TEST 1: Hash Functions

Let S_3 be the set of permutations on the set $\{1,2,3\}$. For each $\pi \in S_3$ let e_{π} be the corresponding bit permutation on B_3 . For each $\pi \in S_3$, determine the number of collisions of the compression function $h_{\pi}(x) = e_{\pi}(x) \oplus x$ where $x \in B_3$.

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- $S_3 = \{123, 132, 213, 231, 312, 321\}$
- $| S_3 | = 3! = 1*2*3 = 6$

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- $S_3 = \{123, 132, 213, 231, 312, 321\}$
- $| S_3 | = 3! = 1*2*3 = 6$
- $B_3 = \{101, 110, 111, 100, 000, 010, 011, 001\}$
- $|B_3| = 2^3 = 8$

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- $S_3 = \{123, 132, 213, 231, 312, 321\}$
- $B_3 = \{101, 110, 111, 100, 000, 010, 011, 001\}$
- $\pi = 213$ and x = 101, $e_{\pi}(x) = e_{(213)}(101) = 011$
- $h_{\pi}(x) = e_{\pi}(x) \oplus x$: $e_{(213)}(101) \oplus 101 = 011 \oplus 101 = 110$

For each $\pi \in S_3$, determine the number of collisions of the compression function $h_{\pi}(x) = e_{\pi}(x) \oplus x$ where $x \in B_3$.

•
$$B_3 = \{101, 110, 111, 100, 000, 010, 011, 001\}, \pi = (123)$$

$$h_{(123)}(x) = e_{(123)}(x) \oplus x$$

•
$$h_{(123)}(101) = e_{(123)}(101) \oplus 101 = 101 \oplus 101 = 000$$

For each $\pi \in S_3$, determine the number of collisions of the compression function $h_{\pi}(x) = e_{\pi}(x) \oplus x$ where $x \in B_3$.

• $B_3 = \{101, 110, 111, 100, 000, 010, 011, 001\}, \pi = (123)$

$$h_{(123)}(x) = e_{(123)}(x) \oplus x$$

- $h_{(123)}(101) = e_{(123)}(101) \oplus 101 = 101 \oplus 101 = 000$
- $h_{(123)}(110) = e_{(123)}(110) \oplus 110 = 110 \oplus 110 = 000$

For each $\pi \in S_3$, determine the number of collisions of the compression function $h_{\pi}(x) = e_{\pi}(x) \oplus x$ where $x \in B_3$.

• $B_3 = \{101, 110, 111, 100, 000, 010, 011, 001\}, \pi = (123)$

$$h_{(123)}(x) = e_{(123)}(x) \oplus x$$

- $h_{(123)}(101) = e_{(123)}(101) \oplus 101 = (101) \oplus (101) = 000$
- $h_{(123)}(110) = e_{(123)}(110) \oplus 110 = (110) \oplus (110) = 000$
- ...
- $h_{(123)}(001) = e_{(123)}(001) \oplus 001 = (001) \oplus (001) = 000$

$$B_3 = \{101, 110, 111, 100, 000, 010, 011, 001\}, \pi = (123)$$

- $h_{(123)}(101) = e_{(123)}(101) \oplus 101 = 101 \oplus 101 = 000$
- $h_{(123)}(110) = e_{(123)}(110) \oplus 110 = 110 \oplus 110 = 000$
- ...
- $h_{(123)}(001) = e_{(123)}(001) \oplus 001 = 001 \oplus 001 = 000$

The number of collisions:

- A. 8
- B. 12
- C. 28

$$B_3 = \{101, 110, 111, 100, 000, 010, 011, 001\}, \pi = (123)$$

- $h_{(123)}(101) = e_{(123)}(101) \oplus 101 = (101) \oplus (101) = 000$
- $h_{(123)}(110) = e_{(123)}(110) \oplus 110 = (110) \oplus (110) = 000$
- ...
- $h_{(123)}(001) = e_{(123)}(001) \oplus 001 = (001) \oplus (001) = 000$

