与 $\mathbf{b}_{eq}$

## Minimum Snap ---- getAbeq()分析

```
#####
% p,v,a,j constraint in start.
Aeq_start = zeros(4, n_all_poly);
beq_start = zeros(4, 1);
% STEP 2.1: write expression of Aeq_start and beq_start
T=0;
for k=0:3 % p,v,a,j
    for i=k:n_order %i>=k
        Aeq_start(k+1,i+1)=factorial(i)/factorial(i-k)*(T^(i-k));
    end
end
beq_start=start_cond';
#####
% p,v,a,j constraint in end
Aeq_end = zeros(4, n_all_poly);
beq_end = zeros(4, 1);
% STEP 2.2: write expression of Aeq_end and beq_end
T=ts(end);
for k=0:3 % p,v,a,j
    for i=k:n_order %i>=k
        Aeq_start(k+1,i+1)=factorial(i)/factorial(i-k)*(T^(i-k));
    end
end
beq_end=end_cond';
#####
% position constrain in all middle waypoints
Aeq_wp = zeros(n_seg-1, n_all_poly);
beq_wp = zeros(n_seg-1, 1);
% STEP 2.3: write expression of Aeq_wp and beq_wp
%对每段轨迹的起点进行约束
for i=1:n_seg-1
    Aeq_wp(i,i*8+1)=1;
    beq_wp(i,1)=waypoints(i+1);
end
```

→ 起点状态约束

→ 终点状态约束

→ 中间点位置约束

- Derivative constraint for one polynomial segment
  - Also models waypoint constraint ( $0^{th}$  order derivative)

$$\begin{aligned}
 f_j^{(k)}(T_j) &= x_j^{(k)} \\
 \Rightarrow \sum_{i \geq k} \frac{i!}{(i-k)!} T_j^{i-k} p_{j,i} &= x_{T,j}^{(k)} \\
 \Rightarrow \begin{bmatrix} \dots & \frac{i!}{(i-k)!} T_j^{i-k} & \dots \end{bmatrix} \begin{bmatrix} \vdots \\ p_{j,i} \\ \vdots \end{bmatrix} &= x_{T,j}^{(k)} \\
 \Rightarrow \begin{bmatrix} \dots & \frac{i!}{(i-k)!} T_{j-1}^{i-k} & \dots \end{bmatrix} \begin{bmatrix} \vdots \\ p_{j,i} \\ \vdots \end{bmatrix} &= \begin{bmatrix} x_{0,j}^{(k)} \\ x_{T,j}^{(k)} \end{bmatrix} \\
 \Rightarrow A_j p_j &= d_j
 \end{aligned}$$

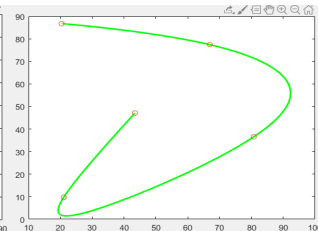
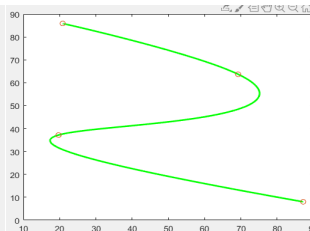
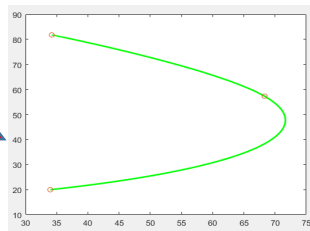
## Minimum Snap ---- getAbeq()分析

```
#####
% position continuity constrain between each 2 segments
Aeq_con_p = zeros(n_seg-1, n_all_poly);
beq_con_p = zeros(n_seg-1, 1);
% STEP 2.4: write expression of Aeq_con_p and beq_con_p
k=0;%k=0,1,2,3
for n=0:n_seg-1-1
    T=ts(n+1);
    idx=(n)*(n_order+1);
    for i=k:n_order
        Aeq_con_p(n+1, idx+i+1)=factorial(i)/factorial(i-k)*(T^(i-k));
    end
    T=0;
    idx=(n+1)*(n_order+1);
    for i=k:n_order
        Aeq_con_p(n+1, idx+i+1)=-factorial(i)/factorial(i-k)*(T^(i-k));
    end
end
% STEP 3: get the coefficients of i-th segment of both x-axis
% and y-axis
Pxi = [];
Pyi = [];
Pxi=flipud(poly_coef_x(i*8+1:i*8+8));
Pyi=flipud(poly_coef_y(i*8+1:i*8+8));
for t = 0:tstep:ts(i+1)
    X_n(k) = polyval(Pxi, t);
    Y_n(k) = polyval(Pyi, t);
    k = k + 1;
end
```

