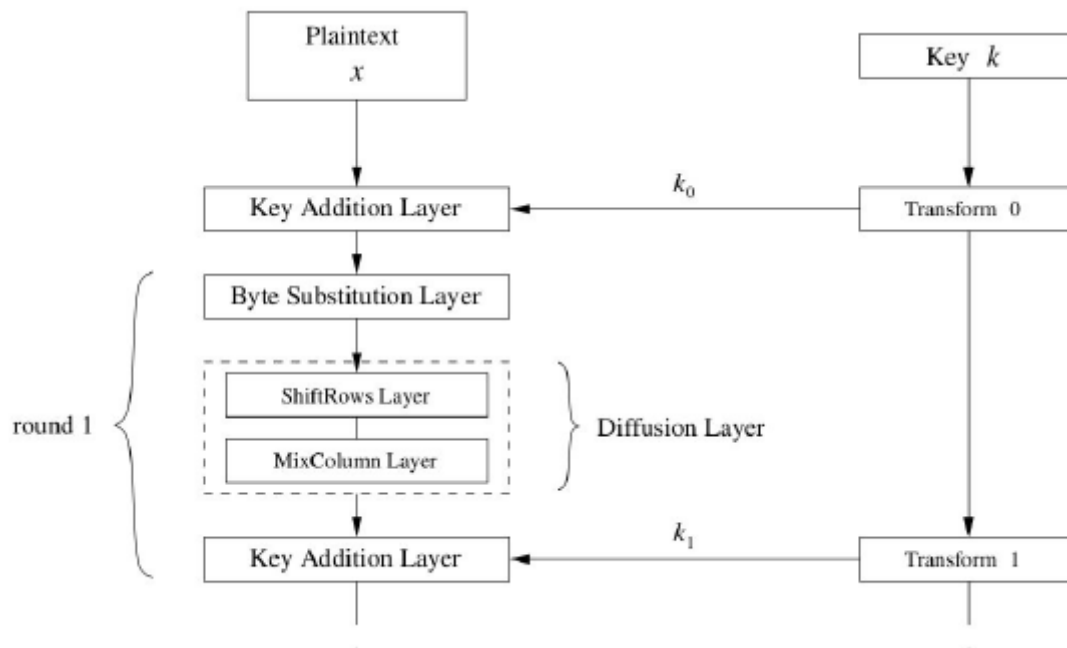


# 3. Symmetric ciphers II (AES)

## 3.1 Standard

Standard for NIST; key: 128/192/256 bytes

input: block of 128 bytes.

