

**SAPIENZA Università di Roma – MSc. in Engineering in Computer Science Formal
Methods - Final Test – June, 26 2020**

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(Time to complete the test: 2:30 hours)

Exercise 1. Express the following UML class diagram in FOL:

**Alphabet: Course(x), Student(y), BScStudent(x), MScStudent(x), Professor(z),
Exam(x, y, z), Mark(x, y, z, w), Supervises(x, y)**

Axioms:

Forall x. BScStudent(x) implies Student(x) ISA

Forall x. MScStudent(x) implies Student(x) ISA

Forall x. BScStudent(x) implies not MScStudent(x) DISJOINTNESS

Forall x. Student(x) implies BScStudent(x) OR MScStudent(x) COMPLETENESS

Forall x, y. Supervises(x, y) implies Professor(x) and MScStudent(y)

Forall x. Professor(x) implies $1 \leq \#\{y \mid \text{supervises}(x, y)\}$

Forall x. MScStudent(x) implies $0 \leq \#\{y \mid \text{supervises}(y, x)\} \leq 1$

Forall x, y, z. Exam(x, y, z) implies Course(x) and Student(y) and Professor(z)

Forall x, y, z, w. Mark(x, y, z, w) implies Exam(x, y, z) and Integer(w)

Forall x, y, z, z'. Exam(x, y, z) and Exam(x, y, z') implies $z=z'$ KEY

Forall x, y, z. Exam(x, y, z) implies $1 \leq \#\{w \mid \text{Mark}(x, y, z, w)\} \leq 1$

Exists Exam(x, y, z, w) and (Forall w, w'. Exam(x, y, z) and Exam(x, y, z, w') implies $w=w'$) MULTIPLICITY EXPLICIT FORM

**Exercise 2. Consider the above UML class diagram and the following (partial)
instantiation:**

**1. Check whether the above instantiation, once completed, is correct, and explain
why it is or it is not.**

**The above instantiation is not correct, it does not contain Student table. In order
to fix the instantiation, we must add the Student table and insert all the instances
of BScStudent and MScStudent, because there is an ISA relation between these
classes. The resulting table is the following:**

Student:= {sb1, sb2, sm1, sm2, sm3}

All other constraints are not violated. The instantiation is now complete.

2. Express in FOL the following queries and evaluate them over the completed instantiation:

Exercise 3.

1. Model check the Mu-Calculus formula $\nu X. \mu Y. ((a \wedge [\text{next}]X) \vee \langle \text{next} \rangle Y)$

2. Model check (by translating in Mu-Calculus) the CTL formula: $EG(a \wedge AFa)$

$\Phi = \nu X. \mu Y. ((a \wedge [\text{next}]X) \vee \langle \text{next} \rangle Y)$

$[[X_0]] = \{0, 1, 2, 3, 4\}$

$[[X_1]] = [[\mu Y. ((a \wedge [\text{next}]X_0) \vee \langle \text{next} \rangle Y)]] =$

$[[Y_0]] = \{ \}$

$[[Y_1]] = [[(a \wedge [\text{next}]X_0) \vee \langle \text{next} \rangle Y_0]] =$

$[[a]] \text{ inter } \text{PreA}(\text{next}, [[X_0]]) \cup \text{PreE}(\text{next}, [[Y_0]]) =$

$\{0, 1, 4\} \text{ inter } \{1, 2, 3, 4\} \cup \{ \} = \{1, 4\}$

$[[Y_2]] = [[(a \wedge [\text{next}]X_0) \vee \langle \text{next} \rangle Y_1]] =$

$[[a]] \text{ inter } \text{PreA}(\text{next}, [[X_0]]) \cup \text{PreE}(\text{next}, [[Y_1]]) =$

$\{0, 1, 4\} \text{ inter } \{1, 2, 3, 4\} \cup \{0, 3\} = \{0, 1, 3, 4\}$

$[[Y_3]] = [[(a \wedge [\text{next}]X_0) \vee \langle \text{next} \rangle Y_2]] =$

$[[a]] \text{ inter } \text{PreA}(\text{next}, [[X_0]]) \cup \text{PreE}(\text{next}, [[Y_2]]) =$

$\{0, 1, 4\} \text{ inter } \{1, 2, 3, 4\} \cup \{0, 3, 4\} = \{0, 1, 3, 4\}$

Found a LFP $\rightarrow [[Y_2]] = [[Y_3]] = \{0, 1, 3, 4\}$

$[[X_2]] = [[\mu Y. ((a \wedge [\text{next}]X_1) \vee \langle \text{next} \rangle Y)]] =$