K=0,1,2,3分别表示位置、速度、加速度、加加速度

连续性约束

可视化



- Continuity constraint between two segments
- Ensures continuity between trajectory segments when no specific derivatives are given

$$f_j^{(k)}(T_j) = f_{j+1}^{(k)}(T_j)$$

$$\Rightarrow \sum_{i=k}^l \frac{i!}{(i-k)!} T_j^{i-k} p_{j,i} - \sum_{i=k}^l \frac{i!}{(i-k)!} T_j^{i-k} p_{j+1,i} = 0$$

$$\Rightarrow \begin{bmatrix} \dots & \frac{i!}{(i-k)!} T_j^{i-k} & \dots & -\frac{l!}{(l-k)!} T_j^{l-k} & \dots \end{bmatrix} \begin{bmatrix} p_{j,i} \\ \vdots \\ p_{j+1,i} \end{bmatrix} = 0$$

$$\Rightarrow [A_j \quad -A_{j+1}] \begin{bmatrix} p_j \\ p_{j+1} \end{bmatrix} = 0$$



## Minimum Snap ---- getM()分析（闭式解法）

- We have  $\mathbf{M}_j \mathbf{p}_j = \mathbf{d}_j$ , where  $\mathbf{M}_j$  is a mapping matrix that maps polynomial coefficients to derivatives

$$J = \begin{bmatrix} \mathbf{p}_1 \\ \vdots \\ \mathbf{p}_M \end{bmatrix}^T \begin{bmatrix} \mathbf{Q}_1 & 0 & 0 \\ 0 & \ddots & 0 \\ 0 & 0 & \mathbf{Q}_M \end{bmatrix} \begin{bmatrix} \mathbf{p}_1 \\ \vdots \\ \mathbf{p}_M \end{bmatrix} \quad J = \begin{bmatrix} \mathbf{d}_1 \\ \vdots \\ \mathbf{d}_M \end{bmatrix}^T \left( \begin{bmatrix} \mathbf{M}_1 & 0 & 0 \\ 0 & \ddots & 0 \\ 0 & 0 & \mathbf{M}_M \end{bmatrix} \right)^{-T} \begin{bmatrix} \mathbf{Q}_1 & 0 & 0 \\ 0 & \ddots & 0 \\ 0 & 0 & \mathbf{Q}_M \end{bmatrix} \left( \begin{bmatrix} \mathbf{M}_1 & 0 & 0 \\ 0 & \ddots & 0 \\ 0 & 0 & \mathbf{M}_M \end{bmatrix} \right)^{-1} \begin{bmatrix} \mathbf{d}_1 \\ \vdots \\ \mathbf{d}_M \end{bmatrix}$$

这里主要使用映射矩阵 $\mathbf{M}$ 将多项式系数映射为导数，提高求解数值稳定性。

## Minimum Snap ---- getM()分析（闭式解法）

```
M = [];
for n = 1:n_seg
    M_k = zeros(n_order+1,n_order+1);
    % STEP 1.1: calculate M_k of the k-th segment
    T=0;
    for k=0:3
        for i=k:n_order
            M_k(k+1,i+1)=factorial(i)/factorial(i-k)*(T^(i-k));
        end
    end
    T=ts(n);
    for k=0:3
        for i=k:n_order
            M_k(4+k+1,i+1)=factorial(i)/factorial(i-k)*(T^(i-k));
        end
    end
    M = blkdiag(M, M_k);
end
```