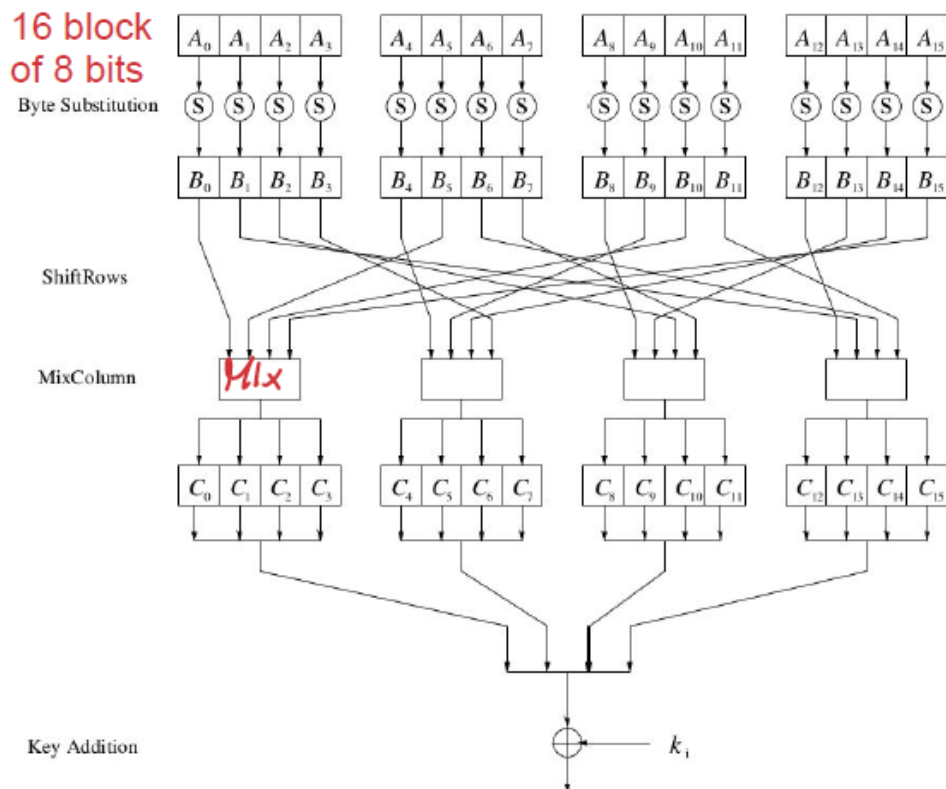


10/12/14 rounds

## 3.2 Layers

In the last round we skip MixColumn Layer

1. ByteSub → CONFUSION
2. ShiftRow → DIFFUSION
3. MixColumn → DIFFUSION
4. Key Addition → KEY WHITENING



### 3.3 Internal structure

- **Byte-oriented** cipher.
- Based on a **substitution-permutation network**
- A ( $A_0, \dots, A_{15}$ , 16-byte (128bit) input) can be draw as a 4x4 matrix

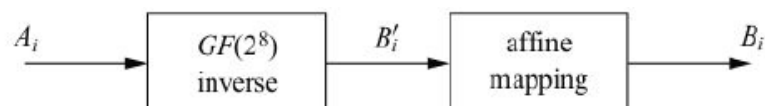
#### 3.3.1 Byte substitution Layer

Independent for every byte.

16 S-Boxes identical, non linear and bijective (S-Box can be uniquely reversed)

**Confusion:** 1 bit in  $A_i$  can affect 3/4 bits in  $B_i$

The S-Box perform two operations:



$$A_i = 1100\ 0010 \Rightarrow A_i(x) = x^7 + x^6 + x$$

The first step computes the inverse (which provides the non linearity in AES):

$$B'_i(x) = A(x)^{-1} \quad P(x) = x^8 + x^4 + x^3 + x + 1$$

such that:  $B'_i(x) \cdot A(x)^{-1} \equiv 1 \pmod{P(x)}$  AES irreducible polynomial

The second step computed in the S-Box is an affine mapping (this is done to destroy some algebraic properties that could be exploited by an attacker):

$$B_i(x) \equiv \begin{pmatrix} b_0 \\ b_1 \\ b_2 \\ b_3 \\ b_4 \\ b_5 \\ b_6 \\ b_7 \end{pmatrix} \equiv \begin{pmatrix} 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 & 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 \end{pmatrix} \begin{pmatrix} b'_0 \\ b'_1 \\ b'_2 \\ b'_3 \\ b'_4 \\ b'_5 \\ b'_6 \\ b'_7 \end{pmatrix} + \begin{pmatrix} 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} \pmod{2}.$$

```
#include <stdint.h>

#define ROTL8(x,shift) (((uint8_t) ((x) << (shift))) | \
#define ((x) >> (8 - (shift))))

void initialize_aes_sbox(uint8_t sbox[256]) {
    uint8_t p = 1, q = 1;

    /* loop invariant: p * q == 1 in the Galois field */
    do {
        /* multiply p by 3 */
        p = p ^ (p << 1) ^ (p & 0x80 ? 0x1B : 0);

        /* divide q by 3 (equals multiplication by 0xf6) */
        q ^= q << 1;
        q ^= q << 2;
        q ^= q << 4;
        q ^= q & 0x80 ? 0x09 : 0;

        /* compute the affine transformation */
        uint8_t xformed = q ^ ROTL8(q, 1) ^ ROTL8(q, 2) ^ ROTL8(q, 3)
            ^ ROTL8(q, 4);

        sbox[p] = xformed ^ 0x63;
    } while (p != 1);

    /* 0 is a special case since it has no inverse */
    sbox[0] = 0x63;
}
```