The number of collisions: $C^2_8 = (7*8)/2 = 28$

$$B_3 = \{101, 110, 111, 100, 000, 010, 011, 001\}, \pi = (132)$$

- $h_{(132)}(101) = e_{(132)}(101) \oplus 101 = 110 \oplus 101 = 011$
- $h_{(132)}(110) = e_{(132)}(110) \oplus 110 = 101 \oplus 101 = 011$
- $h_{(132)}(111) = e_{(132)}(111) \oplus 111 = 111 \oplus 111 = 000$
- $h_{(132)}(100) = e_{(132)}(100) \oplus 100 = 100 \oplus 100 = 000$
- $h_{(132)}(000) = e_{(132)}(000) \oplus 000 = 000 \oplus 000 = 000$
- $h_{(132)}(010) = e_{(132)}(010) \oplus 010 = 001 \oplus 010 = 011$
- $h_{(132)}(011) = e_{(132)}(011) \oplus 011 = 011 \oplus 011 = 000$
- $h_{(132)}(001) = e_{(132)}(001) \oplus 010 = 010 \oplus 101 = 011$

$$B_3 = \{101, 110, 111, 100, 000, 010, 011, 001\}, \pi = (132)$$

- $h_{(132)}(101) = e_{(132)}(101) \oplus 101 = 110 \oplus 101 = 011$
- $h_{(132)}(110) = e_{(132)}(110) \oplus 110 = 101 \oplus 101 = 011$
- $h_{(132)}(111) = e_{(132)}(111) \oplus 111 = 111 \oplus 111 = 000$
- $h_{(132)}(100) = e_{(132)}(100) \oplus 100 = 100 \oplus 100 = 000$
- $h_{(132)}(000) = e_{(132)}(000) \oplus 000 = 000 \oplus 000 = 000$
- $h_{(132)}(010) = e_{(132)}(010) \oplus 010 = 001 \oplus 010 = 011$
- $h_{(132)}(011) = e_{(132)}(011) \oplus 011 = 011 \oplus 011 = 000$
- $h_{(132)}(001) = e_{(132)}(001) \oplus 010 = 010 \oplus 101 = 011$

The number of collisions: $C^2_4 + C^2_4 = 12$

$$S_3 = \{123, 132, 213, 231, 312, 321\}$$

 $B_3 = \{101, 110, 111, 100, 000, 010, 011, 001\}$

- $\pi = 213$: 12
- $\pi = 231$: $4*C^2_2 = 4$
- $\pi = 312: 4$
- $\pi = 321$: 12

Consider the hash function h: $B^* \to B^*$ given by $k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor$ where a bitstring k is identified with the natural number represented by k and r $\mod 1 = r - \lfloor r \rfloor$, for a positive real number r. Moreover, images are padded by leading zeroes to the maximal possible length of all images of h.

- Determine the maximal length in bits of the images.
- Give a collision of this hash function.

$$k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor$$

r mod $1 = r - \lfloor r \rfloor$

Example k = 1

$$h(1) = \lfloor 10000((1(1+\sqrt{5})/2) \mod 1) \rfloor$$

```
\begin{array}{l} \mathsf{k} \mapsto \lfloor 10000((\mathit{k}(1+\sqrt{5})/2) \mod 1) \rfloor \\ \mathsf{r} \mod 1 = \mathsf{r} - \lfloor \mathsf{r} \rfloor \end{array}
```

