

## Tutorial - 2

Q1 Write time complexity of below code

$$j = 1$$

$$i = 1$$

$$j = 2$$

$$i = i + 2$$

$$j = 3$$

$$L = 1 + 2 + 3$$

for (i)

$$\therefore 1 + 2 + 3 + \dots + < n$$

$$1 + 2 + 3 + \dots + m < n$$

$$\frac{m(m+1)}{2} < n$$

$$m \approx \sqrt{n}$$

By summation method,

$$\sum_{i=1}^m = 1 + 1 + \dots + \sqrt{n} \text{ times}$$

$$T(n) = \sqrt{n}$$

$$T.C = O(\sqrt{n})$$

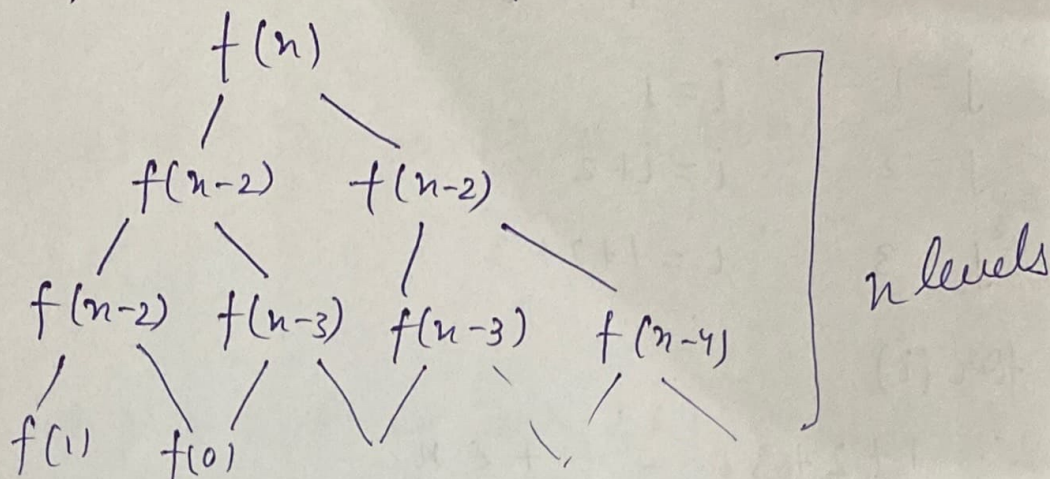


Q2

For fibonacci series

$$f(n) = f(n-1) + f(n-2) \quad f(0) = 0$$

$$f(1) = 1$$



$\therefore$  At every function call we get 2 function calls.

$\therefore$  for n levels

we have  $= 2 \times 2 \dots n$  times

$$\therefore T(n) = 2^n$$

### Space Complexity

Considering recursive

Stack :

no of cells maximum size

For each call we have S.C  $O(1)$

$$\therefore T(n) = O(n)$$

without considering recursive stack each call  
we have time complexity  $O(1)$

$$\therefore T(n) = O(1)$$

Q3

i) void quicksort (int arr[], int l, int n)

```
{  
    if (l < n)  
    {  
        int p = partition (arr, l, n);  
        quicksort (arr, l, p-1);  
        quicksort (arr, p+1, n);  
    }  
}
```

int partition (int arr[], int l, int n)

```
{  
    int pivot = arr[n];  
    int i = l-1;  
    for (j = l; j <= n-1; j++)  
    {  
        if (arr[j] < pivot)  
        {  
            i++;  
            swap (&arr[i], &arr[j]);  
        }  
    }  
}
```

}



ii)  $n^3$

for ( $i=0; i < n; i++$ )

for ( $j=0; j < n; j++$ )

for ( $k=0; k < n; k++$ )

$res[i][j] += arr[i][k] * b[k][j];$

iii)  $\log(\log n)$

for ( $i=2; i < n; i=i*i$ )

count++;

Q5

for	$i$	$j$
	1	1
	2	1+3+5
	3	1+4+7
	⋮	⋮
	$n$	1+5+9

$j = (n-1)$  times

$$\sum_{i=1}^n \left( \frac{n-1}{i} \right)$$

$$T(n) = n-1 + \frac{n-1}{2} + \frac{(n-1)}{3} + \dots + \frac{n-1}{n}$$

$$T(n) = n \left[ 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} \right] - 1 \times \left[ 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} \right]$$

$$= n \log n - \log n$$

$$T(n) = O(n \log n)$$

	$i$	where
for	$2^1$	
	$2^k$	$2^k < m$
	$2^{k^2}$	$k^m = \log_2 n$
	$2^{k^3}$	$m = \log^k \log_2 n$
	$2^{k^n}$	

$$\therefore \sum_{i=1}^m 1 + 1 + 1 + \dots m \text{ times}$$

$$T(n) = O(\log_k \log n)$$

Q8

a)  $100 < \log \log n < \log n < (\log n)^2 < \sqrt{n} < 8n < n \log n < \log(n!) < n^2 < 2^n < n^n < 2^{2n}$

b)  $1 < \log \log n < \sqrt{\log n} < \log n < \log 2n < 2 \log n < n < n \log n < 2n < 4n < \log(n!) < n^2 < n! < 2^{2n}$

c)  $96 < \log_8 n < \log_2 n < 5n < n \log_8(n) < n \log_2 n < \log(n!) < 8n^2 < 7n^3 < n! < 8^{2n}$