### 3.3.2 Diffusion Layer

**Diffusion:** given a byte with some bit flips, it will spread the effect on 32 bits from the state.

**ShiftRows Sublayer:** Permutation of the data on a byte level (shift bit on the left cyclically on every row)

**MixColumn Sublayer:** Matrix operation which combines ("mixes") blocks of four bytes performs a linear operation on state matrices A, B, i.e.,

$$\begin{pmatrix} C_0 \\ C_1 \\ C_2 \\ C_3 \end{pmatrix} = \begin{pmatrix} 02 & 03 & 01 & 01 \\ 01 & 02 & 03 & 01 \\ 01 & 01 & 02 & 03 \\ 03 & 01 & 01 & 02 \end{pmatrix} \cdot \begin{pmatrix} B_0 \\ B_5 \\ B_{10} \\ B_{15} \end{pmatrix}$$

$$C_0 = 02 \cdot B_0 + 03 \cdot B_5 + 01 \cdot B_{10} + 01 \cdot B_{15}$$

$$C_0 = x \cdot B_0 + (x+1) \cdot B_5 + 1 \cdot B_{10} + 1 \cdot B_{15}$$

E.g.,  $B = (25, \dots, 25)$

$$\begin{aligned} 02 \cdot 25 &= x \cdot (x^5 + x^2 + 1) \\ &= x^6 + x^3 + x, \\ 03 \cdot 25 &= (x+1) \cdot (x^5 + x^2 + 1) \\ &= (x^6 + x^3 + x) + (x^5 + x^2 + 1) \\ &= x^6 + x^5 + x^3 + x^2 + x + 1. \end{aligned}$$

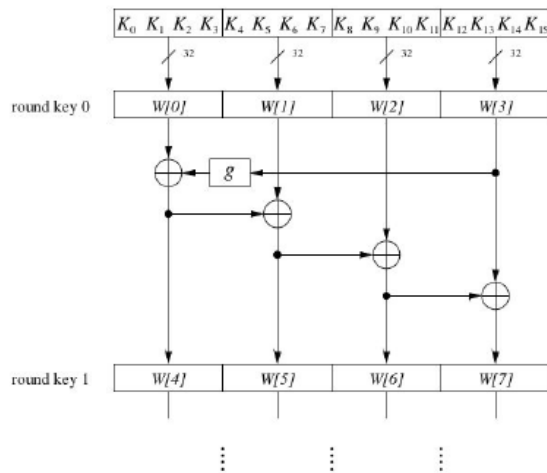
$$\begin{array}{rcl} 01 \cdot 25 &= & x^5 + x^2 + 1 \\ 01 \cdot 25 &= & x^5 + x^2 + 1 \\ 02 \cdot 25 &= & x^6 + x^3 + x \\ 03 \cdot 25 &= & x^6 + x^5 + x^3 + x^2 + x + 1 \\ \hline C_i &= & x^5 + x^2 + 1, \end{array}$$

### 3.3.3 Key Addition Layer

- 16-byte state matrix C and 16-byte subkey  $k_i$ 
  - Output:  $C \oplus k_i$
  - The subkeys are generated in the key schedule recursively from the original

Each round has 1 subkey, plus 1 subkey at the beginning of AES

**Key whitening:** Subkey is used both at the input and output of AES



For 128 bits

- Word-oriented: 1 word = 32 bits
- 11 subkeys are stored in  $W[0]...W[3]$ ,  $W[4]...W[7]$ , ...,  $W[40]...W[43]$
- First subkey  $W[0]...W[3]$  is the original AES key

Function  $g$  rotates its four input bytes and performs a byte-wise S-Box substitution  $\Rightarrow$  nonlinearity

RC (Round coefficient) is only added to the leftmost byte and varies from round to round (equals the number of the round in binary)

### 3.4 Decryption

All layers must be inverted for decryption.