Example k = 1

```
\begin{array}{l} \mathsf{h}(1) = \lfloor 10000((1(1+\sqrt{5})/2) \mod 1) \rfloor \\ = \lfloor 10000(((1+\sqrt{5})/2) \mod 1) \rfloor \\ = \lfloor 10000(1.61803398875 \mod 1) \rfloor \end{array}
```

```
\begin{array}{l} \mathsf{k} \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor \\ \mathsf{r} \mod 1 = \mathsf{r} - \lfloor \mathsf{r} \rfloor \\ \\ \mathsf{Example} \ \mathsf{k} = 1 \\ \mathsf{h}(1) = \lfloor 10000((1(1+\sqrt{5})/2) \mod 1) \rfloor \\ = \lfloor 10000(((1+\sqrt{5})/2) \mod 1) \rfloor \\ = \lfloor 10000(1.61803398875 \mod 1) \rfloor \\ = \lfloor 10000(1.61803398875 - \rfloor 1.61803398875 \rfloor) \  \, \vert \end{array}
```

```
k \mapsto |10000((k(1+\sqrt{5})/2) \mod 1)|
   mod 1 = r - |r|
Example k = 1
```

```
h(1) = |10000((1(1+\sqrt{5})/2) \mod 1)|
= |10000(((1+\sqrt{5})/2) \mod 1)|
= |10000(1.61803398875 \mod 1)|
= |10000(1.61803398875 - |1.61803398875 |)|
= |10000(1.61803398875 - 1)|
```

```
\begin{array}{l} \mathsf{k} \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor \\ \mathsf{r} \mod 1 = \mathsf{r} - \lfloor \mathsf{r} \rfloor \\ \\ \mathsf{Example} \ \mathsf{k} = 1 \\ \mathsf{h}(1) = \lfloor 10000((1(1+\sqrt{5})/2) \mod 1) \rfloor \\ = \lfloor 10000((1+\sqrt{5})/2) \mod 1) \rfloor \\ = \lfloor 10000(1.61803398875 \mod 1) \rfloor \\ = \lfloor 10000(1.61803398875 - \lfloor 1.61803398875 \rfloor) \rfloor \\ = \lfloor 10000(1.61803398875 - 1) \rfloor \\ = \lfloor 10000(0.61803398875) \rfloor \end{array}
```

```
\mathsf{k} \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor
r mod 1 = \mathsf{r} - \lfloor \mathsf{r} \rfloor
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Example k = 1

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\begin{array}{l} \mathsf{h}(1) = \lfloor 10000((1(1+\sqrt{5})/2) \mod 1) \rfloor \\ = \lfloor 10000(((1+\sqrt{5})/2) \mod 1) \rfloor \\ = \lfloor 10000(1.61803398875 \mod 1) \rfloor \\ = \lfloor 10000(1.61803398875 - \lfloor 1.61803398875 \rfloor) \rfloor \\ = \lfloor 10000(1.61803398875 - 1) \rfloor \\ = \lfloor 10000(0.61803398875) \rfloor \\ = \lfloor 6180.3398875 \rfloor = 6180 \end{array}
```

```
k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor
r mod 1 = r - \lfloor r \rfloor
```

Example k = 1

```
\begin{array}{l} \mathsf{h}(1) = \lfloor 10000((1(1+\sqrt{5})/2) \mod 1) \rfloor \\ = \lfloor 10000(((1+\sqrt{5})/2) \mod 1) \rfloor \\ = \lfloor 10000(1.61803398875 \mod 1) \rfloor \\ = \lfloor 10000(1.61803398875 - \lfloor 1.61803398875 \rfloor) \rfloor \\ = \lfloor 10000(1.61803398875 - 1) \rfloor \\ = \lfloor 10000(0.61803398875) \rfloor \\ = \lfloor 6180.3398875 \rfloor = 6180 \end{array}
```

$$k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor$$

r mod $1 = r - |r|$

• Determine the maximal length in bits of the images.

$$k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor$$

r mod $1 = r - |r|$

• Determine the maximal length in bits of the images.

k - preimages and h(k) - images

$$\begin{array}{ll} \mathsf{k} \mapsto \lfloor 10000((\mathit{k}(1+\sqrt{5})/2) \mod 1) \rfloor \\ \mathsf{r} \mod 1 = \mathsf{r} - \mid \mathsf{r} \mid \end{array}$$

Determine the maximal length in bits of the images.

The maximal length in bits of the images \rightarrow The maximal of h(k)

$$k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor$$

r mod $1 = r - |r|$

Determine the maximal length in bits of the images.

$$r \mod 1 = \leq r - |r|$$

$$k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor$$

r mod $1 = r - |r|$

Determine the maximal length in bits of the images.

$$\begin{array}{ll} r & \text{mod} & 1 = \leq r \text{-} \lfloor r \rfloor \\ 0 \leq r \text{-} \mid r \mid \leq 0.9999999... \end{array}$$

$$k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor$$

r mod $1 = r - |r|$

• Determine the maximal length in bits of the images.

$$0 \leq r \mod 1 \leq 0.9999999...$$

$$\lfloor\ 10000((\mathsf{k}\ (1+\sqrt{5})/2)\ \mathsf{mod}\ 1)\rfloor$$

$$\begin{array}{ll} \mathsf{k} \mapsto \lfloor 10000((\mathit{k}(1+\sqrt{5})/2) \mod 1) \rfloor \\ \mathsf{r} \mod 1 = \mathsf{r} - \lfloor \mathsf{r} \rfloor \end{array}$$

Determine the maximal length in bits of the images.

 $0 \leq r \mod 1 \leq 0.9999999...$

```
 \lfloor 10000((\mathsf{k}\ (1+\sqrt{5})/2)\ \mathsf{mod}\ 1)\rfloor \\ \lfloor 10000(0)\rfloor \leq \lfloor 10000((\mathsf{k}\ (1+\sqrt{5})/2)\ \mathsf{mod}\ 1)\rfloor \leq \lfloor 10000(0.99999999...)\rfloor
```

```
k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor
r \mod 1 = r - \lfloor r \rfloor
```

Determine the maximal length in bits of the images.

