

Fault-tolerant design techniques

Key ingredients

- Error processing
- Error detection
- Error diagnosis
- Error recovery
- Backward & forward
- Fault treatment
- Fault diagnosis
- Fault isolation
- Reconfiguration

Strategies

Tolerance achieved through redundancy by:

- Fault masking
- Reconfiguration
 - Fault detection
 - Fault location
 - Fault containment
 - Fault recovery

Redundancy

Addition of information, resources, time

Hardware

Static

Fault masking: passive redundancy, no action required

Dynamic

Active redundancy: Detect faults and perform some action

Hybrid

Combine features of both: NMR with spares; self-purging redundancy

Voter
Triple modular redundancy
Formulas
N-Modular redundancy
Pros: if $N > 2f \rightarrow f$ faults can be tolerated.
Cons: higher cost
Hardware voting: faster, cost hw
Software voting: slow, no addition. Hw

Sift-out MR: N modules using comparators (compare each module's output), detectors (determine which disagreements are reported by the comparator produces 0 or 1), collector (produce system's output) etc.

Standby sparing

One module working, the others in standby; if a fault is detected, that module removed

Hot standby sparing

Cold standby sparing

Key advantage: fewer power tha N parallel modules

Pair-and-a-spare

Combine features of standby sparing

Two modules parallel, results compared to provide the error detection

Watchdog timers

Concept: lack of an action is an indicative of fault

Must be reset

Assumptions: system is fault free if it can perform repetitively an action

Can be used to detect faults

Software

Consistency checks: verify the correctness

Capability checks: verify possess the expected capabilities

ALU Tests: proc. can execute instructions and compare results known in ROM.

Testing of communication: achieved by sending messages from one processor to another.

2 approaches:

N-version prog.

Masks fault, relies on voting, concurrent

Design diversity; can tolerate hw and sew faults, not correlated ones; important to decide the #versions required.

Recovery blocks

Backward scheme, relies on acceptance test, serial

Multiple alternates (backups); primary task executes first, output checked by acceptance test; if not accept. Another is executed

Acceptance test are based on acceptable ranges (as inferential statistics);

Is also called primary-backup approach

SW rejuvenation

Information

Guarantee data consistency by exploiting additional info to achieve redundant encoding; it permit to detect and correct bits

2 types of codes: Error detection codes, Error correction code

Redundant codes

Binary codes: Hamming distance

EDC: Intro redundancy in the info; an error generates a word not belonging to the code; error weight: # corrupted bits tolerated; using hamming we will calculate the distance as the number of different positions of bits; with parity code, we calculate the difference between the words obtained adding the parity bit(double/even errors are unnoticeable)

ECC: a code havin minimum distance d can correct errors with weight $\leq (d-1)/2$ per difetto; codici hamming;

Circuito di EDAC (Error detection and correction)

Time

Transient Fault detection: computations repeated in different points in time, no extra hw.

Permanent fault detection: 1st comp: operands used; 2nd comp: operands econded

RAID (Redundant Array of Inexpensive Disks)

Architecture

Multple small, inexpensive disks drives into a group; appear as a single virtual drive; support fault-tolerance (redundant in disks); data striping for better performance.

Issues: access concurrency; throughput; data striping

Levels

RAID-0

No redundancy \rightarrow no fault tolerance; high i/o perf; storage efficiency

RAID-1

Disk mirroring; best read perf; poor write perf; good fault tolerance

RAID-2

Bit level striping; uses Hamming (ECC); tolerate failures ($\# \text{red disks} \sim O(\log(\text{tot disks}))$); better stop. Eff. Than 1; high throughput but no conc.; exp. write.

RAID-3

Subtopic 1

RAID-4

Block Level striping; intro trade-off; parity disk is a bottleneck in small write; no prob for reads.

RAID-5

BLS with Distributed parity; reduces parity bottleneck; best small and large read; best large write; costly for small write

Limits: most employed; the larger the number of disks, the better perf. But the larger the prob of disk failure

RAID-6

BLS with Dual Distributed parity; 2sets of parity; better fault tolerance (faster data reconstr. Than RAID-5); write worse than 5 due to overhead; better read perf. Cause data and parity are spread into more disks.

Error prop in DS & Rollback Error recovery Techniques

Systems consist of N processes which comm through messages and cooperate to execute Dist.app.

ROLLBACK

Treats DS as a coll of processes that comm through a net.; fault tolerance achieved by storage saving states; failed process restarts from the saved state (checkpoint)

Checkpoints

Uncoordinated: process' autonomy in deciding checkpoints

Coordinated: orchestration of checkpoints in order to form a consistent global state; (coordinator-processes communication to coordinate)

Communication-induced: avoids domino; local & forced checkpoints

Protocols:

Global checkpoint: set of N local checkpoints, one from each process, forming a CSS. The most recent is called recovery line.

Consistent System state

If a process's state reflects a message receipt, then the state of the corr sender reflects sending that message (fundamental goal for rollback-recovery protocol)