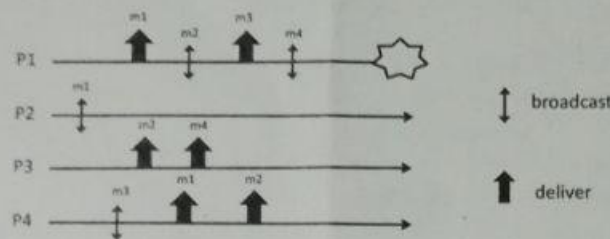


Distributed Systems 23/01/2019
Corso di Laurea Magistrale in Ingegneria Informatica

Family Name _____ Name _____ Student ID _____

Ex 1: Consider the execution depicted in the Figure



Answer to the following questions:

1. Provide all the delivery sequences that satisfy both causal order and total order
2. Complete the execution in order to have a run satisfying TO(NUA, WNUTO), FIFO order but not causal order

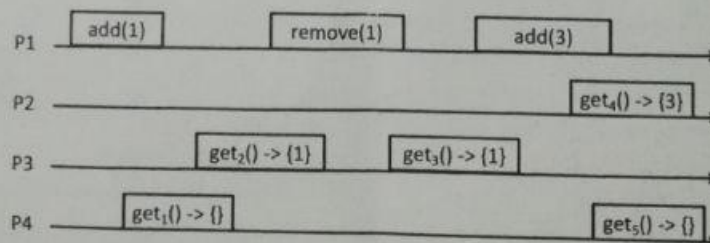
Ex 2: Consider a set object that can be accessed by a set of processes. Processes may invoke the following operations on the object:

- `add(v)`: it adds the value `v` in to the set
- `remove(v)` it removes the value `v` from the set
- `get()`: it returns the content of the set.

Informally, every `get()` operation returns all the values that have been added before its invocation and that have not been removed by any `remove()`.

For the sake of simplicity, assume that a value can be added/removed just once in the execution.

Consider the distributed execution depicted in the Figure



Answer to the following questions:

1. Is the proposed execution linearizable? Motivate your answer with examples.
2. Consider now the following property: "every `get()` operation returns all the values that have been added before its invocation and that have not been removed by any `remove()`. If an `add(v)/remove(v)` operation is concurrent with the `get`, the value `v` may or may be not returned by the `get()`".
 Provide an execution that satisfy get validity and that is not linerizable.

Ex 3: Consider a distributed system composed of N processes p_1, p_2, \dots, p_N , each having a unique identifier $myID$. Initially, all processes are correct (i.e. $correct = \{p_1, p_2, \dots, p_N\}$). Consider the following algorithm:

```

upon event  $xbroadcast(m)$ 
     $mynsn = mysn + 1;$ 
     $\forall p \in correct$ 
         $pp2pSend("MSG", m, mysn, myID);$ 

upon event  $pp2pReceive("MSG", m, sn, i)$ 
     $mynsn = mysn + 1;$ 
    if ( $m \notin delivered$ )
        trigger  $xDeliver(m);$ 
         $delivered = delivered \cup \{m\};$ 

upon event  $crash(p_i)$ 
     $correct = correct / \{p_i\}$ 
    
```

Let us assume that: (i) links are perfect, (ii) the failure detector is perfect and (iii) initially local variables are initialized as follows $mynsn=0$ e $delivered = \emptyset$.

Answer to the following questions:

1. Does the $xbroadcast()$ primitive implement a Reliable Broadcast, a Best Effort Broadcast or none of the two?
2. Considering only the ordering property of broadcast communication primitives discussed during the lectures (FIFO, Causal, Total), explain which ones can be satisfied by the $xbroadcast()$ implementation.

Provide examples to justify your answers.

Ex 4: Consider a distributed system constituted by n processes $\Pi = \{p_1, p_2, \dots, p_n\}$ with unique identifiers that exchange messages through FIFO perfect point-to-point links and are structured through a line (i.e., each process p_i can exchange messages only with processes p_{i-1} and p_{i+1} when they exist). Processes may crash and each process is equipped with a perfect oracle (having the interface $new_right(p)$ and $new_left(p)$) reporting a new neighbor when the previous one is failing. Write the pseudo-code of an algorithm implementing a Perfect failure detector primitive.

According to the Italian law 675 of the 31/12/96, I authorize the instructor of the course to publish on the web site of the course results of the exams.

Signature: _____