实现

$$M_j p_j = d_j$$

$$f_j^{(k)}(T_j) = x_j^{(k)} \\ \Rightarrow \sum_{i \geq k} \frac{i!}{(i-k)!} T_j^{i-k} p_{j,i} = x_{T,j}^{(k)}$$

凝练

定义  $P_{\text{total}} = \begin{pmatrix} p_1 \\ p_2 \\ \vdots \\ p_M \end{pmatrix}^T$ ,  $d_{\text{total}} = \begin{bmatrix} d_1 \\ d_2 \\ \vdots \\ d_K \end{bmatrix}$ ,  $\begin{bmatrix} p_{0,1} \\ p_{1,1} \\ \vdots \\ p_{7,1} \\ p_{0,K} \\ p_{1,K} \\ \vdots \\ p_{7,K} \end{bmatrix} = \begin{bmatrix} d_1 \\ d_2 \\ \vdots \\ d_K \end{bmatrix} = \begin{bmatrix} p_0 \\ v_0 \\ a_0 \\ j_0 \\ p_1 \\ v_1 \\ a_1 \\ j_1 \\ \vdots \\ p_K \\ v_K \\ a_K \\ j_K \end{bmatrix}$ , 由  $M_j p_j = d_j$ , 可得

$$\begin{aligned} p_{0,j} &= p_{j-1} \\ p_{1,j} &= v_{j-1} \\ 2p_{2,j} &= a_{j-1} \\ 6p_{3,j} &= j_{j-1} \\ p_{0,j} + p_{1,j}t_j + p_{2,j}t_j^2 + p_{3,j}t_j^3 + p_{4,j}t_j^4 + p_{5,j}t_j^5 + p_{6,j}t_j^6 + p_{7,j}t_j^7 &= p_j \\ p_{1,j} + 2p_{2,j}t_j + 3p_{3,j}t_j^2 + 4p_{4,j}t_j^3 + 5p_{5,j}t_j^4 + 6p_{6,j}t_j^5 + 7p_{7,j}t_j^6 &= v_j \\ 2p_{2,j} + 6p_{3,j}t_j + 12p_{4,j}t_j^2 + 20p_{5,j}t_j^3 + 30p_{6,j}t_j^4 + 42p_{7,j}t_j^5 &= a_j \\ 6p_{3,j} + 24p_{4,j}t_j + 60p_{5,j}t_j^2 + 120p_{6,j}t_j^3 + 210p_{7,j}t_j^4 &= j_j \end{aligned}$$

## Minimum Snap ---- getCt()分析（闭式解法）

- Use a selection matrix  $\mathbf{C}$  to separate free ( $\mathbf{d}_p$ ) and constrained ( $\mathbf{d}_F$ ) variables
  - Free variables : derivatives unspecified, only enforced by continuity constraints

$$\mathbf{C}^T \begin{bmatrix} \mathbf{d}_F \\ \mathbf{d}_p \end{bmatrix} = \begin{bmatrix} \mathbf{d}_1 \\ \vdots \\ \mathbf{d}_M \end{bmatrix} \quad \longrightarrow \quad J = \begin{bmatrix} \mathbf{d}_F \\ \mathbf{d}_p \end{bmatrix}^T \underbrace{\mathbf{C}\mathbf{M}^{-T}\mathbf{Q}\mathbf{M}^{-1}\mathbf{C}^T}_{\mathbf{R}} \begin{bmatrix} \mathbf{d}_F \\ \mathbf{d}_p \end{bmatrix} = \begin{bmatrix} \mathbf{d}_F \\ \mathbf{d}_p \end{bmatrix}^T \begin{bmatrix} \mathbf{R}_{FF} & \mathbf{R}_{FP} \\ \mathbf{R}_{PF} & \mathbf{R}_{PP} \end{bmatrix} \begin{bmatrix} \mathbf{d}_F \\ \mathbf{d}_p \end{bmatrix}$$

- Turned into an unconstrained quadratic programming that can be solved in closed form:

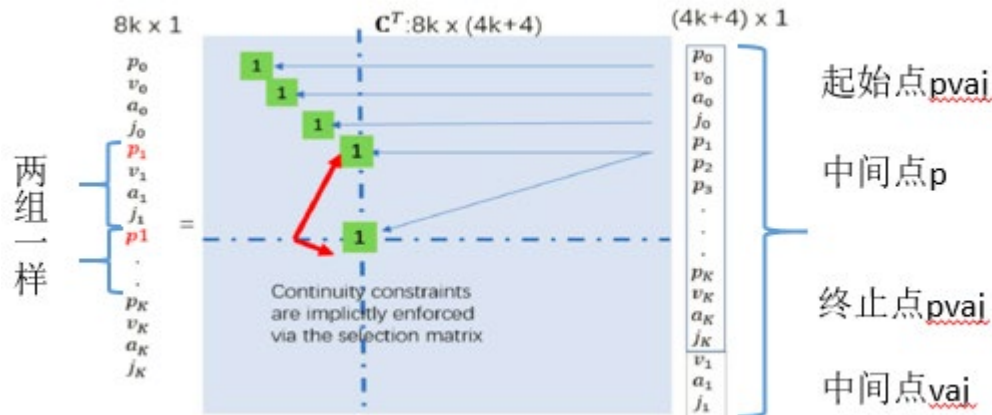
$$J = \mathbf{d}_F^T \mathbf{R}_{FF} \mathbf{d}_F + \mathbf{d}_F^T \mathbf{R}_{FP} \mathbf{d}_p + \mathbf{d}_p^T \mathbf{R}_{PF} \mathbf{d}_F + \mathbf{d}_p^T \mathbf{R}_{PP} \mathbf{d}_p$$

$$\mathbf{d}_p^* = -\mathbf{R}_{PP}^{-1} \mathbf{R}_{FP}^T \mathbf{d}_F$$

这里通过使用选择矩阵 $\mathbf{C}$ 分离自由变量 $\mathbf{d}_p$ 和固定变量 $\mathbf{d}_F$ ，转化为无约束二次规划以闭式求解。