 $0 \leq r \mod 1 \leq 0.9999999...$

```
\begin{array}{l} \mathsf{k} \mapsto \lfloor 10000 ((\mathit{k}(1+\sqrt{5})/2) \mod 1) \rfloor \\ \mathsf{r} \mod 1 = \mathsf{r} - \lfloor \mathsf{r} \rfloor \end{array}
```

Determine the maximal length in bits of the images.

 $0 \leq r \mod 1 \leq 0.9999999...$

```
\begin{array}{l} \mathsf{k} \mapsto \lfloor 10000((\mathit{k}(1+\sqrt{5})/2) \mod 1) \rfloor \\ \mathsf{r} \mod 1 = \mathsf{r} - \lfloor \mathsf{r} \rfloor \end{array}
```

• Determine the maximal length in bits of the images.

```
\begin{array}{ll} r \mod 1 = \leq r - \lfloor \, r \, \rfloor \\ \\ 0 \leq r - \lfloor \, r \, \rfloor \leq 0.99999999... \ \to 0 \leq r \mod 1 \leq 0.99999999... \end{array}
```

```
 \begin{array}{l} \lfloor \ 10000((\mathsf{k}\ (1+\sqrt{5})/2)\ \ \mathsf{mod}\ 1) \rfloor \\ \lfloor 10000(0) \rfloor \leq \lfloor 10000((\mathsf{k}\ (1+\sqrt{5})/2)\ \ \mathsf{mod}\ 1) \rfloor \leq \lfloor \\ 10000(0.99999999...) \rfloor \\ \lfloor \ 0 \ \rfloor \leq \lfloor 10000((\mathit{k}(1+\sqrt{5})/2)\ \ \mathsf{mod}\ 1) \rfloor \leq \lfloor 9999.9999... \rfloor \\ 0 \leq \lfloor 10000((\mathit{k}(1+\sqrt{5})/2)\ \ \mathsf{mod}\ 1) \rfloor \leq 9999 \end{array}
```

The maximum of h(k) = 9999

$$k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor$$

r mod $1 = r - |r|$

• Determine the maximal length in bits of the images.

$$0 \le \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor \le 9999$$

The maximum of $h(k) = 9999 \rightarrow$ the maximal length in bits ?

- A. 4
- B. 8
- C. 14

$$\begin{array}{l} \mathsf{k} \mapsto \lfloor 10000((\mathit{k}(1+\sqrt{5})/2) \mod 1) \rfloor \\ \mathsf{r} \mod 1 = \mathsf{r} - \mid \mathsf{r} \mid \end{array}$$

• Determine the maximal length in bits of the images.

$$0 \leq \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor \leq 9999$$

The maximum of h(k) = 9999 \rightarrow the maximal length in bits ? 9999 \cong $2^{13.287568103} \rightarrow$ $2^{13} <$ h(k) < $2^{14} \rightarrow$ the maximal length in bits = 14

$$k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor$$

$$r \mod 1 = r - |r|$$

• Give a collision of this hash function.

$$k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor$$

$$r \mod 1 = r - \mid r \mid$$

• Give a collision of this hash function.

To find k_1 and k_2 such as $\lfloor 10000((k_1(1+\sqrt{5})/2) \mod 1) \rfloor = \lfloor 10000((k_2(1+\sqrt{5})/2) \mod 1) \rfloor$

$$k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor$$

r mod $1 = r - |r|$

• Give a collision of this hash function.

Do you know this number?

$$(1+\sqrt{5})/2$$

- A. yes
- B. no

Golden Ratio

In mathematics, two quantities are in the **golden ratio** if their ratio is the same as the ratio of their sum to the larger of the two quantities. The figure on the right illustrates the geometric relationship. Expressed algebraically, for quantities a and b with a > b > 0,

$$rac{a+b}{a}=rac{a}{b}\stackrel{ ext{def}}{=}arphi,$$

where the Greek letter phi $(\varphi \circ \phi)$ represents the golden ratio.^(a) It is an irrational number that is a solution to the quadratic equation $x^2-x-1=0$, with a value of:

$$arphi = rac{1+\sqrt{5}}{2} = 1.6180339887\dots$$
[1]





Fibonacci numbers

In mathematics, the **Fibonacci numbers**, commonly denoted F_n , form a sequence, called the **Fibonacci sequence**, such that each number is the sum of the two preceding ones, starting from 0 and 1. That is,^[1]

$$F_0 = 0$$
, $F_1 = 1$.

and

$$F_n = F_{n-1} + F_{n-2}$$

for n > 1.