$[[Y_{00}]] = \{ \}$

$[[Y_{01}]] = [[(a \wedge [\text{next}]X_1) \vee \langle \text{next} \rangle Y_{00}]] =$

$[[a]] \text{ inter } \text{PreA}(\text{next}, [[X_1]]) \cup \text{PreE}(\text{next}, [[Y_{00}]]) =$

$\{0, 1, 4\} \text{ inter } \{3, 4\} \cup \{ \} = \{4\}$

$$\begin{aligned}
[|Y_{02}|] &= [| (a \wedge [next]X_1) \vee \langle next \rangle Y_{01}) |] = \\
& [|a|] \text{ inter PreA(next, [|X_1|]) } \cup \text{PreE(next, [|Y_{01}|])} = \\
&\{0, 1, 4\} \text{ inter } \{3, 4\} \cup \{3\} = \{3, 4\}
\end{aligned}$$

$$\begin{aligned}
[|Y_{03}|] &= [| (a \wedge [next]X_1) \vee \langle next \rangle Y_{02}) |] = \\
& [|a|] \text{ inter PreA(next, [|X_1|]) } \cup \text{PreE(next, [|Y_{02}|])} = \\
&\{0, 1, 4\} \text{ inter } \{3, 4\} \cup \{0, 3\} = \{0, 3, 4\}
\end{aligned}$$

$$\begin{aligned}
[|Y_{04}|] &= [| (a \wedge [next]X_1) \vee \langle next \rangle Y_{03}) |] = \\
& [|a|] \text{ inter PreA(next, [|X_1|]) } \cup \text{PreE(next, [|Y_{03}|])} = \\
&\{0, 1, 4\} \text{ inter } \{3, 4\} \cup \{0, 3, 4\} = \{0, 3, 4\}
\end{aligned}$$

$$\text{Found a LFP} \rightarrow [|Y_{03}|] = [|Y_{04}|] = \{0, 3, 4\}$$

$$\begin{aligned}
[|X_3|] &= [| \mu Y. ((a \wedge [next]X_2) \vee \langle next \rangle Y) |] = \\
& [|Y_{10}|] = \{ \} \\
[|Y_{11}|] &= [| (a \wedge [next]X_2) \vee \langle next \rangle Y_{11}) |] = \\
& [|a|] \text{ inter PreA(next, [|X_2|]) } \cup \text{PreE(next, [|Y_{11}|])} = \\
&\{0, 1, 4\} \text{ inter } \{3, 4\} \cup \{ \} = \{4\}
\end{aligned}$$

$$\begin{aligned}
[|Y_{12}|] &= [| (a \wedge [next]X_2) \vee \langle next \rangle Y_{12}) |] = \\
& [|a|] \text{ inter PreA(next, [|X_2|]) } \cup \text{PreE(next, [|Y_{12}|])} = \\
&\{0, 1, 4\} \text{ inter } \{3, 4\} \cup \{3\} = \{3, 4\}
\end{aligned}$$

$$\begin{aligned}
[|Y_{13}|] &= [| (a \wedge [next]X_2) \vee \langle next \rangle Y_{13}) |] = \\
& [|a|] \text{ inter PreA(next, [|X_2|]) } \cup \text{PreE(next, [|Y_{13}|])} = \\
&\{0, 1, 4\} \text{ inter } \{3, 4\} \cup \{0, 3\} = \{0, 3, 4\}
\end{aligned}$$

$$\begin{aligned}
[|Y_{14}|] &= [| (a \wedge [next]X_2) \vee \langle next \rangle Y_{14}) |] = \\
& [|a|] \text{ inter PreA(next, [|X_2|]) } \cup \text{PreE(next, [|Y_{14}|])} = \\
&\{0, 1, 4\} \text{ inter } \{3, 4\} \cup \{0, 3, 4\} = \{0, 3, 4\}
\end{aligned}$$

$$\text{Found LFP } [|X_2|] = [|x_3|] = \{0, 3, 4\}$$

$$\Phi = \{ 0, 3, 4 \}$$

1 in $[[\Phi]]$? No, initial state of transition system is not present in the extension of Φ , Hence the formula is False in this transition system.