## Minimum Snap ---- getCt()分析（闭式解法）

选择矩阵Ct构造规则：

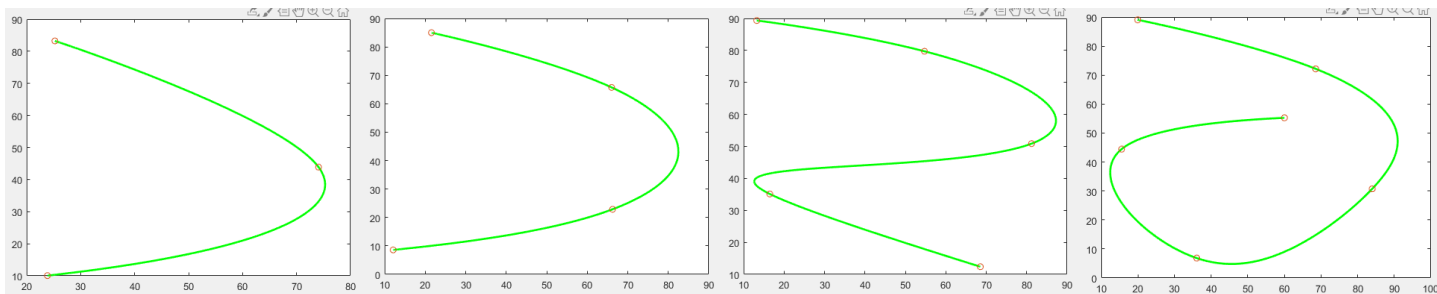


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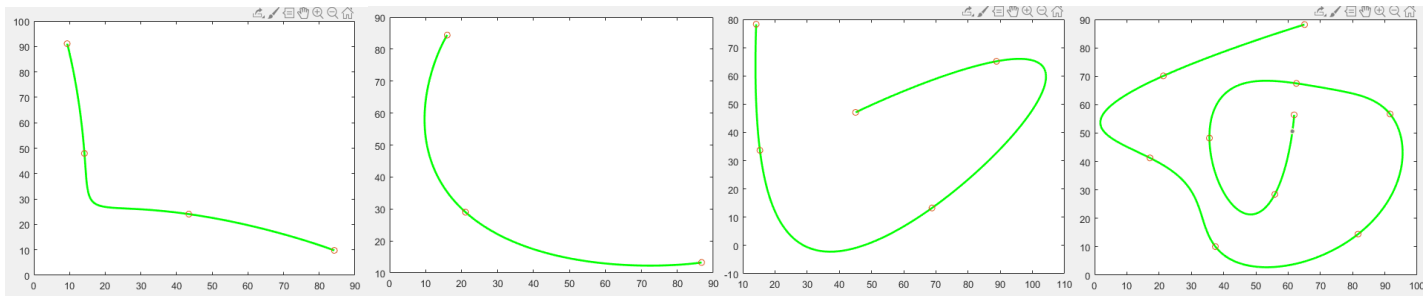
#####
% STEP 2.1: finish the expression of Ct
%fixed derivatives
%起点状态: p0,v0,a0,j0
for i=1:4
    Ct(i,i)=1;
end
%中间点位置: p1,p2,...,p(n-1)
idx_df = size(Ct,2);
for i=1:n_seg-1
    idx=8*i;
    Ct(idx-4+1,idx_df+i)=1;
    Ct(idx+1,idx_df+i)=1;
end
%终点状态: pn,vn,an,jn
idx_df = size(Ct,2);
for i=1:4
    idx=(n_seg-1)*8+4;
    Ct(idx+i,idx_df+i)=1;
end
%free derivatives
%v1,a1,j1,...,v(n-1),a(n-1),j(n-1)
idx_df=size(Ct,2);
for i=1:n_seg-1
    idx=8*i;
    idx_df2=idx_df+(i-1)*3;
    Ct(idx-4+2,idx_df2+1)=1;
    Ct(idx-4+3,idx_df2+2)=1;
    Ct(idx-4+4,idx_df2+3)=1;
    Ct(idx+2,idx_df2+1)=1;
    Ct(idx+3,idx_df2+2)=1;
    Ct(idx+4,idx_df2+3)=1;
end
    
```

## Minimum Snap ---- 效果展示

➤ **MinimumSnapQPSolver()**求解结果:

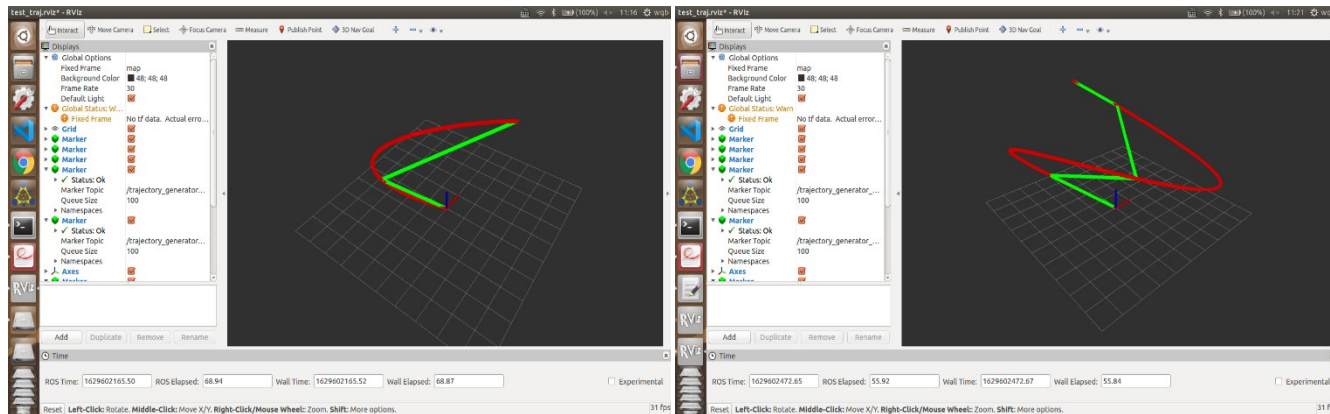


➤ **MinimumSnapCloseformSolver()**求解结果:



## Minimum Snap ---- 效果展示

ROS环境下仍旧可以按照**求解器**与**闭式求解**两种方法进行，大家可以按照所述Minimum Snap 方法进行实现，下面不在赘述具体过程，直接给出运行效果。



**感谢各位聆听 !**  
Thanks for Listening