The beginning of the sequence is thus:

$$0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, ...$$

Fibonacci numbers

$$\lim_{n o\infty}rac{F_{n+1}}{F_n}=arphi.$$

Fibonacci numbers

	Α	В	c
		Fibonacci	Ratio of Adjacent
1	n	Numbers	Terms
2	0	1	-
3	1	1	1
4	2	2	2
5	3	2	1.5
6	4	5	1.666666667
7	- 5	8	1.6
8	6	13	1.625
9	- 7	21	1.615384615
10	8	34	1.619047619
11	9	55	1.617647059
12	10	89	1.618181818
13	11	144	1.617977528
14	12	233	1.618055556
15	13	377	1.618025751
16	14	610	1.618037135
17	15	987	1.618032787
18	16	1597	1.618034448
19	17	2584	1.618033813
20	18	4181	1.618034056
21	19	6765	1.618033963
22	20	10946	1.618033999
23	21	17711	1.618033985
24	22	28657	1.61803399
25	23	46368	1.618033988

$$k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor$$

$$r \mod 1 = r - \mid r \mid$$

• Give a collision of this hash function.

To find k_1 and k_2 such as $\lfloor 10000((k_1(1+\sqrt{5})/2) \mod 1) \rfloor = \lfloor 10000((k_2(1+\sqrt{5})/2) \mod 1) \rfloor$

$$k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor$$

$$r \mod 1 = r - |r|$$

• Give a collision of this hash function.

To find k_1 and k_2 such as $\lfloor 10000((k_1(1+\sqrt{5})/2) \mod 1) \rfloor = \lfloor 10000((k_2(1+\sqrt{5})/2) \mod 1) \rfloor$

$$k \mapsto \lfloor 10000((k(1+\sqrt{5})/2) \mod 1) \rfloor$$

$$r \mod 1 = r - |r|$$

• Give a collision of this hash function.

To find k_1 and k_2 such as $\lfloor 10000((k_1(1+\sqrt{5})/2) \mod 1) \rfloor = \lfloor 10000((k_2(1+\sqrt{5})/2) \mod 1) \rfloor$ Considering four numbers after decimal

• Give a collision of this hash function.

```
Choose k_1 = 1 \lfloor 10000((1+\sqrt{5})/2) \mod 1) \rfloor = \lfloor 10000((k_2(1+\sqrt{5})/2) \mod 1) \rfloor k_2?
```

```
\frac{a}{b} and \frac{a}{b} *(b + 1) = a + \frac{a}{b} If a is a nature number, \frac{a}{b} - \lfloor \frac{a}{b} \rfloor = a + \frac{a}{b} - \lfloor a + \frac{a}{b} \rfloor ( \frac{a}{b} mod 1 = (a + \frac{a}{b}) mod 1)
```

Fibonacci numbers

	Α	В	c l
		Fibonacci	Ratio of Adjacent
1	n	Numbers	Terms
2	0	1	-
3	1	1	1
4	2	2	2
5	3	3	1.5
6	4	5	1.666666667
7	- 5	8	1.6
8	6	13	1.625
9	- 7	21	1.615384615
10	8	34	1.619047619
11	9	55	1.617647059
12	10	89	1.618181818
13	11	144	1.617977528
14	12	233	1.618055556
15	13	377	1.618025751
16	14	610	1.618037135
17	15	987	1.618032787
18	16	1597	1.618034448
19	17	2584	1.618033813
20	18	4181	1.618034056
21	19	6765	1.618033963
22	20	10946	1.618033999
23	21	17711	1.618033985
24	22	28657	1.61803399
25	23	46368	1.618033988

• Give a collision of this hash function.

Choose
$$k_1 = 1$$
 $\lfloor 10000((1+\sqrt{5})/2) \mod 1) \rfloor = \lfloor 10000((k_2(1+\sqrt{5})/2) \mod 1) \rfloor$

k_2 ?