Decompose the CTL formula: $EG(a \wedge AFa)$

$$\text{Alpha} = AFa$$

$$\text{Beta} = a \wedge \text{alpha}$$

$$\text{Gamma} = EG(\text{Beta})$$

Translation of CTL formula:

$$T(\text{Alpha}) = \mu X. a \vee [\text{next}]X$$

$$T(\text{Beta}) = a \wedge T(\text{Alpha})$$

$$T(\text{Gamma}) = \nu X. T(\text{Beta}) \wedge \langle \text{Next} \rangle X$$

$$[[\text{Alpha}]] = [[\mu X. a \vee [\text{next}]X]] =$$

$$[[X_0]] = \{ \}$$

$$[[X_1]] = [[a \vee [\text{next}]X_0]] = [[a]] \cup \text{PreA}(\text{next}, [[X_0]]) = \{0, 1, 4\} \cup \{ \} = \{0, 1, 4\}$$

$$[[X_2]] = [[a \vee [\text{next}]X_1]] = [[a]] \cup \text{PreA}(\text{next}, [[X_1]]) = \{0, 1, 4\} \cup \{3, 4\} = \{0, 1, 3, 4\}$$

$$[[X_3]] = [[a \vee [\text{next}]X_2]] = [[a]] \cup \text{PreA}(\text{next}, [[X_2]]) = \{0, 1, 4\} \cup \{3, 4\} = \{0, 1, 3, 4\}$$

Found a fixpoint $\rightarrow [[X_2]] = [[X_3]] = \{0, 1, 3, 4\}$

$$[[\text{Beta}]] = [[a]] \wedge [[\text{alpha}]] = \{0, 1, 4\} \cap \{0, 1, 3, 4\} = \{0, 1, 4\}$$

$$[[\text{Gamma}]] = [[\nu X. \text{Beta} \wedge \langle \text{Next} \rangle X]] =$$

$$[[X_0]] = \{0, 1, 2, 3, 4\}$$

$$[[X_1]] = [[\text{Beta}]] \cap \text{PreE}(\text{next}, [[X_0]]) = \{0, 1, 4\} \cap \{0, 1, 2, 3, 4\} = \{0, 1, 4\}$$

$$[[X_2]] = [[\text{Beta}]] \cap \text{PreE}(\text{next}, [[X_1]]) = \{0, 1, 4\} \cap \{0, 3, 4\} = \{0, 4\}$$

$$[[X_3]] = [[\text{Beta}]] \cap \text{PreE}(\text{next}, [[X_2]]) = \{0, 1, 4\} \cap \{3, 4\} = \{4\}$$

$$[|X_4|] = [|Beta|] \text{ inter } PreE(next, [|X_3|]) = \\ = \{0, 1, 4\} \text{ inter } \{3\} = \{ \}$$

$$[|X_5|] = [|Beta|] \text{ inter } PreE(next, [|X_4|]) = \\ = \{0, 1, 4\} \text{ inter } \{ \} = \{ \}$$

$$[|Gamma|] = \{ \}$$

CTL Formula = $\{ \}$

Is 1 in $[|Gamma|]$? No, initial state of TS is not present in the extension of Gamma, hence, the formula is not valid in the TS. CTL formula is false in this TS

Exercise 4. Compute the *weakest precondition* for getting $x=y$ executing the following program:

Evaluation semantics:

Computing the weakest precondition, regressing the Post condition Q from the bottom to the top of the execution of program:

$$\{P\} S \{Q\} \text{ iff } P \Rightarrow WP(S, Q)$$

Note, Hoare logic is concerned only about partial correctness, it does not consider termination. If termination and post condition is achieved, producing the right result, then total correctness is achieved. Termination in this case is irrelevant, because there are no while loops.

$$WP(\Delta, Q) \rightarrow WP(\Delta, \{x=y\}) \Rightarrow \{y=5\}$$

$$\{x=y\} \rightarrow \\ \{y:=5;\}$$

$X:=10;$

$$\{ (y > 10 \text{ and } x+5=0) \text{ OR } (y \leq 10 \text{ and } x=2y) \} \rightarrow \\ \rightarrow \{ (y > 10 \text{ and } 10+5 = 0) \text{ OR } (y \leq 10 \text{ and } 10=2y) \} \\ \rightarrow \{ (y > 10 \text{ and } 15=0) \text{ OR } (y \leq 10 \text{ and } y=5) \} \\ \rightarrow \text{FALSE OR TRUE}$$

If $(y > 10)$ then {

$$\{x=y\} \rightarrow \{x + y + 5 = y\} \rightarrow \{x + 5 = y - y\} \rightarrow \{x + 5 = 0\} \text{ FALSE}$$

$$x = x+y;$$

$$\{x=y\} \rightarrow \{x = y-5\} \rightarrow \{y = x + 5\}$$

$$y = y-5;$$

$\{x=y\}$

}

$\{x=y\} \rightarrow \{x-y=y\} \rightarrow \{x = 2y\}$

else $x = x-y$;

$\{x=y\}$

Alpha implies not (Beta)