- $\frac{17711}{10946} \cong \frac{1+\sqrt{5}}{2}$
- $\frac{17711}{10946} \mod 1 = (10946 + 1) * \frac{17711}{10946} \mod 1 = (1 + \frac{17711}{10946}) \mod 1$

• Give a collision of this hash function.

```
Choose k_1 = 1
\lfloor 10000((1+\sqrt{5})/2) \mod 1) \rfloor = \lfloor 10000((k_2(1+\sqrt{5})/2) \mod 1) \rfloor
k_2 \mod 1
\lfloor 10000(10947(1+\sqrt{5})/2) \mod 1) \rfloor = \lfloor 10000(10947\frac{17711}{10946} \mod 1) \rfloor
= \lfloor 10000((10946+1)\frac{17711}{10946} \mod 1) \rfloor = \lfloor 10000((1+\frac{17711}{10946}) \mod 1) \rfloor
= \lfloor 10000(\frac{17711}{10946} \mod 1) \rfloor
\cong \lfloor 10000((1+\sqrt{5})/2 \mod 1) \rfloor
```

$$\begin{array}{ll} \mathsf{k} \mapsto \lfloor 10000((\mathit{k}(1+\sqrt{5})/2) \mod 1) \rfloor \\ \mathsf{r} \mod 1 = \mathsf{r} - \mid \mathsf{r} \mid \end{array}$$

• Give a collision of this hash function.

a collision
$$(k_1, k_2) = (1, 10947)$$

Explain the construction of a hash function from a compression function for the concrete case of r=1.

Merkle-Damgaard procedure

m = n + r and $B^* \rightarrow B^m$

Let $f: B^m \to B^n$ be a compression function and let $r = m - n \ge 2$.

The goal is to construct a hash function $h: B^* \to B^n$ from f.

Merkle-Damgaard procedure

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m = n + r and $B^* \rightarrow B^m$

Preprocessing Step 1

Given $x \in B^*$, prepend the minimal number $0 \le k < r$ of zeroes such that the new length is a multiple of r and append 0^r . Result: $x' = 0^k \|x\| 0^r$

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Preprocessing Step 2

Calculate the binary representation b of the original length of x and prepend zeroes such that its length is divisible by r-1. Starting at the beginning insert 1 at every r-1st position of the resulting string. The length of the resulting string b' is a multiple of r.

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$$m = n + r$$
 and $B^* \rightarrow B^m$

Preprocessing Step 1

Given $x \in B^*$, prepend the minimal number $0 \le k < r$ of zeroes such that the new length is a multiple of r and append 0^r . Result: $x' = 0^k ||x|| 0^r$

Preprocessing Step 2

Calculate the binary representation b of the original length of x and prepend zeroes such that its length is divisible by r-1. Starting at the beginning insert 1 at every r-1st position of the resulting string. The length of the resulting string b' is a multiple of r.

Preprocessing Step 3

Prepend b' to obtain a string $b'||0^k||x||0^r$ of length $t \cdot r$. Decompose into

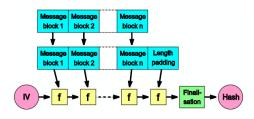
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Constructing the hash function

Definition

Define $h(x) = H_t$ where

- $x_1 || x_2 || \dots || x_t$ with $x_i \in B^r$ is the result of preprocessing $x \in B^*$.
- $H_0 = 0^n$ (or a different, but fixed initialization vector).
- $H_i = f(H_{i-1}||x_i)$ for $1 \le i \le t$.



Preprocessing Step 1

Given $x \in B^*$, prepend the minimal number $0 \le k < r$ of zeroes such that the new length is a multiple of r and append 0^r . Result: $x' = 0^k ||x|| 0^r$

x || 0

Preprocessing Step 2

Calculate the binary representation b of the original length of x and prepend zeroes such that its length is divisible by r-1. Starting at the beginning insert 1 at every r-1st position of the resulting string. The length of the resulting string b' is a multiple of r.

Preprocessing Step 3

Prepend b' to obtain a string $b'\|0^k\|x\|0^r$ of length $t \cdot r$. Decompose into $x_1\|x_2\|\dots\|x_t$ with $x_i \in B^r$.

b $|| \times || 0$. Decompose into $x_1 || x_2 || \dots || x_t$ with $x_i \in B^1$. t = len(x) + len(b) + 1

Preprocessing Step 3

Prepend b' to obtain a string $b'\|0^k\|x\|0^r$ of length $t \cdot r$. Decompose into $x_1\|x_2\|\dots\|x_t$ with $x_i \in B^r$.

b || x || 0. Decompose into $x_1 || x_2 || \dots || x_t$ with $x_i \in B^1$. t = len(x) + len(b) + 1

Constructing the hash function

Definition

Define $h(x) = H_t$ where

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